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PAGE NO.

DATE:

Day 24 of DSA

Task

CheckBox Bit Manipulation

☒ Decimal to Binary (Theory) → Bitwise operators

☒ Binary to Decimal

☒ Solving Bit Manipulation problems.

* Decimal to Binary :-

Approach :-

Pseudocode :-

```

int ans = 0;
int i = 0;
while (n != 0)
{
    int bit = n & 1;
    ans = (bit * pow(10, i) + ans);
    n = n >> 1;
    i++;
}

```

$n = 5$

$i = 1 \quad \text{bit} = 101 \& 1 = 1$

$\text{ans} = (1 * 10^0) + 0 = 10$

$n = 101 >> 1 = 010$

$i++$

$i = 2 \quad \text{bit} = 010 \& 1 = 0$

$\text{ans} = (0 * 10^2) + 10 = 10$

$n = 010 >> 1 = 001$

$i++$

$i = 3 \quad \text{bit} = 100 \& 1 = 1$

$\text{ans} = (1 * 10^3) + 10 = 110$

$n = 000 >> 1 = 0 \quad 100 + 1 = 101$

break

* Binary to Decimal

Approach:

```
int i = 0, ans = 0;
```

```
while (n != 0)
```

```
{
```

```
    int digit = n % 10;
```

```
    if (digit == 1)
```

```
    {
```

```
        ans = ans + pow(2, i);
```

```
    }
```

```
    n = n / 10;
```

```
    i++;
```

```
}
```

$n = 5101$

$i = 0, ans = 0$

$digit = 1015 \% 10 = 1$

$(1 == 1) \Rightarrow ans = 0 + 2^0 = 1$

$i = 1, ans = 1$

$digit = 1015 \% 10 = 0$

$(0 == 1) \times$

$i = 2, ans = 1$

$digit = 1015 \% 10 = 1$

$(1 == 1)$

$ans = 1 + (2^2)$

$= 1 + 4 = 5$

4 Another Approach For Decimal to Binary.

pseudocode:-

```
String deciToBin(int n)
```

```
{  
    String res = "";
```

```
    while (n != 0)
```

```
{
```

```
        if (n % 2 == 1) res += "1";
```

```
        else res += "0";
```

```
        n = n / 2;
```

```
}
```

```
    reverse(res);
```

```
    return res;
```

```
}
```

Time Complexity = $O(\log_2 n)$

Space complexity = $O(\log_2 n)$

Q Another Approach For Binary to Decimal :

Approach:-

```
int BinToDeci ( string x)
```

```
{
```

```
    int len = x.length();
```

```
    int p2 = 1, num = 0;
```

```
    for( i = len-1, i >= 0; i-- )
```

```
    {
```

```
        if (x[i] == '1')
```

```
            num += p2;
```

```
            p2 *= 2;
```

```
    }
```

```
    return num;
```

```
}
```

