

Caching systems: attacks and countermeasures - A Survey

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1 Introduction

Responsibility: Radhika

In Multi-Access Edge Computing (MEC) networks, servers are collocated with Base Stations (BSs) at the edge of the network, to reduce latency incurred by end users, such as Mobile Stations (MSs). These servers are not as powerful as the cloud in terms of RAM, CPU or storage. When end users request for content (images, videos etc.), the requests

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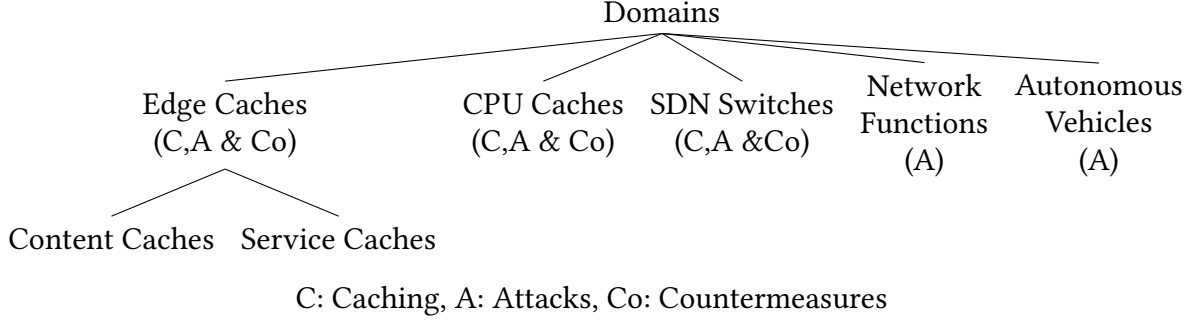


Fig. 1. Domains surveyed

arrive at the edge and are forwarded to the cloud for the content to be downloaded to the end device. Forwarding all requests to the cloud and getting a response can cause high latencies. Moreover, if there are multiple users, the links from the edge to the cloud can get overloaded, causing further delays. Therefore, each time content is downloaded, it can be cached at the edge.

Services are applications or processes that respond to requests from the user. For example, an Alternate Reality (AR) or Virtual Reality (VR) game requires an MS to send data from the device to the edge, which forwards it to a AR/VR game service running in the cloud. When data is sent such a service, a response is sent to the user and then the user can make the next move. To reduce latency, such a service and its associated database can be downloaded to the edge and cached. This differs from content caching in two ways: 1) A service request requires only a response, which is typically much smaller than content. 2) Content has to be always downloaded to the end device, whereas the service is never downloaded to the end device. 3) The edge can decide when to download a service based on an *admission policy*, while it always downloads content in response to a request.

This survey examines attack methods and corresponding countermeasures for systems that employ caching mechanisms, with particular attention to attacks and defenses relevant to edge caching systems. We do not explore any cryptographic methods, or attacks and countermeasures that are not relevant to edge caching systems. We intend to explore attack methods that may be relevant to caching systems, but the targets of which may not necessarily be caching systems.

Caches are ubiquitous in computer systems. As discussed above, there are caches at the edge, caching content and services. In addition, we explore attacks and countermeasures of two other systems that use caching: CPUs and SDN switches. To explore attacks that are relevant to edge caching, we explore Network Functions and Autonomous Vehicle Simulators. Fig. 1 illustrates the domains we survey.

CPU Caches: In the CPU microarchitecture, L1, L2 and L3 caches reside between the processor and the main memory. Typically, L1 and L2 caches are private to each core in a multiprocessor architecture, whereas L3 is common across cores. These caches are vulnerable to various types of attacks. [Expand: \[10\], \[17\]](#)

SDN Switches: In computer networks, the forwarding plane or the data plane forwards packets arriving at switches and routers according to the rules installed on them, while the control plane decides which rules to install. In Software Defined Networks (SDN), these planes are separated: the control plane is centralized in a controller, which programmatically controls the forwarding planes present in SDN switches. SDN switches store rules in Ternary Content Accessable Memory (TCAM) and Static Random Access Memory (SRAM). TCAMs are fast, but expensive, while SRAMs are slow and comparatively less expensive. Therefore, TCAMs may be considered caching flow rules. Therefore, the vulnerabilities, attack methods, and countermeasures are directly relevant.

