

ACM ICPC TEAM REFERENCE - CONTENTS

Universidad de La Habana - UH

Iván Galbán Smith

CONTENTS

1. Misc	4
1.1. Template	4
1.2. istringstream	4
1.3. printf scanf	4
1.4. Cube	5
1.5. Josephus	6
1.6. Partition	6
1.7. Random	6
1.8. Useful	6
2. BitMask	8
2.1. Amount of Hamiltonian Walks	8
2.2. Existence of Hamiltonian Cycle	8
2.3. Existence of Hamiltonian Walk	9
2.4. Finding the Number of Simple Paths	9
2.5. Finding the Shortest Hamiltonian Cycle	10
2.6. Number of Hamiltonian Cycles	11
2.7. Number of Simple Cycles	11
2.8. Shortest Hamiltonian Walk	12
2.9. Subset Subset (3^n)	13
3. Data Structures	14
3.1. AVL Tree	14
3.2. Big Integer	15
3.3. Binary Heap	17
3.4. Disjoint Set	18
3.5. Fenwick Tree 1D	19
3.6. Fenwick Tree 2D	20
3.7. Fraction	20

3.8. Kd-Tree	21
3.9. Longest Common Ancestor. Sparse Table	23
3.10. Polynomial	24
3.11. Range Minimum Query Fast	25
3.12. Range Minimum Query	25
3.13. Range Minimum Sum Segment Query	26
3.14. Segment Tree Lazy Propagation	28
3.15. Segment Tree-1D Query	28
3.16. Segment Tree-2D	29
3.17. Treap	31
3.18. Treap Implicit Key	32
3.19. Trie	34
4. Dynamic Programming	36
4.1. Convex Hull Trick	36
4.2. Longest Increasing Subsequence	36
4.3. Matrix Chain	37
5. Geometry	38
5.1. Basic Operation	38
5.2. Circles	39
5.3. Closest Pair Points	40
5.4. Convex Cut	41
5.5. Convex Hull 3D	41
5.6. Lines	44
5.7. Minkowski	45
5.8. Point 3D	45
5.9. Polygon Triangulation	46
5.10. Rectilinear MST	47
5.11. Rotating Calipers	48
5.12. Segment Intersect	49

5.13. Segments	51	8.5. Tree Stern-Brocot	85
5.14. Semiplane Intersection	52	9. Numeric Methods	86
5.15. Triangles	53	9.1. Fast Fourier Transform	86
6. Graphs	55	9.2. Goldsection Search	86
6.1. Articulation Point And Bridge	55	9.3. Linear Recursion	87
6.2. Bellman-Ford	55	9.4. Romberg	87
6.3. Biconnected Components	56	9.5. Roots Newton	88
6.4. Bipartite Matching	57	9.6. Simplex	88
6.5. Centroid Decomposition	57	9.7. Simpson	89
6.6. Dijkstra	58	10. Parsing	91
6.7. Dominator Tree	58	10.1. Shunting Yard	91
6.8. Flow with Lower Bound	59	11. Sorting-Searching	93
6.9. Floyd Warshall	61	11.1. Ternary Searh	93
6.10. Gabow Edmonds	61	12. String	94
6.11. Gomory hu Tree	62	12.1. Aho Corasick	94
6.12. Hopcroft Karp	63	12.2. Knuth-Morris-Pratt	95
6.13. Hungarian	64	12.3. Longest Common Subsequence	95
6.14. Kruskal	64	12.4. Longest Palindrome Substring	96
6.15. Max Flow Dinic	65	12.5. Manacher	97
6.16. Max Flow push relabel	66	12.6. Maximal Suffix	97
6.17. Min Cost Max Flow	67	12.7. Minimum Rotation	97
6.18. Prim	68	12.8. Palindromic Tree	98
6.19. Satisfiability Two SAT	69	12.9. Substring Palindrome	99
6.20. Strongly Connected Components	70	12.10. Suffix Array	99
6.21. SCC Gabow	71	12.11. Suffix Automaton	100
6.22. Stoer Wagner	71	12.12. Z Function	101
6.23. Tree Isomorphism	72	13. Mathematical facts	103
7. Matrix	74	13.1. Números de Catalán	103
7.1. Gauss	74	13.2. Números de Stirling de primera clase	103
7.2. Gauss Modulo 2	74	13.3. Números de Stirling de segunda clase	103
7.3. Matrix Template	75	13.4. Números de Bell	103
8. Number Theory	78	13.5. Derangement	103
8.1. Binomial Coefficient	78	13.6. Números armónicos	103
8.2. Divisibility	78	13.7. Número de Fibonacci	103
8.3. ALL Number Theory	78	13.8. Sumas de combinatorios	104
8.4. Prime	83	13.9. Funciones generatrices	104

13.10.	The twelvefold way	104		13.13.	Ángulo entre dos vectores	105
13.11.	Teorema de Euler	105		13.14.	Proyección de un vector	105
13.12.	Burnside's Lemma	105				

1. Misc

1.1. Template.

```
#include <bits/stdc++.h>

using namespace std;

////////////////////////////////////
typedef long long Int;
typedef pair<int, int> pii;
typedef vector<int> vi;
////////////////////////////////////
#define REP(i,n) \
    for(int i=0;i<(int)n;++i)
#define FOR(i,n) \
    for(int i=1;i<=(int)n;++i)
#define ITR(c) __typeof((c).begin())
#define foreach(i,c) \
    for (ITR(c) i=(c).begin();i!=(c).end();++i)
#define ALL(c) \
    (c).begin(), (c).end()
#define DB(x) \
    cout << #x << "="_" << x << endl

#define endl '\n'
#define F first
#define S second
```

1.2. istringstream.

```
int main()
{
    int test;
    scanf("%d",&test);
    getchar();
    string line;
    for(int i=0; i<test; i++){
        getline(cin , line);
```

1.3. printf scanf.

```
char s[100];
```

```
#define pb push_back
#define mp make_pair
#define LEFT(n) ((n<<1) + 1)
#define RIGHT(n) ((n<<1) + 2)
#define BIT(n) (1<<n)
#define ONES(n) __builtin_popcount(n)
#define rightZero(n) __builtin_ctz(n); // trailing zeros
#define leftZero(n) __builtin_clz(n); // leading zeros

////////////////////////////////////
const double EPS = 1e-15;
const int oo = (1<<30);
const double PI = M_PI;
const int MOD = 1000000000 + 7;
////////////////////////////////////

int main()
{
    ios_base::sync_with_stdio(false);
    cin.tie(0);

    return 0;
}
```

```
    istringstream in(line);

    while(in>>line)
        cout<<line<<endl;
    }
    return 0;
}
```

```
scanf("%[aeiou]", s); //solo lee las vocales
```

```
scanf("%[^aeiou]", s); //solo lee las letras

scanf("%[^\n]",z); //funciona igual q gets()

//eliminar rayita de la fecha (5-29-2014) o (5/29/2014)
scanf("%d_%c_%d_%c_%d",&m,&d,&y);

printf("%09d\n",f); //imprime el entero f y rellena con 9 ceros

printf("%G\n",c); //imprime c sin ceros finales (convierte a E)

printf("%g\n",c); //imprime c sin ceros finales (convierte a e)

printf("%x\n",x); //imprime x en hexadecimal (Letras minusculas)

printf("%X\n",x); //imprime x en hexadecimal (Letras mayusculas)

printf("%o\n",o); //imprime o como octal unsigned

printf("%e\n",cient); //imprime el # en notacion cientifica (e minuscula)

printf("%E\n",cient); //imprime el # en notacion cientifica (E mayuscula)

//muestra un valor de apuntador en forma de puesta en marcha definida
printf("El_valor_de_Ptr_es_%p\n",ptr);

//Almacena el # char almacenados en el printf.
printf("Total_de_char_impresos_en_esta_linea_es:%n",&cant);
printf("_%d\n\n",cant);

printf("%%\n"); //muestra el caracter de porciento
```

```
printf("\\\n"); //muestra el caracter \

printf("\'\n"); //muestra el caracter '

printf("\n\n"); //muestra el caracter "

printf("\?\n"); //muestra el caracter ?

printf("\n\n\n"); //muestra el caracter \n

printf("%11d\n",123); //justifica a la derecha en 11

//7: ancho del campo 2:precision, valor 98.74 justificado derecha
printf("%.2f\n", 7 , 2 , 98.736);
//if precision < 0 ---> justificado izquierda

/*sprintf*/
char numstr[100];
int num=1200;

sprintf(numstr,"%d",num); // a decimal
printf("%s\n",numstr);

sprintf(numstr,"%X",num); // a hexadecimal en mayuscula
printf("%s\n",numstr);
//-----

int base=8;
cout<<setbase(base)<<8<<endl; //pone a cout a imprimir en base (0,8,10,16)

/* istringstream
```

1.4. Cube.

```
template<class T>
struct cube
{
    T F, U, D, L, R, B;

    void rotX()
    {
        swap(D, B);
        swap(B, U);
        swap(U, F);
    }
};
```

```
} // FUBD -> DFUB

void rotY()
{
    swap(D, R);
    swap(R, U);
    swap(U, L);
} // LURD -> DLUR

void rotZ()
{
    swap(F, B);
    swap(B, L);
    swap(L, R);
    swap(R, U);
    swap(U, D);
    swap(D, F);
}
```

```

swap(B, R);
swap(R, F);
swap(F, L);

```

1.5. Josephus.

```

/*
    Tested: ??????
*/

// n-cantidad de personas, m es la longitud del salto.
// comienza en la k-esima persona.
ll josephus(ll n, ll m, ll k)
{
    ll x = -1;
    for (ll i = n - k + 1; i <= n; ++i)
        x = (x + m) % i;
    return x;
}

```

1.6. Partition.

```

typedef long long ll;

ll partition(ll n)
{
    vector<ll> dp(n + 1);
    dp[0] = 1;
    for (int i = 1; i <= n; i++)

```

1.7. Random.

```

std::default_random_engine generator;

```

1.8. Useful.

```

// TIME
for (int a = 0; ; ++a) {
    if (clock() >= 2.5 * CLOCKS_PER_SEC) break;
    // It will stop when 2.5 seconds have passed
}

```

```

    } // LFRB -> BLFR
};

```

```

}

ll josephus_inv(ll n, ll m, ll x)
{
    for (ll i = n; i--;)
    {
        if (x == i)
            return n - i;
        x = (x - m % i + i) % i;
    }
    return -1;
}

```

```

        for (int j = 1, r = 1; i - (3 * j * j - j) / 2 >= 0; j++, r *= -1)
        {
            dp[i] += dp[i - (3 * j * j - j) / 2] * r;
            if (i - (3 * j * j + j) / 2 >= 0)
                dp[i] += dp[i - (3 * j * j + j) / 2] * r;
        }
        return dp[n];
}

```

```

std::uniform_real_distribution<double> distribution(0.0, 1.0);

```

```

// LAMBDA
function<bool(int, int)> add_edge = [&](int u, int v)
{
    // code here...
    return true;
}

```

```
};
```

```
// RANDOM DISTRIBUTIONS
```

```
std::default_random_engine generator;
```

```
std::uniform_real_distribution<double> distribution(0.0,1.0);
```

2. BITMASK

2.1. Amount of Hamiltonian Walks.

```

/*
    task: Finding the number of Hamiltonian walks in the
           unweighted and directed graph  $G = (V, E)$ .
    complexity:  $O(2^n * n^2)$ 
    notes: Let  $dp[mask][v]$  be the amount of Hamiltonian walks
           on the subgraph generated by vertices in  $mask$  that
           end in the vertex  $v$ .
*/
#define BIT(n) (1 << n)
const int MAXN = 20;

int n, m, u, v, g[MAXN], dp[BIT(MAXN)][MAXN], ans;

int main()
{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;
        g[u] |= BIT(v);
    }
}

```

```

for (int i = 0; i < n; ++i)
    dp[BIT(i)][i] = 1;

for (int msk = 1; msk < BIT(n); ++msk)
{
    for (int i = 0; i < n; ++i) if (msk & BIT(i))
    {
        int tmsk = msk ^ BIT(i);

        for (int j = 0; tmsk && j < n; ++j)
        {
            if (g[j] & BIT(i))
                dp[msk][i] += dp[tmsk][j];
        }
    }
}

for (int i = 0; i < n; ++i)
    ans += dp[BIT(n) - 1][i];
cout << ans << endl;
return 0;
}

```

2.2. Existence of Hamiltonian Cycle.

```

/*
    task: Check for existence of Hamiltonian cycle in a
           directed graph  $G = \langle V, E \rangle$ .

    complexity:  $O(2^n * n)$ 

    notes: Let  $dp[mask]$  be the mask of the subset consisting
           of those vertices  $j$  such that exist a Hamiltonian
           walk over the subset  $mask$  beginning in vertex 0 and
           ending in  $j$ .
*/
#define BIT(n) (1 << n)
const int MAXN = 20;
int n, m, u, v, g[MAXN], dp[BIT(MAXN)];

int main()

```

```

{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;
        g[v] |= BIT(u);
    }

    dp[1] = 1;

    for (int msk = 2; msk < BIT(n); ++msk)
    {
        for (int i = 0; i < n; ++i)
        {
            if ((msk & BIT(i)) && (dp[msk ^ BIT(i)] & g[i]))

```



```

        dp[msk] |= BIT(i);
    }
}

```

```

    cout << ((dp[BIT(n) - 1] & g[0]) != 0) << endl;
    return 0;
}

```

2.3. Existence of Hamiltonian Walk.

```

/*
    task: Check for existence of Hamiltonian walk in the
           directed graph  $G = \langle V, E \rangle$ .

    complexity:  $O(2^n * n)$ 

    notes: Let  $dp[msk]$  be the mask of the subset consisting
           of those vertices  $v$  for which exist a Hamiltonian
           walk over the subset  $msk$  ending in  $v$ .
*/

#define BIT(n) (1 << n)
const int MAXN = 20;

int n, m, u, v, g[MAXN], dp[BIT(MAXN)];

int main()
{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)

```

```

{
    cin >> u >> v;
    g[v] |= BIT(u);
}

for (int i = 0; i < n; ++i)
    dp[BIT(i)] = BIT(i);

for (int msk = 1; msk < BIT(n); ++msk)
{
    for (int i = 0; i < n; ++i)
    {
        if ((msk & BIT(i)) && (dp[msk ^ BIT(i)] & g[i]))
            dp[msk] |= BIT(i);
    }
}

cout << (dp[BIT(n) - 1] != 0) << endl;
return 0;
}

```

2.4. Finding the Number of Simple Paths.

```

/*
    task: Finding the number of simple paths in the
           directed graph  $G = \langle V, E \rangle$ .

    complexity:  $O(f)$ 

    notes: Let  $dp[msk][v]$  be the number of Hamiltonian
           walks in the subgraph generated by vertices
           in  $msk$  that end in  $v$ .
*/

#define BIT(n) (1 << n)
const int MAXN = 20;

```

```

int n, m, u, v, ans, g[MAXN], dp[BIT(MAXN)][MAXN];

int main()
{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;
        g[u] |= BIT(v);
    }

    for (int i = 0; i < n; ++i)
        dp[BIT(i)][i] = 1;

```

```

for (int msk = 1; msk < BIT(n); ++msk)
{
    for (int i = 0; i < n; ++i) if (BIT(i) & msk)
    {
        int tmsk = msk ^ BIT(i);

        for (int j = 0; tmsk && j < n; ++j) if (g[j] & BIT(i))

```

```

        dp[msk][i] += dp[tmsk][j];

        ans += dp[msk][i];
    }
}
cout << ans - n << endl;
return 0;
}

```

2.5. Finding the Shortest Hamiltonian Cycle.

```

/*
    task: Search for the shortest Hamiltonian cycle.
           Let the directed graph  $G = (V, E)$  have
            $n$  vertices, and each
           edge have weight  $d(i, j)$ . We want to find a Hamiltonian
           cycle for which the sum of weights of its edges
           is minimal.

    complexity:  $O(2^n * n^2)$ 

    notes: Let  $dp[msk][v]$  be the length of the shortest Hamiltonian
           walk on the subgraph generated by vertices in  $msk$ 
           beginning in vertex 0 and ending in vertex  $v$ .
*/

#define BIT(n) (1 << n)

using namespace std;

const int MAXN = 20,
        INF = 0xffffffff;

int n, m, u, v, w, g[MAXN][MAXN], dp[BIT(MAXN)][MAXN], ans = INF;

int main()
{
    cin >> n >> m;

    for (int i = 0; i < n; ++i)
    {
        for (int j = 0; j < n; ++j)
            g[i][j] = INF;
    }

```

```

    for (int i = 0; i < BIT(n); ++i)
    {
        for (int j = 0; j < n; ++j)
            dp[i][j] = INF;
    }

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;
        cin >> g[u][v];
    }

    dp[1][0] = 0;

    for (int msk = 2; msk < BIT(n); ++msk)
    {
        for (int i = 0; i < n; ++i) if (msk & BIT(i))
        {
            int tmsk = msk ^ BIT(i);

            for (int j = 0; tmsk && j < n; ++j)
                dp[msk][i] = min(dp[msk][i], dp[tmsk][j] + g[j][i]);
        }
    }

    for (int i = 1; i < n; ++i)
        ans = min(ans, dp[BIT(n) - 1][i] + g[i][0]);

    cout << ans << endl;
    return 0;
}

```

2.6. Number of Hamiltonian Cycles.

```

/*
    task: Finding the number of Hamiltonian cycles in
           the unweighted and directed graph  $G = (V, E)$ .

    complexity:  $O(2^n * n^2)$ 

    notes: Let  $dp[msk][v]$  be the amount of Hamiltonian walks
            on the subgraph generated by vertices in  $msk$ 
            that begin in vertex 0 and end in vertex  $v$ .
*/

#define BIT(n) (1 << n)
const int MAXN = 20;

int n, m, u, v, ans, g[MAXN], dp[BIT(MAXN)][MAXN];

int main()
{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;

```

```

        g[u] |= (1 << v);
    }

    dp[1][0] = 1;

    for (int msk = 2; msk < BIT(n); ++msk)
    {
        for (int i = 0; i < n; ++i) if (msk & BIT(i))
        {
            int tmsk = msk ^ BIT(i);

            for (int j = 0; tmsk && j < n; ++j) if (g[j] & BIT(i))
                dp[msk][i] += dp[tmsk][j];
        }
    }

    for (int i = 1; i < n; ++i) if (g[i] & 1)
        ans += dp[BIT(n) - 1][i];

    cout << ans << endl;
    return 0;
}

```

2.7. Number of Simple Cycles.

```

/*
    task: Finding the number of simple cycles in a
           directed graph  $G = \langle V, E \rangle$ .

    complexity:  $O(2^n * n^2)$ 

    notes: Let  $dp[msk][v]$  be the number of Hamiltonian
            walks in the subgraph generated by vertices
            in  $msk$  that begin in the lowest vertex in  $msk$ 
            and end in vertex  $v$ .
*/

#define BIT(n) (1 << n)
#define ONES(n) __builtin_popcount(n)

const int MAXN = 20;

```

```

int n, m, u, v, g[MAXN];
long long dp[BIT(MAXN)][MAXN], ans;

int main()
{
    cin >> n >> m;

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v;
        g[u] |= BIT(v);
    }

    for (int i = 0; i < n; ++i)
        dp[BIT(i)][i] = 1;

    for (int msk = 1; msk < BIT(n); ++msk)

```

```

{
    for (int i = 0; i < n; ++i)
    {
        if ((msk & BIT(i)) && !(msk & -msk & BIT(i)))
        {
            int tmsk = msk ^ BIT(i);

            for (int j = 0; tmsk && j < n; ++j) if (g[j] & BIT(i))
                dp[msk][i] += dp[tmsk][j];
        }
    }
}

```

```

        if (ONES(msk) > 2 && (g[i] & msk & -msk))
            ans += dp[msk][i];
    }
}
cout << ans << endl;
return 0;
}

```

2.8. Shortest Hamiltonian Walk.

```

/*
    task: Search for the shortest Hamiltonian walk.
           Let the directed graph  $G = (V, E)$  have  $n$ 
           vertices, and each edge have weight  $d(i, j)$ .
           We want to find a Hamiltonian walk for which
           the sum of weights of its edges is minimal.

    complexity:  $O(2^n * n^2)$ 

    notes: Let  $dp[msk][v]$  be the length of the shortest
           Hamiltonian walk on the subgraph generated by
           vertices in  $msk$  that end in vertex  $v$ .
*/

#define MAXN 20
#define INF 0xffffffff
#define BIT(n) (1 << n)

using namespace std;

int n, m, ans = INF, d[MAXN][MAXN], u, v, w, dp[1 << MAXN][MAXN];

int main()
{
    cin >> n >> m;

    for (int i = 0; i < n; ++i)
    {
        for (int j = 0; j < n; ++j)
            d[i][j] = INF;
    }
}

```

```

    for (int i = 0; i < BIT(n); ++i)
    {
        for (int j = 0; j < n; ++j)
            dp[i][j] = INF;
    }

    for (int i = 0; i < m; ++i)
    {
        cin >> u >> v >> w;
        d[u][v] = w;
    }

    for (int i = 0; i < n; ++i)
        dp[1 << i][i] = 0;

    for (int msk = 1; msk < (1 << n); ++msk)
    {
        for (int i = 0; i < n; ++i) if (msk & BIT(i))
        {
            int tmsk = msk ^ BIT(i);

            for (int j = 0; tmsk && j < n; ++j)
                dp[msk][i] = min(dp[tmsk][j] + d[j][i], dp[msk][i]);
        }
    }

    for (int i = 0; i < n; ++i)
        ans = min(ans, dp[BIT(n) - 1][i]);

    cout << ans << endl;
    return 0;
}

```

2.9. Subset Subset (3^n).

```
/*  
    Computing all subset of subset.  
    Time:  $3^n$   
*/  
#include <bits/stdc++.h>  
using namespace std;  
  
int main()  
{  
    int N = 4;  
    for(int i=0; i<(1<<N); ++i){
```

```
        bitset<8> n(i);  
        cout<<"MASK:␣"<<n<<endl;  
        cout<<"SUBMASK:"<<endl;  
        for(int j = i; j; j = (j-1) & i){  
            bitset<8> p(j);  
            cout<<p<<endl;  
        }  
        cout<<endl;  
    }  
    return 0;  
}
```

3. DATA STRUCTURES

3.1. AVL Tree.

```

/*
    Coding an AVL Tree

    Remarks: Assuming keys are integers. The data structure does
             not allows duplicate keys.
    Performance:
        Insert: O(log n)
        Erase: O(log n)
        Contains: O(log n)
        Find minimum: O(log n)
        Find maximum: O(log n)
        Find k-th: O(log n)
*/

struct AVL_Tree
{
    struct node
    {
        int key;
        int size, height;
        node *ch[2];

        int balance_factor() { return ch[1]->height - ch[0]->height; }
        void update()
        {
            height = 1 + max( ch[0]->height, ch[1]->height );
            size = ch[0]->size + ch[1]->size + 1;
        }
    } *root, *null;

    int key;

    node* new_node( const int& key ) {
        node *x = new node();
        x->key = key;
        x->height = x->size = 1;
        x->ch[0] = x->ch[1] = null;
        return x;
    }

    node* rotate( node *x, bool b ) {

        if ( x == null || x->ch[ !b ] == null ) return x;

```

```

        node *y = x->ch[ !b ];
        x->ch[ !b ] = y->ch[b];
        y->ch[b] = x;

        x->update();
        y->update();

        return y;
    }

    node* balance( node *x ) {

        x->update();

        if ( x->balance_factor() > 1 ) {
            if ( x->ch[1]->balance_factor() <= 0 )
                x->ch[1] = rotate( x->ch[1], 1 );
            x = rotate( x, 0 );
        } else
            if ( x->balance_factor() < -1 ) {
                if ( x->ch[0]->balance_factor() >= 0 )
                    x->ch[0] = rotate( x->ch[0], 0 );
                x = rotate( x, 1 );
            }

        x->update();
        return x;
    }

    node* insert( node *x, const int& key ) {

        if ( x == null ) x = new_node( key );
        else {

            if ( key == x->key ) return x;

            bool b = !( key < x->key );
            x->ch[b] = insert( x->ch[b], key );

            x = balance( x );
        }

```

```

        return x = balance( x );
    }

    node* erase( node *x, int key ) {

        if ( x == null ) return x;

        int tmp = x->key;
        if ( tmp == x->key ) {
            if ( x->ch[0] == null || x->ch[1] == null )
                return x->ch[ x->ch[0] == null ];
            else {
                node *p = x->ch[0];
                while ( p->ch[1] != null ) p = p->ch[1];
                x->key = p->key;
                key = p->key;
            }
        }

        bool b = !( key < tmp );
        x->ch[b] = erase( x->ch[b], key );

        return x = balance( x );
    }

    bool contains( node *root, const int& key ) {
        node *x = root;
        for ( ;; ) {
            if ( x == null ) return 0;
            if ( key == x->key ) return 1;
            x = x->ch[ !( key < x->key ) ];
        }
    }

    int get_extreme( bool b ) {
        assert( root != null );

```

3.2. Big Integer.

```

typedef long long Int;
const Int B = 10; // base (power of 10)
const int BW = 1; // log B
const int MAXDIGIT = 100; // it can represent 4 * MAXDIGIT digits (in base 10)

struct BigNum {
    Int digit [MAXDIGIT];

```

```

        node *x = root;
        while ( x->ch[b] != null ) x = x->ch[b];
        return x->key;
    }

    int find_kth( node *root, int k ) {

        assert( root->size >= k );

        node *x = root;
        for ( ;; ) {
            int rank = x->ch[0]->size + 1;
            if ( rank == k ) return x->key;
            if ( k < rank )
                x = x->ch[0];
            else x = x->ch[1], k -= rank;
        }
    }

    /* "Public" methods */

    void insert( int x ) { root = insert( root, key = x ); }
    void erase( int x ) { root = erase( root, x ); }
    bool contains( int x ) { return contains( root, key = x ); }
    int get_min() { return get_extreme( 0 ); }
    int get_max() { return get_extreme( 1 ); }
    int find_kth( int k ) { return find_kth( root, k ); }

    AVL_Tree()
    {
        null = new node();
        null->height = null->size = 0;
        null->ch[0] = null->ch[1] = 0;
        root = null;
    }
};

```

```

    int size;
    BigNum (int size = 1, Int a = 0): size (size) {
        memset (digit, 0, sizeof (digit));
        digit [0] = a;
    }
};
const BigNum ZERO (1, 0), ONE (1, 1);

