

1 **Caching systems: attacks and**
2
3 **countermeasures - A Survey**
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6 Write the abstract here.
7

8 CCS Concepts: • **Do Not Use This Code → Generate the Correct Terms for Your Paper;**
9 *Generate the Correct Terms for Your Paper; Generate the Correct Terms for Your Paper; Generate*
10 *the Correct Terms for Your Paper.*
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13 Paper
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15 **ACM Reference Format:**

16 . 2025. Caching systems: attacks and countermeasures - A Survey. In *Proceedings of Make sure*
17 *to enter the correct conference title from your rights confirmation email (Conference acronym 'XX).*
18 ACM, New York, NY, USA, 16 pages. <https://doi.org/XXXXXXX.XXXXXXX>
19

20
21 **1 Introduction**

22 **Responsibility:** Radhika
23

24 In Multi-Access Edge Computing (MEC) networks, servers are collocated with Base
25 Stations (BSs) at the edge of the network, to reduce latency incurred by end users, such as
26 Mobile Stations (MSs). These servers are not as powerful as the cloud in terms of RAM,
27 CPU or storage. When end users request for content (images, videos etc.), the requests

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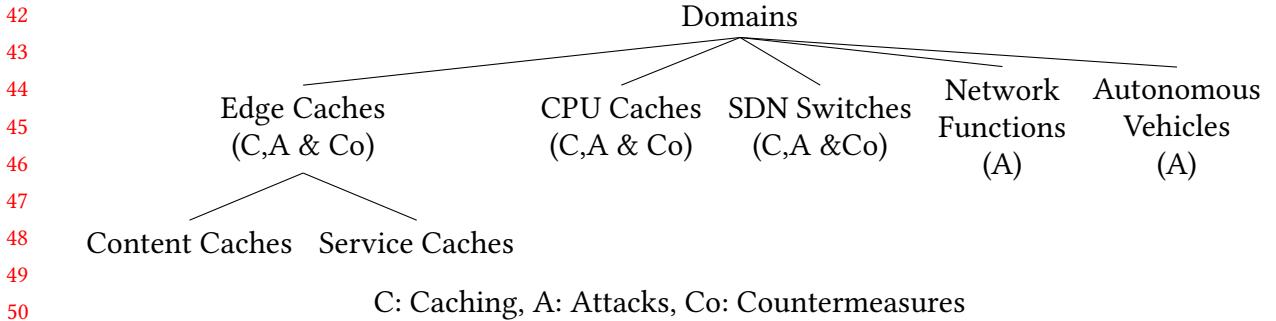


Fig. 1. Domains surveyed

arrive at the edge and are forwarded to the cloud for the content to be downloaded to the end device. Forwarding all requests to the cloud and getting a response can cause high latencies. Moreover, if there are multiple users, the links from the edge to the cloud can get overloaded, causing further delays. Therefore, each time content is downloaded, it can be cached at the edge.

Services are applications or processes that respond to requests from the user. For example, an Alternate Reality (AR) or Virtual Reality (VR) game requires an MS to send data from the device to the edge, which forwards it to a AR/VR game service running in the cloud. When data is sent such a service, a response is sent to the user and then the user can make the next move. To reduce latency, such a service and its associated database can be downloaded to the edge and cached. This differs from content caching in two ways: 1) A service request requires only a response, which is typically much smaller than content. 2) Content has to be always downloaded to the end device, whereas the service is never downloaded to the end device. 3) The edge can decide when to download a service based on an *admission policy*, while it always downloads content in response to a request.

This survey examines attack methods and corresponding countermeasures for systems that employ caching mechanisms, with particular attention to attacks and defenses relevant to edge caching systems. We do not explore any cryptographic methods, or attacks and countermeasures that are not relevant to edge caching systems. We intend to explore attack methods that may be relevant to caching systems, but the targets of which may not necessarily be caching systems.

