

1 **Caching systems: attacks and**
2 **countermeasures - A Survey**

6 Write the abstract here.
7

8 CCS Concepts: • **Do Not Use This Code → Generate the Correct Terms for Your Paper;**
9 *Generate the Correct Terms for Your Paper; Generate the Correct Terms for Your Paper; Generate*
10 *the Correct Terms for Your Paper.*
11

12 Additional Key Words and Phrases: Do, Not, Use, This, Code, Put, the, Correct, Terms, for, Your,
13 Paper
14

15 **ACM Reference Format:**

16 . 2025. Caching systems: attacks and countermeasures - A Survey. In *Proceedings of Make sure*
17 *to enter the correct conference title from your rights confirmation email (Conference acronym 'XX).*
18 ACM, New York, NY, USA, 17 pages. <https://doi.org/XXXXXXX.XXXXXXX>
19

20 **1 Introduction**

21 **Responsibility: Radhika**

22 In Multi-Access Edge Computing (MEC) networks, servers are collocated with Base
23 Stations (BSs) at the edge of the network, to reduce latency incurred by end users, such as
24 Mobile Stations (MSs). These servers are not as powerful as the cloud in terms of RAM,
25 CPU or storage. When end users request for content (images, videos etc.), the requests

26 Author's Contact Information:
27

28 Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted
29 without fee provided that copies are not made or distributed for profit or commercial advantage and that
30 copies bear this notice and the full citation on the first page. Copyrights for components of this work owned
31 by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise,
32 or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee.
33 Request permissions from permissions@acm.org.
34

35 © 2025 Copyright held by the owner/author(s). Publication rights licensed to ACM.
36

37 Manuscript submitted to ACM
38

39 Manuscript submitted to ACM
40

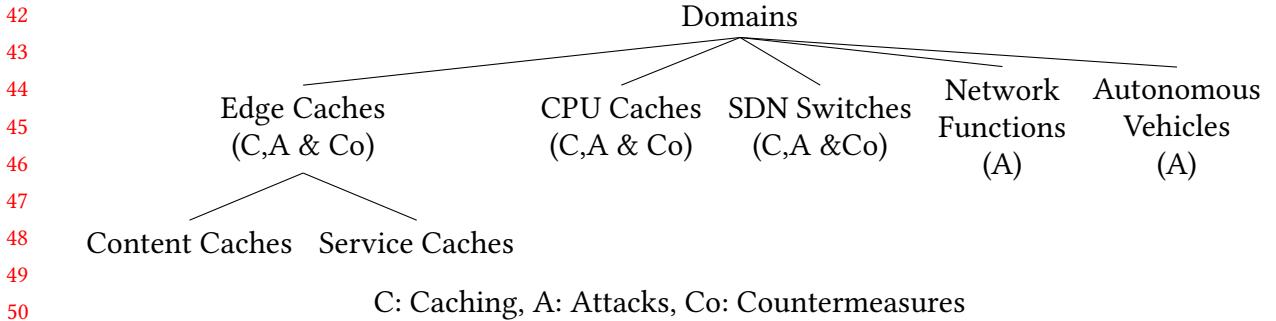


Fig. 1. Domains surveyed

arrive at the edge and are forwarded to the cloud for the content to be downloaded to the end device. Forwarding all requests to the cloud and getting a response can cause high latencies. Moreover, if there are multiple users, the links from the edge to the cloud can get overloaded, causing further delays. Therefore, each time content is downloaded, it can be cached at the edge.

Services are applications or processes that respond to requests from the user. For example, an Alternate Reality (AR) or Virtual Reality (VR) game requires an MS to send data from the device to the edge, which forwards it to a AR/VR game service running in the cloud. When data is sent such a service, a response is sent to the user and then the user can make the next move. To reduce latency, such a service and its associated database can be downloaded to the edge and cached. This differs from content caching in two ways: 1) A service request requires only a response, which is typically much smaller than content. 2) Content has to be always downloaded to the end device, whereas the service is never downloaded to the end device. 3) The edge can decide when to download a service based on an *admission policy*, while it always downloads content in response to a request.

This survey examines attack methods and corresponding countermeasures for systems that employ caching mechanisms, with particular attention to attacks and defenses relevant to edge caching systems. We do not explore any cryptographic methods, or attacks and countermeasures that are not relevant to edge caching systems. We intend to explore attack methods that may be relevant to caching systems, but the targets of which may not necessarily be caching systems.

Caches are ubiquitous in computer systems. As discussed above, there are caches at the edge, caching content and services. In addition, we explore attacks and countermeasures of two other systems that use caching: CPUs and SDN switches. To explore attacks that are relevant to edge caching, we explore Network Functions and Autonomous Vehicle Simulators. Fig. 1 illustrates the domains we survey.