```

```
// Comparators
bool operator < (BigNum x, BigNum y) {
    if (x.size != y.size) return x.size < y.size;
    for (int i = x.size-1; i >= 0; --i)
        if (x.digit[i] != y.digit[i]) return x.digit[i] < y.digit[i];
    return false;
}
bool operator > (BigNum x, BigNum y) {return y < x;}
bool operator <= (BigNum x, BigNum y) {return !(y < x);}
bool operator >= (BigNum x, BigNum y) {return !(x < y);}
bool operator != (BigNum x, BigNum y) {return x < y || y < x;}
bool operator == (BigNum x, BigNum y) {return !(x < y) && !(y < x);}

BigNum normal (BigNum x) {
    Int c = 0;
    for (int i = 0; i < x.size; ++i) {
        while (x.digit[i] < 0)
            x.digit[i + 1] -= 1, x.digit[i] += B;
        Int a = x.digit[i] + c;
        x.digit[i] = a % B;
        c = a / B;
    }
    for (; c > 0; c /= B) x.digit[x.size++] = c % B;
    while (x.size > 1 && x.digit[x.size-1] == 0) --x.size;
    return x;
}

BigNum convert (Int a) {
    return normal (BigNum (1, a));
}

BigNum convert (const string & s) {
    BigNum x;
    int i = s.size() % BW;
    if (i > 0) i -= BW;
    for (; i < (int)s.size(); i += BW) {
        Int a = 0;
        for (int j = 0; j < BW; ++j)
            a = 10 * a + (i + j >= 0 ? s[i + j] - '0' : 0);
        x.digit[x.size++] = a;
    }
    reverse (x.digit, x.digit + x.size);
    return normal (x);
}

// Input / Output
```

```
ostream& operator <<(ostream &os, BigNum x) {
    os << x.digit[x.size-1];
    for(int i = x.size-2; i >= 0; --i)
        os << setw(BW) << setfill ('0') <<x.digit[i];
    return os;
}
istream& operator >>(istream &is, BigNum &x) {
    string s; is>> s;
    x = convert(s);
    return is;
}

// Basic Operations
BigNum operator + (BigNum x, BigNum y) {
    if (x.size < y.size) x.size = y.size;
    for (int i = 0; i < y.size; ++i)
        x.digit[i] += y.digit[i];
    return normal(x);
}

BigNum operator - (BigNum x, BigNum y) {
    assert(x >= y);
    for (int i = 0; i < y.size; ++i)
        x.digit[i] -= y.digit[i];
    return normal(x);
}

BigNum operator * (BigNum x, BigNum y) {
    BigNum z (x.size + y.size);
    for (int i = 0; i < x.size; ++i)
        for (int j = 0; j < y.size; ++j)
            z.digit[i + j] += x.digit[i] * y.digit[j];
    return normal(z);
}

BigNum operator * (BigNum x, Int a) {
    for (int i = 0; i < x.size; ++i)
        x.digit[i] *= a;
    return normal (x);
}

pair <BigNum, Int> divmod (BigNum x, Int a) {
    Int c = 0, t;
    for (int i = x.size-1; i >= 0; --i) {
        t = B * c + x.digit[i];
        x.digit[i] = t / a;
        c = t % a;
    }
}
```



```

    }
    return pair <BigNum, Int> (normal (x), c);
}

BigNum operator / (BigNum x, Int a) {
    return divmod (x, a).first;
}

Int operator % (BigNum x, Int a) {
    return divmod (x, a).second;
}

pair <BigNum, BigNum> divmod (BigNum x, BigNum y) {
    if (x.size < y.size) return pair <BigNum, BigNum> (ZERO, x);
    int F = B / (y.digit[y.size-1] + 1); // multiplying good-factor
    x = x * F; y = y * F;
    BigNum z (x.size - y.size + 1);
    for (int k = z.size - 1, i = x.size - 1; k >= 0; --k, --i) {
        z.digit[k] = (i + 1 < x.size ? x.digit[i + 1] : 0) * B + x.digit[i];
        z.digit[k] /= y.digit[y.size-1];
        BigNum t (k + y.size);
        for (int m = 0; m < y.size; ++m)
            t.digit[k + m] = z.digit[k] * y.digit[m];
        t = normal(t);
        while (x < t) {
            z.digit[k] -= 1;
            for (int m = 0; m < y.size; ++m)
                t.digit[k + m] -= y.digit[m];
            t = normal (t);
        }
        x = x - t;
    }
    return pair <BigNum, BigNum> (normal(z), x / F);
}

BigNum operator / (BigNum x, BigNum y) {

```

```

    return divmod (x, y).first;
}

BigNum operator % (BigNum x, BigNum y) {
    return divmod (x, y).second;
}

// Advanced Operations
BigNum shift (BigNum x, int k) {
    if (x.size == 1 && x.digit[0] == 0) return x;
    x.size += k;
    for (int i = x.size - 1; i >= k; --i) x.digit[i] = x.digit [i + k];
    for (int i = k-1; i >= 0; --i) x.digit[i] = 0;
    return x;
}

BigNum sqrt (BigNum x) { // verified UVA 10023
    const BigNum _20 = convert(2 * B);
    BigNum odd = ZERO;
    BigNum rem (2, 0);
    BigNum ans = ZERO;
    for (int i = 2 * ((x.size-1) / 2); i >= 0; i -= 2) {
        int group = (i + 1 < x.size ? x.digit [i + 1]: 0) * B + x.digit [i];
        odd = _20 * ans + ONE;
        rem = shift (rem, 2) + convert (group);
        int count = 0;
        while (rem >= odd) {
            count = count + 1;
            rem = rem - odd;
            odd.digit[0] += 2;
            odd = normal(odd);
        }
        ans = shift (ans, 1) + convert (count);
    }
    return ans;
}

```

3.3. Binary Heap.

```

int oo = (1<<30);
int N, heap_size;

//O(log n)
void max_heapify(int *A, int i)
{
    int l, r, largest = i;

```

```

do{
    i = largest;
    l = (i<<1) + 1;
    r = (i<<1) + 2;

    if(l < heap_size && A[l] > A[largest]) largest = l;
    if(r < heap_size && A[r] > A[largest]) largest = r;

```

```

        swap(A[largest], A[i]);
    }while(largest != i);
}

//O(1)
int parent(int i)
{
    return (i-1) / 2;
}

//O(log n)
void max_heapifyUp(int *A, int i)
{
    while(i >= 0 && A[i] > A[parent(i)])
    {
        swap(A[i], A[parent(i)]);
        i = parent(i);
    }
}

//O(n)
void build_max_heap(int *A)
{
    heap_size = N;
    for(int i = N/2; i >= 0; --i)
        max_heapify(A, i);
}

//O(1)

```

3.4. Disjoint Set.

```

int N;
int parent[N], cont[N];

void initSet()
{
    for(int i = 0; i < N; ++i){
        parent[i] = i;
        cont[i] = 1;
    }
}

int SetOf(int x)
{

```

```

int max_heap(int *A)
{
    return A[0];
}

//O(log n)
int heap_extract_max(int *A)
{
    if(heap_size < 1)
        return 0;

    int max = A[0];

    swap(A[0], A[heap_size-1]);
    --heap_size;

    max_heapify(A, 0);

    return max;
}

//O(log n)
void heap_increase_key(int *A, int i, int key)
{
    if(key <= A[i])
        return;

    A[i] = key;
    max_heapifyUp(A, i);
}

```

```

        return (x == parent[x]) ? x : parent[x] = SetOf(parent[x]);
    }

void Merge(int x, int y)
{
    x = SetOf(x);
    y = SetOf(y);

    if(x == y)
        return;

    if(cont[x] < cont[y])
        swap(x, y);

```

```
parent[y] = x;
```

3.5. Fenwick Tree 1D.

```
/*
 * Performance:
 * 0-based
 * To start the index on 1
 * lowbit --> O(1)
 * query --> O(log N)
 * update --> O(log N)
 */
template<class T> struct abi{

    vector<T> ft;
    abi(int n):ft(n+1, 0){}

    int lowbit(int x) {return x & -x;}

    //item[pos] += val
    void update(int pos, T val)
    {
        for(; pos <= (int)ft.size(); pos += lowbit(pos))
            ft[pos] += val;
    }

    // Give sum[0...pos]
    T query(int pos)
    {
        T sum = 0;
        for(; pos > 0; pos -= lowbit(pos))
            sum += ft[pos];
        return sum;
    }

    // Give sum[l...r]
    T query(int l, int r)
    {
        l = (l > 0) ? l-1 : 0;
        return query(r) - query(l);
    }
};
```

```
cont[x] += cont[y];
}
```

```
int highestOneBit(int n)
{
    int shift = 31-(__builtin_clz(n));
    int ans = 1;
    ans <= shift;
    return ans;
}

//Return min(p/sum[0,p]>=sum)
int lower_bound(int sum)
{
    --sum;
    int pos = 0;
    for(int blockSize = highestOneBit(ft.size()); blockSize; blockSize >>= 1)
    {
        int nextPos = pos + blockSize;
        if(nextPos <= (int)ft.size() && sum >= ft[nextPos])
        {
            sum -= ft[nextPos];
            pos = nextPos;
        }
    }
    return pos + 1;
}

// number of free places in [0, x]
int getZeros(int x) {
    return x < 0 ? 0 : x + 1 - query(x);
}

int getZeros(int x1, int x2) {
    int s = getZeros(x2) - getZeros(x1 - 1);
    return x1 <= x2 ? s : s + getZeros(ft.size() - 1);
}

};
```

3.6. Fenwick Tree 2D.

```

/*
 * Performance:
 * 0-based
 * To start the index on 1
 * lowbit --> O(1)
 * query --> O( log (N+M) )
 * update --> O( log (N+M) )
 */

//Tested 1904 - Again Making Queries III COJ
#define MOD 10000
#define MaxN 4005
int N, U, Q;
int ft[MaxN][MaxN];

int lowbit(int x){return x & -x;}

bool Valid(int r, int c)
{
    if(r < 1 || r > N) return false;
    if(c < 1 || c > N) return false;
    return true;
}

```

3.7. Fraction.

```

template<class T>
struct fraction{
    T n, d;

    fraction()
    {
        n = 0;
        d = 1;
    }

    fraction(T _n, T _d)
    {
        n = _n;
        d = _d;
    }
};

template<class T>

```

```

}

void update(int r, int c, int val)
{
    if(!Valid(r, c)) return;
    for(int i = r; i <= N; i += lowbit(i))
        for(int j = c; j <= N; j += lowbit(j))
            ft[i][j] += val;
}

int query(int r, int c)
{
    if(!Valid(r, c)) return 0;
    int sum = 0;
    for(int i = r; i > 0; i -= lowbit(i))
        for(int j = c; j > 0; j -= lowbit(j))
            sum += ft[i][j];
    return sum;
}

int query(int r, int c, int R, int C)
{
    return query(R, C) - query(R, c-1) - query(r-1, C) + query(r-1, c-1);
}

```

```

fraction<T> operator +(fraction<T> &a, fraction<T> &b)
{
    T mcm = a.d * b.d;
    return fraction<T>(mcm/a.d*a.n + mcm/b.d*b.n, mcm);
}

template<class T>
fraction<T> operator *(fraction<T> &a, fraction<T> &b)
{
    return fraction<T>(a.n*b.n, a.d*b.d);
}

template<class T>
istream& operator >>(istream &in, fraction<T> &frac)
{
    in >> frac.n >> frac.d;
    return in;
}

```

```
template<class T>
ostream& operator <<(ostream &out, fraction<T> &frac)
{
    out << frac.n << "/" << frac.d;
    return out;
}
```

3.8. Kd-Tree.

```
/*
TASK : Coding a kd-tree

Remarks: The data structure is used in this code to
answer 2D range queries on a set of n 2D
points of the type "report all points inside
a rectangle [a,b]x[c,d]". The points' coordinates
are assumed to be integers.

Performance:
Build kd-tree: O(n log n)*
Query: O(sqrt(n) + k)
k: number of points inside query region

* expected
*/
#define MAXN 10000
#define oo 10000000000

struct point { int x, y; };
struct region { int xlo, xhi, ylo, yhi; };

struct node {
    point p;
    node *l, *r;
    region R;
    node( point p, node *l, node *r, int xlo, int xhi, int ylo, int yhi ) :
        p(p), l(l), r(r) { R = ( region ) { xlo, xhi, ylo, yhi }; }
} *root;

int N, Q;
int xlo, ylo;
int xhi, yhi;
region R;

point p[MAXN];
```

```
template<class T>
bool operator <(fraction<T> a, fraction<T> b)
{
    return a.n*b.d < a.d*b.n;
}
```

```
inline bool leaf( node *x ) { return !x->l && !x->r; }
inline bool less_than( const point& a, const point& b, bool byX ) {
    return byX ? a.x < b.x : a.y < b.y;
}

void partition( point a[], int lo, int hi, const int& k, bool byX ) {

    int l = lo, r = hi - 1, mid = ( lo + hi ) >> 1;

    if ( less_than( a[mid], a[lo], byX ) ) swap( a[mid], a[lo] );
    if ( less_than( a[hi], a[lo], byX ) ) swap( a[hi], a[lo] );
    if ( less_than( a[hi], a[mid], byX ) ) swap( a[hi], a[mid] );

    if ( hi - lo + 1 <= 3 ) return;

    swap( a[mid], a[ hi - 1 ] );
    point pivot = a[ hi - 1 ];

    for ( ;; ) {
        while ( less_than( a[ ++l ], pivot, byX ) );
        while ( less_than( pivot, a[ --r ], byX ) );
        if ( l < r ) swap( a[l], a[r] );
        else break;
    }

    swap( a[l], a[ hi - 1 ] );

    if ( k < l ) partition( a, lo, l - 1, k, byX );
    if ( k > l ) partition( a, l + 1, hi, k, byX );
}

node* build_kd_tree( point p[], int len, int depth,
                    int xlo, int xhi, int ylo, int yhi ) {

    if ( len == 1 )
        return new node( p[0], 0, 0, p[0].x, p[0].x, p[0].y, p[0].y );
```

```

int mid = ( len - 1 ) / 2;
partition( p, 0, len - 1, mid, !( depth & 1 ) );

int c1 = 0, c2 = 0;
point p1[MAXN], p2[MAXN];
for ( int i = 0; i <= mid; i++ ) p1[ c1++ ] = p[i];
for ( int i = mid + 1; i < len; i++ ) p2[ c2++ ] = p[i];

int xlo1 = xlo, xhi1 = xhi, ylo1 = ylo, yhi1 = yhi,
    xlo2 = xlo, xhi2 = xhi, ylo2 = ylo, yhi2 = yhi;

if ( !( depth & 1 ) )
    xhi1 = p[mid].x, xlo2 = p[mid].x + 1;
else yhi1 = p[mid].y, yhi2 = p[mid].y + 1;

node *left = build_kd_tree( p1, mid + 1, depth + 1,
                           xlo1, xhi1, ylo1, yhi1 );
node *right = build_kd_tree( p2, len - mid - 1, depth + 1,
                             xlo2, xhi2, ylo2, yhi2 );

return new node( p[mid], left, right, xlo, xhi, ylo, yhi );
}

void report( node *t ) {
    if ( !t ) return;
    if ( leaf( t ) )
        printf( "(%d,%d)_", t->p.x, t->p.y );
    else {
        report( t->l );
        report( t->r );
    }
}

region make_region( node *t ) {
    return ( region ) { t->R.xlo, t->R.xhi, t->R.ylo, t->R.yhi };
}

bool contained( const region& a, const region& b ) {
    return ( b.xlo <= a.xlo && a.xlo <= b.xhi &&
            b.xlo <= a.xhi && a.xhi <= b.xhi &&
            b.ylo <= a.ylo && a.ylo <= b.yhi &&
            b.ylo <= a.yhi && a.yhi <= b.yhi );
}

```

```

bool intersect( const region& a, const region& b ) {
    bool okX = ( ( a.xlo <= b.xlo && b.xlo <= a.xhi ) ||
                ( a.xlo <= b.xhi && b.xhi <= a.xhi ) );
    bool okY = ( ( a.ylo <= b.ylo && b.ylo <= a.yhi ) ||
                ( a.ylo <= b.yhi && b.yhi <= a.yhi ) );
    return okX && okY;
}

void query( node *t, const region& R ) {

    if ( leaf( t ) ) {
        if ( contained( t->R, R ) ) report( t );
    }
    else {

        region lc = make_region( t->l );
        if ( contained( lc, R ) ) report( t->l );
        else if ( intersect( lc, R ) ) query( t->l, R );

        region rc = make_region( t->r );
        if ( contained( rc, R ) ) report( t->r );
        else if ( intersect( rc, R ) ) query( t->r, R );
    }
}

int main() {

    scanf( "%d", &N );
    for ( int i = 0; i < N; i++ )
        scanf( "%d_%d", &p[i].x, &p[i].y );

    root = build_kd_tree( p, N, 0, -oo, oo, -oo, oo );

    for ( scanf( "%d", &Q ); Q--; ) {
        scanf( "%d_%d_%d_%d", &xlo, &ylo, &xhi, &yhi );
        R = ( region ) { xlo, xhi, ylo, yhi };
        query( root, R );
        printf( "\n" );
    }

    return 0;
}

```

3.9. Longest Common Ancestor. Sparse Table.

```

/*
    TASK : LCA Problem using DP
    Performance:
        Preprocess logarithms --> O(V)
        Build tree --> O(V)
        buildSparseTable --> O(V log V)
        queryLCA --> O(log V)
*/
#define LOGV 16
#define MAXV 1 << LOGV

using namespace std;

struct Node {
    int v, next;
} L[MAXV];

int V;
int P[MAXV];
int level[MAXV], parent[MAXV];
int LCA[MAXV][LOGV];

void readTree()
{
    for(int i = 0; i < V-1; ++i)
    {
        int u, v; cin >> u >> v;
        --u;--v;

        L[2 * i] = (Node) {v, P[u]};
        P[u] = 2 * i;

        L[2 * i + 1] = (Node) {u, P[v]};
        P[v] = 2 * i + 1;
    }
}

void buildSparseTable()
{
    queue<int> Q;
    level[0] = 0; parent[0] = -1;

    for(Q.push(0); !Q.empty(); Q.pop())
    {

```

```

        int u = Q.front();
        for(int i = P[u]; i != -1; i = L[i].next)
        {
            int v = L[i].v;
            if(v == parent[u]) continue;

            parent[v] = u;
            level[v] = level[u] + 1;
            Q.push(v);

            //DP
            LCA[v][0] = u;
            for(int j = 1; j <= __lg(level[v]); ++j)
                LCA[v][j] = LCA[ LCA[v][j - 1] ][j - 1];
        }
    }

    int queryLCA(int u, int v)
    {
        if(level[u] < level[v]) swap(u, v);

        if(level[u] != level[v])
            for(int i = __lg(level[u]); i >= 0; --i)
                if(level[u] - (1 << i) >= level[v])
                    u = LCA[u][i];

        if(u == v) return u;

        for(int i = __lg(level[u]); i >= 0; --i)
            if(level[u] - (1 << i) >= 0 && LCA[u][i] != LCA[v][i]){
                u = LCA[u][i];
                v = LCA[v][i];
            }
        return parent[u];
    }

    void init()
    {
        memset(P, -1, sizeof(P));
        readTree();
        buildSparseTable();
    }
}

```

3.10. Polynomial.

```

template<class T>
struct polynomial{
    int deg;
    vector<T> coef;

    polynomial(){}

    polynomial(int _deg)
    {
        deg = _deg;
        coef = vector<T>(deg + 1, 0);
    }

    polynomial(int _deg, vector<T> _coef)
    {
        deg = _deg;
        coef = _coef;
    }

    T eval(double x)
    {
        T y = 0;
        double pow = 1;
        for(int i = 0; i <= deg; ++i)
        {
            y = y + coef[i]*pow;
            pow = pow * x;
        }
        return y;
    }
};

template<class T>
istream& operator >>(istream& in, polynomial<T> &pol)
{
    in >> pol.deg;
    pol.coef = vector<T>(pol.deg + 1);

    for(int i = 0; i <= pol.deg; ++i)
        in >> pol.coef[i];
    return in;
}

void literal(ostream& out, int i)

```

```

{
    if(i == 0)
        return;
    if(i == 1)
    {
        out << "x";
        return;
    }
    out << "x^" << i;
}

template<class T>
ostream& operator <<(ostream& out, polynomial<T> pol)
{
    bool first = true;
    for(int i = pol.deg; i > 0; --i)
    {
        if(pol.coef[i] != 0)
        {
            if(first)
            {
                if(pol.coef[i] != 1 && pol.coef[i] != -1)
                    out << pol.coef[i];
                else if(pol.coef[i] == -1)
                    out << "-";
            }
            else
            {
                if(pol.coef[i] == 1)
                    out << "+";
                else if(pol.coef[i] == -1)
                    out << "-";
                else if(pol.coef[i] > 0)
                    out << "+" << pol.coef[i];
                else
                    out << pol.coef[i];
            }

            if(i == 1)
                out << "x";
            else if(i > 1)
                out << "x^" << i;

            first = false;
        }
    }
}

```



```

    }

    if(first)
    {
        out << pol.coef[0];
        return out;
    }
    else
    {
        if(pol.coef[0] != 0)
        {
            if(pol.coef[0] > 0)
                out << "+";
            out << pol.coef[0];
        }

        return out;
    }

    template<class T>
    polynomial<T> operator +(polynomial<T> &a, polynomial<T> &b)

```

3.11. Range Minimum Query Fast.

```

struct RMQ{
    vector<int> rmq;int n;
    RMQ(vector<int> &a){n=a.size();buildRMQ(a);}
    void buildRMQ(vector<int> &a){
        int logn=1;
        for(int k=1;k<n;k<=<=1)
            ++logn;
        rmq=vector<int>(n*logn);
        vector<int>::iterator b=rmq.begin();
        copy(ALL(a),b);
        for(int k=1;k<n;k<=<=1){
            copy(b,b+n,b+n);b+=n;

```

3.12. Range Minimum Query.

```

/*
    Start in 0.
    TASK : Range Minimum Query Problem: Given a sequence S of real numbers,
    RMQ(i,j) returns the index of element in S[i...j] with
    smallest value.

```

```

{
    polynomial<T> sum;

    if(a.deg >= b.deg)
        sum = a;
    else
        sum = b;

    for(int i = 0; i <= min(a.deg, b.deg); ++i)
        sum.coef[i] = a.coef[i] + b.coef[i];
    return sum;
}

template<class T>
polynomial<T> operator *(polynomial<T> &p1, polynomial<T> &p2)
{
    polynomial<T> mult(p1.deg + p2.deg);
    for(int i = 0; i <= p1.deg; ++i)
        for(int j = 0; j <= p2.deg; ++j)
            mult.coef[i + j] = mult.coef[i + j] + p1.coef[i] * p2.coef[j];
    return mult;
}

REP(i,n-k)b[i]=min(b[i],b[i+k]);}
}
int minimum(int x,int y){
    int z=y-x,k=0,e=1,s;//y-x>=e*2^k k up to a
    s=((z&0xffff0000)!=0)<<4;z>=s;e<=s;k|=s;
    s=((z&0x0000ffff)!=0)<<3;z>=s;e<=s;k|=s;
    s=((z&0x000000ff)!=0)<<2;z>=s;e<=s;k|=s;
    s=((z&0x0000000c)!=0)<<1;z>=s;e<=s;k|=s;
    s=((z&0x00000002)!=0)<<0;z>=s;e<=s;k|=s;
    return min(rmq[x+n*k],rmq[y+n*k-e+1]);
}
};

```

```

Preprocess Sparse Table --> O(N log N)
Answer query --> O(1)
*/

```

```
//Tested 1651 - Finding Minimum COJ
// 1082 - Array Queries Lightoj

const int Max = 10005,
        MaxLog = 15;

int N;
int rmq[Max][MaxLog], array[Max];

void build()
{
    for(int i = 0; i < N; ++i)
        rmq[i][0] = i;
    for(int i = 1; (1<<i) <= N; ++i)
```

```
        for(int j = 0; j + (1<<i) <= N; ++j){
            if(array[ rmq[j][i-1] ] < array[ rmq[j + (1 << (i-1))][i-1] ] )
                rmq[j][i] = rmq[j][i-1];
            else
                rmq[j][i] = rmq[j + (1 << (i-1))][i-1];
        }
    }

int query(int l, int r)
{
    int k = __lg(r-l+1);
    return array[rmq[l][k]] < array[ rmq[r - (1<<k) + 1][k] ] ? rmq[l][k] :
        rmq[r - (1<<k) + 1][k];
}
```

3.13. Range Minimum Sum Segment Query.

```
/*
    TASK : -Range Minimum-Sum Segment Query Problem
           -With two intervals too.
    Compute arrays C, P and M --> O(N)
    Preprocess RMQ --> O(N log N)
    Answer RMSQ queries --> O(1)
*/
#define MAXN 50005
#define LGN 16

int A[MAXN];
int C[MAXN], P[MAXN], M[MAXN], L[MAXN];
int RMQ[MAXN][LGN][2];

//for two intervals
int rmqMAXC[MAXN][LGN];

int N;

// Compute arrays C, P, L and M
//C[i] = sum(A[1]...A[i])
//L[i] = max{k | C[k] >= C[i] k[1, i-1]}
// {0 otherwise }
//P[i] = max{k | k[L[i]+1, i] and C[k-1] <= C[1] forall l[L[i], i-1]}
//M[i] = sum(P[i], i)
void buildCLPM()
{
    for (int i = 1; i <= N; ++i ) {
        C[i] = C[ i - 1 ] + A[i];
```

```
        L[i] = i - 1; P[i] = i;
        while ( C[ L[i] ] < C[i] && L[i] ) {
            if ( C[ P[ L[i] ] - 1 ] < C[ P[i] - 1 ] )
                P[i] = P[ L[i] ];
            L[i] = L[ L[i] ];
        }
        M[i] = C[i] - C[ P[i] - 1 ];
    }
}

// Preprocess array C for RMQmin and array M for RMQmax
// RMQ[i][j][0] holds the minimum, while RMQ[i][j][1] holds
// the maximum
void buildRMQ()
{
    for (int i = 0; i <= N; ++i )
        RMQ[i][0][0] = RMQ[i][0][1] = i;

    for (int j = 1; j <= __lg( N + 1 ); ++j )
        for (int i = 0; i + ( 1 << j ) - 1 <= N + 1; ++i ) {
            if ( C[ RMQ[i][ j - 1 ][0] ] <= C[ RMQ[ i +
                ( 1 << ( j - 1 ) ) ][ j - 1 ][0] ] )
                RMQ[i][j][0] = RMQ[i][ j - 1 ][0];
            else RMQ[i][j][0] = RMQ[ i + ( 1 << ( j - 1 ) ) ]
                [ j - 1 ][0];
            if ( M[ RMQ[i][ j - 1 ][1] ] >= M[ RMQ[ i +
                ( 1 << ( j - 1 ) ) ][ j - 1 ][1] ] )
                RMQ[i][j][1] = RMQ[i][ j - 1 ][1];
            else RMQ[i][j][1] = RMQ[ i +
```

```

        ( 1 << ( j - 1 ) ) ][ j - 1 ][1];
    }

    int queryRMQ( int l, int r, int b ) {

        int k = __lg( r - l + 1 );

        //For two Intervals
        if(b == 2) return max(rmqMAXC[l][k], rmqMAXC[r - (1<<k) + 1][k]);

        if ( !b ) return C[ RMQ[l][k][b] ] <= C[ RMQ[ r - ( 1 << k ) + 1 ][k][b] ] ?
            RMQ[l][k][b] : RMQ[ r - ( 1 << k ) + 1 ][k][b];
        else return M[ RMQ[l][k][b] ] >= M[ RMQ[ r - ( 1 << k ) + 1 ][k][b] ] ?
            RMQ[l][k][b] : RMQ[ r - ( 1 << k ) + 1 ][k][b];
    }

    pair<int, int> queryRMSQ( int l, int r )
    {
        int x = queryRMQ( l, r, 1 );
        if ( P[x] < 1 ) {
            int y = queryRMQ( x + 1, r, 1 );
            int z = queryRMQ( l - 1, x - 1, 0 ) + 1;
            if ( C[x] - C[ z - 1 ] < M[y] )
                return pair<int, int>(P[y], y);
            return pair<int, int>(z, x);
        }
        return pair<int, int>(P[x], x);
    }

    //RMSQ with two intervals
    //Return i <= x <= j, k <= y <= l
    // max{ Sum(x, y) }

    void buildRMSQ2()
    {
        //Apply RMSQ preprocessing to A
        //Apply RMQmin and RMQmax preprocessing to C[]
        for(int i = 0; i <= N; ++i)
            rmqMAXC[i][0] = i;
        for (int j = 1; j <= __lg( N + 1 ); ++j )
            for (int i = 0; i + ( 1 << j ) - 1 <= N + 1; ++i ) {
                if(C[rmqMAXC[i][j-1]] >= C[rmqMAXC[i + (1 << (j-1))][j-1]])
                    rmqMAXC[i][j] = rmqMAXC[i][j-1];
                else
                    rmqMAXC[i][j] = rmqMAXC[i + (1 << (j-1))][j-1];
            }
    }

