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Data Structures (2)

dsu.hpp

0e8ecd, 33 lines

```
struct DSU
{
    int n;
    VI par, siz;

    DSU (int _n = 0)
    {
        n = _n;
        par.resize(n);
        iota(all(par), 0);
        siz.assign(n, 1);
    }

    int find(int v)
    {
        if (v == par[v])
            return v;
        return par[v] = find(par[v]);
    }

    bool unite(int a, int b)
    {
        a = find(a);
        b = find(b);
        if (a != b)
        {
            if (siz[a] < siz[b])
                swap(a, b);
            par[b] = a;
            siz[a] += siz[b];
            return true;
        }
        return false;
    }
};
```

fenwick.hpp

65d3d2, 44 lines

```
// methods work in 0-indexing
struct Fenwick
{
    int n;
    VL t;

    Fenwick (int _n = 0): n(_n + 1), t(n) {}
    Fenwick (const VL& ar)
    {
        n = sz(ar) + 1;
        t.assign(n, 0);
        FOR (i, 1, n)
        {
            t[i] += ar[i - 1];
            int x = i + (i & -i);
            if (x < n)
                t[x] += t[i];
        }
    }

    void upd(int x, ll val)
    {
        x++;
        while (x < n)
        {
            t[x] += val;
            x += x & -x;
        }
    }

    ll getSum(int x)
```

dsu fenwick fenwick-lower-bound segtree pers-segtree

```
{
    ll res = 0;
    while (x)
    {
        res += t[x];
        x -= x & -x;
    }
    return res;
}

// [l, r]
ll query(int l, int r)
{
    return getSum(r + 1) - getSum(l);
}
```

fenwick-lower-bound.hpp

edeaca, 15 lines

```
// returns first index p such that sum on [0, p] >= val or n if
// not found
int lower_bound(ll val)
{
    ll sm = 0;
    int pos = 0;
    for (int i = 1 << (31 - __builtin_clz(n)); i; i >>= 1)
    {
        if (pos + i < n && sm + t[pos + i] < val)
        {
            sm += t[pos + i];
            pos += i;
        }
    }
    return pos;
}
```

Minimum on a Segment

Maintain two Fenwick trees with $n = 2^k$ — one for the original array and the other for the reversed array. Use: $\mathbf{n} = \text{__bit_ceil}(n)$.

Add on a Segment

Maintain two Fenwick trees: \mathbf{tMult} and \mathbf{tAdd} .

To add x on the segment $[l, r]$, perform:

```
tMult.upd(l, x),
tMult.upd(r + 1, -x),
tAdd.upd(l, -x * l),
tAdd.upd(r + 1, x * (r + 1)).
```

Then, the sum on $[l, r]$ is:

$$\begin{aligned} \text{sum}(l, r) &= (r + 1) \mathbf{tMult.getSum}(r + 1) + \mathbf{tAdd.getSum}(r + 1) \\ &\quad - (l \cdot \mathbf{tMult.getSum}(l) + \mathbf{tAdd.getSum}(l)). \end{aligned}$$

segtree.hpp

a37959, 56 lines

```
struct Node
{
};
```

Node op(const Node& l, const Node& r)

{

}

struct SegTree

{ int n;
vector<Node> t;

SegTree(unsigned int _n)

{ n = bit_ceil(_n);
t.assign(2 * n, Node()); }

SegTree(const vector<Node>& ar)

{ n = bit_ceil(ar.size());
t.assign(2 * n, Node());
FOR (i, 0, sz(ar))
 t[i + n] = ar[i];
RFOR (i, n, 1)
 t[i] = op(t[i << 1], t[(i << 1) | 1]); }

void upd(int p, Node val)

{ p += n;
t[p] = val;
while (p >>= 1)
{
 t[p] = op(t[p << 1], t[(p << 1) | 1]);
}

// [l, r]
Node query(int l, int r)

{ l += n;
r += n;

Node resL, resR;
while (l < r)
{
 if (l & 1) resL = op(resL, t[l++]);
 if (r & 1) resR = op(t[--r], resR);
 l >>= 1;
 r >>= 1;
}
return op(resL, resR);
};

min= and sum with Segment Tree

Store in each node: max, cntMax, max2, sum.

In update check l, r conditions and:

- if ($\text{val} \geq \text{max}$) return;
- else if ($\text{val} > \text{max2}$) update this node;
- else go to left and right

You can do $\text{max}=$ and $+=$ on segment at the same time. Time: $O(\log n)$. Each extra descent decreases number of different elements in segment.

pers-segtree.hpp

2af912, 61 lines

```

struct Node
{
    int sm;
    Node* l;
    Node* r;

    Node(int v = 0): sm(v), l(nullptr), r(nullptr) {}
    Node(Node* _l, Node* _r): sm(_l->sm + _r->sm), l(_l), r(_r) {}

};

struct PersSegTree
{
    int n;
    vector<Node*> t;

    PersSegTree(int _n)
    {
        n = _n;
        t.pb(build(0, n - 1));
    }

    Node* build(int l, int r)
    {
        if (l == r)
            return new Node();
        int m = (l + r) >> 1;
        return new Node(build(l, m), build(m + 1, r));
    }

    Node* upd(Node* v, int l, int r, int pos, int val)
    {
        if (l == r)
            return new Node(val);
        int m = (l + r) >> 1;
        if (pos <= m)
            return new Node(upd(v->l, l, m, pos, val), v->r);
        else
            return new Node(v->l, upd(v->r, m + 1, r, pos, val));
    }

    Node query(Node* v, int l, int r, int ql, int qr)
    {
        if (ql > r || qr < l)
            return Node();
        if (ql <= l && r <= qr)
            return *v;
        int m = (l + r) >> 1;
        Node L = query(v->l, l, m, ql, qr);
        Node R = query(v->r, m + 1, r, ql, qr);
        return Node(&L, &R);
    }

    int upd(int ver, int pos, int val)
    {
        Node* nd = upd(t[ver], 0, n - 1, pos, val);
        t.pb(nd);
        return sz(t) - 1;
    }

    // [l, r]
    Node query(int ver, int l, int r)
    {
        return query(t[ver], 0, n - 1, l, r);
    }
};

```

lazysegtree.hpp

Description: Supports everything related to seg trees

99ad4c, 106 lines

```

template<class S, S (*op)(S, S), S (*e)(),
        class F, S (*mapping)(F, S),

```

lazysegtree segtree-usage

```

        F (*composition)(F, F), F (*id)()>
struct LazySegTree
{
    int n, size, log;
    vector<S> t;
    vector<F> lz;

    LazySegTree(int _n = 0) : LazySegTree(vector<S>(_n, e())) {}
    LazySegTree(const vector<S>& v)
    {
        n = sz(v);
        size = 1;
        while (size < n)
            size <= 1;
        log = __builtin_ctz(size);
        t.assign(2 * size, e());
        lz.assign(size, id());
        FOR(i, 0, n) t[size + i] = v[i];
        RFOR(i, size, 1) update(i);
    }

    void update(int k) { t[k] = op(t[k <> 1], t[k <> 1 | 1]); }

    void all_apply(int k, F f)
    {
        t[k] = mapping(f, t[k]);
        if (k < size) lz[k] = composition(f, lz[k]);
    }

    void push(int k)
    {
        all_apply(k << 1, lz[k]);
        all_apply(k << 1 | 1, lz[k]);
        lz[k] = id();
    }

    void set(int p, S x)
    {
        p += size;
        RFOR(i, log + 1, 1) push(p >> i);
        t[p] = x;
        FOR(i, 1, log + 1) update(p >> i);
    }

    S get(int p)
    {
        p += size;
        RFOR(i, log + 1, 1) push(p >> i);
        return t[p];
    }

    // [l, r]
    S prod(int l, int r)
    {
        if (l == r) return e();
        l += size; r += size;
        RFOR(i, log + 1, 1)
        {
            if (((l >> i) << i) != l) push(l >> i);
            if (((r >> i) << i) != r) push((r - 1) >> i);
        }
        S sml = e(), smr = e();
        while (l < r)
        {
            if (l & 1) sml = op(sml, t[l++]);
            if (r & 1) smr = op(t[--r], smr);
            l >>= 1; r >>= 1;
        }
    }

```

```

        return op(sml, smr);
    }

    S all_prod() { return t[1]; }

    void apply(int p, F f)
    {
        p += size;
        RFOR(i, log + 1, 1) push(p >> i);
        t[p] = mapping(f, t[p]);
        FOR(i, 1, log + 1) update(p >> i);
    }

    void apply(int l, int r, F f)
    {
        if (l == r) return;
        l += size; r += size;
        RFOR(i, log + 1, 1)
        {
            if (((l >> i) << i) != l) push(l >> i);
            if (((r >> i) << i) != r) push((r - 1) >> i);
        }
        int l2 = l, r2 = r;
        while (l < r)
        {
            if (l & 1) all_apply(l++, f);
            if (r & 1) all_apply(--r, f);
            l >>= 1; r >>= 1;
        }
        FOR(i, 1, log + 1)
        {
            if (((l2 >> i) << i) != l2) update(l2 >> i);
            if (((r2 >> i) << i) != r2) update((r2 - 1) >> i);
        }
    }
};

```

segtree-usage.hpp

add5af, 32 lines

```

// Example of (Sum + Range Add) with Lazy Segment Tree
struct S
{
    long long sum;
    int len;
};

using F = long long;

S op(S a, S b)
{
    return {a.sum + b.sum, a.len + b.len};
}

S e()
{
    return {0, 0};
}

S mapping(F f, S x)
{
    return {x.sum + f * x.len, x.len};
}

F composition(F f, F g)
{
    return f + g;
}

F id()
{
    return 0;
}

```

```

vector<S> v(n, {0, 1}); // each segment length = 1 initially
LazySegTree<S, op, e, F, mapping, composition, id> seg(v);

segtree-minleft-maxright.hpp e21244, 67 lines
// 34a7ec for min_left

// If f is monotone, this is the maximum r that satisfies
// f(op(a[l], a[l + 1], ..., a[r - 1])) = true
template<class G>
int max_right(int l, G g)
{
    if (l == n) return n;
    assert(g(e()));
    l += size;
    RFOR(i, log + 1, 1) push(l >> i);
    S sm = e();
    do
    {
        while ((l & 1) == 0) l >>= 1;
        if (!g(op(sm, t[l])))
        {
            while (l < size)
            {
                push(l);
                l = (l << 1);
                if (g(op(sm, t[l])))
                {
                    sm = op(sm, t[l]);
                    l++;
                }
            }
            return l - size;
        }
        sm = op(sm, t[l]);
        l++;
    } while ((l & -l) != l);
    return n;
}

// If f is monotone, this is the minimum l that satisfies
// f(op(a[l], a[l + 1], ..., a[r - 1])) = true
//template<class G>
//int min_left(int r, G g)
//{
//    if (r == 0) return 0;
//    assert(g(e()));
//    r += size;
//    RFOR(i, log + 1, 1) push((r - 1) >> i);
//    S sm = e();
//    do
//    {
//        //r--;
//        //while (r > 1 && (r & 1)) r >>= 1;
//        //if (!g(op(t[r], sm)))
//        //{
//            //push(r);
//            //r = (r << 1) / 1;
//            //if (g(op(t[r], sm)))
//            //{
//                //sm = op(t[r], sm);
//                //r--;
//            //}
//        //}
//        //return r + 1 - size;
//    } while (sm = op(t[r], sm));
}

```

segtree-minleft-maxright sparse-table LCA treap

```

    //} while ((r & -r) != r);
    //return 0;
//}

sparse-table.hpp e666cf, 19 lines
Description: Sparse table for minimum on the range [l, r), l < r. You can push back an element in O(LOG) and query anytime.

struct SparseTable
{
    VI t[LOG];
    void push_back(int v)
    {
        int i = sz(t[0]);
        t[0].pb(v);
        FOR(j, 0, LOG - 1)
            t[j + 1].pb(min(t[j][i], t[j][max(0, i - (1 << j))]));
    }
    // [l, r)
    int query(int l, int r)
    {
        assert(l < r && r <= sz(t[0]));
        int i = 31 - __builtin_clz(r - l);
        return min(t[i][r - 1], t[i][l + (1 << i) - 1]);
    }
};

LCA.hpp c9f639, 43 lines
struct LCA
{
    int n;
    VI I; // v -> po(v)
    VI RI;
    VI M; // to index mapping
    VI D;
    SparseTable st;
    LCA(const vector<vector<int>>& adj, int root)
    {
        n = sz(adj);
        I = VI(n);
        RI = VI(n);
        D = VI(n, -1);
        M = VI(2 * n, -1);
        int ctr = 0;
        vector<int> a;
        function<void(int, int, int)> preorder = [&](int v, int pr, int d)
        {
            I[v] = ctr++;
            RI[I[v]] = v;
            a.pb(I[v]);
            D[v] = d;
            for(auto to: adj[v])
            {
                if(to != pr)
                {
                    preorder(to, v, d + 1);
                    a.pb(I[v]);
                }
            }
        };
        preorder(root, -1, 0);
        FOR(i, 0, sz(a)) st.pb(a[i]);
        FOR(i, 0, sz(a)) M[a[i]] = i;
    }
    int lca(int u, int v)

```

```

    {
        return RI[st.query(min(M[I[u]], M[I[v]]), max(M[I[u]], M[I[v]]) + 1)];
    }
};

treap.hpp 215374, 144 lines
Description: uncomment in split for explicit key or in merge for implicit priority. Minimum and reverse queries.

mt19937 rng;

struct Node
{
    int l, r;
    int x, y;
    int cnt, par;
    int rev, mn;
    Node(int value)
    {
        l = r = -1;
        x = value;
        y = rng();
        cnt = 1;
        par = -1;
        rev = 0;
        mn = value;
    }
};

struct Treap
{
    vector<Node> t;
    int getCnt(int v)
    {
        if (v == -1)
            return 0;
        return t[v].cnt;
    }
    int getMn(int v)
    {
        if (v == -1)
            return INF;
        return t[v].mn;
    }
    int newNode(int val)
    {
        t.pb({val});
        return sz(t) - 1;
    }
    void upd(int v)
    {
        if (v == -1)
            return;
        // important!
        t[v].cnt = getCnt(t[v].l) + getCnt(t[v].r) + 1;
        t[v].mn = min(t[v].x, min(getMn(t[v].l), getMn(t[v].r)));
    }
    void reverse(int v)
    {
        if (v == -1)
            return;
        t[v].rev ^= 1;
    }
    void push(int v)
    {

```

```

if (v == -1 || t[v].rev == 0)
    return;
reverse(t[v].l);
reverse(t[v].r);
swap(t[v].l, t[v].r);
t[v].rev = 0;
}
pii split(int v, int cnt)
{
    if (v == -1)
        return {-1, -1};
    push(v);
    int left = getCnt(t[v].l);
    pii res;
    // elements t[v].x == val will be in right part
    // if (val <= t[v].x)
    if (cnt <= left)
    {
        if (t[v].l != -1)
            t[t[v].l].par = -1;
        // res = split(t[v].l, val);
        res = split(t[v].l, cnt);
        t[v].l = res.y;
        if (res.y != -1)
            t[res.y].par = v;
        res.y = v;
    }
    else
    {
        if (t[v].r != -1)
            t[t[v].r].par = -1;
        // res = split(t[v].r, val);
        res = split(t[v].r, cnt - left - 1);
        t[v].r = res.x;
        if (res.x != -1)
            t[res.x].par = v;
        res.x = v;
    }
    upd(v);
    return res;
}
int merge(int v, int u)
{
    if (v == -1) return u;
    if (u == -1) return v;
    int res;
    // if ((int)(rng() % (getCnt(v) + getCnt(u))) < getCnt(v))
    if (t[v].y > t[u].y)
    {
        push(v);
        if (t[v].r != -1)
            t[t[v].r].par = -1;
        res = merge(t[v].r, u);
        t[v].r = res;
        if (res != -1)
            t[res].par = v;
        res = v;
    }
    else
    {
        push(u);
        if (t[u].l != -1)
            t[t[u].l].par = -1;
        res = merge(v, t[u].l);
        t[u].l = res;
        if (res != -1)
            t[res].par = u;
        res = u;
    }
}