Domain Name Systems: Domain Name Systems (DNS) translate human readable domain names into Internet Protocol (IP) addresses. A DNS cache stores the result of a name lookup at the browser level, at the Operating System (OS) level or at the level of the Internet Service Provider (ISP). Afek et al. [1] provide a history of DNS poisoning attacks.

Content Delivery Networks (CDNs): In a typical content delivery workflow, a user request is first resolved via DNS. The CDN-operated authoritative DNS service selects an appropriate edge server based on factors such as the location of the resolver, network conditions and the load. The DNS response directs the client to a specific CDN edge node, which is often physically colocated within an ISP point of presence, but remains under the operational control of the CDN. The user's HTTP(S) request is then sent directly to this edge server. If the requested content is present in the edge cache, it is served immediately; otherwise, the edge retrieves the content from the origin server (or an upstream parent cache), forwards it to the user, and caches it according to cache-control policies. Commercial agreements between ISPs and CDNs typically govern colocation,

connectivity, and traffic locality, while content placement, DNS-based request steering, and caching decisions are centrally managed by the CDN. CDNs cache objects and small state-less functions, unlike an edge cache, which can cache services too. [Read \[8\] for more on this.](#)

Note: [Caching in the World Wide Web and Information Centric Networking](#) may also be studied. See [15] and associated papers.

Network Functions: Network Functions are software modules implemented on servers that perform packet processing for reasons other than forwarding or routing, such as securing the network, optimizing traffic, etc. Intrusion Detection and Prevention Systems (IDPS), Firewalls and Wide Area Network (WAN) Optimizers are some examples. It is useful to learn the Attacks and Countermeasures employed in these systems with a view to using them in edge caches.

Autonomous Vehicles: Autonomous Vehicles (AV) aim to revolutionize driving in urban areas and highways. Simulators are used to test the safety of such vehicles. The techniques used to develop a homologation framework (the official certification process that ensures that vehicles comply with government requirements) for critical scenarios for such frameworks could be used for edge caching systems [4]. An AV planner is the part of an autonomous driving system that decides what the vehicle should do next and automated stress testing of planners using generated scenarios could be used for similar purposes for edge caching systems [11].

Our Contributions: We study the domains listed in Fig.1 (and explained above) for ideas relevant to reconnaissance, attacks, and countermeasures for edge service caching systems. Moreover, we propose taxonomies for attacks and countermeasures and summarize findings from the recent literature according to these taxonomies. We also discuss open challenges and future research directions.

We organize prior work along the following axes: attack methods and countermeasures, and domains surveyed. We examine attack methods according to their:

- method of probing (reconnaissance), if applicable or if any and what is probed (replacement algorithm, size of cache, etc.)
- conditions for probing (do they consider network jitter, link delays, link failures, etc.)

- resource consumption of the attack (low, medium, high)
- the extent of knowledge of the system required by the attackers (black box, gray box, or white box)
- relevance to edge service caching
- existence of countermeasures, if any
- verification method (tested on real networks or simulators, whether the datasets used, if any, are realistic)
- limitations (assumptions on the system by the attackers)

We examine countermeasures according to their:

- method (heuristic, AI etc.)
- impact on performance
- impact on resource consumption,
- tested on real networks or simulators, the datasets used, if any
- relevance to edge service caching
- limitations (are there strong assumptions, are the datasets realistic, etc.)

1.1 Comparison with existing surveys

Responsibility: TBD

1.2 Why this survey is different

Responsibility: TBD

Add comparison with existing surveys on caching systems, Low-rate attacks in SDN, Blackbox attacks etc.