Caches are ubiquitous in computer systems. As discussed above, there are caches at the edge, caching content and services. In addition, we explore attacks and countermeasures of two other systems that use caching: CPUs and SDN switches. To explore attacks that are relevant to edge caching, we explore Network Functions and Autonomous Vehicle Simulators. Fig. 1 illustrates the domains we survey.

CPU Caches: In the CPU microarchitecture, L1, L2 and L3 caches reside between the processor and the main memory. Typically, L1 and L2 caches are private to each core in a multiprocessor architecture, whereas L3 is common across cores. These caches are vulnerable to various types of attacks. [Expand: \[10\], \[17\]](#)

SDN Switches: In computer networks, the forwarding plane or the data plane forwards packets arriving at switches and routers according to the rules installed on them, while the control plane decides which rules to install. In Software Defined Networks (SDN), these planes are separated: the control plane is centralized in a controller, which programmatically controls the forwarding planes present in SDN switches. SDN switches store rules in Ternary Content Accessible Memory (TCAM) and Static Random Access Memory (SRAM). TCAMs are fast, but expensive, while SRAMs are slow and comparatively less expensive. Therefore, TCAMs may be considered caching flow rules. Therefore, the vulnerabilities, attack methods, and countermeasures are directly relevant.

Domain Name Systems: Domain Name Systems (DNS) translate human readable domain names into Internet Protocol (IP) addresses. A DNS cache stores the result of a name lookup at the browser level, at the Operating System (OS) level or at the level of the Internet Service Provider (ISP). Afek et al. [1] provide a history of DNS poisoning attacks.

Content Delivery Networks (CDNs): In a typical content delivery workflow, a user request is first resolved via DNS. The CDN-operated authoritative DNS service selects an appropriate edge server based on factors such as the location of the resolver, network conditions and the load. The DNS response directs the client to a specific CDN edge node, which is often physically colocated within an ISP point of presence, but remains under the operational control of the CDN. The user's HTTP(S) request is then sent directly to this edge server. If the requested content is present in the edge cache, it is served immediately; otherwise, the edge retrieves the content from the origin server (or an upstream parent cache), forwards it to the user, and caches it according to cache-control policies. Commercial agreements between ISPs and CDNs typically govern colocation,

124 connectivity, and traffic locality, while content placement, DNS-based request steering,
125 and caching decisions are centrally managed by the CDN. CDNs cache objects and small
126 state-less functions, unlike an edge cache, which can cache services too. [Read \[8\] for more](#)
127 [on this](#).

128

129 Note: Caching in the World Wide Web and Information Centric Networking may also
130 be studied. See [15] and associated papers.

131 **Network Functions:** Network Functions are software modules implemented on servers
132 that perform packet processing for reasons other than forwarding or routing, such as
133 securing the network, optimizing traffic, etc. Intrusion Detection and Prevention Systems
134 (IDPS), Firewalls and Wide Area Network (WAN) Optimizers are some examples. It is
135 useful to learn the Attacks and Countermeasures employed in these systems with a view
136 to using them in edge caches.

137

138 **Autonomous Vehicles:** Autonomous Vehicles (AV) aim to revolutionize driving in
139 urban areas and highways. Simulators are used to test the safety of such vehicles. The
140 techniques used to develop a homologation framework (the official certification process
141 that ensures that vehicles comply with government requirements) for critical scenarios
142 for such frameworks could be used for edge caching systems [4]. An AV planner is the
143 part of an autonomous driving system that decides what the vehicle should do next and
144 automated stress testing of planners using generated scenarios could be used for similar
145 purposes for edge caching systems [11].

146 **Our Contributions:** We study the domains listed in Fig.1 (and explained above) for
147 ideas relevant to reconnaissance, attacks, and countermeasures for edge service caching
148 systems. Moreover, we propose taxonomies for attacks and countermeasures and summa-
149 rize findings from the recent literature according to these taxonomies. We also discuss
150 open challenges and future research directions.

