

# The Art of Optimization



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# Preface

This book is an introduction to the *art of optimization*. I have deliberately used the word *art* in the title and placed an artwork in the title page to express that optimization requires imagination, skills and vision. The artwork has been painted by my wife, Elisabetta, and illustrates an optimization problem: one of those problems that we solve every day without even realizing their “technical” nature. I leave it to the reader to guess what the actual problem is. However, if you have attended my Optimization Lectures on a Friday morning in the past 15 years, you have probably solved the problem every week.

This book is the result of the lectures I have given at Imperial College London for Undergraduate, Master and Ph.D. students of all engineering departments (and also of the Mathematics and Physics Departments). I am not an expert in optimization, in the sense that my research activity has only seldom touched upon optimization problems, but I do believe that understanding optimization is essential for all engineers, practitioners and for everyday life. My research work is focused on systems and control: this is why some of the exercises contain a *systems and control* perspective of optimization problems. It is not hard to see that notions such as stationary points and equilibria, convergence and stability, speed of convergence and convergence rate (the former from optimization, the latter from systems and control) are fundamentally identical and one could borrow ideas and tools from systems and control theory to understand optimization problems and design optimization algorithms. Whenever possible, and in particular in the exercises, I have made this connection. Clearly, there are much deeper connections and relations which I have not discussed.

I conclude the preface with two observations. The first one is that the most difficult step in the art of optimization is the formulation of the problem. Only problems which are carefully formulated and in which all physical and engineering insight is captured yield underlying optimization problems for which one can attempt to find a meaningful solution. The second one is that optimization is a much wider *art* than that described in these books: my objective is to stimulate the interest of the reader and open their eyes to a continuously expanding body of knowledge.



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## Chapter 1

# Introduction

## 1.1 Introduction

Optimization is the act of achieving the best possible result under given circumstances.

In design, construction, maintenance, ..., engineers have to take decisions. The goal of all such decisions is either to minimize effort or to maximize benefit.

The effort or the benefit can be usually expressed as a function of certain design variables. Hence, optimization is the process of finding the conditions that give the maximum or the minimum value of a function.

It is obvious that if a point  $x^*$  corresponds to the minimum value of a function  $f(x)$ , the same point corresponds to the maximum value of the function  $-f(x)$ . Thus, optimization can be taken to be minimization.

There is no single method available for solving all optimization problems efficiently. Hence, a number of methods have been developed for solving different types of problems.

Optimum seeking methods are also known as mathematical programming techniques, which are a branch of operations research. Operations research is *coarsely* composed of the following areas.

- Mathematical programming methods. These are useful in finding the minimum of a function of several variables under a prescribed set of constraints.
- Stochastic process techniques. These are used to analyze problems which are described by a set of random variables of known distribution.
- Statistical methods. These are used in the analysis of experimental data and in the construction of empirical models.

These lecture notes deal mainly with the theory and applications of mathematical programming methods. Mathematical programming is a vast area of mathematics and engineering. It includes

- calculus of variations and optimal control;
- linear, quadratic and non-linear programming;
- geometric programming;
- integer programming;
- network methods (PERT);
- game theory.

The foundations of optimization can be traced back to Newton, Lagrange and Cauchy. The development of differential methods for optimization was possible because of the contribution of Newton and Leibnitz. The foundations of the calculus of variations were laid by Bernoulli, Euler, Lagrange and Weierstrasse. Constrained optimization was first studied by Lagrange and the notion of descent was introduced by Cauchy.



Despite these early contributions, very little progress was made till the 20th century, when computer power made the implementation of optimization procedures possible and this in turn stimulated further research methods.

The major developments in the area of numerical methods for unconstrained optimization have been made in the UK. These include the development of the simplex method (Dantzig, 1947), the principle of optimality (Bellman, 1957), necessary and sufficient conditions of optimality (Kuhn and Tucker, 1951).

Optimization in its broadest sense can be applied to solve any engineering problem, *e.g.*

- design of aircraft for minimum weight;
- optimal (minimum time) trajectories for space missions;
- minimum weight design of structures for earthquake;
- optimal design of electric networks;
- optimal production planning, resources allocation, scheduling;
- shortest route;
- design of optimum pipeline networks;
- minimum processing time in production systems;
- optimal control.

## 1.2 Statement of an optimization problem

An optimization, or a mathematical programming problem can be stated as follows.

Find

$$x = (x^1, x^2, \dots, x^n)$$

which minimizes

$$f(x)$$

subject to the constraints

$$g_j(x) \leq 0 \tag{1.1}$$

for  $j = 1, \dots, m$ , and

$$l_j(x) = 0 \tag{1.2}$$

for  $j = 1, \dots, p$ .

The variable  $x$  is called the design vector,  $f(x)$  is the objective function,  $g_j(x)$  are the inequality constraints and  $l_j(x)$  are the equality constraints. The number of variables  $n$  and the number of constraints  $p + m$  need not be related. If  $p + m = 0$  the problem is called an unconstrained optimization problem.

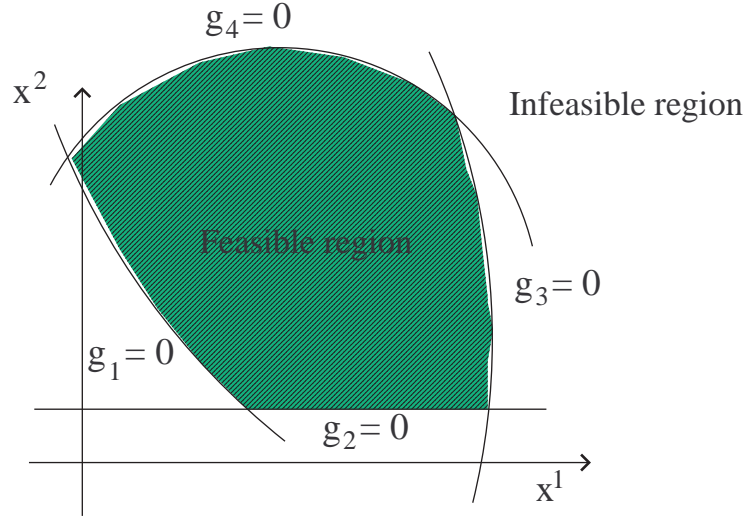


Figure 1.1: Feasible region in a two-dimensional design space. Only inequality constraints are present.

### 1.2.1 Design vector

Any system is described by a set of quantities, some of which are viewed as variables during the design process, and some of which are preassigned parameters or are imposed by the *environment*. All the quantities that can be treated as variables are called design or decision variables, and are collected in the design vector  $x$ .

### 1.2.2 Design constraints

In practice, the design variables cannot be selected arbitrarily, but have to satisfy certain requirements. These restrictions are called design constraints. Design constraints may represent limitation on the performance or behaviour of the system or physical limitations. Consider, for example, an optimization problem with only inequality constraints, *i.e.*  $g_j(x) \leq 0$ . The set of values of  $x$  that satisfy the equations  $g_j(x) = 0$  forms a hypersurface in the design space, which is called constraint surface. In general, if  $n$  is the number of design variables, the constraint surface is an  $n - 1$  dimensional surface. The constraint surface divides the design space into two regions: one in which  $g_j(x) < 0$  and one in which  $g_j(x) > 0$ . The points  $x$  on the constraint surface satisfy the constraint critically, whereas the points  $x$  such that  $g_j(x) > 0$ , for some  $j$ , are infeasible, *i.e.* are unacceptable, see Figure 1.1.

### 1.2.3 Objective function

The classical design procedure aims at finding an acceptable design, *i.e.* a design which satisfies the constraints. In general there are several acceptable designs, and the purpose

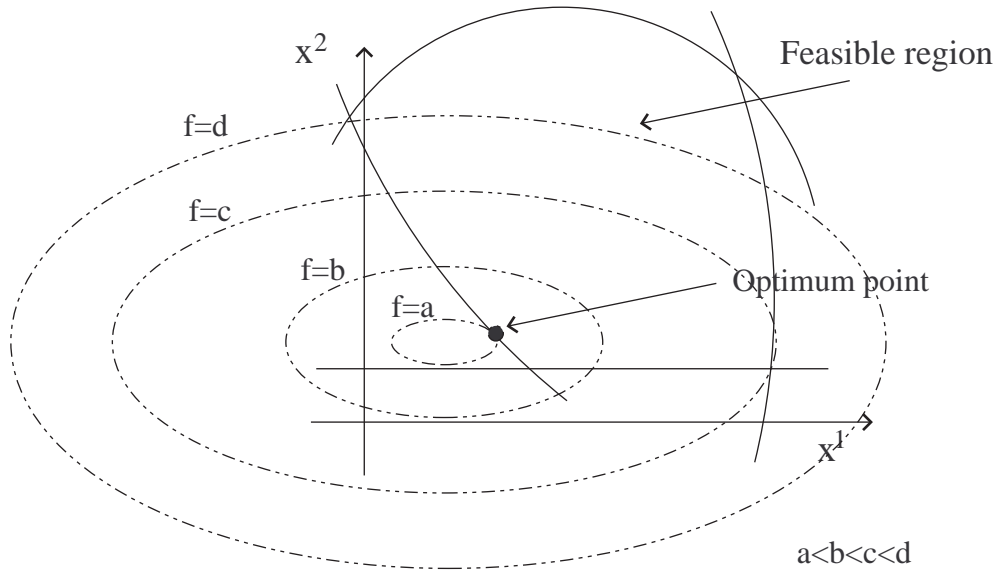


Figure 1.2: Design space, objective functions surfaces, and optimum point.

of the optimization is to single out the best possible design. Thus, a criterion has to be selected for comparing different designs. This criterion, when expressed as a function of the design variables, is known as objective function. The objective function is in general specified by physical or economical considerations. However, the selection of an objective function is not trivial, because what is the optimal design with respect to a certain criterion may be unacceptable with respect to another criterion. Typically there is a trade off performance–cost, or performance–reliability, hence the selection of the objective function is one of the most important decisions in the whole design process. If more than one criterion has to be satisfied we have a multiobjective optimization problem, that may be approximately solved considering a cost function which is a weighted sum of several objective functions.

Given an objective function  $f(x)$ , the locus of all points  $x$  such that  $f(x) = c$  forms a hypersurface. For each value of  $c$  there is a different hypersurface. The set of all these surfaces are called objective function surfaces.

Once the objective function surfaces are drawn, together with the constraint surfaces, the optimization problem can be easily solved, at least in the case of a two dimensional decision space, as shown in Figure 1.2. If the number of decision variables exceeds two or three, this graphical approach is not viable and the problem has to be solved as a mathematical problem. Note however that more general problems have similar geometrical properties of two or three dimensional problems.

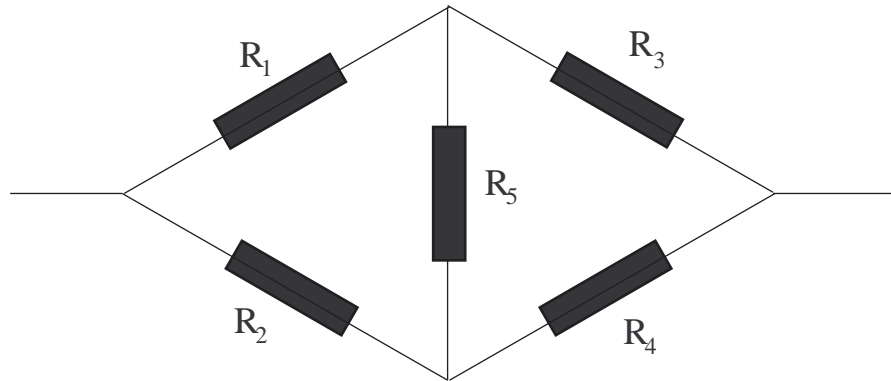


Figure 1.3: Electrical bridge network.

### 1.3 Classification of optimization problems

Optimization problem can be classified in several ways.

- Existence of constraints. An optimization problem can be classified as a constrained or an unconstrained one, depending upon the presence or not of constraints.
- Nature of the equations. Optimization problems can be classified as linear, quadratic, polynomial, non-linear depending upon the nature of the objective functions and the constraints. This classification is important, because computational methods are usually selected on the basis of such a classification, *i.e.* the nature of the involved functions dictates the type of solution procedure.
- Admissible values of the design variables. Depending upon the values permitted for the design variables, optimization problems can be classified as integer or real valued, and deterministic or stochastic.

### 1.4 Examples

**Example 1** A travelling salesman has to visit  $n$  towns. He plans to start from a particular town numbered 1, visit each one of the other  $n - 1$  towns, and return to the town 1. The distance between town  $i$  and  $j$  is given by  $d_{ij}$ . How should he select the sequence in which the towns are visited to minimize the total distance travelled?

**Example 2** The bridge network in Figure 1.3 consists of five resistors  $R_i$ ,  $i = 1, \dots, 5$ . Let  $I_i$  be the current through the resistance  $R_i$ , find the values of  $R_i$  so that the total dissipated power is minimum. The current  $I_i$  can vary between the lower limit  $\underline{I}_i$  and the upper limit  $\bar{I}_i$  and the voltage drop  $V_i = R_i I_i$  must be equal to a constant  $c_i$ .

**Example 3** A manufacturing firm produces two products, A and B, using two limited resources, 1 and 2. The maximum amount of resource 1 available per week is 1000 and the

Article type	$w_i$	$v_i$	$c_i$
1	4	9	5
2	8	7	6
3	2	4	3

Table 1.1: Properties of the articles to load.

maximum amount of resource 2 is 250. The production of one unit of A requires 1 unit of resource 1 and  $1/5$  unit of resource 2. The production of one unit of B requires  $1/2$  unit of resource 1 and  $1/2$  unit of resource 2. The unit cost of resource 1 is  $1 - 0.0005u_1$ , where  $u_1$  is the number of units of resource 1 used. The unit cost of resource 2 is  $3/4 - 0.0001u_2$ , where  $u_2$  is the number of units of resource 2 used. The selling price of one unit of A is

$$2 - 0.005x_A - 0.0001x_B$$

and the selling price of one unit of B is

$$4 - 0.002x_A - 0.01x_B,$$

where  $x_A$  and  $x_B$  are the number of units of A and B sold. Assuming that the firm is able to sell all manufactured units, maximize the weekly profit.

**Example 4** A cargo load is to be prepared for three types of articles. The weight,  $w_i$ , volume,  $v_i$ , and value,  $c_i$ , of each article is given in Table 1.1.

Find the number of articles  $x_i$  selected from type  $i$  so that the total value of the cargo is maximized. The total weight and volume of the cargo cannot exceed 2000 and 2500 units respectively.

**Example 5** There are two types of gas molecules in a gaseous mixture at equilibrium. It is known that the Gibbs free energy

$$G(x) = c_1x^1 + c_2x^2 + x^1\log(x^1/x_T) + x^2\log(x^2/x_T),$$

with  $x_T = x^1 + x^2$  and  $c_1, c_2$  known parameters depending upon the temperature and pressure of the mixture, has to be minimum in these conditions. The minimization of  $G(x)$  is also subject to the mass balance equations:

$$x^1a_{i1} + x^2a_{i2} = b_i,$$

for  $i = 1, \dots, m$ , where  $m$  is the number of atomic species in the mixture,  $b_i$  is the total weight of atoms of type  $i$ , and  $a_{ij}$  is the number of atoms of type  $i$  in the molecule of type  $j$ . Show that the problem of determining the equilibrium of the mixture can be posed as an optimization problem.



## Chapter 2

# Unconstrained optimization

## 2.1 Introduction

Several engineering, economic and planning problems can be posed as optimization problems, *i.e.* as the problem of determining the points of minimum of a function (possibly in the presence of conditions on the decision variables). Moreover, also numerical problems, such as the problem of solving systems of equations or inequalities, can be posed as an optimization problem.

We start with the study of optimization problems in which the decision variables are defined in  $\mathbb{R}^n$ : unconstrained optimization problems. More precisely we study the problem of determining local minimizers for differentiable functions. Although these methods are seldom used in applications, as in real problems the decision variables are subject to constraints, the techniques of unconstrained optimization are instrumental to solve more general problems: the knowledge of good methods for local unconstrained minimization is a necessary pre-requisite for the solution of constrained and global minimization problems. The methods that will be studied can be classified from various points of view. The most interesting classification is based on the information available on the function to be optimized, namely

- methods without derivatives (direct search, finite differences);
- methods based on the knowledge of the first derivatives (gradient, conjugate directions, quasi-Newton);
- methods based on the knowledge of the first and second derivatives (Newton).

## 2.2 Definitions and existence conditions

Consider the following general optimization problem.

**Problem 1** *Minimize*

$$f(x) \quad \text{subject to } x \in \mathcal{F}$$

*in which*  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  *and*<sup>1</sup>  $\mathcal{F} \subset \mathbb{R}^n$ .

With respect to this problem we introduce the following definitions.

**Definition 1** *A point*  $x \in \mathcal{F}$  *is a global minimizer for the Problem 1 if*

$$f(x) \leq f(y)$$

*for all*  $y \in \mathcal{F}$ .

*A point*  $x \in \mathcal{F}$  *is a strict (or isolated) global minimizer for the Problem 1 if*

$$f(x) < f(y)$$

---

<sup>1</sup>The set  $\mathcal{F}$  may be specified by equations of the form (1.1) and/or (1.2).



for all  $y \in \mathcal{F}$  and  $y \neq x$ .

A point  $x \in \mathcal{F}$  is a local minimiser for the Problem 1 if there exists  $\rho > 0$  such that

$$f(x) \leq f(y)$$

for all  $y \in \mathcal{F}$  such that  $\|y - x\| < \rho$ .

A point  $x \in \mathcal{F}$  is a strict (or isolated) local minimizer for the Problem 1 if there exists  $\rho > 0$  such that

$$f(x) < f(y)$$

for all  $y \in \mathcal{F}$  such that  $\|y - x\| < \rho$  and  $y \neq x$ .

**Definition 2** If  $x \in \mathcal{F}$  is a local minimizer for the Problem 1 and if  $x$  is in the interior of  $\mathcal{F}$  then  $x$  is an unconstrained local minimizer of  $f$  in  $\mathcal{F}$ .

The following result provides a sufficient, but not necessary, condition for the existence of a global minimum for Problem 1.

**Proposition 1** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a continuous function and let  $\mathcal{F} \subset \mathbb{R}^n$  be a compact set<sup>2</sup>. Then there exists a global minimum of  $f$  in  $\mathcal{F}$ .

In unconstrained optimization problems the set  $\mathcal{F}$  coincides with  $\mathbb{R}^n$ , hence the above statement cannot be used to establish the existence of global minima. To address the existence problem it is necessary to consider the structure of the level sets of the function  $f$ . See also Section 1.2.3.

**Definition 3** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ . A level set of  $f$  is any non-empty set described by

$$\mathcal{L}(\alpha) = \{x \in \mathbb{R}^n : f(x) \leq \alpha\},$$

with  $\alpha \in \mathbb{R}$ .

For convenience, if  $x_0 \in \mathbb{R}^n$  we denote with  $\mathcal{L}_0$  the level set  $\mathcal{L}(f(x_0))$ . Using the concept of level sets it is possible to establish a simple sufficient condition for the existence of global solutions for an unconstrained optimization problem.

**Proposition 2** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a continuous function. Assume there exists  $x_0 \in \mathbb{R}^n$  such that the level set  $\mathcal{L}_0$  is compact. Then there exists a point of global minimum of  $f$  in  $\mathbb{R}^n$ .

*Proof.* By Proposition 1 there exists a global minimizer  $x_*$  of  $f$  in  $\mathcal{L}_0$ , i.e.  $f(x_*) \leq f(x)$  for all  $x \in \mathcal{L}_0$ . However, if  $x \notin \mathcal{L}_0$  then  $f(x) > f(x_0) \geq f(x_*)$ , hence  $x_*$  is a global minimizer of  $f$  in  $\mathbb{R}^n$ .  $\triangleleft$

It is obvious that the structure of the level sets of the function  $f$  plays a fundamental role in the solution of Problem 1. The following result provides a necessary and sufficient condition for the compactness of all level sets of  $f$ .

---

<sup>2</sup>A compact set is a bounded and closed set.

**Proposition 3** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a continuous function. All level sets of  $f$  are compact if and only if for any sequence  $\{x_k\}$  one has*

$$\lim_{k \rightarrow \infty} \|x_k\| = \infty \quad \Rightarrow \quad \lim_{k \rightarrow \infty} f(x_k) = \infty.$$

*Remark.* In general  $x_k \in \mathbb{R}^n$ , namely

$$x_k = \begin{bmatrix} x_k^1 \\ x_k^2 \\ \vdots \\ x_k^n \end{bmatrix},$$

*i.e.* we use superscripts to denote components of a vector. ◇

A function that satisfies the condition of the above proposition is said to be radially unbounded.

*Proof.* We only prove the necessity. Suppose all level sets of  $f$  are compact. Then, proceeding by contradiction, suppose there exist a sequence  $\{x_k\}$  such that  $\lim_{k \rightarrow \infty} \|x_k\| = \infty$  and a number  $\gamma > 0$  such that  $f(x_k) \leq \gamma < \infty$  for all  $k$ . As a result

$$\{x_k\} \subset \mathcal{L}(\gamma).$$

However, by compactness of  $\mathcal{L}(\gamma)$  it is not possible that  $\lim_{k \rightarrow \infty} \|x_k\| = \infty$ . ◁

**Definition 4** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ . A vector  $d \in \mathbb{R}^n$  is said to be a descent direction for  $f$  in  $x_*$  if there exists  $\delta > 0$  such that*

$$f(x_* + \lambda d) < f(x_*),$$

*for all  $\lambda \in (0, \delta)$ .*

If the function  $f$  is differentiable it is possible to give a simple condition guaranteeing that a certain direction is a descent direction.