Our contributions:

2 Background and system model

Responsibility: TBD

The domains surveyed in this paper are shown in Fig. 1.

Explain 1) The system model for each domain 2) Definitions of common terms and notations. In the rest of the paper you don't need to define any term again.

2.1 SDN Switches

Responsibility: Robin.

Write briefly about how an SDN works (Switches, Controller, Openflow, TCAM, caching etc.).

Software Defined Networking (SDN) introduces a paradigm shift in network design by decoupling the control plane from the data plane. In traditional networks, forwarding devices such as switches and routers independently make packet forwarding decisions based on locally stored routing and forwarding logic. This tightly coupled architecture makes network management complex and inflexible.

In SDN, the control plane is logically centralized in an SDN controller, while the data plane consists of multiple SDN switches that perform packet forwarding. The controller acts as the network's decision-making entity, maintaining a global view of the network and dynamically installing forwarding rules on switches using standardized southbound interfaces such as OpenFlow.

SDN switches maintain flow tables that store flow rules defining how packets should be handled. These rules typically match on packet header fields such as source and destination IP addresses, transport-layer ports, and protocol identifiers. For performance reasons, flow rules are commonly stored in Ternary Content Addressable Memory (TCAM), which enables fast parallel lookups. However, TCAM is expensive and has limited capacity, making flow table space a scarce resource.

When a packet arrives at a switch, the switch performs a lookup in its flow table. If a matching rule is found, the packet is forwarded accordingly at line rate. If no matching rule exists, the packet (or its header) is forwarded to the controller via a Packet-In message. The controller then determines the appropriate forwarding action and may install a new flow rule in the switch. This behavior effectively treats the flow table as a cache, where cache hits result in fast forwarding and cache misses trigger interaction with the controller.

Due to the limited capacity of TCAM and the reactive nature of rule installation, SDN switches are vulnerable to attacks that aim to exhaust flow table space or overload the control channel. These characteristics make SDN switches a relevant domain for studying caching-related attacks and countermeasures.

2.2 Network Functions

Responsibility: Robin What are Network Functions?

2.3 Relevance to edge service caching

Explain why attacks and countermeasures in domains other than caching are relevant. An example is: The cache itself can be attacked. Attacks can be triggered such that the links from edges to the cloud are saturated. If a TCAM in an SDN switch is considered to be a cache, a cache miss for a packet would result in traffic to the controller until a flow rule is downloaded to the cache.

3 Attack taxonomy

Refer Fig.3 and 4 for the attack taxonomy.

Need to decide which of the ACAs given the figure are low-resource and which are high-resource. Currently, all are grouped under low-resource.

Need to add more survey papers related to attacks on caches (pollution attacks etc.).

4 Attacks and countermeasures

4.1 Bayesian Optimization (High resource attacks)

4.1.1 Introduction to Bayesian Optimization. Responsibility: Vasudeva

This should be 2 pages long and must describe the basic idea with references. Use mathematical notation wherever required.

Bayesian Optimization (BO) is one of the powerful strategies for finding extrema of objective functions. Generally, an optimization problem is also formulated as maximizing function $f(x)$ over a domain $A \subset \mathbb{R}^d$, i.e., $\max_{x \in A \subset \mathbb{R}^d} f(x)$. Particularly, a few characteristics of $f(x)$ make BO the appropriate choice of optimization strategy such as costly to evaluate, unknown/no mathematical representation, no access to derivatives, non-convex, limited observations, noisy observations, etc. [5].

BO is a strategy combining ideas from Bayesian inference (Bayes' theorem), Surrogate modeling (mean function, covariance/kernel function), Sequential decision-making (acquisition functions like expected improvement).