151 We organize prior work along the following axes: attack methods and countermeasures,
152 and domains surveyed. We examine attack methods according to their:
153

- 154
- 155 • method of probing (reconnaissance), if applicable or if any and what is probed
156 (replacement algorithm, size of cache, etc.)
- 157 • conditions for probing (do they consider network jitter, link delays, link failures,
158 etc.)

- 165 • resource consumption of the attack (low, medium, high)
- 166 • the extent of knowledge of the system required by the attackers (black box, gray
- 167 box, or white box)
- 168 • relevance to edge service caching
- 169 • existence of countermeasures, if any
- 170 • verification method (tested on real networks or simulators, whether the datasets
- 171 used, if any, are realistic)
- 172 • limitations (assumptions on the system by the attackers)

173
174 We examine countermeasures according to their:

- 175 • method (heuristic, AI etc.)
- 176 • impact on performance
- 177 • impact on resource consumption,
- 178 • tested on real networks or simulators, the datasets used, if any
- 179 • relevance to edge service caching
- 180 • limitations (are there strong assumptions, are the datasets realistic, etc.)

186 1.1 Comparison with existing surveys

187 Responsibility: TBD

190 1.2 Why this survey is different

191 Responsibility: TBD

192 Add comparison with existing surveys on caching systems, Low-rate attacks in SDN, Blackbox
193 attacks etc.

194 Our contributions:

198 2 Background and system model

199 Responsibility: TBD

200 The domains surveyed in this paper are shown in Fig. 1.

201 Explain 1) The system model for each domain 2) Definitions of common terms and
202 notations. In the rest of the paper you don't need to define any term again.

206 **2.1 SDN Switches**

207 Responsibility: Robin.

209 Write briefly about how an SDN works (Switches, Controller, Openflow, TCAM, caching
210 etc.).

211

212 **2.2 Network Functions**

214 Responsibility: Robin What are Network Functions?

215

216 **2.3 Relevance to edge service caching**

218 Explain why attacks and countermeasures in domains other than caching are relevant.
219 An example is: The cache itself can be attacked. Attacks can be triggered such that the
220 links from edges to the cloud are saturated. If a TCAM in an SDN switch is considered to
221 be a cache, a cache miss for a packet would result in traffic to the controller until a flow
222 rule is downloaded to the cache.

224

225

226 **3 Attack taxonomy**

227 Refer Fig.3 and 4 for the attack taxonomy.

228

229 Need to decide which of the ACAs given the figure are low-resource and which are
230 high-resource. Currently, all are grouped under low-resource.

231

232 Need to add more survey papers related to attacks on caches (pollution attacks etc.).

233

234

235 **4 Attacks and countermeasures**

236 **4.1 Bayesian Optimization (High resource attacks)**

237 *4.1.1 Introduction to Bayesian Optimization.* Responsibility: Vasudeva

238

239 This should be 2 pages long and must describe the basic idea with references. Use
240 mathematical notation wherever required.

241

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243

244

245

246 Bayesian Optimization (BO) is one of the powerful strategies for finding extrema of
247 objective functions. Generally, an optimization problem is also formulated as maximizing
248 function $f(x)$ over a domain $A \subset \mathbb{R}^d$, i.e., $\max_{x \in A \subset \mathbb{R}^d} f(x)$. Particularly, a few character-
249 stics of $f(x)$ make BO the appropriate choice of optimization strategy such as costly to
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247 evaluate, unknown/no mathematical representation, no access to derivatives, non-convex,
248 limited observations, noisy observations, etc. [5].

249 BO is a strategy combining ideas from Bayesian inference (Bayes' theorem), Surro-
250 gate modeling (mean function, covariance/kernel function), Sequential decision-making
251 (acquisition functions like expected improvement).

