

12. Extending Event-B Models

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- ▶ Understanding **abstraction** and **refinement**
- ▶ **Extension refinement** by an example

Readings

- ▶ **Hoang [2013]** (in particular section on superposition refinement)

Abstraction and Refinement
On Event-B Refinement
A Secure Database - Requirements
Abstract Level
Concrete Level

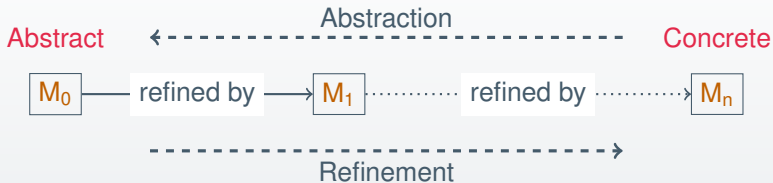
Summary

Abstraction

Abstraction can be viewed as a process of **simplifying our understanding** of a system.

- ▶ The simplification should
 - ▶ focus on the **intended purpose** of the system
 - ▶ **ignore details of how** that purpose is achieved.
- ▶ The modeller needs to make **judgements** about what they believe to be the **key features** of the system.

On Refinement (1/3)



- **Refinement** is a process of enriching or modifying a model to
 - **augment** the functionality being modelled, or
 - **explain how** some purpose is achieved

On Refinement (2/3)



- ▶ M_1 is a refinement of M_0
- ▶ M_0 is an abstraction of M_1

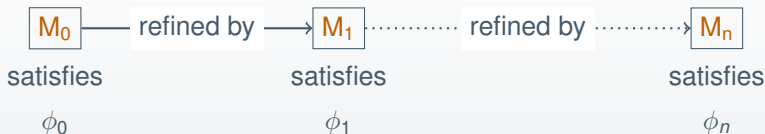
Facilitates abstraction

We can **postpone** treatment of some system features to later refinement steps

Coping with System Complexity

Abstraction and refinement together should allow us to **manage system complexity** in the design process

On Refinement (3/3)

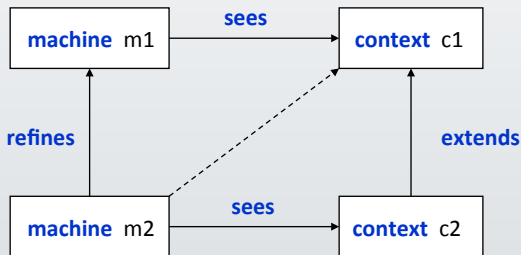
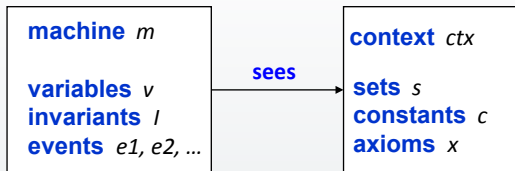


Properties Preservation

Properties are **preserved** during refinement

- ▶ Preserve **safety** (e.g., invariants) properties.
- ▶ We use proofs to **verify** the consistency of a refinement step
- ▶ Failing proofs help **identify inconsistencies** in a refinement step

Modelling Components and Refinement



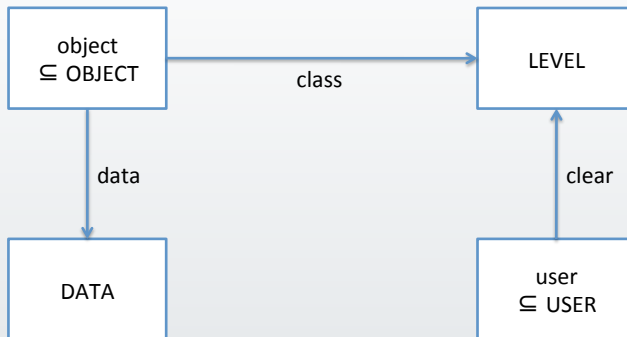
Requirements

- REQ 1 A **new user** can be **added** with a **clearance level** (1 .. 10)
- REQ 2 An existing user can add a **new object**
(hence becomes the owner of that object).
- REQ 3 An **existing object** can be **removed** by its owner.
- REQ 4 An **existing user** owning no objects can be **removed**.
- REQ 5 Each **object** is classified by some **access level** (1 .. 10)
- REQ 6 An object's **classification** is **at most** its owner's **clearance** level.
- REQ 7 Each **object** is associated with some **data**
- REQ 8 An object can be **accessed (read/written)** by any user
whose clearance level at least the object's classification
- REQ 9 The user **clearance** level can be **adjusted**.
- REQ 10 The object **classification** level can be **adjusted**.