```

```

    }

    pair<int, int> queryRMSQ(int i, int j, int k, int l)
    {
        if(j <= k)
            return pair<int, int>(queryRMQ(i-1, j-1, 0) + 1, queryRMQ(k, l, 2));

        int x[4], y[4];

        x[1] = queryRMQ(i-1, k-1, 0) + 1;
        y[1] = queryRMQ(k, l, 2);

        x[2] = queryRMQ(k, j-1, 0) + 1;
        y[2] = queryRMQ(j, l, 2);

        pair<int, int> tmp = queryRMSQ(k, j);
        x[3] = tmp.first;
        y[3] = tmp.second;

        int maxSum = max(C[x[1]] - C[y[1]-1], max(C[x[2]] - C[y[2]-1], C[x[3]] - C[y[3]-1]));

        if(C[x[1]] - C[y[1]-1] == maxSum)
            return pair<int, int>(x[1], y[1]);
        if(C[x[2]] - C[y[2]-1] == maxSum)
            return pair<int, int>(x[2], y[2]);
        return pair<int, int>(x[3], y[3]);
    }

    int main() {

        cin >> N;
        for(int i = 1; i <= N; ++i)
            cin >> A[i];

        buildCLPM();
        buildRMQ();
        buildRMSQ2();

        int q; cin >> q;
        for(int i = 0; i < q; ++i)
        {
            /*int l, r; cin >> l >> r;

```

```

pair<int, int> ans = queryRMSQ(l, r);
cout << ans.first << " " << ans.second << endl;*/
int a, b, c, d; cin >> a >> b >> c >> d;
pair<int, int> ans = queryRMSQ(a, b, c, d);

```

3.14. Segment Tree Lazy Propagation.

```

/*
    In this example:
    update item[l...r] + val
    query sum(item[l...r])
*/

#define MaxN 1000
#define Left(x) ((x<<1) + 1)
#define Right(x) ((x<<1) + 2)

int st[4*MaxN], lazy[4*MaxN];

void push(int node, int nodeL, int nodeR)
{
    int m = (nodeL + nodeR) / 2;

    lazy[Left(node)] += lazy[node];
    lazy[Right(node)] += lazy[node];

    st[Left(node)] += (m - nodeL + 1) * lazy[node];
    st[Right(node)] += (nodeR - m) * lazy[node];

    lazy[node] = 0;
}

void update(int node, int nodeL, int nodeR, int l, int r, int val)
{

```

3.15. Segment Tree-1D Query.

```

/*
    In this example update is in a position and the query is
    the sum of interval. item[N], st[4*N]
*/
#define Left(x) ((x<<1) + 1)
#define Right(x) ((x<<1) + 2)
#define MaxN 1000

```

```

        cout << ans.first << "_" << ans.second << endl;
    }
    return 0;
}

```

```

    if(l > nodeR || r < nodeL)
        return;
    if(nodeL >= l && nodeR <= r)
    {
        st[node] += (nodeR - nodeL + 1) * val;
        lazy[node] += val;
        return;
    }
    push(node, nodeL, nodeR);

    int m = (nodeL + nodeR) / 2;
    update(Left(node), nodeL, m, l, r, val);
    update(Right(node), m+1, nodeR, l, r, val);
    st[node] = st[Left(node)] + st[Right(node)];
}

int query(int node, int nodeL, int nodeR, int l, int r)
{
    if(l > nodeR || r < nodeL)
        return 0;
    if(nodeL >= l && nodeR <= r)
        return st[node];
    push(node, nodeL, nodeR);
    int m = (nodeL + nodeR) / 2;
    return query(Left(node), nodeL, m, l, r) +
           query(Right(node), m+1, nodeR, l, r);
}

```

```

int item[MaxN];

void build(int *st, int node, int nodeL, int nodeR)
{
    if(nodeL == nodeR)
    {
        st[node] = item[nodeL];
        return;
    }

```

```

    }
    int m = (nodeL + nodeR) / 2;
    build(st, Left(node), nodeL, m);
    build(st, Right(node), m+1, nodeR);
    st[node] = st[Left(node)] + st[Right(node)];
}

void update(int *st, int node, int nodeL, int nodeR, int pos, int val)
{
    if(nodeL == nodeR)
    {
        st[node] = val;
        return;
    }
    int m = (nodeL + nodeR) / 2;
    if(pos <= m)
        update(st, Left(node), nodeL, m, pos, val);

```

3.16. Segment Tree-2D.

```

/*
    TASK : -Range Minimum-Sum Segment Query Problem
           -With two intervals too.
    Compute arrays C, P and M --> O(N)
    Preprocess RMQ --> O(N log N)
    Answer RMSQ queries --> O(1)
*/
#define MAXN 50005
#define LGN 16

int A[MAXN];
int C[MAXN], P[MAXN], M[MAXN], L[MAXN];
int RMQ[MAXN][LGN][2];

//for two intervals
int rmqMAXC[MAXN][LGN];

int N;

// Compute arrays C, P, L and M
//C[i] = sum(A[1]...A[i])
//L[i] = max{k | C[k] >= C[i] k[1, i-1]}
// {0 otherwise }
//P[i] = max{k | k[L[i]+1, i] and C[k-1] <= C[1] forall l[L[i], i-1]}
//M[i] = sum(P[i], i)
void buildCLPM()

```

```

    else
        update(st, Right(node), m+1, nodeR, pos, val);
    st[node] = st[Left(node)] + st[Right(node)];
}

int query(int *st, int node, int nodeL, int nodeR, int l, int r)
{
    if(nodeL == l && nodeR == r)
        return st[node];
    int m = (nodeL + nodeR) / 2;
    if(r <= m)
        return query(st, Left(node), nodeL, m, l, r);
    if(l > m)
        return query(st, Right(node), m+1, nodeR, l, r);
    return query(st, Left(node), nodeL, m, l, m) +
           query(st, Right(node), m+1, nodeR, m+1, r);
}

```

```

{
    for (int i = 1; i <= N; ++i) {
        C[i] = C[i-1] + A[i];
        L[i] = i-1; P[i] = i;
        while ( C[ L[i] ] < C[i] && L[i] ) {
            if ( C[ P[ L[i] ] - 1 ] < C[ P[i] - 1 ] )
                P[i] = P[ L[i] ];
            L[i] = L[ L[i] ];
        }
        M[i] = C[i] - C[ P[i] - 1 ];
    }
}

// Preprocess array C for RMQmin and array M for RMQmax
// RMQ[i][j][0] holds the minimum, while RMQ[i][j][1] holds
// the maximum
void buildRMQ()
{
    for (int i = 0; i <= N; ++i)
        RMQ[i][0][0] = RMQ[i][0][1] = i;

    for (int j = 1; j <= __lg( N + 1 ); ++j)
        for (int i = 0; i + ( 1 << j ) - 1 <= N + 1; ++i ) {
            if ( C[ RMQ[i][ j - 1 ][0] ] <= C[ RMQ[ i +
                ( 1 << ( j - 1 ) ) ][ j - 1 ][0] ] )
                RMQ[i][j][0] = RMQ[i][ j - 1 ][0];

```

```

        else RMQ[i][j][0] = RMQ[ i + ( 1 << ( j - 1 ) ) ]
                          [ j - 1 ][0];
        if ( M[ RMQ[i][ j - 1 ][1] ] >= M[ RMQ[ i +
        ( 1 << ( j - 1 ) ) ][ j - 1 ][1] ] )
            RMQ[i][j][1] = RMQ[i][ j - 1 ][1];
        else RMQ[i][j][1] = RMQ[ i +
        ( 1 << ( j - 1 ) ) ][ j - 1 ][1];
    }
}

int queryRMQ( int l, int r, int b ) {

    int k = __lg( r - l + 1 );

    //For two Intervals
    if(b == 2) return max(rmqMAXC[l][k], rmqMAXC[r - (1<<k) + 1][k]);

    if ( !b ) return C[ RMQ[l][k][b] ] <= C[ RMQ[ r - ( 1 << k ) + 1 ][k][b] ] ?
        RMQ[l][k][b] : RMQ[ r - ( 1 << k ) + 1 ][k][b];
    else return M[ RMQ[l][k][b] ] >= M[ RMQ[ r - ( 1 << k ) + 1 ][k][b] ] ?
        RMQ[l][k][b] : RMQ[ r - ( 1 << k ) + 1 ][k][b];
}

pair<int, int> queryRMSQ( int l, int r )
{
    int x = queryRMQ( l, r, 1 );
    if ( P[x] < 1 ) {
        int y = queryRMQ( x + 1, r, 1 );
        int z = queryRMQ( l - 1, x - 1, 0 ) + 1;
        if ( C[x] - C[ z - 1 ] < M[y] )
            return pair<int, int>(P[y], y);
        return pair<int, int>(z, x);
    }
    return pair<int, int>(P[x], x);
}

//RMSQ with two intervals
//Return i <= x <= j, k <= y <= l
// max{ Sum(x, y) }

void buildRMSQ2()
{
    //Apply RMSQ preprocessing to A
    //Apply RMQmin and RMQmax preprocessing to C[]
    for(int i = 0; i <= N; ++i)
        rmqMAXC[i][0] = i;

```

```

        for (int j = 1; j <= __lg( N + 1 ); ++j )
            for (int i = 0; i + ( 1 << j ) - 1 <= N + 1; ++i ) {
                if(C[rmqMAXC[i][j-1]] >= C[rmqMAXC[i + (1 << (j-1))][j-1]])
                    rmqMAXC[i][j] = rmqMAXC[i][j-1];
                else
                    rmqMAXC[i][j] = rmqMAXC[i + (1 << (j-1))][j-1];
            }
    }

pair<int, int> queryRMSQ(int i, int j, int k, int l)
{
    if(j <= k)
        return pair<int, int>(queryRMQ(i-1, j-1, 0) + 1, queryRMQ(k, l, 2));

    int x[4], y[4];

    x[1] = queryRMQ(i-1, k-1, 0) + 1;
    y[1] = queryRMQ(k, l, 2);

    x[2] = queryRMQ(k, j-1, 0) + 1;
    y[2] = queryRMQ(j, l, 2);

    pair<int, int> tmp = queryRMSQ(k, j);
    x[3] = tmp.first;
    y[3] = tmp.second;

    int maxSum = max(C[x[1]] - C[y[1]-1], max(C[x[2]] -
        C[y[2]-1], C[x[3]] - C[y[3]-1]));

    if(C[x[1]] - C[y[1]-1] == maxSum)
        return pair<int, int>(x[1], y[1]);
    if(C[x[2]] - C[y[2]-1] == maxSum)
        return pair<int, int>(x[2], y[2]);
    return pair<int, int>(x[3], y[3]);
}

int main() {

    cin >> N;
    for(int i = 1; i <= N; ++i)
        cin >> A[i];

    buildCLPM();
    buildRMQ();

```

```

buildRMSQ2();

int q; cin >> q;
for(int i = 0; i < q; ++i)
{
    /*int l, r; cin >> l >> r;
    pair<int, int> ans = queryRMSQ(l, r);

```

3.17. Treap.

```

/*
TASK : Coding a treap

Remarks: Assuming keys are integers. Using Max Heap
Performance:
    Insert:  $O(\log n)$ *
    Erase:  $O(\log n)$ *
    Find:  $O(\log n)$ *
    Find k-th:  $O(\log n)$ *
    * expected

*/

struct generator {

    static const int A = 48271;
    static const int M = 2147483647;
    static const int Q = M / A;
    static const int R = M % A;

    int state;

    generator() {
        srand( time( 0 ) );
        state = rand() + 1;
    }
    int pseudo_random() {
        state = A * ( state % Q ) - R * ( state / Q );
        return state > 0 ? state : state += M;
    }
} g;

struct treap {

    #define SIZE( x ) ( (x) ? (x)->size : 0 )
    #define RESIZE( x ) ( SIZE( (x)->ch[0] ) +

```

```

        cout << ans.first << " " << ans.second << endl;*/
        int a, b, c, d; cin >> a >> b >> c >> d;
        pair<int, int> ans = queryRMSQ(a, b, c, d);
        cout << ans.first << " " << ans.second << endl;
    }
    return 0;
}

```

```

        SIZE( (x)->ch[1] ) + (x)->cnt )

struct node {
    int key, p, size, cnt;
    node *ch[2];
    node( int key ) : key( key ), p( g.pseudo_random() ),
        size( 1 ), cnt( 1 )
    {
        ch[0] = ch[1] = 0;
    }
} *root;

int key;

node* rotate( node *x, bool b ) {
    node *y = x->ch[ !b ];
    x->ch[ !b ] = y->ch[b];
    y->ch[b] = x;
    x->size = RESIZE( x );
    y->size = RESIZE( y );
    return y;
}

node* insert( node *t, const int& key ) {

    if ( !t ) return new node( key );

    if ( key == t->key ) t->cnt++, t->size++;
    else {
        bool b = !( key < t->key );
        t->ch[b] = insert( t->ch[b], key );
        t->size = RESIZE( t );
        if ( t->ch[b]->p > t->p ) t = rotate( t, !b );
    }
    return t;
}

```

```

node* erase( node *t, const int& key ) {

    if ( !t ) return 0;

    if ( key != t->key ) {
        bool b = !( key < t->key );
        t->ch[b] = erase( t->ch[b], key );
        t->size = RESIZE( t );
    }
    else {
        if ( t->cnt > 1 ) t->cnt--, t->size--;
        else {
            if ( !t->ch[0] && !t->ch[1] ) {
                delete t;
                return 0;
            }
            else if ( !t->ch[0] ) t = rotate( t, 0 );
            else if ( !t->ch[1] ) t = rotate( t, 1 );
            else t = rotate( t, t->ch[0]->p > t->ch[1]->p );
            t = erase( t, key );
        }
    }

    return t;
}

```

```

/* "Public" methods */

void insert( int x ) { root = insert( root, key = x ); }
void erase( int x ) { root = erase( root, key = x ); }
int size() { return SIZE( root ); }
bool find( int x ) {
    node *t = root;
    while ( t ) {
        if ( x == t->key ) return 1;
        t = t->ch[ !( x < t->key ) ];
    }
    return 0;
}

int find_kth( int k ) { /* assuming k <= SIZE( root ) */
    node *t = root;
    while ( 1 ) {
        int lo_rank = SIZE( t->ch[0] ) + 1,
            hi_rank = SIZE( t->ch[0] ) + t->cnt;
        if ( lo_rank <= k && k <= hi_rank ) return t->key;
        else if ( k < lo_rank ) t = t->ch[0];
        else { k -= hi_rank; t = t->ch[1]; }
    }
}

treap() : root(0) { }
};

```

3.18. Treap Implicit Key.

```

typedef long long ptype;

ptype seed = 47;

ptype my_rand(){
    seed = (seed * 279470273) % 4294967291LL;
    return seed;
}

struct ImplicitTreap{

    struct item{
        int value;
        ptype prior;
        item *l,*r;
    }
}

```

```

    int sons;

    bool rev;
    long long sum;

    item(){
        item(int value) : value(value), l(0), r(0),
            sons(0), rev(0), sum(value)
        {
            prior = my_rand();
        }
    } *root;

    void fix(item* t){
        if(!t) return;
        t->sons = (t->l ? t->l->sons + 1 : 0) + (t->r ? t->r->sons + 1 : 0);
    }
}

```



```

    t->sum = (t->l ? t->l->sum : 0) + (t->r ? t->r->sum : 0) + t->value;
}

void push(item* it){
    if(it && it->rev){
        it->rev = 0;
        swap(it->l, it->r);
        if(it->l) it->l->rev ^= 1;
        if(it->r) it->r->rev ^= 1;
    }
}

void merge(item* &t, item* l, item* r){
    push(l); push(r);

    if(!l || !r)
        t = l ? l : r;
    else if(l->prior > r->prior)
        merge(l->r, l->r, r), t = l;
    else
        merge(r->l, l, r->l), t = r;

    fix(t);
}

void split(item* t, item* &l, item* &r, int pos, int add = 0){
    if(!t) l = r = NULL;
    else{
        push(t);
        int cur_pos = add + (t->l ? l + t->l->sons : 0);

        if(pos <= cur_pos)
            split(t->l, l, t->l, pos, add), r = t;
        else
            split(t->r, t->r, r, pos, cur_pos + 1), l = t;

        fix(t);
    }
}

void insert(item* &t, item* &it, int pos, int add = 0) {
    if(!t) t = it;
    else{
        push(t);
        int cur_pos = add + (t->l ? l + t->l->sons : 0);

        if(it->prior > t->prior){

```

```

            split(t, it->l, it->r, pos, add), t = it;
        }
        else{
            if(pos <= cur_pos) insert(t->l, it, pos, add);
            else insert(t->r, it, pos, cur_pos + 1);
        }
        fix(t);
    }
}

void remove(item* &t, int pos, int add = 0){
    int cur_pos = add + (t->l ? l + t->l->sons : 0);

    if(cur_pos == pos) merge(t, t->l, t->r);
    else if(pos < cur_pos) remove(t->l, pos, add);
    else remove(t->r, pos, cur_pos + 1);

    fix(t);
}

void reverse(item* t, int l, int r){
    item *t1,*t2,*t3;

    split(t, t1, t2, l);
    split(t2, t2, t3, r-l+1);

    t2->rev ^= 1;
    merge(t, t1, t2);
    merge(t, t, t3);
}

long long sum(item* &t, int lo, int hi, int a, int b, int add = 0){
    if(!t || lo > b || hi < a) return 0;

    if(a <= lo && hi <= b) return t->sum;

    if(t->rev) push(t);
    int cur_key = add + (t->l ? l + t->l->sons : 0);

    long long ret = (a <= cur_key && cur_key <= b ? t->value : 0);

    ret += sum(t->l, lo, cur_key - 1, a, b, add);
    ret += sum(t->r, cur_key + 1, hi, a, b, cur_key + 1);

    return ret;
}

```

```

void print (item* t) {
    if (!t) return;
    push (t);
    print (t->l);
    printf ("%d_", t->value);
    print (t->r);
}

void clear(item* t)
{
    if (!t) return;
    clear (t->l);
    delete t;
    clear (t->r);
}

ImplicitTreap(){
    root = 0;
}

/*Public Methods*/

void print(){
    print(root);
}

```

```

int size(){
    return root->sons + 1;
}

long long sum(int l, int r){
    return sum(root, 0, this->size() - 1, l, r);
}

void reverse(int l, int r){
    reverse(root, l, r);
}

void remove(int pos){
    remove(root, pos);
}

void insert(int pos, int val){
    item* node = new item(val);
    insert(root, node, pos);
}

void clear(){
    clear(root);
}

};

```

3.19. Trie.

```

/*
    TASK : Given a set P of strings and a string S, count how many
           elements of P contain S as a prefix, and how many p(i), for
           some i, have |p(i)| < |S|.

    Remarks: Using English alphabet (|S| = 26)
    Performance:
        Insert: O(|p|)
        Count: O(|p|)
        p: string processed
*/

#define MAXLEN 20000

struct Trie {
    struct node {

```

```

        int partial, full;
        node *edge[26];
        node() : partial(0), full(0) { memset( edge, 0, sizeof( edge ) ); }
    } *root;

    Trie() { root = new node(); }

    void insert( char s[], int len ) {

        node *t = root;
        for ( int i = 0; i < len; i++ ) {
            char c = s[i] - 'a';
            if ( !t->edge[c] ) t->edge[c] = new node();
            t = t->edge[c];
            t->partial++;
        }
        t->full++;
    }

```

```
    }  
  
    int count( char s[], int len ) {  
  
        if ( !root ) return 0;  
  
        node *t = root;  
        int ret = 0;  
  
        for ( int i = 0; i < len; i++ ) {
```

```
            if ( !t ) break;  
            ret += t->full;  
            t = t->edge[ s[i] - 'a' ];  
        }  
        if ( t ) ret += t->partial;  
  
        return ret;  
    }  
};
```

4. DYNAMIC PROGRAMMING

4.1. Convex Hull Trick.

```

typedef pair<int, int> pii;
typedef long long ll;

struct line{
    ll m, b;
    line(ll m, ll b): m(m), b(b){}
};

struct ConvexHullTrick{

    int len, ptr;
    vector<line> r;
    ConvexHullTrick(int n)
    {
        r.assign(n, line(0, 0));
        ptr = len = 0;
    }

    bool bad(line l1, line l2, line l3)
    {

```

```

        return (l3.b - l1.b) * (l1.m - l2.m) <
               (l2.b - l1.b) * (l1.m - l3.m);
    }

    void add(line x)
    {
        while(len >= 2 && bad(r[len-2], r[len-1], x))
            --len;
        r[len++] = x;
    }

    ll query(int x)
    {
        ptr = min(ptr, len-1);
        while(ptr + 1 < len && r[ptr+1].m * x +
            r[ptr+1].b < r[ptr].m * x + r[ptr].b)
            ++ptr;
        return r[ptr].m * x + r[ptr].b;
    }
};

```

4.2. Longest Increasing Subsequence.

```

const int oo = 99999999;

#define index_of(as, x)\
    distance(as.begin(), lower_bound(as.begin(), as.end(), x))

/*
    Tested: LISTA
    Contest 3 COCI 2006-2007
*/
vector<int> lis_fast(const vector<int> &a)
{
    const int n = a.size();
    vector<int> A(n, oo), id(n);

    for(int i = 0; i < n; ++i)

```

```

    {
        id[i] = index_of(A, a[i]);
        A[id[i]] = a[i];
    }

    int m = *max_element(id.begin(), id.end());
    vector<int> b(m+1);

    for(int i = n-1; i >= 0; --i)
        if(id[i] == m)
            b[m--] = a[i];

    return b;
}

```

4.3. Matrix Chain.

```
const int oo = 1 << 30;

int matrix_chain(const vector<int> &p)
{
    int n = p.size()-1;
    int dp[n+1][n+1];

    for(int i = 1; i <= n; ++i)
        dp[i][i] = 0;
```

```
    for(int len = 2; len <= n; ++len) {
        for(int i = 1, j = i+len-1; j <= n; ++i, ++j) {
            dp[i][j] = oo;
            for(int k = i; k < j; ++k)
                dp[i][j] = min(dp[i][j], dp[i][k] + dp[k+1][j] + p[i-1] *
                                p[k] + p[j]);
        }
    }

    return dp[1][n];
}
```

5. GEOMETRY

5.1. Basic Operation.

```

#define x(c) real(c)
#define y(c) imag(c)
#define NEXT(i) ((i) + 1) % n

const double EPS = 1e-7;
const int oo = (1<<30);
typedef complex<double> point;

int cmp_double(double x, double y=0)
{
    return (x <= y + EPS) ? (x + EPS < y) ? -1 : 0 : 1;
}

bool cmp_point(const point &a, const point &b)
{
    return (a.x() != b.x()) ? ( cmp_double(a.x() , b.x()) == -1 ) :
        ( cmp_double(a.y() , b.y()) == -1 );
}

bool operator <(const point &a, const point &b)
{
    return cmp_point(a, b);
}

//a1*b2 - a2*b1 = axb = |a||b|*sin()
double cross(const point &a, const point &b)
{
    return imag(conj(a)*b);
}

//a1*b1 + a2*b2 = a.b = |a||b|*cos(a,b)
double dot(const point &a, const point &b)
{
    return real(conj(a)*b);
}

int ccw (point a, point b, point c)
{
    b -= a; c -= a;
    if (cross (b, c) > 0) return + 1; // counter clockwise
    if (cross (b, c) < 0) return - 1; // clockwise
    if (dot (b, c) < 0) return + 2; // c - a - b on line

```

```

        if(cmp_double(norm(b), norm(c)) == -1) return - 2; // a - b - c on line
        return 0; // a - c - b on line;
    }

    int cw (point a, point b, point c)
    {
        return -ccw(a, b, c);
    }

    double sq(double x)
    {
        return x*x;
    }

    double dist2(const point &a, const point &b)
    {
        return sq(a.x() - b.x()) + sq(a.y() - b.y());
    }

    double dist(const point &a, const point &b)
    {
        return abs(a-b);
    }

    /*
    Compares to 2D points by angle
    Angle -90 is the first
    Tested: LightOJ 1292
    */
    bool polar_cmp(point a, point b)
    {
        if(a.x() >= 0 && b.x() < 0) return true;
        if(a.x() < 0 && b.x() >= 0) return false;
        if(a.x() == 0 && b.x() == 0)
        {
            if(a.y() > 0 && b.y() < 0) return false;
            if(a.y() < 0 && b.y() > 0) return true;
        }
        return cross(a, b) > 0;
    }

    //p-q-r: clockwise

```

```
double angle(point p, point q, point r)
{
    point u = p-q, v = r-q;
    return atan2(cross(u,v), dot(u,v));
}

point rotateCCW90(point p)
{

```

5.2. Circles.

```
struct circle{
    point center;
    double ratio;

    circle(point center, double ratio) : center(center) , ratio(ratio){}
};

//Tested [BAPC 2010 Clocks]
vector<point> circles_intersection(const circle &c1, const circle &c2)
{
    vector<point> ret;
    double d = dist(c1.center, c2.center);
    if (d > c1.ratio + c2.ratio || d + min(c1.ratio, c2.ratio) < max(c1.ratio, c2.ratio))
        return ret;
    double x = ( d*d - c2.ratio*c2.ratio + c1.ratio*c1.ratio ) / ( 2*d );
    double y = sqrt( c1.ratio * c1.ratio - x*x );
    point v = (c2.center - c1.center) / d;
    ret.push_back(c1.center + v*x + rotateCCW90(v) * y);
    if (y > 0)
        ret.push_back(c1.center + v*x - rotateCCW90(v) * y);
    return ret;
}

//Interseccion Linea-Circulo
vector<point> intersectLC(line l, circle c)
{
    point a = l[0], b = l[1];
    vector<point> ret;
    b = b-a;
    a = a - c.center;
    double A = dot(b, b);
    double B = dot(a, b);
    double C = dot(a, a) - c.ratio*c.ratio;
    double D = B*B - A*C;