```

```

    upd(res);
    return res;
}
// returns index of element [0, n)
int getIdx(int v, int from = -1)
{
    if (v == -1)
        return 0;
    int x = getIdx(t[v].par, v);
    push(v);
    if (from == -1 || t[v].r == from)
        x += getCnt(t[v].l) + (from != -1);
    return x;
}
};

lct.hpp
Description: Link-Cut Tree. Calculate any path queries. Change upd to maintain what you need. Don't use upd in push(). Calculate non commutative functions in both ways and swap them in push. cnt – number of nodes in current splay tree. Don't touch rev, sub, vsub. v->access() brings v to the top and pushes it; its left subtree will be the path from v to the root and its right subtree will be empty. Only then sub will be the number of nodes in the connected component of v and vsub will be the number of nodes under v. Change upd to calc sum in subtree of other functions. Use makeRoot for arbitrary path queries.
Usage: FOR (i, 0, n) LCT[i] = new snode(i); link(LCT[u], LCT[v]);
Time:  $\mathcal{O}(\log n)$  788027, 159 lines

typedef struct Snode* sn;
struct Snode
{
    sn p, c[2]; // parent, children
    bool rev = false; // subtree reversed or not (internal usage)
    int val, cnt; // value in node, # nodes in splay subtree
    int sub, vsub = 0; // vsub stores sum of virtual children

    Snode(int _val): val(_val)
    {
        p = c[0] = c[1] = 0;
        upd();
    }
    friend int getCnt(sn v)
    {
        return v ? v->cnt : 0;
    }
    friend int getSub(sn v)
    {
        return v ? v->sub : 0;
    }
    void push()
    {
        if (!rev)
            return;
        swap(c[0], c[1]);
        rev = false;
        FOR (i, 0, 2)
            if (c[i])
                c[i]->rev ^= 1;
    }
    void upd()
    {
        FOR (i, 0, 2)
            if (c[i])
                c[i]->push();
        cnt = 1 + getCnt(c[0]) + getCnt(c[1]);
        sub = 1 + getSub(c[0]) + getSub(c[1]) + vsub;
    }
    int dir()
    {
        if (!p) return -2;
        FOR (i, 0, 2)

```

```

            if (p->c[i] == this)
                return i;
            // p is path-parent pointer
            // --> not in current splay tree
            return -1;
        }

        // checks if root of current splay tree
        bool isRoot()
        {
            return dir() < 0;
        }
        friend void setLink(sn p, sn v, int d)
        {
            if (v)
                v->p = p;
            if (d >= 0)
                p->c[d] = v;
        }
        void rot()
        {
            assert(!isRoot());
            int d = dir();
            sn pa = p;
            setLink(pa->p, this, pa->dir());
            setLink(pa, c[d ^ 1], d);
            setLink(this, pa, d ^ 1);
            pa->upd();
        }
        void splay()
        {
            while (!isRoot() && !p->isRoot())
            {
                p->p->push();
                p->push();
                push();
                dir() == p->dir() ? p->rot() : rot();
            }
            if (!isRoot())
                p->push(), push(), rot();
            push();
            upd();
        }
        // bring this to top of tree, propagate
        void access()
        {
            for (sn v = this, pre = 0; v; v = v->p)
            {
                v->splay();
                if (pre)
                    v->vsub -= pre->sub;
                if (v->c[1])
                    v->vsub += v->c[1]->sub;
                v->c[1] = pre;
                v->upd();
                pre = v;
            }
            splay();
            assert(!c[1]);
        }
        void makeRoot()
        {
            access();
            rev ^= 1;
            access();
            assert(!c[0] && !c[1]);
        }
        friend sn lca(sn u, sn v)

```

```

{
    if (u == v)
        return u;
    u->access();
    v->access();
    if (!u->p)
        return 0;
    u->splay();
    return u->p ? u->p : u;
}

friend bool connected(sn u, sn v)
{
    return lca(u, v);
}

void set(int v)
{
    access();
    val = v;
    upd();
}

friend void link(sn u, sn v)
{
    assert(!connected(u, v));
    v->makeRoot();
    u->access();
    setLink(v, u, 0);
    v->upd();
}

// cut v from it's parent in LCT
// make sure about root or better use next function
friend void cut(sn v)
{
    v->access();
    assert(v->c[0]); // assert if not a root
    v->c[0]->p = 0;
    v->c[0] = 0;
    v->upd();
}

// u, v should be adjacent in tree
friend void cut(sn u, sn v)
{
    u->makeRoot();
    v->access();
    assert(v->c[0] == u && !u->c[0] && !u->c[1]);
    cut(v);
}

```

ordered-set.hpp

16 lines

```

#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
using namespace std;
typedef tree<int, null_type, less<int>, rb_tree_tag,
            tree_order_statistics_node_update> ordered_set;

ordered_set s;
s.insert(47);
// Returns the number of elements less than k
s.order_of_key(k);
// Returns iterator to the k-th element or s.end()
s.find_by_order(k);
// Does not exist
s.count();
// Doesn't trigger RE. Returns 0 if compiled using F8
*s.end();

```

ordered-set convex-hull-trick bellman-ford-moore monge-shortest-path

convex-hull-trick.hpp
Description: add(a, b) adds a straight line $y = ax + b$. getMaxY(p) finds the maximum y at $x = p$.

```

struct Line
{
    ll a, b, xLast;
    Line() {}
    Line(ll _a, ll _b): a(_a), b(_b) {}
    bool operator<(const Line& l) const
    {
        return MP(a, b) < MP(l.a, l.b);
    }
    bool operator<(int x) const
    {
        return xLast < x;
    }
    __int128 getY(__int128 x) const
    {
        return a * x + b;
    }
    ll intersect(const Line& l) const
    {
        assert(a < l.a);
        ll dA = l.a - a, dB = b - l.b, x = dB / dA;
        if (dB < 0 && dB % dA != 0)
            x--;
        return x;
    }
};

struct ConvexHull: set<Line, less<>>
{
    bool needErase(iterator it, const Line& l)
    {
        ll x = it->xLast;
        if (it->getY(x) > l.getY(x))
            return false;
        if (it == begin())
            return it->a >= l.a;
        x = prev(it)->xLast + 1;
        return it->getY(x) < l.getY(x);
    }
    void add(ll a, ll b)
    {
        Line l(a, b);
        auto it = lower_bound(l);
        if (it != end())
        {
            ll x = it == begin() ? -LINF :
                prev(it)->xLast;
            if ((it == begin())
                || prev(it)->getY(x) >= l.getY(x))
                && it->getY(x + 1) >= l.getY(x + 1))
                return;
        }
        while (it != end() && needErase(it, l))
            it = erase(it);
        while (it != begin() && needErase(prev(it), l))
            erase(prev(it));
        if (it != begin())
        {
            auto itP = prev(it);
            Line itL = *itP;
            itL.xLast = itP->intersect(l);
            erase(itP);
            insert(itL);
        }
        l.xLast = it == end() ? LINF : l.intersect(*it);
        insert(l);
    }
}
```

```

} ll getMaxY(ll p)
{
    return lower_bound(p)->getY(p);
}

```

Graphs (3)

Shortest paths

bellman-ford-moore.hpp

Description: Computes shortest paths from a single source vertex to all of the other vertices in a weighted directed graph.

Time: $\mathcal{O}(nm)$

eb281b, 35 lines

```

VL spfa(const vector<vector<pair<int, ll>>> &g, int n, int s)
{
    VL dist(n, LINF);
    dist[s] = 0;
    queue<int> q;
    q.push(s);
    VI inQueue(n);
    inQueue[s] = true;
    VI cnt(n);
    bool negCycle = false;
    while (!q.empty())
    {
        int v = q.front();
        q.pop();
        cnt[v]++;
        negCycle |= cnt[v] > n;
        inQueue[v] = false;
        for (auto [to, w] : g[v])
        {
            ll newDist = dist[v] + w;
            if (newDist < dist[to])
            {
                dist[to] = newDist;
                if (!inQueue[to])
                {
                    q.push(to);
                    inQueue[to] = true;
                }
            }
        }
        if (negCycle)
            break;
    }
    return dist;
}

```

monge-shortest-path.hpp

Description: Finds shortest paths from the vertex 0 to all vertices in a DAG with n vertices, where the edges weights $c(i, j)$ satisfy the Monge property: $\forall i, j, k, l, \quad 0 \leq i < j < k < l < n \implies c(i, l) + c(j, k) \geq c(i, k) + c(j, l)$.

Time: $\mathcal{O}(n \log n)$

540e92, 34 lines

```

template<typename F>
VL mongeShortestPath(int n, const F& cost)
{
    VL dist(n, LINF);
    VI amin(n);
    dist[0] = 0;

    auto update = [&](int i, int k)
    {
        ll nd = dist[k] + cost(k, i);
        if (nd < dist[i])
        {

```

```

dist[i] = nd;
amin[i] = k;
};

function<void(int, int)> solve = [&](int l, int r)
{
    if (r - l == 1)
        return;
    int m = (l + r) / 2;
    FOR(k, amin[l], min(m, amin[r] + 1))
        update(m, k);
    solve(l, m);
    FOR(k, l + 1, m + 1)
        update(r, k);
    solve(m, r);
};

update(n - 1, 0);
solve(0, n - 1);
return dist;
}

```

Decompositions

centroid.hpp

8c3d24, 51 lines

```

VI g[N];
int siz[N];
bool usedC[N];

int dfsSZ(int v, int par)
{
    siz[v] = 1;
    for (auto to : g[v])
    {
        if (to != par && !usedC[to])
            siz[v] += dfsSZ(to, v);
    }
    return siz[v];
}

void build(int u)
{
    dfsSZ(u, -1);
    int szAll = siz[u];
    int pr = u;
    while (true)
    {
        int v = -1;
        for (auto to : g[u])
        {
            if (to == pr || usedC[to])
                continue;
            if (siz[to] * 2 > szAll)
            {
                v = to;
                break;
            }
        }
        if (v == -1)
            break;
        pr = u;
        u = v;
    }
    int cent = u;
    usedC[cent] = true;
}

// here calculate f(cent)

```

centroid hld biconnected-components

```

for (auto to : g[cent])
{
    if (!usedC[to])
    {
        build(to);
    }
}

hld.hpp
Description: Run dfsSZ(root, -1, 0) and dfsHLD(root, -1, root) to build the HLD. Each vertex  $v$  has an index  $\text{tin}[v]$ . To update on the path, use the process as defined in get(). The values are stored in the vertices. dc0437, 67 lines
VI g[N];
int siz[N];
int h[N];
int p[N];
int top[N];
int tin[N];
int tout[N];
int t = 0;

void dfsSZ(int v, int par, int hei)
{
    siz[v] = 1;
    h[v] = hei;
    p[v] = par;
    for (auto& to : g[v])
    {
        if (to == par)
            continue;
        dfsSZ(to, v, hei + 1);
        siz[v] += siz[to];
        if (g[v][0] == par || siz[g[v][0]] < siz[to])
            swap(g[v][0], to);
    }
}

void dfsHLD(int v, int par, int tp)
{
    tin[v] = t++;
    top[v] = tp;
    FOR(i, 0, sz(g[v]))
    {
        int to = g[v][i];
        if (to == par)
            continue;
        if (i == 0)
            dfsHLD(to, v, tp);
        else
            dfsHLD(to, v, to);
    }
    tout[v] = t - 1;
}

ll get(int u, int v)
{
    ll res = 0;
    while(true)
    {
        int tu = top[u];
        int tv = top[v];
        if (tu == tv)
        {
            int t1 = tin[u];
            int t2 = tin[v];
            if (t1 > t2)
                swap(t1, t2);
            // query [t1, t2] both inclusive
            // res += query(t1, t2);
            break;
        }
    }
}

```

```

    }

    if (h[tu] < h[tv])
    {
        swap(tu, tv);
        swap(u, v);
    }
    //res += query(tin[tu], tin[u]);
    u = p[tu];
}

return res;
}

```

biconnected-components.hpp

Description: Colors the edges so that the vertices, connected with the same color are still connected if you delete any vertex.

Time: $\mathcal{O}(m)$

7d48ce, 117 lines

```

struct Graph
{
    int n, m;
    vector<pii> edges;
    vector<VI> g;

    VI used, par;
    VI tin, low, inComp;
    int t = 0, c = 0;
    VI st;

    // components of vertices
    // a vertex can be in several components
    vector<VI> verticesCol;
    // components of edges
    vector<VI> components;
    // col[i] – component of the i-th edge
    VI col;

    Graph(int _n = 0, int _m = 0): n(_n), m(_m), edges(m), g(n),
    used(n), par(n, -1), tin(n), low(n), inComp(n), col(m, -1) {}

    void addEdge(int a, int b, int i)
    {
        assert(0 <= a && a < n);
        assert(0 <= b && b < n);
        assert(0 <= i && i < m);

        edges[i] = MP(a, b);
        g[a].pb(i);
        g[b].pb(i);
    }

    void addComp()
    {
        unordered_set<int> s;
        s.reserve(7 * sz(components[c]));
        for (auto e : components[c])
        {
            s.insert(edges[e].x);
            s.insert(edges[e].y);
            inComp[edges[e].x] = true;
            inComp[edges[e].y] = true;
        }
        verticesCol.pb(VI(all(s)));
    }

    void dfs(int v, int p = -1)
    {
        used[v] = 1;
        par[v] = p;
        low[v] = tin[v] = t++;
        int cnt = 0;

```

```

for (auto e : g[v])
{
    int to = edges[e].x;
    if (to == v)
        to = edges[e].y;

    if (p == to) continue;
    if (!used[to])
    {
        cnt++;
        st.pb(e);
        dfs(to, v);

        low[v] = min(low[v], low[to]);
    }

    if ((par[v] == -1 && cnt > 1) || (par[v] != -1 && low[to] >= tin[v]))
    {
        components.pb({});
        while (st.back() != e)
        {
            components[c].pb(st.back());
            col[st.back()] = c;

            st.pop_back();
        }
        components[c].pb(st.back());
        addComp();
        col[st.back()] = c++;

        st.pop_back();
    }
}
else
{
    low[v] = min(low[v], tin[to]);
    if (tin[to] < tin[v])
        st.pb(e);
}
}

void build()
{
    FOR (i, 0, n)
    {
        if (used[i]) continue;
        dfs(i, -1);
        if (st.empty()) continue;
        components.pb({});
        while (!st.empty())
        {
            int e = st.back();
            col[e] = c;
            components[c].pb(e);
            st.pop_back();
        }
        addComp();
        c++;
    }
    FOR (i, 0, n)
        if (!inComp[i])
            verticesCol.pb(VI(1, i));
}

```

scc.hpp

118635, 58 lines

vector<bool> vis;

scc hierholzer

```

void dfs(int v, vector<VI> const& adj, VI &output)
{
    vis[v] = true;
    for (auto u : adj[v])
        if (!vis[u])
            dfs(u, adj, output);
    output.pb(v);
}

/*
 * Kosaraju SCC + condensation graph
 * comps - strongly connected components in topological order
 * adj_cond - condensation DAG, each scc represented by vertex
 *             with the smallest index inside it
 */
void scc(vector<VI> const& adj, vector<VI> &comps, vector<VI> &adj_cond)
{
    int n = sz(adj);
    comps.clear(), adj_cond.clear();

    VI ord;
    vis.assign(n, false);

    FOR (i, 0, n)
        if (!vis[i])
            dfs(i, adj, ord);

    vector<VI> adj_rev(n);
    FOR (v, 0, n)
        for (int u : adj[v])
            adj_rev[u].pb(v);

    vis.assign(n, false);
    reverse(all(ord));

    VI roots(n, 0);

    for (auto v : ord)
    {
        if (!vis[v])
        {
            VI comp;
            dfs(v, adj_rev, comp);
            comps.pb(comp);
            int root = *min_element(all(comp));
            for (auto u : comp)
                roots[u] = root;
        }
    }

    adj_cond.assign(n, {});
    FOR (v, 0, n)
        for (auto u : adj[v])
            if (roots[v] != roots[u])
                adj_cond[roots[v]].pb(roots[u]);
}