253 BO is grounded in Bayes' theorem. Bayes' theorem broadly states *posterior* probability
254 of a model (or hypothesis) f_* given dataset (or observations) D equals *likelihood* of f_*
255 given D multiplied by *prior* probability of f_* and divided by the probability of D , i.e.,
256

$$\frac{p(f_* | D)}{p(D)} = \frac{p(D | f_*) p(f_*)}{p(D)}$$

259 since $p(D)$ does not depend on f_* , it is also represented as $p(f_* | D) \propto p(D | f_*) p(f_*)$
260 in BO literature. Along the lines of Bayes' theorem, BO assumes a prior belief for the
261 unknown objective function f , uses the initial samples in D to fit f_* (referred to as
262 *posterior*) as a candidate for f . In order to efficiently sample more data points, BO uses
263 acquisition functions to validate the search space of x_t to acquire a new sample, evaluate
264 it for y_t using f , and refit f_* , making the new f_* (*posterior*) closer to f . BO continues
265 to sample more data until a stopping criteria. The f is modeled as a random function
266 drawn from a probability distribution over functions. The prior belief in BO represents
267 the space of possible objective functions and is an inductive bias encoding assumptions
268 about the nature of f , such as smoothness, continuity, noise, etc, that make some possible
269 functions more plausible. The prior belief also includes the choice of kernel function and
270 its hyper-parameters, the choice of acquisition function and its hyper-parameters, the
271 choice of other hyper-parameters, among others. For this reason, the prior affects the
272 sampling efficiency in acquisition functions, and overall convergence in BO and thus the
273 choice of prior is crucial.

278 does GP assume lipschitz-continuous? BO conditions on a dataset D with initial samples,
279 where D is a set of initial i pairs of x, y and $x \in \mathbb{R}^d, y \in \mathbb{R}$, i.e., $D = \{x_{1:i}, y_{1:i}\}$ and f
280 evaluates x to y . A prior belief, also referred to as a surrogate model, is assumed based on
281 characteristics of the objective function. In BO literature, the standard surrogate model
282 for f is a Gaussian Process (GP) defined in the Eq. below. A GP is intuitively understood
283 as a distribution over functions; that is, each sample drawn from a GP corresponds to
284

288 a mathematical function. The mean function and the kernel function (also referred to
 289 as the pair-wise covariance function of samples in D) parameterize the characteristics
 290 of this distribution. If $x_t \in \mathbb{R}^d$, the GP is d-dimensional, and a 0-dimensional GP can be
 291 understood as a Gaussian distribution. See an example BO run in the figure ref2. Examples
 292 of other surrogate models are Student-t Process, etc.
 293

$$294 \quad f_* \sim \mathcal{GP}(\mu(x), k(x, x'))$$

296 where $\mu(x)$ represents mean function and $k(x, x')$ is kernel function of the GP. Applying
 297 Bayes' theorem on GP leads to inference rules for mean function ($\mu_t(x)$) and kernel
 298 function (σ_t^2) after t samples are appended to the initial i in D . Each refit as shown in Eq.
 299 (1) increases the probability that a sample f_* , drawn from the GP, is close to the actual
 300 f . Generally, the $\mu = 0$ and the $k(x, x) = 1$ at the start. A few choices for kernel function
 301 include Squared Exponential (RBF) kernel, Matérn kernel, Rational Quadratic kernel, etc.
 302
 303

$$304 \quad f_* = p(f(x) | \mathcal{D}_t) \quad (1)$$

306 The kernel matrix K , calculated using D , and y represent a fit of f_* . At the start f_* is fit
 307 using the initial i samples in D (the i is a hyper-parameter and is problem-specific). The
 308 f_* is refit with the new sample acquired after a few candidates are collected, validated
 309 by an acquisition function, evaluated using the actual objective function f , and finally,
 310 appended to the D , see Eqs. (2)-(5).

$$313 \quad \mu_t(x_{t_c}) = \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{y}_{1:i+t-1} \quad (2.1)$$

