

# BANGALORE INSTITUTE OF TECHNOLOGY

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Department of Computer Science & Engineering

## Automata Theory and Computability Project (18CS54)

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$L = \{w \in \{a, b\}^* : \text{no two consecutive characters are the same}\}.$

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Construct a Turing Machine that can accept the set of all even palindromes over  $\{0,1\}$ .

3. Grammar

For the given Grammar  $G = S \rightarrow aS \mid B, B \rightarrow bB \mid [ L = \{w \in \{a,b\}^*, \#a(w) = \#b(w)\}$ .  
Construct a parse tree and derivation for the following string  $aaaabb$ .

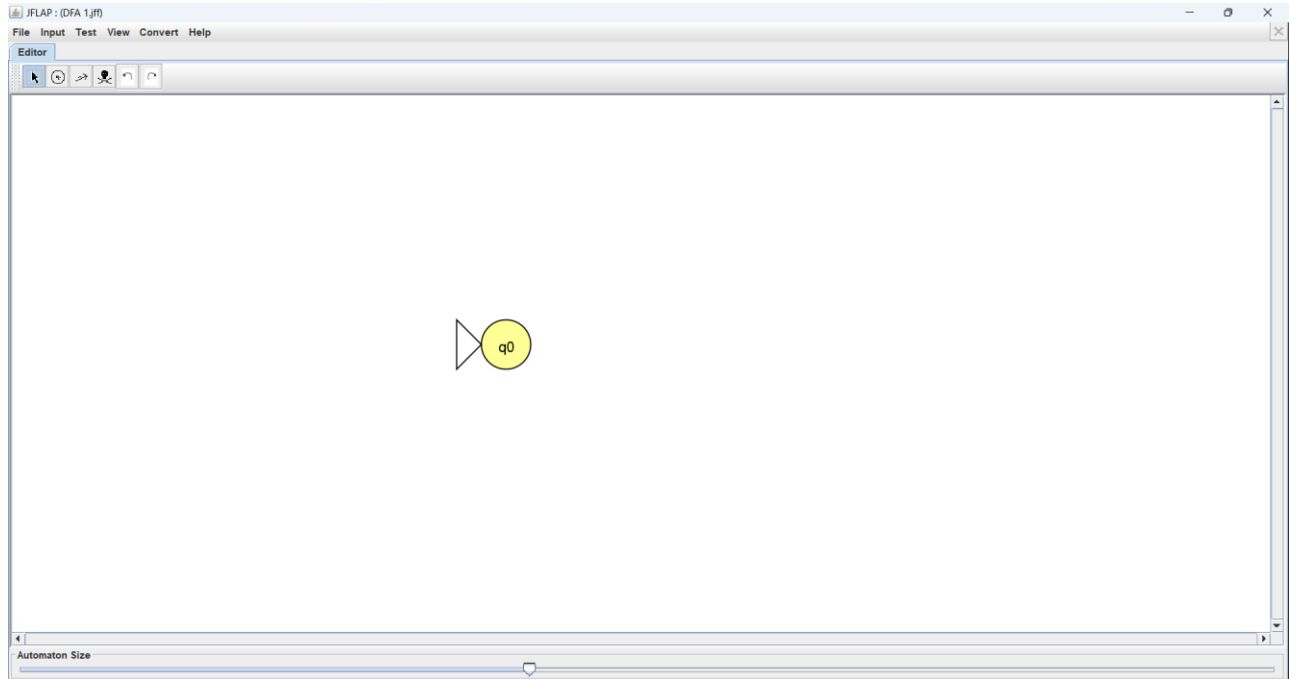
4. Push Down Automata

Construct a push down automata for the language  $L = \{a^n b^n : n \geq 0\}$ .

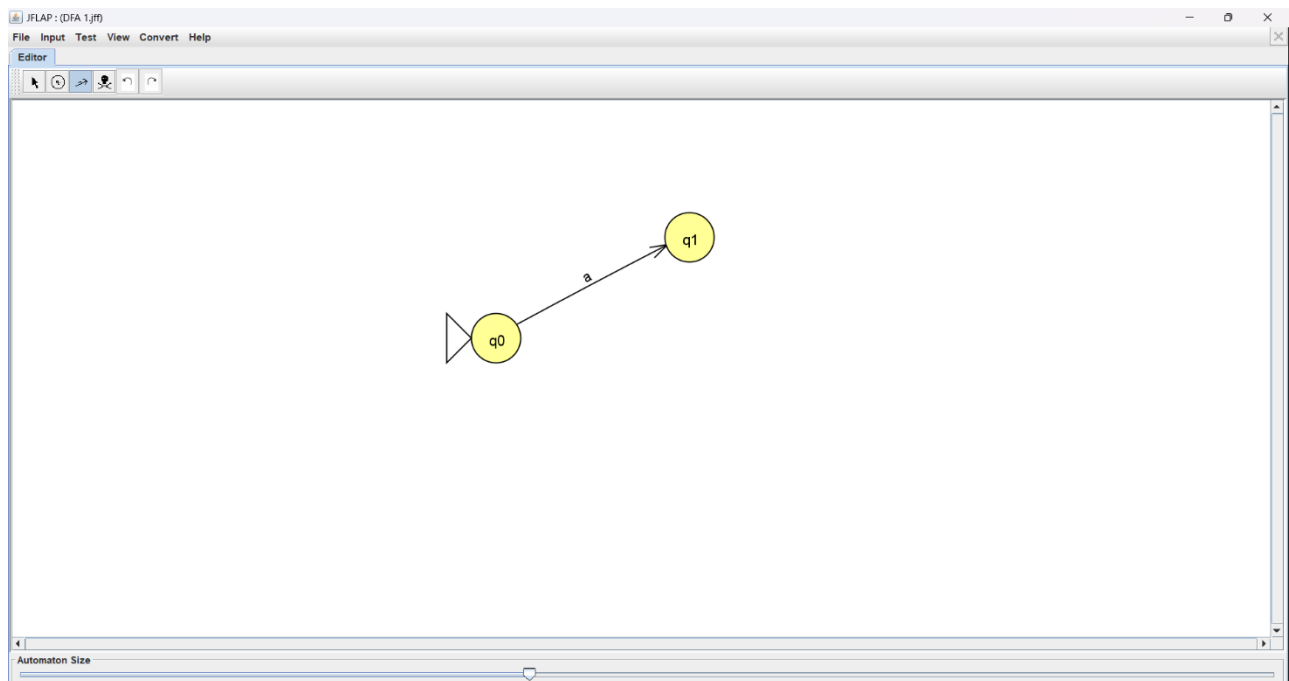
## DETERMINISTIC FINITE AUTOMATA:

$L = \{w \in \{a, b\}^* : \text{no two consecutive characters are the same}\}.$

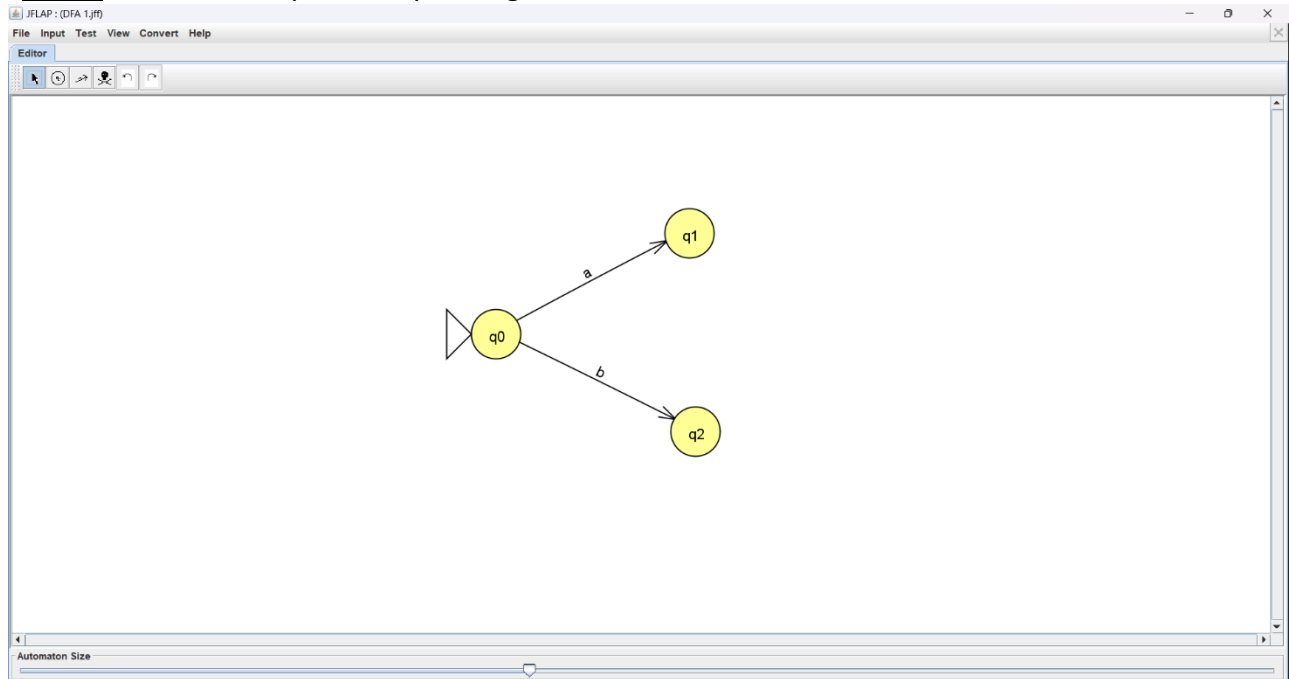
Step 1:  $q_0$  is the initial state of the machine.



