

## 2019 MSOC Assignment 2

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### 0. Generator / Coroutine

#### a. Complete program execution flow

[Python execution flow]

f0	g0	clk	f1	g1	clk	f2	g-2	clk	f3	g-fin	clk	f4	g-fill	clk	f-fin
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[Concept level execution flow]

f0	clk	f1	clk	f2	clk	f3	clk	f4	clk	f-fin
g0		g1		g2		g-fin		--		--

#### b. Explain how *zip\_iterator* works

As shown in the code, for each of the passed in iterables, *zip\_iterator* will make an iterator that aggregates their elements. For *zip\_longest* function used in the code, it will fill a *fillvalue* for iterables with shorter length (in this case, no outputs will be generated if the iterable is shorter).

### 1. Producer / Consumer Revisit

#### a. What if we use Consumer(11) ?

The process will never stop, as the Consumer can never get the 11<sup>th</sup> item due to the limited item (in this case, 10) produced by the producer.

#### b. What if we swap the order in *main\_loop()* ?

##### i. Is there any difference ?

Yes, swapping results in different results. From the table above we can see that after swapping, the printed pattern shows different routine. A notable fact is that it never put and get the item in the same cycle.

Before swap	=== clk === === clk === === clk === Put an item === clk === Get an item === clk === === clk === Put an item Get an item === clk === === clk === === clk === Put an item === clk === Get an item === clk === === clk === Put an item Get an item	After swap	=== clk === === clk === === clk === Put an item === clk === Get an item === clk === === clk === Put an item === clk === === clk === Get an item === clk === Put an item === clk === Get an item === clk === === clk === Put an item === clk ===

ii. Is it reasonable ?

Yes. After swapping, the execution of request changes so that the *Consumer* will always check *n\_item* first, causing the increase of *n\_item* (push item) and reduce of *n\_item* (get item) occur in different cycles.

2. Event (Implement all TODOs in 2\_2\_event\_fifo.py)

3. dont\_initialize

a. What does *dont\_initialize* mean in SystemC ?

For simple process construction in SystemC, the declared process will always run once during the construction. If a *dont\_initialize()* function is set, the process will not be executed until the clock signal starts.

b. How do you implement it ?

The core concept of *dont\_initialize()* is that every process should not be waiting any events other than **INIT\_EVENT** before the clock first triggered, i.e. the first process handled by the triggered **CLOCK\_EVENT** should always be the *Clk()* process. Thus, I implement *dont\_initialize* by simply adding a yield **INIT\_EVENT** in every process.

c. Give an example about the difference with/without *dont\_initialize*.

The following table shows a simple difference of the two scenarios. Without *dont\_initialize*, something strange like putting and getting item in the same cycle may occur, while this never happened with *dont\_initialize*. This is because at the first triggered **CLOCK\_EVENT**, the first handled process is the *Producer()* process, causing the **WRITE\_EVENT** triggering will be pushed to the *event\_pending list* earlier than the next **CLOCK\_EVENT**, resulting in the handling of get item ("Get an item") process before the next clock edge ("handling 7"). On contrast, the with *dont\_initialize* scenario will always trigger the **CLOCK\_EVENT** earlier than any process triggering this cycle, causing the result that there must always exist a clock edge ("handling 7") between putting and getting an item.

Without dont_initialize	With dont_initialize
handling 5 handling 7 handling 7 handling 7 Put an item handling 11 Get an item handling 7 handling 7	handling 5 handling 7 handling 7 handling 7 Put an item handling 7 handling 11 Get an item handling 7