BO is grounded in Bayes' theorem. Bayes' theorem broadly states *posterior* probability of a model (or hypothesis) f_* given dataset (or observations) D equals *likelihood* of f_* given D multiplied by *prior* probability of f_* and divided by the probability of D , i.e.,

$$p(f_* | D) = \frac{p(D | f_*) p(f_*)}{p(D)}$$

since $p(D)$ does not depend on f_* , it is also represented as $p(f_* | D) \propto p(D | f_*) p(f_*)$ in BO literature. Along the lines of Bayes' theorem, BO assumes a prior belief for the unknown objective function f , uses the initial samples in D to fit f_* (referred to as *posterior*) as a candidate for f . In order to efficiently sample more data points, BO uses acquisition functions to validate the search space of x_t to acquire a new sample, evaluate it for y_t using f , and refit f_* , making the new f_* (*posterior*) closer to f . BO continues to sample more data until a stopping criteria. The f is modeled as a random function drawn from a probability distribution over functions. The prior belief in BO represents the space of possible objective functions and is an inductive bias encoding assumptions about the nature of f , such as smoothness, continuity, noise, etc, that make some possible functions more plausible. The prior belief also includes the choice of kernel function and its hyper-parameters, the choice of acquisition function and its hyper-parameters, the choice of other hyper-parameters, among others. For this reason, the prior affects the sampling efficiency in acquisition functions, and overall convergence in BO and thus the choice of prior is crucial.

does GP assume lipschitz-continuous? BO conditions on a dataset D with initial samples, where D is a set of initial i pairs of x, y and $x \in \mathbb{R}^d$, $y \in \mathbb{R}$, i.e., $D = \{x_{1:i}, y_{1:i}\}$ and f evaluates x to y . A prior belief, also referred to as a surrogate model, is assumed based on characteristics of the objective function. In BO literature, the standard surrogate model for f is a Gaussian Process (GP) defined in the Eq. below. A GP is intuitively understood as a distribution over functions; that is, each sample drawn from a GP corresponds to a mathematical function. The mean function and the kernel function (also referred to as the pair-wise covariance function of samples in D) parameterize the characteristics of this distribution. If $x_t \in \mathbb{R}^d$, the GP is d -dimensional, and a 0-dimensional GP can be understood as a Gaussian distribution. See an example BO run in the figure ref2. Examples

of other surrogate models are Student-t Process, etc.

$$f_* \sim \mathcal{GP}(\mu(x), k(x, x'))$$

where $\mu(x)$ represents mean function and $k(x, x')$ is kernel function of the GP. Applying Bayes' theorem on GP leads to inference rules for mean function ($\mu_t(x)$) and kernel function (σ_t^2) after t samples are appended to the initial i in D . Each refit as shown in Eq. (1) increases the probability that a sample f_* , drawn from the GP, is close to the actual f . Generally, the $\mu = 0$ and the $k(x, x) = 1$ at the start. A few choices for kernel function include Squared Exponential (RBF) kernel, Matérn kernel, Rational Quadratic kernel, etc.

$$f_* = p(f(x) \mid \mathcal{D}_t) \quad (1)$$

The kernel matrix K , calculated using D , and y represent a fit of f_* . At the start f_* is fit using the initial i samples in D (the i is a hyper-parameter and is problem-specific). The f_* is refit with the new sample acquired after a few candidates are collected, validated by an acquisition function, evaluated using the actual objective function f , and finally, appended to the D , see Eqs. (2)-(5).

$$\mu_t(x_{t_c}) = \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{y}_{1:i+t-1} \quad (2.1)$$

Where $\mathbf{k}_*(x_t) \in \mathbb{R}^{i+t-1}$ represents the covariance vector of x_t with the other $i + t - 1$ samples in D , $\mathbf{K}_{t-1} \in \mathbb{R}^{i+t-1 \times i+t-1}$ represents the pair-wise covariance matrix of the other $i + t - 1$ samples, and $\mathbf{y}_{1:i+t-1} \in \mathbb{R}^{i+t-1}$ represents the y values vector of all $i + t - 1$ samples in D . Intuitively, the pairwise covariance of the $x_1 : i + t - 1$ data points is projected onto the y space and weighted by the covariance vector of the new point x_t .

$$\sigma_t^2(x_t) = k(x_t, x_t) - \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{k}_*(x_t) \quad (2.2)$$

Similar to Eq. 2.1 Eq. (2.1), the second-term represents information gain based on kernel function values subtracted from variance $k(x, x)$.

