## **CS311 - Computer Architecture Lab**

## **Assignment 5 - Report**

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## 1. Statistics Table:

- Total no. of cycles taken by each benchmark program.
- Total no. of instructions executed.
- The throughput in terms of instructions per cycle

All the calculations for the lab assignment are tabulated below.

|            | Total No. of Cycles | Total No. of instructions | Throughput  |
|------------|---------------------|---------------------------|-------------|
| descending | 11433               | 365                       | 0.03192513  |
| evenorodd  | 249                 | 6                         | 0.024096385 |
| fibonacci  | 3157                | 94                        | 0.029775104 |
| palindrome | 1983                | 56                        | 0.02824004  |
| prime      | 1175                | 34                        | 0.02893617  |

## 2. Interpretation and Comments:

For <u>odd-even</u>, code has no control hazards because the conditional branch was not taken as the value checked was odd. There are a few RAW hazards due to which we see a small number of data hazards.

For <u>fibonacci</u>, there is a control hazard in the loop. Every time it goes in the loop, it is a wrong branch, till the count reaches 10. So, we see a few control hazards. Similarly, there are a few data hazards outside the loop and a few inside. Hence, we also see some stalls in the pipeline.

For <u>descending</u>, code has a lot of branching, 3 nested loops and data hazards. Hence, we see that the total number of hazards increases by a large amount and so does the cycles executed.

For <u>Prime</u>, the loop runs from n=2 to n=5, i.e. 4 times. All these 4 times, we get a wrong branch taken for the jmp statements. Then it takes a wrong branch when it is supposed to enter prime/not prime. Hence, a total of 10. Similarly, we are getting a few data hazards.

We may conclude that that every time the code misses the jump to loop, there is a wrong branch taken, thus we get a control hazard. But we have more data hazards per loop. Hence, we see a larger number of data hazards than control hazards. We additionally add NOP instructions to avoid hazards and reorder the code. Here in all cases we can commonly observe that throughput is almost the same as it is basically a proportion of instruction per cycle. Here, it is in between between 0.024 and 0.032.