CPU Caches: In the CPU microarchitecture, L1, L2 and L3 caches reside between the processor and the main memory. Typically, L1 and L2 caches are private to each core in a multiprocessor architecture, whereas L3 is common across cores. These caches are vulnerable to various types of attacks. [Expand: \[10\], \[17\]](#)

SDN Switches: In computer networks, the forwarding plane or the data plane forwards packets arriving at switches and routers according to the rules installed on them, while the control plane decides which rules to install. In Software Defined Networks (SDN), these planes are separated: the control plane is centralized in a controller, which programmatically controls the forwarding planes present in SDN switches. SDN switches store rules in Ternary Content Accessible Memory (TCAM) and Static Random Access Memory (SRAM). TCAMs are fast, but expensive, while SRAMs are slow and comparatively less expensive. Therefore, TCAMs may be considered caching flow rules. Therefore, the vulnerabilities, attack methods, and countermeasures are directly relevant.

Domain Name Systems: Domain Name Systems (DNS) translate human readable domain names into Internet Protocol (IP) addresses. A DNS cache stores the result of a name lookup at the browser level, at the Operating System (OS) level or at the level of the Internet Service Provider (ISP). Afek et al. [1] provide a history of DNS poisoning attacks.

Content Delivery Networks (CDNs): In a typical content delivery workflow, a user request is first resolved via DNS. The CDN-operated authoritative DNS service selects an appropriate edge server based on factors such as the location of the resolver, network conditions and the load. The DNS response directs the client to a specific CDN edge node, which is often physically colocated within an ISP point of presence, but remains under the operational control of the CDN. The user's HTTP(S) request is then sent directly to this edge server. If the requested content is present in the edge cache, it is served immediately; otherwise, the edge retrieves the content from the origin server (or an upstream parent cache), forwards it to the user, and caches it according to cache-control policies. Commercial agreements between ISPs and CDNs typically govern colocation,

124 connectivity, and traffic locality, while content placement, DNS-based request steering,
125 and caching decisions are centrally managed by the CDN. CDNs cache objects and small
126 state-less functions, unlike an edge cache, which can cache services too. [Read \[8\] for more](#)
127 [on this](#).

128

129 Note: Caching in the World Wide Web and Information Centric Networking may also
130 be studied. See [15] and associated papers.

131 **Network Functions:** Network Functions are software modules implemented on servers
132 that perform packet processing for reasons other than forwarding or routing, such as
133 securing the network, optimizing traffic, etc. Intrusion Detection and Prevention Systems
134 (IDPS), Firewalls and Wide Area Network (WAN) Optimizers are some examples. It is
135 useful to learn the Attacks and Countermeasures employed in these systems with a view
136 to using them in edge caches.

137

138 **Autonomous Vehicles:** Autonomous Vehicles (AV) aim to revolutionize driving in
139 urban areas and highways. Simulators are used to test the safety of such vehicles. The
140 techniques used to develop a homologation framework (the official certification process
141 that ensures that vehicles comply with government requirements) for critical scenarios
142 for such frameworks could be used for edge caching systems [4]. An AV planner is the
143 part of an autonomous driving system that decides what the vehicle should do next and
144 automated stress testing of planners using generated scenarios could be used for similar
145 purposes for edge caching systems [11].

146 **Our Contributions:** We study the domains listed in Fig.1 (and explained above) for
147 ideas relevant to reconnaissance, attacks, and countermeasures for edge service caching
148 systems. Moreover, we propose taxonomies for attacks and countermeasures and summa-
149 rize findings from the recent literature according to these taxonomies. We also discuss
150 open challenges and future research directions.

151 We organize prior work along the following axes: attack methods and countermeasures,
152 and domains surveyed. We examine attack methods according to their:
153

- 154
- 155 • method of probing (reconnaissance), if applicable or if any and what is probed
156 (replacement algorithm, size of cache, etc.)
- 157 • conditions for probing (do they consider network jitter, link delays, link failures,
158 etc.)

- 165 • resource consumption of the attack (low, medium, high)
- 166 • the extent of knowledge of the system required by the attackers (black box, gray
- 167 box, or white box)
- 168 • relevance to edge service caching
- 169 • existence of countermeasures, if any
- 170 • verification method (tested on real networks or simulators, whether the datasets
- 171 used, if any, are realistic)
- 172 • limitations (assumptions on the system by the attackers)

173
174 We examine countermeasures according to their:

- 175 • method (heuristic, AI etc.)
- 176 • impact on performance
- 177 • impact on resource consumption,
- 178 • tested on real networks or simulators, the datasets used, if any
- 179 • relevance to edge service caching
- 180 • limitations (are there strong assumptions, are the datasets realistic, etc.)