**Proposition 4** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume<sup>3</sup>  $\nabla f$  exists and is continuous. Let  $x_*$  and  $d$  be given. Then, if  $\nabla f(x_*)'d < 0$  the direction  $d$  is a descent direction for  $f$  at  $x_*$ .*

*Proof.* Note that  $\nabla f(x_*)'d$  is the directional derivative of  $f$  (which is differentiable by hypothesis) at  $x_*$  along  $d$ , *i.e.*

$$\nabla f(x_*)'d = \lim_{\lambda \rightarrow 0^+} \frac{f(x_* + \lambda d) - f(x_*)}{\lambda},$$

---

<sup>3</sup>We denote with  $\nabla f$  the gradient of the function  $f$ , *i.e.*  $\nabla f = [\frac{\partial f}{\partial x^1}, \dots, \frac{\partial f}{\partial x^n}]'$ . Note that  $\nabla f$  is a column vector.

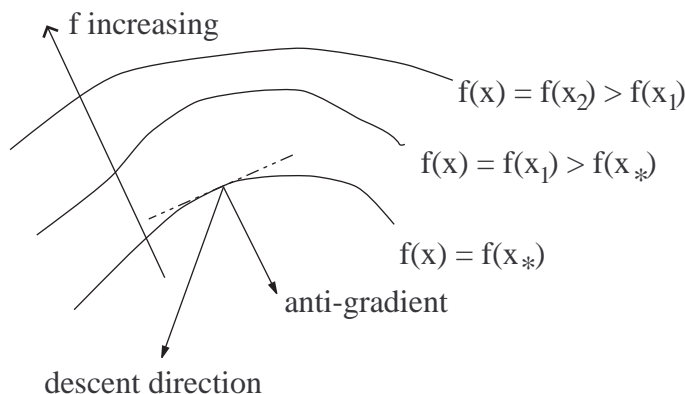


Figure 2.1: Geometrical interpretation of the anti-gradient.

and this is negative by hypothesis. As a result, for  $\lambda > 0$  and sufficiently small

$$f(x_* + \lambda d) - f(x_*) < 0,$$

hence the claim.  $\triangleleft$

The proposition establishes that if  $\nabla f(x_*)'d < 0$  then for sufficiently small positive displacements along  $d$  and starting at  $x_*$  the function  $f$  is decreasing. It is also obvious that if  $\nabla f(x_*)'d > 0$ ,  $d$  is a direction of *ascent*, i.e. the function  $f$  is increasing for sufficiently small positive displacements from  $x_*$  along  $d$ . If  $\nabla f(x_*)'d = 0$ ,  $d$  is orthogonal to  $\nabla f(x_*)$  and it is not possible to establish, without further knowledge on the function  $f$ , what is the nature of the direction  $d$ .

From a geometrical point of view (see also Figure 2.1), the sign of the *directional derivative*  $\nabla f(x_*)'d$  gives information on the angle between  $d$  and the direction of the gradient at  $x_*$ , provided  $\nabla f(x_*) \neq 0$ . If  $\nabla f(x_*)'d > 0$  the angle between  $\nabla f(x_*)$  and  $d$  is acute. If  $\nabla f(x_*)'d < 0$  the angle between  $\nabla f(x_*)$  and  $d$  is obtuse. Finally, if  $\nabla f(x_*)'d = 0$ , and  $\nabla f(x_*) \neq 0$ ,  $\nabla f(x_*)$  and  $d$  are orthogonal. Note that the gradient  $\nabla f(x_*)$ , if it is not identically zero, is a direction orthogonal to the level surface  $\{x : f(x) = f(x_*)\}$  and it is a direction of ascent, hence the anti-gradient  $-\nabla f(x_*)$  is a descent direction.

*Remark.* The scalar product  $x'y$  between the two vectors  $x$  and  $y$  can be used to define the angle between  $x$  and  $y$ . For, define the angle between  $x$  and  $y$  as the number  $\theta \in [0, \pi]$  such that<sup>4</sup>

$$\cos \theta = \frac{x'y}{\|x\|_E \|y\|_E}.$$

If  $x'y = 0$  one has  $\cos \theta = 0$  and the vectors are orthogonal, whereas if  $x$  and  $y$  have the same direction, i.e.  $x = \lambda y$  with  $\lambda > 0$ ,  $\cos \theta = 1$ .  $\diamond$

---

<sup>4</sup> $\|x\|_E$  denotes the Euclidean norm of the vector  $x$ , i.e.  $\|x\|_E = \sqrt{x'x}$ .

We are now ready to state and prove some necessary conditions and some sufficient conditions for a local minimizer.

**Theorem 1** *[First order necessary condition] Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume  $\nabla f$  exists and is continuous. The point  $x_*$  is a local minimizer of  $f$  only if*

$$\nabla f(x_*) = 0.$$

*Remark.* A point  $x_*$  such that  $\nabla f(x_*) = 0$  is called a stationary point of  $f$ .  $\diamond$

*Proof.* If  $\nabla f(x_*) \neq 0$  the direction  $d = -\nabla f(x_*)$  is a descent direction. Therefore, in a neighborhood of  $x_*$  there is a point  $x_* + \lambda d = x_* - \lambda \nabla f(x_*)$  such that

$$f(x_* - \lambda \nabla f(x_*)) < f(x_*),$$

and this contradicts the hypothesis that  $x_*$  is a local minimizer.  $\triangleleft$

**Theorem 2** *[Second order necessary condition] Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume<sup>5</sup>  $\nabla^2 f$  exists and is continuous. The point  $x_*$  is a local minimizer of  $f$  only if*

$$\nabla f(x_*) = 0$$

and

$$x' \nabla^2 f(x_*) x \geq 0$$

for all  $x \in \mathbb{R}^n$ .

*Proof.* The first condition is a consequence of Theorem 1. Note now that, as  $f$  is two times differentiable, for any  $x \neq x_*$  one has

$$f(x_* + \lambda x) = f(x_*) + \lambda \nabla f(x_*)' x + \frac{1}{2} \lambda^2 x' \nabla^2 f(x_*) x + \beta(x_*, \lambda x),$$

where

$$\lim_{\lambda \rightarrow 0} \frac{\beta(x_*, \lambda x)}{\lambda^2 \|x\|^2} = 0,$$

or what is the same (note that  $x$  is fixed)

$$\lim_{\lambda \rightarrow 0} \frac{\beta(x_*, \lambda x)}{\lambda^2} = 0.$$

---

<sup>5</sup>We denote with  $\nabla^2 f$  the Hessian matrix of the function  $f$ , i.e.

$$\begin{bmatrix} \frac{\partial^2 f}{\partial x^1 \partial x^1} & \cdots & \frac{\partial^2 f}{\partial x^1 \partial x^n} \\ \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x^n \partial x^1} & \cdots & \frac{\partial^2 f}{\partial x^n \partial x^n} \end{bmatrix}.$$

Note that  $\nabla^2 f$  is a square matrix and that, under suitable regularity conditions, the Hessian matrix is symmetric.

Moreover, the condition  $\nabla f(x_\star) = 0$  yields

$$\frac{f(x_\star + \lambda x) - f(x_\star)}{\lambda^2} = \frac{1}{2}x' \nabla^2 f(x_\star)x + \frac{\beta(x_\star, \lambda x)}{\lambda^2}. \quad (2.1)$$

However, as  $x_\star$  is a local minimizer, the left hand side of equation (2.1) must be non-negative for all  $\lambda$  sufficiently small, hence

$$\frac{1}{2}x' \nabla^2 f(x_\star)x + \frac{\beta(x_\star, \lambda x)}{\lambda^2} \geq 0,$$

and

$$\lim_{\lambda \rightarrow 0} \left( \frac{1}{2}x' \nabla^2 f(x_\star)x + \frac{\beta(x_\star, \lambda x)}{\lambda^2} \right) = \frac{1}{2}x' \nabla^2 f(x_\star)x \geq 0,$$

which proves the second condition.  $\triangleleft$

**Theorem 3 (Second order sufficient condition)** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume  $\nabla^2 f$  exists and is continuous. The point  $x_\star$  is a strict local minimizer of  $f$  if*

$$\nabla f(x_\star) = 0$$

and

$$x' \nabla^2 f(x_\star)x > 0$$

for all non-zero  $x \in \mathbb{R}^n$ .

*Proof.* To begin with, note that as  $\nabla^2 f(x_\star) > 0$  and  $\nabla^2 f$  is continuous, then there is a neighborhood  $\Omega$  of  $x_\star$  such that for all  $y \in \Omega$

$$\nabla^2 f(y) > 0.$$

Consider now the Taylor series expansion of  $f$  around the point  $x_\star$ , i.e.

$$f(y) = f(x_\star) + \nabla f(x_\star)'(y - x_\star) + \frac{1}{2}(y - x_\star)' \nabla^2 f(\xi)(y - x_\star),$$

where  $\xi = x_\star + \theta(y - x_\star)$ , for some  $\theta \in [0, 1]$ . By the first condition one has

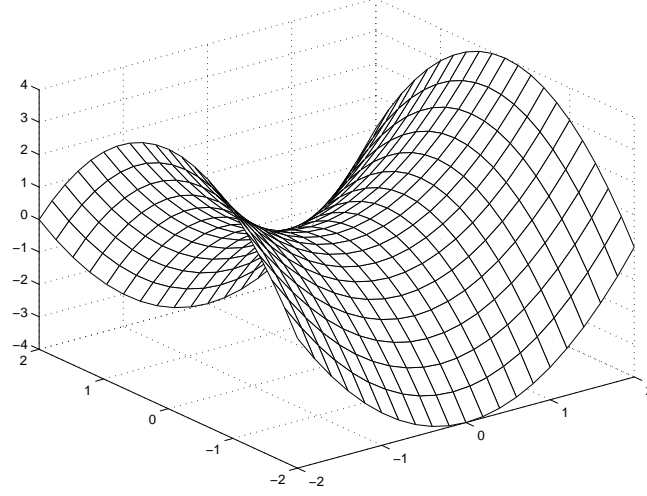
$$f(y) = f(x_\star) + \frac{1}{2}(y - x_\star)' \nabla^2 f(\xi)(y - x_\star),$$

and, for any  $y \in \Omega$  such that  $y \neq x_\star$ ,

$$f(y) > f(x_\star),$$

which proves the claim.  $\triangleleft$

The above results can be easily modified to derive necessary conditions and sufficient conditions for a local maximizer. Moreover, if  $x_\star$  is a stationary point and the Hessian matrix

Figure 2.2: A saddle point in  $\mathbb{R}^2$ .

$\nabla^2 f(x_*)$  is indefinite, the point  $x_*$  is neither a local minimizer neither a local maximizer. Such a point is called a saddle point (see Figure 2.2 for a geometrical illustration).

If  $x_*$  is a stationary point and  $\nabla^2 f(x_*)$  is semi-definite it is not possible to draw any conclusion on the point  $x_*$  without further knowledge on the function  $f$ . Nevertheless, if  $n = 1$  and the function  $f$  is infinitely times differentiable it is possible to establish the following necessary and sufficient condition.

**Proposition 5** *Let  $f : \mathbb{R} \rightarrow \mathbb{R}$  and assume  $f$  is infinitely times differentiable. The point  $x_*$  is a local minimizer if and only if there exists an even integer  $r > 1$  such that*

$$\frac{d^k f(x_*)}{dx^k} = 0$$

for  $k = 1, 2, \dots, r - 1$  and

$$\frac{d^r f(x_*)}{dx^r} > 0.$$

Necessary and sufficient conditions for  $n > 1$  can be only derived if further hypotheses on the function  $f$  are added, as shown for example in the following fact.

**Proposition 6 (Necessary and sufficient condition for convex functions)** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume  $\nabla f$  exists and it is continuous. Suppose  $f$  is convex, i.e.*

$$f(y) - f(x) \geq \nabla f(x)'(y - x) \quad (2.2)$$

for all  $x \in \mathbb{R}^n$  and  $y \in \mathbb{R}^n$ . The point  $x_*$  is a global minimizer if and only if  $\nabla f(x_*) = 0$ .

*Remark.* The convexity condition (2.2) can be rewritten as

$$f(y) \geq f(x) + \nabla f(x)'(y - x).$$

This reveals that the tangent plane for  $f$  at  $x$  is always below the graph of the function, *i.e.* the function is supported from below by the tangent planes.  $\diamond$

*Proof.* The necessity is a consequence of Theorem 1. For the sufficiency note that, by equation (2.2), if  $\nabla f(x_*) = 0$  then

$$f(y) \geq f(x_*),$$

for all  $y \in \mathbb{R}^n$ .  $\triangleleft$

From the above discussion it is clear that to establish the property that  $x_*$ , satisfying  $\nabla f(x_*) = 0$ , is a global minimizer it is enough to assume that the function  $f$  has the following property: for all  $x$  and  $y$  such that

$$\nabla f(x)'(y - x) \geq 0$$

one has

$$f(y) \geq f(x).$$

A function  $f$  satisfying the above property is said pseudo-convex. Note that a differentiable convex function is also pseudo-convex, but the opposite is not true. For example, the function  $x + x^3$  is pseudo-convex but it is not convex. Finally, if  $f$  is strictly convex or strictly pseudo-convex the global minimizer (if it exists) is also unique.

## 2.3 General properties of minimization algorithms

Consider the problem of minimizing the function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and suppose that  $\nabla f$  and  $\nabla^2 f$  exist and are continuous. Suppose that such a problem has a solution, and moreover that there exists  $x_0$  such that the level set

$$\mathcal{L}(f(x_0)) = \{x \in \mathbb{R}^n : f(x) \leq f(x_0)\}$$

is compact.

General unconstrained minimization algorithms allow only to determine stationary points of  $f$ , *i.e.* to determine points in the set

$$\Omega = \{x \in \mathbb{R}^n : \nabla f(x) = 0\}.$$

Moreover, for almost all algorithms, it is possible to exclude that the points of  $\Omega$  yielded by the algorithm are local maximizer. Finally, some algorithms yield points of  $\Omega$  that satisfy also the second order necessary conditions.

### 2.3.1 General unconstrained minimization algorithm

An *algorithm* for the solution of the considered minimization problem is a sequence  $\{x_k\}$ , obtained starting from an initial point  $x_0$ , having some convergence properties in relation with the set  $\Omega$ . Most of the algorithms that will be studied in this notes can be described in the following general way.

- a) Fix a point  $x_0 \in \mathbb{R}^n$  and set  $k = 0$ .
- b) If  $x_k \in \Omega$  STOP.
- c) Compute a direction of research  $d_k \in \mathbb{R}^n$ .
- d) Compute a step  $\alpha_k \in \mathbb{R}$  along  $d_k$ .
- e) Let  $x_{k+1} = x_k + \alpha_k d_k$ . Set  $k = k + 1$  and go back to 2.

The existing algorithms differ in the way the direction of research  $d_k$  is computed and on the criteria used to compute the step  $\alpha_k$ . However, independently from the particular selection, it is important to study the following issues:

- the existence of accumulation points for the sequence  $\{x_k\}$ ;
- the behavior of such accumulation points in relation with the set  $\Omega$ ;
- the speed of convergence of the sequence  $\{x_k\}$  to the points of  $\Omega$ .

### 2.3.2 Existence of accumulation points

To make sure that any subsequence of  $\{x_k\}$  has an accumulation point it is necessary to assume that the sequence  $\{x_k\}$  remains bounded, *i.e.* that there exists  $M > 0$  such that  $\|x_k\| < M$  for any  $k$ . If the level set  $\mathcal{L}(f(x_0))$  is compact, the above condition holds if  $\{x_k\} \in \mathcal{L}(f(x_0))$ . This property, in turn, is guaranteed if

$$f(x_{k+1}) < f(x_k),$$

for any  $k$  such that  $x_k \notin \Omega$ . The algorithms that satisfy this property are denominated descent methods. For such methods, if  $\mathcal{L}(f(x_0))$  is compact and if  $\nabla f$  is continuous one has

- $\{x_k\} \in \mathcal{L}(f(x_0))$  and any subsequence of  $\{x_k\}$  admits a subsequence converging to a point of  $\mathcal{L}(f(x_0))$ ;
- the sequence  $\{f(x_k)\}$  has a limit, *i.e.* there exists  $\bar{f} \in \mathbb{R}$  such that

$$\lim_{k \rightarrow \infty} f(x_k) = \bar{f};$$

- there always exists an element of  $\Omega$  in  $\mathcal{L}(f(x_0))$ . In fact, as  $f$  has a minimizer in  $\mathcal{L}(f(x_0))$ , this minimizer is also a minimizer of  $f$  in  $\mathbb{R}^n$ . Hence, by the assumptions of  $\nabla f$ , such a minimizer must be a point of  $\Omega$ .



*Remark.* To guarantee the descent property it is necessary that the research directions  $d_k$  be directions of descent. This is true if

$$\nabla f(x_k)'d_k < 0,$$

for all  $k$ . Under this condition there exists an interval  $(0, \alpha_*]$  such that

$$f(x_k + \alpha d_k) < f(x_k),$$

for any  $\alpha \in (0, \alpha_*]$ .  $\diamond$

*Remark.* The existence of accumulation points for the sequence  $\{x_k\}$  and the convergence of the sequence  $\{f(x_k)\}$  do not guarantee that the accumulation points of  $\{x_k\}$  are local minimizers of  $f$  or stationary points. To obtain this property it is necessary to impose further restrictions on the research directions  $d_k$  and on the steps  $\alpha_k$ .  $\diamond$

### 2.3.3 Condition of angle

The condition which is in general imposed on the research directions  $d_k$  is the so-called condition of angle, that can be stated as follows.

**Condition 1** *There exists  $\epsilon > 0$ , independent from  $k$ , such that*

$$\nabla f(x_k)'d_k \leq -\epsilon \|\nabla f(x_k)\| \|d_k\|,$$

*for any  $k$ .*

From a geometric point of view the above condition implies that the cosine of the angle between  $d_k$  and  $-\nabla f(x_k)$  is larger than a certain quantity. This condition is imposed to avoid that, for some  $k$ , the research direction is orthogonal to the direction of the gradient. Note moreover that, if the angle condition holds, and if  $\nabla f(x_k) \neq 0$  then  $d_k$  is a descent direction. Finally, if  $\nabla f(x_k) \neq 0$ , it is always possible to find a direction  $d_k$  such that the angle condition holds. For example, the direction  $d_k = -\nabla f(x_k)$  is such that the angle condition is satisfied with  $\epsilon = 1$ .

*Remark.* Let  $\{B_k\}$  be a sequence of matrices such that

$$mI \leq B_k \leq MI,$$

for some  $0 < m < M$ , and for any  $k$ , and consider the directions

$$d_k = -B_k \nabla f(x_k).$$

Then a simple computation shows that the angle condition holds with  $\epsilon = m/M$ .  $\diamond$

The angle condition imposes a constraint only on the research directions  $d_k$ . To make sure that the sequence  $\{x_k\}$  converges to a point in  $\Omega$  it is necessary to impose further conditions on the step  $\alpha_k$ , as expressed in the following statements.

**Theorem 4** *Let  $\{x_k\}$  be the sequence obtained by the algorithm*

$$x_{k+1} = x_k + \alpha_k d_k,$$

*for  $k \geq 0$ . Assume that*

*(H1)  $\nabla f$  is continuous and the level set  $\mathcal{L}(f(x_0))$  is compact.*

*(H2) There exists  $\epsilon > 0$  such that*

$$\nabla f(x_k)' d_k \leq -\epsilon \|\nabla f(x_k)\| \|d_k\|,$$

*for any  $k \geq 0$ .*

*(H3)  $f(x_{k+1}) < f(x_k)$  for any  $k \geq 0$ .*

*(H4) The property*

$$\lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} = 0$$

*holds.*

*Then*

*(C1)  $\{x_k\} \in \mathcal{L}(f(x_0))$  and any subsequence of  $\{x_k\}$  has an accumulation point.*

*(C2)  $\{f(x_k)\}$  is monotonically decreasing and there exists  $\bar{f}$  such that*

$$\lim_{k \rightarrow \infty} f(x_k) = \bar{f}.$$

*(C3)  $\{\nabla f(x_k)\}$  is such that*

$$\lim_{k \rightarrow \infty} \|\nabla f(x_k)\| = 0.$$

*(C4) Any accumulation point  $\bar{x}$  of  $\{x_k\}$  is such that  $\nabla f(\bar{x}) = 0$ .*

*Proof.* Conditions (C1) and (C2) are a simple consequence of (H1) and (H3). Note now that (H2) implies

$$\epsilon \|\nabla f(x_k)\| \leq \frac{|\nabla f(x_k)' d_k|}{\|d_k\|},$$

for all  $k$ . As a result, and by (H4),

$$\lim_{k \rightarrow \infty} \epsilon \|\nabla f(x_k)\| \leq \lim_{k \rightarrow \infty} \frac{|\nabla f(x_k)' d_k|}{\|d_k\|} = 0$$

hence (C3) holds. Finally, let  $\bar{x}$  be an accumulation point of the sequence  $\{x_k\}$ , *i.e.* there is a subsequence that converges to  $\bar{x}$ . For such a subsequence, and by continuity of  $f$ , one has

$$\lim_{k \rightarrow \infty} \nabla f(x_k) = \nabla f(\bar{x}),$$

and, by (C3),

$$\nabla f(\bar{x}) = 0,$$

which proves (C4).  $\triangleleft$

*Remark.* Theorem 4 does not guarantee the convergence of the sequence  $\{x_k\}$  to a unique accumulation point. Obviously  $\{x_k\}$  has a unique accumulation point if either  $\Omega \cap \mathcal{L}(f(x_0))$  contains only one point or  $x, y \in \Omega \cap \mathcal{L}(f(x_0))$ , with  $x \neq y$  implies  $f(x) \neq f(y)$ . Finally, if the set  $\Omega \cap \mathcal{L}(f(x_0))$  contains a finite number of points, a sufficient condition for the existence of a unique accumulation point is

$$\lim_{k \rightarrow \infty} \|x_{k+1} - x_k\| = 0.$$

$\diamond$

*Remark.* The angle condition can be replaced by the following one. There exists  $\eta > 0$  and  $q > 0$ , both independent from  $k$ , such that

$$\nabla f(x_k)' d_k \leq -\eta \|\nabla f(x_k)\|^q \|d_k\|.$$

$\diamond$

The result illustrated in Theorem 4 requires the fulfillment of the angle condition or of a similar one, *i.e.* of a condition involving  $\nabla f$ . In many algorithms that do not make use of the gradient it may be difficult to check the validity of the angle condition, hence it is necessary to use different conditions on the research directions. For example, it is possible to replace the angle condition with a property of linear independence of the research directions.