```

Hierholzer's algorithm

hierholzer.hpp

Description: Finds an Eulerian path in a directed or undirected graph. g is a graph with n vertices. $g[u]$ is a vector of pairs $(v, \text{edge_id})$. m is the number of edges in the graph. The vertices are numbered from 0 to $n-1$, and the edges - from 0 to $m-1$. If there is no Eulerian path, returns $\{\{-1\}, \{-1\}\}$. Otherwise, returns the path in the form (vertices, edges) with vertices containing $m+1$ elements and edges containing m elements. If you need an Eulerian cycle, check $\text{vertices}[0] = \text{vertices.back()}$.

fa6dc3, 101 lines

// f14a40 for undirected

```

tuple<bool, int, int> checkDirected(vector<vector<pii>> &g)
{
    int n = sz(g), v1 = -1, v2 = -1;
    bool bad = false;
    VI degIn(n);
    FOR(u, 0, n)
        for (auto [v, e] : g[u])
            degIn[v]++;
    FOR(u, 0, n)
    {
        bad |= abs(degIn[u] - sz(g[u])) > 1;
        if (degIn[u] < sz(g[u]))
        {
            bad |= v2 != -1;
            v2 = u;
        }
        else if (degIn[u] > sz(g[u]))
        {
            bad |= v1 != -1;
            v1 = u;
        }
    }
    return {bad, v1, v2};
}

/*tuple<bool, int, int> checkUndirected(vector<vector<pii>> &g)
{
    int n = sz(g), v1 = -1, v2 = -1;
    bool bad = false;
    FOR(u, 0, n)
    {
        if (sz(g[u]) & 1)
        {
            bad |= v2 != -1;
            if (v1 == -1)
                v1 = u;
            else
                v2 = u;
        }
    }
    return {bad, v1, v2};
}*/
```

pair<VI, VI> hierholzer(**vector**<**vector**<pii>> g, **int** m)
{
 // checkUndirected if undirected
 auto [bad, v1, v2] = checkDirected(g);
 if (bad)
 return {{-1}, {-1}};
 if (v1 != -1)
 {
 g[v1].pb({v2, m});
 // uncomment if undirected
 // g[v2].PB({v1, m});
 m++;
 }
 deque<pii> d;
 VI used(m);
 int v = 0, k = 0;
 while (m > 0 && g[v].empty())
 v++;
 while (sz(d) < m)
 {
 while (k < m)
 {
 while (!g[v].empty() && used[g[v].back().y])
 g[v].pop_back();
 if (!g[v].empty())
 break;
 }
 }
}

```

d.push_front(d.back());
d.pop_back();
v = d.back().x;
k++;
}
if (k == m)
    return {{-1}, {-1}};
d.pb(g[v].back());
used[g[v].back().y] = true;
g[v].pop_back();
v = d.back().x;
}
while (v1 != -1 && d.back().y != m - 1)
{
    d.push_front(d.back());
    d.pop_back();
    v = d.back().x;
}
VI vertices = {v}, edges;
for (auto [u, e] : d)
{
    vertices.pb(u);
    edges.pb(e);
}
if (v1 != -1)
{
    vertices.pop_back();
    edges.pop_back();
}
return {vertices, edges};
}

```

Maximum matching

kuhn.hpp

Description: mateFor is -1 or mate. addEdge([0, L), [0, R)).

Time: 0.6s for $L, R \leq 10^5$, $|E| \leq 2 \cdot 10^5$

930365, 76 lines

mt19937 rng;

```

struct Graph
{
    int szL, szR;
    // edges from the left to the right, 0-indexed
    vector<VI> g;
    VI mateForL, usedL, mateForR;

    Graph(int L = 0, int R = 0): szL(L), szR(R), g(L),
        mateForL(L), usedL(L), mateForR(R) {}

    void addEdge(int from, int to)
    {
        assert(0 <= from && from < szL);
        assert(0 <= to && to < szR);

        g[from].pb(to);
    }

    int iter;
    bool kuhn(int v)
    {
        if (usedL[v] == iter) return false;
        usedL[v] = iter;
        shuffle(all(g[v]), rng);
        for (int to : g[v])
        {
            if (mateForR[to] == -1)
            {
                mateForR[to] = v;
                mateForL[v] = to;
            }
        }
    }
}

```

```

        return true;
    }

    for (int to : g[v])
    {
        if (kuhn(mateForR[to]))
        {
            mateForR[to] = v;
            mateForL[v] = to;
            return true;
        }
    }
    return false;
}

int doKuhn()
{
    fill(all(mateForR), -1);
    fill(all(mateForL), -1);
    fill(all(usedL), -1);

    int res = 0;
    iter = 0;

    while (true)
    {
        iter++;

        bool ok = false;
        FOR(v, 0, szL)
        {
            if (mateForL[v] == -1)
            {
                if (kuhn(v))
                {
                    ok = true;
                    res++;
                }
            }
            if (!ok) break;
        }
        return res;
    }
}

```

edmonds-blossom.hpp

Description: Finds the maximum matching in a graph.

Time: $\mathcal{O}(n^2m)$

```

int t = mate[v];
mate[v] = w;
if (mate[t] != v)
    return;
if (label[v] <= n)
{
    mate[t] = label[v];
    augmentPath(label[v], t);
    return;
}
int x = label[v] / (n + 1);
int y = label[v] % (n + 1);
augmentPath(x, y);
augmentPath(y, x);
}

int findMaxMatching()
{
    FOR(i, 0, n + 1)
        assert(mate[i] == 0);
    int mt = 0;
    DSU dsu(n + 1);
    FOR(u, 1, n + 1)
    {
        if (mate[u] != 0)
            continue;
        fill(all(label), -1);
        iota(all(fir), 0);
        label[u] = 0;
        dsu.unite(u, 0);
        queue<int> q;
        q.push(u);
        while (!q.empty())
        {
            int x = q.front();
            q.pop();
            for (int y : g[x])
            {
                if (mate[y] == 0 && y != u)
                {
                    mate[y] = x;
                    augmentPath(x, y);
                    while (!q.empty())
                        q.pop();
                    mt++;
                    break;
                }
                if (label[y] < 0)
                {
                    int v = mate[y];
                    if (label[v] < 0)
                    {
                        label[v] = x;
                        dsu.unite(v, y);
                        q.push(v);
                    }
                }
                else
                {
                    int r = fir[dsu.find(x)], s = fir[dsu.find(y)];
                    if (r == s)
                        continue;
                    int edgeLabel = (n + 1) * x + y;
                    label[r] = label[s] = -edgeLabel;
                    int join;
                    while (true)
                    {
                        if (s != 0)
                            swap(r, s);
                        r = fir[dsu.find(label[mate[r]])];

```

ffc914, 124 lines

```

    if (label[r] == -edgeLabel)
    {
        join = r;
        break;
    }
    label[r] = -edgeLabel;
}
for (int z: {x, y})
{
    for (int v = fir[dsu.find(z)];
        v != join;
        v = fir[dsu.find(label[mate[v]])])
    {
        label[v] = edgeLabel;
        if (dsu.unite(v, join))
            fir[dsu.find(join)] = join;
        q.push(v);
    }
}
return mt;
}
int getMate(int v)
{
    assert(0 <= v && v < n);
    v++;
    int u = mate[v];
    assert(u == 0 || mate[u] == v);
    u--;
    return u;
}
};

```

Tutte matrix

Given an undirected graph $G = (V, E)$, its Tutte matrix is:

$$T_{ij} = \begin{cases} x_{ij} & \text{if } i < j \text{ and } (i, j) \in E \\ -x_{ji} & \text{if } i > j \text{ and } (i, j) \in E \\ 0 & \text{otherwise.} \end{cases}$$

$\det(T) \neq 0$ if and only if G has a perfect matching.

Flows

dinic.hpp

Description: Finds the maximum flow in a network.

Time: $\mathcal{O}(n^2m)$. If all capacities are less than c , then the complexity of the Dinic

is bounded by $\mathcal{O}\left(\min(n^{\frac{2}{3}}, \sqrt{cm}) \cdot cm\right)$.

bc6418, 87 lines

```

struct Graph
{
    struct Edge
    {
        int from, to;
        ll cap, flow;
    };

    int n;
    vector<Edge> edges;
    vector<VI> g;
    VI d, p;

    Graph(int _n): n(_n), g(n), d(n), p(n) {}
    void addEdge(int from, int to, ll cap)
    {
        edges.push_back({from, to, cap, 0});
        g[from].pb(to);
        g[to].pb(from);
    }
};

```

dinic successive-shortest-path

```

assert(0 <= from && from < n);
assert(0 <= to && to < n);
assert(0 <= cap);
g[from].pb(sz(edges));
edges.pb({from, to, cap, 0});
g[to].pb(sz(edges));
edges.pb({to, from, 0, 0});
}

int bfs(int s, int t)
{
    fill(all(d), -1);
    d[s] = 0;
    queue<int> q;
    q.push(s);
    while (!q.empty())
    {
        int v = q.front();
        q.pop();
        for (int e : g[v])
        {
            int to = edges[e].to;
            if (edges[e].flow < edges[e].cap && d[to] == -1)
            {
                d[to] = d[v] + 1;
                q.push(to);
            }
        }
    }
    return d[t];
}

ll dfs(int v, int t, ll flow)
{
    if (v == t || flow == 0)
        return flow;
    for (; p[v] < sz(g[v]); p[v]++)
    {
        int e = g[v][p[v]], to = edges[e].to;
        ll c = edges[e].cap, f = edges[e].flow;
        if (f < c && (to == t || d[to] == d[v] + 1))
        {
            ll push = dfs(to, t, min(flow, c - f));
            if (push > 0)
            {
                edges[e].flow += push;
                edges[e ^ 1].flow -= push;
                return push;
            }
        }
    }
    return 0;
}

ll flow(int s, int t)
{
    assert(0 <= s && s < n);
    assert(0 <= t && t < n);
    assert(s != t);
    ll flow = 0;
    while (bfs(s, t) != -1)
    {
        fill(all(p), 0);
        while (true)
        {
            ll f = dfs(s, t, LINF);
            if (f == 0)
                break;
            flow += f;
        }
    }
    return flow;
}

```

}

};

successive-shortest-path.hpp

Description: Finds the minimum cost maximum flow in a network. If the network contains negative-cost edges, uncomment initPotentials.
Time: $\mathcal{O}(|F| \cdot m \log n)$ without negative-cost edges, and $\mathcal{O}(|F| \cdot m \log n + nm)$ with negative-cost edges.

a220bb, 103 lines

```

struct Graph
{
    struct Edge
    {
        int from, to;
        int cap, flow;
        ll cost;
    };

    int n;
    vector<Edge> edges;
    vector<VI> g;
    VL pi, d;
    VI pred;

    Graph(int _n = 0): n(_n), g(n), pi(n), d(n), pred(n) {}

    void addEdge(int from, int to, int cap, ll cost)
    {
        assert(0 <= from && from < n);
        assert(0 <= to && to < n);
        assert(0 <= cap);
        g[from].pb(sz(edges));
        edges.pb({from, to, cap, 0, cost});
        g[to].pb(sz(edges));
        edges.pb({to, from, 0, 0, -cost});
    }

    /*void initPotentials(int s)
    {
        vector<vector<pair<int, ll>>> gr(n);
        FOR(v, 0, n)
        {
            for (int e : g[v])
            {
                const Edge& edge = edges[e];
                if (edge.flow < edge.cap)
                    gr[v].pb({edge.to, edge.cost});
            }
        }
        pi = spfa(gr, n, s);
    }*/

    pair<int, ll> flow(int s, int t)
    {
        assert(0 <= s && s < n);
        assert(0 <= t && t < n);
        assert(s != t);
        //initPotentials(s);
        int flow = 0;
        ll cost = 0;
        for (int it = 0; ; it++)
        {
            fill(all(d), LINF);
            fill(all(pred), -1);
            d[s] = 0;
            priority_queue<pair<ll, int>> q;
            q.push({0, s});
            while (!q.empty())
            {
                auto [dv, v] = q.top();
                q.pop();
                if (it > 0 && v == t)

```

```

break;
if (-dv != d[v])
    continue;
for (int i : g[v])
{
    if (edges[i].flow == edges[i].cap)
        continue;
    int to = edges[i].to;
    ll nd = d[v] + edges[i].cost + pi[v] - pi[to];
    if (nd < d[to])
    {
        d[to] = nd;
        pred[to] = i;
        q.push({-nd, to});
    }
}
if (d[t] == LINF)
    break;
int curFlow = INF;
for (int v = t; v != s;)
{
    int i = pred[v];
    curFlow = min(curFlow, edges[i].cap - edges[i].flow);
    v = edges[i].from;
}
for (int v = t; v != s;)
{
    int i = pred[v];
    edges[i].flow += curFlow;
    edges[i ^ 1].flow -= curFlow;
    v = edges[i].from;
}
flow += curFlow;
cost += (d[t] + pi[t] - pi[s]) * curFlow;
FOR(u, 0, n)
    if (it == 0 || d[u] <= d[t])
        pi[u] += d[u] - d[t];
}
return {flow, cost};
}

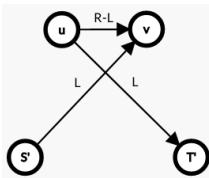
```

Maximum flow with minimum capacities

On the resulting graph, accumulate maximum flow in the following order:

- from S' to T'
- from S' to T
- from S to T'
- from S to T .

An $S - T$ flow that satisfies the minimum capacities exists if and only if, for all outgoing edges from S' and incoming edges to T' , the flow and capacity are equal.



3-cycles 4-cycles twosat

Quadratic supermodular pseudoboolean optimization

$$\sum_i a_i x_i + \sum_i b_i \bar{x}_i + \sum_{i,j} c_{ij} x_i \bar{x}_j \rightarrow \min$$

$$c_{ij} x_i x_j = c_{ij} x_i - c_{ij} x_i \bar{x}_j$$

If $a_i \leq b_i$, add an edge from S to i of capacity $b_i - a_i$ and add a_i to the answer.

Otherwise, add an edge from i to T of capacity $a_i - b_i$ and add b_i to the answer.

Add an edge from i to j of capacity c_{ij} .

Add the $S - T$ minimum cut to the answer.

Matching tricks

Minimum cut

To find the min-cut, search from vertex S on unsaturated edges. Original edges from used vertices to unused ones are in the min-cut.

Minimum vertex cover

The vertex cover problem is not NP-complete in bipartite graphs. The minimum number of vertices required to cover all edges is equal to the size of the maximum matching. To reconstruct the minimum vertex cover, create a directed graph:

- matched edges from the right part to the left part
- unmatched edges from the left part to the right part.

Start traversal from unmatched vertices in the left part. The cover includes vertices from the matching:

- unvisited vertices in the left part
- visited vertices in the right part.

Maximum independent set

The independent set problem is not NP-complete in bipartite graphs. It is the complement of the minimum vertex cover.

Minimum edge cover

A minimum edge cover can be found in any graph. The minimum number of edges required to cover all vertices can only be determined in graphs without isolated vertices. By utilizing one edge in the matching, we cover two vertices, while any other vertices are covered using one edge for each.

DAG paths

In a DAG, you can find the minimum number of non-intersecting paths that cover all vertices. Duplicate vertices and create a bipartite graph with edges $u_L \rightarrow v_R$. Edges in the matching correspond to edges in the paths.

Dominating set

A dominating set for a graph is a subset D of V such that any vertex is in D , or has a neighbor in D . The dominating set problem is NP-complete even on bipartite graphs. It can be found greedily on a tree.

Sqrt problems

3-cycles.hpp

Description: Finds all triangles in a graph. Each triangle (v, u, w) increments the cnt.