```

```
        return point(-p.imag(),p.real());
    }

    point rotate_by(const point &p, const point &about, double radians)
    {
        return (p - about) * exp(point(0, radians)) + about;
    }

    if (cmp(D) < 0) return ret;
    ret.push_back(c.center + a + b*(-B + sqrt(D+EPS)) / A);
    if (cmp(D) > 0) ret.push_back(c.center + a + b * (-B - sqrt(D)) / A);

    return ret;
}

/*
    Area of the intersection of a circle with a polygon
    Circle's center lies in (0,0)
    Polygon must be given counterclockwise
    Tested [Light OJ 1358]
*/

#define xx(_t) (xa+(_t)*a)
#define yy(_t) (ya+(_t)*b)

double radian(double xa,double ya,double xb,double yb)
{
    return atan2(xa*yb - xb*ya, xa*xb + ya*yb);
}

double part(double xa,double ya,double xb,double yb,double r)
{
    double l = sqrt((xa-xb) * (xa-xb) + (ya-yb) * (ya-yb));
    double a = (xb-xa) / l, b = (yb - ya) / l, c = a*xa + b*ya;
    double d = 4.0 * (c*c - xa*xa - ya*ya + r*r);
    if(d < EPS) return radian(xa,ya,xb,yb) * r * r * 0.5;
    else
    {
        d = sqrt(d) * 0.5;
        double s = -c-d, t = -c+d;
        if(s < 0.0) s = 0.0;

```

```

        else if(s > 1) s = 1;
        if(t < 0.0) t = 0.0;
        else if(t > 1) t = 1;
        return (xx(s)*yy(t) - xx(t)*yy(s) + (radian(xa,ya,xx(s),yy(s))
            + radian(xx(t),yy(t),xb,yb))*r*r) * 0.5;
    }
}

double area_intersectionPC(polygon P, double r)
{
    double s = 0.0;
    int n = (int)P.size();
    P.push_back(P[0]);
    for(int i = 0; i < n; ++i)
        s += part(P[i].x(), P[i].y(), P[NEXT(i)].x(), P[NEXT(i)].y(), r);
    return fabs(s);
}

// circle tangents through point
vector<point> tangent(point p, circle C)
{
    // not tested enough

    double D = abs(p - C.p);

    if (D + eps < C.r) return {};
    point t = C.p - p;

    double theta = asin( C.r / D );
    double d = cos(theta) * D;

    t = t / abs(t) * d;
    if ( abs(D - C.r) < eps ) return {p + t};
    point rot( cos(theta), sin(theta) );
    return {p + t * rot, p + t * conj(rot)};
}

bool incircle(point a, point b, point c, point p)
{
    a -= p; b -= p; c -= p;

```

5.3. Closest Pair Points.

```

/*
    Compute distance between closest points.

```

```

        return norm(a) * cross(b, c)
            + norm(b) * cross(c, a)
            + norm(c) * cross(a, b) >= 0;
        // < : inside, = cocircular, > outside
    }

point three_point_circle(point a, point b, point c)
{
    point x = 1.0 / conj(b - a), y = 1.0 / conj(c - a);
    return (y - x) / (conj(x) * y - x * conj(y)) + a;
}

/*
    Get the center of the circles that pass through p0 and p1
    and has ratio r.

    Be careful with epsilon.
*/
vector<point> two_point_ratio_circle(point p0, point p1, double r){
    if (abs(p1 - p0) > 2 * r + eps) // Points are too far.
        return {};

    point pm = (p1 + p0) / 2.01;
    point pv = p1 - p0;

    pv = point(-pv.imag(), pv.real());

    double x1 = p1.real(), y1 = p1.imag();
    double xm = pm.real(), ym = pm.imag();
    double xv = pv.real(), yv = pv.imag();

    double A = (sqr(xv) + sqr(yv));
    double C = sqr(xm - x1) + sqr(ym - y1) - sqr(r);
    double D = sqrt( - 4 * A * C );
    double t = D / 2.0 / A;

    if (abs(t) <= eps)
        return {pm};

    return {c1, c2};
}

```

Tested: AIZU(judge.u-aizu.ac.jp) CGL.5A


```

        Complexity:  $O(n \log n)$ 
    */

double closest_pair_points(vector<point> &P)
{
    auto cmp = [](point a, point b)
    {
        return make_pair(a.imag(), a.real())
            < make_pair(b.imag(), b.real());
    };

    int n = P.size();
    sort(P.begin(), P.end(), cmp_point);

    set<point, decltype(cmp)> S(cmp);
    const double oo = 1e9; // adjust
    double ans = oo;

```

5.4. Convex Cut.

```

/*
    Cut a convex polygon by a line and
    return the part to the left of the line

    Tested: AIZU(judge.u-aizu.ac.jp) CGL.4C
    Complexity:  $O(n)$ 
*/

polygon convex_cut(const polygon &P, const line &l)
{

```

5.5. Convex Hull 3D.

```

// TODO: Change vec3 to use point3d from team reference

template<typename vtype>
struct vec3 {
    vec3() { X[0] = X[1] = X[2] = 0; }
    vec3(vtype x, vtype y, vtype z) { X[0] = x; X[1] = y; X[2] = z; }

    /* 3D cross product */
    vec3 operator*(const vec3& v) const {
        return vec3(X[1] * v.X[2] - X[2] * v.X[1],
            X[2] * v.X[0] - X[0] * v.X[2],

```

```

        for (int i = 0, ptr = 0; i < n; ++i)
        {
            while (ptr < i && abs(P[i].real() - P[ptr].real()) >= ans)
                S.erase(P[ptr++]);

            auto lo = S.lower_bound(point(-oo, P[i].imag() - ans - eps));
            auto hi = S.upper_bound(point(-oo, P[i].imag() + ans + eps));

            for (decltype(lo) it = lo; it != hi; ++it)
                ans = min(ans, abs(P[i] - *it));

            S.insert(P[i]);
        }

        return ans;
    }
}

```

```

polygon Q;
for (int i = 0, n = P.size(); i < n; ++i)
{
    point A = P[i], B = P[(i + 1) % n];
    if (ccw(l.p, l.q, A) != -1) Q.push_back(A);
    if (ccw(l.p, l.q, A) * ccw(l.p, l.q, B) < 0)
        Q.push_back(crosspoint((line){ A, B }, l));
}
return Q;
}

```

```

        X[0] * v.X[1] - X[1] * v.X[0]);
    }

    vec3 operator-(const vec3& v) const {
        return vec3(X[0] - v.X[0], X[1] - v.X[1], X[2] - v.X[2]);
    }

    vec3 operator+(const vec3& v) const {
        return vec3(X[0] + v.X[0], X[1] + v.X[1], X[2] + v.X[2]);
    }

    vec3 operator-() const {

```

```

    return vec3(-X[0], -X[1], -X[2]);
}

vec3 operator*(vtype d) const{
    return vec3(X[0] * d, X[1] * d, X[2] * d);
}

vtype dot(const vec3& v) const {
    return X[0] * v.X[0] + X[1] * v.X[1] + X[2] * v.X[2];
}

bool operator !=(const vec3 v){
    return X[0] != v.X[0] || X[1] != v.X[1] || X[2] != v.X[2];
}

void print(){
    cout << X[0] << " " << X[1] << " " << X[2] << endl;
}

bool zero(){
    return abs(X[0]) < eps && abs(X[1]) < eps && abs(X[2]) < eps;
}

bool notZero(){
    return abs(X[0]) > eps || abs(X[1]) > eps || abs(X[2]) > eps;
}

vtype X[3];
};

typedef vec3<double> point;

struct face{
    int idx[3];

    face(){}
    face(int i, int j, int k){
        idx[0] = i, idx[1] = j, idx[2] = k;
    }

    int& operator[](int u) { return idx[u]; }
};

vector<point> read(){
    int n; cin >> n;
    vector<point> P(n);

```

```

        for (int i = 0; i < n; ++i){
            double x, y, z; cin >> x >> y >> z;
            P[i] = point(x, y, z);
        }

        return P;
    }

vector<face> convex_hull( vector<point> &cloud ){
    // bad

    int n = (int)cloud.size();

    point a = cloud[0], b = cloud[1];

    for (int i = 2; i < n; ++i){
        point nr = (b - a) * (cloud[i] - a);

        if (nr.notZero()){
            swap(cloud[i], cloud[2]);
            break;
        }
    }

    point c = (b - a) * (cloud[2] - a);

    for (int i = 3; i < n; ++i){
        if (abs( c.dot( cloud[i] - a ) ) > eps){
            swap(cloud[i], cloud[3]);
            break;
        }
    }

    vector<face> faces;

    function<point(face&)> normal = [&](face &f){
        point a = cloud[ f[1] ] - cloud[ f[0] ];
        point b = cloud[ f[2] ] - cloud[ f[0] ];
        return a * b;
    };

    function<void(int,int,int)> add_face = [&](int x, int y, int z){
        point a = cloud[x] * n, b = cloud[y] * n, c = cloud[z] * n;

        point nr = (b - a) * (c - a);

        for (int i = 0; i < n; ++i){

```

```

        point d = cloud[i] - a;

        auto value = d.dot( nr );

        if (abs(value) > eps){
            if (value > 0) swap(y, z);
            break;
        }
    }

    faces.push_back( face(x, y, z) );
};

for (int i = 0; i < 4; ++i)
    for (int j = i + 1; j < 4; ++j)
        for (int k = j + 1; k < 4; ++k)
            add_face(i, j, k);

for (int i = 4; i < n; ++i){
    point x = cloud[i];

    vector<vi> seen(n, vi(n));
    vector<face> next_faces;

    for (auto f : faces){
        if ( (x - cloud[ f[0] ]).dot( normal(f) ) > eps ){
            for (int u = 0; u < 3; ++u)
                for (int v = 0; v < 3; ++v)
                    seen[ f[u] ][ f[v] ]++;
        }
        else
            next_faces.push_back( f );
    }

    faces.swap( next_faces );

    for (int j = 0; j < i; ++j)
        for (int k = j + 1; k < i; ++k){
            if (seen[j][k] == 1)
                add_face(i, j, k);
        }
}

return faces;
}

```

```

int L[ 100 ];

vector<face> convex_hull_slow( vector<point> &cloud ){
    // good  $O(n^4)$ 

    int n = (int)cloud.size();

    vector<face> faces;

    for (int i = 0; i < n; ++i)
        for (int j = i + 1; j < n; ++j)
            for (int k = j + 1; k < n; ++k){

                point a = cloud[i], b = cloud[j], c = cloud[k];
                point nr = (b - a) * (c - a);

                int pnt = 0;
                L[ pnt++ ] = j;
                L[ pnt++ ] = k;

                bool proc = true;

                int v = 0, V = 0;

                for (int l = 0; l < n && proc; ++l){
                    if (l == i || l == j || l == k) continue;

                    double t = nr.dot( cloud[l] - a );
                    if ( abs(t) < eps){
                        if (l < k) proc = false;
                        else L[ pnt++ ] = l;
                    }
                    else{
                        if (t < 0) v = -1;
                        else V = +1;
                    }
                }

                if (!proc || v * V == -1) continue;

                // cout << "tri: " << i << " " << j << " " << k << endl;
                // for (int l = 0; l < pnt; ++l)
                //     cout << L[ l ] << " ";
                // cout << endl;

                function<bool(int,int)> compare = [&](int u, int v){
                    return nr.dot( (cloud[u] - a) * (cloud[v] - a) ) > 0

```

```

        };

        sort(L, L + pnt, compare);

        for (int l = 0; l + 1 < pnt; ++l)
            faces.push_back( face(i, L[l], L[l + 1]) );
    }

    return faces;
}

void mass_center( vector<point> &cloud, vector<face> &faces ){
    point pivot = cloud[0];

    double x = 0, y = 0, z = 0, v = 0;

    for (auto f : faces){
        auto value = ( cloud[f[0]] - pivot ).dot(
            ( cloud[f[1]] - pivot ) * (cloud[f[2]] - pivot)

```

```

        );

        point sum = cloud[f[0]] + cloud[f[1]] + cloud[f[2]] + pivot;
        double cvol = abs( 1. * value / 6 );

        v += cvol;

        cvol /= 4;

        x += cvol * sum.X[0];
        y += cvol * sum.X[1];
        z += cvol * sum.X[2];
    }

    x /= v, y /= v, z /= v;

    // Mass center of a polyhedron at (x, y, z)
}

```

5.6. Lines.

```

struct line : public vector<point>{
    line(const point &a, const point &b){
        if(a < b){
            push_back(a);
            push_back(b);
        }
        else{
            push_back(b);
            push_back(a);
        }
    }
};

bool intersectLL (const line &l, const line &m)
{
    return abs (cross (l[1] - l[0], m[1] - m[0])) > EPS || // non-parallel
        abs (cross (l[1] - l[0], m[0] - l[0])) < EPS; // same line
}

bool intersectLP (const line &l, const point &p)
{
    return abs (cross (l[1] - p, l[0] - p)) < EPS;
}

```

```

point projectionPL (const point &p, const line &l)
{
    double t = dot (p - l[0], l[0] - l[1]) / norm (l[0] - l[1]);
    return l[0] + t * (l[0] - l[1]);
}

point reflectPL(const point &p, const line &l)
{
    point z = p - l[0];
    point w = l[1] - l[0];
    return conj(z / w) * w + l[0];
}

double distancePL (const point &p, const line &l)
{
    return abs (p - projectionPL (p, l));
}

double distanceLL (const line &l, const line &m)
{
    return intersectLL (l, m) ? 0 : distancePL (m[0], l);
}

```

```
//Punto interseccion recta recta
point crosspoint(const line &l, const line &m)
{
    double A = cross( l[1] - l[0], m[1] - m[0]);
    double B = cross( l[1] - l[0], l[1] - m[0]);
    if (abs(A) < EPS && abs(B) < EPS) return m[0]; //Same line
    if (abs(A) < EPS) return point(0,0); //parallels
    return m[0] + B / A * (m[1] - m[0]);
}

bool parallelLL(const line &l, const line &m)
```

5.7. Minkowski.

```
/*
    Minkowski sum of two convex polygons.  $O(n + m)$ 

    Note: Polygons MUST be counterclockwise
*/

polygon minkowski(polygon &A, polygon &B){
    int na = (int)A.size(), nb = (int)B.size();

    if (A.empty() || B.empty()) return polygon();

    rotate(A.begin(), min_element(A.begin(), A.end()), A.end());
    rotate(B.begin(), min_element(B.begin(), B.end()), B.end());

    int pa = 0, pb = 0;
```

5.8. Point 3D.

```
const double pi = acos(-1.0);

// Construct a point on a sphere with center on the origin and radius R
// TESTED [COJ-1436]
struct point3d
{
    double x, y, z;

    point3d(double x = 0, double y = 0, double z = 0) : x(x), y(y), z(z) {}

    double operator*(const point3d &p) const
```

```
{
    return !cmp_double(cross(l[1] - l[0], m[0] - m[1]));
}

bool collinearLL(const line &l, const line &m)
{
    return parallelLL(l, m)
        && !cmp_double(cross(l[0] - l[1], l[0] - m[0]))
        && !cmp_double(cross(m[0] - m[1], m[0] - l[0]));
}
```

```
    polygon M;

    while (pa < na && pb < nb){
        M.push_back(A[pa] + B[pb]);
        double x = cross(A[(pa + 1) % na] - A[pa],
                        B[(pb + 1) % nb] - B[pb]);

        if (x <= eps) pb++;
        if (-eps <= x) pa++;
    }

    while (pa < na) M.push_back(A[pa++] + B[0]);
    while (pb < nb) M.push_back(B[pb++] + A[0]);

    return M;
}
```

```
{
    return x * p.x + y * p.y + z * p.z;
}

point3d operator-(const point3d &p) const
{
    return point3d(x - p.x, y - p.y, z - p.z);
}

double abs(point3d p)
{

```

```

    return sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
}

point3d from_polar(double lat, double lon, double R)
{
    lat = lat / 180.0 * pi;
    lon = lon / 180.0 * pi;
    return point3d(R * cos(lat) * sin(lon),
                   R * cos(lat) * cos(lon), R * sin(lat));
}

struct plane
{
    double A, B, C, D;
};

double euclideanDistance(point3d p, point3d q)
{
    return abs(p - q);
}

/*
Geodesic distance between points in a sphere
R is the radius of the sphere

```

5.9. Polygon Triangulation.

```

typedef vector <point> triangle;

triangle make_triangle (const point &a, const point &b, const point &c)
{
    triangle ret (3);
    ret[0] = a; ret[1] = b; ret[2] = c;
    return ret;
}

bool triangle_contains (const triangle &tri, const point &p)
{
    return ccw (tri[0], tri[1], p) >= 0 &&
           ccw (tri[1], tri[2], p) >= 0 &&
           ccw (tri[2], tri[0], p) >= 0;
}

bool ear_Q (int i, int j, int k, const polygon &P) {
    triangle tri = make_triangle (P[i], P[j], P[k]);
    if (ccw (tri[0], tri[1], tri[2]) <= 0) return false;

```

```

*/
double geodesic_distance(point3d p, point3d q, double r)
{
    return r * acos(p * q / r / r);
}

const double eps = 1e-9;

// Find the rect of intersection of two planes on the space
// The rect is given parametrical
// TESTED [TIMUS 1239]
void planePlaneIntersection(plane p, plane q)
{
    if (abs(p.C * q.B - q.C * p.B) < eps)
        return; // Planes are parallel

    double mz = (q.A * p.B - p.A * q.B) / (p.C * q.B - q.C * p.B);
    double nz = (q.D * p.B - p.D * q.B) / (p.C * q.B - q.C * p.B);

    double my = (q.A * p.C - p.A * q.C) / (p.B * q.C - p.C * q.B);
    double ny = (q.D * p.C - p.D * q.C) / (p.B * q.C - p.C * q.B);

    // parametric rect: (x, my * x + ny, mz * x + nz)
}

```

```

    for (int m = 0; m < (int)P.size(); ++m)
        if (m != i && m != j && m != k)
            if (triangle_contains (tri, P[m]))
                return false;

    return true;
}

void triangulate (const polygon &P, vector <triangle> &t)
{
    const int n = P.size();
    vector <int> l, r;

    for (int i = 0; i < n; ++i) {
        l.push_back ((i-1 + n)% n);
        r.push_back ((i + 1 + n)% n);
    }

    int i = n-1;
    while ((int)t.size() < n-2)

```

```

{
    i = r[i];
    if (ear_Q (l[i], i, r[i], P))
    {
        t.push_back (make_triangle (P[ l[i] ], P[i], P[ r[i] ]));
        l[ r[i] ] = l[i];
        r[ l[i] ] = r[i];
    }
}

/*
    Perturbative deformation of a polygon.
    Each side of the polygon in counterclockwise
    polygon len making just the right translation.
*/

```

5.10. Rectilinear MST.

```

/*
    Tested: USACO OPEN08 (Cow Neighborhoods)
    Complexity: O(n log n)
*/

typedef long long ll;
typedef complex<ll> point;

ll rectilinear_mst(vector<point> ps)
{
    vector<int> id(ps.size());
    iota(id.begin(), id.end(), 0);

    struct edge
    {
        int src, dst;
        ll weight;
    };

    vector<edge> edges;
    for (int s = 0; s < 2; ++s)
    {
        for (int t = 0; t < 2; ++t)
        {
            sort(id.begin(), id.end(), [&](int i, int j)
            {

```

```

#define curr(P, i) P[i]
#define prev(P, i) P [ ( i - 1 ) + P.size() ] % P.size()
#define next(P, i) P[ ( i + 1 ) % P.size() ]

polygon shrink_polygon (const polygon &P, double len)
{
    polygon res;
    for (int i = 0; i < (int)P.size(); ++i) {
        point a = prev (P, i), b = curr (P, i), c = next(P, i);
        point u = (b - a) / abs (b - a);
        double th = arg((c - b) / u) * 0.5;
        point tmp(-sin (th), cos (th));
        res.push_back (b + u * tmp * len / cos(th));
    }
    return res;
}

```

```

        return real(ps[i] - ps[j]) < imag(ps[j] - ps[i]);
    });

    map<ll, int> sweep;

    for (int i : id)
    {
        for (auto it = sweep.lower_bound(-imag(ps[i]));
             it != sweep.end(); sweep.erase(it++))
        {
            int j = it->second;
            if (imag(ps[j] - ps[i]) < real(ps[j] - ps[i]))
                break;
            ll d = abs(real(ps[i] - ps[j]))
                + abs(imag(ps[i] - ps[j]));
            edges.push_back({ i, j, d });
        }
        sweep[-imag(ps[i])] = i;
    }

    for (auto &p : ps)
        p = point(imag(p), real(p));

    for (auto &p : ps)
        p = point(-real(p), imag(p));
}

```

```

    }

    ll cost = 0;
    sort(edges.begin(), edges.end(), [](edge a, edge b)
    {
        return a.weight < b.weight;
    });

```

5.11. Rotating Calipers.

```

/*
    Gets all the antipodal pair of points
    Time: O(n)
*/
#define NEXT(i) (((i) + 1) % n)
double area (point a, point b, point c) //2 * area
{
    return abs(cross(b - a, c - a));
}

vector<pair<int, int> > antipodal_pairs (polygon &P)
{
    vector<pair<int, int> > ans;
    int n = P.size();

    if (P.size() == 2)
        ans.push_back(make_pair(0, 1));

    if (P.size() < 3)
        return ans;

    int q0 = 0;

    while (area(P[n - 1], P[0], P[NEXT(q0)]) >
           area(P[n - 1], P[0], P[q0]))
        ++q0;

    for (int q = q0, p = 0; q != 0 && p <= q0; ++p)
    {
        ans.push_back(make_pair(p, q));

        while (area(P[p], P[NEXT(p)], P[NEXT(q)]) >
               area(P[p], P[NEXT(p)], P[q]))
        {
            q = NEXT(q);

```

```

        union_find uf(ps.size());
        for (edge e : edges)
            if (uf.join(e.src, e.dst))
                cost += e.weight;

        return cost;
    }

```

```

        if (p == q0 && q == 0)
            return ans;

        ans.push_back(make_pair(p, q));
    }

    if (area(P[p], P[NEXT(p)], P[NEXT(q)]) ==
        area(P[p], P[NEXT(p)], P[q]))
    {
        if (p != q0 || q != n - 1)
            ans.push_back(make_pair(p, NEXT(q)));
        else
            ans.push_back(make_pair(NEXT(p), q));
    }

    return ans;
}

/*
    Gets the farthest pair of points of the given points.
    (maybe TLE using double)
    TESTED [POJ 2187]
*/
pair<point, point> farthest_pair (polygon &P)
{
    P = convex_hull(P);
    vector<pair<int, int> > pairs = antipodal_pairs(P);

    double best = 0;
    pair<point, point> ans;

    for (int i = 0; i < (int)pairs.size(); ++i)

```



```

{
    point p1 = P[pairs[i].first];
    point p2 = P[pairs[i].second];

    double dist = norm(p1-p2);

    if (dist > best)
    {
        best = dist;
        ans = make_pair(p1, p2);
    }
}

return ans;
}

/*
    Gets the minimum distance between parallel lines of
    support of the convex polygon P
    Time: O(n)
*/

double check (int a, int b, int c, int d, polygon &P)
{
    for (int i = 0; i < 4 && a != c; ++i)
    {
        if (i == 1)
            swap(a, b);
        else

```

```

            swap(c, d);
        }
        if (a == c) //a admits a support line parallel to bd
        {
            //assert(b != d)
            double A = area(P[a], P[b], P[d]); //double of the triangle area
            double base = abs(P[b] - P[d]); //base of the triangle abd
            return A / base;
        }
        return oo;
    }
}

double polygon_width (polygon &P)
{
    if (P.size() < 3)
        return 0;

    vector<pair<int, int>> pairs = antipodal_pairs(P);
    double best = oo;
    int n = pairs.size();

    for (int i = 0; i < n; ++i)
    {
        double tmp = check(pairs[i].first, pairs[i].second,
                           pairs[NEXT(i)].first,
                           pairs[NEXT(i)].second, P);
        best = min(best, tmp);
    }
    return best;
}

```

5.12. Segment Intersect.

```

#define _GLIBCXX_DEBUG
#include <stdio.h>
#include <iostream>
#include <string>
#include <string.h>
#include <vector>
#include <set>
#include <map>
#include <queue>
#include <stack>
#include <complex>
#include <algorithm>

```

```

using namespace std;
#define REP(i,n) for(int i=0;i<(int)n;++i)
#define FOR(i,c) for(__typeof((c).begin()) i=(c).begin();i!=(c).end();++i)
#define ALL(c) (c).begin(), (c).end()
#define Y(c) imag(c)
#define X(c) real(c)
#define INF 1000000000
//Graph Only
typedef int Weight;
struct Edge {
    int src, dst;
    Weight weight;

```

```

Edge( int src, int dst, Weight weight) :
    src(src), dst(dst), weight(weight) {}
};

bool operator <( const Edge &e, const Edge &f) {
    return e.weight!=f.weight?e.weight>f.weight:
        e.src!=f.src?e.src<f.src:e.dst<f.dst;
}

typedef vector<Edge> Edges;
typedef vector<Edges> Graph;
typedef vector<Weight> Array;
typedef vector<Array> Matrix;
#define P complex<double>
typedef vector<P> Pol;
bool operator<(const P &a,const P &b){
    return X(a)!=X(b)?X(a)<X(b):Y(a)<Y(b);}
struct L: public vector<P>{
    L (const P &a,const P &b){
        if(a<b){push_back(a);push_back(b);}
        else{push_back(b);push_back(a);}}};
const double EPS = 1e-8, oo = 1e12;

bool op_min(const P &a,const P &b){
    return X(a)!=X(b)?X(a)<X(b):Y(a)<Y(b);}

double cross(P a, P b){ return Y(conj(a)*b);}
double dot(P a, P b){ return X(conj(a)*b);}

int ccw(P a, P b, P c){ //Orientacion de 3 puntos
    b-=a; c-=a;
    if (cross(b,c) > 0) return +1; //counter clockwise
    if (cross(b,c) < 0) return -1; //clockwise
    if (dot(b,c) < 0) return +2; //c - a - b line
    if (norm(b)<norm(c)) return -2; //a - b - c line
    return 0;}

bool intersectSS (L s, L t){ //Inters de 2 segm
    if (abs(s[0]-t[0])<EPS||abs(s[0]-t[1])< EPS||
        abs(s[1]-t[0])<EPS||abs(s[1]-t[1])< EPS)
        return 1; //Puntos Iguales
    return ccw(s[0],s[1],t[0])*ccw(s[0],s[1],t[1])<=0
        && ccw(t[0],t[1],s[0]) * ccw(t[0],t[1],s[1]) <= 0;}

P crosspoint(L l, L m){ //Punto inters /2 rectas
    double A = cross( l[1]-l[0], m[1]-m[0]);
    double B = cross( l[1]-l[0], l[1]-m[0]);
    if (abs(A) < EPS && abs(B) < EPS)