314 Where $\mathbf{k}_*(x_t) \in \mathbb{R}^{i+t-1}$ represents the covariance vector of x_t with the other $i + t - 1$
 315 samples in D , $\mathbf{K}_{t-1} \in \mathbb{R}^{i+t-1 \times i+t-1}$ represents the pair-wise covariance matrix of the other
 316 $i + t - 1$ samples, and $\mathbf{y}_{1:i+t-1} \in \mathbb{R}^{i+t-1}$ represents the y values vector of all $i + t - 1$ samples
 317 in D . Intuitively, the pairwise covariance of the $x_1 : i + t - 1$ data points is projected onto
 318 the y space and weighted by the covariance vector of the new point x_t .
 319
 320

$$321 \quad \sigma_t^2(x_t) = k(x_t, x_t) - \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{k}_*(x_t) \quad (2.2)$$

323 Similar to Eq. 2.1 Eq. (2.1), the second-term represents information gain based on kernel
 324 function values subtracted from variance $k(x, x)$.
 325

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Using the Eqs. (2.1) and (2.2), the mean function and kernel function inferred based on the *posterior* f_* for a candidate x_{t_c} of acquisition function α . Generally, candidates are sampled across the search space of x , and gradient based optimization techniques like L-BFGS-B are used to estimate the maxima as x_t .

$$x_{t+1} = \arg \max_{x \in \mathcal{X}} \alpha(x \mid \mathcal{D}_t) \quad (3)$$

The Eq. 3 is the acquisition function α to efficiently select the next sample x_t to evaluate f for y_t . An ideal choice of α should efficiently balance exploration and exploitation. A few choices for the acquisition function include Expected Improvement, Thompson Sampling, Probability of Improvement, Upper Confidence Bound, etc.

$$y_{t+1} = f(x_{t+1}) \quad (4)$$

In some cases, a Gaussian noise term $\varepsilon \sim \mathcal{N}(0, \sigma^2)$ is added to function evaluation as $f(x) + \varepsilon$ to make BO robust and not overfit.

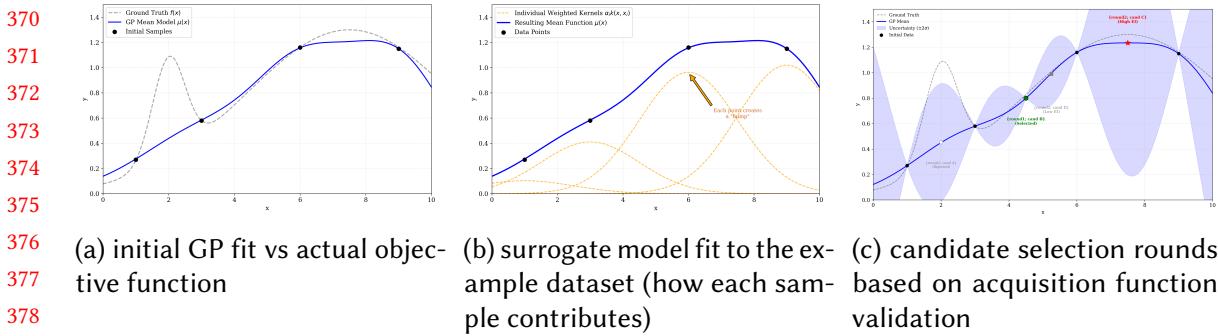
$$\mathcal{D}_{t+1} = \mathcal{D}_t \cup \{(x_{t+1}, y_{t+1})\} \quad (5)$$

Until a problem-specific goal or stopping-criteria is achieved, more samples are acquired and the GP is refit following Eqs. (2) - (5).

Applications of BO include hyper-parameter tuning in neural networks, kriging in mining, etc.