Using the Eqs. (2.1) and (2.2), the mean function and kernel function inferred based on the *posterior* f_* for a candidate x_{t_c} of acquisition function α . Generally, candidates are sampled across the search space of x , and gradient based optimization techniques like

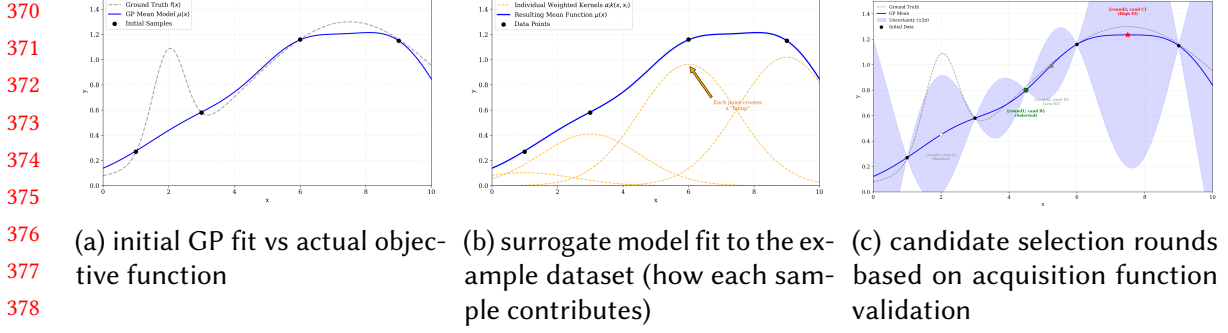


Fig. 2. AI summary (to add acquisition functions used, their equations from google docs, explanation of initial dataset used): Bayesian optimization workflow on an example dataset using a Gaussian Process surrogate. Starting from an initial set of sampled inputs, the surrogate model captures predictive mean and uncertainty over the objective function (Figure 2a). As observations from the example dataset are incorporated, the model progressively approximates the underlying function (Figure 2b). An acquisition function is then evaluated over the input space to balance exploration and exploitation, and the maximizer of this acquisition determines the next sample to be evaluated, shown as the selected candidate in Figure 2c.

L-BFGS-B are used to estimate the maxima as x_t .

$$x_{t+1} = \arg \max_{x \in \mathcal{X}} \alpha(x \mid \mathcal{D}_t) \quad (3)$$

The Eq. 3 is the acquisition function α to efficiently select the next sample x_t to evaluate f for y_t . An ideal choice of α should efficiently balance exploration and exploitation. A few choices for the acquisition function include Expected Improvement, Thompson Sampling, Probability of Improvement, Upper Confidence Bound, etc.

$$y_{t+1} = f(x_{t+1}) \quad (4)$$

In some cases, a Gaussian noise term $\varepsilon \sim \mathcal{N}(0, \sigma^2)$ is added to function evaluation as $f(x) + \varepsilon$ to make BO robust and not overfit.

$$\mathcal{D}_{t+1} = \mathcal{D}_t \cup \{(x_{t+1}, y_{t+1})\} \quad (5)$$

Until a problem-specific goal or stopping-criteria is achieved, more samples are acquired and the GP is refit following Eqs. (2) - (5).

Applications of BO include hyper-parameter tuning in neural networks, kriging in mining, etc.

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4.1.2 Attacks and Countermeasures.