186 1.1 Comparison with existing surveys

187 Responsibility: TBD

190 1.2 Why this survey is different

191 Responsibility: TBD

192 Add comparison with existing surveys on caching systems, Low-rate attacks in SDN, Blackbox
193 attacks etc.

194 Our contributions:

198 2 Background and system model

199 Responsibility: TBD

200 The domains surveyed in this paper are shown in Fig. 1.

201 Explain 1) The system model for each domain 2) Definitions of common terms and
202 notations. In the rest of the paper you don't need to define any term again.

206 2.1 SDN Switches

207 Responsibility: Robin.

209 Write briefly about how an SDN works (Switches, Controller, Openflow, TCAM, caching
210 etc.).

211 Software Defined Networking (SDN) introduces a paradigm shift in network design by
212 decoupling the control plane from the data plane. In traditional networks, forwarding
213 devices such as switches and routers independently make packet forwarding decisions
214 based on locally stored routing and forwarding logic. This tightly coupled architecture
215 makes network management complex and inflexible.

217 In SDN, the control plane is logically centralized in an SDN controller, while the data
218 plane consists of multiple SDN switches that perform packet forwarding. The controller
219 acts as the network's decision-making entity, maintaining a global view of the network
220 and dynamically installing forwarding rules on switches using standardized southbound
221 interfaces such as OpenFlow.

223 SDN switches maintain flow tables that store flow rules defining how packets should be
224 handled. These rules typically match on packet header fields such as source and destination
225 IP addresses, transport-layer ports, and protocol identifiers. For performance reasons,
226 flow rules are commonly stored in Ternary Content Addressable Memory (TCAM), which
227 enables fast parallel lookups. However, TCAM is expensive and has limited capacity,
228 making flow table space a scarce resource.

231 When a packet arrives at a switch, the switch performs a lookup in its flow table. If a
232 matching rule is found, the packet is forwarded accordingly at line rate. If no matching rule
233 exists, the packet (or its header) is forwarded to the controller via a Packet-In message. The
234 controller then determines the appropriate forwarding action and may install a new flow
235 rule in the switch. This behavior effectively treats the flow table as a cache, where cache
236 hits result in fast forwarding and cache misses trigger interaction with the controller.

238 Due to the limited capacity of TCAM and the reactive nature of rule installation, SDN
239 switches are vulnerable to attacks that aim to exhaust flow table space or overload the
240 control channel. These characteristics make SDN switches a relevant domain for studying
241 caching-related attacks and countermeasures.

2.2 Network Functions

248 Responsibility: Robin What are Network Functions?

249

251 2.3 Relevance to edge service caching

252 Explain why attacks and countermeasures in domains other than caching are relevant.
253 An example is: The cache itself can be attacked. Attacks can be triggered such that the
254 links from edges to the cloud are saturated. If a TCAM in an SDN switch is considered to
255 be a cache, a cache miss for a packet would result in traffic to the controller until a flow
256 rule is downloaded to the cache.
257

258

260 3 Attack taxonomy

261 Refer Fig.3 and 4 for the attack taxonomy.

262 Need to decide which of the ACAs given the figure are low-resource and which are
263 high-resource. Currently, all are grouped under low-resource.

264 Need to add more survey papers related to attacks on caches (pollution attacks etc.).
265

266

268 4 Attacks and countermeasures**269 4.1 Bayesian Optimization (High resource attacks)**

270 *4.1.1 Introduction to Bayesian Optimization.* Responsibility: Vasudeva

271 This should be 2 pages long and must describe the basic idea with references. Use
272 mathematical notation wherever required.

273 Bayesian Optimization (BO) is one of the powerful strategies for finding extrema of
274 objective functions. Generally, an optimization problem is also formulated as maximizing
275 function $f(x)$ over a domain $A \subset \mathbb{R}^d$, i.e., $\max_{x \in A \subset \mathbb{R}^d} f(x)$. Particularly, a few character-
276 stics of $f(x)$ make BO the appropriate choice of optimization strategy such as costly to
277 evaluate, unknown/no mathematical representation, no access to derivatives, non-convex,
278 limited observations, noisy observations, etc. [5].

279 BO is a strategy combining ideas from Bayesian inference (Bayes' theorem), Surro-
280 gate modeling (mean function, covariance/kernel function), Sequential decision-making
281 (acquisition functions like expected improvement).

