

## Department of Computer Engineering

Academic Term : Jan-May 23-24

**Class : T.E. (Computer)**

**Subject Name : System Programming and Compiler Construction**

**Subject Code : (CPC601)**

<b>Practical No:</b>	4
<b>Title:</b>	To generate an Intermediate code.
<b>Date of Performance:</b>	
<b>Date of Submission:</b>	
<b>Roll No:</b>	9601
<b>Name of the Student:</b>	Ivan Dsouza

### Evaluation:

Sr. No	Rubric	Grade
1	Time Line (2)	
2	Output(3)	
3	Code optimization (2)	
4	Postlab (3)	

**Signature of the Teacher :**

Experiment No 4

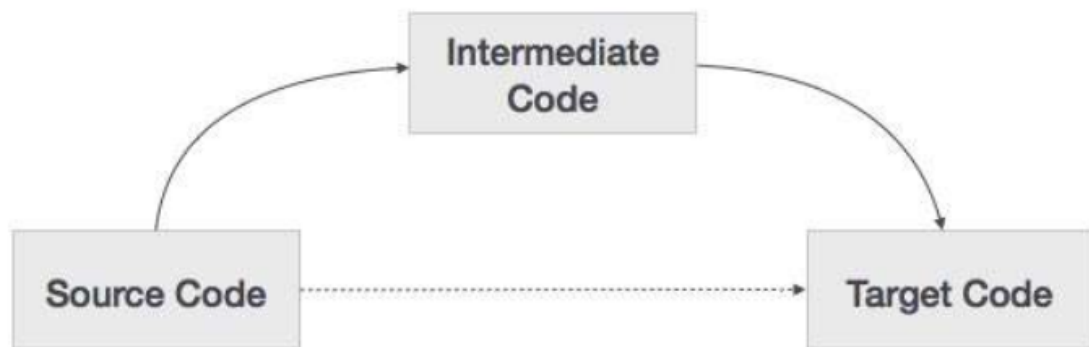
- Ivan Dsouza

- 9601

- T.E. Comps A (Batch C)

**Aim :** To generate an Intermediate code.

**Description:**



- If a compiler translates the source language to its target machine language without having the option for generating intermediate code, then for each new machine, a full native compiler is required.
- Intermediate code eliminates the need of a new full compiler for every unique machine by keeping the analysis portion same for all the compilers.
- The second part of compiler, synthesis, is changed according to the target machine.
- It becomes easier to apply the source code modifications to improve code performance by applying code optimization techniques on the intermediate code.

## • Three-Address Code

Intermediate code generator receives input from its predecessor phase, semantic analyzer, in the form of an annotated syntax tree. That syntax tree then can be converted into a linear representation, e.g., postfix notation. Intermediate code tends to be machine independent code. Therefore, code generator assumes to have unlimited number of memory storage (register) to generate code.

- **For example:**

- `a = b + c * d;`

- The intermediate code generator will try to divide this expression into sub-expressions and then generate the corresponding code.

- `r1 = c * d;`
- `r2 = b + r1;`
- `a = r2`

System Programming and Compiler Construction

VI Semester ( Computer) Academic Year: 23-24

- r being used as registers in the target program.
- A three-address code has at most three address locations to calculate the expression.  
A three-address code can be represented in two forms : quadruples and triples.

## Quadruples

Each instruction in quadruples presentation is divided into four fields: operator, arg1, arg2, and result. The above example is represented below in quadruples format:

Op	arg <sub>1</sub>	arg <sub>2</sub>	result
*	c	d	r1
+	b	r1	r2
+	r2	r1	r3
=	r3		a

## Triples

Each instruction in triples presentation has three fields : op, arg1, and arg2. The results of respective sub-expressions are denoted by the position of expression. Triples represent similarity with DAG and syntax tree. They are equivalent to DAG while representing expressions.

Op	arg <sub>1</sub>	arg <sub>2</sub>
*	c	d
+	b	(0)
+	(1)	(0)
=	(2)	

## System Programming and Compiler Construction

VI Semester ( Computer) Academic Year: 23-24

Triples face the problem of code immovability while optimization, as the results are positional and changing the order or position of an expression may cause problems.

### Indirect Triples

This representation is an enhancement over triples representation. It uses pointers instead of position to store results. This enables the optimizers to freely re-position the sub-expression to produce an optimized code.

#### Postlab Question

1. Write the intermediate code generated for ---- while ( a<b ) do

If ( c< d) then

X= y+z

Else

X= y-z

2. Write the intermediate code generated for ---- switch E

Begin

case  $V_1$  :  $S_1$

case  $V_2$  :  $S_2$

....

default:  $S_n$

end

#### Code:

```
import re
```

```
op = set('+ - / *')
```

```
address = 100
```

```
count = 0
```

```
def arithmetic(exp):
```

```
    global count
```

```
    symbols = []
```

```
operators = []

for i in exp:
    if i in op:
        operators.append(i)
    else:
        symbols.insert(0, i)

if "=" in symbols:
    while True:
        s = symbols.pop()
        if s == "=":
            break
        symbols.insert(0, s)

for i in operators:
    count += 1
    e = "temp{0} = {1} {2} {3}".format(
        count, symbols.pop(), i, symbols.pop())
    symbols.append("temp{}".format(count))
    print(e)

if len(symbols) != 2:
    return

temp = symbols.pop()
print("{} = {}".format(symbols.pop(), temp))
```

```
def relation(exp):

    global address

    tokens = re.split(r">|=|<|=|>|<", exp)

    operators = re.findall(r">|=|<|=|>|<", exp)

    print("{0} if {2} {3} {1} goto {4}".format(
        address, tokens.pop(), tokens.pop(), operators.pop(), address +
3))

    print("{} T := 0 ".format(address + 1))

    print("{} goto {}".format(address+2, address+4))

    print("{} T := 1".format(address + 3))

    address += 4

    print(address)


if __name__ == "__main__":

    while True:

        option = input(

            "1 Assignment\n2 Arithmetic\n3 Relation\n4 Exit\nEnter choice
: ")

        if option == "1":

            exp = input("Enter an expression : ")

            arithmetic(exp)

        if option == "2":

            exp = input("Enter an expression : ")

            arithmetic(exp)

        if option == "3":

            exp = input("Enter an expression : ")

            relation(exp)

        if option == "4":
```

System Programming and Compiler Construction

VI Semester ( Computer) Academic Year: 23-24

`break`

`print()`

```
● PS C:\Users\ivana\Desktop\College\Third Year\SEM 6\SPCC Pracs> python -u
  "c:\Users\ivana\Desktop\College\Third Year\SEM 6\SPCC Pracs\Experiment4
  \9601_e4.py"
1 Assignment
2 Arithmetic
3 Relation
4 Exit
Enter choice : 1
Enter an expression : x=i+v+a+n
temp1 = i + v
temp2 = temp1 + a
temp3 = temp2 + n
x = temp3

1 Assignment
2 Arithmetic
3 Relation
4 Exit
Enter choice : 2
Enter an expression : a=b+c*d
temp4 = b + c
temp5 = temp4 * d
a = temp5

1 Assignment
2 Arithmetic
3 Relation
4 Exit
Enter choice : 3
Enter an expression : i1<=i2
100 if i1 <= i2 goto 103
101 T := 0
102 goto 104
103 T := 1
104
```

Postlab