Time: $\mathcal{O}(m \cdot \sqrt{m})$

e5e996, 22 lines

int triangles(int n)

```

{
    vector<VI> ng(n);
    FOR (v, 0, n)
        for (auto u : adj[v])
            if (MP(sz(adj[v]), v) < MP(sz(adj[u]), u))
                ng[v].pb(u);
    int cnt = 0;
    VI used(n, 0);
    FOR (v, 0, n)
    {
        for (auto u : ng[v])
            used[u] = 1;
        for (auto u : ng[v])
            for (auto w : ng[u])
                if (used[w])
                    cnt++;
        for (auto u : ng[v])
            used[u] = 0;
    }
    return cnt;
}

```

4-cycles.hpp

Description: Sort d and add breaks to speed up. With breaks works 0.5s for $m = 5 \cdot 10^5$.

Time: $\mathcal{O}\left(\sum_{uv \in E} \min(\deg(u), \deg(v))\right) = \mathcal{O}(m \cdot \sqrt{m})$

73a48f, 20 lines

ll rect(int n)

```

ll cnt4 = 0;
vector<pii> d(n);
FOR (v, 0, n) d[v] = MP(sz(adj[v]), v);
VI L(n);
FOR (v, 0, n)
{
    for (auto u : adj[v])
        if (d[u] < d[v])
            for (auto y : adj[u])
                if (d[y] < d[v])
                    cnt4 += L[y], L[y]++;
    for (auto u : adj[v])
        if (d[u] < d[v])
            for (auto y : adj[u])
                L[y] = 0;
}
return cnt4;
}

```

Two sat

twosat.hpp

d5c678, 43 lines

```

// ! remove adj_cond from scc
struct TwoSat
{

```

```

    int n, N;
    vector<VI> g;
    vector<VI> comps;
    vector<bool> val;

```

```

    TwoSat(int _n) : n(_n), N(2 * n), g(N), val(n, false) {}

    // add clause (x_a OR x_b). na, nb - whether to negate
    void add(int a, bool na, int b, bool nb)

```

```

{
    int A = 2 * a ^ na;
    int B = 2 * b ^ nb;
    int NA = A ^ 1;
    int NB = B ^ 1;
    g[NA].pb(B);
    g[NB].pb(A);
}

// returns true if the formula is satisfiable. in this case,
// val contains a valid assignment
bool solve()
{
    comps.clear();
    scc(g, comps);

    VI id(N, -1);
    FOR (i, 0, sz(comps))
        for (int v : comps[i])
            id[v] = i;

    FOR (i, 0, n)
        if (id[2 * i] == id[2 * i + 1])
            return false;

    val.assign(n, false);
    FOR (i, 0, n)
        val[i] = (id[2 * i] > id[2 * i + 1]);

    return true;
}

```

Strings (4)

aho-corasick.hpp

e59836, 64 lines

```

const int AL = 26;

struct Node
{
    int p;
    int c;
    int g[AL];
    int nxt[AL];
    int link;
};

Node(int _c, int _p)
{
    c = _c;
    p = _p;
    fill(g, g + AL, -1);
    fill(nxt, nxt + AL, -1);
    link = -1;
};

struct AC
{
    vector<Node> a;
    AC(): a(1, {-1, -1}) {}

    int addStr(const string& s)
    {
        int v = 0;
        FOR (i, 0, sz(s))
        {
            // change to [0 AL)

```

aho-corasick suffix-automaton suffix-array

```

int c = s[i] - 'a';
if (a[v].nxt[c] == -1)
{
    a[v].nxt[c] = sz(a);
    a.pb(Node(c, v));
}
v = a[v].nxt[c];
}
return v;
}

int go(int v, int c)
{
    if (a[v].g[c] != -1)
        return a[v].g[c];

    if (a[v].nxt[c] != -1)
        a[v].g[c] = a[v].nxt[c];
    else if (v != 0)
        a[v].g[c] = go(getLink(v), c);
    else
        a[v].g[c] = 0;

    return a[v].g[c];
}

int getLink(int v)
{
    if (a[v].link != -1)
        return a[v].link;
    if (v == 0 || a[v].p == 0)
        return 0;
    return a[v].link = go(getLink(a[v].p), a[v].c);
}

```

suffix-automaton.hpp

183478, 57 lines

```

const int AL = 26;

struct Node
{
    int g[AL];
    int link;
    int len;
    int cnt;
    Node(): link(-1), len(0), cnt(1)
    {
        fill(g, g + AL, -1);
    }
};

struct Automaton
{
    vector<Node> a;
    int head;
    Automaton(): a(1), head(0) {}
    void add(char c)
    {
        // change to [0 AL)
        int ch = c - 'a';
        int nhead = sz(a);
        a.pb(Node());
        a[nhead].len = a[head].len + 1;
        int cur = head;
        head = nhead;
        while (cur != -1 && a[cur].g[ch] == -1)
        {
            a[cur].g[ch] = head;
            cur = a[cur].link;
        }
    }
};

```

```

if (cur == -1)
{
    a[head].link = 0;
    return;
}
int p = a[cur].g[ch];
if (a[p].len == a[cur].len + 1)
{
    a[head].link = p;
    return;
}
int q = sz(a);
a.pb(Node());
a[q] = a[p];
a[q].cnt = 0;
a[q].len = a[cur].len + 1;
a[p].link = a[head].link = q;
while (cur != -1 && a[cur].g[ch] == p)
{
    a[cur].g[ch] = q;
    cur = a[cur].link;
}
};

suffix-array.hpp

```

Description: Cast your string to vector. Don't forget about delimiters. No need to add anything at the end. sa represents permutations of positions if you sort all suffixes.

Time: $\mathcal{O}(n \log n)$

aa241e, 59 lines

```

void countSort(VI& p, const VI& c)
{
    int n = sz(p);
    VI cnt(n);
    FOR (i, 0, n)
        cnt[c[i]]++;
    VI pos(n);
    FOR (i, 1, n)
        pos[i] = pos[i - 1] + cnt[i - 1];
    VI p2(n);
    for (auto x : p)
    {
        int i = c[x];
        p2[pos[i]++] = x;
    }
    p = p2;
}

```

VI suffixArray(VI s)

```

{
    // strictly smaller than any other element
    s.pb(-1);
    int n = sz(s);
    VI p(n), c(n);
    iota(all(p), 0);
    sort(all(p), [&](int i, int j)
    {
        return s[i] < s[j];
    });
    int x = 0;
    c[p[0]] = 0;
    FOR (i, 1, n)
    {
        if (s[p[i]] != s[p[i - 1]])
            x++;
        c[p[i]] = x;
    }
    int k = 0;
    while ((1 << k) < n)

```

```

FOR (i, 0, n)
    p[i] = (p[i] - (1 << k) + n) % n;
countSort(p, c);
VI c2(n);
pii pr = {c[p[0]], c[(p[0] + (1 << k)) % n]};
FOR (i, 1, n)
{
    pii nx = {c[p[i]], c[(p[i] + (1 << k)) % n]};
    c2[p[i]] = c2[p[i - 1]];
    if (pr != nx)
        c2[p[i]]++;
    pr = nx;
}
c = c2;
k++;
}
p.erase(p.begin());
return p;
}

lcp.hpp
Description: queryLcp returns the longest common prefix of substrings starting
at i and j.
911c8c, 49 lines

```

```

struct LCP
{
    int n;
    VI s, sa, rnk, lcp;
    SparseTable st;

    LCP(VI _s): n(sz(_s)), s(_s)
    {
        sa = suffixArray(s);
        rnk.resize(n);
        FOR (i, 0, n)
            rnk[sa[i]] = i;
        lcpArray();
        FOR (i, 0, n - 1)
            st.pb(lcp[i]);
    }

    void lcpArray()
    {
        lcp.resize(n - 1);
        int h = 0;
        FOR (i, 0, n)
        {
            if (h > 0)
                h--;
            if (rnk[i] == 0)
                continue;
            int j = sa[rnk[i] - 1];
            for (; j + h < n && i + h < n; h++)
            {
                if (s[j + h] != s[i + h])
                    break;
            }
            lcp[rnk[i] - 1] = h;
        }
    }

    int queryLcp(int i, int j)
    {
        if (i == n || j == n)
            return 0;
        assert(i != j); // return n - i ??????
        i = rnk[i];
        j = rnk[j];
        if (i > j)
            swap(i, j);
    }
}

```

lcp run-enumerate suffix-tree

```

    // query (i, j)
    return st.query(i, j);
}

run-enumerate.hpp
Description: Enumerate all tuples  $(t, l, r)$  with  $t$  being the minimum period of  $s[l, r]$  and  $r - l \geq 2 \cdot t$ .  $l$  and  $r$  are maximal. In other words  $(t, l - 1, r)$  and  $(t, l, r + 1)$  do not satisfy the previous condition.
The number of runs is  $\leq |s|$ . Other properties are stated at the end of the function.
Time:  $\mathcal{O}(n \log n)$ , where  $n = |s|$ .
f9baf1, 62 lines

struct Run
{
    int t, l, r;
    bool operator<(const Run& p) const
    {
        return make_tuple(t, l, r) < make_tuple(p.t, p.l, p.r);
    }
    bool operator==(const Run& p) const
    {
        return !(*this < p) && !(p < *this);
    }
};

vector<Run> runEnumerate(VI s)
{
    int n = sz(s);
    LCP lcp(s); reverse(all(s));
    LCP rev(s); reverse(all(s));

    vector<Run> runs;
    FOR(inv, 0, 2)
    {
        VI st = {n};
        auto pop = [&](int i)
        {
            int j = st.back();
            int dist = j - i;
            int distPrev = st[sz(st) - 2] - j;
            int distMn = min(dist, distPrev);

            int len = lcp.queryLcp(i, j);
            if((len >= distMn && dist < distPrev) ||
               (len < distMn && ((s[i + len] < s[j + len]) ^ inv)))
                return true;
            return false;
        };

        RFOR(i, n, 0)
        {
            while(sz(st) > 1 && pop(i))
                st.pop_back();
            int j = st.back();
            int dist = j - i;
            st.pb(i);

            int x = rev.queryLcp(n - i, n - j);
            int y = lcp.queryLcp(i, j);
            if(x < dist && x + y >= dist)
                runs.pb((dist, i - x, j + y));
        }
    }
    sort(all(runs));
    runs.resize(unique(all(runs)) - runs.begin());
}

// ll sumLen = 0, sumCnt = 0, sum = 0;
// for(auto [len, l, r] : runs)
//     sumLen += len, sumCnt += (r - l) / len, sum += r - l;
// assert(sz(runs) <= sz(s));
// assert(sumLen <= LOG * sz(s));

```

```

//assert(sumCnt <= 2 * sz(s));
//assert(sum <= 2 * LOG * sz(s));
return runs;
}
```

suffix-tree.hpp

Description: Ukkonen's algorithm for building a suffix tree. Cast your string to vector. Don't forget about delimiters. $a[v].g[c]$ is a transition in format (u, l, r) , that goes from v to u and the string spelled out by this transition is the substring $s_{l..r}$. For transitions that go to leaves, $r = \text{INF}$. For the root node which has number 0, link == -1. For leaves, link == -2. For all other nodes, link is maintained explicitly.

Time: $\mathcal{O}(n \log |\Sigma|)$, where Σ is an alphabet

4aa61c, 85 lines

```

struct SuffixTree
{
    struct Transition
    {
        int u, l, r;
    };
    struct Node
    {
        map<int, Transition> g;
        int link;
        Node(): link(-2) {}
    };
    VI s;
    vector<Node> a;
    pair<bool, int> testAndSplit(int v, int l, int r, int c)
    {
        if (v == -1)
            return {true, -1};
        if (l <= r)
        {
            auto [nv, nl, nr] = a[v].g[s[l]];
            if (c == s[nl + r - l + 1])
                return {true, nv};
            int newNode = sz(a);
            a.pb(Node());
            a[v].g[s[l]] = {newNode, nl, nl + r - l};
            a[newNode].g[s[nl + r - l + 1]] = {nv, nl + r - l + 1, nr};
        }
        return {false, newNode};
    }
    return {a[v].g.count(c), v};
}

pii canonize(int v, int l, int r)
{
    if (v == -1 && l <= r)
    {
        v = 0;
        l++;
    }
    if (r < l)
        return {v, l};
    Transition cur = a[v].g[s[l]];
    while (cur.r - cur.l <= r - l)
    {
        l += cur.r - cur.l + 1;
        v = cur.u;
        if (l <= r)
            cur = a[v].g[s[l]];
    }
    return {v, l};
}

pii update(int v, int l, int r)
{
    int oldu = 0;
    auto [endPoint, u] = testAndSplit(v, l, r - 1, s[r]);
    while (!endPoint)

```

```

{
    int newNode = sz(a);
    a.pb(Node());
    a[u].g[s[r]] = {newNode, r, INF};
    if (oldu != 0)
        a[oldu].link = u;
    oldu = u;
    tie(v, l) = canonize(a[v].link, l, r - 1);
    tie(endPoint, u) = testAndSplit(v, l, r - 1, s[r]);
}
if (oldu != 0)
    a[oldu].link = v;
return {v, l};
}

SuffixTree(const VI& _s)
{
    s = _s;
    // Add the symbol that was not present in 's'
    s.pb(-1);
    a.reserve(2 * sz(s));
    a = {Node()};
    a[0].link = -1;
    int v = 0, l = 0;
    FOR(i, 0, sz(s))
    {
        tie(v, l) = update(v, l, i);
        tie(v, l) = canonize(v, l, i);
    }
}
};

z.hpp

```

9da7e8, 23 lines

```

VI zFunction(const string& s)
{
    int n = sz(s);
    VI z(n);

    int l = 0;
    int r = 0;
    FOR (i, 1, n)
    {
        z[i] = 0;
        if (i <= r)
            z[i] = min(r - i + 1, z[i - l]);

        while(i + z[i] < n && s[i + z[i]] == s[z[i]])
            z[i]++;
        if(i + z[i] - 1 > r)
        {
            r = i + z[i] - 1;
            l = i;
        }
    }
    return z;
}

prefix.hpp

```

5b81c4, 16 lines

```

VI prefixFunction(const string& s)
{
    int n = sz(s);
    VI p(n);
    p[0] = 0;
    FOR (i, 1, n)
    {
        int j = p[i - 1];
        while(j != 0 && s[i] != s[j])
            j = p[j - 1];
    }
}

```

z prefix minimal-cyclic-shift manacher palindromic-tree point

```

        if (s[i] == s[j]) j++;
        p[i] = j;
    }
    return p;
}

```

minimal-cyclic-shift.hpp

Description: $s_{shift}, s_{shift+1}, \dots$ is lexicographically smallest cyclic shift. If more than one answer it finds the minimum value of $shift$.
Time: $\mathcal{O}(n)$ time and memory complexity.

d4d30a, 29 lines

```

int minimalCyclicShift(VI s)
{
    int n = sz(s);
    s.resize(2 * n);
    FOR(i, 0, n)
        s[n + i] = s[i];

    int shift = 0;
    VI f(2 * n);
    FOR(i, 1, 2 * n)
    {
        int j = f[i - 1 - shift];
        while(j > 0 && s[shift + j] != s[i])
        {
            if(s[shift + j] > s[i])
                shift = i - j;
            j = f[j - 1];
        }
        if(j == 0 && s[shift] != s[i])
        {
            if(s[shift] > s[i])
                shift = i;
        }
        else
            j++;
        f[i - shift] = j;
    }
    return shift;
}

```

manacher.hpp

Description: $s[i - d0_i, i + d0_i - 1], s[i - d1_i + 1, i + d1_i - 1]$ are palindromes.