288 BO is grounded in Bayes' theorem. Bayes' theorem broadly states *posterior* probability
 289 of a model (or hypothesis) f_* given dataset (or observations) D equals *likelihood* of f_*
 290 given D multiplied by *prior* probability of f_* and divided by the probability of D , i.e.,
 291

$$292 \quad p(f_* | D) = \frac{p(D | f_*) p(f_*)}{p(D)} \\ 293$$

294 since $p(D)$ does not depend on f_* , it is also represented as $p(f_* | D) \propto p(D | f_*) p(f_*)$
 295 in BO literature. Along the lines of Bayes' theorem, BO assumes a prior belief for the
 296 unknown objective function f , uses the initial samples in D to fit f_* (referred to as
 297 *posterior*) as a candidate for f . In order to efficiently sample more data points, BO uses
 298 acquisition functions to validate the search space of x_t to acquire a new sample, evaluate
 299 it for y_t using f , and refit f_* , making the new f_* (*posterior*) closer to f . BO continues
 300 to sample more data until a stopping criteria. The f is modeled as a random function
 301 drawn from a probability distribution over functions. The prior belief in BO represents
 302 the space of possible objective functions and is an inductive bias encoding assumptions
 303 about the nature of f , such as smoothness, continuity, noise, etc, that make some possible
 304 functions more plausible. The prior belief also includes the choice of kernel function and
 305 its hyper-parameters, the choice of acquisition function and its hyper-parameters, the
 306 choice of other hyper-parameters, among others. For this reason, the prior affects the
 307 sampling efficiency in acquisition functions, and overall convergence in BO and thus the
 308 choice of prior is crucial.
 309

310 does GP assume lipschitz-continuous? BO conditions on a dataset D with initial samples,
 311 where D is a set of initial i pairs of x, y and $x \in \mathbb{R}^d, y \in \mathbb{R}$, i.e., $D = \{x_{1:i}, y_{1:i}\}$ and f
 312 evaluates x to y . A prior belief, also referred to as a surrogate model, is assumed based on
 313 characteristics of the objective function. In BO literature, the standard surrogate model
 314 for f is a Gaussian Process (GP) defined in the Eq. below. A GP is intuitively understood
 315 as a distribution over functions; that is, each sample drawn from a GP corresponds to
 316 a mathematical function. The mean function and the kernel function (also referred to
 317 as the pair-wise covariance function of samples in D) parameterize the characteristics
 318 of this distribution. If $x_t \in \mathbb{R}^d$, the GP is d-dimensional, and a 0-dimensional GP can be
 319 understood as a Gaussian distribution. See an example BO run in the figure ref2. Examples
 320

321

322 Manuscript submitted to ACM

323

329 of other surrogate models are Student-t Process, etc.

330 $f_* \sim \mathcal{GP}(\mu(x), k(x, x'))$

332 where $\mu(x)$ represents mean function and $k(x, x')$ is kernel function of the *GP*. Applying
 333 Bayes' theorem on GP leads to inference rules for mean function ($\mu_t(x)$) and kernel
 334 function (σ_t^2) after t samples are appended to the initial i in D . Each refit as shown in Eq.
 335 (1) increases the probability that a sample f_* , drawn from the *GP*, is close to the actual
 336 f . Generally, the $\mu = 0$ and the $k(x, x) = 1$ at the start. A few choices for kernel function
 337 include Squared Exponential (RBF) kernel, Matérn kernel, Rational Quadratic kernel, etc.
 338

340 $f_* = p(f(x) | \mathcal{D}_t) \quad (1)$

342 The kernel matrix K , calculated using D , and y represent a fit of f_* . At the start f_* is fit
 343 using the initial i samples in D (the i is a hyper-parameter and is problem-specific). The
 344 f_* is refit with the new sample acquired after a few candidates are collected, validated
 345 by an acquisition function, evaluated using the actual objective function f , and finally,
 346 appended to the D , see Eqs. (2)-(5).

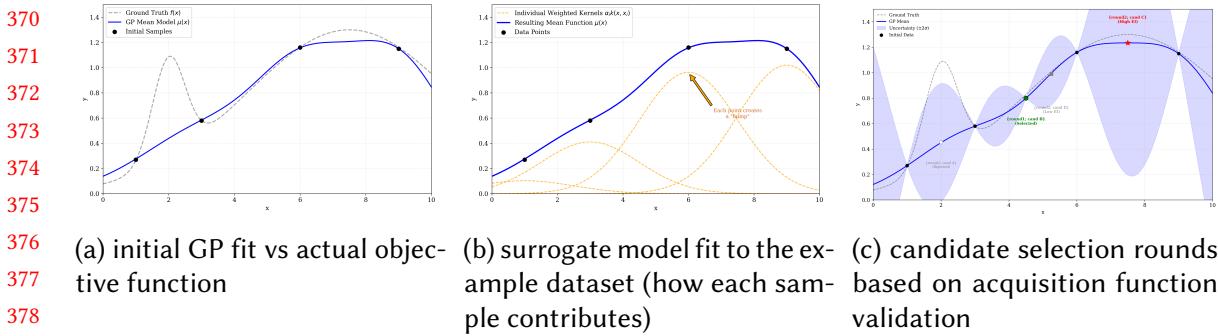
347 $\mu_t(x_{t_c}) = \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{y}_{1:i+t-1} \quad (2.1)$

348 Where $\mathbf{k}_*(x_t) \in \mathbb{R}^{i+t-1}$ represents the covariance vector of x_t with the other $i + t - 1$
 349 samples in D , $\mathbf{K}_{t-1} \in \mathbb{R}^{i+t-1 \times i+t-1}$ represents the pair-wise covariance matrix of the other
 350 $i + t - 1$ samples, and $\mathbf{y}_{1:i+t-1} \in \mathbb{R}^{i+t-1}$ represents the y values vector of all $i + t - 1$ samples
 351 in D . Intuitively, the pairwise covariance of the $x_1 : i + t - 1$ data points is projected onto
 352 the y space and weighted by the covariance vector of the new point x_t .