```

```

        return m[0]; //Same L
    if (abs(A) < EPS) return P(0,0); //parallels
    return m[0] + B / A * (m[1] - m[0]);}

struct event {
    double x;int type;L seg;
    event(double x,int type,const L& seg):
        x(x),type(type),seg(seg){}
    bool operator<(const event &e)const{
        return x!=e.x?x>e.x:type>e.type;}};
struct segComp{
    bool operator()(const L &a,const L &b){
        if(a[0]<b[0])return true;
        if(a[1]<b[1])return true;
        return false;}};
int segment_intersects( const vector<L>& segs, vector<P> &out) {
    priority_queue<event> Q;
    for(int i=0;i<segs.size();++i){
        double x1=real(segs[i][0]),x2=real(segs[i][1]);
        Q.push(event(min(x1,x2),0,segs[i]));
        Q.push(event(max(x1,x2),1,segs[i]));}
    int count=0;
    set<L,segComp> T;
    while(!Q.empty()){
        event e=Q.top();Q.pop();
        if(e.type==0){
            for (set<L,segComp>::iterator itr=T.begin();
                itr != T.end();++itr)
                if (intersectSS(*itr, e.seg)) {
                    out.push_back(crosspoint(*itr, e.seg));
                    ++count;}
            T.insert(e.seg);
        } else T.erase(e.seg);}
    return count;}

bool merge_if_able (L &s, L t) {
    if (abs(cross(s[1]-s[0],t[1]-t[0]))>EPS)return false;
    if(ccw(s[0],t[0],s[1])==+1||
        ccw(s[0],t[0],s[1])==+1)return false; //not on the same line
    if(ccw(s[0],s[1],t[0])==+2||
        ccw(t[0],t[1],s[0])==+2)return false; //separated
    s=L(min(s[0],t[0],op_min),max(s[1],t[1],op_min));
    return true;}

void merge_segments(vector<L>& segs) {
    for(int i=0;i<segs.size();++i)
        for(int j=i+1;j<segs.size();++j)
            if(merge_if_able(segs[i],segs[j]))

```

```

        segs[j--]=segs.back(),segs.pop_back();}

pair<P,P> closestPair(vector<P>&p){
    int n=p.size(),s=0,t=1,m=2,S[n];S[0]=0,S[1]=1;
    sort(ALL(p),op_min); //"p<q"<=>"px<qx"
    double d=norm(p[s]-p[t]);
    for(int i=2;i<n;S[m++]=i++)REP(j,m){
        if(norm(p[S[j]]-p[i])<d)d=norm(p[s=S[j]]-p[t=i]);
        if(real(p[S[j]])<real(p[i])-d)S[j--]=S[--m];}
    return make_pair(p[s],p[t]);}

int main(){
    P p1(0,1);

```

5.13. Segments.

```

//Interseccion recta y segmento
bool intersectLS (const line &l, const line &s)
{
    return cross(l[1]-l[0], s[0]-l[0]) * //s[0] is left of l
           cross(l[1]-l[0], s[1]-l[0]) <EPS; //s[1] is right of l
}

bool intersectSS (const line &s, const line &t)
{
    return ccw (s[0], s[1], t[0]) * ccw (s[0], s[1], t[1]) <= 0 &&
           ccw (t[0], t[1], s[0]) * ccw (t[0], t[1], s[1]) <= 0;
}

bool intersectPS (const point &p, const line &s)
{
    return abs(s[0]-p) + abs(s[1]-p) - abs(s[1]-s[0]) < EPS; // triangle inequality
}

double distanceLS (const line &l, const line &s)
{
    if (intersectLS (l, s)) return 0;
    return min(distancePL (s[0], l), distancePL (s[1], l));
}

double distancePS (const point &p, const line &s)
{
    const point r = projectionPL(p, s);
    if (intersectPS(r, s)) return abs(r - p);
    return min (abs (s[0] - p), abs (s[1] - p));
}

```

```

P p2(-1,0);
P p3(0,2);
P p4(4,2);
P p5(1,0);
P p6(-1,0);
vector<L> segs; vector<P> out;
segs.push_back(L(p1,p2));
segs.push_back(L(p3,p4));
segs.push_back(L(p5,p6));
cout << segment_intersects(segs,out)<< endl;
FOR(i,out)cout<<*i<<endl;
return 0;
}

```

```

}

double distanceSS (const line &s, const line &t)
{
    if (intersectSS(s, t)) return 0;
    return min (min (distancePS (t[0], s), distancePS (t[1], s)),
                min (distancePS (s[0], t), distancePS (s[1], t)));
}

point projectionPS(const point &p, const line &l)
{
    double r = dot(l[1] - l[0], l[1]-l[0]);
    if (cmp_double(r,0) == 0) return l[0];
    r = dot(p-l[0], l[1]-l[0]) / r;
    if (r < 0) return l[0];
    if (r > 1) return l[1];
    return l[0] + (l[1] - l[0]) * r;
}

bool merge_if_able (line &s, line t)
{
    if (abs( cross( s[1]-s[0], t[1]-t[0]) ) > EPS )
        return false;

    if( ccw(s[0], t[0], s[1]) == +1 || ccw(s[0], t[0], s[1]) == -1)
        return false; //nsame line
}

```

```

if(ccw(s[0], s[1], t[0]) == -2 || ccw(t[0], t[1], s[0]) == -2)
    return false; //separated

s = line(min(s[0], t[0], cmp_point), max(s[1], t[1], cmp_point));

return true;
}

/*
    Tested: STRAZA
    Contest 2 - COCI 2006-2007
*/
void merge_segments(vector<line>& segs)
{
    bool changed = true;
    while(changed)

```

```

{
    changed = false;
    for(int i = 0; i < (int)segs.size(); ++i)
        for(int j = i+1; j < (int)segs.size(); ++j)
        {
            line a = segs[i], b = segs[j];
            if(merge_if_able(segs[i], segs[j]))
            {
                changed = true;
                segs.erase(segs.begin() + j);
                break;
            }
        }
    }
}

```

5.14. Semiplane Intersection.

```

/*
    Check wether there is a point in the intersection of
    several semi-planes. if p lies in the border of some
    semiplane it is considered to belong to the semiplane.

    Expected Running time: linear

    Tested on Triathlon [Cuban Campament Contest]
*/
bool intersect( vector<line> semiplane ){

    function<bool(line&,point&)> side = [](line &l, point &p){
        // IMPORTANT: point p belongs to semiplane defined by l
        // iff p it's clockwise respect to segment < l.p, l.q >
        // i.e. (non negative cross product)

        return cross( l.q - l.p, p - l.p ) >= 0;
    };

    function<bool(line&, line&, point&)> crosspoint = [](const line &l, const line &m, point &x){
        double A = cross(l.q - l.p, m.q - m.p);
        double B = cross(l.q - l.p, l.q - m.p);
        if (abs(A) < eps) return false;
        x = m.p + B / A * (m.q - m.p);
        return true;
    };

```

```

int n = (int)sempiplane.size();

random_shuffle( semiplane.begin(), semiplane.end() );

point cent(0, 1e9);

for (int i = 0; i < n; ++i){
    line &S = semiplane[ i ];

    if (side(S, cent)) continue;

    point d = S.q - S.p; d /= abs( d );

    point A = S.p - d * 1e8, B = S.p + d * 1e8;

    for (int j = 0; j < i; ++j){
        point x;
        line &T = semiplane[j];

        if ( crosspoint(T, S, x) ){
            int cnt = 0;

            if (!side(T, A)){
                A = x;
                cnt++;
            }
        }
    }
}

```

```

    }

    if (!side(T, B)){
        B = x;
        cnt++;
    }

    if (cnt == 2)
        return false;
    }
    else{

```

5.15. Triangles.

```

double area_heron(double const &a, double const &b, double const &c)
{
    double s=(a+b+c)/2;
    return sqrt(s*(s-a)*(s-b)*(s-c));
}

double circumradius(const double &a, const double &b, const double &c)
{
    return a*b*c/4/area_heron(a,b,c);
}

double inradius(const double &a, const double &b, const double &c)
{
    return 2*area_heron(a,b,c)/(a+b+c);
}

/*
Center of the circumference of a triangle
[Tested COJ 1572 - Joining the Centers]
*/
point circumference_center(point a, point b, point c)
{
    point x = 1.0 / conj(b - a), y = 1.0 / conj(c - a);
    return (y - x) / (conj(x) * y - x * conj(y)) + a;
}

bool circumference_center(point &a, point &b, point &c, point &r)
{
    double d = (a.x() * (b.y() - c.y()) + b.x() * (c.y()
        - a.y()) + c.x() * (a.y() - b.y())) * 2.0;

    if(fabs(d) < EPS)
        return false;

```

```

        if (!side(T, A)) return false;
    }

    if (imag(B) > imag(A)) swap(A, B);
    cent = A;
}

return true;
}

```

```

r.x() = ((a.x() * a.x() + a.y() * a.y()) * (b.y() - c.y())
    + (b.x() * b.x() + b.y() * b.y()) * (c.y() - a.y()) +
    (c.x() * c.x() + c.y() * c.y()) * (a.y() - b.y())) / d;
r.y() = -((a.x() * a.x() + a.y() * a.y()) * (b.x() - c.x()) +
    (b.x() * b.x() + b.y() * b.y()) * (c.x() - a.x())
    + (c.x() * c.x() + c.y() * c.y()) * (a.x() - b.x())) / d;

return true;
}

/*
//Interseccion de las bisectrices
double incenter(vect &a,vect &b,vect &c,vect &r)
{
    double u=(b-c).length(),v=(c-a).length(),w=(a-b).length(),s=u+v+w;
    if(s<EPS) {r=a;return 0.0;}
    r.x=(a.x*u+b.x*v+c.x*w)/s;
    r.y=(a.y*u+b.y*v+c.y*w)/s;
    return sqrt((v+w-u)*(w+u-v)*(u+v-w)/s)*0.5;
}

//Interseccion de las alturas
bool orthocenter(vect &a,vect &b,vect &c,vect &r)
{
    double d=a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y);
    if(fabs(d)<EPS) return false;
    r.x=((c.x*b.x+c.y*b.y)*(c.y-b.y)+(a.x*c.x+a.y*c.y)*(a.y-c.y)
        + (b.x*a.x+b.y*a.y)*(b.y-a.y))/d;
    r.y=-((c.x*b.x+c.y*b.y)*(c.x-b.x)+(a.x*c.x+a.y*c.y)*(a.x-c.x)
        + (b.x*a.x+b.y*a.y)*(b.x-a.x))/d;
    return true;
}
*/

```

```
double signed_area(const point &p1, const point &p2, const point &p3)
{
    return cross(p2-p1,p3-p1);
}
```

```
double triangle_area(const point &a, const point &b,const point &c)
{
    return 0.5* abs( cross(b-a,c-a) );
}
```

6. GRAPHS

6.1. Articulation Point And Bridge.

```

const int
    MaxV = 10005;

enum { White, Gray, Black };

vi g[MaxV];
int d[MaxV], low[MaxV], pi[MaxV];
int step = 0;
bool puntoArticulacion[MaxV];
set<pii> aristaPuente;
int dfsRoot, rootChildren;

int n, m;

void DFS(int u)
{
    low[u] = d[u] = ++step;
    REP(i, g[u].size())
    {
        int v = g[u][i];
        if(d[v] == White)
        {
            pi[v] = u;
            if(u == dfsRoot)
                ++rootChildren;
        }
    }
}

```

```

DFS(v);

if(low[v] >= d[u]) //for articulation point
    puntoArticulacion[u] = true;
if(low[v] > d[u]) //for bridge
    aristaPuente.insert(pii(u, v));

low[u] = min(low[u], low[v]);
}
else if(v != pi[u])
    low[u] = min(low[u], d[v]);
}
}

void articulationPointAndBridge()
{
    step = 0;
    REP(i, n-1){
        if(d[i] == White){
            dfsRoot = i;
            rootChildren = 0;
            DFS(i);
            puntoArticulacion[dfsRoot] = (rootChildren > 1);
        }
    }
}

```

6.2. Bellman-Ford.

```

int n, m;

const int
    MaxN = 1000;

vector<pii> g[MaxN];

struct Edge{

    int src, dst, weight;
    Edge(int a, int b, int c) :
        src(a), dst(b), weight(c) {}
}

```

```

};

vector<Edge> edges;
vector<int> dist;

bool Bellman_Ford(int s)
{
    dist = vector<int>(n, oo);
    dist[s] = 0;

    for(int i = 0; i < n-1; ++i)
    {

```

```

    foreach(e, edges)
    {
        int u = e -> src;
        int v = e -> dst;
        int w = e -> weight;

        if(dist[u] + w < dist[v])
            dist[v] = dist[u] + w;
    }

    foreach(e, edges)

```

```

{
    int u = e -> src;
    int v = e -> dst;
    int w = e -> weight;

    if(dist[u] + w < dist[v])
        return false; //negative_cycle_exist
}
return true;
}

```

6.3. Biconnected Components.

```

const int
    MaxN = 10000;

int n, m;
vector<int> g[MaxN];
int d[MaxN], low[MaxN], pi[MaxN];
int step;

stack<pii> bicon;

void BiconComp(int u)
{
    d[u] = low[u] = ++step;
    REP(i, g[u].size())
    {
        int w = g[u][i];
        if(w != pi[u] && d[w] < d[u]) //forward edge
        {
            bicon.push(pii(u, w));
            if(d[w] == 0)
            {
                BiconComp(w);
                low[u] = min(low[u], low[w]);
                if(low[w] >= d[u])
                {
                    printf("New_Biconnected_Component:\n");

```

```

                pii tmp;
                do{
                    tmp = bicon.top(); bicon.pop();
                    printf("%d_%d\n", tmp.F, tmp.S);
                }
                while(!(tmp.F == u && tmp.S == w));
                printf("\n");
            }
        }
        else if(w != pi[u]) //back edge
            low[u] = min(low[u], d[w]);
    }
}

void init()
{
    step = 0;
    REP(i, n)
    {
        d[i] = low[i] = 0;
        pi[i] = -1;
    }
}

```


6.4. Bipartite Matching.

```

/*
   Tested: AIZU(judge.u-aizu.ac.jp) GRL_7_A
   Complexity: O(nm)
*/

struct graph
{
    int L, R;
    vector<vector<int>> adj;

    graph(int L, int R) : L(L), R(R), adj(L + R) {}

    void add_edge(int u, int v)
    {
        adj[u].push_back(v + L);
        adj[v + L].push_back(u);
    }

    int maximum_matching()
    {
        vector<int> visited(L), mate(L + R, -1);
        function<bool(int)> augment = [&](int u)
        {
            if (visited[u]) return false;

```

```

            visited[u] = true;
            for (int w : adj[u])
            {
                int v = mate[w];
                if (v < 0 || augment(v))
                {
                    mate[u] = w;
                    mate[w] = u;
                    return true;
                }
            }
            return false;
        };
        int match = 0;
        for (int u = 0; u < L; ++u)
        {
            fill(visited.begin(), visited.end(), 0);
            if (augment(u))
                ++match;
        }
        return match;
    }
};

```

6.5. Centroid Decomposition.

```

/*
   Centroid decomposition of a tree.
   Find the centroid of the subtree that contains node c.

   Nodes availables are those which aren't marked, i.e mk[u] == False
*/

vi adj[maxn];
bool mk[maxn];
int q[maxn], p[maxn], sz[maxn], mc[maxn];

int centroid(int c){
    int b = 0, e = 0;
    q[e++] = c, p[c] = -1, sz[c] = 1, mc[c] = 0;

    while (b < e){

```

```

        int u = q[b++];
        for (auto v : adj[u]) if (v != p[u] && !mk[v])
            p[v] = u, sz[v] = 1, mc[v] = 0, q[e++] = v;
    }

    for (int i = e - 1; ~i; --i){
        int u = q[i];
        int bc = max(e - sz[u], mc[u]);
        if (2 * bc <= e) return u;
        sz[p[u]] += sz[u], mc[p[u]] = max(mc[p[u]], sz[u]);
    }

    assert(false);
    return -1;
}

```

6.6. Dijkstra.

```

const int
    MaxN = 1000;

int n, m;

vector<pii> g[MaxN];
int pi[MaxN];

priority_queue<pii, vector<pii>, greater<pii>> > pq;

vector<int> Dijkstra(int s)
{
    vector<int> d(n, oo);
    pq = priority_queue<ii, vector<ii>, greater<ii>> >();

    d[s] = 0;
    pq.push(ii(0, s));

    while(!pq.empty())

```

```

{
    int dist = pq.top().F, u = pq.top().S;
    pq.pop();

    if(dist == d[u])
    {
        for(int i = 0; i < (int)g[u].size(); ++i)
        {
            int v = g[u][i].S;
            int w = g[u][i].F;
            if(d[u] + w < d[v])
            {
                d[v] = d[u] + w;
                pq.push(ii(d[v], v));
            }
        }
    }
}

return d;
}

```

6.7. Dominator Tree.

```

/*
    Dominator Tree (Lengauer-Tarjan)

    Tested: SPOJ EN
    Complexity: O(m log n)
*/

struct graph
{
    int n;
    vector<vector<int>> adj, radj;

    graph(int n) : n(n), adj(n), radj(n) {}

    void add_edge(int src, int dst)
    {
        adj[src].push_back(dst);
        radj[dst].push_back(src);
    }
}

```

```

vector<int> rank, semi, low, anc;

int eval(int v)
{
    if (anc[v] < n && anc[anc[v]] < n)
    {
        int x = eval(anc[v]);
        if (rank[semi[low[v]]] > rank[semi[x]])
            low[v] = x;
        anc[v] = anc[anc[v]];
    }
    return low[v];
}

vector<int> prev, ord;

void dfs(int u)
{
    rank[u] = ord.size();
    ord.push_back(u);
}

```

```

        for (auto v : adj[u])
        {
            if (rank[v] < n)
                continue;
            dfs(v);
            prev[v] = u;
        }
    }

    vector<int> idom; // idom[u] is an immediate dominator of u

    void dominator_tree(int r)
    {
        idom.assign(n, n);
        prev = rank = anc = idom;
        semi.resize(n);
        iota(semi.begin(), semi.end(), 0);
        low = semi;
        ord.clear();
        dfs(r);

        vector<vector<int>> dom(n);
        for (int i = (int) ord.size() - 1; i >= 1; --i)
        {
            int w = ord[i];
            for (auto v : radj[w])
            {
                int u = eval(v);
                if (rank[semi[w]] > rank[semi[u]])

```

```

                semi[w] = semi[u];
            }
            dom[semi[w]].push_back(w);
            anc[w] = prev[w];
            for (int v : dom[prev[w]])
            {
                int u = eval(v);
                idom[v] = (rank[prev[w]] > rank[semi[u]]
                    ? u : prev[w]);
            }
            dom[prev[w]].clear();
        }

        for (int i = 1; i < (int) ord.size(); ++i)
        {
            int w = ord[i];
            if (idom[w] != semi[w])
                idom[w] = idom[idom[w]];
        }
    }

    vector<int> dominators(int u)
    {
        vector<int> S;
        for (; u < n; u = idom[u])
            S.push_back(u);
        return S;
    }
};

```

6.8. Flow with Lower Bound.

```

/*
    Flow with lower bound

    Tested: ZOJ 3229
    Complexity:  $O(n^2 m)$ 
*/

template<typename T>
struct dinic
{
    struct edge
    {
        int src, dst;
        T low, cap, flow;
    };

```

```

        int rev;
    };

    int n;
    vector<vector<edge>> adj;

    dinic(int n) : n(n), adj(n + 2) {}

    void add_edge(int src, int dst, T low, T cap)
    {
        adj[src].push_back({ src, dst, low, cap, 0, (int) adj[dst].size() });
        if (src == dst)
            adj[src].back().rev++;
        adj[dst].push_back({ dst, src, 0, 0, 0, (int) adj[src].size() - 1 });
    }

```

```

}

vector<int> level, iter;

T augment(int u, int t, T cur)
{
    if (u == t)
        return cur;
    for (int &i = iter[u]; i < (int) adj[u].size(); ++i)
    {
        edge &e = adj[u][i];
        if (e.cap - e.flow > 0 && level[u] > level[e.dst])
        {
            T f = augment(e.dst, t, min(cur, e.cap - e.flow));
            if (f > 0)
            {
                e.flow += f;
                adj[e.dst][e.rev].flow -= f;
                return f;
            }
        }
    }
    return 0;
}

int bfs(int s, int t)
{
    level.assign(n + 2, n + 2);
    level[t] = 0;
    queue<int> Q;
    for (Q.push(t); !Q.empty(); Q.pop())
    {
        int u = Q.front();
        if (u == s)
            break;
        for (edge &e : adj[u])
        {
            edge &erev = adj[e.dst][e.rev];
            if (erev.cap - erev.flow > 0
                && level[e.dst] > level[u] + 1)
            {
                Q.push(e.dst);
                level[e.dst] = level[u] + 1;
            }
        }
    }
    return level[s];
}

```

```

}

const T oo = numeric_limits<T>::max();

T max_flow(int source, int sink)
{
    vector<T> delta(n + 2);

    for (int u = 0; u < n; ++u) // initialize
        for (auto &e : adj[u])
        {
            delta[e.src] -= e.low;
            delta[e.dst] += e.low;
            e.cap -= e.low;
            e.flow = 0;
        }

    T sum = 0;
    int s = n, t = n + 1;

    for (int u = 0; u < n; ++u)
    {
        if (delta[u] > 0)
        {
            add_edge(s, u, 0, delta[u]);
            sum += delta[u];
        }
        else if (delta[u] < 0)
            add_edge(u, t, 0, -delta[u]);
    }

    add_edge(sink, source, 0, oo);
    T flow = 0;

    while (bfs(s, t) < n + 2)
    {
        iter.assign(n + 2, 0);
        for (T f; (f = augment(s, t, oo)) > 0;)
            flow += f;
    }

    if (flow != sum)
        return -1; // no solution

    for (int u = 0; u < n; ++u)
        for (auto &e : adj[u])
        {

```

```

        e.cap += e.low;
        e.flow += e.low;
        edge &erev = adj[e.dst][e.rev];
        erev.cap -= e.low;
        erev.flow -= e.low;
    }

    adj[sink].pop_back();
    adj[source].pop_back();

```

6.9. Floyd Warshall.

```

const int
    MaxN = 10000;

int n, m;
int g[MaxN][MaxN];
int dist[MaxN][MaxN];

void Floyd_Warshall()
{
    for(int k = 0; k < n; ++k)
        for(int i = 0; i < n; ++i)
            for(int j = 0; j < n; ++j)
                dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j]);
}

```

6.10. Gabow Edmonds.

```

/*
    Tested: Timus 1099
    Complexity:  $O(n^3)$ 
*/

struct graph
{
    int n;
    vector<vector<int>>> adj;

    graph(int n) : n(n), adj(n) {}

    void add_edge(int u, int v)
    {

```

```

        while (bfs(source, sink) < n + 2)
        {
            iter.assign(n + 2, 0);
            for (T f; (f = augment(source, sink, oo)) > 0;)
                flow += f;
        } // level[u] == n + 2 ==> s-side

        return flow;
    }
};

```

```

int init()
{
    REP(i, n) REP(j, n)
    {
        if(dist[i][j] == 0)
        {
            dist[i][j] = oo;
            g[i][j] = oo;
        }
    }

    for(int i = 0; i < n; ++i)
        dist[i][i] = 0;
}

```

```

        adj[u].push_back(v);
        adj[v].push_back(u);
    }

    queue<int> q;
    vector<int> label, mate, cycle;

    void rematch(int x, int y)
    {
        int m = mate[x];
        mate[x] = y;
        if (mate[m] == x)
        {
            if (label[x] < n)

```

```

        rematch(mate[m] = label[x], m);
    else
    {
        int s = (label[x] - n) / n, t = (label[x] - n) % n;
        rematch(s, t);
        rematch(t, s);
    }
}

void traverse(int x)
{
    vector<int> save = mate;
    rematch(x, x);
    for (int u = 0; u < n; ++u)
        if (mate[u] != save[u])
            cycle[u] ^= 1;
    save.swap(mate);
}

void relabel(int x, int y)
{
    cycle = vector<int>(n, 0);
    traverse(x);
    traverse(y);
    for (int u = 0; u < n; ++u)
    {
        if (!cycle[u] || label[u] >= 0)
            continue;
        label[u] = n + x + y * n;
        q.push(u);
    }
}

int augment(int r)
{

```

```

    label.assign(n, -2);
    label[r] = -1;
    q = queue<int>();
    for (q.push(r); !q.empty(); q.pop())
    {
        int x = q.front();
        for (int y : adj[x])
        {
            if (mate[y] < 0 && r != y)
            {
                rematch(mate[y] = x, y);
                return 1;
            }
            else if (label[y] >= -1)
                relabel(x, y);
            else if (label[mate[y]] < -1)
            {
                label[mate[y]] = x;
                q.push(mate[y]);
            }
        }
    }
    return 0;
}

int maximum_matching()
{
    mate.assign(n, -2);
    int matching = 0;
    for (int u = 0; u < n; ++u)
        if (mate[u] < 0)
            matching += augment(u);
    return matching;
}
};

```

6.11. Gomory hu Tree.

```

/*
    Gomory-Hu tree

    Tested: SPOJ MCQUERY
    Complexity: O(n-1) max-flow call
*/

```

```

template<typename flow_type>
struct edge
{
    int src, dst;
    flow_type cap;
};

```

```

template<typename flow_type>
vector<edge<flow_type>> gomory_hu(dinic<flow_type> &adj)
{
    int n = adj.n;

    vector<edge<flow_type>> tree;
    vector<int> parent(n);

    for (int u = 1; u < n; ++u)

```

```

{
    tree.push_back({ u, parent[u], adj.max_flow(u, parent[u]) });
    for (int v = u + 1; v < n; ++v)
        if (adj.level[v] == -1 && parent[v] == parent[u])
            parent[v] = u;
}

return tree;
}

```

6.12. Hopcroft Karp.

```

/*
    Tested: SPOJ MATCHING
    Complexity:  $O(m \cdot n^{0.5})$ 
*/

struct graph
{
    int L, R;
    vector<vector<int>> adj;

    graph(int L, int R) : L(L), R(R), adj(L + R) {}

    void add_edge(int u, int v)
    {
        adj[u].push_back(v + L);
        adj[v + L].push_back(u);
    }

    int maximum_matching()
    {
        vector<int> level(L), mate(L + R, -1);

        function<bool(void)> levelize = [&]()
        {
            queue<int> Q;
            for (int u = 0; u < L; ++u)
            {
                level[u] = -1;
                if (mate[u] < 0)
                {
                    level[u] = 0;
                    Q.push(u);
                }
            }
        }
    }
}

```

```

while (!Q.empty())
{
    int u = Q.front(); Q.pop();
    for (int w : adj[u])
    {
        int v = mate[w];
        if (v < 0) return true;
        if (level[v] < 0)
        {
            level[v] = level[u] + 1;
            Q.push(v);
        }
    }
}

return false;
};

function<bool(int)> augment = [&](int u)
{
    for (int w : adj[u])
    {
        int v = mate[w];
        if (v < 0 || (level[v] > level[u] && augment(v)))
        {
            mate[u] = w;
            mate[w] = u;
            return true;
        }
    }
    return false;
};

int match = 0;
while (levelize())
    for (int u = 0; u < L; ++u)