664138, 25 lines

```

vector<VI> manacher(const string& s)
{
    int n = sz(s);
    vector<VI> d(2);
    FOR (t, 0, 2)
    {
        d[t].resize(n);
        int l = -1;
        int r = -1;
        FOR (i, 0, n)
        {
            if (i <= r)
                d[t][i] = min(r - i + 1, d[t][l + (r - i) + 1 - t]);
            while (i + d[t][i] < n && i + t - d[t][i] - 1 >= 0
                   && s[i + d[t][i]] == s[i + t - d[t][i] - 1])
                d[t][i]++;
            if (i + d[t][i] - t > r)
            {
                r = i + d[t][i] - 1;
                l = i - d[t][i] + t;
            }
        }
        return d;
    }
}

```

palindromic-tree.hpp

62993e, 54 lines

```

const int AL = 26;

struct Node
{
    int to[AL];
    int link;
    int len;
    Node(int _link, int _len)
    {
        fill(to, to + AL, -1);
        link = _link;
        len = _len;
    }
};

struct PalTree
{
    string s;
    vector<Node> a;
    int last;

    PalTree(string t = "") : s(t), a({{-1, -1}, {0, 0}}), last(1)
    {
    }

    void add(int idx)
    {
        // change to /0, AL)
        int ch = s[idx] - 'a';

        int cur = last;
        while (cur != -1)
        {
            int pos = idx - a[cur].len - 1;
            if (pos >= 0 && s[pos] == s[idx])
                break;
            cur = a[cur].link;
        }
        if (a[cur].to[ch] == -1)
        {
            a[cur].to[ch] = sz(a);
            int link = a[cur].link;
            while (link != -1)
            {
                int pos = idx - a[link].len - 1;
                if (pos >= 0 && s[pos] == s[idx])
                    break;
                link = a[link].link;
            }
            if (link == -1)
                link = 1;
            else
                link = a[link].to[ch];
            a.pb(Node(link, a[cur].len + 2));
        }
        last = a[cur].to[ch];
    }
};

```

Geometry (5)

point.hpp

1a2063, 91 lines

```

struct Pt
{
    db x, y;
    Pt operator+(const Pt& p) const
    {
        return {x + p.x, y + p.y};
    }
}

```

```

} Pt operator-(const Pt& p) const
{
    return {x - p.x, y - p.y};
}
Pt operator*(db d) const
{
    return {x * d, y * d};
}
Pt operator/(db d) const
{
    return {x / d, y / d};
}
db sq(const Pt& p)
{
    return p.x * p.x + p.y * p.y;
}
db abs(const Pt& p)
{
    return sqrt(sq(p));
}
int sgn(db x)
{
    return (EPS < x) - (x < -EPS);
}
// Returns 'p' rotated counter-clockwise by 'a'
Pt rot(const Pt& p, db a)
{
    db co = cos(a), si = sin(a);
    return {p.x * co - p.y * si,
            p.x * si + p.y * co};
}
// Returns 'p' rotated counter-clockwise by 90 degrees
Pt perp(const Pt& p)
{
    return {-p.y, p.x};
}
db dot(const Pt& p, const Pt& q)
{
    return p.x * q.x + p.y * q.y;
}
// Returns the angle between 'p' and 'q' in [0, pi]
db angle(const Pt& p, const Pt& q)
{
    return acos(clamp(dot(p, q) / abs(p) /
        abs(q), (db)-1.0, (db)1.0));
}
db cross(const Pt& p, const Pt& q)
{
    return p.x * q.y - p.y * q.x;
}
// Positive if R is on the left side of PQ,
// negative on the right side,
// and zero if R is on the line containing PQ
db orient(const Pt& p, const Pt& q, const Pt& r)
{
    return cross(q - p, r - p) / abs(q - p);
}
// Checks if argument of 'p' is in [-pi, 0)
bool half(const Pt& p)
{
    assert(sgn(p.x) != 0 || sgn(p.y) != 0);
    return sgn(p.y) == -1 ||
        (sgn(p.y) == 0 && sgn(p.x) == -1);
}
void polarSortAround(const Pt& o, vector<Pt>& v)
{
    sort(all(v), [o](Pt p, Pt q)

```

```

{
    p = p - o;
    q = q - o;
    bool hp = half(p), hq = half(q);
    if (hp != hq)
        return hp < hq;
    int s = sgn(cross(p, q));
    if (s != 0)
        return s == 1;
    return sq(p) < sq(q);
});
ostream& operator<<(ostream& os, const Pt& p)
{
    return os << "(" << p.x << ", " << p.y << ")";
}

line.hpp
83c9af, 50 lines
struct Line
{
    // Equation of the line is dot(n, p) + c = 0
    Pt n;
    db c;
    Line(const Pt& _n, db _c): n(_n), c(_c) {}
    // n is the normal vector to the left of PQ
    Line(const Pt& p, const Pt& q):
        n(perp(q - p)), c(-dot(n, p)) {}
    // The "positive side": dot(n, p) + c > 0
    // The "negative side": dot(n, p) + c < 0
    db side(const Pt& p) const
    {
        return dot(n, p) + c;
    }
    db dist(const Pt& p) const
    {
        return abs(side(p)) / abs(n);
    }
    db sqDist(const Pt& p) const
    {
        return side(p) * side(p) / (db)sq(n);
    }
    Line perpThrough(const Pt& p) const
    {
        return {p, p + n};
    }
    bool cmpProj(const Pt& p, const Pt& q) const
    {
        return sgn(cross(p, n) - cross(q, n)) < 0;
    }
    Pt proj(const Pt& p) const
    {
        return p - n * side(p) / sq(n);
    }
    Pt reflect(const Pt& p) const
    {
        return p - n * 2 * side(p) / sq(n);
    }
};

bool parallel(const Line& l1, const Line& l2)
{
    return sgn(cross(l1.n, l2.n)) == 0;
}
Pt inter(const Line& l1, const Line& l2)
{
    db d = cross(l1.n, l2.n);
    assert(sgn(d) != 0);
    return perp(l2.n * l1.c - l1.n * l2.c) / d;
}

```

687634, 39 lines

```

segment.hpp
bool inDisk(const Pt& a, const Pt& b, const Pt& p)
{
    return sgn(dot(a - p, b - p)) <= 0;
}
// Checks if 'p' lies on segment [ab]
bool onSegment(const Pt& a, const Pt& b, const Pt& p)
{
    return sgn(orient(a, b, p)) == 0 && inDisk(a, b, p);
}
// Checks if the segments [ab] and [cd] intersect
// properly (their intersection is one point
// which is not an endpoint of either segment)
bool properInter(const Pt& a, const Pt& b, const Pt& c, const Pt& d)
{
    db oa = orient(c, d, a);
    db ob = orient(c, d, b);
    db oc = orient(a, b, c);
    db od = orient(a, b, d);
    return sgn(oa) * sgn(ob) == -1 && sgn(oc) * sgn(od) == -1;
}
// Returns the distance between [ab] and 'p'
db segPt(const Pt& a, const Pt& b, const Pt& p)
{
    Line l(a, b);
    assert(sgn(sq(l.n)) != 0);
    if (l.cmpProj(a, p) && l.cmpProj(p, b))
        return l.dist(p);
    return min(abs(p - a), abs(p - b));
}
// Returns the distance between [ab] and [cd]
db segSeg(const Pt& a, const Pt& b, const Pt& c, const Pt& d)
{
    if (properInter(a, b, c, d))
        return 0;
    return min({segPt(a, b, c), segPt(a, b, d),
                segPt(c, d, a), segPt(c, d, b)});
}

polygon.hpp
d2cc47, 67 lines
bool isConvex(const vector<Pt>& v)
{
    bool hasPos = false, hasNeg = false;
    int n = sz(v);
    FOR(i, 0, n)
    {
        int s = sgn(orient(v[i], v[(i + 1) % n], v[(i + 2) % n]));
        hasPos |= s > 0;
        hasNeg |= s < 0;
    }
    return !(hasPos && hasNeg);
}
db areaTriangle(const Pt& a, const Pt& b, const Pt& c)
{
    return abs(cross(b - a, c - a)) / 2.0;
}
db areaPolygon(const vector<Pt>& v)
{
    db area = 0.0;
    int n = sz(v);
    FOR(i, 0, n)
        area += cross(v[i], v[(i + 1) % n]);
    return abs(area) / 2.0;
}
// Checks if point 'a' is inside the convex

```

```
// polygon 'v'. Returns true if on the boundary.
// 'v' must not contain duplicated vertices.
// Time: O(log n)
bool inConvexPolygon(const vector<Pt>& v, const Pt& a)
{
    assert(sz(v) >= 2);
    if (sz(v) == 2)
        return onSegment(v[0], v[1], a);
    if (sgn(orient(v.back(), v[0], a)) < 0
        || sgn(orient(v[0], v[1], a)) < 0)
        return false;
    int i = lower_bound(v.begin() + 2, v.end(), a,
    [&](const Pt& p, const Pt& q)
    {
        return sgn(orient(v[0], p, q)) > 0;
    }) - v.begin();
    return sgn(orient(v[i - 1], v[i], a)) >= 0;
}
bool above(const Pt& a, const Pt& p)
{
    return sgn(p.y - a.y) >= 0;
}
bool crossesRay(const Pt& a, const Pt& p,
    const Pt& q)
{
    return sgn((above(a, q) - above(a, p))
        * orient(a, p, q)) == 1;
}
// Checks if point 'a' is inside the polygon
// If 'strict', false when 'a' is on the boundary
bool inPolygon(const vector<Pt>& v, const Pt& a, bool strict =
    true)
{
    int numCrossings = 0;
    int n = sz(v);
    FOR(i, 0, n)
    {
        if (onSegment(v[i], v[(i + 1) % n], a))
            return !strict;
        numCrossings += crossesRay(a, v[i], v[(i + 1) % n]);
    }
    return numCrossings & 1;
}
```

convex-hull.hpp 67f952, 27 lines

```
vector<Pt> convexHull(vector<Pt> v, bool include_collinear =
    false)
{
    if (sz(v) <= 1)
        return v;
    sort(all(v), [](const Pt& p, const Pt& q)
    {
        int dx = sgn(p.x - q.x);
        if (dx != 0)
            return dx < 0;
        return sgn(p.y - q.y) < 0;
    });
    vector<Pt> lower, upper;
    for (const Pt& p : v)
    {
        while (sz(lower) > 1 &&
            sgn(orient(lower[sz(lower) - 2], lower.back(), p)) < (
                include_collinear ? 0 : 1))
            lower.pop_back();
        while (sz(upper) > 1 &&
            sgn(orient(upper[sz(upper) - 2], upper.back(), p)) > (
                include_collinear ? 0 : -1))
            upper.pop_back();
    }
```

convex-hull tangents-to-convex-polygon minkowski-sum ear-clipping

```
lower.pb(p);
upper.pb(p);
}
reverse(all(upper));
lower.insert(lower.end(), next(upper.begin()), prev(upper.end()
()));
return lower;
}
```

tangents-to-convex-polygon.hpp

Description: Returns the indices of tangent points from p . p must be strictly outside the polygon.

```
vector<Pt> res;
int n1 = sz(v1), n2 = sz(v2),
j1 = 0, j2 = 0;
while (j1 < n1 || j2 < n2)
{
    const Pt& p1 = v1[(i1 + j1) % n1];
    const Pt& q1 = v1[(i1 + j1 + 1) % n1];
    const Pt& p2 = v2[(i2 + j2) % n2];
    const Pt& q2 = v2[(i2 + j2 + 1) % n2];
    if (sz(res) >= 2 && onSegment(res[sz(res) - 2], p1 + p2,
        res.back()))
        res.pop_back();
    res.pb(p1 + p2);
    int s = sgn(cross(q1 - p1, q2 - p2));
    if (j1 < n1 && (j2 == n2 || s > 0
        || (s == 0 && (sz(res) < 2
        || sgn(dot(res.back() - res[sz(res) - 2],
            q1 + p2 - res.back())) > 0))))
        j1++;
    else
        j2++;
}
if (sz(res) > 2 && onSegment(res[sz(res) - 2], res[0], res.
back()))
    res.pop_back();
return res;
}
```

ear-clipping.hpp

Description: Finds an arbitrary triangulation of a simple polygon with no three collinear vertices.

```
0252d5, 55 lines
vector<tuple<int, int, int>> earClipping(const vector<Pt>& v)
{
    int n = sz(v);
    vector<tuple<int, int, int>> res;
    VI indices(n), ear(n), reflex(n);
    iota(all(indices), 0);
    auto updReflexStatus = [&](int i)
    {
        int sz = sz(indices),
            pos = find(all(indices), i) - indices.begin();
        int iPrev = indices[(pos + sz - 1) % sz],
            iNext = indices[(pos + 1) % sz];
        reflex[i] = orient(v[iPrev], v[i], v[iNext]) < 0;
    };
    auto updEarStatus = [&](int i)
    {
        if (reflex[i])
        {
            ear[i] = 0;
            return;
        }
        int sz = sz(indices),
            pos = find(all(indices), i) - indices.begin();
        int iPrev = indices[(pos + sz - 1) % sz],
            iNext = indices[(pos + 1) % sz];
        ear[i] = 1;
        for (int j : indices)
        {
            if (j != iPrev && j != i && j != iNext && reflex[j])
                && inConvexPolygon({v[iPrev], v[i], v[iNext]}, v[j]))
            {
                ear[i] = 0;
                break;
            }
        }
    };
    FOR(i, 0, n)
```

```

updReflexStatus(i);
FOR(i, 0, n)
    updEarStatus(i);
RFOR(sz, n + 1, 3)
{
    int i = 0;
    while (!ear[indices[i]])
        i++;
    int iPrev = indices[(i + sz - 1) % sz], iNext = indices[(i + 1) % sz];
    res.pb({iPrev, indices[i], iNext});
    indices.erase(indices.begin() + i);
    updReflexStatus(iPrev);
    updReflexStatus(iNext);
    updEarStatus(iPrev);
    updEarStatus(iNext);
}
return res;
}

```

halfplane-intersection.hpp

Description: Returns the counter-clockwise ordered vertices of the half-plane intersection. Returns empty if the intersection is empty. Adds a bounding box to ensure a finite area.

cf6d03, 47 lines

```

vector<Pt> hplaneInter(vector<Line> lines)
{
    const db C = 1e9;
    lines.pb({{-C, C}, {-C, -C}});
    lines.pb({{-C, -C}, {C, -C}});
    lines.pb({{C, -C}, {C, C}});
    lines.pb({{C, C}, {-C, C}});
    sort(all(lines), [](const Line& l1, const Line& l2)
    {
        bool h1 = half(l1.n), h2 = half(l2.n);
        if (h1 != h2)
            return h1 < h2;
        int p = sgn(cross(l1.n, l2.n));
        if (p != 0)
            return p > 0;
        return sgn(l1.c / abs(l1.n) - l2.c / abs(l2.n)) < 0;
    });
    lines.erase(unique(all(lines), parallel), lines.end());
    deque<pair<Line, Pt>> d;
    for (const Line& l : lines)
    {
        while (sz(d) > 1 && sgn(l.side((d.end() - 1)->y)) < 0)
            d.pop_back();
        while (sz(d) > 1 && sgn(l.side((d.begin() + 1)->y)) < 0)
            d.pop_front();
        if (!d.empty() && sgn(cross(d.back().x.n, l.n)) <= 0)
            return {};
        if (sz(d) < 2 || sgn(d.front().x.side(inter(l, d.back().x)))
            ) >= 0)
        {
            Pt p;
            if (!d.empty())
            {
                p = inter(l, d.back().x);
                if (!parallel(l, d.front().x))
                    d.front().y = inter(l, d.front().x);
            }
            d.pb({l, p});
        }
    }
    vector<Pt> res;
    for (auto [l, p] : d)
    {
        if (res.empty() || sgn(sq(p - res.back())) > 0)
            res.pb(p);
    }
}

```

halfplane-intersection circle tangents welzl closest-pair

```

    }
    return res;
}

circle.hpp
// Returns the circumcenter of triangle abc.
// The circumcircle of a triangle is a circle that passes
// through all three vertices.
Pt circumCenter(const Pt& a, Pt b, Pt c)
{
    b = b - a;
    c = c - a;
    assert(sgn(cross(b, c)) != 0);
    return a + perp(b * sq(c) - c * sq(b)) / cross(b, c) / 2;
}

// Returns circle-line intersection points
vector<Pt> circleLine(const Pt& o, db r, const Line& l)
{
    db h2 = r * r - l.sqDist(o);
    if (sgn(h2) == -1)
        return {};
    Pt p = l.proj(o);
    if (sgn(h2) == 0)
        return {p};
    Pt h = perp(l.n) * sqrt(h2) / abs(l.n);
    return {p - h, p + h};
}

// Returns circle-circle intersection points
vector<Pt> circleCircle(const Pt& o1, db r1, const Pt& o2, db r2)
{
    Pt d = o2 - o1;
    db d2 = sq(d);
    if (sgn(d2) == 0)
    {
        // assuming the circles don't coincide
        assert(sgn(r2 - r1) != 0);
        return {};
    }
    db pd = (d2 + r1 * r1 - r2 * r2) / 2;
    db h2 = r1 * r1 - pd * pd / d2;
    if (sgn(h2) == -1)
        return {};
    Pt p = o1 + d * pd / d2;
    if (sgn(h2) == 0)
        return {p};
    Pt h = perp(d) * sqrt(h2 / d2);
    return {p - h, p + h};
}

tangents.hpp
Description: Finds common tangents (outer or inner) to two circles. If there are
two tangents, returns the pairs of tangency points on each circle ( $p_1, p_2$ ). If there
is one tangent, the circles are tangent to each other at some point  $p$ , res contains  $p$ 
four times, and the tangent line can be found as line( $o_1, p$ ).perpThrough( $p$ ). The
same code can be used to find the tangent to a circle through a point by setting  $r_2$ 
to 0 (in which case inner doesn't matter).