353 $\sigma_t^2(x_t) = k(x_t, x_t) - \mathbf{k}_*(x_t)^\top (\mathbf{K}_{t-1})^{-1} \mathbf{k}_*(x_t) \quad (2.2)$

354 Similar to Eq. 2.1 Eq. (2.1), the second-term represents information gain based on kernel
 355 function values subtracted from variance $k(x, x)$.

356 Using the Eqs. (2.1) and (2.2), the mean function and kernel function inferred based on
 357 the *posterior* f_* for a candidate x_{t_c} of acquisition function α . Generally, candidates are
 358 sampled across the search space of x , and gradient based optimization techniques like
 359



411 4.1.2 *Attacks and Countermeasures.*

412
413 *High Resource Attack:* A high resource attack also referred to as resource exhaustion at-
414 tack is a type of attack in which an adversary aims to intentionally diminish the target's de-
415 fined computational resources, including cpu cycles, memory allocation, bandwidth, or disk
416 space leading the system or service unavailable to authentic users. Such attacks aim to over-
417 whelm the target by provoking disproportionate demand for these resources, frequently
418 recurring actions that exploit vulnerabilities in resource management. Salient character-
419 istics of resource high resource attacks comprise their non-destructive nature to data
420 integrity and confidentiality hence their primary goal is to disrupt the service availability
421 , degrade performance without altering or stealing information.

422
423 The core principle behind resource exhaustion attacks resource asymmetry leads attack-
424 ers to incur minimal cost while forcing the target system to perform disproportionately
425 expensive operations. This imbalance gives adversaries an advantage to exploit and induce
426 significant resource consumption using relatively low-rate or low-volume inputs. The on-
427 demand resource allocation nature of modern systems particularly makes them vulnerable
428 to resource exhaustion attack. Dynamically provisioned resources, such as buffers, threads,
429 and connections, can accumulate when malicious requests exceed reclamation capacity,
430 leading to a denial of service. Systems that lack adequate validation, rate limiting, or
431 allocation controls are especially susceptible, enabling rapid exhaustion even at moderate
432 attack rates.

433
434 These attacks are aggravated by persistence and amplification mechanisms. Persistence
435 prolongs resource occupancy, while amplification enables low-cost inputs to trigger
436 disproportionately expensive system responses.

437
438 Depending on the targeted resource, resource exhaustion attacks can be broadly classi-
439 fied as CPU exhaustion, memory exhaustion, bandwidth exhaustion, connection exhaus-
440 tion, and disk exhaustion attacks.

441
442 *Bayesian Optimization as High Resource Attack:* Bayesian optimization has become a
443 prominent technique in adversarial machine learning, particularly when dealing with the
444 more realistic black-box setting that requires an attacker to find an adversarial perturba-
445 tion without any knowledge of the architecture, parameters, or training data of the target
446 model. In such cases, information about the model can only be obtained through queries,

452 i.e. supplying an input to the model and receiving the corresponding output. In these situations,
453 an attacker can treat interactions with the target model as a sequential optimization problem,
454 in which like the model's response latency, predicted labels, confidence scores
455 and resource utilization , serve as noisy measurements of some unknown underlying goal
456 the attacker is trying to achieve.
457

458 When applied to high-resource attacks, BO enables an adversary to efficiently identify
459 inputs that result in maximal computational cost during model inference. By approximating
460 the behavior of the target model with a probabilistic surrogate and using acquisition
461 functions to guide the selection of the query, Bayesian optimization allows an attacker to
462 progressively identify inputs that induce execution paths of the worst-case.
463

464 Triggering such inputs in a loop may result in extensive resource exhaustion, leading to
465 system degradation rendering the failure of service or denial of service. This trait of BO
466 outperforms random search and heuristic approaches and makes it particularly effective
467 for these attacks by achieving greater impact in fewer iterations. Despite their effectiveness,
468 these attacks are constrained by the presence of resource control measures, including rate
469 limiting, timeouts, and input validation. Furthermore, targeting well-provisioned systems
470 typically requires distributed attack infrastructures, increasing the risk of detection. Hence,
471 some mitigation strategies like limiting resource quotas, load balancing, anomaly detection
472 and adaptive throttling, are critical to minimizing system vulnerability.
473

474

475

476

477

478

479

480

481

Countermeasures: Despite the fact that their effectiveness, these attacks are constrained
 by the presence of resource control measures, including rate limiting, timeouts, and input
 validation. Moreover, targeting well-resourced systems usually requires distributed attack
 infrastructures thereby increasing the risk of detection.