4.1.2 Attacks and Countermeasures.

High Resource Attack: A high resource attack also referred to as resource exhaustion attack is a type of attack in which an adversary aims to intentionally diminish the target's defined computational resources, including cpu cycles, memory allocation, bandwidth, or disk space leading the system or service unavailable to authentic users. Such attacks aim to overwhelm the target by provoking disproportionate demand for these resources, frequently recurring actions that exploit vulnerabilities in resource management. Salient characteristics of resource high resource attacks comprise their non-destructive nature to data integrity and confidentiality hence their primary goal is to disrupt the service availability, degrade performance without altering or stealing information.



The core principle behind resource exhaustion attacks resource asymmetry leads attackers to incur minimal cost while forcing the target system to perform disproportionately expensive operations. This imbalance gives adversaries an advantage to exploit and induce significant resource consumption using relatively low-rate or low-volume inputs. The on-demand resource allocation nature of modern systems particularly makes them vulnerable to resource exhaustion attack. Dynamically provisioned resources, such as buffers, threads, and connections, can accumulate when malicious requests exceed reclamation capacity, leading to a denial of service. Systems that lack adequate validation, rate limiting, or allocation controls are especially susceptible, enabling rapid exhaustion even at moderate attack rates.

These attacks are aggravated by persistence and amplification mechanisms. Persistence prolongs resource occupancy, while amplification enables low-cost inputs to trigger disproportionately expensive system responses.

Depending on the targeted resource, resource exhaustion attacks can be broadly classified as CPU exhaustion, memory exhaustion, bandwidth exhaustion, connection exhaustion, and disk exhaustion attacks.

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411 Bayesian Optimization as High Resource Attack: Bayesian optimization has become a
412 prominent technique in adversarial machine learning, particularly when dealing with the
413 more realistic black-box setting that requires an attacker to find an adversarial perturba-
414 tion without any knowledge of the architecture, parameters, or training data of the target
415 model. In such cases, information about the model can only be obtained through queries,
416 i.e. supplying an input to the model and receiving the corresponding output. In these situa-
417 tions, an attacker can treat interactions with the target model as a sequential optimization
418 problem, in which like the model's response latency, predicted labels, confidence scores
419 and resource utilization , serve as noisy measurements of some unknown underlying goal
420 the attacker is trying to achieve.

*421 When applied to high-resource attacks, BO enables an adversary to efficiently identify
422 inputs that result in maximal computational cost during model inference. By approximat-
423 ing the behavior of the target model with a probabilistic surrogate and using acquisition
424 functions to guide the selection of the query, Bayesian optimization allows an attacker to
425 progressively identify inputs that induce execution paths of the worst-case.*

*426 Triggering such inputs in a loop may result in extensive resource exhaustion, leading to
427 system degradation rendering the failure of service or denial of service. This trait of BO
428 outperforms random search and heuristic approaches and makes it particularly effective
429 for these attacks by achieving greater impact in fewer iterations. Despite their effectiveness,
430 these attacks are constrained by the presence of resource control measures, including rate
431 limiting, timeouts, and input validation. Furthermore, targeting well-provisioned systems
432 typically requires distributed attack infrastructures, increasing the risk of detection. Hence,
433 some mitigation strategies like limiting resource quotas, load balancing, anomaly detection
434 and adaptive throttling, are critical to minimizing system vulnerability.*

*435
436
437
438
439
440 Countermeasures:* Despite the fact that their effectiveness, these attacks are constrained
441 by the presence of resource control measures, including rate limiting, timeouts, and input
442 validation. Moreover, targeting well-resourced systems usually requires distributed attack
443 infrastructures thereby increasing the risk of detection.

