

**Solution to Problem 1.**

(a)

- i A vector  $\mathbf{v} \in \mathbb{R}^n$  is an eigenvector of  $\mathbf{M}$  if  $\mathbf{M}\mathbf{v} = \lambda\mathbf{v}$  where  $\lambda \in \mathbb{R}$  is the corresponding eigenvalue. [1]
- ii The  $i^{\text{th}}$  eigenvector and eigenvalue of  $\mathbf{B} = \mathbf{A}\mathbf{A}^T$  are given by  $\mathbf{u}_i$  and  $\sigma_i^2$  respectively. This can be justified by

$$\begin{aligned} \mathbf{B}\mathbf{u}_i &= \mathbf{A}\mathbf{A}^T\mathbf{u}_i = \mathbf{U}\Sigma\mathbf{V}^T\mathbf{V}\Sigma^T\mathbf{U}^T\mathbf{u}_i \\ &= \mathbf{U}\Sigma\Sigma^T\mathbf{U}^T\mathbf{u}_i = \mathbf{U}\Sigma\Sigma^T\mathbf{e}_i \\ &= \sigma_i^2\mathbf{U}\mathbf{e}_i = \sigma_i^2\mathbf{u}_i, \end{aligned}$$

where  $\mathbf{e}_i$  is the standard basis vector. [2]

- iii Let  $r$  be the rank of  $\mathbf{A}$ . Define  $\mathbf{U}_r = [\mathbf{u}_1, \dots, \mathbf{u}_r]$ ,  $\Sigma_r = \text{diag}([\sigma_1, \dots, \sigma_r])$ , and  $\mathbf{V}_r = [\mathbf{v}_1, \dots, \mathbf{v}_r]$ . Then  $\mathbf{A}^\dagger = \mathbf{V}_r\Sigma_r^{-1}\mathbf{U}_r^T$ . [1]
- iv  $\text{proj}(\mathbf{x}, \mathbf{A}) = \mathbf{A}\mathbf{A}^\dagger\mathbf{x}$ . [1]
- v The orthogonality can be verified as

$$\begin{aligned} \mathbf{A}^T\mathbf{x}_r &= \mathbf{A}^T\mathbf{x} - \mathbf{A}^T(\mathbf{A}\mathbf{A}^\dagger)\mathbf{x} = (\mathbf{A}^T - \mathbf{A}^T(\mathbf{A}\mathbf{A}^\dagger))\mathbf{x} \\ &= (\mathbf{V}_r\Sigma_r\mathbf{U}_r^T - \mathbf{V}_r\Sigma_r\mathbf{U}_r^T(\mathbf{U}_r\mathbf{U}_r^T))\mathbf{x} \\ &= (\mathbf{V}_r\Sigma_r\mathbf{U}_r^T - \mathbf{V}_r\Sigma_r\mathbf{U}_r^T)\mathbf{x} = \mathbf{0}. \end{aligned}$$

[2]

vi Note that

$$\begin{aligned} \|\mathbf{x} - \mathbf{v}\|_2^2 &= \|\mathbf{x} - \mathbf{x}_p + \mathbf{x}_p - \mathbf{v}\|_2^2 = \|\mathbf{x}_r + \mathbf{x}_p - \mathbf{v}\|_2^2 \\ &= \|\mathbf{x}_r\|_2^2 + 2\mathbf{x}_r^T(\mathbf{x}_p - \mathbf{v}) + \|\mathbf{x}_p - \mathbf{v}\|_2^2 \\ &= \|\mathbf{x}_r\|_2^2 + \|\mathbf{x}_p - \mathbf{v}\|_2^2, \end{aligned}$$

where the last equality comes from the orthogonality between  $\mathbf{x}_r$  and  $\mathbf{x}_p - \mathbf{v}$ . Since  $\|\mathbf{x}_p - \mathbf{v}\|_2^2 \geq 0$ , it is clear that  $\|\mathbf{x} - \mathbf{v}\|_2 \geq \|\mathbf{x}_r\|_2$ . [3]

