

X. Hashing

→

1	2	1	3	2
---	---	---	---	---

key	Frequency
1	2
2	2
3	1
10	0
4	0

• Bruteforce approach:-

→ Pseudocode:-

```
int fun(numbers, a[])
```

```
{
```

```
    int count = 0;
```

```
    for (i = 0; i < n; i++)
```

```
    {
```

```
        if (arr[i] == numbers)
```

```
            count += 1;
```

```
    }
```

```
    return count;
```

```
}
```

→ Time complexity:- $O(N)$

If we have 'Q' inputs then $Q \times O(N)$

i.e. Time complexity :- $O(Q \times N)$

Suppose if $Q = 10^5$ and $N = 10^5$

then $O(10^5 \times 10^5)$

$O(10^{10})$

$O(10^8) \approx 1 \text{ second}$

i.e. Program takes 1 second.

$10^8 \rightarrow 1 \text{ sec}$

$10^{10} \rightarrow \frac{1}{10^8} \times 10^{10} \approx 10^2$ i.e., 100 seconds.
So, it's not good.

So, we use Hashing for these type of problems

→ Hashing :- Storing and Fetching.

i.e. we store some values and fetch those values whenever needed.

→ Number Hashing

1 2 1 3 2

→ Suppose we can take a map (12)

1	2
3	1
4	0
2	2
10	0
⋮	⋮
12	0

→ Hash Array

0	2	2	1	0	0	0	0	0	0	0	0	0
0	1	2	3	4	5	6	7	8	9	10	11	12

→ Code 1 -

```
#include <-->
```

```
using ---
```

```
int main() {
```

```
    int n;
```

```
    cin >> n;
```

```
    int arr[n];
```

```
    for (int i = 0; i < n; i++) {
```

```
        cin >> arr[i];
```

```
    }
```

```
// Precompute
```

```
int hash[13] = {0};
```

```
for (int i = 0; i < n; i++) {
```

```
    hash[arr[i]] += 1;
```

```
}
```

```
int k;
```

```
cin >> k;
```

```
while (k-- > 0) {
```

```
    int number;
```

```
    cin >> number;
```



```
// Fetch
cout << hash[numbers] << endl;
}
return 0;
}
```

[NOTE:-

→ Size of Hash array allowed inside
main funcn :- arr[10⁶]
↳ in case of 'int'

if size > 10⁶

then it will give segmentation fault.

→ Size of Hash array allowed outside main
funcn or Globally :- arr[10⁷]

↳ in case of 'int'

i.e. Array Hashing cannot be done if size of
array is > 10⁷.

→
character
hashing

S = "abcdabefc"

↑

string

• Bruteforce approach :-

→ Pseudocode :-

```
int fun(char ch, string s){
    int count = 0;
    for (i=0; i<n; i++){
        if (s[i] == ch)
            count++;
    }
    return count;
}
```


→ Time Complexity:- $O(8 \times N)$

g-inputs $\left\{ \begin{array}{l} a \rightarrow 2 \\ b \rightarrow 2 \\ c \rightarrow 2 \\ \vdots \\ z \rightarrow 0 \end{array} \right.$

→ In above question, it is given for lower case characters.

so,

Hash
Array

2	2	2	1	-	-	-	0	0
0	1	2	3	-	-	-	24	25
a	b	c	d	-	-	-	y	z

abcdabefc

*
}

→ Observe that here we are restricted for lowercase characters, ~~if~~ there is no condition of upper/lower case then we can make a HashArray[] of size '256'.

Hash
Array

			-	-		-	-	-	
0	1	2			65		97		255
					(A)		(a)		

↓
since,
there
are
256
characters.