High Resource Attack: A high resource attack also referred to as resource exhaustion attack is a type of attack in which an adversary aims to intentionally diminish the target's defined computational resources, including cpu cycles, memory allocation, bandwidth, or disk space leading the system or service unavailable to authentic users. Such attacks aim to overwhelm the target by provoking disproportionate demand for these resources, frequently recurring actions that exploit vulnerabilities in resource management. Salient characteristics of resource high resource attacks comprise their non-destructive nature to data integrity and confidentiality hence their primary goal is to disrupt the service availability, degrade performance without altering or stealing information.

The core principle behind resource exhaustion attacks resource asymmetry leads attackers to incur minimal cost while forcing the target system to perform disproportionately expensive operations. This imbalance gives adversaries an advantage to exploit and induce significant resource consumption using relatively low-rate or low-volume inputs. The on-demand resource allocation nature of modern systems particularly makes them vulnerable to resource exhaustion attack. Dynamically provisioned resources, such as buffers, threads, and connections, can accumulate when malicious requests exceed reclamation capacity, leading to a denial of service. Systems that lack adequate validation, rate limiting, or allocation controls are especially susceptible, enabling rapid exhaustion even at moderate attack rates.

These attacks are aggravated by persistence and amplification mechanisms. Persistence prolongs resource occupancy, while amplification enables low-cost inputs to trigger disproportionately expensive system responses.

Depending on the targeted resource, resource exhaustion attacks can be broadly classified as CPU exhaustion, memory exhaustion, bandwidth exhaustion, connection exhaustion, and disk exhaustion attacks.

Bayesian Optimization as High Resource Attack: Bayesian optimization has become a prominent technique in adversarial machine learning, particularly when dealing with the more realistic black-box setting that requires an attacker to find an adversarial perturbation without any knowledge of the architecture, parameters, or training data of the target model. In such cases, information about the model can only be obtained through queries,

i.e. supplying an input to the model and receiving the corresponding output. In these situations, an attacker can treat interactions with the target model as a sequential optimization problem, in which like the model's response latency, predicted labels, confidence scores and resource utilization, serve as noisy measurements of some unknown underlying goal the attacker is trying to achieve.

When applied to high-resource attacks, BO enables an adversary to efficiently identify inputs that result in maximal computational cost during model inference. By approximating the behavior of the target model with a probabilistic surrogate and using acquisition functions to guide the selection of the query, Bayesian optimization allows an attacker to progressively identify inputs that induce execution paths of the worst-case.

Triggering such inputs in a loop may result in extensive resource exhaustion, leading to system degradation rendering the failure of service or denial of service. This trait of BO outperforms random search and heuristic approaches and makes it particularly effective for these attacks by achieving greater impact in fewer iterations. Despite their effectiveness, these attacks are constrained by the presence of resource control measures, including rate limiting, timeouts, and input validation. Furthermore, targeting well-provisioned systems typically requires distributed attack infrastructures, increasing the risk of detection. Hence, some mitigation strategies like limiting resource quotas, load balancing, anomaly detection and adaptive throttling, are critical to minimizing system vulnerability.

Countermeasures: Despite the fact that their effectiveness, these attacks are constrained by the presence of resource control measures, including rate limiting, timeouts, and input validation. Moreover, targeting well-resourced systems usually requires distributed attack infrastructures thereby increasing the risk of detection.

Responsibility: Amit Singh

Fill Tables 1 and 2.

Write a few paragraphs on attacks and countermeasures here summarising each paper along the columns of Tables 1 and 2. Add relevant information, if any, not mentioned in the table. For example, how the probing and later the attack happens, whether there are trade-offs, etc. State the dominant features across the papers that you have studied that you want to highlight.