(b)

i

$$\begin{aligned}\operatorname{tr}(AB) &= \sum_i (AB)_{i,i} = \sum_i \sum_j A_{i,j} B_{j,i} \\ &= \sum_j \sum_i B_{j,i} A_{i,j} = \sum_j (BA)_{j,j} = \operatorname{tr}(BA).\end{aligned}$$

[2]

ii

$$\begin{aligned}\|A\|_F^2 &= \sum_{i,j} A_{i,j}^2 = \operatorname{tr}(A^T A) = \operatorname{tr}(V \Sigma^2 V^T) \\ &= \operatorname{tr}(\Sigma^2 V^T V) = \operatorname{tr}(\Sigma^2) = \sum_i \sigma_i^2.\end{aligned}$$

[2]

iii  $\|A\|_2 = \sigma_{\max} = \sigma_1$ . This can be proved as follows.

$$\begin{aligned}\|Ax\|_2^2 &= x^T A^T A x = x^T V \Sigma^2 V^T x \\ &= \sum_i \sigma_i^2 (x^T v_i)^2 \\ &\leq \sigma_1^2 \sum_i (x^T v_i)^2 = \sigma_1^2 x^T V V^T x \\ &= \sigma_1^2 x^T x = \sigma_1^2.\end{aligned}$$

[3]

iv Since  $A = A^T$  and  $A \geq 0$ , the singular value decomposition of  $A$  can be written as  $A = U \Sigma U^T$ . Then

$$\begin{aligned}\operatorname{tr}(A) &= \operatorname{tr}(U \Sigma U^T) = \operatorname{tr}(\Sigma U^T U) = \operatorname{tr}(\Sigma) \\ &= \sum_{i=1}^{\min(m,n)} \sigma_i = \|A\|_*.\end{aligned}$$

[3]

**Solution to Problem 2.**

(a)

i

A. The soft-thresholding function is of the form

$$x^* = \eta(z; \lambda) = \begin{cases} z - \lambda & \text{if } z \geq \lambda, \\ 0 & \text{if } -\lambda < z < \lambda, \\ z + \lambda & \text{if } z \leq -\lambda. \end{cases} \quad [1]$$

B. The IST algorithm is an iterative algorithm where in the  $k^{\text{th}}$  iteration the variable  $\mathbf{x}^k$  is updated by

$$\mathbf{x}^k = \eta(\mathbf{x}^{k-1} + t_k \mathbf{A}^T (\mathbf{y} - \mathbf{A} \mathbf{x}^{k-1}); \lambda t_k),$$

and  $t_k > 0$  is an appropriately chosen step size. [1]

ii

A. Given an input vector  $\mathbf{z}$ , the hard thresholding function  $H_S(\mathbf{z})$  sets all but the largest (in magnitude)  $S$  elements of  $\mathbf{z}$  to zero. It is designed to solve the non-convex optimisation problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x} - \mathbf{z}\|_2^2 \text{ subject to } \|\mathbf{x}\|_0 \leq S. \quad [2]$$

B. The IHT algorithm is an iterative algorithm where in the  $k^{\text{th}}$  iteration the variable  $\mathbf{x}^k$  is updated by

$$\mathbf{x}^k = H_S(\mathbf{x}^{k-1} + t_k \mathbf{A}^T (\mathbf{y} - \mathbf{A} \mathbf{x}^{k-1})),$$

and  $t_k > 0$  is an appropriately chosen step size. It is designed to solve the non-convex optimisation problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{y} - \mathbf{A} \mathbf{x}\|_2^2 \text{ subject to } \|\mathbf{x}\|_0 \leq S. \quad [2]$$