```

```

        if (mate[u] < 0 && augment(u))
            ++match;
    return match;

```

6.13. Hungarian.

```

/*
Maximum assignment (Kuhn-Munkres)

Description:
- We are given a cost table of size n times m with n <= m.
- It finds a maximum cost assignment, i.e.,
    max sum_{i,j} c(i,j) x(i,j)
    where sum_{i in [n]} x(i,j) = 1,
    sum_{j in [n]} x(i,j) <= 1.

Complexity: O(n^3)

Tested: http://www.spoj.com/problems/SCITIIES/
*/

template<typename T>
T max_assignment(const vector<vector<T>> &a)
{
    int n = a.size(), m = a[0].size();
    assert(n <= m);

    vector<int> x(n, -1), y(m, -1);
    vector<T> px(n, numeric_limits<T>::min()), py(m, 0);

    for (int u = 0; u < n; ++u)
        for (int v = 0; v < m; ++v) px[u] = max(px[u], a[u][v]);

    for (int u = 0, p, q; u < n; )
    {
        vector<int> s(n + 1, u), t(m, -1);

        for (p = q = 0; p <= q && x[u] < 0; ++p)
            for (int k = s[p], v = 0; v < m && x[u] < 0; ++v)

```

```

    }
};

```

```

        if (px[k] + py[v] == a[k][v] && t[v] < 0)
        {
            s[++q] = y[v], t[v] = k;
            if (s[q] < 0)
                for (p = v; p >= 0; v = p)
                    y[v] = k = t[v], p = x[k], x[k] = v;
        }

        if (x[u] < 0)
        {
            T delta = numeric_limits<T>::max();

            for (int i = 0; i <= q; ++i)
                for (int v = 0; v < m; ++v) if (t[v] < 0)
                    delta = min(delta, px[s[i]] + py[v] - a[s[i]][v]);

            for (int i = 0; i <= q; ++i)
                px[s[i]] -= delta;

            for (int v = 0; v < m; ++v)
                py[v] += (t[v] < 0 ? 0 : delta);
        }
        else ++u;
    }

    T cost = 0;

    for (int u = 0; u < n; ++u)
        cost += a[u][x[u]];

    return cost;
}

```

6.14. Kruskal.

```

struct Edge{
    int src, dst, weight;
    Edge(int a, int b, int c):

```

```

        src(a), dst(b), weight(c){}
};

```



```

const int
    MaxN = 10000;

vector<Edge> mst;
vector<Edge> edge;

bool cmp(Edge x, Edge y)
{
    return x.weight < y.weight;
}

int cost = 0;
void Kruskal()
{
    mst.clear();
    initDisjointSet();

```

```

sort(ALL(edge), cmp);

for(int i = 0; i < (int)edge.size(); ++i)
{
    int u = edge[i].src;
    int v = edge[i].dst;
    if(SetOf(u) != SetOf(v))
    {
        cost += edge[i].weight;
        Merge(u, v);
    }
}

```

6.15. Max Flow Dinic.

```

/*
    Maximum Flow (Dinitz)

    Complexity:  $O(n^2 m)$  but very fast in practice

    Tested: http://www.spoj.com/problems/FASTFLOW/
*/

template<typename flow_type>
struct dinic
{
    struct edge
    {
        size_t src, dst, rev;
        flow_type flow, cap;
    };

    int n;
    vector<vector<edge>> adj;

    dinic(int n) : n(n), adj(n), level(n), q(n), it(n) {}

    void add_edge(size_t src, size_t dst, flow_type cap, flow_type rcap = 0)
    {
        adj[src].push_back({src, dst, adj[dst].size(), 0, cap});
        if (src == dst) adj[src].back().rev++;
        adj[dst].push_back({dst, src, adj[src].size() - 1, 0, rcap});
    }

```

```

}

vector<int> level, q, it;

bool bfs(int source, int sink)
{
    fill(level.begin(), level.end(), -1);
    for (int qf = level[q[0] = sink] = 0, qb = 1; qf < qb; ++qf)
    {
        sink = q[qf];
        for (edge &e : adj[sink])
        {
            edge &r = adj[e.dst][e.rev];
            if (r.flow < r.cap && level[e.dst] == -1)
                level[q[qb++] = e.dst] = 1 + level[sink];
        }
    }
    return level[source] != -1;
}

flow_type augment(int source, int sink, flow_type flow)
{
    if (source == sink) return flow;
    for (; it[source] != adj[source].size(); ++it[source])
    {
        edge &e = adj[source][it[source]];
        if (e.flow < e.cap && level[e.dst] + 1 == level[source])

```

```

        {
            flow_type delta = augment(e.dst, sink,
                                     min(flow, e.cap - e.flow));
            if (delta > 0)
            {
                e.flow += delta;
                adj[e.dst][e.rev].flow -= delta;
                return delta;
            }
        }
        return 0;
    }

    flow_type max_flow(int source, int sink)
    {

```

```

        for (int u = 0; u < n; ++u)
            for (edge &e : adj[u]) e.flow = 0;
        flow_type flow = 0;
        flow_type oo = numeric_limits<flow_type>::max();

        while (bfs(source, sink))
        {
            fill(it.begin(), it.end(), 0);
            for (flow_type f; (f = augment(source, sink, oo)) > 0;)
                flow += f;

            // level[u] = -1 => source side of min cut
            return flow;
        }
    };

```

6.16. Max Flow push relabel.

```

/*
    Maximum Flow (Goldberg-Tarjan)

    Complexity:  $O(n^3)$  faster than Dinic in most cases

    Tested: http://www.spoj.com/problems/FASTFLOW/
*/

template<typename flow_type>
struct goldberg_tarjan
{
    struct edge
    {
        size_t src, dst, rev;
        flow_type flow, cap;
    };

    int n;
    vector<vector<edge>> adj;

    goldberg_tarjan(int n) : n(n), adj(n) {}

    void add_edge(size_t src, size_t dst, flow_type cap, flow_type rcap = 0)
    {
        adj[src].push_back({ src, dst, adj[dst].size(), 0, cap });
        if (src == dst) adj[src].back().rev++;
        adj[dst].push_back({ dst, src, adj[src].size() - 1, 0, rcap });
    }

```

```

}

flow_type max_flow(int source, int sink)
{
    vector<flow_type> excess(n);
    vector<int> dist(n), active(n), count(2 * n);
    queue<int> q;
    auto enqueue = [&](int v)
    {
        if (!active[v] && excess[v] > 0)
        {
            active[v] = true;
            q.push(v);
        }
    };
    auto push = [&](edge &e)
    {
        flow_type f = min(excess[e.src], e.cap - e.flow);
        if (dist[e.src] <= dist[e.dst] || f == 0) return;
        e.flow += f;
        adj[e.dst][e.rev].flow -= f;
        excess[e.dst] += f;
        excess[e.src] -= f;
        enqueue(e.dst);
    };
    dist[source] = n;
    active[source] = active[sink] = true;

```

```

count[0] = n - 1;
count[n] = 1;
for (int u = 0; u < n; ++u)
    for (edge &e : adj[u]) e.flow = 0;
for (edge &e : adj[source])
{
    excess[source] += e.cap;
    push(e);
}
for (int u; !q.empty(); q.pop())
{
    active[u = q.front()] = false;
    for (auto &e : adj[u]) push(e);
    if (excess[u] > 0)
    {
        if (count[dist[u]] == 1)
        {
            int k = dist[u]; // Gap Heuristics
            for (int v = 0; v < n; v++)
            {
                if (dist[v] < k)
                    continue;
                count[dist[v]]--;
                dist[v] = max(dist[v], n + 1);
            }
        }
    }
}

```

```

        count[dist[v]]++;
        enqueue(v);
    }
}
else
{
    count[dist[u]]--; // Relabel
    dist[u] = 2 * n;
    for (edge &e : adj[u])
        if (e.cap > e.flow)
            dist[u] = min(dist[u],
                           dist[e.dst] + 1);
    count[dist[u]]++;
    enqueue(u);
}
}
}
flow_type flow = 0;
for (edge e : adj[source])
    flow += e.flow;
return flow;
}
};

```

6.17. Min Cost Max Flow.

```

/*
    Minimum Cost Flow (Tomizawa, Edmonds-Karp)

    Complexity:  $O(F m \log n)$ , where  $F$  is the amount of maximum flow

    Tested: Codeforces [http://codeforces.com/problemset/problem/717/G]
*/

template<typename flow_type, typename cost_type>
struct min_cost_max_flow
{
    struct edge
    {
        size_t src, dst, rev;
        flow_type flow, cap;
        cost_type cost;
    };

    int n;

```

```

    vector<vector<edge>> adj;

    min_cost_max_flow(int n) : n(n), adj(n), potential(n), dist(n), back(n) {}

    void add_edge(size_t src, size_t dst, flow_type cap, cost_type cost)
    {
        adj[src].push_back({src, dst, adj[dst].size(), 0, cap, cost});
        if (src == dst)
            adj[src].back().rev++;
        adj[dst].push_back({dst, src, adj[src].size() - 1, 0, 0, -cost});
    }

    vector<cost_type> potential;

    inline cost_type rcost(const edge &e)
    {
        return e.cost + potential[e.src] - potential[e.dst];
    }
}

```

```

void bellman_ford(int source)
{
    for (int k = 0; k < n; ++k)
        for (int u = 0; u < n; ++u)
            for (edge &e : adj[u])
                if (e.cap > 0 && rcost(e) < 0)
                    potential[e.dst] += rcost(e);
}

const cost_type oo = numeric_limits<cost_type>::max();

vector<cost_type> dist;
vector<edge*> back;

cost_type dijkstra(int source, int sink)
{
    fill(dist.begin(), dist.end(), oo);

    typedef pair<cost_type, int> node;
    priority_queue<node, vector<node>, greater<node>> pq;

    for (pq.push({dist[source] = 0, source}); !pq.empty();)
    {
        node p = pq.top(); pq.pop();

        if (dist[p.second] < p.first) continue;
        if (p.second == sink) break;

        for (edge &e : adj[p.second])
            if (e.flow < e.cap &&
                dist[e.dst] > dist[e.src] + rcost(e))
            {
                back[e.dst] = &e;
                pq.push({dist[e.dst] = dist[e.src] + rcost(e),
                        e.dst});
            }
    }
}

```

```

    return dist[sink];
}

pair<flow_type, cost_type> max_flow(int source, int sink)
{
    flow_type flow = 0;
    cost_type cost = 0;

    for (int u = 0; u < n; ++u)
        for (edge &e : adj[u]) e.flow = 0;

    potential.assign(n, 0);
    dist.assign(n, 0);
    back.assign(n, nullptr);

    bellman_ford(source); // remove negative costs

    while (dijkstra(source, sink) < oo)
    {
        for (int u = 0; u < n; ++u)
            if (dist[u] < dist[sink])
                potential[u] += dist[u] - dist[sink];

        flow_type f = numeric_limits<flow_type>::max();

        for (edge *e = back[sink]; e; e = back[e->src])
            f = min(f, e->cap - e->flow);
        for (edge *e = back[sink]; e; e = back[e->src])
            e->flow += f, adj[e->dst][e->rev].flow -= f;

        flow += f;
        cost += f * (potential[sink] - potential[source]);
    }
    return {flow, cost};
};

```

6.18. Prim.

```

const int
    MaxN = 10000;

int n, m;
typedef pair<int, pii> par;

```

```

priority_queue<par, vector<par>, greater<par>> pq;
vi taken;
vector<pii> g[MaxN];
int mstCost;
vector<pii> mstEdge;

```

```

void process(int u)
{
    taken[u] = 1;
    for(int i = 0; i < (int)g[u].size(); ++i)
    {
        pii v = g[u][i];
        if(!taken[v.S])
            pq.push(par(v.F, pii(u, v.S)));
    }
}

void Prim(int s)
{
    taken.assign(n, 0);
    pq = priority_queue<par, vector<par>, greater<par>>>();

    process(s);
    mstCost = 0;
}

```

```

while(!pq.empty())
{
    par top = pq.top(); pq.pop();
    pii node = top.S;

    int w = top.F;
    int u = node.F;
    int v = node.S;

    if(!taken[v])
    {
        mstCost += w;
        mstEdge.pb(pii(u, v));
        process(v);
    }
}
}

```

6.19. Satisfiability Two SAT.

```

/*
    Two-Sat

    Complexity: O(n)

    Tested: POI (Gates)
*/

struct satisfiability_twosat
{
    int n;
    vector<vector<int>> imp;

    satisfiability_twosat(int n) : n(n), imp(2 * n) {}

    void add_edge(int u, int v)
    {
        imp[u].push_back(v);
    }

    int neg(int u) { return (n << 1) - u - 1; }

    void implication(int u, int v)
    {

```

```

        add_edge(u, v);
        add_edge(neg(v), neg(u));
    }

    vector<bool> solve()
    {
        int size = 2 * n;
        vector<int> S, B, I(size);

        function<void(int)> dfs = [&](int u)
        {
            B.push_back(I[u] = S.size());
            S.push_back(u);

            for (int v : imp[u])
                if (!I[v]) dfs(v);
                else while (I[v] < B.back()) B.pop_back();

            if (I[u] == B.back())
                for (B.pop_back(), ++size; I[u] < S.size(); S.pop_back())
                    I[S.back()] = size;
        };

        for (int u = 0; u < 2 * n; ++u)

```

```

        if (!I[u]) dfs(u);

vector<bool> values(n);

for (int u = 0; u < n; ++u)
    if (I[u] == I[neg(u)]) return {};

```

6.20. Strongly Connected Components.

```

const int
    MaxN = 10000;

struct edge{

    int src, dst, w;
    edge(int a, int b, int c): src(a), dst(b), w(c){}

};

typedef vector<edge> Graph;
int n, m;
Graph g[MaxN];
Graph gt[MaxN];
int order[MaxN], mk[MaxN];
int scc[MaxN];
int vcount[MaxN];
int cur;
int cur_scc;

void dfs(int u)
{
    mk[u] = true;
    for(int i = 0; i < (int)g[u].size(); ++i)
    {
        int v = g[u][i].dst;
        if(!mk[v])
            dfs(v);
    }
    order[n-1-cur++] = u;
}

void dfs_rev(int u)
{
    scc[u] = cur_scc;
    ++vcount[cur_scc];
    mk[u] = true;

```

```

        else values[u] = I[u] < I[neg(u)];

    return values;
}
};

```

```

    for(int i = 0; i < (int)gt[u].size(); ++i)
    {
        int v = gt[u][i].dst;
        if(!mk[v])
            dfs_rev(v);
    }

void make_scc()
{
    cur = 0;
    memset(mk, 0, sizeof(mk));
    for(int i = 0; i < n; ++i)
        if(!mk[i])
            dfs(i);

    cur_scc = 0;
    memset(mk, 0, sizeof(mk));

    for(int i = 0; i < n; ++i)
    {
        int v = order[i];
        if(!mk[v])
        {
            dfs_rev(v);
            ++cur_scc;
        }
    }

void init()
{
    for(int i = 0; i < n; ++i)
    {
        g[i].clear();
        gt[i].clear();

```

```

        vcount[i] = 0;
    }

```

6.21. SCC Gabow.

```

/*
    Gabow's strongly connected component

    Complexity:  $O(n + m)$ 

    Tested: http://www.spoj.com/problems/CAPCITY/
*/

struct graph
{
    int n;
    vector<vector<int>>> adj;

    graph(int n) : n(n), adj(n) {}

    void add_edge(int u, int v)
    {
        adj[u].push_back(v);
    }

    vector<int>& operator[](int u) { return adj[u]; }
};

vector<vector<int>>> scc_gabow(graph &adj)
{
    int n = adj.n;

    vector<vector<int>>> scc;

```

6.22. Stoer Wagner.

```

/*
    Tested: ZOJ 2753
    Complexity:  $O(n^3)$ 
*/

template<typename T>
pair<T, vector<int>>> stoer_wagner(vector<vector<T>>> &weights)
{
    int n = weights.size();

```

```

    }

```

```

vector<int> S, B, I(n);

function<void(int)> dfs = [&](int u)
{
    B.push_back(I[u] = S.size());
    S.push_back(u);

    for (int v : adj[u])
        if (!I[v]) dfs(v);
        else while (I[v] < B.back()) B.pop_back();

    if (I[u] == B.back())
    {
        scc.push_back({});
        for (B.pop_back(); I[u] < S.size(); S.pop_back())
        {
            scc.back().push_back(S.back());
            I[S.back()] = n + scc.size();
        }
    }
};

for (int u = 0; u < n; ++u)
    if (!I[u]) dfs(u);

return scc; // in reverse topological order
}

```

```

vector<int> used(n), cut, best_cut;
T best_weight = -1;

for (int phase = n - 1; phase >= 0; --phase)
{
    vector<T> w = weights[0];
    vector<int> added = used;
    int prev, last = 0;

```

```

for (int i = 0; i < phase; ++i)
{
    prev = last;
    last = -1;
    for (int j = 1; j < n; ++j)
        if (!added[j] && (last == -1 || w[j] > w[last]))
            last = j;

    if (i == phase - 1)
    {
        for (int j = 0; j < n; ++j)
            weights[prev][j] += weights[last][j];
        for (int j = 0; j < n; ++j)
            weights[j][prev] = weights[prev][j];

        used[last] = true;
        cut.push_back(last);
    }
}

```

```

    if (best_weight == -1 || w[last] < best_weight)
    {
        best_cut = cut;
        best_weight = w[last];
    }
    else
    {
        for (int j = 0; j < n; ++j)
            w[j] += weights[last][j];
        added[last] = true;
    }
}

return make_pair(best_weight, best_cut);
}

```

6.23. Tree Isomorphism.

```

/*
    Tested: SPOJ TREEISO
    Complexity: O(n log n)
*/

#define all(c) (c).begin(), (c).end()

struct tree
{
    int n;
    vector<vector<int>> adj;

    tree(int n) : n(n), adj(n) {}

    void add_edge(int src, int dst)
    {
        adj[src].push_back(dst);
        adj[dst].push_back(src);
    }

    vector<int> centers()
    {
        vector<int> prev;
        int u = 0;
        for (int k = 0; k < 2; ++k)

```

```

        {
            queue<int> q;
            prev.assign(n, -1);
            for (q.push(prev[u] = u); !q.empty(); q.pop())
            {
                u = q.front();
                for (auto v : adj[u])
                {
                    if (prev[v] >= 0)
                        continue;
                    q.push(v);
                    prev[v] = u;
                }
            }
        }

        vector<int> path = { u };
        while (u != prev[u])
            path.push_back(u = prev[u]);

        int m = path.size();
        if (m % 2 == 0)
            return {path[m/2-1], path[m/2]};
        else
            return {path[m/2]};
    }
}

```



```

    }

    vector<vector<int>> layer;
    vector<int> prev;

    int levelize(int r)
    {
        prev.assign(n, -1);
        prev[r] = n;
        layer = {{r}};
        while (1)
        {
            vector<int> next;
            for (int u : layer.back())
                for (int v : adj[u])
                {
                    if (prev[v] >= 0)
                        continue;
                    prev[v] = u;
                    next.push_back(v);
                }

            if (next.empty())
                break;
            layer.push_back(next);
        }
        return layer.size();
    }
};

bool isomorphic(tree S, int s, tree T, int t)
{
    if (S.n != T.n)
        return false;
    if (S.levelize(s) != T.levelize(t))
        return false;

    vector<vector<int>> longcodeS(S.n + 1), longcodeT(T.n + 1);
    vector<int> codeS(S.n), codeT(T.n);
    for (int h = (int) S.layer.size() - 1; h >= 0; --h)

```

```

    {
        map<vector<int>, int> bucket;
        for (int u : S.layer[h])
        {
            sort(all(longcodeS[u]));
            bucket[longcodeS[u]] = 0;
        }
        for (int u : T.layer[h])
        {
            sort(all(longcodeT[u]));
            bucket[longcodeT[u]] = 0;
        }

        int id = 0;
        for (auto &p : bucket)
            p.second = id++;
        for (int u : S.layer[h])
        {
            codeS[u] = bucket[longcodeS[u]];
            longcodeS[S.prev[u]].push_back(codeS[u]);
        }
        for (int u : T.layer[h])
        {
            codeT[u] = bucket[longcodeT[u]];
            longcodeT[T.prev[u]].push_back(codeT[u]);
        }
    }

    return codeS[s] == codeT[t];
}

bool isomorphic(tree S, tree T)
{
    auto x = S.centers(), y = T.centers();
    if (x.size() != y.size())
        return false;
    if (isomorphic(S, x[0], T, y[0]))
        return true;
    return x.size() > 1 && isomorphic(S, x[1], T, y[0]);
}

```

7. MATRIX

7.1. Gauss.

```

/*
[TESTED COJ 2536 05/11/2014]
*/
const int MAXN = 110;
const int oo = (1<<30);
const double EPS = 1e-6;

double a[MAXN][MAXN];
double ans[MAXN];

int n; //ecuations
int m; //variables

void init(int _n, int _m)
{
    n = _n;
    m = _m;
    memset(a, 0, sizeof a);
    memset(ans, 0, sizeof ans);
}

int solve()
{
    vector<int> where (m, -1);
    for (int col = 0, row = 0; col < m && row < n; ++col)
    {
        int sel = row;
        for (int i = row; i < n; ++i)
            if (abs (a[i][col]) > abs (a[sel][col]))
                sel = i;

        if (abs (a[sel][col]) < EPS)
            continue;
    }
}

```

7.2. Gauss Modulo 2.

```

/*
[TESTED: SPOJ XMAX, LightOJ 1272,1288
Matrix: (2) sol|x1 x2...xn
Answer: ans[vars-1...0]

```

```

    for (int i = col; i <= m; ++i)
        swap (a[sel][i], a[row][i]);

    where[col] = row;

    for (int i = 0; i < n; ++i)
    {
        if (i != row)
        {
            double c = a[i][col] / a[row][col];
            for (int j = col; j <= m; ++j)
                a[i][j] -= a[row][j] * c;
        }
        ++row;
    }

    for (int i = 0; i < m; ++i)
        if (where[i] != -1)
            ans[i] = a[where[i]][m] / a[where[i]][i];

    for (int i = 0; i < n; ++i)
    {
        double sum = 0;
        for (int j = 0; j < m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
            return 0;
    }

    for (int i = 0; i < m; ++i)
        if (where[i] == -1)
            return oo;

    return 1;
}

```

```

*/
const int MAXN = 110;
const int MAXR = 70;

```

```

bitset<MAXN> row[MAXR];

int ans[MAXN];
int first[MAXR];
int vars;
int rows;

void init(int _vars)
{
    vars = _vars;
    rows = 0;
}

bool add(bitset<MAXN> cur)
{
    for(int i = 0; i < rows; i++)
    {
        if(cur[first[i]] != 0)
        {
            cur ^= row[i];
        }
    }
    first[rows] = 0;
    while(first[rows] < vars && !cur[first[rows]])
        first[rows]++;

    /*remove if want to add always the equation*/
    if(first[rows] == vars && cur[vars]) return false;
    row[rows++] = cur;
    return true;
}

void solve()
{
    memset(ans, 0, sizeof ans);

```

```

        for(int i = rows-1; i >= 0; i--)
        {
            int aux = row[i][vars];
            for(int j = first[i]; j < vars; j++)
                aux ^= (ans[j] * row[i][j]);
            ans[first[i]] = aux;
        }
    }

int main() {
    init(3);

    bitset<MAXN> eq1(14), eq2(3), eq3(4);
    /*
        1/1 1 0
        0/0 1 1
        0/1 0 0
        -----
        Ans:0 1 1
    */
    cout << add(eq1);
    cout << add(eq2);
    cout << add(eq3) << endl;

    solve();

    for(int i = vars-1; i >= 0; --i)
        cout << ans[i] << "_";

    return 0;
}

```

7.3. Matrix Template.

```

#define maxn 500

template<class T>
struct Matrix{
    vector< vector<T> > data;
    int m, n;

```

```

Matrix(int m, int n)
{
    this->m = m;
    this->n = n;
    data = vector< vector<T> >(m);
    for(int i = 0; i < m; ++i)
        data[i] = vector<T>(n, 0);
}

```

```

void ident()
{
    for(int i = 0; i < m; ++i)
        data[i][i] = 1;
}

Matrix<T> operator *(Matrix<T> &mtx)
{
    Matrix<T> ans(m, mtx.n);
    for(int i = 0; i < ans.m; ++i)
        for(int j = 0; j < ans.n; ++j)
            for(int k = 0; k < n; ++k)
                ans.data[i][j] += data[i][k] * mtx.data[k][j];

    return ans;
}

Matrix<T> operator ^(int exp)
{
    Matrix<T> ret(m, n);
    Matrix<T> a = *this;

    ret.ident();

    if(exp == 0) return ret;
    if(exp == 1) return a;

    while(exp)
    {
        if(exp & 1)
            ret = ret*a;
        a = (a*a);
        exp >>= 1;
    }
    return ret;
}

};

template<class T>
istream& operator >>( istream &in, Matrix<T> &mtx )
{
    for( int i = 0; i < mtx.m; ++i )
        for(int j = 0; j < mtx.n; ++j)
            in >> mtx.data[i][j];
    return in;
}

```

```

template<class T>
ostream& operator <<( ostream &out, Matrix<T> &mtx )
{
    for( int i = 0; i < mtx.m; ++i )
    {
        for( int j = 0; j < mtx.n; ++j )
        {
            if( j ) out << "_";
            out << mtx.data[i][j];
        }
        out << endl;
    }
    return out;
}

const double eps = 1e-7;

//Determinante
template<class T>
double det(Matrix<T> M0)
{
    double ans = 1;
    int size = M0.m;

    for(int i = 0, r = 0; i < size; ++i){
        bool found = false;

        for(int j = r; j < size; ++j)
            if(fabs(M0.data[j][i]) > eps){
                found = true;

                if(j>r) ans = -ans;
                else break;

                for(int k = 0; k < size; ++k)
                    swap(M0.data[r][k], M0.data[j][k]);
                break;
            }
        if(found){
            for(int j = r + 1; j < size; ++j){
                double aux = M0.data[j][i] / M0.data[r][i];
                for(int k = i; k < size; ++k)
                    M0.data[j][k] -= aux * M0.data[r][k];
            }
            r++;
        }
    }
}

```

```
    }  
    else return 0;  
}
```

```
for(int i = 0; i < size; ++i)  
    ans *= M0.data[i][i];  
return ans;  
}
```

8. NUMBER THEORY

8.1. Binomial Coefficient.

```

/*
  CALCULA COMBINATORIA DE n en k
  USANDO EL TRIANGULO DE PASCAL
*/
#include <cstdio>
#include <iostream>

#define MAX 10000

using namespace std;

int C[MAX][MAX];

void Pascal(int level){
    for(int n = 0; n <= level; ++n){