Table 1. Comparison of Probing-Based Analysis Techniques

<i>Paper (Domain)</i>	<i>Method</i>	<i>Conditions</i>	<i>Resources</i>	<i>System Knowledge</i>	<i>Countermeasures?</i>	<i>Verification Method</i>	<i>Limitations</i>	<i>Relevance</i>	<i>Notes</i>
Dummy row remove this[3](NFs)	Packets - sent	Network jitter	High CPU, mod- erate memory	Black- box	Yes	Synthetic datasets	Limited scalabil- ity	High	Seminal

Table 2. Comparison of Countermeasures

<i>Paper (Domain)</i>	<i>Method</i>	<i>Performance Impact</i>	<i>Resource Overhead</i>	<i>Test Method</i>	<i>Limitations</i>	<i>Relevance</i>	<i>Remarks</i>
Dummy row remove this[2](NFs)	BO -	Latency increase under peak load	Moderate CPU and memory overhead	Real work	net- Reduced effective- ness under adaptive attacks	High	Dataset not shared

While you do this, if there is something that you find applicable to the section on open challenges and future research directions, add that there.

4.2 Low Rate Flow-table overflow attacks (Low Resource)

Responsibility: Robin

Copy Tables.1 and 2 here.

Write a few paragraphs on attacks and countermeasures here summarising each paper along the columns of Tables 1 and 2. Add relevant information, if any, not mentioned in the table. For example, how the probing and later the attack happens, whether there are trade-offs, etc. State the dominant features across the papers that you have studied that you want to highlight.

While you do this, if there is something that you find applicable to the section on open challenges and future research directions, add that there.

5 Open Challenges and Future Research Directions

5.1 Bayesian Optimization

5.1.1 Open Challenges.

5.1.2 Future Directions.

5.2 Low Rate Flow-table overflow attacks

5.2.1 Open Challenges.

5.2.2 Future Directions.

5.3

6 Conclusions

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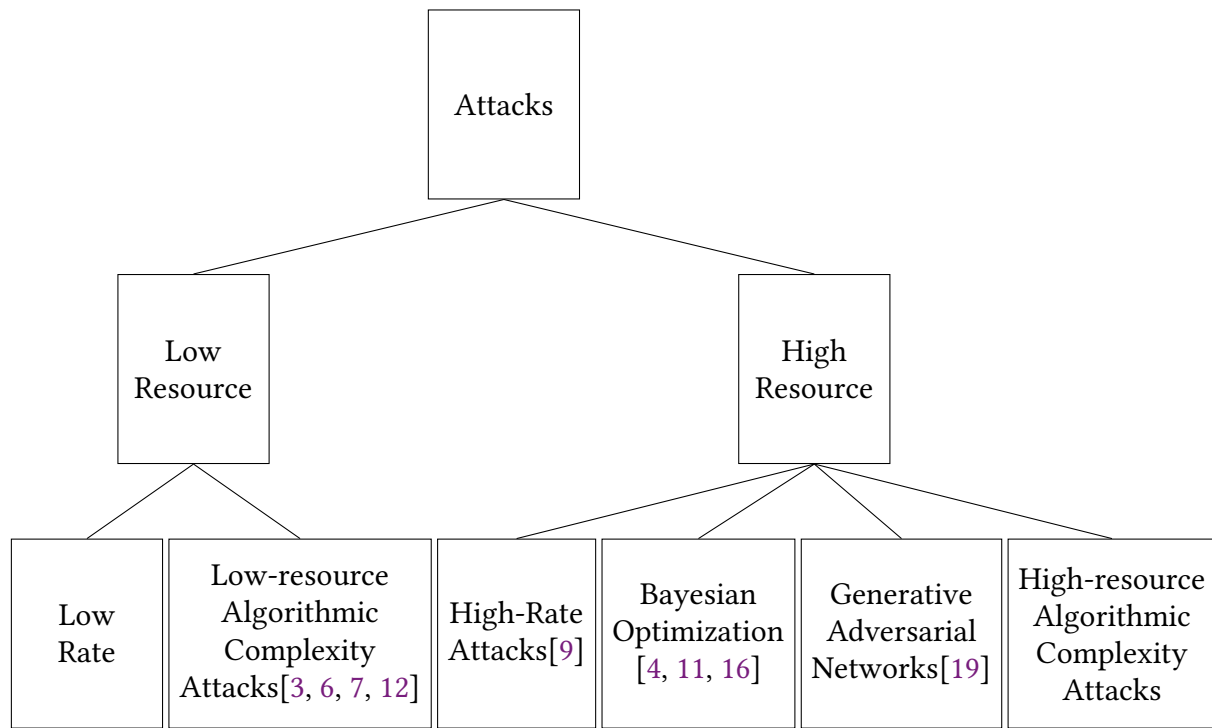


