

1. A Framework for Test and Analysis

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Testing e verifica del software

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Goals of Software Testing and Verification

- to assess software **qualities**

examples of sw qualities

- my program never crashes
- my program works
- my program is useful

- to make it possible to improve the software by finding **defects**

example of sw defects

- the pointer is not null
- when the user presses OK button, the application

Section 1

Validation and Verification

Validation & Verification

Validation

Does the software system meet the user's real needs? are we building the **right** software?

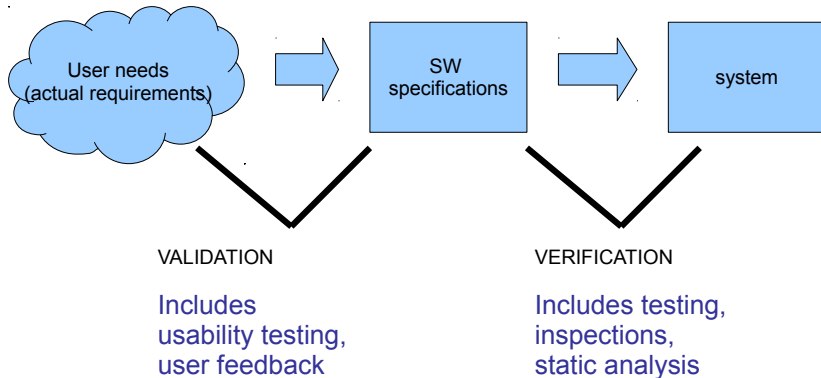
Specification

A statement (document) about a particular proposed solution to a problem.

Verification

Does the software system meet the requirements specifications? are we building the software **right**?

Validation & Verification



Validation & Verification - standard definitions

IEEE standard in its 4th edition defines the two terms as follows:

Validation. The assurance that a product, service, or system meets the needs of the customer and other identified stakeholders. It often involves acceptance and suitability with external customers. Contrast with verification.

Verification. The evaluation of whether or not a product, service, or system complies with a regulation, requirement, specification, or imposed condition. It is often an internal process. Contrast with validation.

ISO 9001 standard defines them this way :

Verification is the conformation that a product meets identified specifications.

Validation is the conformation that a product appropriately meets its design function or the intended use

Validation & Verification - standard definitions

Capability Maturity Model (CMMI-SW v1.1):

Software Verification: The process of evaluating software to determine whether the products of a given development phase **satisfy the conditions imposed at the start of that phase.**

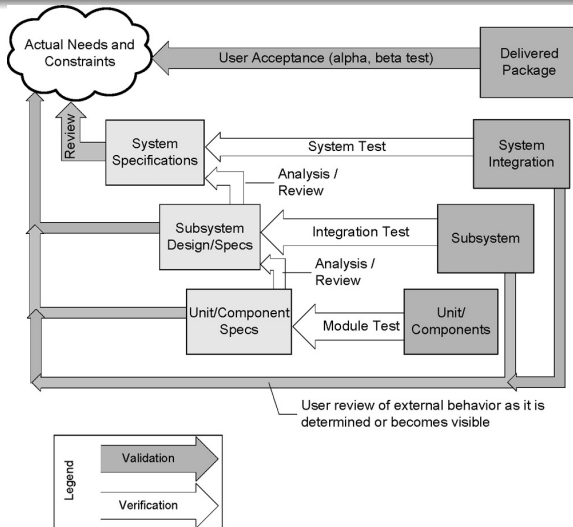
Software Validation: The process of evaluating software during or at the end of the development process to determine whether it **satisfies specified requirements.**

Boehm succinctly expressed the difference between:

Software Verification: Are we building the **product right?**

Software Validation: Are we building the **right product?**

Validation & Verification



Verification

- Verification generally compares two or more **artifacts**
- Verification can consist in checking for **self-consistency** and **well-formedness** **one** artifact.
 - For example, we can certainly determine that some programs are "incorrect" because they are **ill-formed**.
 - We may likewise determine that a specification itself is **ill-formed** because it is inconsistent (requires two properties that cannot both be true) or ambiguous (can be interpreted to require some property or not),
 - or because it does not satisfy some other **well-formedness** constraint that we impose, such as adherence to a standard imposed by a regulatory agency.