```

```

vector<pair<Pt, Pt>> tangents(const Pt& o1, db r1,
                                    const Pt& o2, db r2, bool inner)
{
    if (inner)
        r2 = -r2;
    Pt d = o2 - o1;
    db dr = r1 - r2, d2 = sq(d), h2 = d2 - dr * dr;
    if (sgn(d2) == 0 || sgn(h2) < 0)
    {
        assert(sgn(h2) != 0);
        return {};
    }
}

```

```

    }
    vector<pair<Pt, Pt>> res;
    for (db sign : {-1, 1})
    {
        Pt v = (d * dr + perp(d) * sqrt(h2) * sign) / d2;
        res.pb({o1 + v * r1, o2 + v * r2});
    }
    return res;
}

```

welzl.hpp

Description: Returns the smallest enclosing circle of points in v
Time: $\mathcal{O}(n)$ (expected)

f6000c, 36 lines

```

pair<Pt, db> welzl(vector<Pt> v)
{
    int n = sz(v), k = 0, idxes[2];
    mt19937 rng;
    shuffle(all(v), rng);
    Pt c = v[0];
    db r = 0;
    while (true)
    {
        FOR(i, k, n)
        {
            if (sgn(abs(v[i] - c) - r) > 0)
            {
                swap(v[i], v[k]);
                if (k == 0)
                    c = v[0];
                else if (k == 1)
                    c = (v[0] + v[1]) / 2;
                else
                    c = circumCenter(v[0], v[1], v[2]);
                r = abs(v[0] - c);
                if (k < i)
                {
                    if (k < 2)
                        idxes[k++] = i;
                    shuffle(v.begin() + k, v.begin() + i + 1, rng);
                    break;
                }
            }
        }
        while (k > 0 && idxes[k - 1] == i)
            k--;
        if (i == n - 1)
            return {c, r};
    }
}

```

closest-pair.hpp

Description: Returns the distance between the closest points
Time: $\mathcal{O}(n \log n)$

678ecf, 23 lines

```

db closestPair(vector<Pt> v)
{
    sort(all(v), [](const Pt& p, const Pt& q)
    {
        return sgn(p.x - q.x) < 0;
    });
    set<pd> s;
    int n = sz(v), ptr = 0;
    db h = 1e18;
    FOR(i, 0, n)
    {
        for (auto it = s.lower_bound(MP(v[i].y - h, v[i].x));
              it != s.end() && sgn(it->x - (v[i].y + h)) <= 0; it++)
        {
            Pt q = {it->y, it->x};
            if (sq(p - q) < h)
                h = sq(p - q);
        }
    }
}

```

```

    h = min(h, abs(v[i] - q));
}
for (; sgn(v[ptr].x - (v[i].x - h)) <= 0; ptr++)
    s.erase({v[ptr].y, v[ptr].x});
s.insert({v[i].y, v[i].x});
}
return h;
}

```

planar-graph.hpp

Description: Finds faces in a planar graph. Use `addVertex()` and `addEdge()` for initializing the graph and `addQueryPoint()` for initializing the queries. After initialization, call `findFaces()` before using other functions. `getIncidentFaces(i)` returns the pair of faces (u, v) (possibly $u = v$) such that the i -th edge lies on the boundary of these faces. `getFaceOfQueryPoint(i)` returns the face where the i -th query point lies.

939539, 169 lines

```

namespace PlanarGraph
{
    struct IndexedPt
    {
        Pt p;
        int index;
        bool operator<(const IndexedPt& q) const
        {
            return p.x < q.p.x;
        }
    };
    struct Edge
    {
        // cross(vertices[j].p - vertices[i].p, l.n) > 0
        int i, j;
        Line l;
    };
    vector<IndexedPt> vertices, queryPoints;
    vector<Edge> edges;
    struct Comparator
    {
        using is_transparent = void;
        static IndexedPt vertex;
        db getY(const Line& l) const
        {
            return -(l.n.x * vertex.p.x + l.c) / l.n.y;
        }
        bool operator()(int i, int j) const
        {
            auto [u1, v1, l1] = edges[i];
            auto [u2, v2, l2] = edges[j];
            if (u1 == vertex.index && u2 == vertex.index)
                return sgn(cross(l1.n, l2.n)) > 0;
            if (v1 == vertex.index && v2 == vertex.index)
                return sgn(cross(l1.n, l2.n)) < 0;
            int dy = sgn(getY(l1) - getY(l2));
            assert(dy != 0);
            return dy < 0;
        }
        bool operator()(int i, const Pt& p) const
        {
            int dy = sgn(getY(edges[i].l) - p.y);
            assert(dy != 0);
            return dy < 0;
        }
    } comparator;
    IndexedPt Comparator::vertex;
    DSU dsu;
    VI upperFace, queryAns;
    void addVertex(const Pt& p)
    {
        vertices.pb({p, sz(vertices)});
    }

```

planar-graph modular-arithmetics

```

    }
    void addEdge(int i, int j, const Line& l)
    {
        assert(0 <= i && i < sz(vertices));
        assert(0 <= j && j < sz(vertices));
        assert(i != j);
        assert(vertices[i].index == i);
        assert(vertices[j].index == j);
        edges.pb({i, j, l});
    }
    void addEdge(int i, int j)
    {
        addEdge(i, j, {vertices[i].p, vertices[j].p});
    }
    void addQueryPoint(const Pt& p)
    {
        queryPoints.pb({p, sz(queryPoints)});
    }
    void findFaces()
    {
        int n = sz(vertices), m = sz(edges);
        const db ROT_ANGLE = 4;
        for (auto& p : vertices)
            p.p = rot(p.p, ROT_ANGLE);
        for (auto& p : queryPoints)
            p.p = rot(p.p, ROT_ANGLE);
        vector<VI> edgesL(n), edgesR(n);
        FOR(k, 0, m)
        {
            auto& [i, j, l] = edges[k];
            l.n = rot(l.n, ROT_ANGLE);
            if (vertices[i].p.x > vertices[j].p.x)
            {
                swap(i, j);
                l.n = l.n * (-1);
                l.c *= -1;
            }
            edgesL[j].pb(k);
            edgesR[i].pb(k);
        }
        sort(all(vertices));
        sort(all(queryPoints));
        // when choosing INF, remember that we rotate the plane
        addVertex({-INF, INF});
        addVertex({INF, INF});
        addEdge(n, n + 1);
        dsu = DSU(m + 1);
        set<int, Comparator> s;
        s.insert(m);
        upperFace.resize(m);
        int ptr = 0;
        queryAns.resize(sz(queryPoints));
        for (const IndexedPt& vertex : vertices)
        {
            int i = vertex.index;
            while (ptr < sz(queryPoints)
                   && (i >= n || queryPoints[ptr] < vertex))
            {
                const auto& [pt, j] = queryPoints[ptr++];
                Comparator::vertex = {pt, -1};
                queryAns[j] = *s.lower_bound(pt);
            }
            if (i >= n)
                break;
            Comparator::vertex = vertex;
            int upper = -1, lower = -1;
            if (!edgesL[i].empty())
            {
                sort(all(edgesL[i]), comparator);

```

```

                auto it = s.lower_bound(edgesL[i][0]);
                lower = edgesL[i][0];
                for (int e : edgesL[i])
                {
                    assert(*it == e);
                    assert(next(it) != s.end());
                    upperFace[e] = *next(it);
                    it = s.erase(it);
                }
                assert(it != s.end());
                upper = *it;
            }
            if (!edgesR[i].empty())
            {
                sort(all(edgesR[i]), comparator);
                if (upper == -1)
                {
                    upper = *s.lower_bound(edgesR[i][0]);
                }
                int prv = -1;
                for (int e : edgesR[i])
                {
                    s.insert(e);
                    if (prv != -1)
                    {
                        upperFace[prv] = e;
                    }
                    prv = e;
                }
                upperFace[edgesR[i].back()] = upper;
                dsu.unite(edgesL[i].empty() ? upper : lower, edgesR[i][0]);
            }
            else if (lower != -1 && upper != -1)
            {
                dsu.unite(upper, lower);
            }
        }
        pii getIncidentFaces(int i)
        {
            return {dsu.find(i), dsu.find(upperFace[i])};
        }
        int getFaceOfQueryPoint(int i)
        {
            return dsu.find(queryAns[i]);
        }
    };

```

Mathematics (6)

Number-theoretic algorithms

```

modular-arithmetics.hpp
const int MOD = 998244353;

int add(int a, int b)
{
    return a + b < MOD ? a + b : a + b - MOD;
}

void updAdd(int& a, int b)
{
    a += b;
    if (a >= MOD)
        a -= MOD;
}

```

6271b9, 67 lines

```

int sub(int a, int b)
{
    return a - b >= 0 ? a - b : a - b + MOD;
}

void updSub(int& a, int b)
{
    a -= b;
    if (a < 0)
        a += MOD;
}

int mult(int a, int b)
{
    return (ll)a * b % MOD;
}

int binPow(int a, ll n)
{
    int res = 1;
    while (n)
    {
        if (n & 1)
            res = mult(res, a);
        a = mult(a, a);
        n >>= 1;
    }
    return res;
}

int inv[N], fact[N], ifact[N];

void init()
{
    inv[1] = 1;
    FOR(i, 2, N)
    {
        inv[i] = mult(MOD - MOD / i, inv[MOD % i]);
    }
    fact[0] = ifact[0] = 1;
    FOR(i, 1, N)
    {
        fact[i] = mult(fact[i - 1], i);
        ifact[i] = mult(ifact[i - 1], inv[i]);
    }
}

int C(int n, int k)
{
    if (k < 0 || k > n)
        return 0;
    return mult(fact[n], mult(ifact[n - k], ifact[k]));
}

gcd.hpp
Description:  $ax + by = d$ ,  $\gcd(a, b) = |d| \rightarrow (d, x, y)$ .
Minimizes  $|x| + |y|$ . And minimizes  $|x - y|$  for  $a > 0, b > 0$ .
bcd80c, 16 lines
tuple<ll, ll, ll> gcdExt(ll a, ll b)
{
    ll x1 = 1, y1 = 0;
    ll x2 = 0, y2 = 1;
    while (b)
    {
        ll k = a / b;
        x1 -= k * x2;
        y1 -= k * y2;
        a %= b;
        swap(a, b);
        swap(x1, x2);
    }
}

```

gcd fast-chinese miller-rabin pollard

```

        swap(y1, y2);
    }
    return {a, x1, y1};
}

fast-chinese.hpp
Description:  $x \% p_i = m_i$ ,  $\text{lcm}(p_i) \leq 10^{18}$ ,  $0 \leq x < \text{lcm}(p_i) \rightarrow x$  or  $-1$ .
Time:  $\mathcal{O}(n \log(\text{lcm}(p_i)))$ 
046449, 24 lines
ll fastChinese(vector<ll> m, vector<ll> p)
{
    assert(sz(m) == sz(p));
    ll aa = p[0];
    ll bb = m[0];
    FOR(i, 1, sz(m))
    {
        ll b = (m[i] - bb % p[i] + p[i]) % p[i];
        ll a = aa % p[i];
        ll c = p[i];

        auto [d, x, y] = gcdExt(a, c);
        if(b % d != 0)
            return -1;
        a /= d;
        b /= d;
        c /= d;
        b = (b * (__int128)x % c + c) % c;

        bb = aa * b + bb;
        aa = aa * c;
    }
    return bb;
}

miller-rabin.hpp
Description: To speed up change candidates to at least 4 random values  $\text{rng}() \% (n - 3) + 2$ . Use  $\text{__int128}$  in mult.
Time:  $\mathcal{O}(|\text{candidates}| \cdot \log n)$ 
2f89bb, 33 lines
VI candidates = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 47};
bool millerRabin(ll n)
{
    if (n == 1)
        return false;
    if (n == 2 || n == 3)
        return true;
    ll d = n - 1;
    int s = __builtin_ctzll(d);
    d >>= s;

    for (ll b : candidates)
    {
        if (b >= n)
            break;
        b = binpow(b, d, n);
        if (b == 1)
            continue;
        bool ok = false;
        FOR (i, 0, s)
        {
            if (b + 1 == n)
            {
                ok = true;
                break;
            }
            b = mult(b, b, n);
        }
        if (!ok)
            return false;
    }
}

```

```

    return true;
}

pollard.hpp
Description: Uses the Miller-Rabin test. rho finds a divisor of  $n$ . Use  $\text{__int128}$  in mult.
Time:  $\mathcal{O}(n^{1/4} \cdot \log n)$ 
69a916, 62 lines
ll f(ll x, ll c, ll n)
{
    return add(mult(x, x, n), c, n);
}

ll rho(ll n)
{
    const int iter = 47 * pow(n, 0.25);
    while (true)
    {
        ll x0 = rng() % n;
        ll c = rng() % n;
        ll x = x0;
        ll y = x0;
        ll g = 1;
        FOR (i, 0, iter)
        {
            x = f(x, c, n);
            y = f(y, c, n);
            y = f(y, c, n);
            g = gcd(abs(x - y), n);
            if (g != 1)
                break;
        }
        if (g > 1 && g < n)
            return g;
    }
}
VI primes = {2, 3, 5, 7, 11, 13, 17, 19, 23};

VL factorize(ll n)
{
    VL ans;

    for (auto p : primes)
    {
        while (n % p == 0)
        {
            ans.pb(p);
            n /= p;
        }
    }
    queue<ll> q;
    q.push(n);

    while (!q.empty())
    {
        ll x = q.front();
        q.pop();
        if (x == 1)
            continue;
        if (millerRabin(x))
            ans.pb(x);
        else
        {
            ll y = rho(x);
            q.push(y);
            q.push(x / y);
        }
    }
    return ans;
}

```

floor-sum.hpp

Description: Computes $\sum_{i=0}^{n-1} \left\lfloor \frac{a \cdot i + b}{m} \right\rfloor$.

Time: $\mathcal{O}(\log m)$.

floor-sum min-mod-linear mod-inequality disLog gaussian hungarian

ll floorSum(ll n, ll m, ll a, ll b)

```
{
    ll ans = 0;
    while (true)
    {
        ans += (a / m) * n * (n - 1) / 2 + (b / m) * n;
        a %= m;
        b %= m;
        if (a == 0)
            return ans;
        ll k = (a * (n - 1) + b) / m;
        b = a * n - m * k + b;
        n = k;
        swap(a, m);
    }
}
```

min-mod-linear.hpp

Description: Finds $\min\{(ax + b) \bmod m \mid 0 \leq x < n\}$.