**Theorem 5** *Let  $\{x_k\}$  be the sequence obtained by the algorithm*

$$x_{k+1} = x_k + \alpha_k d_k,$$

*for  $k \geq 0$ . Assume that*

- $\nabla^2 f$  is continuous and the level set  $\mathcal{L}(f(x_0))$  is compact.
- There exist  $\sigma > 0$ , independent from  $k$ , and  $k_0 > 0$  such that, for any  $k \geq k_0$  the matrix  $P_k$  composed of the columns

$$\frac{d_k}{\|d_k\|}, \frac{d_{k+1}}{\|d_{k+1}\|}, \dots, \frac{d_{k+n-1}}{\|d_{k+n-1}\|},$$

*is such that*

$$|\det P_k| \geq \sigma.$$

- $\lim_{k \rightarrow \infty} \|x_{k+1} - x_k\| = 0$ .

- $f(x_{k+1}) < f(x_k)$  for any  $k \geq 0$ .
- The property

$$\lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} = 0$$

holds.

Then

- $\{x_k\} \in \mathcal{L}(f(x_0))$  and any subsequence of  $\{x_k\}$  has an accumulation point.
- $\{f(x_k)\}$  is monotonically decreasing and there exists  $\bar{f}$  such that

$$\lim_{k \rightarrow \infty} f(x_k) = \bar{f}.$$

- Any accumulation point  $\bar{x}$  of  $\{x_k\}$  is such that  $\nabla f(\bar{x}) = 0$ .

Moreover, if the set  $\Omega \cap \mathcal{L}(f(x_0))$  is composed of a finite number of points, the sequence  $\{x_k\}$  has a unique accumulation point.

### 2.3.4 Speed of convergence

Together with the property of convergence of the sequence  $\{x_k\}$  it is important to study also the speed of convergence. To study such a notion it is convenient to assume that  $\{x_k\}$  converges to a point  $x_*$ .

If there exists a finite  $k$  such that  $x_k = x_*$  then we say that the sequence  $\{x_k\}$  has finite convergence. Note that if  $\{x_k\}$  is generated by an algorithm, there is a stopping condition that has to be satisfied at step  $k$ .

If  $x_k \neq x_*$  for any finite  $k$ , it is possible (and convenient) to study the asymptotic properties of  $\{x_k\}$ . One criterion to estimate the speed of convergence is based on the behavior of the error  $\mathcal{E}_k = \|x_k - x_*\|$ , and in particular on the relation between  $\mathcal{E}_{k+1}$  and  $\mathcal{E}_k$ .

We say that  $\{x_k\}$  has speed of convergence of order  $p$  if

$$\lim_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^p} \right) = C_p$$

with  $p \geq 1$  and  $0 < C_p < \infty$ . Note that if  $\{x_k\}$  has speed of convergence of order  $p$  then

$$\lim_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^q} \right) = 0,$$

if  $1 \leq q < p$ , and

$$\lim_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^q} \right) = \infty,$$

if  $q > p$ . Moreover, from the definition of speed of convergence, it is easy to see that if  $\{x_k\}$  has speed of convergence of order  $p$  then, for any  $\epsilon > 0$  there exists  $k_0$  such that

$$\mathcal{E}_{k+1} \leq (C_p + \epsilon) \mathcal{E}_k^p,$$

for any  $k > k_0$ .

In the cases  $p = 1$  or  $p = 2$  the following terminology is often used. If  $p = 1$  and  $0 < C_1 \leq 1$  the speed of convergence is linear; if  $p = 1$  and  $C_1 > 1$  the speed of convergence is sublinear; if

$$\lim_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k} \right) = 0$$

the speed of convergence is superlinear, and finally if  $p = 2$  the speed of convergence is quadratic.

Of special interest in optimization is the case of superlinear convergence, as this is the kind of convergence that can be established for the *efficient* minimization algorithms. Note that if  $x_k$  has superlinear convergence to  $x_*$  then

$$\lim_{k \rightarrow \infty} \frac{\|x_{k+1} - x_k\|}{\|x_k - x_*\|} = 1.$$

*Remark.* In some cases it is not possible to establish the existence of the limit

$$\lim_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^q} \right).$$

In these cases an estimate of the speed of convergence is given by

$$Q_p = \limsup_{k \rightarrow \infty} \left( \frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^q} \right).$$

◇

## 2.4 Line search

A line search is a method to compute the step  $\alpha_k$  along a given direction  $d_k$ . The choice of  $\alpha_k$  affects both the convergence and the speed of convergence of the algorithm. In any line search one considers the function of one variable  $\phi : \mathbb{R} \rightarrow \mathbb{R}$  defined as

$$\phi(\alpha) = f(x_k + \alpha d_k) - f(x_k).$$

The derivative of  $\phi(\alpha)$  with respect to  $\alpha$  is given by

$$\dot{\phi}(\alpha) = \nabla f(x_k + \alpha d_k)' d_k$$

provided that  $\nabla f$  is continuous. Note that  $\nabla f(x_k + \alpha d_k)' d_k$  describes the slope of the tangent to the function  $\phi(\alpha)$ , and in particular

$$\dot{\phi}(0) = \nabla f(x_k)' d_k$$

coincides with the directional derivative of  $f$  at  $x_k$  along  $d_k$ .

From the general convergence results described, we conclude that the line search has to enforce the following conditions

$$\begin{aligned} f(x_{k+1}) &< f(x_k) \\ \lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} &= 0 \end{aligned}$$

and, whenever possible, also the condition

$$\lim_{k \rightarrow \infty} \|x_{k+1} - x_k\| = 0.$$

To begin with, we assume that the directions  $d_k$  are such that

$$\nabla f(x_k)' d_k < 0$$

for all  $k$ , *i.e.*  $d_k$  is a descent direction, and that it is possible to compute, for any fixed  $x$ , both  $f$  and  $\nabla f$ . Finally, we assume that the level set  $\mathcal{L}(f(x_0))$  is compact.

### 2.4.1 Exact line search

The exact line search consists in finding  $\alpha_k$  such that

$$\phi(\alpha_k) = f(x_k + \alpha_k d_k) - f(x_k) \leq f(x_k + \alpha d_k) - f(x_k) = \phi(\alpha)$$

for any  $\alpha \geq 0$ . Note that, as  $d_k$  is a descent direction and the set

$$\{\alpha \in \mathbb{R}^+ : \phi(\alpha) \leq \phi(0)\}$$

is compact, because of compactness of  $\mathcal{L}(f(x_0))$ , there exists an  $\alpha_k$  that minimizes  $\phi(\alpha)$ . Moreover, for such  $\alpha_k$  one has

$$\dot{\phi}(\alpha_k) = \nabla f(x_k + \alpha_k d_k)' d_k = 0,$$

*i.e.* if  $\alpha_k$  minimizes  $\phi(\alpha)$  the gradient of  $f$  at  $x_k + \alpha_k d_k$  is orthogonal to the direction  $d_k$ . From a geometrical point of view, if  $\alpha_k$  minimizes  $\phi(\alpha)$  then the level surface of  $f$  through the point  $x_k + \alpha_k d_k$  is tangent to the direction  $d_k$  at such a point. (If there are several points of tangency,  $\alpha_k$  is the one for which  $f$  has the smallest value).

The search of  $\alpha_k$  that minimizes  $\phi(\alpha)$  is very *expensive*, especially if  $f$  is not convex. Moreover, in general, the whole minimization algorithm does not gain any special advantage from the knowledge of such *optimal*  $\alpha_k$ . It is therefore more convenient to use approximate methods, *i.e.* methods which are computationally simple and which guarantee particular convergence properties. Such methods are aimed at finding an interval of acceptable values for  $\alpha_k$  subject to the following two conditions

- $\alpha_k$  has to guarantee a sufficient reduction of  $f$ ;
- $\alpha_k$  has to be sufficiently distant from 0, *i.e.*  $x_k + \alpha_k d_k$  has to be sufficiently away from  $x_k$ .

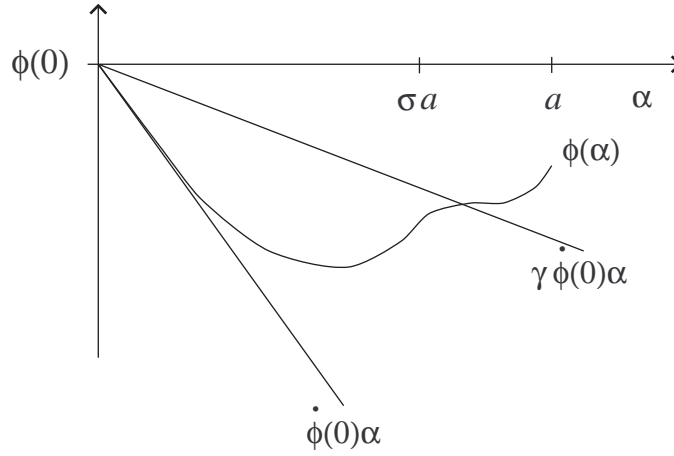


Figure 2.3: Geometrical interpretation of Armijo method.

### 2.4.2 Armijo method

Armijo method was the first non-exact linear search method.

Let  $a > 0$ ,  $\sigma \in (0, 1)$  and  $\gamma \in (0, 1/2)$  be given and define the set of points

$$A = \{\alpha \in \mathbb{R} : \alpha = a\sigma^j, j = 0, 1, \dots\}.$$

Armijo method consists in finding the largest  $\alpha \in A$  such that

$$\phi(\alpha) = f(x_k + \alpha d_k) - f(x_k) \leq \gamma \alpha \nabla f(x_k)' d_k = \gamma \alpha \dot{\phi}(0).$$

Armijo method can be implemented using the following (conceptual) algorithm.

**Step 1.** Set  $\alpha = a$ .

**Step 2.** If

$$f(x_k + \alpha d_k) - f(x_k) \leq \gamma \alpha \nabla f(x_k)' d_k$$

set  $\alpha_k = \alpha$  and STOP. Else go to **Step 3**.

**Step 3.** Set  $\alpha = \sigma \alpha$ , and go to **Step 2**.

From a geometric point of view (see Figure 2.3) the condition in **Step 2** requires that  $\alpha_k$  is such that  $\phi(\alpha_k)$  is below the straight line passing through the point  $(0, \phi(0))$  and with slope  $\gamma \dot{\phi}(0)$ . Note that, as  $\gamma \in (0, 1/2)$  and  $\dot{\phi}(0) < 0$ , such a straight line has a slope smaller than the slope of the tangent at the curve  $\phi(\alpha)$  at the point  $(0, \phi(0))$ .

For Armijo method it is possible to prove the following convergence result.

**Theorem 6** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume  $\nabla f$  is continuous and  $\mathcal{L}(f(x_0))$  is compact. Assume  $\nabla f(x_k)' d_k < 0$  for all  $k$  and there exist  $C_1 > 0$  and  $C_2 > 0$  such that

$$C_1 \geq \|d_k\| \geq C_2 \|\nabla f(x_k)\|^q,$$

for some  $q > 0$  and for all  $k$ .

Then Armijo method yields in a finite number of iterations a value of  $\alpha_k > 0$  satisfying the condition in **Step 2**. Moreover, the sequence obtained setting  $x_{k+1} = x_k + \alpha_k d_k$  is such that

$$f(x_{k+1}) < f(x_k),$$

for all  $k$ , and

$$\lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} = 0.$$

*Proof.* We only prove that the method cannot loop indefinitely between **Step 2** and **Step 3**. In fact, if this is the case, then the condition in **Step 2** will never be satisfied, hence

$$\frac{f(x_k + a\sigma^j d_k) - f(x_k)}{a\sigma^j} > \gamma \nabla f(x_k)' d_k.$$

Note now that  $\sigma^j \rightarrow 0$  as  $j \rightarrow \infty$ , and the above inequality for  $j \rightarrow \infty$  is

$$\nabla f(x_k)' d_k > \gamma \nabla f(x_k)' d_k,$$

which is not possible since  $\gamma \in (0, 1/2)$  and  $\nabla f(x_k)' d_k \neq 0$ .  $\triangleleft$

*Remark.* It is interesting to observe that in Theorem 6 it is not necessary to assume that  $x_{k+1} = x_k + \alpha_k d_k$ . It is enough that  $x_{k+1}$  is such that

$$f(x_{k+1}) \leq f(x_k + \alpha_k d_k),$$

where  $\alpha_k$  is generated using Armijo method. This implies that all acceptable values of  $\alpha$  are those such that

$$f(x_k + \alpha d_k) \leq f(x_k + \alpha_k d_k).$$

As a result, Theorem 6 can be used to prove also the convergence of an algorithm based on the exact line search.  $\diamond$

### 2.4.3 Goldstein conditions

The main disadvantage of Armijo method is in the fact that, to find  $\alpha_k$ , all points in the set  $A$ , starting from the point  $\alpha = a$ , have to be tested till the condition in **Step 2** is fulfilled. There are variations of the method that do not suffer from this disadvantage. A criterion similar to Armijo's, but that allows to find an acceptable  $\alpha_k$  in one step, is based on the so-called Goldstein conditions.

Goldstein conditions state that given  $\gamma_1 \in (0, 1)$  and  $\gamma_2 \in (0, 1)$  such that  $\gamma_1 < \gamma_2$ ,  $\alpha_k$  is any positive number such that

$$f(x_k + \alpha_k d_k) - f(x_k) \leq \alpha_k \gamma_1 \nabla f(x_k)' d_k$$



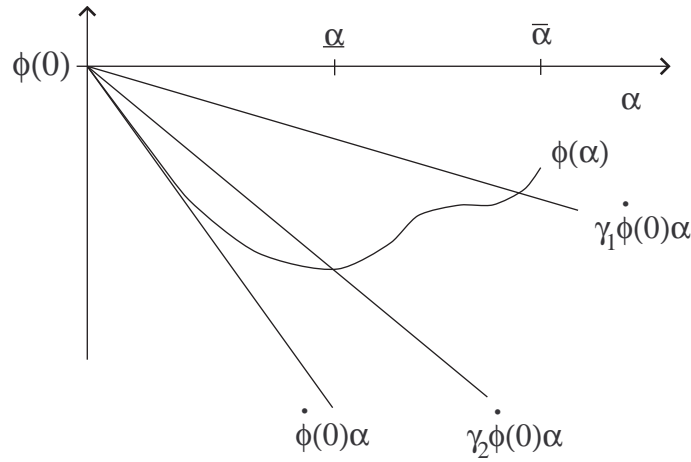


Figure 2.4: Geometrical interpretation of Goldstein method.

*i.e.* there is a sufficient reduction in  $f$ , and

$$f(x_k + \alpha_k d_k) - f(x_k) \geq \alpha_k \gamma_2 \nabla f(x_k)' d_k$$

*i.e.* there is a sufficient distance between  $x_k$  and  $x_{k+1}$ .

From a geometric point of view (see Figure 2.4) this is equivalent to select  $\alpha_k$  as any point such that the corresponding value of  $f$  is included between two straight lines, of slope  $\gamma_1 \nabla f(x_k)' d_k$  and  $\gamma_2 \nabla f(x_k)' d_k$ , respectively, and passing through the point  $(0, \phi(0))$ . As  $0 < \gamma_1 < \gamma_2 < 1$  it is obvious that there exists always an interval  $I = [\underline{\alpha}, \bar{\alpha}]$  such that Goldstein conditions hold for any  $\alpha \in I$ .

Note that, a result similar to Theorem 6, can be also established if the sequence  $\{x_k\}$  is generated using Goldstein conditions.

The main disadvantage of Armijo and Goldstein methods is in the fact that none of them impose conditions on the derivative of the function  $\phi(\alpha)$  in the point  $\alpha_k$ , or what is the same on the value of  $\nabla f(x_{k+1})' d_k$ . Such extra conditions are sometimes useful in establishing convergence results for particular algorithms. However, for simplicity, we omit the discussion of these more general conditions (known as Wolfe conditions).

#### 2.4.4 Line search without derivatives

It is possible to construct methods similar to Armijo's or Goldstein's also in the case that no information on the derivatives of the function  $f$  is available.

Suppose, for simplicity, that  $\|d_k\| = 1$ , for all  $k$ , and that the sequence  $\{x_k\}$  is generated by

$$x_{k+1} = x_k + \alpha_k d_k.$$

If  $\nabla f$  is not available it is not possible to decide *a priori* if the direction  $d_k$  is a descent direction, hence it is necessary to consider also negative values of  $\alpha$ .

We now describe the simplest line search method that can be constructed with the considered hypothesis. This method is a modification of Armijo method and it is known as parabolic search.

Given  $\lambda_0 > 0$ ,  $\sigma \in (0, 1/2)$ ,  $\gamma > 0$  and  $\rho \in (0, 1)$ . Compute  $\alpha_k$  and  $\lambda_k$  such that one of the following conditions hold.

Condition (i)

- $\lambda_k = \lambda_{k-1}$ ;
- $\alpha_k$  is the largest value in the set

$$A = \{\alpha \in \mathbb{R} : \alpha = \pm \sigma^j, j = 0, 1, \dots\}$$

such that

$$f(x_k + \alpha_k d_k) \leq f(x_k) - \gamma \alpha_k^2,$$

or, equivalently,  $\phi(\alpha_k) \leq -\gamma \alpha_k^2$ .

Condition (ii)

- $\alpha_k = 0$ ,  $\lambda_k \leq \rho \lambda_{k-1}$ ;
- $\min(f(x_k + \lambda_k d_k), f(x_k - \lambda_k d_k)) \geq f(x_k) - \gamma \lambda_k^2$ .

At each step it is necessary to satisfy either Condition (i) or Condition (ii). Note that this is always possible for any  $d_k \neq 0$ . Condition (i) requires that  $\alpha_k$  is the largest number in the set  $A$  such that  $f(x_k + \alpha_k d_k)$  is below the parabola  $f(x_k) - \gamma \alpha^2$ . If the function  $\phi(\alpha)$  has a stationary point for  $\alpha = 0$  then there may be no  $\alpha \in A$  such that Condition (i) holds. However, in this case it is possible to find  $\lambda_k$  such that Condition (ii) holds. If Condition (ii) holds then  $\alpha_k = 0$ , i.e. the point  $x_k$  remains unchanged and the algorithm continues with a new direction  $d_{k+1} \neq d_k$ .

For the parabolic search algorithm it is possible to prove the following convergence result.

**Theorem 7** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and assume  $\nabla f$  is continuous and  $\mathcal{L}(f(x_0))$  is compact. If  $\alpha_k$  is selected following the conditions of the parabolic search and if  $x_{k+1} = x_k + \alpha_k d_k$ , with  $\|d_k\| = 1$  then the sequence  $\{x_k\}$  is such that*

$$f(x_{k+1}) \leq f(x_k)$$

for all  $k$ ,

$$\lim_{k \rightarrow \infty} \nabla f(x_k)' d_k = 0$$

and

$$\lim_{k \rightarrow \infty} \|x_{k+1} - x_k\| = 0.$$

*Proof.* (Sketch) Note that Condition (i) implies  $f(x_{k+1}) < f(x_k)$ , whereas Condition (ii) implies  $f(x_{k+1}) = f(x_k)$ . Note now that if Condition (ii) holds for all  $k \geq \bar{k}$ , then  $\alpha_k = 0$  for all  $k \geq \bar{k}$ , i.e.  $\|x_{k+1} - x_k\| = 0$ . Moreover, as  $\lambda_k$  is reduced at each step, necessarily  $\nabla f(x_{\bar{k}})' \bar{d} = 0$ , where  $\bar{d}$  is a limit of the sequence  $\{d_k\}$ .  $\triangleleft$

### 2.4.5 Implementation of a line search algorithm

On the basis of the conditions described so far it is possible to construct algorithms that yield  $\alpha_k$  in a finite number of steps. One such an algorithm can be described as follows. (For simplicity we assume that  $\nabla f$  is known.)