482

483

Responsibility: Amit Singh

484

Fill Tables 1 and 2.

485

486

487

488

489

490

Write a few paragraphs on attacks and countermeasures here summarising each paper
 along the columns of Tables 1 and 2. Add relevant information, if any, not mentioned in
 the table. For example, how the probing and later the attack happens, whether there are
 trade-offs, etc. State the dominant features across the papers that you have studied that
 you want to highlight.

491

Manuscript submitted to ACM

492

Table 1. Comparison of Probing-Based Analysis Techniques

Paper (Domain)	Method	Conditions	Resources	System Knowledge	Countermeasures?	Verification Method	Limitations	Relevance	Notes
Dummy row - sent remove this[3](NFs)	Packets	Network jitter	High CPU, moderate memory	Black-box	Yes	Synthetic datasets	Limited scalability	High	Seminal

Table 2. Comparison of Countermeasures

Paper (Domain)	Method	Performance Impact	Resource Overhead	Test Method	Limitations	Relevance	Remarks
Dummy row - remove this[2](NFs)	BO	Latency increase under peak load	Moderate CPU and memory overhead	Real work	net-effectiveness under adaptive attacks	High	Dataset not shared

While you do this, if there is something that you find applicable to the section on open challenges and future research directions, add that there.

4.2 Low Rate Flow-table overflow attacks (Low Resource)

Responsibility: Robin

Copy Tables.1 and 2 here.

534 Write a few paragraphs on attacks and countermeasures here summarising each paper
 535 along the columns of Tables 1 and 2. Add relevant information, if any, not mentioned in
 536 the table. For example, how the probing and later the attack happens, whether there are
 537 trade-offs, etc. State the dominant features across the papers that you have studied that
 538 you want to highlight.
 539

540 While you do this, if there is something that you find applicable to the section on open
 541 challenges and future research directions, add that there.
 542

543 5 Open Challenges and Future Research Directions

544 5.1 Bayesian Optimization

545 5.1.1 *Open Challenges.*

546 5.1.2 *Future Directions.*
 547

548 5.2 Low Rate Flow-table overflow attacks

549 5.2.1 *Open Challenges.*
 550

551 5.2.2 *Future Directions.*
 552

553 5.3

554 6 Conclusions

555 References

- 556 [1] Yehuda Afek, Harel Berger, and Anat Bremler-Barr. 2025. POPS: From History to Mitigation of DNS
 557 Cache Poisoning Attacks. *arXiv preprint arXiv:2501.13540* (2025).
- 558 [2] Abdussalam Ahmed Alashhab, Mohd Soperi Mohd Zahid, Mohamed A Azim, Muhammad Yunis Daha,
 559 Babangida Isyaku, and Shimhaz Ali. 2022. A survey of low rate DDoS detection techniques based on
 560 machine learning in software-defined networks. *Symmetry* 14, 8 (2022), 1563.
- 561 [3] Nirav Atre, Hugo Sadok, Erica Chiang, Weina Wang, and Justine Sherry. 2022. Surgeprotector:
 562 Mitigating temporal algorithmic complexity attacks using adversarial scheduling. In *Proceedings of the
 563 ACM SIGCOMM 2022 Conference*. 723–738.
- 564 [4] Satyesh Shanker Awasthi, Mohammed Irshadh Ismaaeel Sathyamangalam Imran, Stefano Arrigoni,
 565 and Francesco Braghin. 2025. Bayesian Optimization applied for accelerated Virtual Validation of the
 566 Autonomous Driving Function. *arXiv preprint arXiv:2507.22769* (2025).