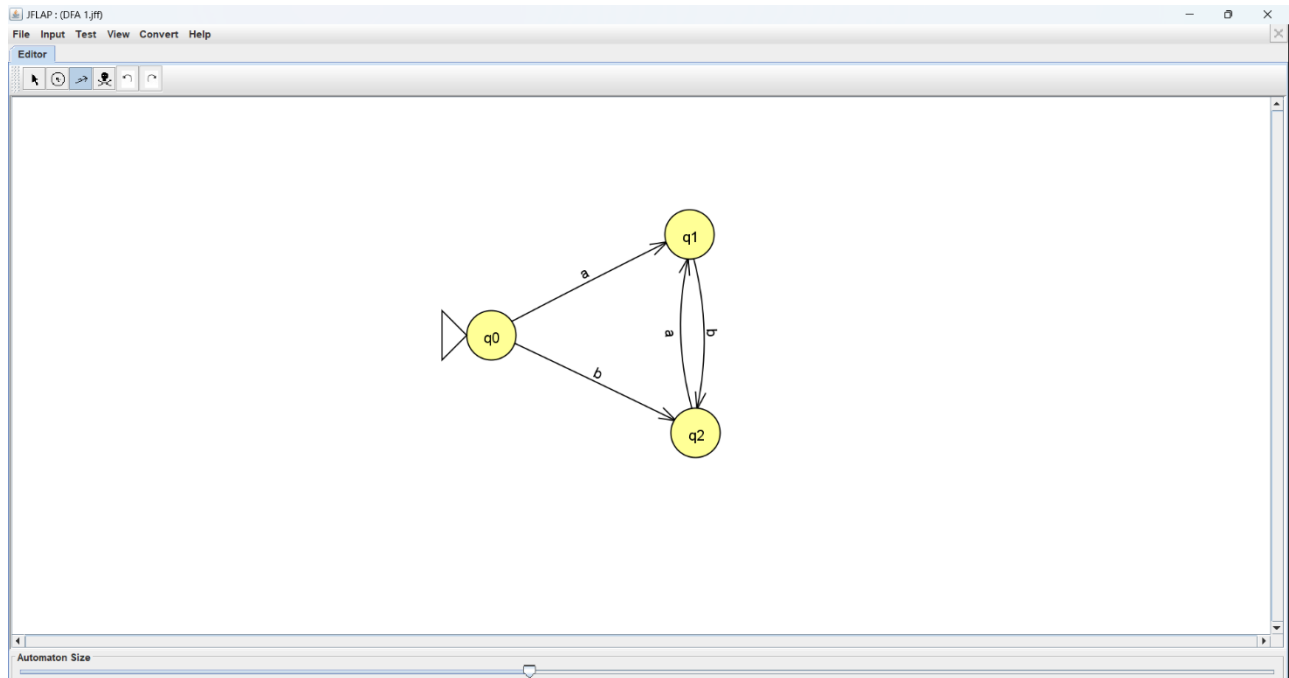
Step 2: Create state  $q_1$  to accept string 'a'.



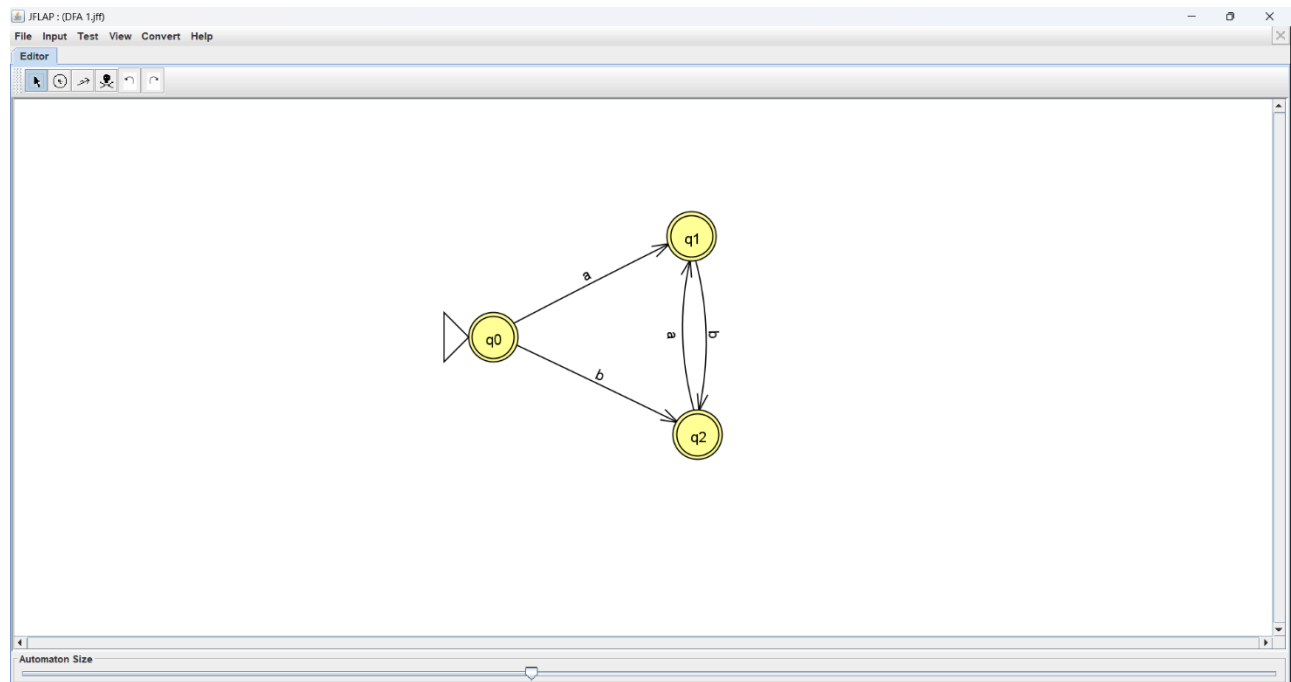
Step 3: Create state q2 to accept string 'b'.



Step 4: State q1 goes to state q2 on accepting string b and state q2 goes to state q1 on accepting string a.

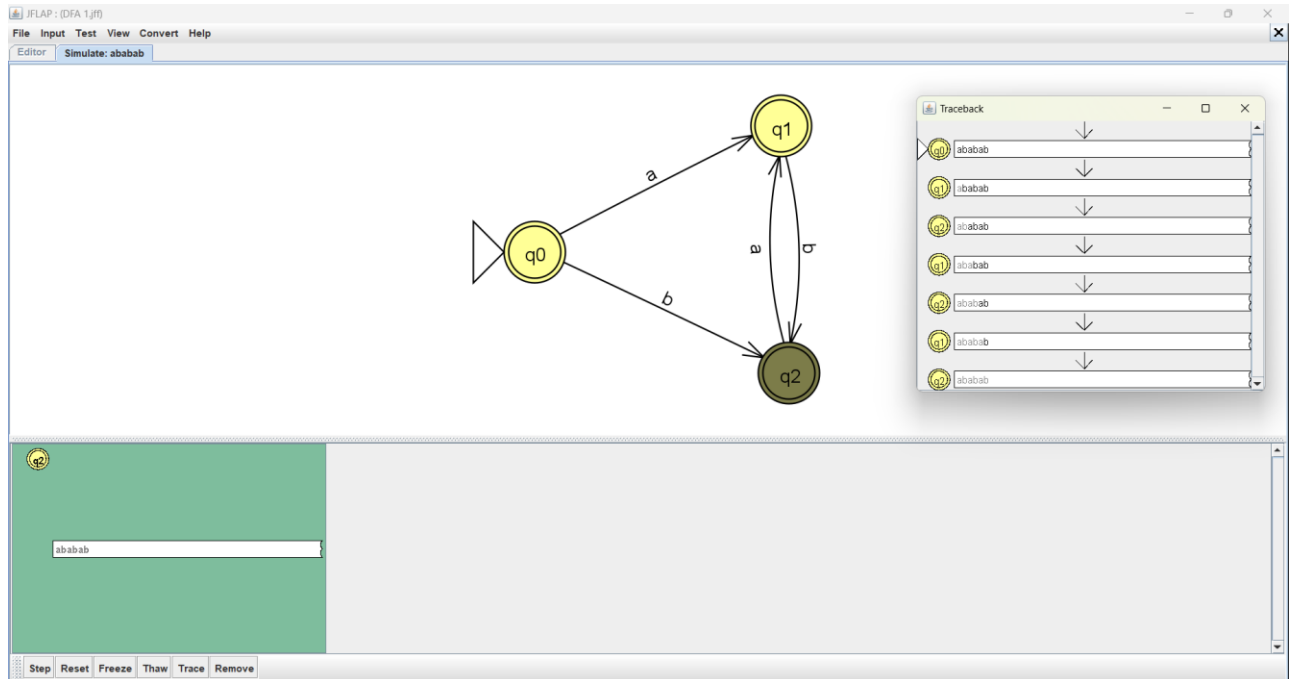


Step 5: State's q0, q1 and q2 are the final state's.

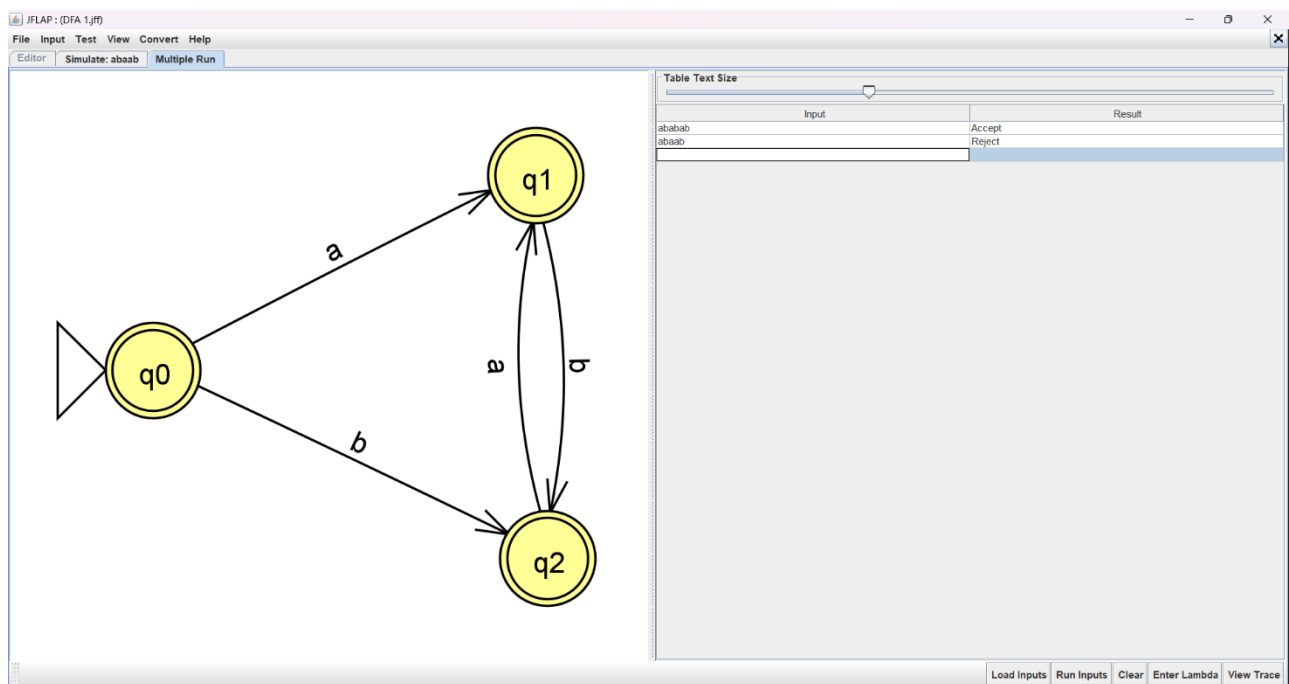


## Trace

Tracing for string 'ababab' we get accepted as it satisfies the above DFSM condition.