- ▶ **SecureDB1**: We focus on
 - ▶ objects and their classification, and
 - ▶ users and their clearance levels.
- ▶ **SecureDB2**: We focus on **objects' ownership**.

Abstract Level SecureDB1 (Recall)

Class diagram for the Secure Database



Abstract Level SecureDB1 (Recall)

sets : *OBJECT DATA USER*

constants : *LEVEL*

axioms :
 $LEVEL = 1..10$

variables : *object, user, data, class, clear*

invariants :
 $object \subseteq OBJECT$
 $user \subseteq USER$
 $data \in object \rightarrow DATA // REQ 7$
 $class \in object \rightarrow LEVEL // REQ 5$
 $clear \in user \rightarrow LEVEL // REQ 1$

Abstract Level SecureDB1 (Recall)

Events for Secure Database

- ▶ AddUser: REQ 1
- ▶ RemoveUser: REQ 4
- ▶ AddObject: REQ 2
- ▶ RemoveObject: REQ 3
- ▶ Read: REQ 8
- ▶ Write: REQ 8
- ▶ ChangeClearance: REQ 9
- ▶ ChangeClass: REQ 10

Extend the database specification with object ownerships

REQ 3 An existing object can be removed by its owner.

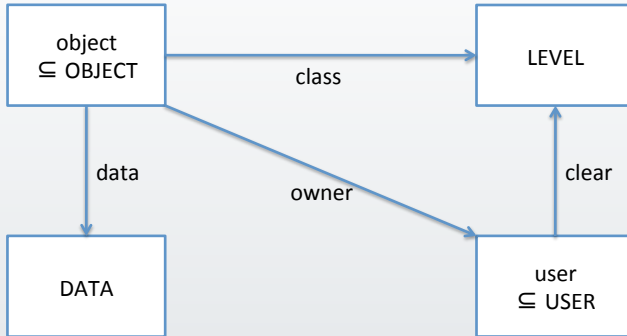
REQ 6 An object's classification is at most its owner's clearance level.

Design consideration

- ▶ What additional variables are required?
- ▶ What events are affected?
 - ▶ Existing events
 - ▶ New events

Concrete Level SecureDB2

Class diagram with ownership



variables : *object, user, data, class, clear, owner*

invariants :

owner \in *object* \rightarrow *user*

$\forall o \cdot o \in \textit{object} \Rightarrow \textit{class}(o) \leq \textit{clear}(\textit{owner}(o))$ // **REQ 6**

Important

- ▶ For extension refinement, we must list **all the variables**:
 - ▶ those from **SecureDB1** that we wish to retain,
 - ▶ new variables, e.g., *owner*
- ▶ We do not repeat invariants of **SecureDB1** in **SecureDB2**

```
AddUser
  any  $u, c$  where
     $u \in USER$ 
     $u \notin user$ 
     $c \in LEVEL$ 
  then
     $user := user \cup \{u\}$ 
     $clear(u) := c$ 
  end
```

Do we need to modify event AddUser?

```
AddObject
  any   $o, d, c$   where
     $o \in OBJECT$ 
     $o \notin object$ 
     $d \in DATA$ 
     $c \in LEVEL$ 
  then
     $object := object \cup \{o\}$ 
     $data(o) := d$ 
     $class(o) := c$ 
  end
```

Do we need to modify event AddObject?

Event AddObject

Event Extension

```
AddObject extended  
refines AddObject  
any  $u$  where  
   $u \in \text{user}$   
   $\text{clear}(u) \geq c$   
then  
   $\text{owner}(o) := u$   
end
```

```
AddObject  
refines AddObject  
any  $o, d, c, u$  where  
   $o \in \text{OBJECT}$   
   $o \notin \text{object}$   
   $d \in \text{DATA}$   
   $c \in \text{LEVEL}$   
   $u \in \text{user}$   
   $\text{clear}(u) \geq c$   
then  
   $\text{object} := \text{object} \cup \{o\}$   
   $\text{data}(o) := d$   
   $\text{class}(o) := c$   
   $\text{owner}(o) := u$   
end
```

Other events to consider

- ▶ Read
- ▶ Write
- ▶ ChangeClass
- ▶ ChangeClearance
- ▶ RemoveUser
- ▶ RemoveObject

Another question

Do we need new events?

(Alternatively) What additional functionality the system should have?

Concluding

- ▶ Abstraction vs. Refinement
- ▶ Event-B Refinement
 - ▶ Extension Refinement

Thai Son Hoang. An introduction to the Event-B modelling method. In
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Springer-Verlag, July 2013. [http://www.springer.com/
computer/swe/book/978-3-642-33169-5](http://www.springer.com/computer/swe/book/978-3-642-33169-5).