- Initial data.  $x_k, f(x_k), \nabla f(x_k), \underline{\alpha}$  and  $\bar{\alpha}$ .
- Initial guess for  $\alpha$ . A possibility is to select  $\alpha$  as the point in which a parabola through  $(0, \phi(0))$  with derivative  $\dot{\phi}(0)$  for  $\alpha = 0$  takes a pre-specified minimum value  $f_*$ . Initially, *i.e.* for  $k = 0$ ,  $f_*$  has to be selected by the designer. For  $k > 0$  it is possible to select  $f_*$  such that

$$f(x_k) - f_* = f(x_{k-1}) - f(x_k).$$

The resulting  $\alpha$  is

$$\alpha_* = -2 \frac{f(x_k) - f_*}{\nabla f(x_k)' d_k}.$$

In some algorithms it is convenient to select  $\alpha \leq 1$ , hence the initial guess for  $\alpha$  will be  $\min(1, \alpha_*)$ .

- Computation of  $\alpha_k$ . A value for  $\alpha_k$  is computed using a line search method. If  $\alpha_k \leq \underline{\alpha}$  the direction  $d_k$  may not be a descent direction. If  $\alpha_k \geq \bar{\alpha}$  the level set  $\mathcal{L}(f(x_k))$  may not be compact. If  $\alpha_k \notin [\underline{\alpha}, \bar{\alpha}]$  the line search fails, and it is necessary to select a new research direction  $d_k$ . Otherwise the line search terminates and  $x_{k+1} = x_k + \alpha_k d_k$ .

## 2.5 The gradient method

The gradient method consists in selecting, as research direction, the direction of the anti-gradient at  $x_k$ , *i.e.*

$$d_k = -\nabla f(x_k),$$

for all  $k$ . This selection is justified noting that the direction<sup>6</sup>

$$-\frac{\nabla f(x_k)}{\|\nabla f(x_k)\|_E}$$

is the direction that minimizes the directional derivative, among all direction with unitary Euclidean norm. In fact, by Schwartz inequality, one has

$$|\nabla f(x_k)' d| \leq \|d\|_E \|\nabla f(x_k)\|_E,$$

and the equality sign holds if and only if  $d = \lambda \nabla f(x_k)$ , with  $\lambda \in \mathbb{R}$ . As a consequence, the problem

$$\min_{\|d\|_E=1} \nabla f(x_k)' d$$

---

<sup>6</sup>We denote with  $\|v\|_E$  the Euclidean norm of the vector  $v$ , *i.e.*  $\|v\|_E = \sqrt{v'v}$ .

has the solution  $d_\star = -\frac{\nabla f(x_k)}{\|\nabla f(x_k)\|_E}$ . For this reason, the gradient method is sometimes called the method of the steepest descent. Note however that the (local) optimality of the direction  $-\nabla f(x_k)$  depends upon the selection of the norm, and that with a proper selection of the norm, any descent direction can be regarded as the steepest descent. The real interest in the direction  $-\nabla f(x_k)$  rests on the fact that, if  $\nabla f$  is continuous, then the former is a continuous descent direction, which is zero only if the gradient is zero, *i.e.* at a stationary point.

The gradient algorithm can be schematized as follows.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute  $\nabla f(x_k)$ . If  $\nabla f(x_k) = 0$  STOP. Else set  $d_k = -\nabla f(x_k)$ .

**Step 3.** Compute a step  $\alpha_k$  along the direction  $d_k$  with any line search method such that

$$f(x_k + \alpha_k d_k) \leq f(x_k)$$

and

$$\lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} = 0.$$

**Step 4.** Set  $x_{k+1} = x_k + \alpha_k d_k$ ,  $k = k + 1$ . Go to **Step 2**.

By the general results established in Theorem 4, we have the following fact regarding the convergence properties of the gradient method.

**Theorem 8** *Consider  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ . Assume  $\nabla f$  is continuous and the level set  $\mathcal{L}(f(x_0))$  is compact. Then any accumulation point of the sequence  $\{x_k\}$  generated by the gradient algorithm is a stationary point of  $f$ .*

To estimate the speed of convergence of the method we can consider the behavior of the method in the minimization of a quadratic function, *i.e.* in the case

$$f(x) = \frac{1}{2} x' Q x + c' x + d,$$

with  $Q = Q' > 0$ . In such a case it is possible to obtain the following estimate

$$\|x_{k+1} - x_\star\| \leq \sqrt{\frac{\lambda_M}{\lambda_m}} \frac{\sqrt{\frac{\lambda_M}{\lambda_m} - 1}}{\sqrt{\frac{\lambda_M}{\lambda_m} + 1}} \|x_k - x_\star\|,$$

where  $\lambda_M \geq \lambda_m > 0$  are the maximum and minimum eigenvalue of  $Q$ , respectively. Note that the above estimate is exact for some initial points  $x_0$ . As a result, if  $\lambda_M \neq \lambda_m$  the gradient algorithm has linear convergence, however, if  $\lambda_M/\lambda_m$  is large the convergence can be very slow (see Figure 2.5).

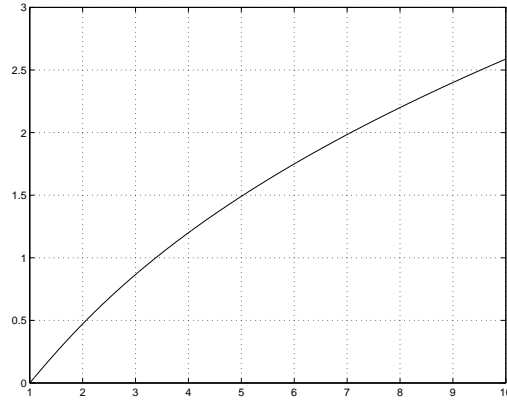


Figure 2.5: The function  $\sqrt{\xi} \frac{\xi - 1}{\xi + 1}$ .

Finally, if  $\lambda_M/\lambda_m = 1$  the gradient algorithm converges in one step. From a geometric point of view the ratio  $\lambda_M/\lambda_m$  expresses the ratio between the lengths of the maximum and the minimum axes of the ellipsoids, that constitute the level surfaces of  $f$ . If this ratio is big there are points from which the gradient algorithm converges very slowly, see *e.g.* Figure 2.6.

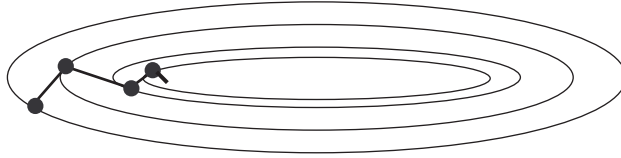


Figure 2.6: Behavior of the gradient algorithm.

In the non-quadratic case, the performance of the gradient method are unacceptable, especially if the level surfaces of  $f$  have high curvature.

## 2.6 Newton's method

Newton's method, with all its variations, is the most important method in unconstrained optimization. Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  be a given function and assume that  $\nabla^2 f$  is continuous. Newton's method for the minimization of  $f$  can be derived assuming that, given  $x_k$ , the point  $x_{k+1}$  is obtained minimizing a quadratic approximation of  $f$ . As  $f$  is two times differentiable, it is possible to write

$$f(x_k + s) = f(x_k) + \nabla f(x_k)'s + \frac{1}{2}s'\nabla^2 f(x_k)s + \beta(x_k, s),$$

in which

$$\lim_{\|s\| \rightarrow 0} \frac{\beta(x_k, s)}{\|s\|^2} = 0.$$

For  $\|s\|$  sufficiently small, it is possible to approximate  $f(x_k + s)$  with its quadratic approximation

$$q(s) = f(x_k) + \nabla f(x_k)'s + \frac{1}{2}s'\nabla^2 f(x_k)s.$$

If  $\nabla^2 f(x_k) > 0$ , the value of  $s$  minimizing  $q(s)$  can be obtained setting to zero the gradient of  $q(s)$ , *i.e.*

$$\nabla q(s) = \nabla f(x_k) + \nabla^2 f(x_k)s = 0,$$

yielding

$$s = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k).$$

The point  $x_{k+1}$  is thus given by

$$x_{k+1} = x_k - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k).$$

Finally, Newton's method can be described by the simple scheme.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute

$$s = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k).$$

**Step 3.** Set  $x_{k+1} = x_k + s$ ,  $k = k + 1$ . Go to **Step 2**.

*Remark.* An equivalent way to introduce Newton's method for unconstrained optimization is to regard the method as an algorithm for the solution of the system of  $n$  non-linear equations in  $n$  unknowns given by

$$\nabla f(x) = 0.$$

For, consider, in general, a system of  $n$  equations in  $n$  unknown

$$F(x) = 0,$$

with  $x \in \mathbb{R}^n$  and  $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ . If the Jacobian matrix of  $F$  exists and is continuous, then one can write

$$F(x + s) = F(x) + \frac{\partial F}{\partial x}(x)s + \gamma(x, s),$$

with

$$\lim_{\|s\| \rightarrow 0} \frac{\gamma(x, s)}{\|s\|} = 0.$$

Hence, given a point  $x_k$  we can determine  $x_{k+1} = x_k + s$  setting  $s$  such that

$$F(x_k) + \frac{\partial F}{\partial x}(x_k)s = 0.$$

If  $\frac{\partial F}{\partial x}(x_k)$  is invertible we have

$$s = - \left[ \frac{\partial F}{\partial x}(x_k) \right]^{-1} F(x_k),$$

hence Newton's method for the solution of the system of equation  $F(x) = 0$  is

$$x_{k+1} = x_k - \left[ \frac{\partial F}{\partial x}(x_k) \right]^{-1} F(x_k), \quad (2.3)$$

with  $k = 0, 1, \dots$ . Note that, if  $F(x) = \nabla f$ , then the above iteration coincides with Newton's method for the minimization of  $f$ .  $\diamond$

To study the convergence properties of Newton's method we can consider the algorithm for the solution of a set of non-linear equations, summarized in equation (2.3). The following local convergence result, providing also an estimate of the speed of convergence, can be proved.

**Theorem 9** *Let  $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$  and assume that  $F$  is continuously differentiable in an open set  $\mathcal{D} \subset \mathbb{R}^n$ . Suppose moreover that*

- *there exists  $x_\star \in \mathcal{D}$  such that  $F(x_\star) = 0$ ;*
- *the Jacobian matrix  $\frac{\partial F}{\partial x}(x_\star)$  is non-singular;*
- *there exists  $L > 0$  such that<sup>7</sup>*

$$\left\| \frac{\partial F}{\partial x}(z) - \frac{\partial F}{\partial x}(y) \right\| \leq L \|z - y\|,$$

*for all  $z \in \mathcal{D}$  and  $y \in \mathcal{D}$ .*

*Then there exists an open set  $\mathcal{B} \subset \mathcal{D}$  such that for any  $x_0 \in \mathcal{B}$  the sequence  $\{x_k\}$  generated by equation (2.3) remains in  $\mathcal{B}$  and converges to  $x_\star$  with quadratic speed of convergence.*

The result in Theorem 9 can be easily recast as a result for the convergence of Newton's method for unconstrained optimization. For, it is enough to note that all hypotheses on  $F$  and  $\frac{\partial F}{\partial x}$  translate into hypotheses on  $\nabla f$  and  $\nabla^2 f$ . Note however that the result is only local and does not allow to distinguish between local minimizers and local maximizers. To construct an algorithm for which the sequence  $\{x_k\}$  does not converge to maxima, and for which global convergence, *i.e.* convergence from points outside the set  $\mathcal{B}$ , holds,

---

<sup>7</sup>This is equivalent to say that  $\frac{\partial F}{\partial x}(x)$  is Lipschitz continuous in  $\mathcal{D}$ .

it is possible to modify Newton's method considering a line search along the direction  $d_k = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$ . As a result, the modified Newton's algorithm

$$x_{k+1} = x_k - \alpha_k [\nabla^2 f(x_k)]^{-1} \nabla f(x_k), \quad (2.4)$$

in which  $\alpha_k$  is computed using any line search algorithm, is obtained. If  $\nabla^2 f$  is uniformly positive definite, and this implies that the function  $f$  is convex, the direction  $d_k = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$  is a descent direction satisfying the condition of angle. Hence, by Theorem 4, we can conclude the (global) convergence of the algorithm (2.4). Moreover, it is possible to prove that, for  $k$  sufficiently large, the step  $\alpha_k = 1$  satisfies the conditions of Armijo method, hence the sequence  $\{x_k\}$  has quadratic speed of convergence.

*Remark.* If the function to be minimized is quadratic, *i.e.*

$$f(x) = \frac{1}{2} x' Q x + c' x + d,$$

and if  $Q > 0$ , Newton's method yields the (global) minimizer of  $f$  in one step.  $\diamond$

In general, *i.e.* if  $\nabla^2 f(x)$  is not positive definite for all  $x$ , Newton's method may be inapplicable because either  $\nabla^2 f(x_k)$  is not invertible, or  $d_k = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$  is not a descent direction. In these cases it is necessary to further modify Newton's method. Diverse criteria have been proposed, most of which rely on the substitution of the matrix  $\nabla^2 f(x_k)$  with a matrix  $M_k > 0$  which is *close in some sense* to  $\nabla^2 f(x_k)$ . A simpler modification can be obtained using the direction  $d_k = -\nabla f(x_k)$  whenever the direction  $d_k = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$  is not a descent direction. This modification yields the following algorithm.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$  and  $\epsilon > 0$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute  $\nabla f(x_k)$ . If  $\nabla f(x_k) = 0$  STOP. Else compute  $\nabla^2 f(x_k)$ . If  $\nabla^2 f(x_k)$  is singular set  $d_k = -\nabla f(x_k)$  and go to **Step 6**.

**Step 3.** Compute Newton direction  $s$  solving the (linear) system

$$\nabla^2 f(x_k) s = -\nabla f(x_k).$$

**Step 4.** If

$$|\nabla f(x_k)' s| < \epsilon \|\nabla f(x_k)\| \|s\|$$

set  $d_k = -\nabla f(x_k)$  and go to **Step 6**.

**Step 5.** If

$$\nabla f(x_k)' s < 0$$

set  $d_k = s$ ; if

$$\nabla f(x_k)' s > 0$$

set  $d_k = -s$ .



**Step 6.** Make a line search along  $d_k$  assuming as initial estimate  $\alpha = 1$ . Compute  $x_{k+1} = x_k + \alpha_k d_k$ , set  $k = k + 1$  and go to **Step 2**.

The above algorithm is such that the direction  $d_k$  satisfies the condition of angle, *i.e.*

$$\nabla f(x_k)' d_k \leq -\epsilon \|\nabla f(x_k)\| \|d_k\|,$$

for all  $k$ . Hence, the convergence is guaranteed by the general result in Theorem 4. Moreover, if  $\epsilon$  is sufficiently small, if the hypotheses of Theorem 9 hold, and if the line search is performed with Armijo method and with the initial guess  $\alpha = 1$ , then the above algorithm has quadratic speed of convergence.

Finally, note that it is possible to modify Newton's method, whenever it is not applicable, without making use of the direction of the anti-gradient. We now briefly discuss two such modifications.

### 2.6.1 Method of the trust region

A possible approach to modify Newton's method to yield global convergence is to set the direction  $d_k$  and the step  $\alpha_k$  in such a way to minimize the quadratic approximation of  $f$  on a sphere centered at  $x_k$  and of radius  $a_k$ . Such a sphere is called *trust region*. This name refers to the fact that, in a small region around  $x_k$  we are confident (we trust) that the quadratic approximation of  $f$  is a *good* approximation.

The method of the trust region consists in selecting  $x_{k+1} = x_k + s_k$ , where  $s_k$  is the solution of the problem

$$\min_{\|s\| \leq a_k} q(s), \quad (2.5)$$

with

$$q(s) = f(x_k) + \nabla f(x_k)' s + \frac{1}{2} s' \nabla^2 f(x_k) s,$$

and  $a_k > 0$  the estimate at step  $k$  of the trust region. As the above (constrained) optimization problem has always a solution, the direction  $s_k$  is always defined. The computation of the estimate  $a_k$  is done, iteratively, in such a way to enforce the condition  $f(x_{k+1}) < f(x_k)$  and to make sure that  $f(x_k + s_k) \approx q(s_k)$ , *i.e.* that the change of  $f$  and the estimated change of  $f$  are *close*.

Using these simple ingredients it is possible to construct the following algorithm.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$  and  $a_0 > 0$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute  $\nabla f(x_k)$ . If  $\nabla f(x_k) = 0$  STOP. Else go to **Step 3**.

**Step 3.** Compute  $s_k$  solving problem (2.5).

**Step 4.** Compute<sup>8</sup>

$$\rho_k = \frac{f(x_k + s_k) - f(x_k)}{q(s_k) - f(x_k)}. \quad (2.6)$$

---

<sup>8</sup>If  $f$  is quadratic then  $\rho_k = 1$  for all  $k$ .

**Step 5.** If  $\rho_k < 1/4$  set  $a_{k+1} = \|s_k\|/4$ . If  $\rho_k > 3/4$  and  $\|s_k\| = a_k$  set  $a_{k+1} = 2a_k$ . Else set  $a_{k+1} = a_k$ .

**Step 6.** If  $\rho_k \leq 0$  set  $x_{k+1} = x_k$ . Else set  $x_{k+1} = x_k + s_k$ .

**Step 7.** Set  $k = k + 1$  and go to **Step 2**.

*Remark.* Equation (2.6) expresses the ratio between the actual change of  $f$  and the estimated change of  $f$ .  $\diamond$

It is possible to prove that, if  $\mathcal{L}(f(x_0))$  is compact and  $\nabla^2 f$  is continuous, any accumulation point resulting from the above algorithm is a stationary point of  $f$ , in which the second order necessary conditions hold.

The update of  $a_k$  is devised to enlarge or shrink the region of confidence on the basis of the number  $\rho_k$ . It is possible to show that if  $\{x_k\}$  converges to a local minimizer in which  $\nabla^2 f$  is positive definite, then  $\rho_k$  converges to one and the direction  $s_k$  coincides, for  $k$  sufficiently large, with the Newton direction. As a result, the method has quadratic speed of convergence.

In practice, the solution of the problem (2.5) cannot be obtained analytically, hence approximate problems have to be solved. For, consider  $s_k$  as the solution of the equation

$$(\nabla^2 f(x_k) + \nu_k I) s_k = -\nabla f(x_k), \quad (2.7)$$

in which  $\nu_k > 0$  has to be determined with proper considerations. Under certain hypotheses, the  $s_k$  determined solving equation (2.7) coincides with the  $s_k$  computed using the method of the trust region.

*Remark.* A potential disadvantage of the method of the trust region is to reduce the step along Newton direction even if the selection  $\alpha_k = 1$  would be feasible.  $\diamond$

## 2.6.2 Non-monotonic line search

Experimental evidence shows that Newton's method gives the best result if the step  $\alpha_k = 1$  is used. Therefore, the use of  $\alpha_k < 1$  along Newton direction, resulting *e.g.* from the application of Armijo method, results in a degradation of the performance of the algorithm. To avoid this phenomenon it has been suggested to relax the condition  $f(x_{k+1}) < f(x_k)$  imposed on Newton algorithm, thus allowing the function  $f$  to increase for a certain number of steps. For example, it is possible to substitute the *reduction* condition of Armijo method with the condition

$$f(x_k + \alpha_k d_k) \leq \max_{0 \leq j \leq M} [f(x_{k-j})] + \gamma \alpha_k \nabla f(x_k)' d_k$$

for all  $k \geq M$ , where  $M > 0$  is a fixed integer independent from  $k$ .

	Gradient method	Newton's method
Information required at each step	$f$ and $\nabla f$	$f$ , $\nabla f$ and $\nabla^2 f$
Computation to find the research direction	$\nabla f(x_k)$	$\nabla f(x_k)$ , $\nabla^2 f(x_k)$ , $-[\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$
Convergence	Global if $\mathcal{L}(f(x_0))$ compact and $\nabla f$ continuous	Local, but may be rendered global
Behavior for quadratic functions	Asymptotic convergence	Convergence in one step
Speed of convergence	Linear for quadratic functions	Quadratic (under proper hypotheses)

Table 2.1: Comparison between the gradient method and Newton's method.

### 2.6.3 Comparison between Newton's method and the gradient method

The gradient method and Newton's method can be compared from different point of views, as described in Table 2.1. From the table, it is obvious that Newton's method has better convergence properties but it is computationally more expensive. There exist methods which preserve some of the advantages of Newton's method, namely speed of convergence faster than the speed of the gradient method and finite convergence for quadratic functions, without requiring the knowledge of  $\nabla^2 f$ . Such methods are

- the conjugate directions methods;
- quasi-Newton methods.

## 2.7 Conjugate directions methods

Conjugate directions methods have been motivated by the need of improving the convergence speed of the gradient method, without requiring the computation of  $\nabla^2 f$ , as required in Newton's method.

A basic characteristic of conjugate directions methods is to find the minimizer of a quadratic function in a finite number of steps. These methods have been introduced for the solution of systems of linear equations and have later been extended to the solution of unconstrained optimization problems for non-quadratic functions.

**Definition 5** Given a matrix  $Q = Q'$ , the vectors  $d_1$  and  $d_2$  are said to be  $Q$ -conjugate if

$$d_1' Q d_2 = 0.$$

*Remark.* If  $Q = I$  then two vectors are  $Q$ -conjugate if they are orthogonal.  $\diamond$

**Theorem 10** Let  $Q \in \mathbb{R}^{n \times n}$  and  $Q = Q' > 0$ . Let  $d_i \in \mathbb{R}^n$ , for  $i = 0, \dots, k$ , be non-zero vectors. If  $d_i$  are mutually  $Q$ -conjugate, i.e.

$$d_i' Q d_j = 0,$$

for all  $i \neq j$ , then the vectors  $d_i$  are linearly independent.