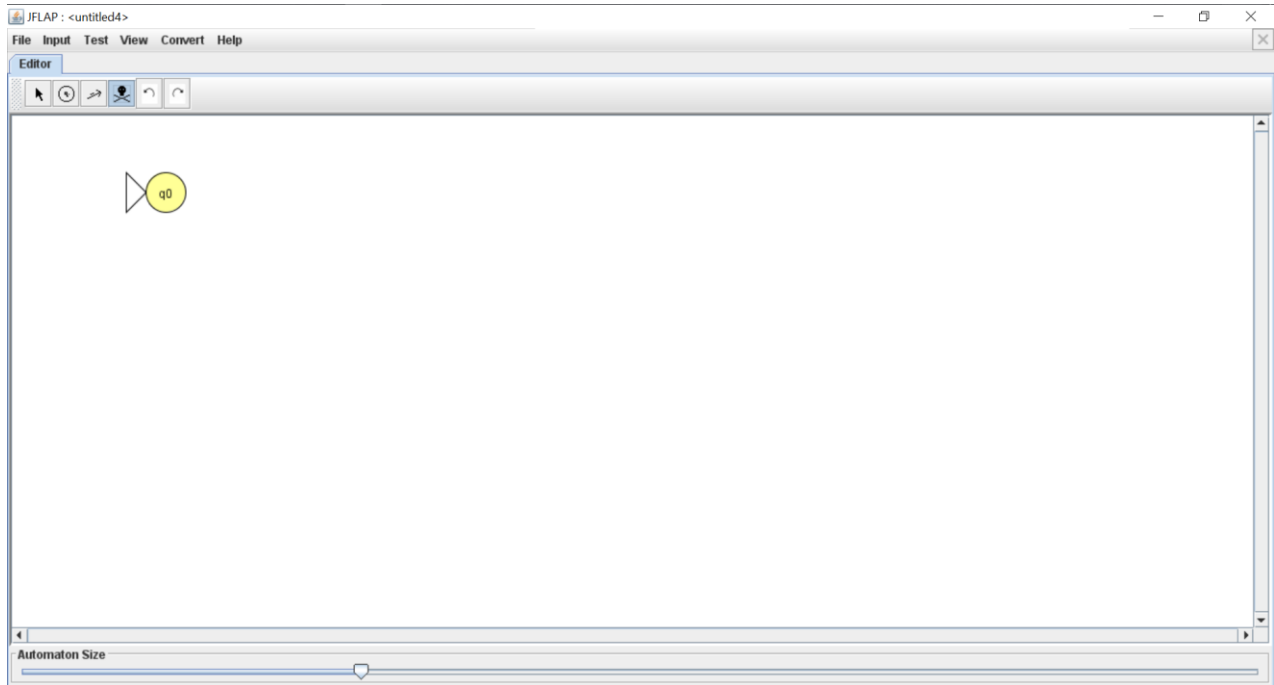
When checking multiple strings, i.e., 'ababab' and 'abaab'. We observe that 'ababab' gets accepted and 'abaab' gets rejected. 'abaab' gets rejected because it doesn't follow the DFSM conditions above.



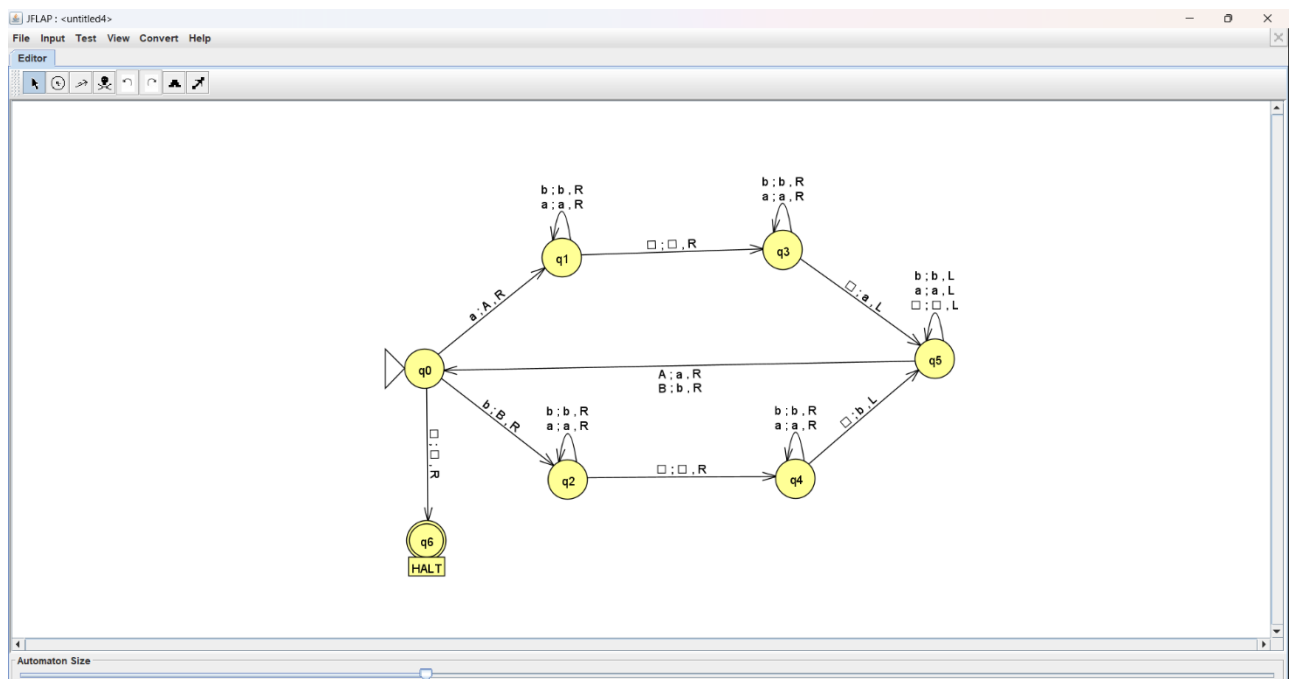
## TURING MACHINE:

**Construct a Turing Machine that can accept the set of all even palindromes over  $\{0,1\}$ .**

**Step 1:** Create a new initial state  $q_0$ . This is not the final state.

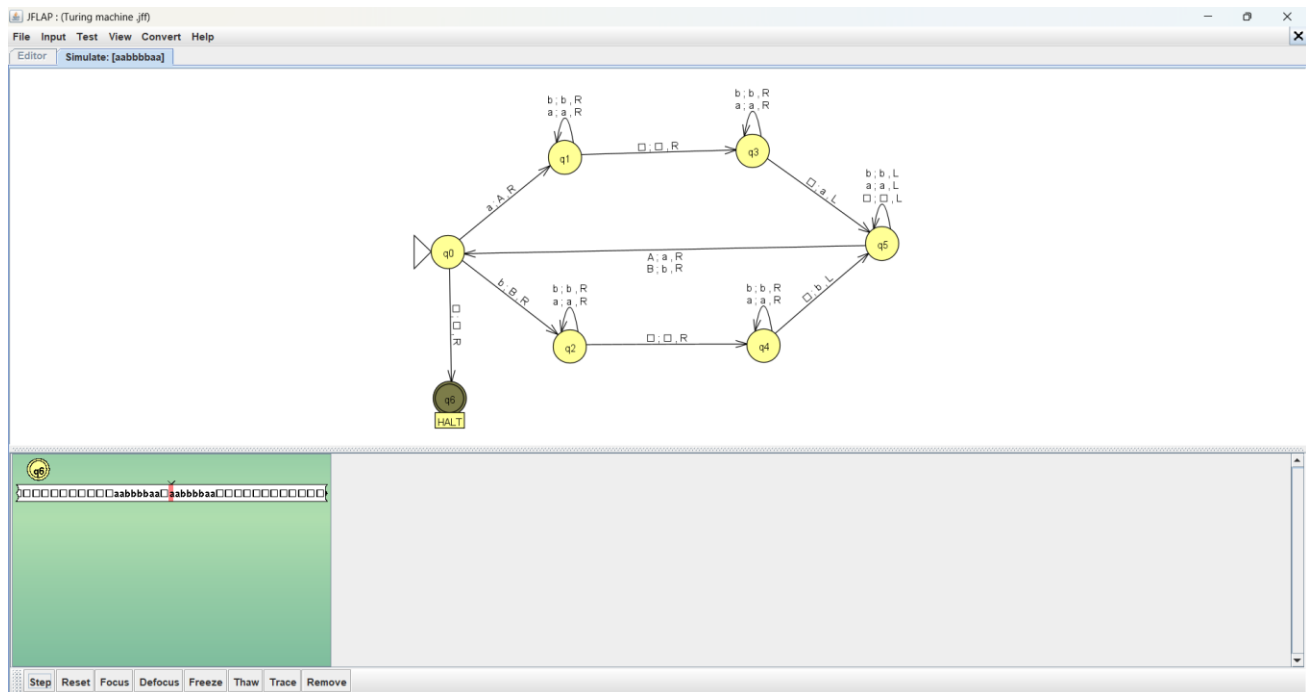
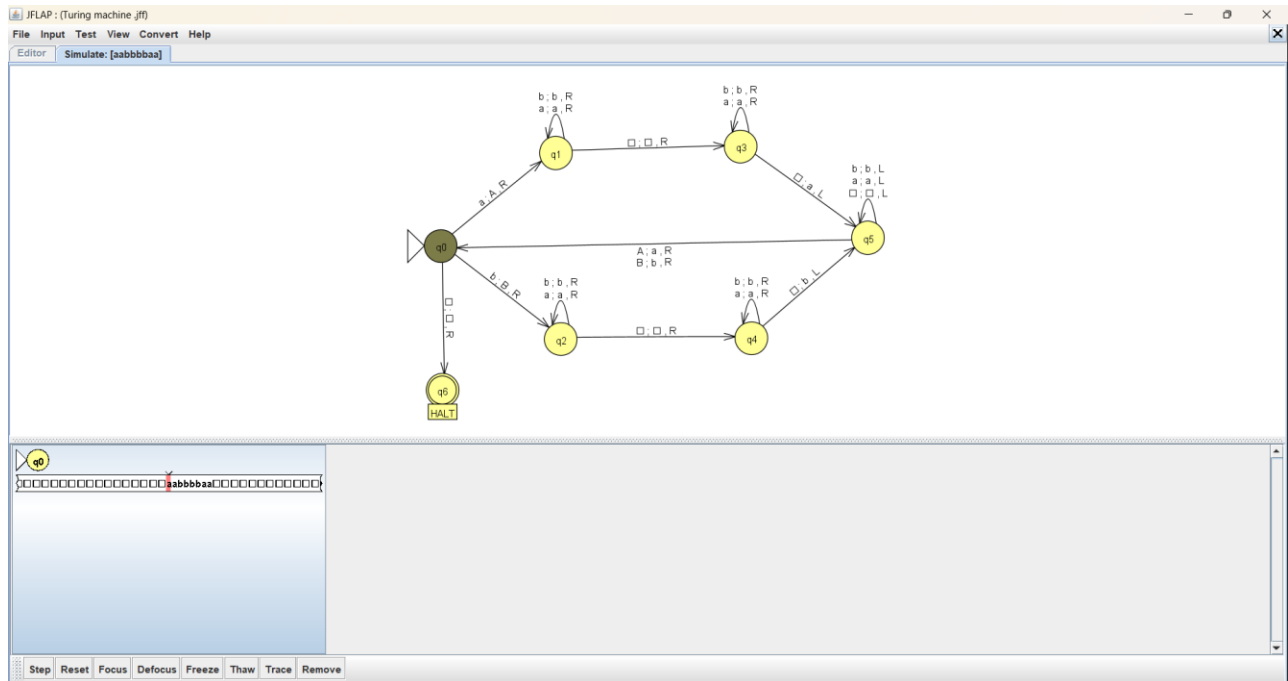