```

```

        C[n][0] = C[n][n] = 1;
        for(int k = 1; k < n; ++k)
            C[n][k] = C[n-1][k] + C[n-1][k-1];
    }
}

int main()
{
    int n,k; cin>>n>>k;
    Pascal(n);
    cout << C[n][k];

    return 0;
}

```

8.2. Divisibility.

```

pair<vector<int>,int> rmatrix(int base,int div)
{
    vector<int> vis(div,-1);
    vector<int> res;
    res.push_back(1);
    vis[1]=0;

    while(vis[(res[res.size() - 1] * base) % div] == -1)
    {
        vis[(res[res.size() - 1] * base) % div] = res.size();
        res.push_back( (res[res.size() - 1] * base) % div);
    }
    return make_pair(res, vis[(res[res.size() - 1] * base) % div]);
}

```

```

bool div(int base,int div,vector<int> &num)//reverse num
{
    pair<vector<int>,int> r = rmatrix(base,div);
    int pp = 0, b = r.second;
    vector<int> a = r.first;
    for(int i = 0; i < num.size(); ++i)
    {
        int kk = num[i];
        if(i < b)
            pp += ( (kk * a[i]) % div );
        else
            pp += ( (kk * a[b + ( (i-b) % (a.size() - b) ) ]) % div);
    }
    return pp % div == 0;
}

```

8.3. ALL Number Theory.

```

/*
  Binary Multiplication
  [Tested Timus 1141,1204]**

```

```

*/
Int mod_mult(Int a, Int b, Int mod)
{

```

```

    Int x = 0;
    while(b){
        if(b & 1) x = (x + a) % mod;
        a = (a << 1) % mod;
        b >>= 1;
    }
    return x;
}

/*
    Binary Exponentiation
    [Tested Timus 1141,1204]**
*/
Int mod_pow(Int a, Int n, Int mod)
{
    Int x = 1;
    while(n){
        if(n & 1) x = mod_mult(x, a, mod);
        a = mod_mult(a, a, mod);
        n >>= 1;
    }
    return x;
}

/*
    Extended Euclidean algorithm
    Solve ax+by = (a,b)
    Works well even for negative numbers
    [Tested Timus 1141,1204]**
*/
int gcd(int a,int b,int &x,int &y)
{
    if(b==0) {x = 1; y = 0; return a;}
    int r = gcd(b, a%b, y, x);
    y -= a/b*x;
    return r;
}

/*
    Euler's function
    phi(p^a) = p^a - p^(a-1)
    (a,b) = 1 => phi(a*b) = phi(a)*phi(b)
    [Tested Timus 1141]*
*/
int phi(int a)
{
    int b = a;

```

```

    for(int i = 2; i*i <= a; ++i)
        if(a % i == 0)
        {
            b = b/i*(i-1);
            do a/=i;
            while(a%i==0);
        }
    if(a > 1) b = b/a*(a-1);
    return b;
}

/*
    Modular Inverse
    (a,m) = 1
    Solves a*x = 1 (m)
    [Tested Timus 1141, 1204]**
*/
int inverse(int a, int m)
{
    int x, y ;
    if(gcd( a, m, x, y ) != 1) return 0;
    return (x%m + m) % m;
}

/*
    Baby-Step-Giant-Step Algorithm
    O(sqrt(m) log(m))
    Solve a^x = b(mod m)
    [TESTED LightOJ 1325 05/11/2014]
*/
Int discrete_log(Int a, Int b, Int m)
{
    map<Int, Int> hash;
    Int n = phi(m), k = sqrt(n);

    for(Int i = 0,t = 1; i < k; i++)
    {
        hash[t] = i;
        t = (t * a) % m;
    }
    Int c = mod_pow(a, n - k, m);
    for(Int i = 0; i * k < n; i++)
    {
        if(hash.find(b) != hash.end())
            return (i * k + hash[b]) % n;

        b = (b * c) % m;
    }
}

```

```

    }
    return -1;
}

/*
    Solves  $a \cdot x = b \pmod{p}$ 
    [Tested CodeChef Quadratic Equations]
*/
long solve_linear(long a, long b, int p)
{
    return (b*inverse(a,p)) % p;
}

/*
    Solve  $x \equiv a_i \pmod{m_i}$ 
    For any  $i$  and  $j$ ,  $(m_i, m_j) \mid a_i - a_j$ .
    Return  $x_0$  in  $[0, M]$ .
     $M = \text{mlm2}..mn$ 
    All solutions are  $x = x_0 + t[M]$ .
*/
int linear_con(int a[], int m[], int n)
{
    int u = a[0], v = m[0], p, q, r, t;
    for(int i = 1; i < n; i++)
    {
        r = gcd(v, m[i], p, q);
        t = v;
        v = v / r * m[i];
        u = ( (a[i] - u) / r * p * t + u ) % v;
    }
    if(u < 0) u += v;
    return u;
}

/*
    Solve  $x \equiv a_i \pmod{m_i}$ 
    For any  $i$  and  $j$ ,  $(m_i, m_j) = 1$ .
    Returns  $x_0$  in  $[0, M]$ .
     $M = \text{mlm2}..mn$ 
    All solutions are  $x = x_0 + tM$ .
*/
int chinese(int a[], int m[], int n)
{
    int s = 1, t, ans = 0, p, q;
    for(int i = 0; i < n; i++) s *= m[i];
    for(int i = 0; i < n; i++){
        t = s / m[i];

```

```

        gcd(t, m[i], p, q);
        ans = (ans + t * p * a[i]) % s;
    }
    if(ans < 0) ans += s;
    return ans;
}

/*
    Kth discrete roots of  $a \pmod{n}$ 
 $x^k = a \pmod{n}$ 
    When  $(k, \phi(n)) = 1$ 
    [Tested Timus 1141]**
*/
int discrete_root(int k, int a, int n)
{
    int _phi = phi(n);
    int s = (int)inverse(k, _phi);
    return (int)mod_pow(a, s, n);
}

/*
    Tonelli Shank's algorithm
    Solves  $x^2 \equiv a \pmod{p}$ 
    [Tested CodeChef Quadratic Equations, Timus 1132]
    Warning: Precompute primes to avoid TLE
*/
int solve_quadratic(int a, int p)
{
    if(a == 0) return 0;
    if(p == 2) return a;
    if(mod_pow(a, (p-1)/2, p) != 1) return -1;

    int phi = p-1;
    int n = 0, k = 0;

    while(phi % 2 == 0)
    {
        phi /= 2;
        n++;
    }

    k = phi;
    int q = 0;

    for(int j = 2; j < p; j++)
        if(mod_pow(j, (p-1)/2, p) == p-1)
        {

```



```

        q = j; break;
    }

    int t = mod_pow(a, (k+1)/2, p);
    int r = mod_pow(a, k, p);

    while(r!=1)
    {
        int i = 0, v = 1;
        while(mod_pow(r, v, p) != 1)
        {
            v*=2;
            i++;
        }

        int e = mod_pow(2, n-i-1, p);
        int u = mod_pow(q, k*e, p);

        t = (t*u)%p;
        r = (r*u*u)%p;
    }

    return t;
}

/*
Solves  $a*x^2 + b*x + c = 0 \pmod{p}$ 
[Tested CodeChef Quadratic Equations]
*/
set<Int> solve_quadratic(Int a, Int b, Int c, int p)
{
    set<Int> ans;
    if(c==0) ans.insert(0L);
    if(a==0) ans.insert(solve_linear((p-b)%p, c, p));
    else if(p==2 && (a+b+c)%2==0) ans.insert(1L);
    else
    {
        Int r = ((b*b)%p - (4*a*c)%p + p)%p;
        Int x = solve_quadratic(r, p);
        if(x == -1) return ans;
        Int w = solve_linear((2*a)%p, (x-b+p)%p, p);
        ans.insert(w);
        w = solve_linear((2*a)%p, (p-x-b+p)%p, p);
        ans.insert(w);
    }
    return ans;
}

```

```

/*
Primitive roots
[Tested Timus 1268]
Warning: Precompute primes to avoid TLE
Only:  $m = 1, p^k, n = 2p^k$  ( $p$  prime  $> 2$ ),
       $m = 2, m = 4$ 
*/
int primitive_root(int m, int p[])
{
    if(m == 1) return 0;
    if(m == 2) return 1;
    if(m == 4) return 3;

    int t = m;
    if((t&1) == 0) t >>= 1;

    for(int i = 0; p[i]*p[i] <= t; ++i)
    {
        if(t % p[i]) continue;
        do t /= p[i]; while(t % p[i] == 0);
        if(t > 1 || p[i] == 2) return 0;
    }

    int f[100];
    int x = phi(m), y = x, n=0;

    for(int i = 0; p[i]*p[i] <= y; ++i)
    {
        if(y % p[i]) continue;
        do y /= p[i]; while(y % p[i] == 0);
        f[n++] = p[i];
    }

    if(y > 1) f[n++] = y;

    for(int i = 1; i < m; ++i)
    {
        if(__gcd(i, m) > 1) continue;
        bool flag = true;

        for(int j = 0; j < n; ++j)
            if(mod_pow(i, x/f[j], m) == 1)
            {
                flag = false;
                break;
            }
    }
}

```

```

    }

    if(flag)
        return i;
    }
    return 0;
}

typedef long long ll;

ll divisor_sigma(ll n)
{
    ll sigma = 0, d = 1;
    for (; d * d < n; ++d)
        if (n % d == 0)
            sigma += d + n / d;
    if (d * d == n)
        sigma += d;
    return sigma;
}

// sigma(n) for all n in [lo, hi)
vector<ll> divisor_sigma(ll lo, ll hi)
{
    vector<ll> ps = primes(sqrt(hi) + 1);
    vector<ll> res(hi - lo), sigma(hi - lo, 1);
    iota(res.begin(), res.end(), lo);
    for (ll p : ps)
        for (ll k = ((lo + (p - 1)) / p) * p; k < hi; k += p)
        {
            ll b = 1;
            while (res[k - lo] > 1 && res[k - lo] % p == 0)
            {
                res[k - lo] /= p;
                b = 1 + b * p;
            }
            sigma[k - lo] *= b;
        }
    for (ll k = lo; k < hi; ++k)
        if (res[k - lo] > 1)
            sigma[k - lo] *= (1 + res[k - lo]);
    return sigma; // sigma[k-lo] = sigma(k)
}

```

```

typedef long long ll;

ll mobius_mu(ll n)
{
    if (n == 0)
        return 0;
    ll mu = 1;
    for (ll x = 2; x * x <= n; ++x)
        if (n % x == 0)
        {
            mu = -mu;
            n /= x;
            if (n % x == 0)
                return 0;
        }
    return n > 1 ? -mu : mu;
}

// phi(n) for all n in [lo, hi)
vector<ll> mobius_mu(ll lo, ll hi)
{
    vector<ll> ps = primes(sqrt(hi) + 1);
    vector<ll> res(hi - lo), mu(hi - lo, 1);
    iota(res.begin(), res.end(), lo);
    for (ll p : ps)
        for (ll k = ((lo + (p - 1)) / p) * p; k < hi; k += p)
        {
            mu[k - lo] = -mu[k - lo];
            if (res[k - lo] % p == 0)
            {
                res[k - lo] /= p;
                if (res[k - lo] % p == 0)
                {
                    mu[k - lo] = 0;
                    res[k - lo] = 1;
                }
            }
        }
    for (ll k = lo; k < hi; ++k)
        if (res[k - lo] > 1)
            mu[k - lo] = -mu[k - lo];
    return mu; // mu[k-lo] = mu(k)
}

```

8.4. Prime.

```

const int N=16000000;
const int sqrtN=sqrt(N);
bool isP[N];

O(N log log N)
void sieve()
{
    fill(isP, isP+N, true);
    isP[0] = isP[1] = false;

    for(int i = 4; i < N; i += 2)
        isP[i] = false;

    for(int i = 3; i < sqrtN; i += 2)
        if(isP[i])
            for(int j = i*i; j < N; j += 2*i)
                isP[j] = false;
}

/*
    Binary Multiplication
    [Tested Timus 1141,1204]**
*/
int mod_mult(Int a, Int b, Int mod)
{
    Int x = 0;
    while(b){
        if(b & 1) x = (x + a) % mod;
        a = (a << 1) % mod;
        b >>= 1;
    }
    return x;
}

/*
    Binary Exponentiation
    [Tested Timus 1141,1204]**
*/
int mod_pow(Int a, Int n, Int mod)
{
    Int x = 1;
    while(n){
        if(n & 1) x = mod_mult(x, a, mod);

```

```

        a = mod_mult(a, a, mod);
        n >>= 1;
    }
    return x;
}

/*
    Miller Rabin
    [Tested SPOJ PON]
*/
bool witness(Int a, Int s, Int d, Int n)
{
    Int x = mod_pow(a, d, n);
    if (x == 1 || x == n - 1) return false;
    for(int i = 0; i < s - 1; i++)
    {
        x = mod_mult(x, x, n);
        if (x == 1) return true;
        if (x == n - 1) return false;
    }
    return true;
}

bool isPrime(Int n)
{
    if (n < 2) return false;
    if (n == 2) return true;
    if (n % 2 == 0) return false;
    Int d = n - 1, s = 0;
    while (d % 2 == 0) ++s, d /= 2;
    Int test[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 0};
    for (int i = 0; test[i] && test[i] < n; ++i)
        if (witness(test[i], s, d, n))
            return false; // composite
    return true; // probably prime
}

/*
    Integer Factorization Pollard's Rho
    */
uint64 pollar_rho(uint64 n) //n shouldn't be prime
{
    if(!(n&1)) return 2;

```

```

while(true)
{
    uint64 x = (uint64) rand() % n, y = x, c = rand() % n;

    if(c == 0 || c == 2)
        c = 1;

    for(int i = 1, k = 2; ; i++)
    {
        x = mod_mult(x, x, n);
        if(x >= c) x -= c;
        else x+=n-c;
        if(x == n) x=0;
        if(x == 0) x = n-1;
        else x--;

        uint64 d = __gcd(x > y ? x - y : y - x, n);

        if(d == n) break;
        if(d != 1) return d;
        if(i == k)
        {
            y = x;
            k <<= 1;
        }
    }
}

//fact primos de n
vector<pair<Int,Int> > fact(Int n)
{
    vector<pair<Int,Int> > fp;
    for(int i = 2; i <= n; ++i)
    {
        pair<Int,Int> pp = make_pair(i,0);
        while(!(n%i))
        {
            n /= i;
            pp.second++;
        }
        if(pp.second)
            fp.push_back(pp);
    }

    if(n > 1)

```

```

        fp.push_back(make_pair(n,1));

    return fp;
}

vector<Int> primes;

//fact primos de n!
vector<pair<Int,Int> > factF(Int n)
{
    vector<pair<Int,Int> > fp;
    Int p;
    for(int i = 0; i < (int) primes.size(); ++i)
    {
        p = primes[i];
        if(p > n)
            break;

        Int k = n;
        pair<Int,Int> pp = make_pair(p,0);
        while(k)
        {
            pp.second += k / p;
            k /= p;
        }
        fp.push_back(pp);
    }

    return fp;
}

/*
    Tested: SPOJ PRIME1, ETFS
    Complexity: O(n log log n)
*/

typedef long long ll;

// primes in [lo, hi)
vector<ll> primes(ll lo, ll hi)
{
    const ll M = 1 << 14, SQR = 1 << 16;
    vector<bool> composite(M), small_composite(SQR);
    vector<pair<ll, ll>> sieve;
    for (ll i = 3; i < SQR; i += 2)
        if (!small_composite[i])

```

```

        {
            ll k = i * i + 2 * i * max(0.0, ceil((lo - i*i)/(2.0*i)));
            sieve.push_back({ 2 * i, k });
            for (ll j = i * i; j < SQR; j += 2 * i)
                small_composite[j] = 1;
        }
        vector<ll> ps;
        if (lo <= 2)
        {
            ps.push_back(2);
            lo = 3;
        }
        for (ll k = lo | 1, low = lo; low < hi; low += M)
        {
            ll high = min(low + M, hi);

```

```

            fill(composite.begin(), composite.end(), 0);
            for (auto &z : sieve)
                for (; z.second < high; z.second += z.first)
                    composite[z.second - low] = 1;
            for (; k < high; k += 2)
                if (!composite[k - low])
                    ps.push_back(k);
        }
        return ps;
    }

    vector<ll> primes(ll hi)
    {
        return primes(0, hi);
    }

```

8.5. Tree Stern-Brocot.

```

/*
Stern-Brocot Tree for enumerating rationals
Enumerating all irreducible rationals ascending order,
Whose sum of N and D is atmost B
*/
void sternBrocot (Int B, Int pl = 0, Int ql = 1,
                  Int pr = 1, Int qr = 0) {

```

```

    Int pm = pl + pr, qm = ql + qr;
    if (pm + qm > B) return;
    sternBrocot (B, pl, ql, pm, qm); // [pl / ql, pm / qm]
    cout <<pm <<"/" <<qm <<endl;
    sternBrocot (B, pm, qm, pr, qr); // [pm / qm, pr / qr]
}

```

9. NUMERIC METHODS

9.1. Fast Fourier Transform.

```

typedef complex < double > base ;

// y[i] = A(w^(dir*i)),
// w = exp(2pi/N) is N-th complex principal root of unity,
// A(x) = a[0] + a[1] x + ... + a[n-1] x^{n-1},
// * N must be a power of 2,
long double PI = 2 * acos(0.0L);

void fft ( vector < base > &a, bool invert )
{
    int n = ( int ) a. size ( ) ;

    for ( int i = 1 , j = 0 ; i < n ; ++i ) {
        int bit = n >> 1 ;
        for ( ; j >= bit ; bit >>= 1 ) j -= bit ;
        j += bit ;
        if ( i < j ) swap ( a [ i ] , a [ j ] ) ;
    }
    for ( int len = 2 ; len <= n ; len <<= 1 )
    {
        double ang = 2 * PI / len * ( invert ? - 1 : 1 ) ;
        base wlen ( cos ( ang ) , sin ( ang ) ) ;

        for ( int i = 0 ; i < n ; i += len )
        {
            base w(1);
            for ( int j = 0 ; j < len / 2 ; ++j )
            {
                base u = a [i+j], v = a[i+j+len / 2 ] * w ;
                a [ i + j ] = u + v ;
                a [ i + j + len / 2 ] = u - v ;
                w *= wlen ;
            }
        }
    }
}

```

9.2. Goldsection Search.

```

/*
    Minimum of unimodal function (goldsection search)

    Tested: COJ 2890 :(
*/

```

```

    }
}

if (invert)
    for ( int i = 0 ; i < n ; ++i )
        a [ i ] /= n ;
}

void convolve(const vector <int> &a, const vector <int> &b, vector
<int> &res)
{
    vector < base > fa ( a. begin ( ) , a. end ( ) ) , fb (
        b. begin ( ) , b. end ( ) ) ;
    size_t n = 1 ;
    while ( n < max ( a. size ( ) , b. size ( ) ) ) n <<= 1 ;
    n <<= 1 ;
    fa. resize ( n ) , fb. resize ( n ) ;
    fft ( fa, false ) , fft ( fb, false ) ;
    for ( size_t i = 0 ; i < n ; ++i )
        fa [ i ] *= fb [ i ] ;
    fft ( fa, true ) ;
    res. resize ( n ) ;
    for ( size_t i = 0 ; i < n ; ++i )
        res [ i ] = int ( fa [i].real( ) + 0.5 );
}

void print(vector<int> a)
{
    cout << a.size()<<endl;
    for(int i = 0; i < (int)a.size(); ++i)
        cout << a[i] << "_";
    cout << endl;
}

```

```

template<class F>
double find_min(F f, double a, double d, double eps = 1e-9)
{
    const int iter = 150;

```

```

const double r = 2 / (3 + sqrt(5.));
double b = a + r * (d - a), c = d - r * (d - a), fb = f(b), fc = f(c);
for (int it = 0; it < iter && d - a > eps; ++it)
{
    // '<': maximum, '>': minimum
    if (fb > fc)
    {
        a = b;
        b = c;
        c = d - r * (d - a);
        fb = fc;
        fc = f(c);
    }
}

```

```

    }
    else
    {
        d = c;
        c = b;
        b = a + r * (d - a);
        fc = fb;
        fb = f(b);
    }
}
return c;
}

```

9.3. Linear Recursion.

```

/*
    Linear Recurrence Solver

    Description: Consider
     $x[i+n] = a[0] x[i] + a[1] x[i+1] + \dots + a[n-1] x[i+n-1]$ 
    with initial solution  $x[0], x[1], \dots, x[n-1]$ 
    We compute  $k$ -th term of  $x$  in  $O(n^2 \log k)$  time.

    Tested: SPOJ REC
    Complexity:  $O(n^2 \log k)$  time,  $O(n \log k)$  space
*/

typedef long long ll;

ll linear_recurrence(vector<ll> a, vector<ll> x, ll k)
{
    int n = a.size();
    vector<ll> t(2 * n + 1);
    function<vector<ll>(ll)> rec = [&](ll k)
    {
        vector<ll> c(n);
        if (k < n) c[k] = 1;
    }
}

```

```

    else
    {
        vector<ll> b = rec(k / 2);
        fill(t.begin(), t.end(), 0);
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < n; ++j)
                t[i+j+(k&1)] += b[i]*b[j];
        for (int i = 2*n-1; i >= n; --i)
            for (int j = 0; j < n; ++j)
                t[i-n+j] += a[j]*t[i];
        for (int i = 0; i < n; ++i)
            c[i] = t[i];
    }
    return c;
};

vector<ll> c = rec(k);
ll ans = 0;
for (int i = 0; i < x.size(); ++i)
    ans += c[i] * x[i];
return ans;
}

```

9.4. Romberg.

```

const double EPS = 1e-6;

// Romberg
// Assume  $F' = f$ 
// input: interval  $[a,b]$  and a function  $f$ 

```

```

// output:  $F(b)-F(a)$ 

inline int cmp(double x, double y=0)
{
    return (x <= y + EPS) ? (x + EPS < y) ? -1 : 0 : 1;
}

```

```

}

int pow(int a, int n)
{
    int x = 1;
    while(n) {
        if(n & 1) x *= a;
        n >>= 1;
        a *= a;
    }
    return x;
}

long double romberg(int a, int b, double(*func)(double))
{
    long double approx[2][50];
    long double *cur = approx[1], *prev = approx[0];

    prev[0] = 1/2.0 * (b-a) * (func(a) + func(b));

```

9.5. Roots Newton.

```

template<class F, class G>
double find_root(F f, G df, double x)
{
    for (int iter = 0; iter < 100; ++iter)
    {
        double fx = f(x), dfx = df(x);

```

9.6. Simplex.

```

/*
    Parametric Self-Dual Simplex method

    Description:
    - Solve a canonical LP:
        min. c x
    s.t. A x <= b
        x >= 0

    Complexity: O(n+m) iterations on average

    Tested: http://codeforces.com/contest/375/problem/E
*/

```

```

for(int it = 1; it < 25; ++it, swap(cur, prev))
{
    if(it > 1 && cmp(prev[it-1], prev[it-2]) == 0)
        return prev[it-1];

    cur[0] = 1/2.0 * prev[0];
    long double div = (b-a)/pow(2, it);

    for(long double sample = a + div; sample < b; sample += 2 * div)
        cur[0] += div * func(a + sample);

    for(int j = 1; j <= it; ++j)
        cur[j] = cur[j-1] + 1 /
            (pow(4, it) - 1)*(cur[j-1] + prev[j-1]);
}

return prev[24];
}

```

```

x -= fx / dfx;
if (fabs(fx) < 1e-12)
    break;
}
return x;
}

```

```

const double eps = 1e-9, oo = numeric_limits<double>::infinity();

typedef vector<double> vec;
typedef vector<vec> mat;

double simplexMethodPD(mat &A, vec &b, vec &c)
{
    int n = c.size(), m = b.size();
    mat T(m + 1, vec(n + m + 1));
    vector<int> base(n + m), row(m);

    for(int j = 0; j < m; ++j)

```



```

{
    for (int i = 0; i < n; ++i)
        T[j][i] = A[j][i];
    T[j][n + j] = 1;
    base[row[j] = n + j] = 1;
    T[j][n + m] = b[j];
}

for (int i = 0; i < n; ++i)
    T[m][i] = c[i];

while (1)
{
    int p = 0, q = 0;
    for (int i = 0; i < n + m; ++i)
        if (T[m][i] <= T[m][p])
            p = i;

    for (int j = 0; j < m; ++j)
        if (T[j][n + m] <= T[q][n + m])
            q = j;

    double t = min(T[m][p], T[q][n + m]);

    if (t >= -eps)
    {
        vec x(n);
        for (int i = 0; i < m; ++i)
            if (row[i] < n) x[row[i]] = T[i][n + m];
        // x is the solution
        return -T[m][n + m]; // optimal
    }

    if (t < T[q][n + m])
    {
        // tight on c -> primal update
        for (int j = 0; j < m; ++j)
            if (T[j][p] >= eps)
                if (T[j][p] * (T[q][n + m] - t) >=
                    T[q][p] * (T[j][n + m] - t))

```

9.7. Simpson.

```

/*
Tested: COJ
2121 - Environment Protection

```

```

        q = j;

        if (T[q][p] <= eps)
            return oo; // primal infeasible
    }
    else
    {
        // tight on b -> dual update
        for (int i = 0; i < n + m + 1; ++i)
            T[q][i] = -T[q][i];

        for (int i = 0; i < n + m; ++i)
            if (T[q][i] >= eps)
                if (T[q][i] * (T[m][p] - t) >=
                    T[q][p] * (T[m][i] - t))
                    p = i;

        if (T[q][p] <= eps)
            return -oo; // dual infeasible
    }

    for (int i = 0; i < m + n + 1; ++i)
        if (i != p) T[q][i] /= T[q][p];

    T[q][p] = 1; // pivot(q, p)
    base[p] = 1;
    base[row[q]] = 0;
    row[q] = p;

    for (int j = 0; j < m + 1; ++j)
        if (j != q)
        {
            double alpha = T[j][p];
            for (int i = 0; i < n + m + 1; ++i)
                T[j][i] -= T[q][i] * alpha;
        }

    return oo;
}

```

```

*/
//METODO DE SIMPSON 1/3 Compuesta
// a,b: intervalo de integracion

```

```
// n = 10000: numero de pasos (ya multiplicado por 2)
double Simpson(int n, double a, double b, double (*f) (double)){
    double s = 0;
    double h = (double) (b - a) / n;
    for (int i = 0; i <= n; ++i) {
```

```
        double x = a + h * i;
        s += f(x) * ((i==0 || i==n) ? 1 : ((i&1)==0) ? 2 : 4);
    }
    return s*(h/3);
}
```

10. PARSING

10.1. Shunting Yard.

```

enum type { op, value, obracekt, cbracekt }; //types
struct token
{
    string text;
    type ttype;
};

template <typename T>
struct operation
{
    int precedence;
    function<void(stack<T> &s)> operate;
};

void mul(stack<string> &s); //operator
void pluss(stack<string> &s);
void poww(stack<string> &s);

unordered_map<string, operation<string>> operations;

bool rpn(const vector<token> &tokens, queue<token> &rpn)
{
    stack<token> operators;
    for (auto &token : tokens)
    {
        if (token.ttype == value)
            rpn.push(token);
        else if (token.ttype == op)
        {
            while (operators.size() > 0 &&
                    operators.top().ttype != obracekt && operations[token.text].precedence > operations[operators.top().text].precedence)
            {
                rpn.push(operators.top());
                operators.pop();
            }
            operators.push(token);
        }
        else if (token.ttype == obracekt)
            operators.push(token);
        else if (token.ttype == cbracekt)
        {
            while (operators.top().ttype != obracekt)
            {