444 Responsibility: Amit Singh

445 Fill Tables 1 and 2.

Table 1. Comparison of Probing-Based Analysis Techniques

Paper (Domain)	Method	Conditions	Resources	System Knowledge	Countermeasures?	Verification Method	Limitations	Relevance	Notes
Dummy row - sent remove this[3](NFs)	Packets	Network jitter	High CPU, moderate memory	Black-box	Yes	Synthetic datasets	Limited scalability	High	Seminal

Table 2. Comparison of Countermeasures

Paper (Domain)	Method	Performance Impact	Resource Overhead	Test Method	Limitations	Relevance	Remarks
Dummy row - remove this[2](NFs)	BO	Latency increase under peak load	Moderate CPU and memory overhead	Real work	net- Reduced effectiveness under adaptive attacks	High	Dataset not shared

Write a few paragraphs on attacks and countermeasures here summarising each paper along the columns of Tables 1 and 2. Add relevant information, if any, not mentioned in the table. For example, how the probing and later the attack happens, whether there are trade-offs, etc. State the dominant features across the papers that you have studied that you want to highlight.

While you do this, if there is something that you find applicable to the section on open challenges and future research directions, add that there.

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493 4.2 Low Rate Flow-table overflow attacks (Low Resource)**494 Responsibility:** Robin495 Copy Tables.**1** and **2** here.496 Write a few paragraphs on attacks and countermeasures here summarising each paper
497 along the columns of Tables **1** and **2**. Add relevant information, if any, not mentioned in
498 the table. For example, how the probing and later the attack happens, whether there are
499 trade-offs, etc. State the dominant features across the papers that you have studied that
500 you want to highlight.501 While you do this, if there is something that you find applicable to the section on open
502 challenges and future research directions, add that there.**503 5 Open Challenges and Future Research Directions****504 5.1 Bayesian Optimization**505 **5.1.1 Open Challenges.**506 **5.1.2 Future Directions.****507 5.2 Low Rate Flow-table overflow attacks**508 **5.2.1 Open Challenges.**509 **5.2.2 Future Directions.**510 **5.3****511 6 Conclusions****512 References**

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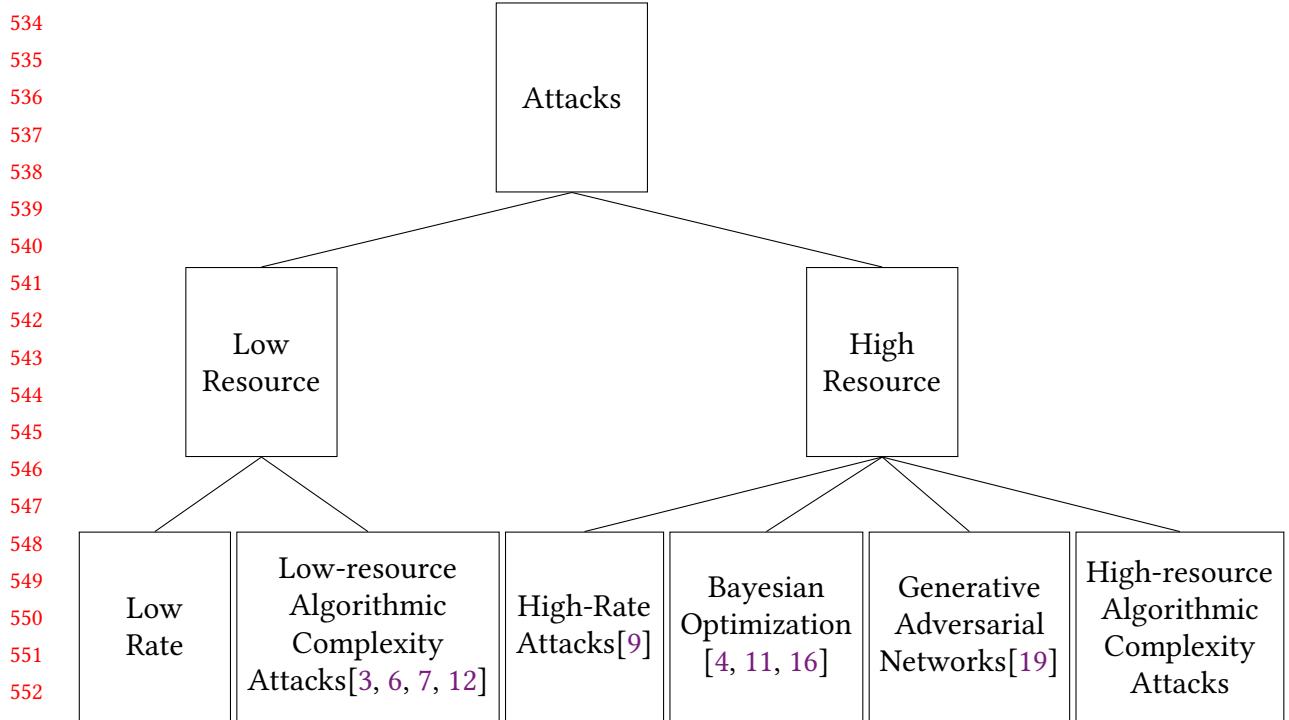


