# Asymmetric Encryption

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#### Functional Requirements



- → The public key K<sub>p</sub> for encryption
- → The private key K<sub>s</sub> for decryption
  - 1.  $D_{ks}(E_{kp}(m))=m$  for every pair  $(K_p, K_s)$
- 2.  $E_{kp}(m)$  is easy to compute (either polynomial or linear)
- 3.  $D_{ks}(C)$  is easy to compute (either polynomial or linear)
- 4.  $p = D_{ks}(C)$  finding m is hard without  $K_s$  (exponential)
- 5. Generating a pair  $(K_p, K_s)$  is easy to compute (polynomial)
- 6. Finding a matching key  $K_s$  for a given  $K_p$  is hard (exponential)

### RSA - Rivest, Shamir and Alderman

Key Size	1024 - 4096
Speed	~ factor of 106 cycles / operation
Mathematical Foundation	Prime number theory

### Number Theory - Prime numbers

#### **Prime Numbers**

- p is prime if 1 and p are its only divisors e.g 3, 5, 7, 11 ...
- p and q are relatively prime (a.k.a. coprime) if gcd(p,q) = 1e.g gcd(4,5) = 1
- There are infinitely many primes

#### **Euler-Fermat Theorem**

```
If n=p. q and z=(p-1).(q-1) and a such that a and n are relative primes Then a^z\equiv 1\ (mod\ n)
```

## Computational Complexity

#### Easy problems with prime numbers

- Generating a prime number p
- Addition, multiplication, exponentiation
- Inversion, solving linear equations

#### Hard problem with prime numbers

• Factoring primes e.g. given n find p and q such that n = p. q

## RSA - generating the key pair

- 1. Pick p and q two large prime numbers and calculate  $n = p \cdot q$  (see primality tests)
- 2. Compute z = (p-1).(q-1)
- 3. Pick a prime number e < z such that e and z are relative primes
- → (e,n) is the public key
- 4. Solve the linear equation  $e * d = 1 \pmod{z}$  to find d
- → (d,n) is the **private key**however p and q must be kept secret too

## RSA - encryption and decryption

Given Kp = (e, n) and Ks = (d,n)

- $\Rightarrow$  Encryption :  $E_{kp}(m) = m^e \mod n = c$
- $\rightarrow$  Decryption :  $D_{ks}(c) = c^d \mod n = m$
- $\rightarrow$  (m<sup>e</sup>)<sup>d</sup> mod n = (m<sup>d</sup>)<sup>e</sup> mod n = m

### The security of RSA

#### RSA Labs Challenge: factoring primes set

Key length	Year	Time
140	1999	I month
155	1999	4 months
160	2003	20 days
200	2005	18 months
768	2009	3 years

Challenges are no longer active

### Key length and Key n-bit security

- RSA has very long keys, 1024, 2048 and 4096 are common
- Is it more secure than asymmetric crypto with key lengths of 56, 128, 192, 256?
- → Key lengths do not compare !

RSA Key length	Effective key length
1,024	80
2,048	112
3,072	128
7,680	192
15,360	256

## Asymmetric vs Symmetric

	Symmetric	Asymmetric
pro	Fast	No key agreement
cons	Key agreement	Very slow

#### The best of both worlds

- → Use RSA to encrypt a shared key
- → Use AES to encrypt message

$$E(m) = RSA_{Kp}(k), AES_k(m)$$

### Other asymmetric cryptography schemes

#### Diffie-Hellman (precursor)

→ No Authentication but good for key-exchange

#### **EI-Gamal**

→ Good properties for homomorphic encryption

#### Elliptic Curve Cryptography (trending nowadays)

→ Fast and small keys (190 bits equivalent to 1024 bits RSA)