567 Manuscript submitted to ACM
 568

569

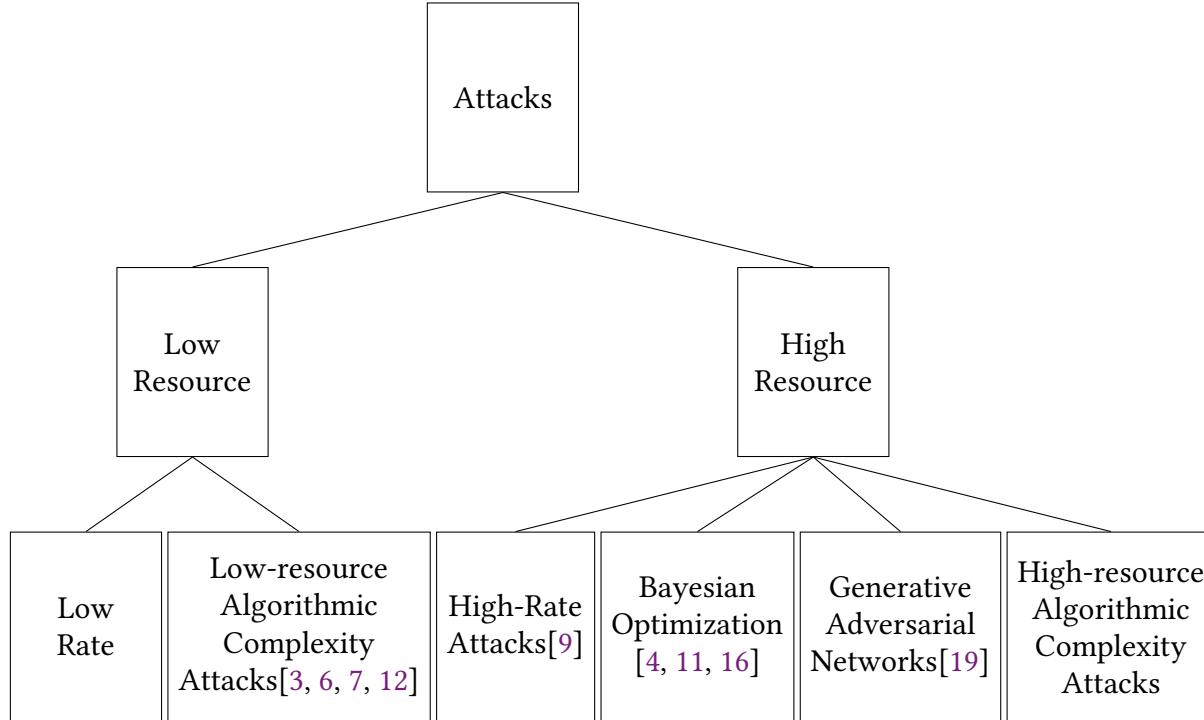


Fig. 3. Attack taxonomy

- [5] Eric Brochu, Vlad M. Cora, and Nando De Freitas. 2010. A tutorial on Bayesian optimization of expensive cost functions, with application to active user modeling and hierarchical reinforcement learning. *arXiv preprint arXiv:1012.2599* (2010). <https://arxiv.org/abs/1012.2599>
- [6] Scott A Crosby and Dan S Wallach. 2003. Denial of service via algorithmic complexity attacks. In *12th USENIX Security Symposium (USENIX Security 03)*.
- [7] Levente Csikor, Dinil Mon Divakaran, Min Suk Kang, Attila Kőrösi, Balázs Sonkoly, Dávid Haja, Dimitrios P Pezaros, Stefan Schmid, and Gábor Rétvári. 2019. Tuple space explosion: A denial-of-service attack against a software packet classifier. In *Proceedings of the 15th International Conference on Emerging Networking Experiments And Technologies*. 292–304.
- [8] Milad Ghaznavi, Elaheh Jalalpour, Mohammad A Salahuddin, Raouf Boutaba, Daniel Migault, and Stere Preda. 2021. Content delivery network security: A survey. *IEEE Communications Surveys & Tutorials* 23, 4 (2021), 2166–2190.
- [9] Suruchi Karnani, Neha Agrawal, and Rohit Kumar. 2024. A comprehensive survey on low-rate and high-rate DDoS defense approaches in SDN: taxonomy, research challenges, and opportunities. *Multimedia Tools and applications* 83, 12 (2024), 35253–35306.