*Proof.* Suppose there exists constants  $\alpha_i$ , with  $\alpha_i \neq 0$  for some  $i$ , such that

$$\alpha_0 d_0 + \dots + \alpha_k d_k = 0.$$

Then, left multiplying with  $Q$  and  $d_j'$  yields

$$\alpha_j d_j' Q d_j = 0,$$

which implies, as  $Q > 0$ ,  $\alpha_j = 0$ . Repeating the same considerations for all  $j \in [0, k]$  yields the claim.  $\triangleleft$

Consider now a quadratic function

$$f(x) = \frac{1}{2} x' Q x + c' x + d,$$

with  $x \in \mathbb{R}^n$  and  $Q = Q' > 0$ . The (global) minimizer of  $f$  is given by

$$x_\star = -Q^{-1}c,$$

and this can be computed using the procedure given in the next statement.

**Theorem 11** Let  $Q = Q' > 0$  and let  $d_0, d_1, \dots, d_{n-1}$  be  $n$  non-zero vectors mutually  $Q$ -conjugate. Consider the algorithm

$$x_{k+1} = x_k + \alpha_k d_k$$

with

$$\alpha_k = -\frac{\nabla f(x_k)' d_k}{d_k' Q d_k} = -\frac{(x_k' Q + c') d_k}{d_k' Q d_k}.$$

Then, for any  $x_0$ , the sequence  $\{x_k\}$  converges, in at most  $n$  steps, to  $x_\star = -Q^{-1}c$ , i.e. it converges to the minimizer of the quadratic function  $f$ .

*Remark.* Note that  $\alpha_k$  is selected at each step to minimize the function  $f(x_k + \alpha d_k)$  with respect to  $\alpha$ , i.e. at each step an exact line search in the direction  $d_k$  is performed.  $\diamond$

In the above statement we have assumed that the directions  $d_k$  have been preliminarily assigned. However, it is possible to construct a procedure in which the directions are computed iteratively. For, consider the quadratic function  $f(x) = \frac{1}{2}x'Qx + c'x + d$ , with  $Q > 0$ , and the following algorithm, known as conjugate gradient method.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$  and the direction

$$d_0 = -\nabla f(x_0) = -(Qx_0 + c).$$

**Step 1.** Set  $k = 0$ .

**Step 2.** Let

$$x_{k+1} = x_k + \alpha_k d_k$$

with

$$\alpha_k = -\frac{\nabla f(x_k)'d_k}{d_k'Qd_k} - \frac{(x_k'Q + c')d_k}{d_k'Qd_k}.$$

**Step 3.** Compute  $d_{k+1}$  as follows

$$d_{k+1} = -\nabla f(x_{k+1}) + \beta_k d_k,$$

with

$$\beta_k = \frac{\nabla f(x_{k+1})'Qd_k}{d_k'Qd_k}.$$

**Step 4.** Set  $k = k + 1$  and go to **Step 2**.

*Remark.* As already observed,  $\alpha_k$  is selected to minimize the function  $f(x_k + \alpha d_k)$ . Moreover, this selection of  $\alpha_k$  is also such that

$$\nabla f(x_{k+1})'d_k = 0. \quad (2.8)$$

In fact,

$$Qx_{k+1} = Qx_k + \alpha_k Qd_k$$

hence

$$\nabla f(x_{k+1}) = \nabla f(x_k) + \alpha_k Qd_k. \quad (2.9)$$

Left multiplying with  $d_k'$  yields

$$d_k' \nabla f(x_{k+1}) = d_k' \nabla f(x_k) + d_k' Qd_k \alpha_k = d_k' \nabla f(x_k) - d_k' Qd_k \frac{\nabla f(x_k)'d_k}{d_k' Qd_k} = 0.$$

$\diamond$

*Remark.*  $\beta_k$  is such that  $d_{k+1}$  is  $Q$ -conjugate with respect to  $d_k$ . In fact,

$$d'_k Q d_{k+1} = d'_k Q \left( -\nabla f(x_{k+1}) + \frac{\nabla f(x_{k+1})' Q d_k}{d'_k Q d_k} d_k \right) = d'_k Q (-\nabla f(x_{k+1}) + \nabla f(x_{k+1})) = 0.$$

Moreover, this selection of  $\beta_k$  yields also

$$\nabla f(x_k)' d_k = -\nabla f(x_k)' \nabla f(x_k). \quad (2.10)$$

◇

For the conjugate gradient method it is possible to prove the following fact.

**Theorem 12** *The conjugate gradient method yields the minimizer of the quadratic function*

$$f(x) = \frac{1}{2} x' Q x + c' x + d,$$

with  $Q = Q' > 0$ , in at most  $n$  iterations, i.e. there exists  $m \leq n - 1$  such that

$$\nabla f(x_{m+1}) = 0.$$

Moreover

$$\nabla f(x_j)' \nabla f(x_i) = 0 \quad (2.11)$$

and

$$d'_j Q d_i = 0, \quad (2.12)$$

for all  $[0, m+1] \ni i \neq j \in [0, m+1]$ .

*Proof.* To prove the (finite) convergence of the sequence  $\{x_k\}$  it is enough to show that the directions  $d_k$  are  $Q$ -conjugate, i.e. that equation (2.12) holds. In fact, if equation (2.12) holds the claim is a consequence of Theorem 11. ◁

The conjugate gradient algorithm, in the form described above, cannot be used for the minimization of non-quadratic functions, as it requires the knowledge of the matrix  $Q$ , which is the Hessian of the function  $f$ . Note that the matrix  $Q$  appears at two levels in the algorithm: in the computation of the scalar  $\beta_k$  required to compute the new direction of research, and in the computation of the step  $\alpha_k$ . It is therefore necessary to modify the algorithm to avoid the computation of  $\nabla^2 f$ , but at the same time it is reasonable to make sure that the modified algorithm coincides with the above one in the quadratic case.

### 2.7.1 Modification of $\beta_k$

To begin with note that, by equation (2.9),  $\beta_k$  can be written as

$$\beta_k = \frac{\nabla f(x_{k+1})' \frac{\nabla f(x_{k+1}) - \nabla f(x_k)}{\alpha_k}}{d'_k \frac{\nabla f(x_{k+1}) - \nabla f(x_k)}{\alpha_k}} = \frac{\nabla f(x_{k+1})' [\nabla f(x_{k+1}) - \nabla f(x_k)]}{d'_k [\nabla f(x_{k+1}) - \nabla f(x_k)]},$$

and, by equation (2.8),

$$\beta_k = -\frac{\nabla f(x_{k+1})' [\nabla f(x_{k+1}) - \nabla f(x_k)]}{d_k' \nabla f(x_k)}. \quad (2.13)$$

Using equation (2.13), it is possible to construct several expressions for  $\beta_k$ , all equivalent in the quadratic case, but yielding different algorithms in the general (non-quadratic) case. A first possibility is to consider equations (2.10) and (2.11) and to define

$$\beta_k = \frac{\nabla f(x_{k+1})' \nabla f(x_{k+1})}{\nabla f(x_k)' \nabla f(x_k)} = \frac{\|\nabla f(x_{k+1})\|^2}{\|\nabla f(x_k)\|^2}, \quad (2.14)$$

which is known as Fletcher-Reeves formula.

A second possibility is to write the denominator as in equation (2.14) and the numerator as in equation (2.13), yielding

$$\beta_k = \frac{\nabla f(x_{k+1})' [\nabla f(x_{k+1}) - \nabla f(x_k)]}{\|\nabla f(x_k)\|^2}, \quad (2.15)$$

which is known as Polak-Ribiere formula. Finally, it is possible to have the denominator as in (2.13) and the numerator as in (2.14), *i.e.*

$$\beta_k = -\frac{\|\nabla f(x_{k+1})\|^2}{d_k' \nabla f(x_k)}. \quad (2.16)$$

### 2.7.2 Modification of $\alpha_k$

As already observed, in the quadratic version of the conjugate gradient method also the step  $\alpha_k$  depends upon  $Q$ . However, instead of using the  $\alpha_k$  given in **Step 2** of the algorithm, it is possible to use a line search along the direction  $\alpha_k$ . In this way, an algorithm for non-quadratic functions can be constructed. Note that  $\alpha_k$ , in the algorithm for quadratic functions, is also such that  $d_k' \nabla f(x_{k+1}) = 0$ . Therefore, in the line search, it is reasonable to select  $\alpha_k$  such that, not only  $f(x_{k+1}) < f(x_k)$ , but also  $d_k$  is approximately orthogonal to  $\nabla f(x_{k+1})$ .

*Remark.* The condition of approximate orthogonality between  $d_k$  and  $\nabla f(x_{k+1})$  cannot be enforced using Armijo method or Goldstein conditions. However, there are more sophisticated line search algorithms, known as Wolfe conditions, which allow to enforce the above constraint.  $\diamond$

### 2.7.3 Polak-Ribiere algorithm

As a result of the modifications discussed in the last sections, it is possible to construct an algorithm for the minimization of general functions. For example, using equation (2.15) we obtain the following algorithm, due to Polak-Ribiere, which has proved to be one of the most efficient among the class of conjugate directions methods.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute  $\nabla f(x_k)$ . If  $\nabla f(x_k) = 0$  STOP. Else let

$$d_k = \begin{cases} -\nabla f(x_0), & \text{if } k = 0 \\ -\nabla f(x_k) + \frac{\nabla f(x_k)' [\nabla f(x_k) - \nabla f(x_{k-1})]}{\|\nabla f(x_{k-1})\|^2} d_{k-1}, & \text{if } k \geq 1. \end{cases}$$

**Step 3.** Compute  $\alpha_k$  performing a line search along  $d_k$ .

**Step 4.** Set  $x_{k+1} = x_k + \alpha_k d_k$ ,  $k = k + 1$  and go to **Step 2**.

*Remark.* The line search has to be sufficiently accurate, to make sure that all directions generated by the algorithm are descent directions. A suitable line search algorithm is the so-called Wolfe method, which is a modification of Goldstein method.  $\diamond$

*Remark.* To guarantee global convergence of a subsequence it is possible to use, every  $n$  steps, the direction  $-\nabla f$ . In this case, it is said that the algorithm uses a *restart* procedure. For the algorithm with restart it is possible to have quadratic speed of convergence in  $n$  steps, *i.e.*

$$\|x_{k+n} - x_\star\| \leq \gamma \|x_k - x_\star\|^2,$$

for some  $\gamma > 0$ .  $\diamond$

*Remark.* It is possible to modify Polak-Ribiere algorithm to make sure that at each step the angle condition holds. In this case, whenever the direction  $d_k$  does not satisfy the angle condition, it is sufficient to use the direction  $-\nabla f$ . Note that, enforcing the angle condition, yields a globally convergent algorithm.  $\diamond$

*Remark.* Even if the use of the direction  $-\nabla f$  every  $n$  steps, or whenever the angle condition is not satisfied, allows to prove global convergence of Polak-Ribiere algorithm, it has been observed in numerical experiments that such modified algorithms do not perform as well as the original one.  $\diamond$

## 2.8 Quasi-Newton methods

Conjugate gradient methods have proved to be more efficient than the gradient method. However, in general, it is not possible to guarantee superlinear convergence. The main advantage of conjugate gradient methods is in the fact that they do not require to construct and store any matrix, hence can be used in large scale problems.



In small and medium scale problems, *i.e.* problems with less than a few hundreds decision variables, in which  $\nabla^2 f$  is not available, it is convenient to use the so-called quasi-Newton methods.

Quasi Newton methods, as conjugate directions methods, have been introduced for quadratic functions. They are described by an algorithm of the form

$$x_{k+1} = x_k - \alpha_k H_k \nabla f(x_k),$$

with  $H_0$  given. The matrix  $H_k$  is an approximation of  $[\nabla^2 f(x_k)]^{-1}$  and it is computed iteratively at each step.

If  $f$  is a quadratic function, the gradient of  $f$  is given by

$$\nabla f(x) = Qx + c,$$

for some  $Q$  and  $c$ , hence for any  $x \in \mathbb{R}^n$  and  $y \in \mathbb{R}^n$  one has

$$\nabla f(y) - \nabla f(x) = Q(y - x),$$

or, equivalently,

$$Q^{-1}[\nabla f(y) - \nabla f(x)] = y - x.$$

It is then natural, in general, to construct the sequence  $\{H_k\}$  such that

$$H_{k+1}[\nabla f(x_{k+1}) - \nabla f(x_k)] = x_{k+1} - x_k. \quad (2.17)$$

Equation (2.17) is known as quasi-Newton equation.

There exist several update methods satisfying the quasi-Newton equation. For simplicity, set

$$\gamma_k = \nabla f(x_{k+1}) - \nabla f(x_k),$$

and

$$\delta_k = x_{k+1} - x_k.$$

As a result, equation (2.17) can be rewritten as

$$H_{k+1}\gamma_k = \delta_k.$$

One of the first quasi-Newton methods has been proposed by Davidon, Fletcher and Powell, and can be summarized by the equations

$$\text{DFP} \begin{cases} H_0 &= I \\ H_{k+1} &= H_k + \frac{\delta_k \delta'_k}{\delta'_k \gamma_k} - \frac{H_k \gamma_k \gamma'_k H_k}{\gamma'_k H_k \gamma_k}. \end{cases} \quad (2.18)$$

It is easy to show that the matrix  $H_{k+1}$  satisfies the quasi-Newton equation (2.17), *i.e.*

$$\begin{aligned} H_{k+1}\gamma_k &= H_k \gamma_k + \frac{\delta_k \delta'_k}{\delta'_k \gamma_k} \gamma_k - \frac{H_k \gamma_k \gamma'_k H_k}{\gamma'_k H_k \gamma_k} \gamma_k \\ &= H_k \gamma_k + \frac{\delta'_k \gamma_k}{\delta'_k \gamma_k} \delta_k - \frac{\gamma'_k H_k \gamma_k}{\gamma'_k H_k \gamma_k} H_k \gamma_k \\ &= \delta_k. \end{aligned}$$

Moreover, it is possible to prove the following fact, which gives conditions such that the matrices generated by DFP method are positive definite for all  $k$ .

**Theorem 13** *Let  $H_k = H'_k > 0$  and assume  $\delta'_k \gamma_k > 0$ . Then the matrix*

$$H_k + \frac{\delta_k \delta'_k}{\delta'_k \gamma_k} - \frac{H_k \gamma_k \gamma'_k H_k}{\gamma'_k H_k \gamma_k}$$

*is positive definite.*

DFP method has the following properties. In the quadratic case, if  $\alpha_k$  is selected to minimize

$$f(x_k - \alpha H_k \nabla f(x_k)),$$

then

- the directions  $d_k = -H_k \nabla f(x_k)$  are mutually conjugate;
- the minimizer of the (quadratic) function is found in at most  $n$  steps, moreover  $H_n = Q^{-1}$ ;
- the matrices  $H_k$  are always positive definite.

In the non-quadratic case

- the matrices  $H_k$  are positive definite (hence  $d_k = -H_k \nabla f(x_k)$  is a descent direction) if  $\delta'_k \gamma_k > 0$ ;
- it is globally convergent if  $f$  is strictly convex and if the line search is exact;
- it has superlinear speed of convergence (under proper hypotheses).

A second, and more general, class of update formulae, including as a particular case DFP formula, is the so-called Broyden class, defined by the equations

$$\text{Broyden} \begin{cases} H_0 &= I \\ H_{k+1} &= H_k + \frac{\delta_k \delta'_k}{\delta'_k \gamma_k} - \frac{H_k \gamma_k \gamma'_k H_k}{\gamma'_k H_k \gamma_k} + \phi v_k v'_k, \end{cases} \quad (2.19)$$

with  $\phi \geq 0$  and

$$v_k = (\gamma'_k H_k \gamma_k)^{1/2} \left( \frac{\delta_k}{\delta'_k \gamma_k} - \frac{H_k \gamma_k}{\gamma'_k H_k \gamma_k} \right).$$

If  $\phi = 0$  then we obtain DFP formula, whereas for  $\phi = 1$  we have the so-called Broyden-Fletcher-Goldfarb-Shanno (BFGS) formula, which is one of the preferred algorithms in applications. From Theorem 13 it is easy to infer that, if  $H_0 > 0$ ,  $\gamma'_k \delta_k > 0$  and  $\phi \geq 0$ , then all formulae in the class of Broyden generate matrices  $H_k > 0$ .

*Remark.* Note that the condition  $\delta'_k \gamma_k > 0$  is equivalent to

$$(\nabla f(x_{k+1}) - \nabla f(x_k))' d_k > 0,$$

and this can be enforced with a sufficiently precise line search.  $\diamond$

For the method based on BFGS formula, a global convergence result, for convex functions and in the case of non-exact (but sufficiently accurate) line search, has been proved. Moreover, it has been shown that the algorithm has superlinear speed of convergence. This algorithm can be summarized as follows.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Compute  $\nabla f(x_k)$ . If  $\nabla f(x_k) = 0$  STOP. Else compute  $H_k$  with BFGS equation and set

$$d_k = -H_k \nabla f(x_k).$$

**Step 3.** Compute  $\alpha_k$  performing a line search along  $d_k$ .

**Step 4.** Set  $x_{k+1} = x_k + \alpha_k d_k$ ,  $k = k + 1$  and go to **Step 2**.

In the general case it is not possible to prove global convergence of the algorithm. However, this can be enforced verifying (at the end of **Step 2**), if the direction  $d_k$  satisfies an angle condition, and if not use the direction  $d_k = -\nabla f(x_k)$ . However, as already observed, this modification improves the convergence properties, but reduces (sometimes drastically) the speed of convergence.

## 2.9 Methods without derivatives

All the algorithms that have been discussed presuppose the knowledge of the derivatives (first and/or second) of the function  $f$ . There are, however, also methods which do not require such a knowledge. These methods can be divided in two classes: direct search methods and methods using finite difference approximations.

Direct search methods are based upon the direct comparison of the values of the function  $f$  in the points generated by the algorithm, without making use of the necessary condition of optimality  $\nabla f = 0$ . In this class, the most interesting methods, *i.e.* the methods for which it is possible to give theoretical results, are those that make use cyclically of  $n$  linearly independent directions. The simplest possible method, known as the method of the coordinate directions, can be described by the following algorithm.

**Step 0.** Given  $x_0 \in \mathbb{R}^n$ .

**Step 1.** Set  $k = 0$ .

**Step 2.** Set  $j = 1$ .

**Step 3.** Set  $d_k = e_j$ , where  $e_j$  is the  $j$ -th coordinate direction.

**Step 4.** Compute  $\alpha_k$  performing a line search without derivatives along  $d_k$ .

**Step 5.** Set  $x_{k+1} = x_k + \alpha_k d_k$ ,  $k = k + 1$ .

**Step 6.** If  $j < n$  set  $j = j + 1$  and go to **Step 3**. If  $j = n$  go to **Step 2**.

It is easy to verify that the matrix

$$P_k = \begin{bmatrix} d_k & d_{k+1} & \cdots & d_{k+n-1} \end{bmatrix}$$

is such that

$$|\det P_k| = 1,$$

hence, if the line search is such that

$$\lim_{k \rightarrow \infty} \frac{\nabla f(x_k)' d_k}{\|d_k\|} = 0$$

and

$$\lim_{k \rightarrow \infty} \|x_{k+1} - x_k\| = 0,$$

convergence to stationary points is ensured by the general result in Theorem 5. Note that, the line search can be performed using the parabolic line search method described in Section 2.4.4.

The method of the coordinate directions is not very efficient, in terms of speed of convergence. Therefore, a series of heuristics have been proposed to improve its performance. One such heuristics is the so-called method of Jeeves and Hooke, in which not only the search along the coordinate directions is performed, but also a search along directions joining pairs of points generated by the algorithm. In this way, the search is performed along what may be considered to be the most promising directions.

An alternative direct search method is the so-called simplex method (which should not be confused with the simplex method of linear programming). The method starts with  $n + 1$  (equally spaced) points  $x_{(i)} \in \mathbb{R}^n$  (these points give a simplex in  $\mathbb{R}^n$ ). In each of these points the function  $f$  is computed and the vertex where the function  $f$  attains the maximum value is determined. Suppose this is the vertex  $x_{(n+1)}$ . This vertex is reflected with respect to the center of the simplex, *i.e.* the point

$$x_c = \frac{1}{n+1} \sum_{i=1}^{n+1} x_{(i)}.$$

As a result, the new vertex

$$x_{(n+2)} = x_c + \alpha(x_c - x_{(n+1)})$$

where  $\alpha > 0$ , is constructed, see Figure 2.7. The procedure is then repeated.

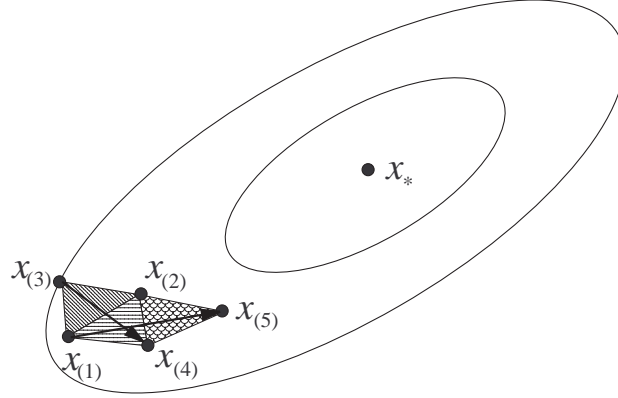


Figure 2.7: The simplex method. The points  $x_{(1)}$ ,  $x_{(2)}$  and  $x_{(3)}$  yields the starting simplex. The second simplex is given by the points  $x_{(1)}$ ,  $x_{(2)}$  and  $x_{(4)}$ . The third simplex is given by the points  $x_{(2)}$ ,  $x_{(4)}$  and  $x_{(5)}$ .