**Step 2:** This is the Transition Diagram for the required Turing Machine that can accept the set of all even palindromes over  $\{0,1\}$ ..

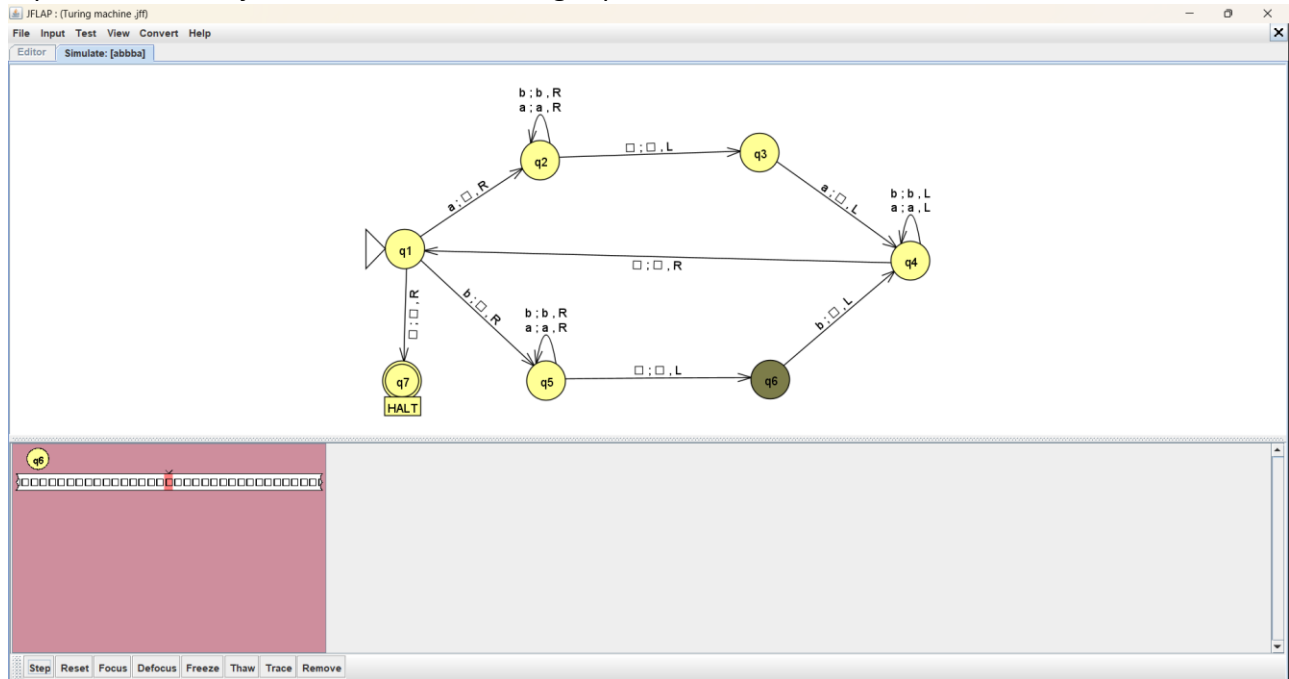


## Trace

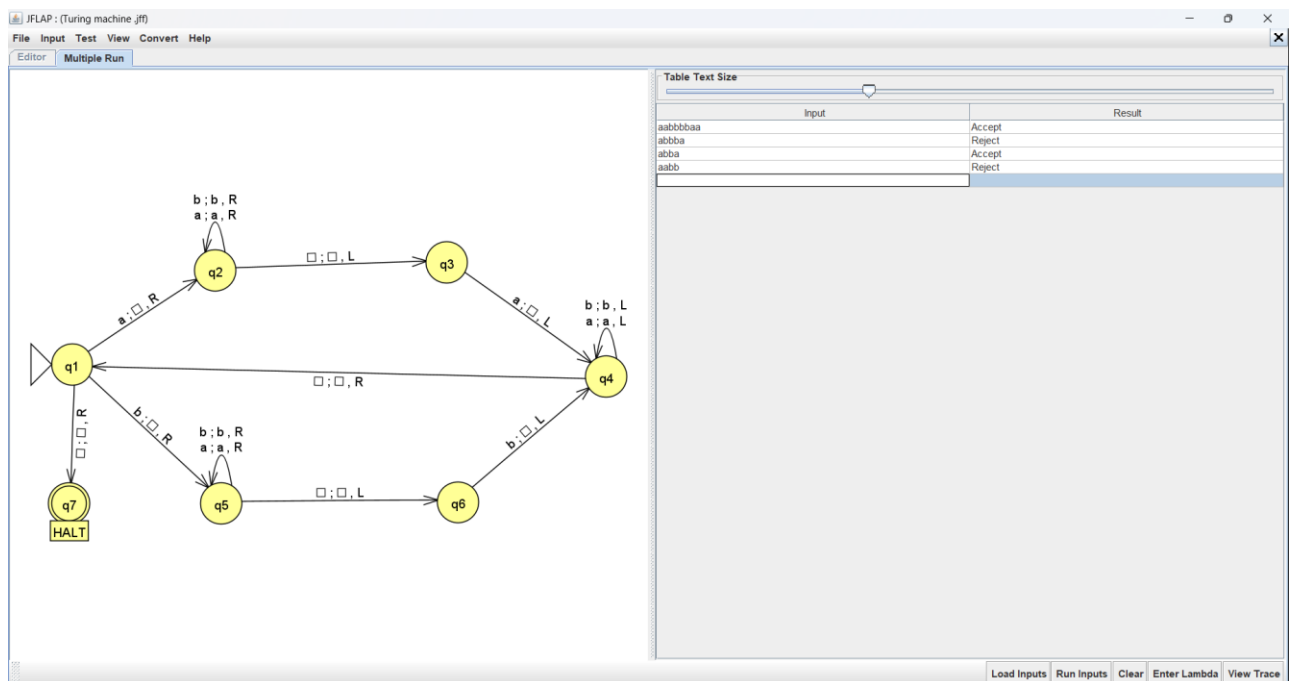
Tracing for 'aabbbbbaa', we get an accepted state. Also, this a valid string according to our conditions.



Input 'abbba' is rejected as it is an odd length palindrome.