(b)

i Define the soft thresholding function  $\eta_\sigma(\mathbf{Z}; \lambda)$  as

$$\mathbf{X} = \eta_\sigma(\mathbf{Z}; \lambda) = \sum_i \mathbf{u}_i \eta(\sigma_i; \lambda) \mathbf{v}_i^T,$$

where  $\sigma_i$  is the  $i^{th}$  singular value of  $\mathbf{Z}$ ,  $\mathbf{u}_i$  and  $\mathbf{v}_i$  are the corresponding singular vectors. The IST algorithm to solve the low-rank matrix recovery problem is an iterative algorithm where in the  $k^{th}$  iteration the matrix  $\mathbf{X}^k$  is updated by

$$\mathbf{X}^k = \eta_\sigma(\mathbf{X}^{k-1} + t_k \mathcal{A}^*(\mathbf{y} - \mathcal{A}(\mathbf{X}^{k-1})); \lambda t_k),$$

and  $t_k > 0$  is an appropriately chosen step size. [2]

ii Define the hard thresholding function  $H_r(\mathbf{Z}; \lambda)$  as

$$\mathbf{X} = H_r(\mathbf{Z}) = \mathbf{U}_r \text{diag}([\sigma_1, \dots, \sigma_r]) \mathbf{V}_r^T,$$

where  $\sigma_i$  is the  $i^{th}$  singular value of  $\mathbf{Z}$ ,  $\mathbf{U}_r = [\mathbf{u}_1, \dots, \mathbf{u}_r]$  and  $\mathbf{V}_r = [\mathbf{v}_1, \dots, \mathbf{v}_r]$ . The IHT algorithm to solve the low-rank matrix recovery problem is an iterative algorithm where in the  $k^{th}$  iteration the matrix  $\mathbf{X}^k$  is updated by

$$\mathbf{X}^k = H_r(\mathbf{X}^{k-1} + t_k \mathcal{A}^*(\mathbf{y} - \mathcal{A}(\mathbf{X}^{k-1}))),$$

and  $t_k > 0$  is an appropriately chosen step size. [2]

iii

A. Since  $\mathbf{X}$  is of rank 1,  $\mathbf{X} = \mathbf{u}\mathbf{v}^T$  and  $\mathbf{u}$  must be of the form  $\mathbf{u} = [1, 1, -2]^T$ . Hence  $\mathbf{v} = [1, -1, -2]^T$  and

$$\mathbf{X} = \begin{bmatrix} 1 & -1 & -2 \\ 1 & -1 & -2 \\ -2 & 2 & 4 \end{bmatrix}.$$

[3]

B.

$$\mathbf{A} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}.$$

$$\mathbf{y} = [1, 1, -1, 2, 4]^T, \text{ and}$$

$$\mathcal{A}^*(\mathbf{y}) = \begin{bmatrix} 1 & 0 & 0 \\ 1 & -1 & 0 \\ 0 & 2 & 4 \end{bmatrix}.$$

[3]

(c) Write

$$\mathbf{y} = \mathcal{A} \left( \begin{bmatrix} h(1) \\ h(2) \\ h(3) \end{bmatrix} \begin{bmatrix} x(1) \\ x(2) \\ x(3) \end{bmatrix}^T \right) = \mathcal{A} \left( \begin{bmatrix} h(1)x(1) & h(1)x(2) & h(1)x(3) \\ h(2)x(1) & h(2)x(2) & h(2)x(3) \\ h(3)x(1) & h(3)x(2) & h(3)x(3) \end{bmatrix} \right).$$

Clearly the matrix presentation of  $\mathcal{A}$  is given by

$$\mathbf{A} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}.$$

Possible convex optimisation formulations for blind deconvolution include

$$\min_{\mathbf{X}} \frac{1}{2} \|\mathbf{y} - \mathcal{A}(\mathbf{X})\|_2^2 + \lambda \|\mathbf{X}\|_*,$$

and

$$\min_{\mathbf{X}} \|\mathbf{X}\|_* \text{ subject to } \mathbf{y} = \mathcal{A}(\mathbf{X}).$$

[4]

### Solution to Problem 3.

(a)

i A set  $\mathcal{S} \subset \mathbb{R}^n$  is convex if

$$\lambda \mathbf{x} + (1 - \lambda) \mathbf{y} \in \mathcal{S},$$

for all  $\mathbf{x}, \mathbf{y} \in \mathcal{S}$  and all  $\lambda \in [0, 1]$ .