It is possible that the vertex that is generated by one step of the algorithm is (again) the one where the function  $f$  has its maximum. In this case, the algorithm cycles, hence the next vertex has to be determined using a different strategy. For example, it is possible to construct the next vertex by reflecting another of the remaining  $n$  vertex, or to shrink the simplex.

As a stopping criterion it is possible to consider the condition

$$\frac{1}{n+1} \sum_{i=1}^{n+1} \left( f(x_{(i)}) - \bar{f} \right)^2 < \epsilon \quad (2.20)$$

where  $\epsilon > 0$  is assigned by the designer, and

$$\bar{f} = \frac{1}{n+1} \sum_{i=1}^{n+1} f(x_{(i)}),$$

*i.e.*  $\bar{f}$  is the mean value of the  $f(x_{(i)})$ . Condition (2.20) implies that the points  $x_{(i)}$  are all in a region where the function  $f$  is *flat*.

As already observed, direct search methods are not very efficient, and can be used only for problems with a few decision variables and when approximate solutions are acceptable. As an alternative, if the derivatives of the function  $f$  are not available, it is possible to resort to numeric differentiation, *e.g.* the entries of the gradient of  $f$  can be computed using the so-called forward difference approximation, *i.e.*

$$\frac{\partial f(x)}{\partial x_i} \approx \frac{f(x + te_i) - f(x)}{t},$$

where  $e_i$  is the  $i$ -th column of the identity matrix of dimension  $n$ , and  $t > 0$  has to be fixed by the user. Note that there are methods for the computation of the optimal value of  $t$ , *i.e.* the value of  $t$  which minimizes the approximation error.

## 2.10 Exercises

This section contains a set of exercises related to the notions, concepts, algorithms and tools discussed in Chapter 2, together with exercises on topics which have not been covered in the book. In these cases, the specific topic is briefly illustrated in the text of the exercise. The objective is to draw the reader's attention to the fact that there are many more ideas, methods and algorithms which have been developed to solve unconstrained optimization problems and which are not covered in the book. The basic principles provided in Chapter 2 should however be sufficient to understand more advanced and involved methods. Not surprisingly, a significant number of exercises is devoted to Newton's method and its modifications: undoubtedly, Newton's method is one of the most important methods in optimization (and numerical analysis). All exercises have a brief worked out solution, which provides guidelines and checkpoints to help the reader assess their level of understanding and familiarity with the content covered. Note that the exercises are not ordered in any particular way: the order is the result of the history of my optimization course and of the exam papers I have set over the years.

**Exercise 1** Consider the problem of minimizing the function

$$f(x) = x_1^2 + x_2^2 - x_1.$$

- Compute the unique stationary point of the function, and show that the function is radially unbounded.
- Using second order sufficient conditions show that the stationary point determined in part a) is a local minimizer. Also show that the point is a global minimizer.
- Consider the minimization of the function  $f$  using the gradient algorithm. Express analytically the form of the generic iteration, *i.e.*

$$p_{k+1} = p_k - \alpha \nabla f,$$

where  $p_i = [x_1^i, x_2^i]'$ .

- Consider the initial point  $p_0 = [1, 1]'$  and apply one step of the gradient algorithm from part c) with exact line search. Verify that  $p_1$  coincides with the stationary point determined in part a).
- It is known that for quadratic functions, such as the function  $f$  above, the gradient algorithm is globally convergent, however the speed of convergence may be very slow. Discuss why, for the function  $f$ , the gradient algorithm with exact line search converges in one step.

### Solution 1

- The stationary points of the function  $f$  are computed solving the equation

$$0 = \nabla f = \begin{bmatrix} 2x_1 - 1 \\ 2x_2 \end{bmatrix},$$

yielding the unique stationary point  $x_1^* = 1/2$  and  $x_2^* = 0$ . The function  $f$  is of the form  $x'Qx + c'x$  with  $Q = \text{diag}(1, 1) > 0$ , hence it is radially unbounded.

- b) Note that  $\nabla^2 f = \text{diag}(2, 2) > 0$ , hence  $x^*$  is a local minimizer. It is also a global minimizer for the following reasons: the function  $f$  is  $C^1$  and radially unbounded, therefore the global minimizer is a stationary point.
- c) The generic iteration of the gradient algorithm for the considered function  $f$  is
- $$x_1^{k+1} = x_1^k - \alpha(2x_1^k - 1) \quad x_2^{k+1} = x_2^k - \alpha(2x_2^k).$$
- d) Let  $x_1^0 = x_2^0 = 1$ . Hence
- $$x_1^1 = 1 - \alpha \quad x_2^1 = 1 - 2\alpha.$$
- Note now that  $f(x_1^1, x_2^1) = 1 - 5\alpha + 5\alpha^2$  and this is minimized by  $\alpha^* = 1/2$ , yielding
- $$x_1^1 = 1 - \alpha^* = 1/2 = x_1^* \quad x_2^1 = 1 - 2\alpha^* = 0 = x_2^*.$$
- e) For the considered function the gradient algorithm with exact line search converges in one step (from any initial point) because the function is quadratic and the minimum and maximum eigenvalues of  $\nabla^2 f$  coincide.

**Exercise 2** Consider the problem of minimizing the function

$$f(x) = x_1^4 + x_1 x_2 + \frac{1}{2} x_2^2.$$

- a) Compute the stationary points of the function.
- b) Using second order sufficient conditions *classify* the stationary points determined in part a), *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point.
- c) Consider the minimization of the function  $f$  using Newton's algorithm. Express analytically the form of the generic iteration, *i.e.*

$$p_{k+1} = p_k - [\nabla^2 f]^{-1} \nabla f,$$

where  $p_i = [x_1^i, x_2^i]'$ .

- d) The equation in part c) defines a nonlinear discrete-time system with equilibria coinciding with the stationary points of the function  $f$ .

Consider the linear approximation of the system in part c) around the equilibrium corresponding to the local minimizer of the function  $f$  with  $x_1 > 0$ , and compute the eigenvalues associated with the linear approximation.

Interpret the result obtained in terms of convergence properties of sequences generated by Newton's algorithm and initialized close to a local minimizer.

- e) Consider the initial point

$$p_0 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

and, using the results in part c), apply four steps of Newton's algorithm to generate the points  $p_1, p_2, p_3, p_4$ . Comment on the speed of convergence of the sequence.

**Solution 2**

- a) The stationary point of the function  $f$  are computed solving the equation

$$0 = \nabla f = \begin{bmatrix} 4x_1^3 + x_2 \\ x_1 + x_2 \end{bmatrix},$$

yielding the stationary points

$$p^* = \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \quad \tilde{p}^* = \begin{bmatrix} -1/2 \\ 1/2 \end{bmatrix}, \quad \hat{p}^* = \begin{bmatrix} 1/2 \\ -1/2 \end{bmatrix}.$$

b) Note that

$$\nabla^2 f(p^*) = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \not\geq 0 \quad \nabla^2 f(\tilde{p}^*) = \nabla^2 f(\hat{p}^*) = \begin{bmatrix} 3 & 1 \\ 1 & 1 \end{bmatrix} > 0.$$

Hence,  $p^*$  is a saddle point, and  $\tilde{p}^*$  and  $\hat{p}^*$  are local minimizers.

c) The generic iteration of Newton's algorithm for the considered function  $f$  is

$$x_1^{k+1} = \frac{8(x_1^k)^3}{12(x_1^k)^2 - 1}, \quad x_2^{k+1} = -\frac{8(x_1^k)^3}{12(x_1^k)^2 - 1}.$$

d) The linear approximation of the above nonlinear discrete-time system around the point  $\hat{p}^*$  is the system

$$p_{k+1} = Ap_k = \begin{bmatrix} 0 & 0 \\ 0 & 0 \end{bmatrix} p_k.$$

This linear system is such that  $p_1 = 0$  for any  $p_0$ , and this explains the local 'fast' speed of convergence of Newton's iteration.

e) A simple computation yields the sequence

$$p_0 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \quad p_1 = \begin{bmatrix} 0.7272727273 \\ -0.7272727273 \end{bmatrix}, \quad p_2 = \begin{bmatrix} 0.57552339460 \\ -0.57552339460 \end{bmatrix},$$

$$p_3 = \begin{bmatrix} 0.51266296461 \\ -0.51266296461 \end{bmatrix}, \quad p_4 = \begin{bmatrix} 0.5004542259 \\ -0.5004542259 \end{bmatrix},$$

and this shows the fast convergence (approximately two exact digits for each iteration) of Newton's algorithm.

**Exercise 3** Consider the function

$$f(x) = x_1^2 + x_1x_2 + (x_1 - x_2)^4.$$

- Compute all stationary points of the function.  
(Hint: obtain first  $(x_1 - x_2)^3$  in terms of  $x_1$  from the necessary conditions of optimality.)
- Using second order sufficient conditions, *classify* the stationary points determined in part a), *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point.
- Consider the point  $p = (0, 0)$  and the direction  $d = \begin{bmatrix} 1 \\ -3 \end{bmatrix}$ . Using the definition of a descent direction, show that  $d$  is a descent direction for  $f$  at  $p$ .
- Perform an exact line search along the direction  $d = \begin{bmatrix} 1 \\ -3 \end{bmatrix}$  starting at  $p = (0, 0)$ . Show that the point obtained as a result of the line search procedure is a local minimizer of the function  $f$ .

**Solution 3**

a) The stationary points of the function  $f$  are computed by solving the equation

$$0 = \nabla f = \begin{bmatrix} 2x_1 + x_2 + 4(x_1 - x_2)^3 \\ x_1 - 4(x_1 - x_2)^3 \end{bmatrix},$$

yielding

$$P_1 = (0, 0), \quad P_2 = \left(-\frac{1}{16}, \frac{3}{16}\right), \quad P_3 = \left(\frac{1}{16}, -\frac{3}{16}\right).$$



- b) Note that

$$\nabla^2 f = \begin{bmatrix} 2 + 12(x_1 - x_2)^2 & 1 - 12(x_1 - x_2)^2 \\ 1 - 12(x_1 - x_2)^2 & 12(x_1 - x_2)^2 \end{bmatrix}.$$

Thus

$$\nabla^2 f(P_1) = \begin{bmatrix} 2 & 1 \\ 1 & 0 \end{bmatrix},$$

which is an indefinite matrix, and

$$\nabla^2 f(P_2) = \nabla^2 f(P_3) = \begin{bmatrix} 11/4 & 1/4 \\ 1/4 & 3/4 \end{bmatrix},$$

which is a positive definite matrix. As a result,  $P_1$  is a saddle point, and  $P_2$  and  $P_3$  are local minimizers.

- c) By definition, a direction
- $d$
- is a descent direction for
- $f$
- at
- $p$
- if there exists
- $\delta > 0$
- such that

$$f(p + \lambda d) < f(p),$$

for all  $\lambda \in (0, \delta)$ . Consider now the given direction and note that  $f(p) = 0$  and that

$$f(p + \lambda d) = -2\lambda^2 + 256\lambda^4.$$

Hence,  $f(p + \lambda d) < f(p)$  for all  $\lambda > 0$  and sufficiently small. (Note that  $\nabla f(p)'d = 0$ , hence this condition cannot be used to decide if  $d$  is a descent direction, or otherwise.)

- d) To perform an exact line search along the direction
- $d$
- , starting from
- $p = (0, 0)$
- , we need to find the minimum of the function

$$\phi(\lambda) = f(p + \lambda d) - f(p) = -2\lambda^2 + 256\lambda^4$$

for  $\lambda > 0$ . Note that

$$\frac{d\phi}{d\lambda} = -4\lambda + 1024\lambda^3,$$

hence the minimum is achieved for  $\lambda = 1/16$ . The resulting point is  $(1/16, -3/16)$  and this coincides with one of the local minimizers determined in part a).

**Exercise 4** Consider the problem of minimizing the function

$$f(x) = x_1^2 + 2x_2^2 + 4x_1 + 4x_2.$$

- a) Compute the stationary points of the function.  
 b) Consider the minimization of the function  $f$  using the gradient algorithm. Express analytically the form of the generic iteration, *i.e.*

$$p_{k+1} = p_k - \alpha \nabla f,$$

where  $p_i = [x_1^i, x_2^i]'$ .

- c) Compute three steps of the gradient algorithm with exact line search from the initial point  $p_0 = [0, 0]'$ , using the fact that, for this  $p_0$  the exact line search parameter  $\alpha$  is equal to  $1/3$  for all  $k$ . Check that indeed  $\alpha^* = 1/3$  for the first iteration.  
 d) Exploit the results of part c) to show that the gradient iteration with exact line search for  $p_0 = [0, 0]'$  gives

$$x_1^{k+1} = \frac{1}{3}x_1^k - \frac{4}{3},$$

$$x_2^{k+1} = -\frac{1}{3}x_2^k - \frac{4}{3},$$

and hence show that

$$(x_1^{k+1} + 2) = \frac{1}{3}(x_1^k + 2),$$

$$(x_2^{k+1} + 1) = -\frac{1}{3}(x_2^k + 1).$$

Hence, deduce that the sequence  $\{p_k\}$  can be written as

$$p_{k+1} = \begin{bmatrix} \frac{2}{3^{k+1}} - 2 \\ \left(-\frac{1}{3}\right)^{k+1} - 1 \end{bmatrix}.$$

Show that the sequence  $\{p_k\}$  converges to the stationary point determined in part a).

#### Solution 4

- a) The stationary points of the function  $f$  are computed by solving the equation

$$0 = \nabla f = \begin{bmatrix} 2x_1 + 4 \\ 4x_2 + 4 \end{bmatrix},$$

yielding the stationary point

$$p^* = (-2, 1).$$

- b) The generic iteration of the gradient algorithm for the considered function  $f$  is

$$x_1^{k+1} = x_1^k - \alpha(2x_1^k + 4), \quad x_2^{k+1} = x_2^k - \alpha(4x_2^k + 4).$$

- c) Setting  $(x_1^0, x_2^0) = (0, 0)$  one has  $(x_1^1, x_2^1) = (-4\alpha, -4\alpha)$  and

$$f(-4\alpha, -4\alpha) - f(0, 0) = 48\alpha^2 - 32\alpha.$$

Minimizing this function yields  $\alpha^* = 1/3$  (as stated). Therefore,  $(x_1^1, x_2^1) = (-4/3, -4/3)$ .

Repeating the same considerations, and setting always  $\alpha = 1/3$ , one has

$$(x_1^2, x_2^2) = (-16/9, -8/9)$$

and

$$(x_1^3, x_2^3) = (-52/27, -28/27).$$

- d) Setting  $\alpha = 1/3$  in the gradient iteration yields

$$x_1^{k+1} = \frac{1}{3}x_1^k - \frac{4}{3}, \quad x_2^{k+1} = -\frac{1}{3}x_2^k - \frac{4}{3},$$

and this can be also written as

$$(x_1^{k+1} + 2) = \frac{1}{3}(x_1^k + 2), \quad (x_2^{k+1} + 1) = -\frac{1}{3}(x_2^k + 1).$$

As a result

$$(x_1^{k+1} + 2) = \left(\frac{1}{3}\right)^k (x_1^0 + 2), \quad (x_2^{k+1} + 1) = \left(-\frac{1}{3}\right)^k (x_2^0 + 1),$$

or, equivalently,

$$x_1^{k+1} = 2\left(\frac{1}{3}\right)^k - 2, \quad x_2^{k+1} = \left(-\frac{1}{3}\right)^k - 1.$$

Finally, as  $k \rightarrow \infty$   $x_1^k \rightarrow -2$  and  $x_2^k \rightarrow -1$ , *i.e.* the sequence converges to the stationary point determined in part a).

#### Exercise 5

- a) An electrical engineer wants to maximize the current  $I$  between two points A and B of a complex network by adjusting the values  $x_1$  and  $x_2$  of two variable resistors. The engineer does not have a model of the network and decides to opt for this procedure.

- Keep the value  $x_2$  fixed and adjust  $x_1$  to maximize  $I$ .
- Keep the value  $x_1$  fixed and adjust  $x_2$  to maximize  $I$ .
- Repeat the above steps until no further improvement can be obtained.

Explain if this approach has sound theoretical basis, *i.e.* discuss under what assumptions the above procedure determines a stationary point of the function  $I$ .

- b) Let  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  be a differentiable function. Suppose that  $x_*$  is a local minimizer of  $f$  along every line that passes through  $x_*$ , *i.e.* the function

$$g(\alpha) = f(x_* + \alpha d)$$

is minimized at  $\alpha = 0$  for all  $d \in \mathbb{R}^n$ .

- Show that  $\nabla f(x_*) = 0$ .
- Is  $x_*$  a local minimizer of  $f$ ?
- Consider the function

$$f(x_1, x_2) = (x_2 - x_1^2)(x_2 - 2x_1^2).$$

Show that the point  $(0, 0)$  is a local minimizer of  $f$  along every line that passes through  $(0, 0)$ .

Show that the point  $(0, 0)$  is not a local minimizer of  $f$ .

(Hint: consider the values of  $f$  for  $x_1 = y$  and  $x_2 = my^2$  and  $m \in \mathbb{R}$ .)

#### Solution 5

- a) The engineer is applying the so-called coordinate directions method with an exact line search (without derivatives), as described in Section 2.9, for the minimization of the function  $-I = -I(x_1, x_2)$ . This approach provides a sequence of points converging to a stationary point of the function  $I$  provided that the initial point is selected inside a compact level set of  $-I(x_1, x_2)$ .

- b) i) Note that, by assumption, the function

$$\frac{dg}{d\alpha} = \nabla f(x_* + \alpha d)' d$$

is zero for  $\alpha = 0$  and for every  $d$ . This means that

$$\nabla f(x_*)' d = 0$$

for every  $d$ , and this implies that

$$\nabla f(x_*) = 0.$$

- Without further information on  $f$  it is not possible to draw any conclusion on  $x_*$ , *i.e.*  $x_*$  is a stationary point of  $f$ , but it may be a local minimizer, a local maximizer or a saddle point.
- Consider a line that goes through zero, namely  $x_2 = \gamma x_1$ , and note that

$$f(x_1, \gamma x_1) = (\gamma x_1 - x_1^2)(\gamma x_1 - 2x_1^2) = \gamma^2 x_1^2 - 3\gamma x_1^3 + 2x_1^4$$

and this shows that for all  $\gamma$  the point  $x_1 = 0$  is a local minimizer of  $f(x_1, \gamma x_1)$ . For completeness we have also to consider the line  $x_1 = 0$  (which corresponds formally to  $\gamma = \infty$ ). Note that

$$f(0, x_2) = x_2^2$$

hence the point  $x_2 = 0$  is a minimizer of  $f(0, x_2)$ .

To show that  $(0, 0)$  is not a local minimizer of  $f$  note first that  $f(0, 0) = 0$  and then let  $x_1 = y$  and  $x_2 = my^2$ . Note that

$$f(y, my^2) = y^4(m - 1)(m - 2).$$

Pick  $m \in (1, 2)$  and note that for such values of  $m$

$$f(y, my^2) = y^4(m - 1)(m - 2) < 0$$

for all  $y \neq 0$ . This shows that close to the point  $(0, 0)$ , where the function is zero, there are points in which the function takes negative values. Hence,  $(0, 0)$  is not a local minimizer of  $f$ .

**Exercise 6** Consider the problem of minimizing the function

$$f(x_1, x_2) = \frac{1}{3}x_1^2 - \alpha x_1^4 + \frac{1}{4}x_1^6 + x_1x_2 + x_2^2,$$

where  $\alpha$  is a constant.

- Compute all stationary points of the function.
- Let  $\alpha = 5/12$ . Using second order sufficient conditions *classify* the stationary points determined in part a), *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point.
- Let  $\alpha = 5/12$ . Show that the function  $f$  is radially unbounded and hence compute the global minimum of  $f$ . Is the global minimizer unique?

**Solution 6**

- The stationary points of the function  $f$  are computed by solving the equation

$$0 = \nabla f = \begin{bmatrix} 2/3x_1 - 4\alpha x_1^3 + 3/2x_1^5 + x_2 \\ x_1 + 2x_2 \end{bmatrix}.$$

Solving the second equation for  $x_2$  yields  $x_2 = -1/2 x_1$ , and upon replacement in the first equation we obtain

$$\frac{1}{6}x_1 - 4\alpha x_1^3 + \frac{3}{2}x_1^5 = 0,$$

yielding

$$x_{1a} = 0, \quad x_{1b} = \frac{1}{3}\sqrt{12\alpha + 3\sqrt{16\alpha^2 - 1}}, \quad x_{1c} = -\frac{1}{3}\sqrt{12\alpha + 3\sqrt{16\alpha^2 - 1}},$$

$$x_{1d} = \frac{1}{3}\sqrt{12\alpha - 3\sqrt{16\alpha^2 - 1}}, \quad x_{1e} = -\frac{1}{3}\sqrt{12\alpha - 3\sqrt{16\alpha^2 - 1}}.$$

- For  $\alpha = 5/12$  we obtain the stationary points

$$P_a = (0, 0), \quad P_b = (1, -1/2), \quad P_c = (-1, 1/2),$$

$$P_d = (1/3, -1/6), \quad P_e = (-1/3, 1/6).$$

Note now that, for  $\alpha = 5/12$ ,

$$\nabla^2 f = \begin{bmatrix} 2/3 - 5x_1^2 + 15/2x_1^4 & 1 \\ 1 & 2 \end{bmatrix}$$

and that

$$\nabla^2 f(P_a) = \begin{bmatrix} 2/3 & 1 \\ 1 & 2 \end{bmatrix} > 0$$

$$\nabla^2 f(P_b) = \nabla^2 f(P_c) = \begin{bmatrix} 19/6 & 1 \\ 1 & 2 \end{bmatrix} > 0$$

$$\nabla^2 f(P_d) = \nabla^2 f(P_e) = \begin{bmatrix} 11/54 & 1 \\ 1 & 2 \end{bmatrix} \not> 0.$$

As a result,  $P_a$ ,  $P_b$  and  $P_c$  are local minimizers, and  $P_d$  and  $P_e$  are saddle points.