For different type of inputs, we get these various answers -

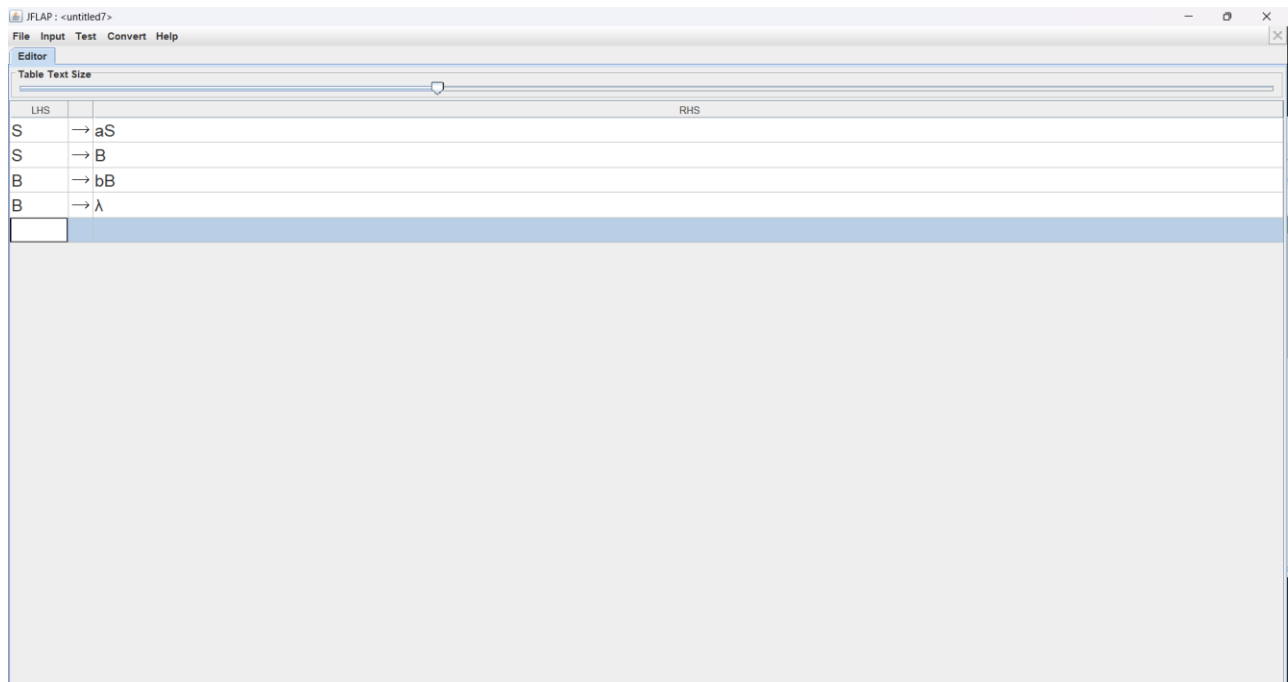




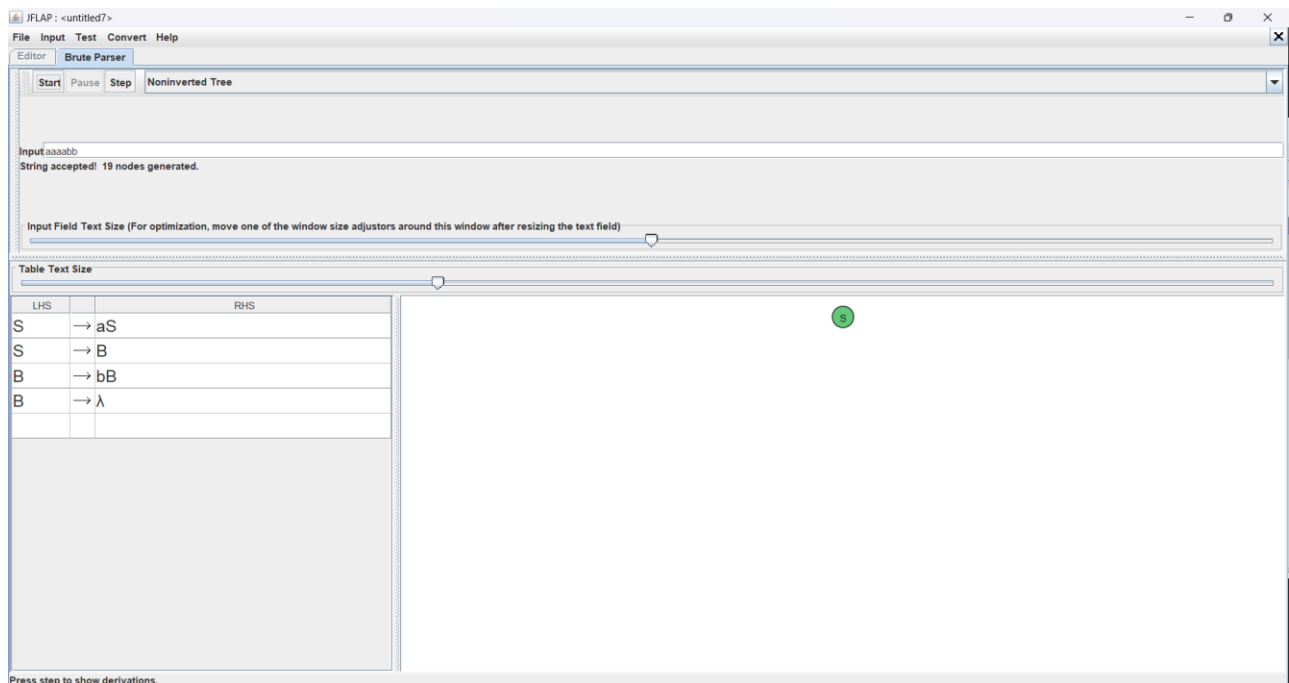
## GRAMMER

For the given Grammar  $G = S \rightarrow aS \mid B, B \rightarrow bB \mid \lambda$  [  $L = \{w = \{a,b\}^*, \#a(w) = \#b(w)\}$  ]. Construct a parse tree and derivation for the following string aaaabb.

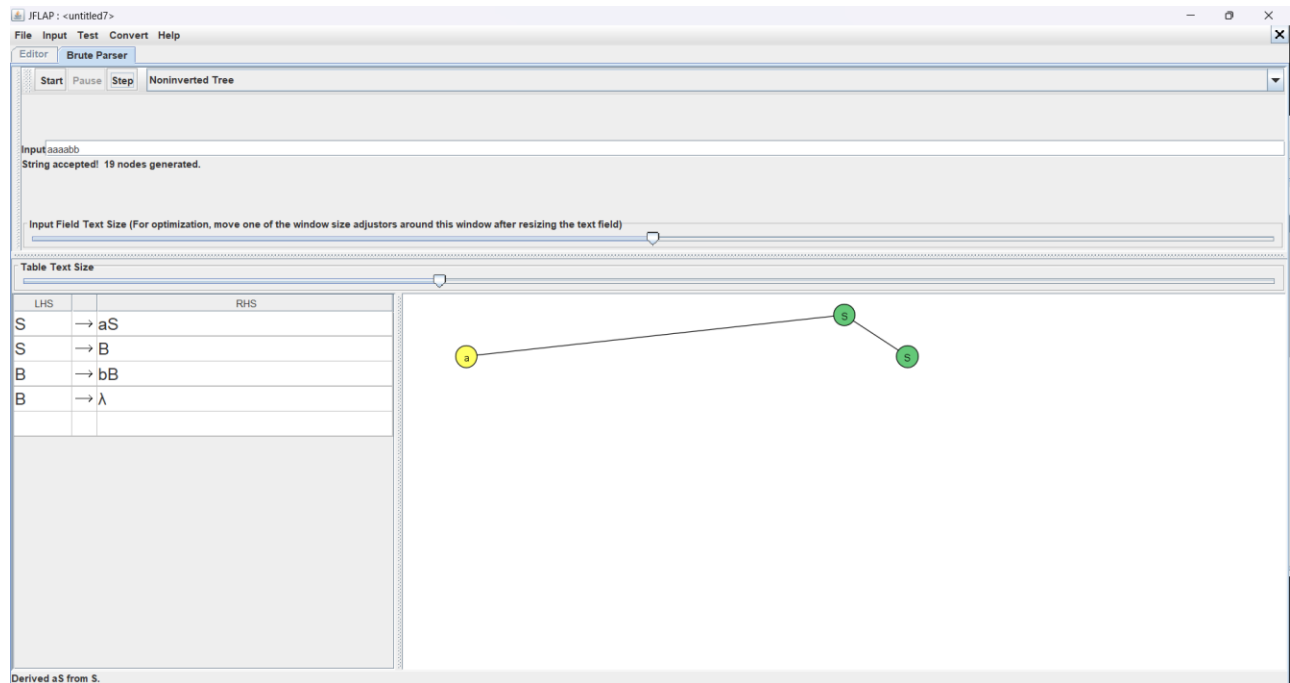
Step 1: The given Grammer is -



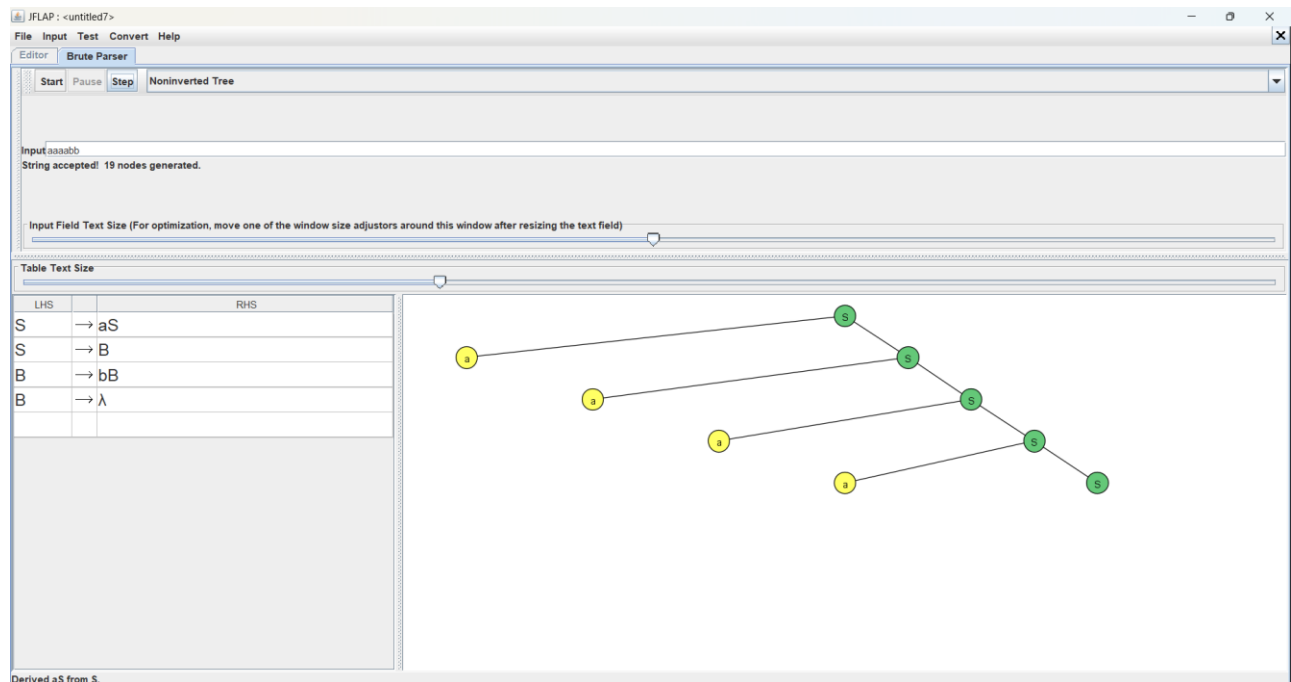
Step 2: S is the root node of the parse tree, which is a non-terminal.