Time: $\mathcal{O}(\log m)$.

9517db, 16 lines

int minModLinear(ll n, ll m, ll a, ll b)

```
{
    ll res = m;
    while (n > 0)
    {
        a %= m;
        b = (b % m + m) % m;
        res = min(res, b);
        n = (a * (n - 1) + b) / m;
        b -= m * n;
        swap(a, m);
    }
    return res;
}
```

03b25c, 14 lines

mod-inequality.hpp

Description: Finds the smallest $x \geq 0$ such that $(ax + b) \bmod m \geq c$. Returns -1 , if the solution does not exist.

Time: $\mathcal{O}(\log m)$.

4a4b4a, 15 lines

```
int modInequality(ll m, ll a, ll b, ll c)
{
    a %= m;
    b %= m;
    if (b >= c)
        return 0;
    if (a == 0)
        return -1;
    if (c + a < m)
        return (c - b + a - 1) / a;
    int k = modInequality(a, m, c - b - 1, c + a - m);
    if (k == -1)
        return -1;
    return (k * m + c - b + a - 1) / a;
}
```

disLog.hpp

a986d8, 23 lines

```
// Returns minimum x for which  $(a ^ x) \% MOD = b \% MOD$ , a and MOD are coprime.
int disLog(int a, int b)
{
    int n = sqrt(MOD) + 1;
    int an = binPow(a, n);
    unordered_map<int, int> vals;
```

```

for (int q = 0, cur = b; q <= n; ++q)
{
    vals[cur] = q;
    cur = mult(cur, a);
}

for (int p = 1, cur = 1; p <= n; ++p)
{
    cur = mult(cur, an);
    if (vals.count(cur))
    {
        return n * p - vals[cur];
    }
}
return -1;
}
```

Matrices

gaussian.hpp

Description: Solves the system $Ax = b$. Returns (v, w) such that every solution x can be represented as $v + c_1w_1 + c_2w_2 + \dots + c_kw_k$, where v is arbitrary solution, c_i are scalars and w is basis. If there is no solution, returns an empty pair. If the solution is unique, then w is empty.

Time: $\mathcal{O}(nm \min(n, m))$

3fa52c, 66 lines

pair<VI, vector<VI>> solveLinearSystem(vector<VI> a, VI b)

```

int n = sz(a), m = sz(a[0]);
assert(sz(b) == n);
FOR(i, 0, n)
{
    assert(sz(a[i]) == m);
    a[i].pb(b[i]);
}
int p = 0;
VI pivots;
FOR(j, 0, m)
{
    // with doubles, abs(a[p]/j) -> max
    if (a[p][j] == 0)
    {
        int l = -1;
        FOR(i, p, n)
            if (a[i][j] != 0)
                l = i;
        if (l == -1)
            continue;
        swap(a[p], a[l]);
    }
    int in = binPow(a[p][j], MOD - 2);
    FOR(i, p + 1, n)
    {
        int c = mult(a[i][j], in);
        FOR(k, j, m + 1)
            updSub(a[i][k], mult(c, a[p][k]));
    }
    pivots.pb(j);
    p++;
    if (p == n)
        break;
}
FOR(i, p, n)
{
    if (a[i].back() != 0)
        return {};
    VI v(m);
    RFOR(i, p, 0)
    {
        int j = pivots[i];
        v[j] = a[i].back();
        FOR(k, j + 1, m)

```

```

        updSub(v[j], mult(a[i][k], v[k]));
        v[j] = mult(v[j], binPow(a[i][j], MOD - 2));
    }
    vector<VI> w;
    FOR(q, 0, m)
    {
        if (find(all(pivots), q) != pivots.end())
            continue;
        VI d(m);
        d[q] = 1;
        RFOR(i, p, 0)
        {
            int j = pivots[i];
            FOR(k, j + 1, m)
                updSub(d[j], mult(a[i][k], d[k]));
            d[j] = mult(d[j], binPow(a[i][j], MOD - 2));
        }
        w.pb(d);
    }
    return {v, w};
}
```

hungarian.hpp

Description: Finds a maximum matching that has the minimum weight in a weighted bipartite graph.

Time: $\mathcal{O}(n^2m)$

792894, 63 lines

ll hungarian(const vector<VL>& a)

```

int n = sz(a), m = sz(a[0]);
assert(n <= m);
VL u(n + 1), v(m + 1);
VI p(m + 1, n), way(m + 1);
FOR(i, 0, n)
{
    p[m] = i;
    int j0 = m;
    VL minv(m + 1, LINF);
    VI used(m + 1);
    while (p[j0] != n)
    {
        used[j0] = true;
        int i0 = p[j0], j1 = -1;
        ll delta = LINF;
        FOR(j, 0, m)
        {
            if (!used[j])
            {
                ll cur = a[i0][j] - u[i0] - v[j];
                if (cur < minv[j])
                {
                    minv[j] = cur;
                    way[j] = j0;
                }
                if (minv[j] < delta)
                {
                    delta = minv[j];
                    j1 = j;
                }
            }
        }
        assert(j1 != -1);
        FOR(j, 0, m + 1)
        {
            if (used[j])
            {
                u[p[j]] += delta;
                v[j] -= delta;
            }
        }
    }
}
```

```

    else
        minv[j] -= delta;
}
j0 = j1;
}
while (j0 != m)
{
    int j1 = way[j0];
    p[j0] = p[j1];
    j0 = j1;
}
VI ans(n + 1);
FOR(j, 0, m)
    ans[p[j]] = j;
ll res = 0;
FOR(i, 0, n)
    res += a[i][ans[i]];
assert(res == -v[m]);
return res;
}

```

Polynomials and FFT

fft.hpp

Description: Number-theoretic transform. If you need complex-valued FFT, use the commented out code.

Time: $\mathcal{O}(n \log n)$

1a18a5, 73 lines

```

const int LEN = 1 << 23;
const int GEN = 31;

/*typedef complex<db> com;
com pw[LEN];
void init()
{
    db phi = (db)2 * PI / LEN;
    FOR(i, 0, LEN)
        pw[i] = com(cos(phi * i), sin(phi * i));
}*/
```

```

void fft(VI& a, bool inverse)
{
    const int IGEN = binPow(GEN, MOD - 2);
    int lg = __builtin_ctz(sz(a));
    FOR(i, 0, sz(a))
    {
        int k = 0;
        FOR(j, 0, lg)
            k |= ((i >> j) & 1) << (lg - j - 1);
        if(i < k)
            swap(a[i], a[k]);
    }
    for(int len = 2; len <= sz(a); len *= 2)
    {
        // int diff = inv ? LEN - LEN / len : LEN / len;
        int ml = binPow(inverse ? IGEN : GEN, LEN / len);
        for(int i = 0; i < sz(a); i += len)
        {
            // int pos = 0;
            int pw = 1;
            FOR(j, 0, len / 2)
            {
                int u = a[i + j];
                int v = mult(a[i + j + len / 2], pw); // * pw[pos]
                a[i + j] = add(u, v);
                a[i + j + len / 2] = sub(u, v);
                // pos = (pos + diff) % LEN;
                pw = mult(pw, ml);
            }
        }
    }
}
```

```

    }
    if (inverse)
    {
        int m = binPow(sz(a), MOD - 2);
        FOR(i, 0, sz(a))
            // a[i] /= SZ(a);
            a[i] = mult(a[i], m);
    }
}

VI mult(VI a, VI b)
{
    int n = sz(a), m = sz(b);
    if (n == 0 || m == 0)
        return {};
    int sz = 1, szRes = n + m - 1;
    while(sz < szRes)
        sz *= 2;
    a.resize(sz);
    b.resize(sz);

    fft(a, false);
    fft(b, false);

    FOR(i, 0, sz)
        a[i] = mult(a[i], b[i]);

    fft(a, true);
    a.resize(szRes);
    return a;
}

```

Newton's method

Usable to find the solution of equation $F(Q) = 0$.

For example $F(Q) = x \cdot Q^2 + A - Q = 0$.

Newton's method approximates the solution of the equation using the formula:

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)}, \text{ where } F' = \frac{dF}{dQ}$$

Example of the derivative: $F'(Q) = 2 \cdot x \cdot Q - 1$.

Keep in mind that $|Q_k| = 2^k$.

FFT tricks

Two-dimensional FFT

The complexity is $\mathcal{O}(nm(\log n + \log m))$. The main problem is to resize the matrix. You must add non-empty vectors.

Divide-and-conquer FFT

Suppose we have the following DP relation:

$f(t) = g(t) - \sum_{0 \leq u < t} f(u)h(t-u)$, where $g(t)$ and $h(t)$ are known and we want to compute $f(t)$. We can apply divide-and-conquer FFT.

Let $m = \lfloor \frac{l+r}{2} \rfloor$. We guarantee the following invariant conditions.

By the time we compute the values for the segment $[l, r]$, the following conditions are already met:

- The values for $[0, l]$ on the DP is already determined.

- The sum of contributions from $[0, l)$ through $[l, r)$ is already applied to the DP in $[l, r)$.

When calculate the values for the segment $[l, r)$ do:

- Calculate the values for the segment $[l, m)$ recursively.
- Calculate the contributions from $[l, m)$ to $[m, r)$.
- Calculate the values for the segment $[m, r)$ recursively.

Properties of the discrete Fourier transform

$$DFT(x)_k = \sum_{n=0}^{N-1} x_n \cdot e^{-i2\pi \frac{k}{N} n}$$

Let $x_n^R = x_{N-n \bmod N}$.

$$DFT(x)^R = \overline{DFT(x)}$$

For real x , $DFT(x)^R = \overline{DFT(x)}$.

Interpolation

When x_0, x_1, \dots, x_d and y_0, y_1, \dots, y_d are given (where x_i are pairwise distinct), a polynomial $f(x)$ of degree no more than d such that $f(x_i) = y_i (i = 0, \dots, d)$ is uniquely determined.

Lagrange polynomial

Lagrange basis polynomial: $L_i(x) = \prod_{j \neq i} \frac{x - x_j}{x_i - x_j}$.

$$f(x) = y_0 L_0(x) + y_1 L_1(x) + \dots + y_d L_d(x).$$

Newton polynomial

Divided differences:

$$[y_i] = y_i$$

$$[y_i, y_{i+1}] = \frac{y_{i+1} - y_i}{x_{i+1} - x_i}$$

$$[y_i, \dots, y_j] = \frac{[y_{i+1}, \dots, y_j] - [y_i, \dots, y_{j-1}]}{x_j - x_i}$$

Newton basis polynomial: $N_i(x) = \prod_{j=0}^{i-1} (x - x_j)$.

$$f(x) = [y_0] N_0(x) + \dots + [y_0, y_1, \dots, y_d] N_d(x).$$

Linear recurrence

berlekamp-massey.hpp

Description: Finds a sequence of d integers c_1, \dots, c_d of the minimum length d such that $a_i = \sum_{j=1}^d c_j a_{i-j}$.

9979fe, 36 lines

VI berlekampMassey(const VI& a)

```

VI c = {1}, bp = {1};
int l = 0, b = 1, x = 1;
FOR(j, 0, sz(a))
{
    assert(sz(c) == l + 1);
    int d = a[j];
    FOR(i, 1, l + 1)
        updAdd(d, mult(c[i], a[j - i]));
    if (d == 0)
    {
        x++;
        continue;
    }
}
```

```

VI t = c;
int coef = mult(d, binPow(b, MOD - 2));
if (sz(bp) + x > sz(c))
    c.resize(sz(bp) + x);
FOR(i, 0, sz(bp))
    updSub(c[i + x], mult(coef, bp[i]));
if (2 * l > j)
{
    x++;
    continue;
}
l = j + 1 - l;
bp = t;
b = d;
x = 1;
}
c.erase(c.begin());
for (int& ci : c)
    ci = mult(ci, MOD - 1);
return c;
}

```

bostan-mori.hpp

Description: Computes the n -th term of a given linearly recurrent sequence $a_i = \sum_{j=1}^d c_j a_{i-j}$. The first d terms a_0, a_1, \dots, a_{d-1} are given.

The problem reduces to determining $[x^n]P(x)/Q(x)$.

$$\frac{P(x)}{Q(x)} = \frac{P(x)Q(-x)}{Q(x)Q(-x)} = \frac{U_e(x^2)}{V(x^2)} + x \cdot \frac{U_o(x^2)}{V(x^2)}.$$

Time: $\mathcal{O}(d \log d \log n)$.

e2a8cf, 25 lines

```

int bostanMori(const VI& c, VI a, ll n)
{
    int k = sz(c);
    assert(sz(a) == k);
    VI q(k + 1);
    q[0] = 1;
    FOR(i, 0, k)
        q[i + 1] = sub(0, c[i]);
    VI p = mult(a, q);
    p.resize(k);
    while (n)
    {
        VI qMinus = q;
        for (int i = 1; i <= k; i += 2)
            qMinus[i] = sub(0, qMinus[i]);
        VI newP = mult(p, qMinus);
        VI newQ = mult(q, qMinus);
        FOR(i, 0, k)
            p[i] = newP[2 * i + (n & 1)];
        FOR(i, 0, k + 1)
            q[i] = newQ[2 * i];
        n >>= 1;
    }
    return mult(p[0], binPow(q[0], MOD - 2));
}

```

Mathematical analysis and numerical methods**Taylor series**

$$f(x) = \sum_{k=0}^n \frac{f^{(k)}(x_0)}{k!} (x - x_0)^k + o((x - x_0)^n)$$

bostan-mori

$$\begin{aligned} e^x &= \sum_{n=0}^{\infty} \frac{x^n}{n!} & \ln(1+x) &= \sum_{n=1}^{\infty} (-1)^{n+1} \frac{x^n}{n} \\ \cos x &= \sum_{n=0}^{\infty} (-1)^n \frac{x^{2n}}{(2n)!} & \sin x &= \sum_{n=0}^{\infty} (-1)^n \frac{x^{2n+1}}{(2n+1)!} \end{aligned}$$

Green's theorem

$$\oint_C (L dx + M dy) = \iint_D \left(\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y} \right) dx dy$$

Runge-Kutta 4th Order

$$\begin{aligned} \frac{dy}{dx} &= f(x, y), y(0) = y_0, x_{i+1} - x_i = h \\ y_{i+1} &= y_i + \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4)h \end{aligned}$$

$$\begin{aligned} k_1 &= f(x_i, y_i) & k_2 &= f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_1 h) \\ k_3 &= f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_2 h) & k_4 &= f(x_i + h, y_i + k_3 h) \end{aligned}$$

List of integrals

$$\begin{aligned} \int \frac{dx}{a^2 + x^2} &= \frac{1}{a} \operatorname{arctg} \frac{x}{a} + C \\ \int \frac{dx}{a^2 - x^2} &= \frac{1}{2a} \ln \left| \frac{x+a}{x-a} \right| + C \\ \int \frac{dx}{\sqrt{a^2 - x^2}} &= \arcsin \frac{x}{a} + C \\ \int \frac{dx}{\sqrt{x^2 + a}} &= \ln \left| x + \sqrt{x^2 + a} \right| + C \\ \int \frac{dx}{\cos^2 x} &= \operatorname{tg} x + C \\ \int \frac{dx}{\sin^2 x} &= -\operatorname{ctg} x + C \end{aligned}$$

Simpson's rule

n — even number, $h = \frac{b-a}{n}$, $x_i = a + ih$

$$\int_a^b f(x) dx \approx \frac{h}{3} \left[f(x_0) + 4 \sum_{i=1}^{\frac{n}{2}} f(x_{2i-1}) + \sum_{i=1}^{\frac{n}{2}-1} f(x_{2i}) + f(x_n) \right]$$

Vandermonde matrix

$$V = V(x_0, x_1, \dots, x_m) = \begin{bmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & x_1^2 & \dots & x_1^n \\ 1 & x_2 & x_2^2 & \dots & x_2^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_m & x_m^2 & \dots & x_m^n \end{bmatrix}$$

$$V_{i,j} = x_i^j, \quad \det(V) = \prod_{0 \leq i < j \leq n} (x_j - x_i).$$

Hadamard matrix

$$H_1 = [1], \quad H_{2^k} = \begin{bmatrix} H_{2^{k-1}} & H_{2^{k-1}} \\ H_{2^{k-1}} & -H_{2^{k-1}} \end{bmatrix}$$

$$\det(H_n) = \pm n^{\frac{n}{2}}$$

For a matrix M such that $|M_{ij}| \leq 1$, holds $|\det(M)| \leq n^{n/2}$.