→ Code:-

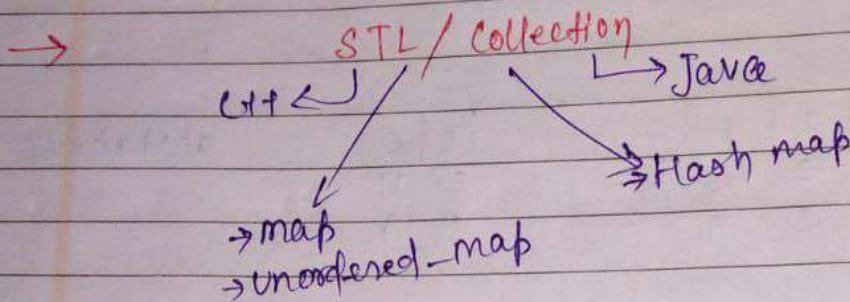
```
int main() {
    string s;
    cin >> s;
    // Pre-Compute
    int hash[26] = {0};
    for (int i = 0; i < s.size(); i++) {
        hash[s[i] - 'a']++;
    }
}
```

```
int k;
cin >> k;
while (k--) {
    char ch;
    cin >> ch;
    // fetch
    cout << hash[ch - 'a'] << endl;
} return 0; }
```


→ Code for all characters :-

```
int main() {
    string s;
    cin >> s;
    // Pre-Complete
    int hash[256] = {0};
    for (int i = 0; i < s.size(); i++) {
        hash[s[i]]++;
    }
    int q;
    cin >> q;
```

```
while (q--) {
    char ch;
    cin >> ch;
    // fetch
    cout << hash[ch] << endl;
}
return 0;
}
```



→ Map

arr = [1 | 2 | 3 | 1 | 3 | 2 | 12]

(12 → 1)
(3 → 2)
(2 → 2)
(1 → 2)

map

★ mpp[arr[i]]++

{ mpp[1] → 0
 mpp[2]++ → 1

• map <int, int>

↓ ↓
 Key Value
re-number Give frequency / count

[NOTE:-

1 2 3 1 3 2 12

In Array Hashing, here we have to make a hash array[] of size 13.

But, in map, we only need to store (1, 2, 3, 12)

Count of these

So, map takes some less memory.

→ Code:-

```
#include <-->
```

```
using ---
```

```
int main() {
```

```
int n;
```

```
cin >> n;
```

```
int arr[n];
```

```
for (int i=0; i<n; i++) {
```

```
cin >> arr[i];
```

```
}
```

```
// Pre-compute
```

```
map<int, int> mpp;
```

```
for (int i=0; i<n; i++) {
```

```
    mpp[arr[i]]++;
```

```
}
```

```
int q;
```

```
cin >> q;
```

```
while (q-- > 0) {
```

```
    int number;
```

```
    cin >> number;
```

```
    // fetch
```

```
    cout << mpp[number]
```

```
    << endl;
```

```
}
```

```
return 0;
```

```
}
```

[NOTE:-

- Values in map are stored in sorted order.

1 2 3 1 3 2 12

for (auto it : mpp) { // iterate in the map.

cout << it.first << " → " << it.second << endl;

}

O/P screen,

print
key

1 → 2
2 → 2
3 → 2
12 → 1

↓

print count/

frequency

- Time complexity of map :-

storing } $\rightarrow (\log n)$; $n \rightarrow$ no. of elements in map.
 Fetching }
 \rightarrow in all cases
 (Best, Avg., Worst)

\rightarrow unordered-map

- Values in unordered-map are not stored in sorted order.
- Time complexity of unordered-map :-

storing } $\rightarrow O(1)$ (Best Avg.)
 Fetching } $\rightarrow O(n)$; n is no. of elements in unordered-map.
 (Worst)

NOTE:-

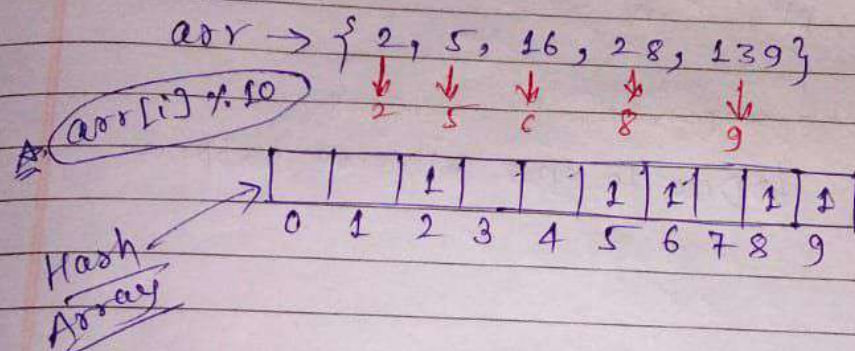
We generally use unordered-map, but if its give time limit exceeded (TLE), then we switch to map.