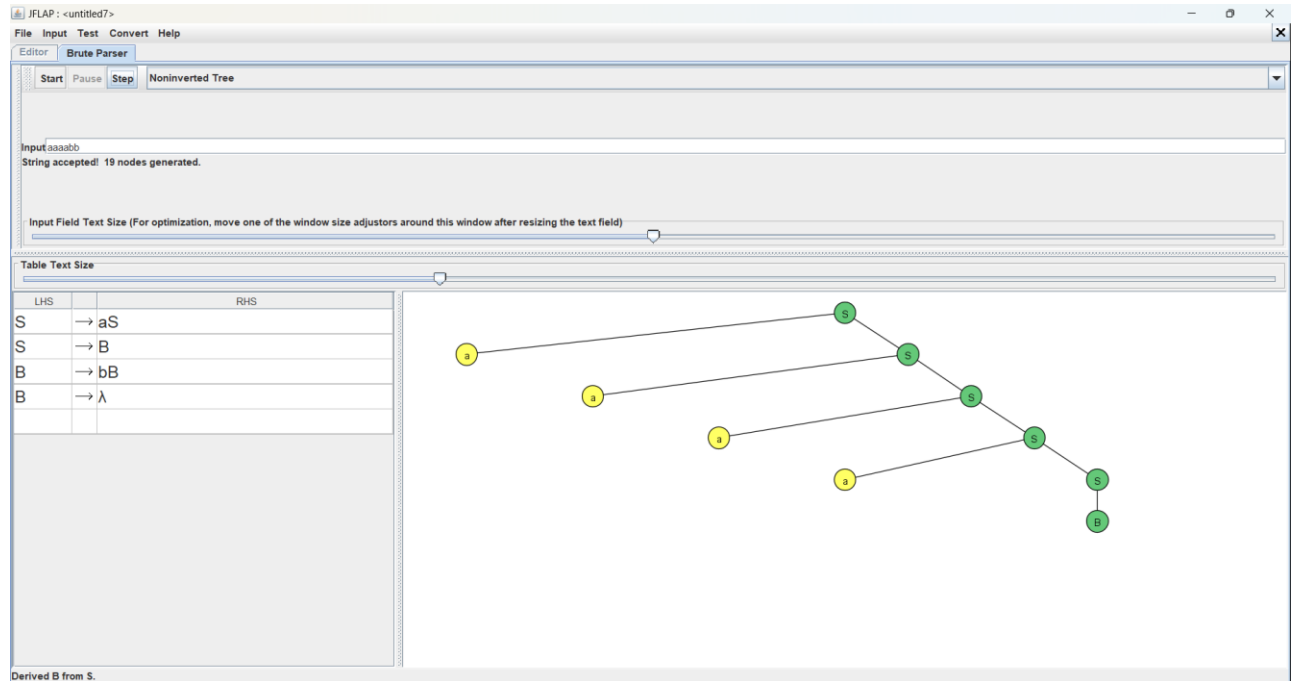
Step 3: The input string aaaabb is split into three sub trees. Hence the root node splits into a, S and S.



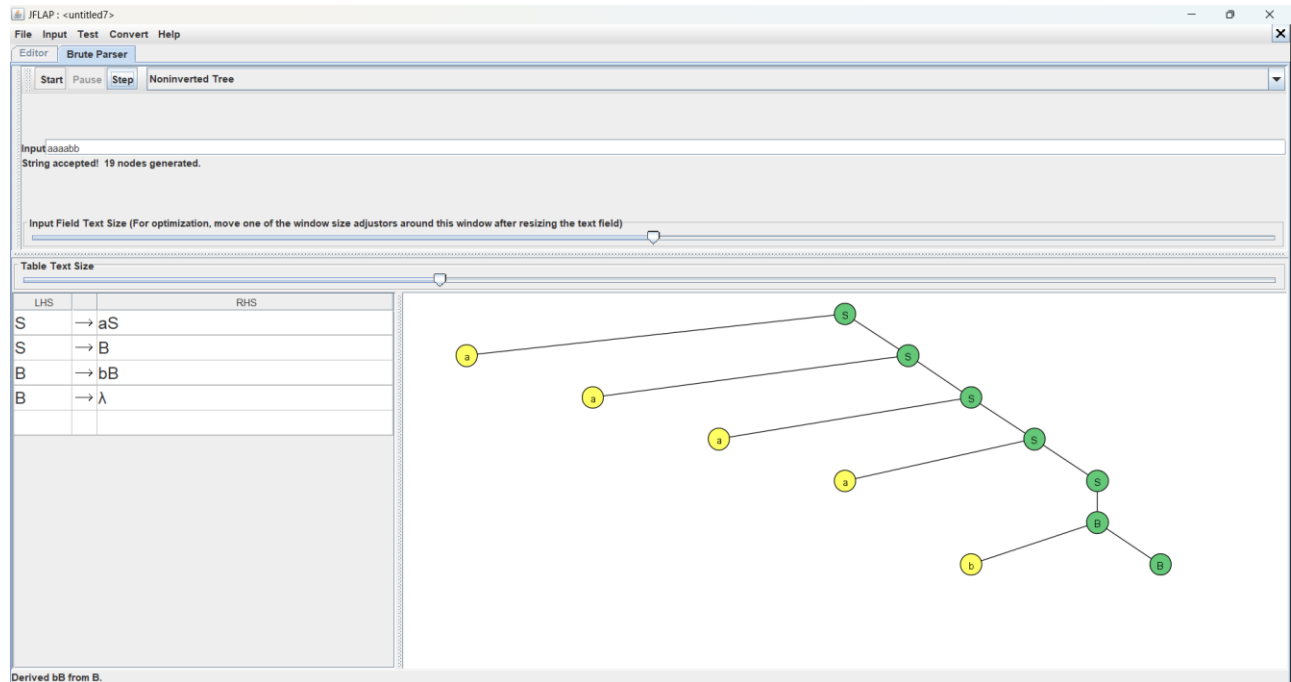
Step 4: The mid tree S which is aaab further splits into aS and hence has three branches towards mid.



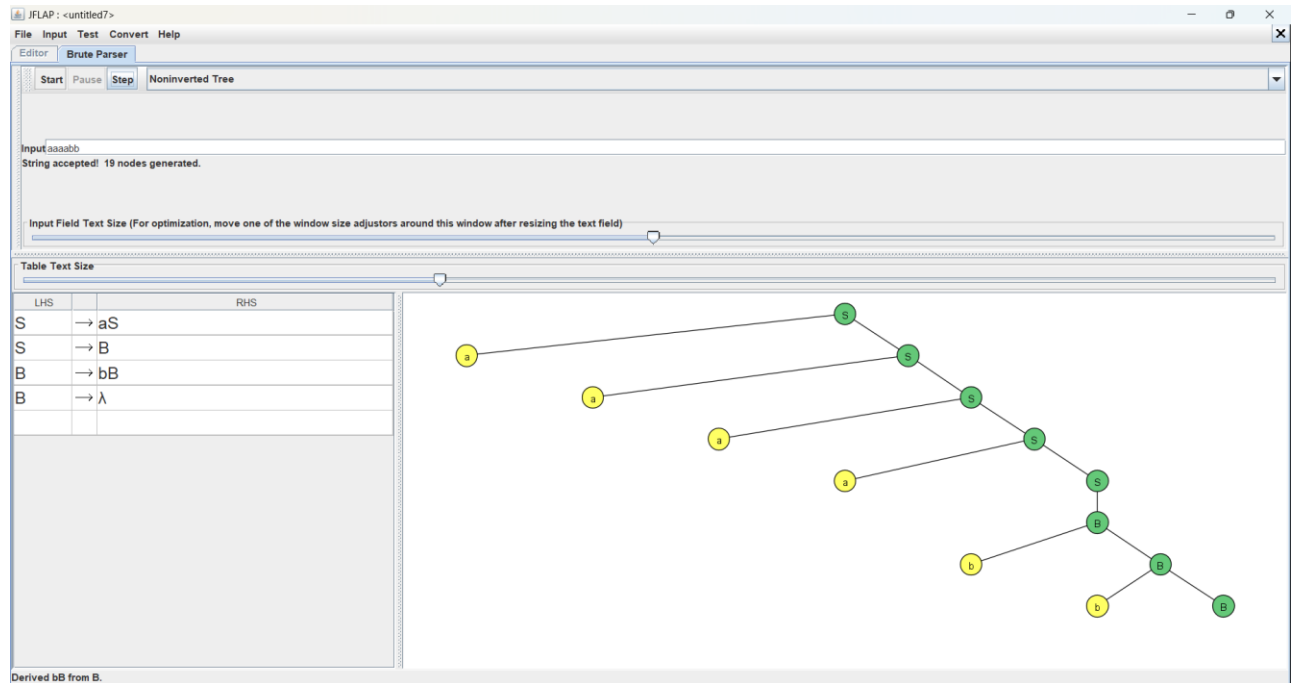
**Step 5:** The S in 4<sup>th</sup> level which has aab splits into SB and hence has one branches B.



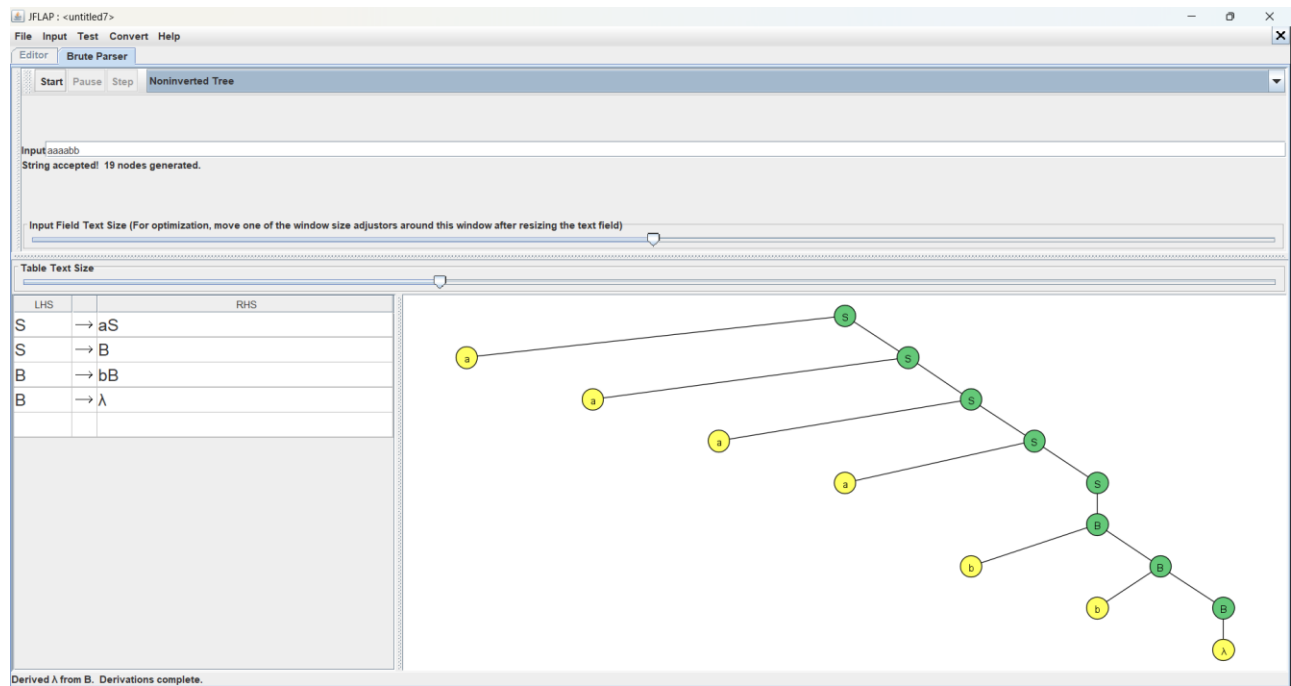
**Step 6:** The B in 5<sup>th</sup> level further splits into bB and hence has two branches b and B.