Fig. 3. Attack taxonomy

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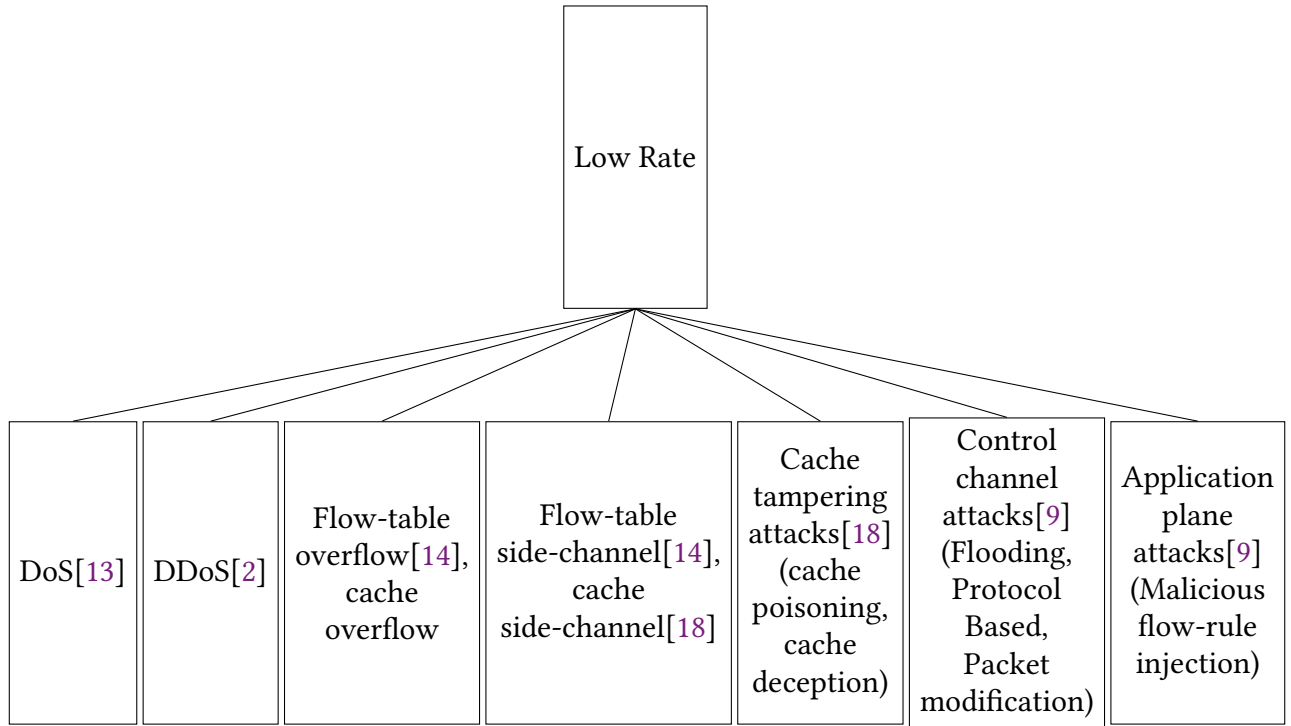


Fig. 4. Taxonomy of low-rate attacks

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