Fig. 3. Attack taxonomy

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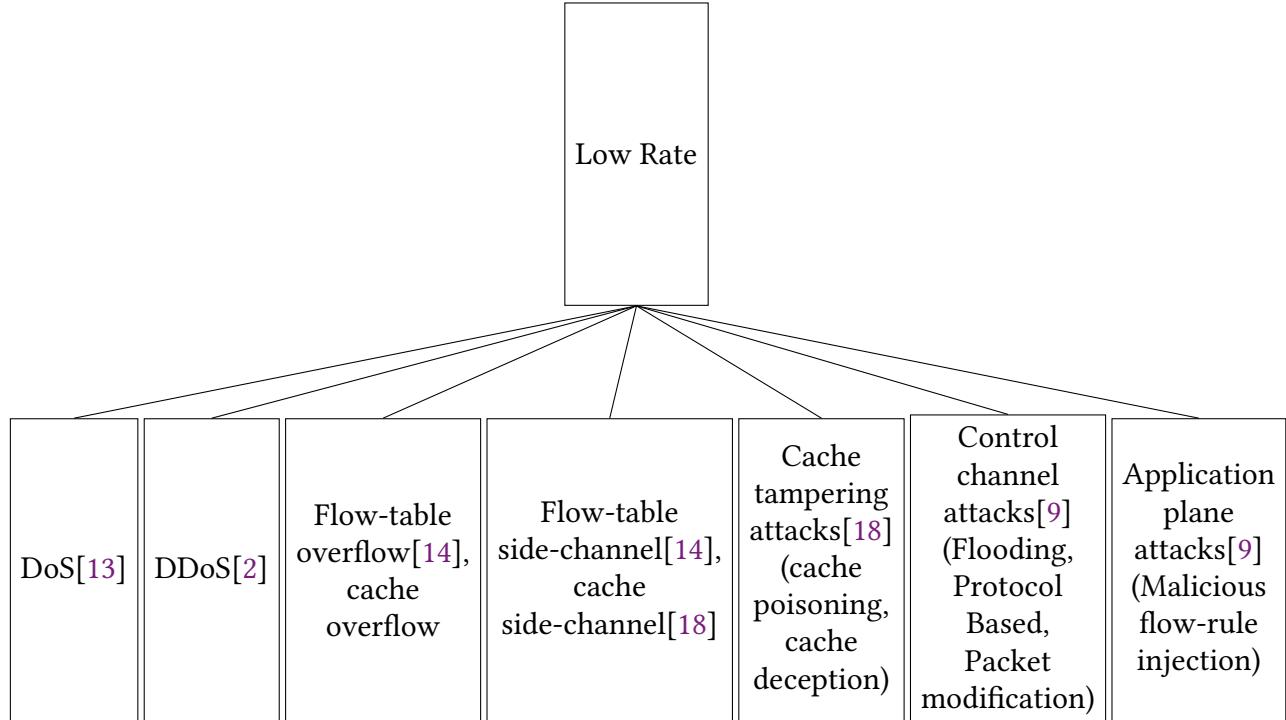


Fig. 4. Taxonomy of low-rate attacks

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