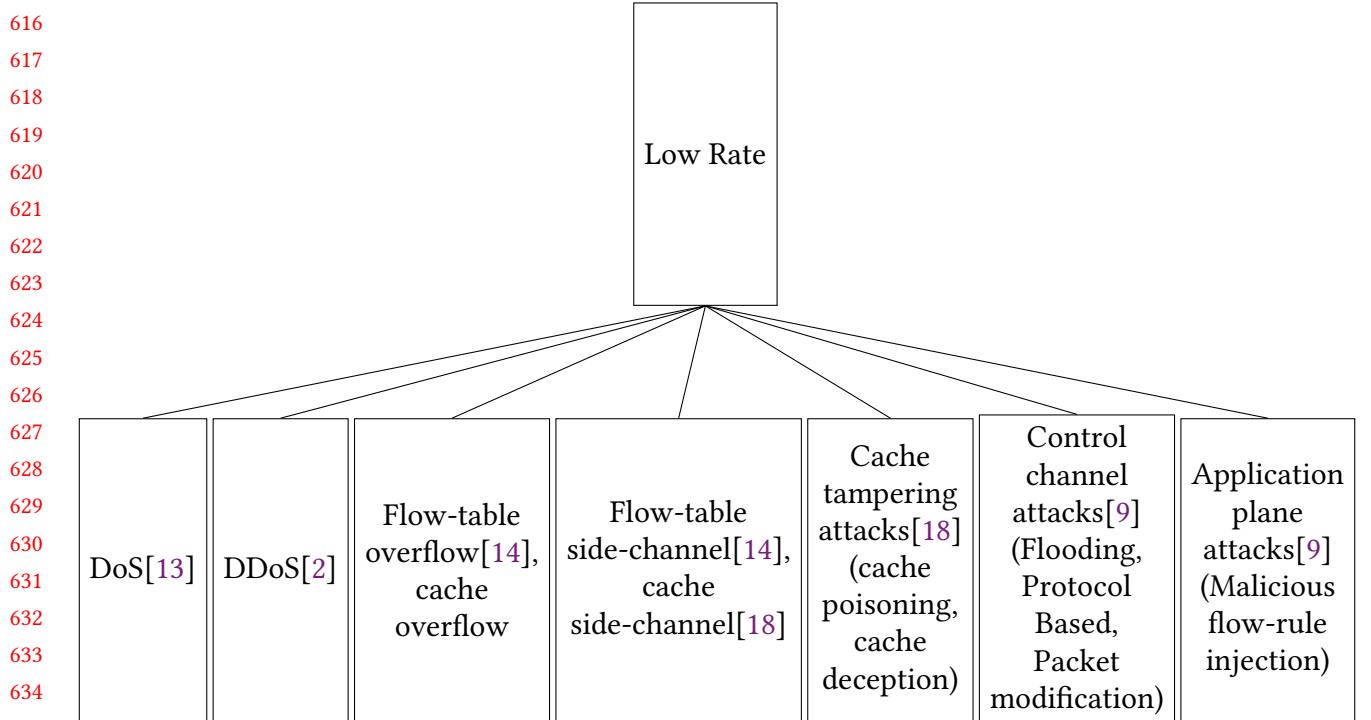


Fig. 4. Taxonomy of low-rate attacks

- [10] Yangdi Lyu and Prabhat Mishra. 2018. A survey of side-channel attacks on caches and countermeasures. *Journal of Hardware and Systems Security* 2, 1 (2018), 33–50.
- [11] Augusto Mondelli, Yueshan Li, Alessandro Zanardi, and Emilio Fazzoli. 2025. Test Automation for Interactive Scenarios via Promptable Traffic Simulation. *arXiv preprint arXiv:2506.01199* (2025).
- [12] Theofilos Petsios, Jason Zhao, Angelos D Keromytis, and Suman Jana. 2017. Slowfuzz: Automated domain-independent detection of algorithmic complexity vulnerabilities. In *Proceedings of the 2017 ACM SIGSAC conference on computer and communications security*. 2155–2168.
- [13] Vinícius De Miranda Rios, Pedro RM Inácio, Damien Magoni, and Mário M Freire. 2022. Detection and mitigation of low-rate denial-of-service attacks: A survey. *IEEE Access* 10 (2022), 76648–76668.
- [14] Dan Tang, Rui Dai, Yudong Yan, Keqin Li, Wei Liang, and Zheng Qin. 2024. When sdn meets low-rate threats: A survey of attacks and countermeasures in programmable networks. *Comput. Surveys* 57, 4 (2024), 1–32.
- [15] Tian Xie. 2024. *Resilient Cache Network Management: Algorithms, Analysis, Experiments*. The Pennsylvania State University.
- [16] Johannes Zerwas, Patrick Kalmbach, Laurenz Henkel, Gábor Rétvári, Wolfgang Kellerer, Andreas Blenk, and Stefan Schmid. 2019. Netboa: Self-driving network benchmarking. In *Proceedings of the* Manuscript submitted to ACM

657 2019 Workshop on Network Meets AI & ML. 8–14.

658 [17] Jiliang Zhang, Congcong Chen, Jinhua Cui, and Keqin Li. 2024. Timing side-channel attacks and
659 countermeasures in CPU microarchitectures. *Comput. Surveys* 56, 7 (2024), 1–40.

660 [18] Xianzhi Zhang, Yipeng Zhou, Di Wu, Quan Z. Sheng, Shazia Riaz, Miao Hu, and Linchang Xiao. 2025.
661 A Survey on Privacy-Preserving Caching at Network Edge: Classification, Solutions, and Challenges.
662 *Comput. Surveys* 57, 5 (2025), 1–38. [doi:10.1145/3706630](https://doi.org/10.1145/3706630)

663 [19] Yiran Zhu, Lei Cui, Zhenquan Ding, Lun Li, Yongji Liu, and Zhiyu Hao. 2022. Black box attack and
664 network intrusion detection using machine learning for malicious traffic. *Computers & Security* 123
665 (2022), 102922.

666
667 Received 20 February 2007; revised 12 March 2009; accepted 5 June 2009

668
669

670

671

672

673

674

675

676

677

678

679

680

681

682

683

684

685

686

687

688

689

690

691

692

693

694

695

696

697