Step 7: The B in the 7<sup>th</sup> level further splits into bB and hence has two branches b and B.



Step 8: The B can now have an empty string.



## Derivation Table

The screenshot shows the JFLAP software interface. The top menu bar includes File, Input, Test, Convert, and Help. Below the menu bar are tabs for Editor, Multiple Run, and Brute Parser. The Brute Parser tab is active, showing a Start button, a Pause button, and a Step button. The main area displays the input string 'aaaabb' and the message 'String accepted! 19 nodes generated.' Below this is a text field for the input string, with a hint: 'Input Field Text Size (For optimization, move one of the window size adjusters around this window after resizing the text field)'. A slider for 'Table Text Size' is also visible. The Derivation Table is shown below the text field, with columns for LHS and RHS. The table contains the following entries:

LHS	RHS
S	→ aS
S	→ B
B	→ bB
B	→ λ

Below the table, a list of strings is shown, representing the derivation process:

Table Text Size
S
aS
aaS
aaaS
aaaaS
aaaaB
aaaaBB
aaaaBBB
aaaaBBB
aaaaBBB
aaaaBBB

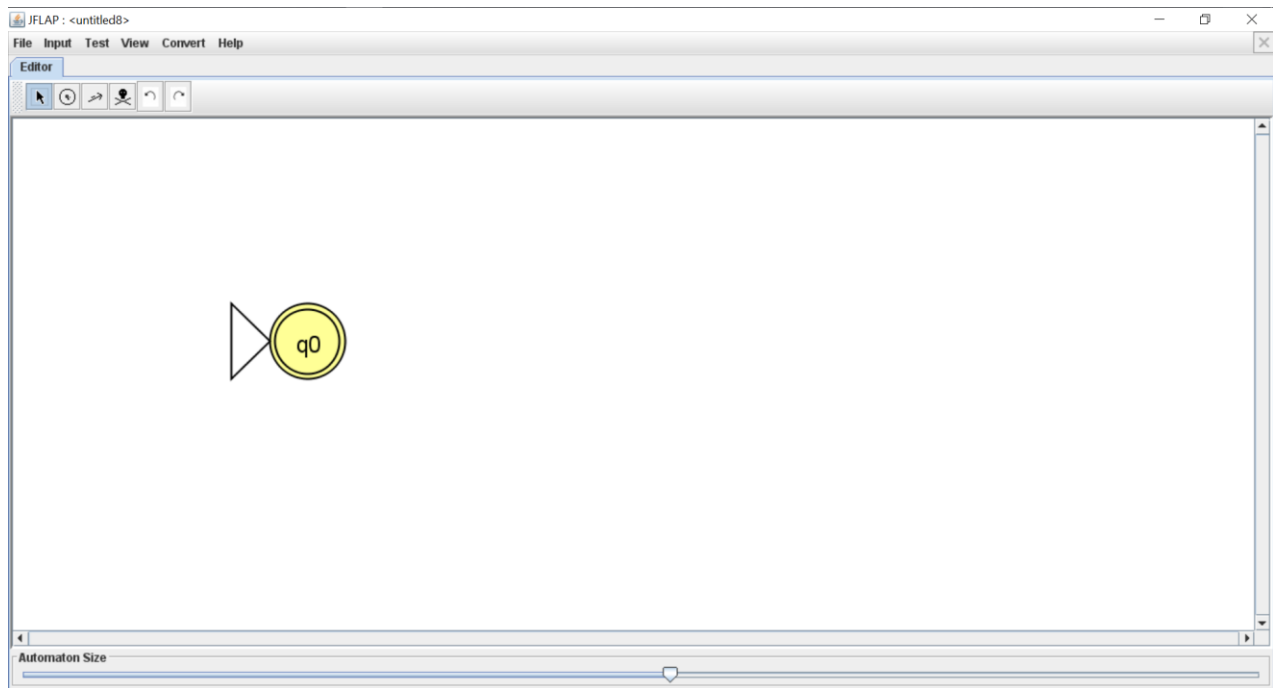
At the bottom of the window, a status bar indicates 'Derived A from B. Derivations complete.'

Since the string could be derived using a parse tree and derivation table, the string **aaaabb** is accepted.

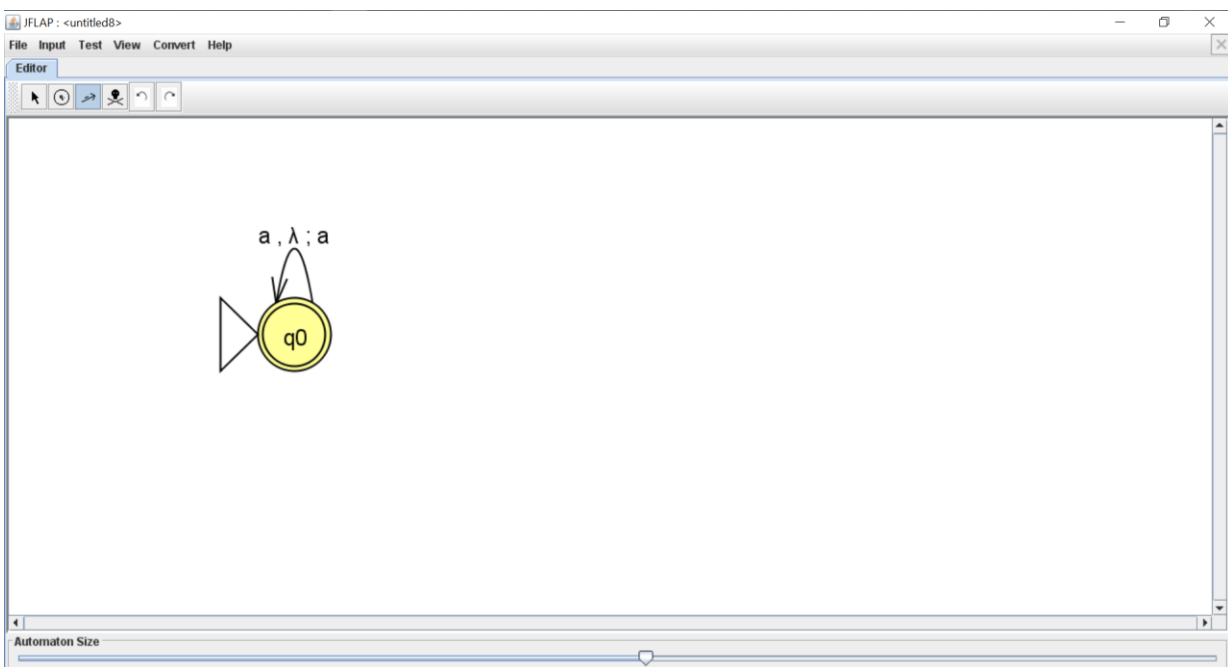
## PUSH DOWN AUTOMATA

Construct a push down automata for the language  $L = \{ a^n b^n : n \geq 0 \}$ .

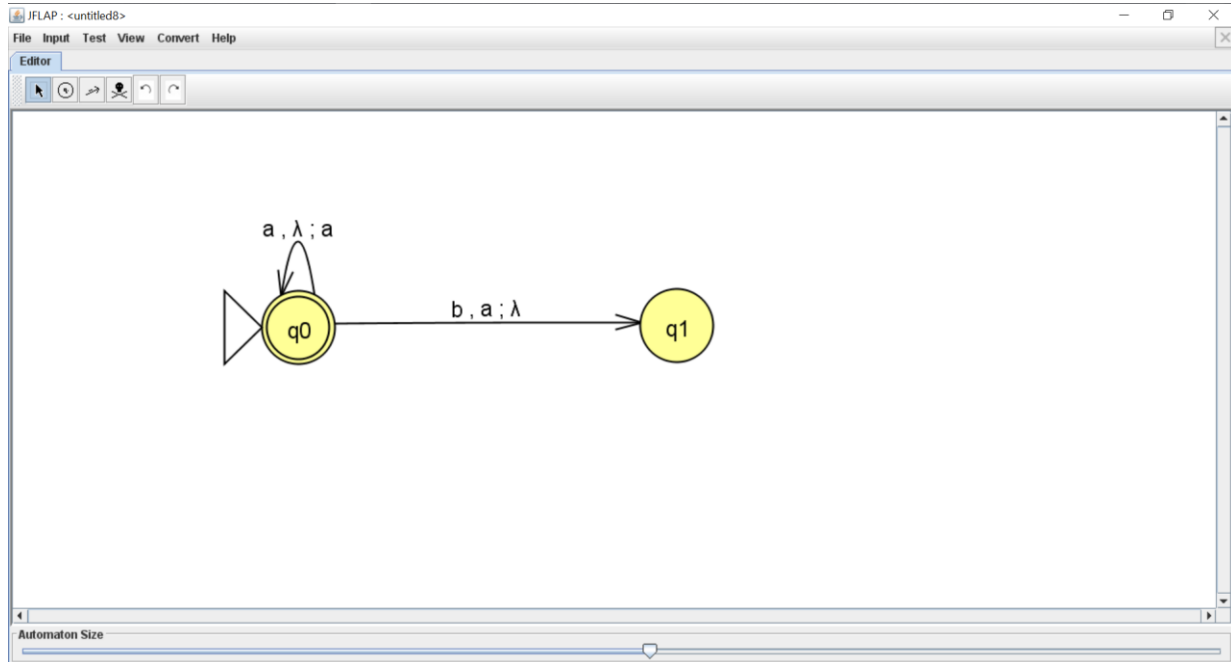
Step 1: Create a new initial state  $q_0$ . It is a final state since the language doesn't accept empty strings.