[2]

ii A function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is convex if

$$f(\lambda \mathbf{x} + (1 - \lambda) \mathbf{y}) \leq \lambda f(\mathbf{x}) + (1 - \lambda) f(\mathbf{y})$$

for all  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^n$  and all  $\lambda \in [0, 1]$ .

[2]

iii The standard form of a convex optimisation problem is

$$\begin{aligned} \min_{\mathbf{x}} f(\mathbf{x}) \\ \text{subject to } h_i(\mathbf{x}) &\leq 0, \quad i \in \{1, \dots, m\} \\ \ell_j(\mathbf{x}) &= 0, \quad j \in \{1, \dots, p\} \end{aligned}$$

where  $f(\mathbf{x})$ ,  $h_i(\mathbf{x})$  are convex and  $\ell_j$  is affine.

[2]

iv The corresponding Lagrangian is given by

$$L(\mathbf{x}, \mathbf{u}) = \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x} + \mathbf{u}^T \mathbf{A} \mathbf{x}.$$

Minimise the Lagrangian with respect to  $\mathbf{x}$  gives

$$\mathbf{Q} \mathbf{x} + \mathbf{c} + \mathbf{A}^T \mathbf{u} = \mathbf{0}.$$

Combine this with the constraints  $\mathbf{A} \mathbf{x} = \mathbf{0}$ . One has

$$\begin{bmatrix} \mathbf{Q} & \mathbf{A}^T \\ \mathbf{A} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \mathbf{u} \end{bmatrix} = \begin{bmatrix} -\mathbf{c} \\ \mathbf{0} \end{bmatrix}.$$

Hence the optimal  $\mathbf{x}$  can be computed by

$$\begin{bmatrix} \mathbf{x} \\ \mathbf{u} \end{bmatrix} = \begin{bmatrix} \mathbf{Q} & \mathbf{A}^T \\ \mathbf{A} & \mathbf{0} \end{bmatrix}^{\dagger} \begin{bmatrix} -\mathbf{c} \\ \mathbf{0} \end{bmatrix}$$

[3]

(b)

i For any  $\mathbf{x}, \mathbf{y} \in \mathcal{C}_\alpha$ , it holds that  $\forall \lambda \in [0, 1]$ ,

$$\begin{aligned} f(\lambda \mathbf{x} + (1 - \lambda) \mathbf{y}) &\leq \lambda f(\mathbf{x}) + (1 - \lambda) f(\mathbf{y}) \\ &\leq \lambda \alpha + (1 - \lambda) \alpha = \alpha. \end{aligned}$$

Hence  $\lambda \mathbf{x} + (1 - \lambda) \mathbf{y} \in \mathcal{C}_\alpha$ . [2]

ii The sublevel set  $\mathcal{C}_0 = \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{x}\|_2 = \sqrt{y}\}$  which gives the shell of the ball with radius  $\sqrt{y}$ . Pick arbitrary two points on the shell. The line segment between these two points is not on the shell (in the ball defined by the shell). Hence the sublevel set  $\mathcal{C}_0$  is not convex and  $f$  is not convex. [2]