```

```

                rpn.push(operators.top());
                operators.pop();
            }
            operators.pop();
        }
    }
    while (operators.size() > 0)
    {
        if (operators.top().ttype == obracekt)
            return false;
        rpn.push(operators.top());
        operators.pop();
    }
    return true;
}

template <typename T>
T eval(queue<token> &rpn, bool &ok)
{
    stack<T> result;
    while (rpn.size() > 0)
    {
        auto t = rpn.front();
        rpn.pop();
        if (t.ttype == value)
            result.push(t.text); //parsear t.text
        if (t.ttype == op)
            operations[t.text].operate(result);
        ok = result.size() == 1;
        return result.top();
    }
}

vector<token> lex(const string &str); //lexer

int main()
{
    operations["*"] = {1, poww};
    operations["."] = {2, mul};
    operations["|"] = {3, pluss};
    string str;

```

```
auto toks = lex(str);  
queue<token> q;  
rpn(toks, q);  
bool ok;
```

```
auto result = eval<string>(q, ok);  
cout << result << '\n';  
return 0;  
}
```

11. SORTING-SEARCHING

11.1. Ternary Search.

```
double TernarySearchMin(double l, double r)
{
    while(r - l > EPS){
        double m1 = (2*l + r) / 3.0;
        double m2 = (l + 2*r) / 3.0;

        if(f(m1) < f(m2)) r = m2;
        else l = m1;
    }
    return (l + r) / 2.0;
}

double TernarySearchMax(double l, double r)
{
    while(r - l > EPS){
        double m1 = (2*l + r) / 3.0;
        double m2 = (l + 2*r) / 3.0;

        if(f1(m1) < f1(m2)) l = m1;
```

```
        else r = m2;
    }
    return (l + r) / 2.0;
}

//Discrete
int SearchMin(vector<int> &y)
{
    int l = 0, r = y.size()-1;
    while(r - l < 3){
        int m1 = (2*l + r) / 3;
        int m2 = (l + 2*r) / 3;

        if(y[m1] < y[m2]) r = m2;
        else l = m1;
    }
    return min_element(y.begin()+l, y.begin()+r) - y.begin();
}
```

12. STRING

12.1. Aho Corasick.

```

#include <bits/stdc++.h>

using namespace std;

#define endl '\n'
#define DB(x) cout << #x << " = " << x << endl;

const int size = 505;
const int MAXS = size * size + 10;
const int MAXC = 26;

struct aho_corasick
{
    vector<string> key;
    vector<bitset<505> > output;
    vector<int> failure;
    vector<vector<int> > gto;

    int buildMachine()
    {
        int states = 1;
        for(int i = 0; i < key.size(); ++i) {
            const string &word = key[i];
            int currentState = 0;

            for(int j = 0; j < word.size(); ++j) {
                int ch = word[j] - 'a';

                if (gto[currentState][ch] == -1)
                    gto[currentState][ch] = states++;

                currentState = gto[currentState][ch];
            }
            output[currentState].set(i);
        }

        for(int ch = 0; ch < MAXC; ++ch)
            if (gto[0][ch] == -1)
                gto[0][ch] = 0;

        queue<int> q;
        for(int ch = 0; ch < MAXC; ++ch) {
            if(gto[0][ch] != 0){

```

```

                failure[gto[0][ch]] = 0;
                q.push(gto[0][ch]);
            }
        }

        while(!q.empty()) {
            int state = q.front(); q.pop();

            for(int ch = 0; ch < MAXC; ++ch) {
                if(gto[state][ch] != -1) {
                    int f = failure[state];
                    while(gto[f][ch] == -1)
                        f = failure[f];

                    f = gto[f][ch];
                    failure[gto[state][ch]] = f;
                    output[gto[state][ch]] |= output[f];
                    q.push(gto[state][ch]);
                }
            }
        }

        return states;
    }

    aho_corasick(const vector<string> &k) : key(k)
    {
        failure = vector<int>(MAXS, -1);
        gto = vector<vector<int> >(MAXS, vector<int>(MAXC, -1));
        output = vector<bitset<505> >(MAXS);

        buildMachine();
    }

    int nextState(int currentState, char nextInput)
    {
        int state = currentState;
        int ch = nextInput - 'a';
        while(gto[state][ch] == -1) state = failure[state];
        return gto[state][ch];
    }

    vector<int> match(const string &text)
    {

```

```

vector<int> ans(key.size());
int currentState = 0;

for(int i = 0; i < text.size(); ++i) {
    currentState = nextState(currentState, text[i]);

    if(output[currentState].any())
        for(int j = 0; j < key.size(); ++j)
            if(output[currentState].test(j))
                ans[j]++;
}
return ans;
};

int main()
{
    int nc; cin >> nc;

```

12.2. Knuth-Morris-Pratt.

```

//pi[1...m]
vector<int> buildFail(string p)
{
    int m = p.size();
    vector<int> pi(m+1, 0);

    int j = pi[0] = -1;

    for (int i = 1; i <= m; ++i)
    {
        while (j >= 0 && p[j] != p[i - 1])
            j = pi[j];
        pi[i] = ++j;
    }
    return pi;
}

//KMP Cuenta la cantidad de veces que aparece una
//sub-cadena (p) en la cadena (t)

```

12.3. Longest Common Subsequence.

```

#define MAX 100
char X[MAX], Y[MAX];

```

```

for(int tc = 1; tc <= nc; ++tc)
{
    int n; cin >> n;
    string t; cin >> t;
    vector<string> key(n);
    for(int i = 0; i < n; ++i)
        cin >> key[i];

    aho_corasick aho(key);

    cout << "Case_" << tc << ":\n";
    vector<int> ans = aho.match(t);
    for(int i = 0; i < ans.size(); ++i)
        cout << ans[i] << endl;
}

return 0;
}

```

```

int match(string t, string p, vector<int> &pi)
{
    int n = t.size(), m = p.size();
    int count = 0;

    for (int i = 0, k = 0; i < n; ++i)
    {
        while (k >= 0 && p[k] != t[i])
            k = pi[k];
        if (++k >= m)
        {
            ++count;
            k = pi[k];
        }
    }
    return count;
}

```

```

int i, j, m, n, c[MAX][MAX], b[MAX][MAX];

```

```

int LCSlength() {
    m=strlen(X);
    n=strlen(Y);

    for (i=1;i<=m;i++) c[i][0]=0;
    for (j=0;j<=n;j++) c[0][j]=0;

    for (i=1;i<=m;i++)
        for (j=1;j<=n;j++) {
            if (X[i-1]==Y[j-1]) {
                c[i][j]=c[i-1][j-1]+1;
                b[i][j]=1; /* from north west */
            }
            else if (c[i-1][j]>=c[i][j-1]) {
                c[i][j]=c[i-1][j];
                b[i][j]=2; /* from north */
            }
            else {
                c[i][j]=c[i][j-1];
                b[i][j]=3; /* from west */
            }
        }

    return c[m][n];
}
    
```

12.4. Longest Palindrome Substring.

```

// Transform S into T.
// For example, S = "abba", T = "^#a#b#b#a#$".
// ^ and $ signs are sentinels appended to each end to avoid bounds checking
string preProcess(string s)
{
    int n = s.length();
    if (n == 0) return "^$";
    string ret = "^";
    for (int i = 0; i < n; i++)
        ret += "#" + s.substr(i, 1);

    ret += "$";
    return ret;
}

//Time: O(n)
string longestPalindrome(string s)
{
    
```

```

void printLCS(int i,int j) {
    if (i==0 || j==0) return;

    if (b[i][j]==1) {
        printLCS(i-1,j-1);
        printf("%c",X[i-1]);
    }
    else if (b[i][j]==2)
        printLCS(i-1,j);
    else
        printLCS(i,j-1);
}

int main() {
    while (1) {
        gets(X);
        if (feof(stdin)) break; /* press ctrl+z to terminate */
        gets(Y);
        printf("LCS_length->_%d\n",LCSlength()); /* count length */
        printLCS(m,n); /* reconstruct LCS */
        printf("\n");
    }
    return 0;
}
    
```

```

string T = preProcess(s);
int n = T.length();
int *P = new int[n];
int C = 0, R = 0;

for (int i = 1; i < n-1; i++) {
    int i_mirror = 2*C-i; // equals to i' = C - (i-C)

    P[i] = (R > i) ? min(R-i, P[i_mirror]) : 0;

    // Attempt to expand palindrome centered at i
    while (T[i + 1 + P[i]] == T[i - 1 - P[i]])
        P[i]++;

    // If palindrome centered at i expand past R,
    // adjust center based on expanded palindrome.
    if (i + P[i] > R) {
        C = i;
    }
}
    
```



```

    R = i + P[i];
}
}

//Find the maximum element in P.
int maxlen = 0;
int centerIndex = 0;
for (int i = 1; i < n-1; i++) {
    if (P[i] > maxlen) {

```

12.5. Manacher.

```

#define MAX 100
int rank[MAX], LCP [MAX];

// ""[ (i-d)/2 , (i+d)/2 ) "" l[i] = d
vector<int> manacher(string text)
{
    int n = text.size(), i, j, k = 0;
    vector<int> rad(n<<1);

    for(i = 0, j = 0; i < (n<<1); i += k, j = max(j-k, 0))

```

12.6. Maximal Suffix.

```

/*
    Complexity: O(n)
*/

int maximal_suffix(const string &s)
{
    int n = s.length(), i = 0, j = 1;

    for (int k = 0; j < n - 1; k = 0)
    {
        while (j + k < n - 1 && s[i + k] == s[j + k]) ++k;

```

12.7. Minimum Rotation.

```

/*
    Complexity: O(n)
*/

```

```

    maxlen = P[i];
    centerIndex = i;
}
}

delete[] P;

return s.substr((centerIndex - 1 - maxlen)/2, maxlen);
}

```

```

{
    while(i-j >= 0 && i+j+1 < (n<<1) && text[(i-j)>>1] == text[(i+j+1)>>1])
        ++j;
    rad[i] = j;
    for(k = 1; i-k >= 0 && rad[i]-k >= 0 && rad[i-k] != rad[i]-k ; ++k)
        rad[i + k] = min(rad[i-k], rad[i]-k);
}
rad.insert(rad.begin(), 0);
return rad;
}

```

```

    if (s[i + k] < s[j + k])
    {
        i += (k / (j - i) + 1) * (j - i);
        j = i + 1;
    }
    else j += k + 1;
}

return i;
}

```

```

int minimum_rotation(const string &s)
{
    int n = s.length(), i = 0, j = 1, k = 0;

```

```

while (i + k < 2 * n && j + k < 2 * n)
{
    char a = i + k < n ? s[i + k] : s[i + k - n];
    char b = j + k < n ? s[j + k] : s[j + k - n];

    if (a > b)
    {
        i += k + 1;
        k = 0;
        if (i <= j)
            i = j + 1;
    }
}

```

```

else if (a < b)
{
    j += k + 1;
    k = 0;
    if (j <= i)
        j = i + 1;
}
else ++k;
}

return min(i, j);
}

```

12.8. Palindromic Tree.

```

/*
    Palindromic Tree

    Complexity: O(n)

    Tested: ??
*/

template<size_t maxlen, size_t alpha>
struct PalindromicTree
{
    int go[maxlen + 2][alpha], slink[maxlen + 2], length[maxlen + 2];
    int s[maxlen], slength, size, last;

    int new_node()
    {
        memset(go[size], 0, sizeof go[size]);
        slink[size] = length[size] = 0;
        return size++;
    }

    PalindromicTree() { reset(); }

    void reset()
    {
        size = slength = 0;
        length[new_node()] = -1;
    }
}

```

```

        last = new_node();
    }

    int get_link(int p)
    {
        for (int i = slength - 1;
             i - 1 - length[p] < 0 || s[i - 1 - length[p]] != s[i];)
            p = slink[p];
        return p;
    }

    int _extend(int c)
    {
        s[slength++] = c;
        int p = get_link(last), np;
        if (go[p][c]) return go[p][c];
        length[np = new_node()] = 2 + length[p];
        go[p][c] = np;
        if (length[np] == 1) return slink[np] = 1, np;
        p = slink[p];
        slink[np] = go[get_link(p)][c];
        return np;
    }

    void extend(int c) { last = _extend(c); }
};

```

12.9. Substring Palindrome.

```

using System;
namespace hash
{
    class Program
    {
        static int MAXN = 100000 + 10;
        static long[] fh, bh, prime;
        static long mod = 1000000009;
        static long x = 1223;
        static string s;
        static int n;

        static void prime_power(int n)
        {
            prime[0] = 1;
            for (int i = 1; i <= n + 5; i++)
                prime[i] = (prime[i - 1] * x) % mod;
        }

        static void compute_hash(string s)
        {
            for (int i = 1, j = n; i <= n; j--, i++)
            {
                fh[i] = (fh[i - 1] + s[i - 1] * prime[i]) % mod;
                bh[j] = (bh[j + 1] + s[j - 1] * prime[i]) % mod;
            }
        }

        static bool subtring_palindrome(int l, int r)
        {
            ++l; ++r;
            long h1 = (fh[r] - fh[l - 1] + mod) % mod;
            long h2 = (bh[l] - bh[r + 1] + mod) % mod;

            if (l <= n - r + 1)

```

```

        {
            int pow = (n - r + 1) - 1;
            h1 = (h1 * prime[pow]) % mod;
        }
        else
        {
            int pow = 1 - (n - r + 1);
            h2 = (h2 * prime[pow]) % mod;
        }
        return h1 == h2;
    }

    static void Main(string[] args)
    {
        fh = new long[MAXN];
        bh = new long[MAXN];
        prime = new long[MAXN];

        string s = Console.ReadLine();
        n = s.Length;
        prime_power(s.Length);
        compute_hash(s);

        int q = int.Parse(Console.ReadLine());
        for (int i = 0; i < q; ++i)
        {
            int[] query = Array.ConvertAll(
                Console.ReadLine().Split(), int.Parse);
            Console.WriteLine("{0}", subtring_palindrome(query[0],
                query[1]) ? "YES" : "NO");
        }
    }
}

```

12.10. Suffix Array.

```

/*
    Suffix array + lcp

    Complexity: O(n log n)

```

```

Tested:
- http://www.spoj.com/problems/SARRAY/
- http://acm.timus.ru/problem.aspx?space=1&num=1393
- http://wcipeg.com/problem/coci092p6
- http://www.spoj.com/problems/LCS/

```

```

    Note: lcp[i] = lcp(s[sa[i-1]...], s[sa[i]...])
*/

template<typename charT>
struct SuffixArray
{
    int n;
    vector<int> sa, rank, lcp;

    SuffixArray(const basic_string<charT> &s) :
        n(s.length() + 1), sa(n), rank(n), lcp(n)
    {
        vector<int> _sa(n), bucket(n);

        iota(sa.rbegin(), sa.rend(), 0);
        sort(next(sa.begin()), sa.end(),
            [&](int i, int j) { return s[i] < s[j]; });

        for (int i = 1, j = 0; i < n; ++i)
        {
            rank[sa[i]] = rank[sa[i - 1]] +
                (i == 1 || s[sa[i - 1]] < s[sa[i]]);
            if (rank[sa[i]] != rank[sa[i - 1]])
                bucket[++j] = i;
        }

        for (int len = 1; len <= n; len += len)
        {

```

```

            for (int i = 0, j; i < n; ++i)
            {
                if ((j = sa[i] - len) < 0) j += n;
                _sa[bucket[rank[j]]++] = j;
            }

            sa[_sa[bucket[0] = 0]] = 0;

            for (int i = 1, j = 0; i < n; ++i)
            {
                if (rank[_sa[i]] != rank[_sa[i - 1]] ||
                    rank[_sa[i] + len] != rank[_sa[i - 1] + len])
                    bucket[++j] = i;

                sa[_sa[i]] = j;
            }

            copy(sa.begin(), sa.end(), rank.begin());
            sa.swap(_sa);

            if (rank[sa[n - 1]] == n - 1) break;
        }

        for (int i = 0, j = rank[lcp[0] = 0], k = 0; i < n - 1; ++i, ++k)
            while (k >= 0 && s[i] != s[sa[j - 1] + k])
                lcp[j] = k--, j = rank[sa[j] + 1];
    }
};

```

12.11. Suffix Automaton.

```

/*
    Generalized Suffix Automaton

    Complexity: O(n)

    Tested:
    - http://codeforces.com/contest/616/problem/F
    - http://codeforces.com/contest/452/problem/E
    - http://codeforces.com/contest/204/problem/E
*/

template<size_t maxlen, size_t alpha>
struct SuffixAutomaton

```

```

{
    int go[2 * maxlen][alpha], slink[2 * maxlen], length[2 * maxlen];
    int size, last;

    int new_node()
    {
        memset(go[size], 0, sizeof go[size]);
        slink[size] = length[size] = 0;
        return size++;
    }

    SuffixAutomaton() { reset(); }

    void reset()

```

```

{
    size = last = 0;
    new_node();
    slink[0] = -1;
}

int _extend(int c)
{
    int p, q, np, nq;
    if (q = go[last][c])
    {
        if (length[q] == 1 + length[last]) return q;
        int nq = new_node();
        length[nq] = 1 + length[last];
        memcpy(go[nq], go[q], sizeof go[q]);
        slink[nq] = slink[q];
        slink[q] = nq;
        for (p = last; p != -1 && go[p][c] == q; p = slink[p])
            go[p][c] = nq;
        return nq;
    }
    np = new_node();
    length[np] = 1 + length[last];
    for (p = last; p != -1 && !go[p][c]; p = slink[p])
        go[p][c] = np;
    if (p == -1) return slink[np] = 0, np;
    if (length[q = go[p][c]] == 1 + length[p]) return slink[np] = q, np;
    nq = new_node();

```

```

    length[nq] = 1 + length[p];
    memcpy(go[nq], go[q], sizeof go[q]);
    slink[nq] = slink[q];
    slink[q] = slink[np] = nq;
    for (; p != -1 && go[p][c] == q; p = slink[p])
        go[p][c] = nq;
    return np;
}

```

```
void extend(int c) { last = _extend(c); }
```

```
int bucket[maxlen + 1], order[2 * maxlen];
```

```

void top_sort()
{
    int maxl = 0;
    for (int e = 0; e < size; ++e)
        maxl = max(maxl, length[e]);
    for (int l = 0; l <= maxl; ++l)
        bucket[l] = 0;
    for (int e = 0; e < size; ++e)
        ++bucket[length[e]];
    for (int l = 1; l <= maxl; ++l)
        bucket[l] += bucket[l - 1];
    for (int e = 0; e < size; ++e)
        order[--bucket[length[e]]] = e;
}
};

```

12.12. Z Function.

```

//Z[i] is the length of the longest substring
//starting from S[i] which is also a prefix of S.
vector<int> z_function (string s)
{
    int n = (int)s.length();
    vector<int> z(n);

    for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (i <= r)
            z[i] = min (r-i+1, z[i-l]);
        while (i+z[i] < n && s[z[i]] == s[i+z[i]])
            ++z[i];
        if (i+z[i]-1 > r)
            l = i, r = i+z[i]-1;
    }
}

```

```

    return z;
}

// suff[i] = length of the longest common suffix of s and s[0..i]
vector<int> suffixes(const string &s)
{
    int n = s.length();

    vector<int> suff(n, n);

    for (int i = n - 2, g = n - 1, f; i >= 0; --i)
    {
        if (i > g && suff[i + n - 1 - f] != i - g)
            suff[i] = min(suff[i + n - 1 - f], i - g);
        else

```

```
{  
    for (g = min(g, f - i); g >= 0 &&  
        s[g] == s[g + n - 1 - f]; --g);  
    suff[i] = f - g;  
}  
|  
    }  
    return suff;  
}
```

13. MATHEMATICAL FACTS

13.1. **Números de Catalán.** están definidos por la recurrencia:

$$C_{n+1} = \sum_{i=0}^n C_i C_{n-i}$$

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1}$$

13.2. **Números de Stirling de primera clase.** son el número de permutaciones de n elementos con exactamente k ciclos disjuntos.

$$\left[\begin{matrix} n \\ k \end{matrix} \right] = (n-1) \left[\begin{matrix} n-1 \\ k \end{matrix} \right] + \left[\begin{matrix} n-1 \\ k-1 \end{matrix} \right]$$

13.3. **Números de Stirling de segunda clase.** son el número de particionar un conjunto de n elementos en k subconjuntos no vacíos.

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = k \left\{ \begin{matrix} n-1 \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n-1 \\ k-1 \end{matrix} \right\}$$

Además:

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

13.4. **Números de Bell.** cuentan el número de formas de dividir n elementos en subconjuntos.

$$\mathcal{B}_{n+1} = \sum_{k=0}^n \binom{n}{k} \mathcal{B}_k$$

x	0	1	2	3	4	5	6	7	8	9	10
\mathcal{B}_x	1	1	2	5	15	52	203	877	4.140	21.147	115.975

13.5. **Derangement.** permutación que no deja ningún elemento en su lugar original

$$!n = (n-1)(!(n-1) + !(n-2)); !1 = 0, !2 = 1$$

$$!n = n! \sum_{k=0}^n \frac{(-1)^k}{k!}$$

13.6. **Números armónicos.**

$$H_n = \sum_{k=1}^n \frac{1}{k}$$

$$\frac{1}{2n+1} < H_n - \ln n - \gamma < \frac{1}{2n}$$

$$\gamma = 0.577215664901532860606512090082402431042159335 \dots$$

13.7. **Número de Fibonacci.** $f_0 = 0, f_1 = 1$:

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^n$$

$$f_{n+1} = f_n * 2 - f_{n-2}$$

$$f_0 + f_1 + f_2 + \dots + f_n = f_{n+2} - 1$$

$$f_0 - f_1 + f_2 - \dots + (-1)^n f_n = (-1)^n f_{n-1} - 1$$

$$f_1 + f_3 + f_5 + \dots + f_{2n-1} = f_{2n}$$

$$f_0 + f_2 + f_4 + \dots + f_{2n} = f_{2n+1} - 1$$

$$f_0^2 + f_1^2 + f_2^2 + \dots + f_n^2 = f_n f_{n+1}$$

$$f_1 f_2 + f_2 f_3 + f_3 f_4 + \cdots + f_{2n-1} f_n = f_{2n}^2$$

$$f_1 f_2 + f_2 f_3 + f_3 f_4 + \cdots + f_{2n} f_{2n+1} = f_{2n+1}^2 - 1$$

$$k \geq 1 \Rightarrow f_{n+k} = f_k f_{n+1} + f_{k-1} f_n \forall n \geq 0$$

Identidad de Cassini: $f_{n+1} f_n - 1 - f_n^2 = (-1)^n$

$$f_{n+1}^2 + f_n^2 = f_{2n+1}$$

$$f_{n+2}^2 - f_n^2 = f_{2n+2}$$

$$f_{n+2}^2 - f_{n+1}^2 = f_n f_{n+3}$$

$$f_{n+2}^3 - f_{n+1}^3 - f_n^3 = f_{3n+3}$$

$$\text{mcd}(f_n, f_m) = f_{\text{mcd}(n, m)}$$

$$f_{n+1} = \sum_{j=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n-j}{j}$$

$$f_{3n} = \sum_{j=0}^n \binom{n}{j} 2^j f_j$$

El último dígito de cada número se repite periódicamente cada 60 números. Los dos últimos, cada 300; a partir de ahí, se repiten cada $15 * 10^{n-1}$ números.

13.8. Sumas de combinatorios.

$$\sum_{i=n}^m \binom{i}{n} = \binom{m+1}{n+1}$$

$$\sum_{i=0}^k \binom{n}{i} \binom{m}{k-i} = \binom{n+m}{k}$$

13.9. **Funciones generatrices.** Una lista de funciones generatrices para secuencias útiles:

$(1, 1, 1, 1, 1, \dots)$	$\frac{1}{1-z}$
$(1, -1, 1, -1, 1, \dots)$	$\frac{1}{1+z}$
$(1, 0, 1, 0, 1, 0, \dots)$	$\frac{1}{1-z^2}$
$(1, 0, \dots, 0, 1, 0, 1, 0, \dots, 0, 1, 0, \dots)$	$\frac{1}{1-z^2}$
$(1, 2, 3, 4, 5, 6, \dots)$	$\frac{1}{(1-z)^2}$
$(1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \dots)$	$\frac{1}{(1-z)^{m+1}}$
$(1, c, \binom{c+1}{2}, \binom{c+2}{3}, \dots)$	$\frac{1}{(1-z)^c}$
$(1, c, c^2, c^3, \dots)$	$\frac{1}{1-cz}$
$(0, 1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \dots)$	$\ln \frac{1}{1-z}$

Truco de manipulación:

$$\frac{1}{1-z} G(z) = \sum_n \sum_{k \leq n} g_k z^n$$

13.10. **The twelvefold way.** ¿Cuántas funciones $f: N \rightarrow X$ hay?

N	X	Any f	Injective	Surjective
dist.	dist.	x^n	$(x)_n$	$x! \binom{n}{x}$
indist.	dist.	$\binom{x+n-1}{n}$	$\binom{x}{n}$	$\binom{n-1}{n-x}$
dist.	indist.	$\{1\}^n + \dots + \{x\}^n$	$[n \leq x]$	$\{k\}^n$
indist.	indist.	$p_1(n) + \dots + p_x(n)$	$[n \leq x]$	$p_x(n)$

Where $\binom{a}{b} = \frac{1}{b!} (a)_b$ and $p_x(n)$ is the number of ways to partition the integer n using x summands.

13.11. **Teorema de Euler.** si un grafo conexo, plano es dibujado sobre un plano sin intersección de aristas, y siendo v el número de vértices, e el de aristas y f la cantidad de caras (regiones conectadas por aristas, incluyendo la región externa e infinita), entonces

$$v - e + f = 2$$

13.12. **Burnside's Lemma.** Si X es un conjunto finito y G es un grupo de permutaciones que actúa sobre X , sean $S_x = \{g \in G : g * x = x\}$ y $Fix(g) = \{x \in X : g * x = x\}$. Entonces el número de órbitas está

dado por:

$$N = \frac{1}{|G|} \sum_{x \in X} |S_x| = \frac{1}{|G|} \sum_{g \in G} |Fix(g)|$$

13.13. **Ángulo entre dos vectores.** Sea α el ángulo entre \vec{a} y \vec{b} :

$$\cos \alpha = \frac{\vec{a} \cdot \vec{b}}{\|\vec{a}\| \|\vec{b}\|}$$

13.14. **Proyección de un vector.** Proyección de \vec{a} sobre \vec{b} :

$$\text{proy}_{\vec{b}} \vec{a} = \left(\frac{\vec{a} \cdot \vec{b}}{\vec{b} \cdot \vec{b}} \right) \vec{b}$$