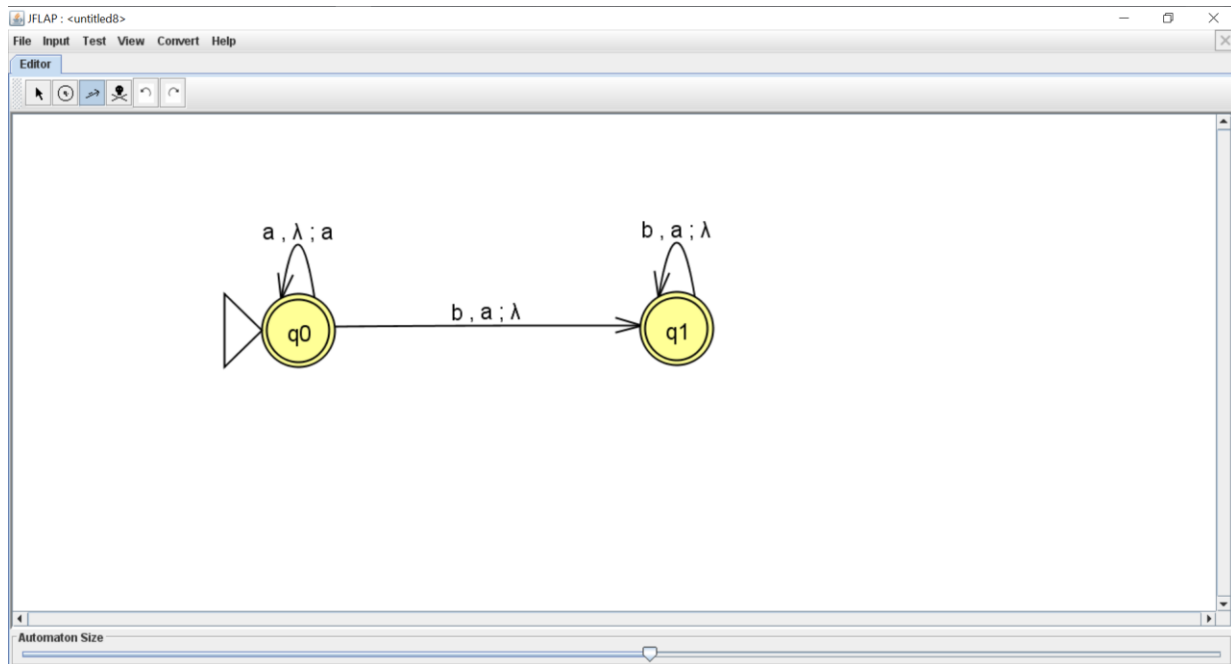
Step 2: Self-transition on  $q_0$  – if the input symbol is  $a$  and the stack is empty,  $a$  is pushed onto the stack.



Step 3: b transition from state q0 to q1.



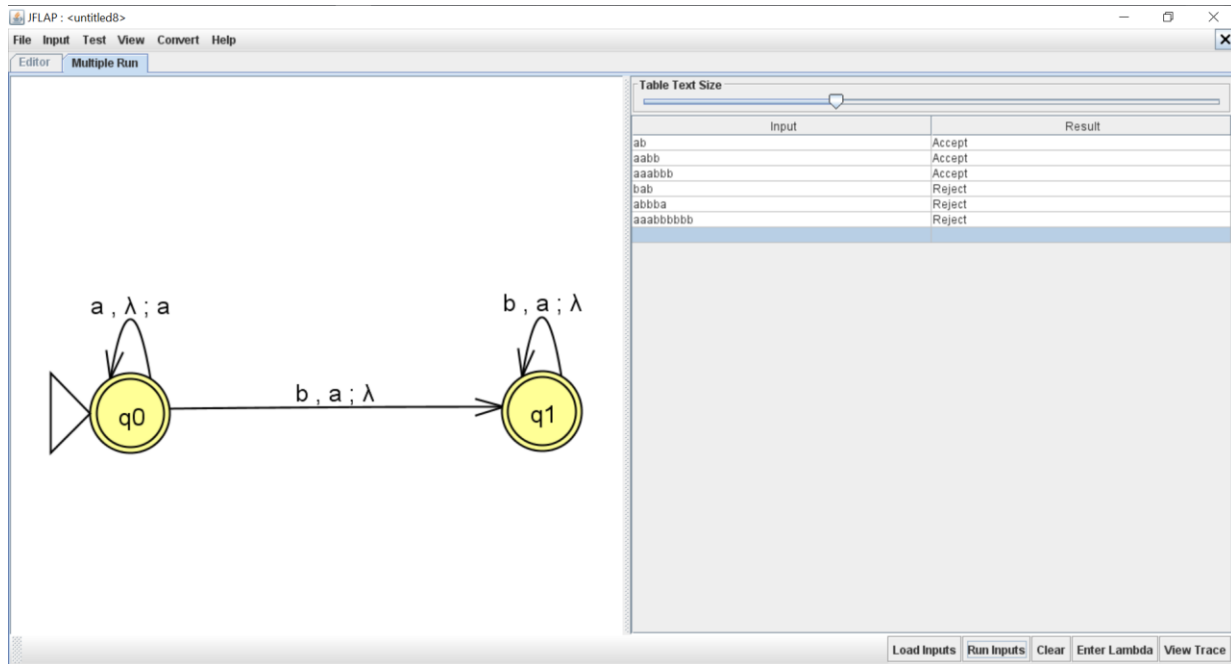
Step 4: State q1 is made the final state. Self-transition on q1 – if the input symbol is b and a is on top of the stack then symbol on top is popped off.



Step 6: Testing the PDA with multiple inputs and verifying the results.

Accepting strings: ab, aabb, aaabbb

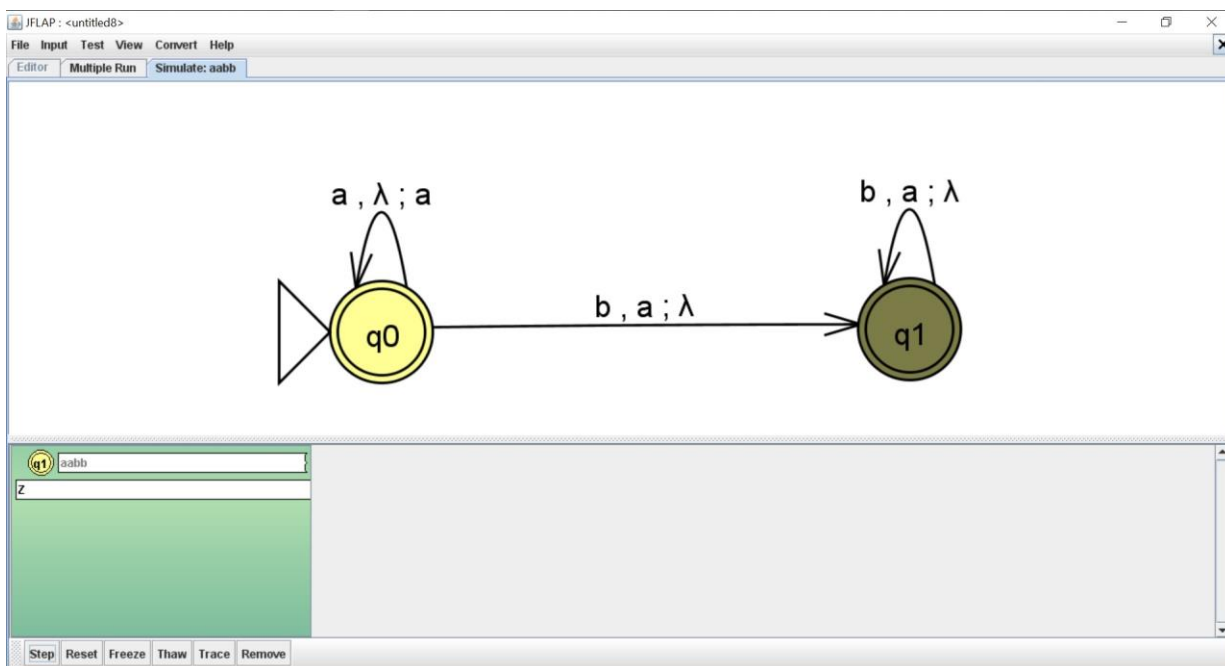
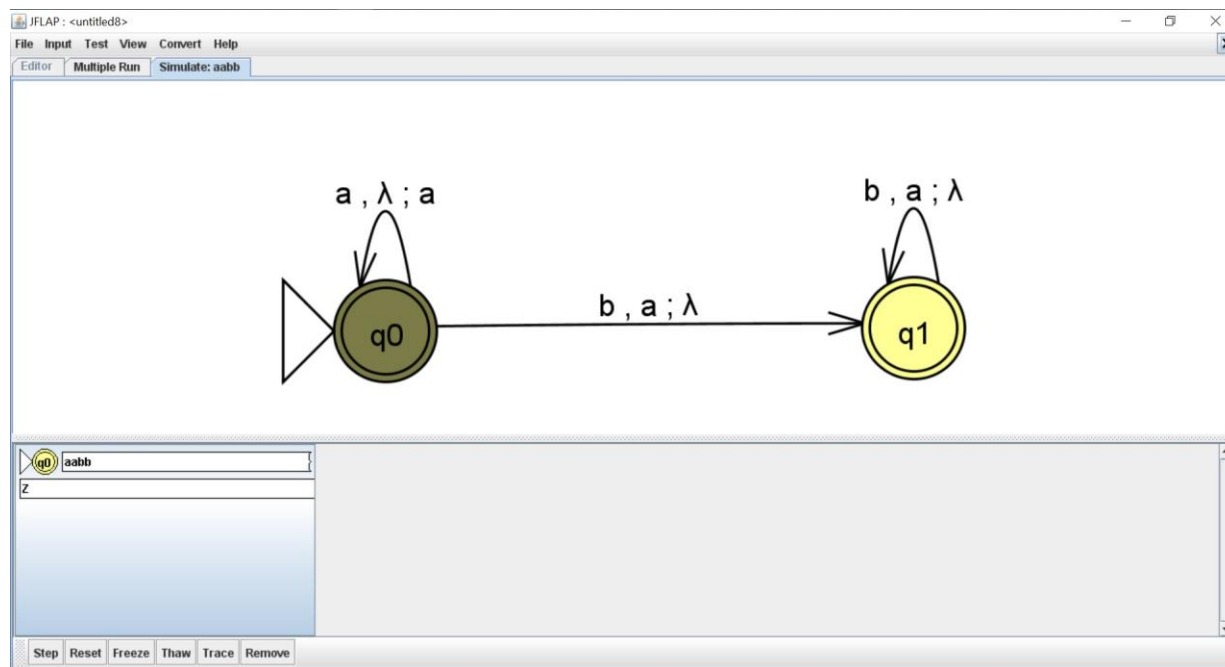
Rejecting strings: bab, abbba, baaab, aaabbbbbbb





## Trace

Accepting String: aabb



Rejecting String: bab