(c)

i Note that  $(\mathbf{a}_i^T \mathbf{x})^2 = (\mathbf{a}_i^T \mathbf{x})(\mathbf{x}^T \mathbf{a}_i) = \mathbf{a}_i^T \mathbf{x} (\mathbf{x}^T \mathbf{a}_i) = \text{tr}(\mathbf{x} (\mathbf{x}^T \mathbf{a}_i) \mathbf{a}_i^T) = \text{tr}(\mathbf{X} \mathbf{A}_i)$ . One has  $f(\mathbf{x}) = \sum_i (y_i - (\mathbf{a}_i^T \mathbf{x})^2)^2 = \sum_i (y_i - \text{tr}(\mathbf{X} \mathbf{A}_i))^2$ . [2]

ii

A. It is sufficient to show that the objective function is convex and the constraint set is convex.

The first term in the objective function is a quadratic function of  $\mathbf{X}$  and hence convex. In fact, it can be written as  $\sum_i (y_i - \mathcal{A}(\mathbf{X}))^2$  where  $\mathcal{A}$  is a linear operator. The second term in the objective function is a linear function of  $\mathbf{X}$  and hence convex.

The first constraint is essentially to say  $x_{i,j} - x_{j,i} = 0$ . It is an equality constraint involving linear functions of  $\mathbf{X}$  and hence defines a convex set.

The only non-trivial part is to verify that the constraint  $\mathbf{X} \geq 0$  gives a convex set. To show that, note that  $\mathbf{X} \geq 0$  if and only if  $\mathbf{v}^T \mathbf{X} \mathbf{v} \geq 0$  for all  $\mathbf{v}$ . Let  $\mathbf{X}_1 \geq 0$  and  $\mathbf{X}_2 \geq 0$ . Then for all  $\lambda \in [0, 1]$ ,

$$\begin{aligned} \mathbf{v}^T (\lambda \mathbf{X}_1 + (1 - \lambda) \mathbf{X}_2) \mathbf{v} &= \lambda \mathbf{v}^T \mathbf{X}_1 \mathbf{v} + (1 - \lambda) \mathbf{v}^T \mathbf{X}_2 \mathbf{v} \\ &\geq 0. \end{aligned}$$

Hence the constraint  $\mathbf{X} \geq 0$  gives a convex set. [3]

B. The goal of using the optimisation problem (3.4) is to find an  $\mathbf{X}$

such that  $\mathbf{X} = \mathbf{x}\mathbf{x}^T$ . The first term in the objective function is for data consistency. The second term in the objective function is to promote low-rank matrices as solutions: the matrix  $\mathbf{x}\mathbf{x}^T$  has rank 1 and  $\text{tr}(\mathbf{X}) = \|\mathbf{X}\|_*$ , under the constraints given in (3.4). The two constraints are motivated by the fact that the matrix  $\mathbf{x}\mathbf{x}^T$  is symmetric and non-negative definite. [2]



Solution to Problem 4.

(a) The mutual coherence is defined as

$$\mu(A) = \max_{i \neq j} |\langle \mathbf{a}_i, \mathbf{a}_j \rangle| = \max_{i \neq j} |\mathbf{a}_i^T \mathbf{a}_j|. \quad [1]$$

(b)

i A matrix  $A \in \mathbb{R}^{m \times n}$  is said to satisfy the RIP with parameters  $(K, \delta)$ , if for all  $\mathcal{T} \subset \{1, \dots, n\}$  such that  $|\mathcal{T}| \leq K$  and for all  $\mathbf{q} \in \mathbb{R}^{|\mathcal{T}|}$ , it holds that

$$(1 - \delta) \|\mathbf{q}\|_2^2 \leq \|A_{\mathcal{T}} \mathbf{q}\|_2^2 \leq (1 + \delta) \|\mathbf{q}\|_2^2.$$