\rightarrow TLE occurs due to worst case. i.e. $O(n)$ of ^{unordered} map.
 And, worst case occurs due to internal collisions.

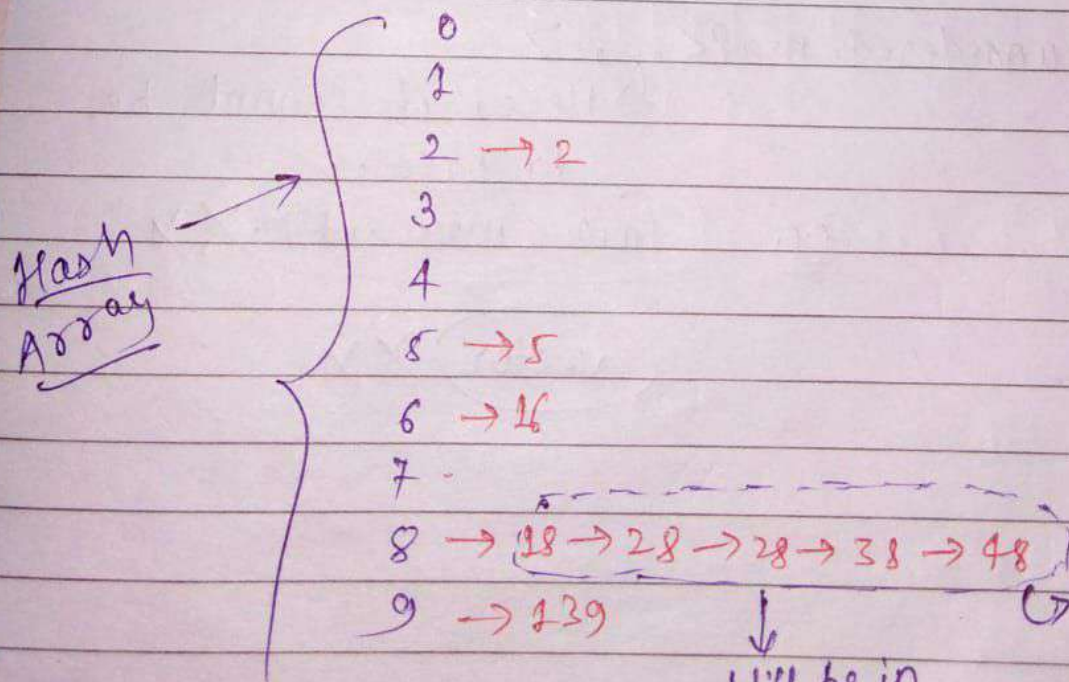
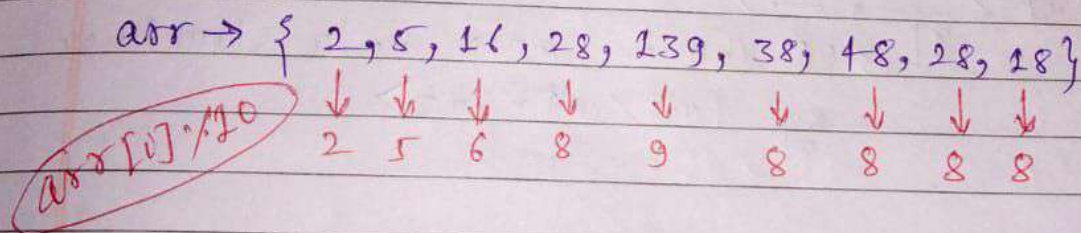
→ How Hashing is done?

- Division method
- Folding method
- mid-square method

→ Division method :-



{ suppose, we have given that array size cannot be greater than 10 }.



Will be in sorted order.

can be implemented using Linked List

→ if we have an array like

arr → {18, 28, 38, 48, ..., 1008}

0
1
2
3
4
5
6
7
8
9

All keys ends at hash
in same index.

8 → 18 → 28 → 38 → 48 → ... → 1008

↳ Due to this collision happens

[NOTE]-

→ map <_, _>

↳ It can be any data structure
e.g. Pair<int, int>

→ unordered_map <_, _>

↳ Here, it cannot be
in pairs.

e.g. Pair<int, int> ❌❌

vector ❌❌