Number theory**Calculation of $a^b \pmod{m}$**

If $b \geq \phi(m)$, then value $a^b \equiv a^{[b \pmod{\phi(m)}] + \phi(m)} \pmod{m}$.

Generators

A generator exists only for $n = 1, 2, 4, p^k, 2p^k$ for odd primes p and positive integers k .

g is a generator modulo n if any number coprime with n can be represented as $[g^i \pmod{n}]$, $0 \leq i < \phi(n)$.

To find a generator:

- find $\phi(n)$ and p_1, \dots, p_m — the prime factors of $\phi(n)$
- g is generator only if $g^{\frac{\phi(n)}{p_j}} \not\equiv 1 \pmod{n}$ for each j
- check $g = 2, 3, 4, \dots, p-1$

Wilson's theorem

p is prime if and only if $(p-1)! \equiv (p-1) \pmod{p}$.

Quadratic residues

q is a quadratic residue modulo p if there exists an integer x such that $x^2 \equiv q \pmod{p}$. If p is odd prime then there exist $\frac{p+1}{2}$ residues (including 0).

Number theory functions

$$n = p_1^{\alpha_1} \cdots p_k^{\alpha_k}$$

$$\phi(n) = \prod p_i^{\alpha_i-1} (p_i - 1) \text{ — the number of coprimes}$$

$$F(n) = \frac{n \cdot \phi(n)}{2} \text{ — the sum of coprimes for } n > 1$$

$$\mu(n) = (-1)^k \text{ if } \max(\alpha_i) = 1, \text{ else } 0$$

$$\sigma_k(n) = \sum_{d|n} d^k$$

$$\sigma_0(n) = \prod (\alpha_i + 1)$$

$$\sigma_{k>0}(n) = \prod \frac{p_i^{(\alpha_i+1) \cdot k} - 1}{p_i^k - 1}$$

Möbius

$$g(n) = \sum_{d|n} f(d) \iff f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right)$$

$$M(n) = \sum_{k=1}^n \mu(k), \quad \sum_{d=1}^n M\left(\left\lfloor \frac{n}{d} \right\rfloor\right) = 1$$

$$\sum_{d|n} \phi(d) = n, \quad \sum_{d|n} \mu(d) = [n = 1]$$

Combinatorics

Binomials

$$\sum_{k=0}^n C_n^k = 2^n$$

$$\sum_{k=0}^m C_{n+k}^k = C_{n+m+1}^m$$

$$\sum_{m=0}^n C_m^k = C_{n+1}^{k+1}$$

$$\sum_{k=0}^n (C_n^k)^2 = C_{2n}^n$$

$$\sum_{j=0}^k C_m^j C_{n-m}^{k-j} = C_n^k$$

$$\sum_{j=0}^m C_m^j C_{n-m}^{k-j} = C_{n+1}^{k+1}$$

$$\sum_{k=0}^n C_{n-k}^k = F_{n+1}$$

Catalan numbers

$$C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} C_{2n}^n = C_{2n}^n - C_{2n}^{n-1}$$

1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786

Fibonacci numbers

$$F_1 = F_2 = 1$$

$$\gcd(F_m, F_n) = F_{\gcd(m,n)}$$

$$F_n = F_{n-1} + F_{n-2}$$

$$F_{n+1} F_{n-1} - F_n^2 = (-1)^n$$

$$F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$$

$$F_{47} \approx 2.9 \cdot 10^9$$

$$F_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n}{\sqrt{5}}$$

$$F_{88} \approx 1.1 \cdot 10^{18}$$

Stirling numbers of the second kind

$S(n, k)$ – the number of ways to divide n elements into k non-empty groups.

$$S(n, n) = 1, n \geq 0$$

$$S(n, 0) = 0, n > 0$$

$$S(n, k) = S(n-1, k-1) + S(n-1, k) \cdot k.$$

$$B_n = \sum_{k=0}^n S(n, k) \text{ from } n = 0:$$

1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597, 2764437, 190899322, 1382958545, 10480142147, 82864869804,...

Burnside's lemma

Let G be a finite group that acts on a set X .

The *orbit* of an element x in X is the set of elements in X to which x can be moved by the elements of G . The orbit of x is denoted by $G \cdot x$:

$$G \cdot x = \{g \cdot x \mid g \in G\}.$$

For each g in G , let X^g denote the set of elements in X that are fixed by g (also said to be left invariant by g), that is, $X^g = \{x \in X \mid g \cdot x = x\}$. Burnside's lemma asserts the following formula for the number of orbits, denoted $|X/G|$:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

Graphs

Prüfer sequence

At step i , remove the leaf with the smallest label and set the i -th element of the Prüfer sequence to be the label of this leaf's neighbour. The Prüfer sequence of a labeled tree is unique and has length $n-2$.

The number of spanning trees of K_n is n^{n-2} .

The number of spanning trees of $K_{L,R}$ number is $L^{R-1} \cdot R^{L-1}$.

Let $T_{n,k}$ be the number of labelled forests on n vertices with k connected components, such that vertices $1, \dots, k$ all belong to different components. $T_{n,k} = k \cdot n^{n-k-1}$.

The number of spanning trees in a complete graph K_n with the fixed degrees d_i is equal to: $\frac{(n-2)!}{\prod (d_i-1)}$

For a forest graph with connected components of sizes s_0, \dots, s_{k-1} , the number of ways to add edges to make a spanning tree is equal to: $n^{k-2} \cdot \prod s_i$

Chromatic polynomial

For a graph G , $\chi(G, \lambda) = \chi(\lambda)$ counts the number of its vertex λ -colorings. There is a unique polynomial $\chi(\lambda)$.

Deletion-contraction:

- The graph G/uv is obtained by merging u and v .
- The graph $G - uv$ is obtained by deleting the edge uv .
- $\chi(G, \lambda) = \chi(G - uv, \lambda) - \chi(G/uv, \lambda)$.

G is tree	$\chi(\lambda) = \lambda(\lambda-1)^{n-1}$
G is cycle C_n	$\chi(\lambda) = (\lambda-1)^n + (-1)^n(\lambda-1)$

Proposition. $\chi(\lambda)$ is equal to the number of pairs (σ, O) , where σ is any map $\sigma : V \rightarrow \{1, \dots, \lambda\}$ and O is an orientation of G , subject to the two conditions:

- The orientation O is acyclic.
- If $u \rightarrow v$ in O , then $\sigma(u) > \sigma(v)$.

Define $\bar{\chi}(\lambda)$ to be the number of pairs (σ, O) , where σ is any map $\sigma : V \rightarrow \{1, \dots, \lambda\}$ and O is an orientation of G , subject to the two conditions:

- The orientation O is acyclic.
- If $u \rightarrow v$ in O , then $\sigma(u) \geq \sigma(v)$.

Theorem. Suppose that $|V| = n$. Then for all non-negative integers λ holds:

$$\bar{\chi}(\lambda) = (-1)^n \chi(-\lambda)$$

Corollary. $(-1)^n \chi(G, -1)$ is equal to the number of acyclic orientations of G .

Kirchhoff's theorem

Let G be a finite graph, allowing multiple edges but not loops.

The laplacian matrix L of G is the $n \times n$ matrix whose (i, j) -entry L_{ij} is given by

$$L_{ij} = \begin{cases} -m_{ij}, & \text{if } i \neq j, m_{ij} \text{ edges between } v_i \text{ and } v_j, \\ \deg(v_i), & \text{if } i = j. \end{cases}$$

Let L_0 denote L with the i -th row and column removed for any i . Then for a connected graph, $\det(L_0)$ equals the number of spanning trees of G .

Karp's minimum mean-weight cycle algorithm

Let $G = (V, E)$ be a directed graph with weight function $w : E \rightarrow \mathbb{R}$, and let $n = |V|$. We define the **mean weight** of a cycle $c = \langle e_1, e_2, \dots, e_k \rangle$ of edges in E to be

$$\mu(c) = \frac{1}{k} \sum_{i=1}^k w(e_i).$$

Let $\mu^* = \min_c \mu(c)$, where c ranges over all directed cycles in G . We call a cycle c for which $\mu(c) = \mu^*$ a **minimum mean-weight cycle**.

Assume without loss of generality that every vertex $v \in V$ is reachable from a source vertex $s \in V$. Let $\delta_k(s, v)$ be the weight of a shortest path from s to v consisting of exactly k edges. If there is no path from s to v with exactly k edges, then $\delta_k(s, v) = \infty$.

$$\mu^* = \min_{v \in V} \max_{0 \leq k \leq n-1} \frac{\delta_n(s, v) - \delta_k(s, v)}{n - k}.$$

This can be computed in time $O(VE)$.

Erdős–Gallai theorem

A sequence of non-negative integers $d_1 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \dots + d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$ holds for every k in $1 \leq k \leq n$.

Planar graph properties

For a simple, connected, planar graph with v vertices, e edges and f faces, the following simple conditions hold for $v \geq 3$:

- Theorem 1. $e \leq 3 \cdot v - 6$.
- Theorem 2. If there are no cycles of length 3, then $e \leq 2 \cdot v - 4$.
- Theorem 3. $f \leq 2 \cdot v - 4$.
- Euler's formula. $v - e + f = 2$.
- Theorem 4. $3 \cdot f \leq 2 \cdot e$.
- Theorem 5. The dual graph is also planar.
- Theorem 6. There exists a vertex v with $\deg(v) \leq 5$.

Dilworth's theorem

A partially ordered set is a set S with a relation \leq on S satisfying:

1. $a \leq a$ for all $a \in S$ (reflexivity);
2. if $a \leq b$ and $b \leq a$, then $a = b$ (antisymmetry);
3. if $a \leq b$ and $b \leq c$, then $a \leq c$ (transitivity).

A chain is a subset of a set where each pair of distinct elements is comparable. An antichain is a subset of a set where every pair of elements is incomparable.

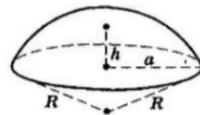
Dilworth's theorem states that, in any finite partially ordered set, the **largest antichain** has the same size as the **smallest chain decomposition**. Here, the size of the antichain is its number of elements, and the size of the chain decomposition is its number of chains.

Geometry

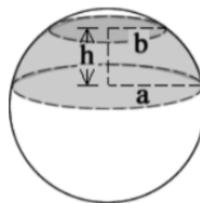
Trigonometry formulas

$$\begin{aligned}\sin(v+w) &= \sin v \cos w + \cos v \sin w \\ \sin(v-w) &= \sin v \cos w - \cos v \sin w \\ \tan(v+w) &= \frac{\tan v + \tan w}{1 - \tan v \tan w} \\ \sin v + \sin w &= 2 \sin \frac{v+w}{2} \cos \frac{v-w}{2} \\ \cos v + \cos w &= 2 \cos \frac{v+w}{2} \cos \frac{v-w}{2}\end{aligned}$$

Ball formulas



$$\begin{aligned}a &= \sqrt{h \cdot (2R - h)} \\ V &= \pi \cdot h^2 (R - \frac{h}{3})\end{aligned}$$



$$\begin{aligned}V &= \frac{1}{6} \pi h (3a^2 + 3b^2 + h^2) \\ R &= \sqrt{\frac{((a-b)^2 + h^2)((a+b)^2 + h^2)}{4h^2}}\end{aligned}$$

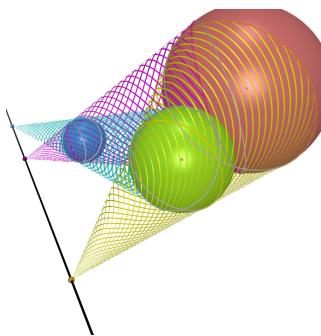
golden-section-search linear-basis

Triangle formulas

$$\begin{aligned}S &= \sqrt{p(p-a)(p-b)(p-c)} = \frac{abc}{4R} \\ m_a^2 &= \frac{2b^2 + 2c^2 - a^2}{4} \text{ (median)} \\ w_a^2 &= \frac{bc((b+c)^2 - a^2)}{(b+c)^2} \text{ (bisector)} \\ \frac{a}{\sin A} &= \frac{b}{\sin B} = \frac{c}{\sin C} = 2R \\ a^2 &= b^2 + c^2 - 2bc \cos A\end{aligned}$$

Monge's theorem

There are three circles(balls) of different radii, for each pair of circles find the point of intersection of the external tangents. All three obtained points **lie on a line**. The point from the pair of the largest and the smallest **lies between** the other two.



Pick's theorem

Suppose that a polygon has integer coordinates for all of its vertices. Let i be the number of integer points inside, and let b be the number of integer points on boundary. Then the area $S = i + \frac{b}{2} - 1$.

Ptolemy's theorem

For a general quadrilateral $ABCD$ holds:
 $AB \cdot CD + AD \cdot BC \geq AC \cdot BD$.

Equality holds if and only if the quadrilateral is cyclic.

Euler line

For a general triangle, the orthocenter H , the centroid G , and the circumcenter O , in this order, lie on the same line (Euler line) and $\frac{|HG|}{|GO|} = \frac{2}{1}$.

Fermat point

In a given triangle $\triangle ABC$ the Fermat point is the point X , which minimizes the sum of distances from A , B , and C , $|AX| + |BX| + |CX|$.

If all angles of the triangle are less than 120° , the the Fermat point is the interior point X from which each side subtends an angle of 120° , i.e., $\angle BXC = \angle CXA = \angle AXB = 120^\circ$.

If any angle of the triangle formed by those points is 120° or more, then the Fermat point is the vertex of that angle.

Various (7)

golden-section-search.hpp

4c0990, 27 lines

```
db goldenSectionSearch(db l, db r)
{
    const db c = (-1 + sqrt(5)) / 2;
    const int M = 474;
    db m1 = r - c * (r - l), fm1 = f(m1),
        m2 = l + c * (r - l), fm2 = f(m2);
    FOR(i, 0, M)
    {
        if (fm1 < fm2)
        {
            r = m2;
            m2 = m1;
            fm2 = fm1;
            m1 = r - c * (r - l);
            fm1 = f(m1);
        }
        else
        {
            l = m1;
            m1 = m2;
            fm1 = fm2;
            m2 = l + c * (r - l);
            fm2 = f(m2);
        }
    }
    return (l + r) / 2;
}
```

linear-basis.hpp

2ff3b8, 45 lines

```
const int MAX_BITS = 64;

struct LinearBasis
{
    bitset<MAX_BITS> basis[MAX_BITS];
    int size;

    LinearBasis()
    {
        size = 0;
    }

    void insert(bitset<MAX_BITS> x)
    {
        for (int i = MAX_BITS - 1; i >= 0; --i)
        {
            if (!x[i]) continue;
            if (basis[i].none())
            {
                basis[i] = x;
                ++size;
                return;
            }
            x ^= basis[i];
        }
    }

    bool canRepresent(bitset<MAX_BITS> x)
    {
        RFOR(i, MAX_BITS, 0)
            if (x[i]) x ^= basis[i];
        return x.none();
    }

    bitset<MAX_BITS> getMaxXOR()
    {
```

```
bitset<MAX_BITS> res;
RFOR(i,MAX_BITS,0)
{
    if ((res ^ basis[i]).to_ullong() > res.to_ullong())
        res ^= basis[i];
}
return res;
};
```