- Note that

$$-\frac{5}{12}x_1^4 + \frac{1}{4}x_1^6 = x_1^4 \left( \frac{1}{4}x_1^2 - \frac{5}{12} \right)$$

is radially unbounded. Hence

$$f(x_1, x_2) = \left( \frac{1}{3}x_1^2 + x_1x_2 + x_2^2 \right) + x_1^4 \left( \frac{1}{4}x_1^2 - \frac{5}{12} \right)$$

is also radially unbounded. The global minimum of  $f$  is also a local minimum of  $f$ . Note that

$$f(P_a) = 0 \quad f(P_b) = f(P_c) = -0.833 \dots$$

Hence,  $P_b$  and  $P_c$  are both global minimizers, therefore the global minimizer is not unique.

**Exercise 7** Consider the problem of minimizing the function

$$f(x) = x - \log x,$$

with  $x > 0$ .

- a) Compute analytically the minimizer of  $f$ .
- b) Write Newton's iteration for the considered problem.
- c) Consider the Newton's iteration in part b) with initial point  $x_0 = 1.99$ . Compute ten steps of the Newton's iteration. Argue that the resulting sequence converges to the minimizer of  $f$ . Show that the sequence converges to the minimizer of  $f$  with quadratic speed of convergence.
- d) Consider the Newton's iteration in part b) with initial point  $x_0 = 2.01$ . Compute five steps of the Newton's iteration. Argue that the resulting sequence diverges.
- e) Consider the Newton's iteration in part b). Show that
  - i) if the initial point  $x_0 = 2$  then  $x_k = 0$ , for all  $k \geq 1$ ;
  - ii) if the initial point  $x_0 = 0$  then  $x_k = 0$ , for all  $k \geq 1$ ;
  - iii) if the initial point  $x_0 > 2$  then  $x_k < 0$ , for all  $k \geq 1$  and the sequence does not converge;
  - iv) if the initial point  $x_0 \in (0, 2)$  then  $x_k \in (0, 2)$ , for all  $k \geq 1$  and the sequence converges to the minimizer determined in part a).

**Solution 7**

- a) The minimizer of  $f$  is obtained solving  $\nabla f = 1 - 1/x = 0$ , yielding  $x = 1$ . Note that  $x = 1$  is indeed a minimizer (a global one), because the function  $f$  is convex for all  $x > 0$ .
- b) The Newton's iteration is

$$x_{k+1} = x_k - \frac{1}{\nabla^2 f(x_k)} \nabla f(x_k) = x_k - x_k^2 \left(1 - \frac{1}{x_k}\right) = (2 - x_k)x_k.$$

- c) Let  $x_0 = 1.99$  then

$$\begin{array}{rcl} x_1 & = & 0.01990 \\ x_2 & = & 0.03940399 \\ x_3 & = & 0.07725530557208 \\ x_4 & = & 0.14854222890512 \\ x_5 & = & 0.27501966404215 \\ x_6 & = & 0.47440351247444 \\ x_7 & = & 0.72374833230079 \\ x_8 & = & 0.92368501609341 \\ x_9 & = & 0.99417602323134 \\ x_{10} & = & 0.99996608129460. \end{array}$$

The sequence is converging to  $x = 1$ , *i.e.* to the local minimizer of  $f$ . To establish quadratic speed of convergence note that

$$\frac{\mathcal{E}_{k+1}}{\mathcal{E}_k^2} = \frac{|x_{k+1} - 1|}{|x_k - 1|^2} = \frac{|(2 - x_k)x_k - 1|}{(x_k - 1)^2} = 1.$$

d) Let  $x_0 = 2.01$  then

$$\begin{aligned} x_1 &= -0.02010 \\ x_2 &= -0.04060401 \\ x_3 &= -0.08285670562808 \\ x_4 &= -0.17257864492369 \\ x_5 &= -0.37494067853109. \end{aligned}$$

We then infer that the sequence is monotonically decreasing and  $\lim_{k \rightarrow \infty} x_k = -\infty$ .

e) The first two points are trivial noting that

$$x_{k+1} = (2 - x_k)x_k$$

and the right hand side of this equation is zero for  $x_k = 0$  or  $x_k = 2$ , that is  $x = 0$  and  $x = 2$  are equilibria of the above discrete-time system. Note now that if  $x_0 > 2$  then  $x_1 < 0$ . Moreover if  $x_k < 0$  then

$$x_{k+1} = (2 - x_k)x_k < x_k,$$

which proves the third claim. Finally, if  $x_k \in (0, 2)$  then it is easy to verify that

$$0 < x_{k+1} = (2 - x_k)x_k < 2.$$

Moreover, if  $x_k = 1$  then  $x_{k+1} = 1$ , hence  $x = 1$  is an equilibrium of the discrete-time system  $x_{k+1} = (2 - x_k)x_k$ . Finally, if  $x_k \in (1, 2)$

$$0 < x_{k+1} < 1,$$

and if  $x_k \in (0, 1)$

$$x_k < x_{k+1} < 1,$$

which shows convergence of the sequence to  $x = 1$ .

**Exercise 8** Consider the problem of minimizing the function

$$f(x_1, x_2) = \frac{1}{2n+2} x_1^{2n+2} - x_1 x_2 + \frac{1}{2} x_2^2,$$

where  $n$  is a positive integer.

- Compute all stationary points of the function.
- Using second order sufficient conditions *classify* the stationary points determined in part a), *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point.
- Show that the function  $f$  is radially unbounded and hence compute the global minimum of  $f$ . Is the global minimizer unique?
- Consider the point  $P_0 = (0, 0)$  and the direction

$$d = \begin{bmatrix} 1 \\ 1 \end{bmatrix}.$$

Show that the direction  $d$  is a descent direction for  $f$  at  $P_0$ .

**Solution 8**

- The stationary points of the function  $f$  are computed by solving the equation

$$0 = \nabla f = \begin{bmatrix} x_1^{2n+1} - x_2 \\ -x_1 + x_2 \end{bmatrix}.$$

The second equation yields  $x_2 = x_1$ , hence the first equation becomes

$$0 = x_1^{2n+1} - x_1 = x_1(x_1^{2n} - 1).$$

The (real) solutions of this equation are  $x_1 = 0$ ,  $x_1 = 1$  and  $x_1 = -1$ . In summary, the function  $f$  has three stationary points

$$P_a = (0, 0), \quad P_b = (1, 1), \quad P_c = (-1, -1).$$

- b) Note that (recall that  $n$  is a positive integer)

$$\nabla^2 f = \begin{bmatrix} (2n+1)x_1^{2n} & -1 \\ -1 & 1 \end{bmatrix}.$$

Hence

$$\nabla^2 f(P_a) = \begin{bmatrix} 0 & -1 \\ -1 & 1 \end{bmatrix}$$

which is an indefinite matrix, and

$$\nabla^2 f(P_b) = \nabla^2 f(P_c) = \begin{bmatrix} 2n+1 & -1 \\ -1 & 1 \end{bmatrix} > 0.$$

As a result  $P_a$  is a saddle point, and  $P_b$  and  $P_c$  are local minimizers.

- c) Note that

$$f = \frac{1}{2n+2}x_1^{2n+2} - x_1x_2 + \frac{1}{2}x_2^2 = \frac{1}{2n+2}x_1^{2n+2} - x_1^2 + \left(x_1^2 - x_1x_2 + \frac{1}{2}x_2^2\right).$$

The function

$$\frac{1}{2n+2}x_1^{2n+2} - x_1^2 = x_1^2 \left( \frac{1}{2n+2}x_1^{2n} - 1 \right)$$

is radially unbounded, as a function of  $x_1$  alone, and the function  $x_1^2 - x_1x_2 + \frac{1}{2}x_2^2$  is radially unbounded as a function of  $x_1$  and  $x_2$ . As a result the global minimum of  $f$  is also a local minimum. Note that (recall again that  $n$  is a positive integer)

$$f(P_b) = f(P_c) = -\frac{1}{2} \frac{n}{n+1} < 0,$$

hence both  $P_b$  and  $P_c$  are global minimizers.

- d) The point  $P_0$  coincides with the saddle point  $P_a$ . The function  $f$  along the direction  $d$  is given by

$$\phi(\alpha) = f(\alpha, \alpha) = \frac{1}{2n+2}\alpha^{2n+2} - \frac{1}{2}\alpha^2.$$

Note that  $\phi(0) = 0$  and that  $\phi(\alpha) < 0$  for  $\alpha > 0$  and sufficiently small (namely for all  $\alpha \in \left(0, (n+1)^{\frac{1}{2n}}\right)$ ), hence  $d$  is a descent direction for  $f$  at  $P_0$ .

(Note that  $\phi(\alpha)$  is negative also for  $\alpha \in \left(-(n+1)^{\frac{1}{2n}}, 0\right)$ , i.e.  $-d$  is also a descent direction for  $f$  at  $P_0$ .)

**Exercise 9** Newton's method for the minimization of a function  $f : \mathbb{R} \rightarrow \mathbb{R}$  is based on a quadratic approximation of the function at a given point. An alternative way to construct a quadratic approximation that does not require the computation of the second derivative is to consider an approximation based on the knowledge of two points  $x_k$  and  $x_{k-1}$  and of the values  $f(x_k)$ ,  $\frac{df(x_k)}{dx}$  and  $\frac{df(x_{k-1})}{dx}$ . Such an approximation is given by

$$q(x) = f(x_k) + \frac{df(x_k)}{dx}(x - x_k) + \frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k} \frac{(x - x_k)^2}{2}.$$

- a) Show that the function  $q(x)$  is such that

$$q(x_k) = f(x_k), \quad \frac{dq(x_k)}{dx} = \frac{df(x_k)}{dx}, \quad \frac{dq(x_{k-1})}{dx} = \frac{df(x_{k-1})}{dx}.$$

- b) Compute the stationary point  $x_*$  of  $q(x)$ .  
 c) Consider the algorithm, known as the method of the false position, obtained by setting  $x_{k+1} = x_*$ , with  $x_*$  as in part b), and argue that this algorithm provides an approximation of Newton's method that does not require the computation of the second derivative of  $f$ .

- d) Show that the method of the false position applied to the minimization of a quadratic function  $f = ax^2 + bx + c$ , with  $a > 0$ , coincides with Newton's method.
- e) Consider the function  $f = \frac{x^4}{4} + x$ . This function has a global minimizer at  $x = -1$ .
- i) Show that the method of the false position yields the iteration

$$x_{k+1} = x_k - (x_k^3 + 1) \frac{1}{x_{k-1}^2 + x_{k-1}x_k + x_k^2}.$$

- ii) Evaluate

$$\frac{|\mathcal{E}_{k+1}|}{\mathcal{E}_k^2} = \frac{|x_{k+1} + 1|}{(x_k + 1)^2}$$

and show that if  $\lim_{k \rightarrow \infty} x_k = -1$  then

$$\lim_{k \rightarrow \infty} \frac{|\mathcal{E}_{k+1}|}{\mathcal{E}_k^2} = 1.$$

Hence, quantify the speed of convergence of the method.

#### Solution 9

- a) Setting  $x = x_k$  in  $q(x)$  yields  $q(x_k) = f(x_k)$ . Note that

$$\frac{dq(x)}{dx} = \frac{df(x_k)}{dx} + \frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k} (x - x_k)$$

hence, setting  $x = x_k$  and  $x = x_{k-1}$  yields

$$\frac{dq(x_k)}{dx} = \frac{df(x_k)}{dx}, \quad \frac{dq(x_{k-1})}{dx} = \frac{df(x_{k-1})}{dx}.$$

- b) The stationary point  $x_*$  of  $q(x)$  is obtained by solving the equation

$$\frac{dq(x)}{dx} = 0,$$

which yields

$$x_* = x_k - \left( \frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k} \right)^{-1} \frac{df(x_k)}{dx}.$$

- c) The method of the false position is therefore given by

$$x_{k+1} = x_k - \left( \frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k} \right)^{-1} \frac{df(x_k)}{dx}.$$

This algorithm is an approximation of Newton's method because the quantity

$$\frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k}$$

is an approximation of  $\frac{d^2f(x)}{dx^2}$  at  $x = x_k$ . Note however that, unlike Newton's method, the method of the false position does not need the computation of the second derivative: it uses an approximation.



d) For quadratic functions one has

$$\frac{d^2 f(x)}{dx^2} = 2a$$

and

$$\frac{\frac{df(x_{k-1})}{dx} - \frac{df(x_k)}{dx}}{x_{k-1} - x_k} = \frac{(2ax_{k-1} + b) - (2ax_k + b)}{x_{k-1} - x_k} = 2a,$$

hence, for such functions, Newton's method and the method of the false position coincide.

e) If  $f = \frac{x^4}{4} + x$  then  $\frac{df(x)}{dx} = x^3 + 1$ , and replacing in the expression of the considered method yields

$$x_{k+1} = x_k - \frac{x_{k-1} - x_k}{(x_{k-1}^3 + 1) - (x_k^3 + 1)}(x_k^3 + 1) = x_k - \frac{x_{k-1} - x_k}{x_{k-1}^3 - x_k^3}(x_k^3 + 1).$$

and, recalling that

$$x_{k-1}^3 - x_k^3 = (x_{k-1} - x_k)(x_{k-1}^2 + x_{k-1}x_k + x_k^2),$$

$$x_{k+1} = x_k - (x_k^3 + 1) \frac{1}{x_{k-1}^2 + x_{k-1}x_k + x_k^2}.$$

Note that

$$\begin{aligned} x_{k+1} + 1 &= x_k + 1 - (x_k^3 + 1) \frac{1}{x_{k-1}^2 + x_{k-1}x_k + x_k^2} \\ &= (x_k + 1)(x_{k-1} + 1) \frac{x_k + x_{k-1} - 1}{x_{k-1}^2 + x_{k-1}x_k + x_k^2}, \end{aligned}$$

hence

$$\frac{|\mathcal{E}_{k+1}|}{\mathcal{E}_k^2} = \left| \frac{x_{k-1} + 1}{x_k + 1} \frac{x_k + x_{k-1} - 1}{x_{k-1}^2 + x_{k-1}x_k + x_k^2} \right|.$$

If  $\lim_{k \rightarrow \infty} x_k = -1$  then also  $\lim_{k \rightarrow \infty} x_{k-1} = -1$ , hence  $\frac{|\mathcal{E}_{k+1}|}{\mathcal{E}_k^2} = 1$ , which shows that the algorithm has quadratic speed of convergence (if it converges).

**Exercise 10** Consider the problem of minimizing the function

$$f(x_1, x_2, \dots, x_n, y) = \frac{1}{4}x_1^4 + \frac{1}{4}x_2^4 + \dots + \frac{1}{4}x_n^4 - (x_1 + x_2 + \dots + x_n)y + \frac{n}{2}y^2,$$

where  $n$  is a positive integer.

- Compute all stationary points of the function.
- Using second order sufficient conditions *classify* the stationary points determined in part a), *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point.
- Show that the function  $f$  is radially unbounded and hence compute the global minimum of  $f$ . Is the global minimizer unique?
- Consider the points  $P_p = (1, 1, \dots, 1, 1)$  and  $P_m = (-1, -1, \dots, -1, -1)$  and the direction  $d$  from  $P_p$  to  $P_m$ . Show that this is an ascent direction for  $f$  at  $P_p$ .

**Solution 10**

- The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} x_1^3 - y \\ x_2^3 - y \\ \vdots \\ x_n^3 - y \\ -x_1 - x_2 - \dots - x_n + ny \end{bmatrix}.$$

The first  $n$  equations yield  $x_i = y^{1/3}$ , hence the last equation becomes

$$0 = -ny^{1/3} + ny = n(y - y^{1/3}).$$

The solutions of this equation are  $y = 0$ ,  $y = 1$  and  $y = -1$ . In summary, the function  $f$  has three stationary points

$$P_a = (0, \dots, 0, 0) \quad P_b = (1, \dots, 1, 1) \quad P_c = (-1, \dots, -1, -1).$$

b) Note that

$$\nabla^2 f = \begin{bmatrix} 3x_1^2 & 0 & \cdots & 0 & -1 \\ 0 & 3x_2^2 & \cdots & 0 & -1 \\ 0 & \vdots & \ddots & \vdots & \vdots \\ 0 & \cdots & 0 & 3x_n^2 & -1 \\ -1 & -1 & \cdots & -1 & n \end{bmatrix}.$$

Hence

$$\nabla^2 f(P_a) = \begin{bmatrix} 0 & 0 & \cdots & 0 & -1 \\ 0 & 0 & \cdots & 0 & -1 \\ 0 & \vdots & \ddots & \vdots & \vdots \\ 0 & \cdots & 0 & 0 & -1 \\ -1 & -1 & \cdots & -1 & n \end{bmatrix},$$

which is an indefinite matrix, hence  $P_a$  is a saddle point. Finally,

$$\nabla^2 f(P_b) = \nabla^2 f(P_c) = \begin{bmatrix} 3I & -v \\ -v' & n \end{bmatrix},$$

where  $v' = [1 \ \cdots \ 1]$ . Exploiting the relation

$$\begin{bmatrix} I & 0 \\ v'/3 & 1 \end{bmatrix} \begin{bmatrix} 3I & -v \\ -v' & n \end{bmatrix} \begin{bmatrix} I & v/3 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 3I & 0 \\ 0 & 2/3n \end{bmatrix},$$

we conclude that  $P_b$  and  $P_c$  are local minimizers.

c) The function  $f$  can be written as

$$f = \frac{1}{4}(x_1^2 - 1)^2 + \cdots + \frac{1}{4}(x_n^2 - 1)^2 + \frac{1}{2}(x_1 - y)^2 + \cdots + \frac{1}{2}(x_n - y)^2 - \frac{n}{4}.$$

Hence  $f + n/4$  is a *sum of squares*, and all variables  $x_1, x_2, \dots, x_n, y$  are present in one of the squares. As a result the function is radially unbounded and the local minimum of  $f$  is also a global minimum. Note that

$$f(P_b) = f(P_c) = -\frac{n}{4} < 0,$$

hence both  $P_b$  and  $P_c$  are global minimizers.

d) The direction from  $P_p$  to  $P_m$  is

$$d = P_m - P_p = -2 \begin{bmatrix} 1 \\ \vdots \\ 1 \\ 1 \end{bmatrix}.$$

The function  $f$  along the direction  $d$  at  $P_p$  is given by

$$\phi(\alpha) = f(1 - 2\alpha, \dots, 1 - 2\alpha, 1 - 2\alpha) = \frac{n}{4}(1 - 2\alpha)^4 - \frac{n}{2}(1 - 2\alpha^2) = -\frac{n}{4} + 4n\alpha^2 + \cdots$$

Note that  $\phi(0) = -n/4$  and that  $\phi(\alpha) > -n/4$  for  $\alpha > 0$  and sufficiently small, hence  $d$  is an ascent direction for  $f$  at  $P_p$ .

**Exercise 11** The problem of minimizing a differentiable function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  can be solved with the so-called heavy ball algorithm, which is a modification of the gradient algorithm, and it is described (in its simplest form) by the iteration

$$x_{k+1} = x_k - \alpha \nabla f(x_k) + \beta(x_k - x_{k-1}),$$

where  $\alpha > 0$  is a constant and  $\beta \in [0, 1)$  is the *heavy ball* parameter.

- a) Assume  $x_{-1} = x_0$ . Show that the iteration of the heavy ball algorithm can be written as

$$x_{k+1} = x_k - \alpha \left( \nabla f(x_k) + \beta \nabla f(x_{k-1}) + \beta^2 \nabla f(x_{k-2}) + \cdots + \beta^k \nabla f(x_0) \right),$$

for  $k \geq 1$ .

- b) Consider the function

$$f(x_1, x_2) = 2x_1^2 + \frac{1}{2}x_2^2,$$

which has a unique (global) minimizer at  $x_1 = x_2 = 0$ .

- i) Consider the heavy ball algorithm with  $\beta = 0$ , *i.e.* the gradient algorithm with a constant line search parameter  $\alpha$ . Show that the sequence  $\{x_k\} = \{(x_{k,1}, x_{k,2})\}$  generated by this algorithm converges to the minimizer of  $f$  if and only if  $0 < \alpha < \frac{1}{2}$ . Select  $\alpha = 1/4$ . Show that  $x_{k,1} = 0$ , for all  $k \geq 1$ , and determine the speed of convergence of the sequence  $\{x_k\}$ .
- ii) Consider the heavy ball algorithm described above with  $x_{-1} = x_0$ ,  $\alpha = 1/4$  and  $\beta = 3/4$ . Show that the sequence  $\{x_k\} = \{(x_{k,1}, x_{k,2})\}$  generated by this algorithm is such that  $x_{k,1} = 0$  for all  $k \geq 1$ . Evaluate  $x_{k,2}$  for  $k = 1, \dots, 4$ . Estimate the speed of convergence of the sequence  $\{x_k\}$ .