Verification

- Validation against actual requirements necessarily involves human judgment
- Verification can be automatized

Section 2

Degrees of Freedom

Validation & Verification

- ① Can we arrive at some logically sound argument or proof that a program satisfies the specified properties?
- ② Alan Turing proved that some problems cannot be solved by any computer program.
- ③ an **undecidable problem** is a decision problem for which it is known to be impossible to construct a single algorithm that always leads to a correct yes-or-no answer.
- ④ for instance the **halting problem**
- ⑤ every interesting property regarding the behavior of computer programs can be shown to "embed" the halting problem,

HALTING PROBLEM

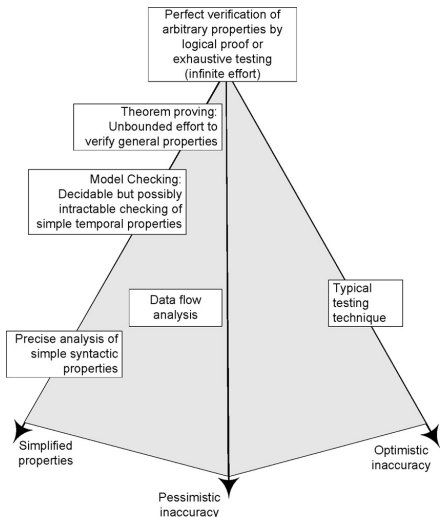
Given the description of an arbitrary program and a finite input, decide whether the program finishes running or will run forever.

Exhaustive testing

```
static int sum(int a, int b) {return a+b;}
```

- 1 Exhaustive testing, that is, executing and checking every possible behavior of a program, would be a "proof by cases," which is a correct way to construct a logical proof. How long would this take? *=> limitare : casi di test (approssimazione)*
- 2 there are only $2^{32} \times 2^{32} = 2^{64} \approx 10^{21}$ different inputs on which the method `Trivial.sum()` need be tested to obtain a proof of its correctness. At one nanosecond (10^{-9} seconds) per test case, this will take approximately 10^{12} seconds, or about 30,000 years.

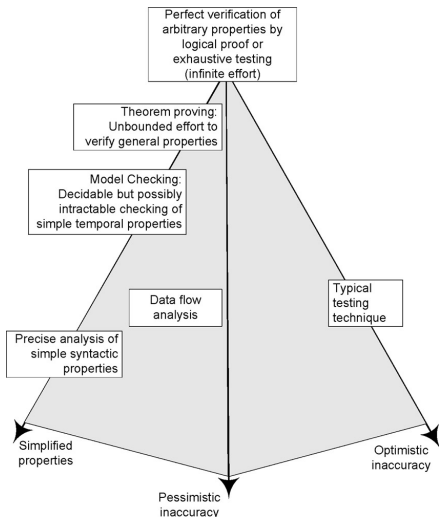
Pessimistic and Optimistic inaccuracy



A (testing/analysis) technique can be approximate:

- 1 **pessimistic**: it is not guaranteed to accept a program even if the program does possess the property being analyzed
- 2 **optimistic**: if it may accept some programs that do not possess the property (i.e., it may not detect all violations)

Simplification/abstraction



- ① we want to verify a property S , but
 - ① we cannot accept the optimistic inaccuracy of testing for S
 - ② precise analysis is too difficult
- ② a simpler property S' is a sufficient, but not necessary, condition for S
- ③ we check S' rather than S
- ④ we require S' to be satisfied

Simplification/abstraction Example

```
int i, sum;  
int first=1;  
for (i=0; i<10; ++i) {  
    if (first) {  
        sum=0; first=0;  
    }  
    sum += i;  
}
```

Property: each variable should be initialized with a value before its value is used in an expression

- 1 P vale??
- 2 in C language?
- 3 in Java ???

Example of simplified property: Unmatched Semaphore Operations

Property: every semaphore it is eventually unlocked

```
if ( .... ) {  
    ...    lock(S);  
}  
...  
if ( ... ) {  
    ...    unlock(S);  
}
```

Static checking for match is
necessarily inaccurate ...

Java solution: synchronized
statements specify the object
that provides the intrinsic lock

```
synchronized(S) {  
    ...  
}
```

It is guaranteed that the lock S
is released.

How to deal with undecidable problems

- ① **optimistic inaccuracy**: we may accept some programs that do not possess the property (i.e., it may not detect all violations).
 - testing
- ② **pessimistic inaccuracy**: it is not guaranteed to accept a program even if the program does possess the property being analyzed
 - automated program analysis techniques
- ③ **simplified properties**: reduce the degree of freedom for simplifying the property to check

Some Terminology

Safe (Sicuro): A **safe** analysis has no optimistic inaccuracy, i.e., it **accepts only correct programs.**

- if a program is “wrong” it is rejected.

Sound (Corretto): An analysis of a program P with respect to a formula F is sound if the analysis returns **true only when the program does satisfy the formula.**

- if a program is accepted, it is correct
- no wrong program is accepted
- there may be correct programs that are not accepted (conservative - pessimistic)
- testing is not sound.

Complete (completo): An analysis of a program P with respect to a formula F is complete if the analysis **always returns true when the program actually does satisfy the formula.**