The RIC  $\delta_K$  is defined as the smallest constant  $\delta$  for which the  $(K, \delta)$ -RIP holds, i.e.,

$$\delta_K = \inf \left\{ \delta : (1 - \delta) \|\mathbf{q}\|_2^2 \leq \|A_{\mathcal{T}} \mathbf{q}\|_2^2 \leq (1 + \delta) \|\mathbf{q}\|_2^2 \right. \\ \left. \forall |\mathcal{T}| \leq K, \forall \mathbf{q} \in \mathbb{R}^{|\mathcal{T}|} \right\}. \quad [3]$$

ii It holds that

$$|\langle A_{\mathcal{I}} \mathbf{a}, A_{\mathcal{J}} \mathbf{b} \rangle| \leq \delta_{k+\ell} \|\mathbf{a}\|_2 \|\mathbf{b}\|_2,$$

that is,  $c = \delta_{k+\ell}$ . [1]

iii Suppose that there exists another  $S$ -sparse vector  $\mathbf{x}' \neq \mathbf{x}_0$  such that  $\mathbf{y} = A\mathbf{x}'$ . Then

$$A(\mathbf{x}_0 - \mathbf{x}') = \mathbf{y} - \mathbf{y} = \mathbf{0}.$$

Note that  $\mathbf{x}_0 - \mathbf{x}'$  has sparsity level at most  $2S$ . By the definition of RIP,

$$\|A(\mathbf{x}_0 - \mathbf{x}')\|_2 \geq (1 - \delta_{2S}) \|\mathbf{x}_0 - \mathbf{x}'\|_2 > 0.$$

We have a contradiction. The  $S$ -sparse solution of  $\mathbf{y} = A\mathbf{x}$  must be unique. [3]

(c) The diagonal elements of  $M$  are 1 and off-diagonal elements are bounded by  $\mu$  by the definition of mutual coherence constant, i.e.,  $|M_{i,j}| < \mu$  for  $i \neq j$ .

By the Gershgorin circle theorem,

$$\lambda(M) \in \left[ 1 - \sum_{j \neq i} |M_{i,j}|, 1 + \sum_{j \neq i} |M_{i,j}| \right] \subset [1 - S\mu, 1 + S\mu].$$

Hence,

$$\delta_S \leq S\mu.$$

[3]

(d)

i Suppose that  $i \notin \mathcal{T}$ .

$$\begin{aligned} |\mathbf{a}_i^T \mathbf{y}| &= |\mathbf{a}_i^T \mathbf{A} \mathbf{x}| = |\mathbf{a}_i^T \mathbf{A}_{\mathcal{T}} \mathbf{x}_{0,\mathcal{T}}| \\ &\leq \delta_{S+1} \|\mathbf{x}_{0,\mathcal{T}}\|_2 = \delta_{S+1} \|\mathbf{x}_0\|_2. \end{aligned}$$

[3]

ii Suppose that  $i \in \mathcal{T}$ .

$$\begin{aligned} |\mathbf{a}_i^T \mathbf{y}| &= \left| \mathbf{a}_i^T \mathbf{a}_i x_{0,i} + \mathbf{a}_i^T \sum_j \mathbf{a}_j x_{0,j} \right| \\ &\geq |x_{0,i}| - |\mathbf{a}_i^T \mathbf{A}_{\mathcal{T} \setminus \{i\}} \mathbf{x}_{0,\mathcal{T} \setminus \{i\}}| \\ &\geq |x_{0,i}| - \delta_{S+1} \|\mathbf{x}_{0,\mathcal{T} \setminus \{i\}}\|_2 \\ &\geq |x_{0,i}| - \delta_{S+1} \|\mathbf{x}_0\|_2. \end{aligned}$$

At the same time,

$$\max_{i \in \mathcal{T}} |x_{0,i}| \geq \frac{1}{\sqrt{S}} \|\mathbf{x}_0\|.$$

Hence

$$\max_{i \in \mathcal{T}} |\mathbf{a}_i^T \mathbf{y}| \geq \left( \frac{1}{\sqrt{S}} - \delta_{S+1} \right) \|\mathbf{x}_0\|_2.$$

[3]

iii To guarantee that  $i^* \in \mathcal{T}$ , one needs

$$\frac{1}{\sqrt{S}} - \delta_{S+1} > \delta_{S+1}.$$

Or equivalently

$$\delta_{S+1} \leq \frac{1}{2\sqrt{S}}.$$

[3]