Time Complexity: $O(n)$

Space Complexity: $O(1)$

* 1's Complement

$$(13) \rightarrow (1101)_2$$

↓
Flip

$$(0010)_2$$

+1

✓

* 2's Complement

1) 1's complement

2) add 1 to it

$$\begin{array}{r} 0010 \\ + 0001 \\ \hline \end{array}$$

$$(0011)_2$$

* AND

All true \rightarrow true

, False \rightarrow False

$$\begin{array}{r} 1101 = 13 \\ 0111 = 7 \\ \hline \end{array}$$

$$0101$$

5

* OR

1. true = true

all false = false

$$\begin{array}{r} 1101 = 13 \\ 0111 = 7 \\ \hline \end{array}$$

$$(1111)$$

15

* XOR

no. of 1s \rightarrow odd \rightarrow 1

no. of 1s \rightarrow even \rightarrow 0

$$\text{XOR } 1101 = 13$$

$$0111 = 7$$

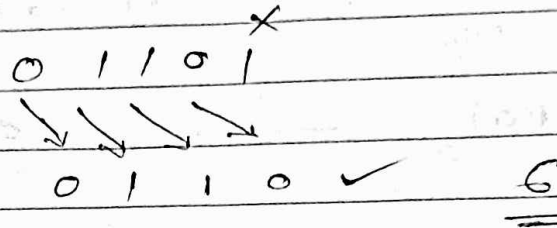
$$(1010)$$

10

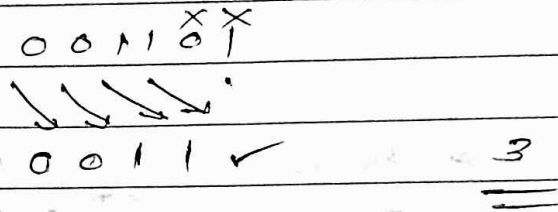
* Right shift: \gg

$$x \gg k = \frac{x}{2^k}$$

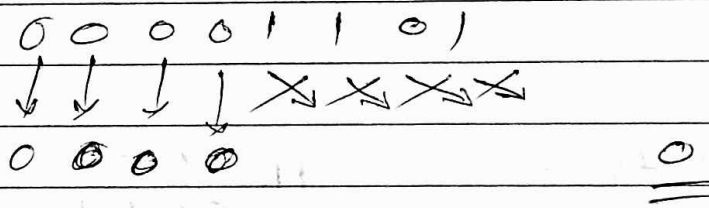
① $13 \gg 1$



② $13 \gg 2$



③ $13 \gg 4$



* largest

$$2^{31}$$

+ve sign $\rightarrow 0$ 1 1 1 1 1 1 1 1

$$(2^{31} - 1) \rightarrow \text{INT-MAX}$$

* smallest

-ve sign \rightarrow 1 0 0 0 0 0 0 0

$$-2^{31}$$

$$\underline{\underline{\text{INT-MIN}}}$$

* left shift : 22

num 4 2^k

12 221

$$\begin{array}{r} \times \\ 01101 \\ \sqrt{1111} \end{array}$$

11010 = 26

26

* Not (N)

1) Flip

2) Check -ve

$$\begin{array}{r} \swarrow \searrow \\ \text{yes} \end{array}$$

25

stop

eg.

$$\begin{array}{r} 0000 \dots 101 \\ 1111 \dots 010 \end{array}$$

0000 101

+ 1

$$\begin{array}{r} \hline 1 \dots 110 \end{array}$$

* Swap two Numbers:-

$$a = 6 \quad b = 5$$

$$(\because 5 \wedge 5 = 0)$$

$$\underline{a} = a \wedge \bar{b}$$

$$\underline{b} = \underline{a} \wedge b = (a \wedge \bar{b}) \wedge b = \underline{a}$$

$$\underline{a} = a \wedge \bar{b} = (\underline{a} \wedge \bar{b}) \wedge \bar{\underline{a}} = \underline{b}$$

* check if the i th bit is set or not.

$$N = 13 \quad i = 2$$

$$\& \quad 1 \ 1 \ 0 \ 1$$

$$1 \ll 2$$

$$0 \ 1 \ 0 \ 0$$

$$0 \ 0 \ 1 \ 0 \ 0$$

$$\leftarrow \text{O's} \quad \text{O's} \rightarrow$$

① $if ((N \& (1 \ll i)) != 0) \text{ Logic}$

{ set

}
else

{ Not set

}

Logic

② $\text{if } ((N \gg i) \& 1 == 0)$

{

Not set

{

else

{

Set

{

* set the ith bit

 $N = 9$ $i = 2$

Logic

or

1 0 0 1

0 1 0 0

1 1 0 1

=

 $\text{if } (N \& (1 \ll i))$

{

Set

{

~~else { Not set }~~

* clear the ith bit

 $N = 13$ $i = 2$

1 1 0 1

1 0 1 1

1 0 0 1

=

 $\sim(1 \ll i)$ $\sim(0100)$

(1011)

Logic = $(N \& \sim(1 \ll i))$ ✓

* Toggle the i th bit.

$$N = 13 \quad i = 2$$

$$\begin{array}{r} A \quad 1 \ 1 \ 0 \ 1 \\ \quad 0 \ 1 \ 0 \ 0 \\ \hline \quad 1 \ 0 \ 0 \ 1 \quad \checkmark \end{array}$$

$$\text{Logic} = N \wedge (1 \ll i) \quad \checkmark$$

* Remove the last set bit (rightmost)

$$N = 40$$

$$\begin{array}{r} N \quad 1 \ 0 \ 1 \ 0 \ 0 \ 0 \\ N-1 \quad 1 \ 0 \ 0 \ 1 \ 1 \ 1 \\ \hline \quad 1 \ 0 \ 0 \ 0 \ 0 \ 0 \quad \checkmark \end{array}$$

$$\text{Logic} = N \& (N-1) \quad \checkmark$$

* check if the number is a power of 2

$$N = 16 = 2^4$$

$$\text{Logic} (N \& N-1 == 0) \quad \checkmark$$

$$\begin{array}{r} 1 \ 0 \ 0 \ 0 \ 0 \\ 0 \ 1 \ 1 \ 1 \ 1 \\ \hline 0 \ 0 \ 0 \ 0 \ 0 \quad \checkmark \end{array}$$

* Count set of bits =

① int CEB(int n)

{ cnt = 0;

while (n > 1)

{ if (n % 2 == 1) cnt += 1;

n /= 2;

if (n == 1) cnt += 1;

return cnt;

odd ko liye

cnt += n % 2

n = n >> 1

↓
n / 2

* ②

cnt = 0

while (N != 0)

{

N = N & N - 1;

cnt ++;

}

Time complexity : $O(\log_2 n)$