#### Solution 11

- a) Setting  $x_{-1} = x_0$  yields

$$k = 0 \quad \Rightarrow \quad x_1 = x_0 - \alpha \nabla f(x_0),$$

$$k = 1 \quad \Rightarrow \quad x_2 = x_1 - \alpha \nabla f(x_1) + \beta(x_1 - x_0) = x_1 - \alpha(\nabla f(x_1) + \beta \nabla f(x_0)),$$

$$k = 2 \quad \Rightarrow \quad x_3 = x_2 - \alpha \nabla f(x_2) + \beta(x_2 - x_1) = x_2 - \alpha(\nabla f(x_2) + \beta \nabla f(x_1) + \beta^2 \nabla f(x_0))$$

from which we deduce the general expression

$$x_{k+1} = x_k - \alpha \left( \nabla f(x_k) + \beta \nabla f(x_{k-1}) + \beta^2 \nabla f(x_{k-2}) + \cdots + \beta^k \nabla f(x_0) \right).$$

- b) i) For the considered function the gradient algorithm with constant  $\alpha$  is described by the iteration

$$x_{k+1,1} = x_{k,1} - \alpha(4x_{k,1}) = (1 - 4\alpha)x_{k,1},$$

$$x_{k+1,2} = x_{k,2} - \alpha(x_{k,2}) = (1 - \alpha)x_{k,2}.$$

Both sequences  $\{x_{k,1}\}$  and  $\{x_{k,2}\}$  converge to 0 if, and only if, the conditions

$$-1 < 1 - 4\alpha < 1, \quad -1 < 1 - \alpha < 1$$

hold simultaneously, which is equivalent to  $\alpha \in (0, 1/2)$ .

Setting  $\alpha = 1/4$  yields

$$x_{k+1,1} = 0, \quad x_{k+1,2} = \frac{3}{4}x_{k,2},$$

hence  $x_{k,1} = 0$ , for all  $k \geq 1$ .

To determine the speed of convergence note that we can consider only the sequence  $\{x_{k,2}\}$ , which is such that (recall that the sequence converges to 0)

$$\frac{x_{k+1,2}}{x_{k,2}} = \frac{3}{4},$$

which shows linear speed of convergence.

- ii) For the considered function and under the stated conditions the heavy ball algorithm is described by the iteration

$$\begin{aligned}x_{k+1,1} &= x_{k,1} - \alpha(4x_{k,1}) + \beta(x_{k,1} - x_{k-1,1}), \\x_{k+1,2} &= x_{k,2} - \alpha(x_{k,2}) + \beta(x_{k,2} - x_{k-1,2}).\end{aligned}$$

The first of the equations above, the condition  $x_{1,0} = x_{1,-1}$ , and  $\alpha = 1/4$  imply  $x_{1,1} = 0$  and  $x_{k,1} = 0$ , for all  $k \geq 1$ .

The second of the equations above and the results in part a) yield

$$x_{k+1,2} = x_{k,2} - \frac{1}{4} \left( x_{k,2} + \frac{3}{4}x_{k-1,2} + \dots \right).$$

Hence

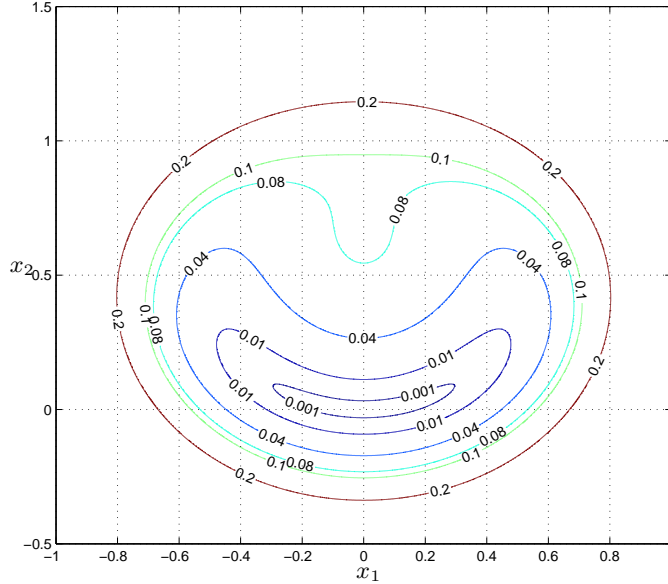
$$\begin{aligned}x_{1,2} &= \frac{3}{4}x_{0,2}, \\x_{2,2} &= x_{1,2} - \frac{1}{4}(x_{1,2} + 3/4x_{0,2}) = \frac{1}{2}x_{2,1}, \\x_{3,2} &= x_{2,2} - \frac{1}{4}(x_{2,2} + \frac{3}{4}x_{1,2} + \frac{9}{16}x_{0,2}) = 0, \\x_{4,2} &= 0,\end{aligned}$$

which shows that the sequence generated by the heavy ball algorithm converges in *finite time*.

**Exercise 12** Consider the problem of minimizing the function

$$f(x_1, x_2) = x_2^2 - \delta x_2(x_1^2 + x_2^2) + (x_1^2 + x_2^2)^2,$$

the level lines of which, for  $\delta = \sqrt{32}/3$  are plotted in the figure below.



- a) Compute all stationary points of the function as a function of  $\delta$ .

- b) Assume  $\delta = \sqrt{32}/3$ .
- Determine the stationary points of the function  $f$ , indicate them the figure, and *classify* the stationary points *i.e.* say which is a local minimizer, or a local maximizer, or a saddle point, without computing the Hessian matrix of  $f$ .
  - Determine, from inspection of the figure, a set of points such that the gradient algorithm with exact line search initialized at such points yields a sequence which converges to the global minimizer in one step. Sketch the obtained set on the figure.
  - Determine, analytically, all points such that the gradient algorithm with exact line search initialized at such points yields a sequence which converges to the global minimizer in one step. Sketch the obtained set on the figure.

**Solution 12**

- a) The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} 2x_1(2x_1^2 - \delta x_2 + 2x_2^2) \\ 2x_2 - \delta x_1^2 - 3\delta x_2^2 + 4x_2x_1^2 + 4x_2^3 \end{bmatrix}.$$

From the first equation we have  $x_1 = 0$  or  $x_1^2 = -x_2^2 + \frac{\delta}{2}x_2$ . Replacing  $x_1 = 0$  in the second equation yields

$$0 = x_2(2 - 3\delta x_2 + 4x_2^2).$$

Replacing  $x_1^2 = -x_2^2 + \frac{\delta}{2}x_2$  in the second equation yields

$$0 = -\frac{1}{2}x_2(\delta - 2)(\delta + 2).$$

In conclusion the function  $f$  has the following stationary points.

- $P_0 = (0, 0)$ , for any value of  $\delta$ .
- $P_1 = (0, \frac{3\delta + \sqrt{9\delta^2 - 32}}{8})$  and  $P_2 = (0, \frac{3\delta - \sqrt{9\delta^2 - 32}}{8})$  if  $\delta^2 \geq \frac{32}{9}$ . Note that if  $\delta = \pm \frac{\sqrt{32}}{3}$  then  $P_1 = P_2$ .
- If  $\delta = \pm 2$  then all points in the set  $x_1^2 + x_2^2 - \frac{\delta}{2}x_2 = x_1^2 + x_2^2 \mp x_2 = 0$  are stationary points.

- b) Consider now the case in which  $\delta = \frac{\sqrt{32}}{3}$ .

- The only stationary points are  $P_0$  and  $P_1 = P_2 = (0, \frac{\sqrt{2}}{2})$ . From the figure we conclude that  $P_0$  is a local minimizer, and  $P_1 = P_2$  is a saddle point. (The Hessian matrix is singular at  $P_0$  and  $P_1$ , hence it cannot be used to classify these points.)
- Note that the gradient of  $f$  on the  $x_2$ -axis is given by

$$\nabla f(0, x_2) = \begin{bmatrix} 0 \\ x_2(2 - \sqrt{32}x_2 + 4x_2^2) \end{bmatrix}.$$

The gradient of  $f$  on the  $x_2$ -axis is a direction of ascent which is parallel to the  $x_2$ -axis. Therefore, the gradient algorithm with exact line search yields the global minimizer in one step for all initial points on the  $x_2$ -axis.

- The set of points such that the gradient algorithm with exact line search yields a sequence which converges to the global minimizer in one step is obtained eliminating  $\alpha$ , *i.e.* the line search parameter, from the equation

$$0 = x - \alpha \nabla f(x).$$

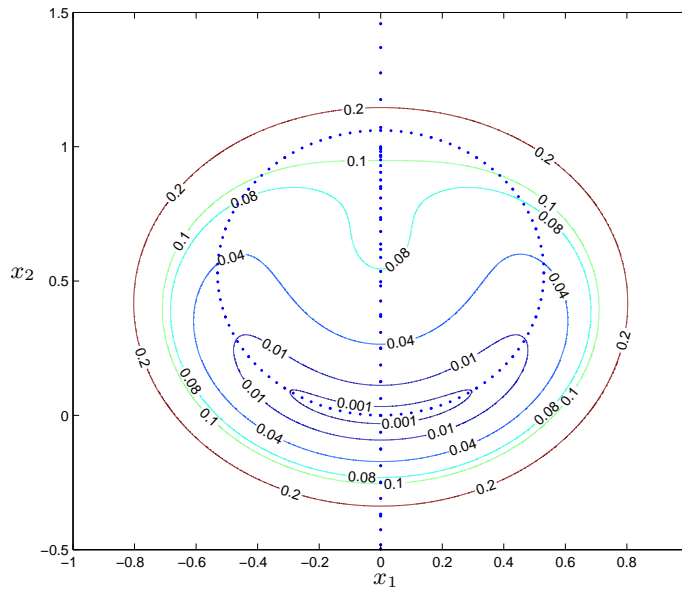
This yields the set of points described by

$$x_1(2\sqrt{2}(x_1^2 + x_2^2) - 3x_2) = 0,$$

i.e. the  $x_2$ -axis and the circle

$$x_1^2 + x_2^2 - \frac{3}{4}\sqrt{2}x_2 = 0,$$

which is a circle centered at  $P = (0, \frac{3}{8}\sqrt{2})$  and with radius equal to  $\frac{3}{8}\sqrt{2}$ . The set of all points with the requested property is indicated on the figure with “dots”.



**Exercise 13** Consider the problem of minimizing the function

$$f(x_1, x_2) = 4x_1^2 - 2x_1^4 + \frac{1}{3}x_1^6 + x_1x_2 + \frac{1}{4}x_2^2.$$

- Compute all stationary points of the function.
- Using second order sufficient conditions *classify* the stationary points determined in part a), i.e. say which is a local minimizer, or a local maximizer, or a saddle point.
- Show that the function  $f$  is radially unbounded and hence compute the global minimum of  $f$ . Is the global minimizer unique?
- Using the results of parts a), b) and c) sketch the level lines of the function  $f$ .

**Solution 13**

- The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} 8x_1 - 8x_1^3 + 2x_1^5 + x_2 \\ x_1 + \frac{1}{2}x_2 \end{bmatrix}.$$

The second equation yields  $x_2 = -2x_1$ , which replaced in the first equation yields

$$0 = 2x_1(x_1 - 1)(x_1 + 1)(x_1^2 - 3).$$

As a result, the function  $f$  has five stationary points

$$P_1 = (0, 0), \quad P_2 = (-1, 2), \quad P_3 = (1, -2), \quad P_4 = (\sqrt{3}, -2\sqrt{3}), \quad P_5 = (-\sqrt{3}, 2\sqrt{3}).$$

b) Note that

$$\nabla^2 f = \begin{bmatrix} 8 - 24x_1^2 + 10x_1^4 & 1 \\ 1 & \frac{1}{2} \end{bmatrix}.$$

As a result

$$\nabla^2 f(P_1) = \begin{bmatrix} 8 & 1 \\ 1 & \frac{1}{2} \end{bmatrix},$$

which is a positive definite matrix, hence  $P_1$  is a local minimizer;

$$\nabla^2 f(P_2) = \nabla^2 f(P_3) = \begin{bmatrix} -6 & 1 \\ 1 & \frac{1}{2} \end{bmatrix},$$

which is an indefinite matrix, hence  $P_2$  and  $P_3$  are saddle points;

$$\nabla^2 f(P_4) = \nabla^2 f(P_5) = \begin{bmatrix} 26 & 1 \\ 1 & \frac{1}{2} \end{bmatrix},$$

which is a positive definite matrix, hence  $P_4$  and  $P_5$  are local minimizers.

c) The function  $f$  can be written as

$$f = \left(x_1 + \frac{1}{2}x_2\right)^2 + \frac{x_1^2}{3}(x_1^2 - 3)^2.$$

Hence  $f$  is a *sum of squares*, and all variables  $x_1$  and  $x_2$  are present in one of the squares. As a result the function is radially unbounded and the local minimum of  $f$  is also a global minimum. Note that

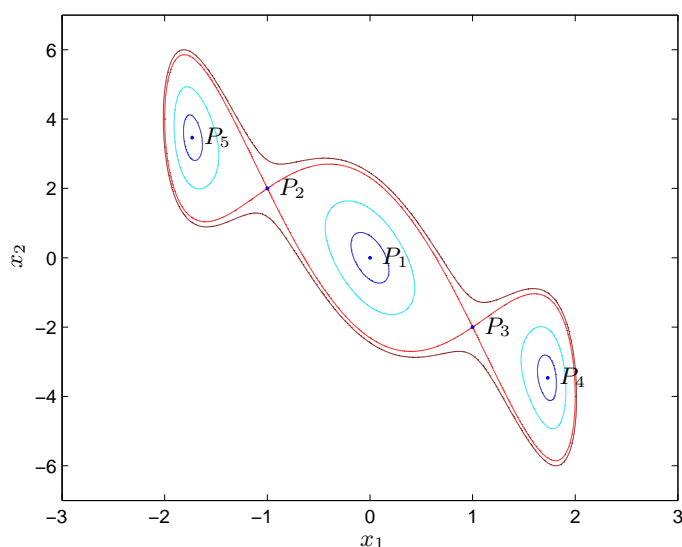
$$f(P_1) = f(P_4) = f(P_5) = 0$$

hence  $P_1$ ,  $P_4$  and  $P_5$  are all global minimizers.

d) The level lines of  $f$  can be sketched using the following considerations.

- Around the minimizers the level lines are closed.
- The value of  $f$  at the saddle points  $P_2$  and  $P_3$  is  $4/3$ . There is a level line that *connects* the saddle points. Close to the saddle points this level line is composed of two curves.

A sketch of the level lines is in the figure below.



**Exercise 14** Consider the problem of minimizing the function

$$f(x_1, x_2) = \frac{1}{2}x_1^2 \left( \frac{1}{6}x_1^2 + 1 \right) + x_2 \arctan x_2 - \frac{1}{2} \ln(x_2^2 + 1).$$

- a) Compute the unique stationary point of the function.
- b) Using second order sufficient conditions show that the stationary point determined in part a) is a local minimizer.
- c) Consider now the minimization of the function using Newton's method.
  - i) Write Newton's iteration for the considered problem.
  - ii) Show that Newton's direction is a descent direction for  $f$  at any point which is not a stationary point.
  - iii) Compute four steps of Newton's algorithm from the initial point  $(1, 0.5)$ . Compute four steps of Newton's algorithm from the initial point  $(1, 2)$ .
  - iv) Discuss why the second sequence computed in part c.iii) does not converge to the global minimizer, despite the fact that Newton's direction is always a descent direction. Propose a simple modification of Newton's iteration that would guarantee global convergence to the minimizer.

**Solution 14**

- a) The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} \frac{1}{3}x_1^3 + x_1 \\ \arctan x_2 \end{bmatrix}.$$

These equations have the unique solution  $x_1 = x_2 = 0$ , which is therefore the unique stationary point of  $f$ .

- b) Note that

$$\nabla^2 f = \begin{bmatrix} x_1^2 + 1 & 0 \\ 0 & \frac{1}{1+x_2^2} \end{bmatrix}.$$

Hence  $\nabla^2 f(0, 0) = \text{diag}(1, 1)$ , which is a positive definite matrix. The stationary point is a local minimizer.

- c) i) Newton's iteration is

$$x_{k+1} = x_k - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$$

hence

$$x_{k+1,1} = \frac{2x_{k,1}^3}{3(x_{k,1}^2 + 1)}, \quad x_{k+1,2} = x_{k,2} - (1 + x_{k,2}^2) \arctan x_{k,2}.$$

- ii) Newton's direction is

$$d = -[\nabla^2 f(x)]^{-1} \nabla f(x).$$

Note that

$$\nabla f' d = -\nabla f' [\nabla^2 f(x)]^{-1} \nabla f(x) < 0,$$

for all points such that  $\nabla f(x) \neq 0$ , since  $\nabla^2 f$  is positive definite. As a result,  $d$  is a descent direction for  $f$  for all  $x \neq 0$ .

- iii) A direct computation yields

$$\begin{aligned} x_0 &= (1, 1/2), & x_1 &= (1/3, -0.079), & x_2 &= (0.022, 0.00033), \\ x_3 &= (0.000007, -2.5 \cdot 10^{-11}), & x_4 &= (2.6 \cdot 10^{-16}, 0), \end{aligned}$$

and

$$\begin{aligned} x_0 &= (1, 2), & x_1 &= (1/3, -3.53), & x_2 &= (0.022, 13.95), \\ x_3 &= (0.000007, -279.34), & x_4 &= (2.6 \cdot 10^{-16}, 1.2 \cdot 10^5). \end{aligned}$$



- iv) The second sequence does not converge since Newton's method guarantee only local convergence properties. To achieve global convergence, since Newton's direction is a descent direction for  $f$  at any  $x \neq 0$ , it is enough to introduce a line search parameter, *i.e.* to consider the iteration

$$x_{k+1} = x_k - \alpha [\nabla^2 f(x_k)]^{-1} \nabla f(x_k),$$

with  $\alpha > 0$ , and determined using a line search algorithm.

**Exercise 15** Consider the problem of minimizing a function of  $n$  variables  $x_1, x_2, \dots, x_n$ , defined as

$$f(x_1, \dots, x_n) = f_1(x_1) f_2(x_2) \cdots f_n(x_n),$$

that is the function  $f$  is the product of the  $n$  functions  $f_i$ , each of the variable  $x_i$  only.

- a) Assume that all functions  $f_i$  are such that

$$f_i(x_i) > 0$$

for all  $x_i$  and that there exist unique  $x_i^*$  such that  $x_i^*$  is a stationary point of  $f_i$ .

- i) Compute the stationary point  $x^*$  of the function  $f$ .
  - ii) Using second order sufficient conditions show that the stationary point  $x^*$  of the function  $f$  is a strict local minimizer if and only if all  $x_i^*$  are strict local minimizers of the functions  $f_i$ .
- b) Assume  $n = 3$ , that is consider the function

$$f(x_1, x_2, x_3) = f_1(x_1) f_2(x_2) f_3(x_3).$$

Assume that the functions  $f_i$  do not have stationary points but that there exists, for  $i = 1, 2, 3$ , a unique point  $x_i^\circ$  such that

$$f_i(x_i^\circ) = 0$$

and

$$f_i(x_i) \neq 0$$

for all  $x_i \neq x_i^\circ$ .

- i) Compute all stationary points of the function  $f$ .
- ii) Show that the Hessian matrix of  $f$  at any stationary point is either identically zero or it has positive and negative eigenvalues. Hence argue that none of the stationary point can be a strict local minimizer.  
(Hint: recall that a symmetric matrix has real eigenvalues and that the trace of a matrix, that is the sum of its diagonal entries, is equal to the sum of its eigenvalues.)

**Solution 15**

- a) Consider the function  $f$ .

- i) The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} f_2 \cdots f_n \\ f_1 \frac{\partial f_2}{\partial x_2} f_3 \cdots f_n \\ \vdots \\ f_1 f_2 \cdots \frac{\partial f_n}{\partial x_n} \end{bmatrix}.$$

Since all  $f_i$ 's are positive, and have a unique stationary point, the only stationary point of  $f$  is the point

$$x^* = (x_1^*, x_2^*, \dots, x_n^*).$$

ii) Note that, for  $i \neq j$ ,

$$\frac{\partial^2 f}{\partial x_i \partial x_j} = \frac{\partial f_i}{\partial x_i} \frac{\partial f_j}{\partial x_j} M$$

where  $M$  is a positive function, hence

$$\frac{\partial^2 f}{\partial x_i \partial x_j}(x^*) = 0.$$

As a result, the Hessian matrix of the function  $f$  at  $x^*$  is

$$\nabla^2 f(x^*) = \text{diag} \left( \frac{\partial^2 f_1}{\partial x_1^2}(x_1^*) f_2(x_2^*) \cdots f_n(x_n^*), \cdots, f_1(x_1^*) f_2(x_2^*) \cdots \frac{\partial^2 f_n}{\partial x_n^2}(x_n^*) \right).$$

This implies that the function  $f$  has a strict local minimizer at  $x^*$  if and only if all functions  $f_i$  have a strict local minimizer at  $x_i^*$ .

b) Consider the function  $f$  with  $n = 3$ .

i) The stationary points of the functions  $f$  are the solution of the equations

$$0 = \nabla f = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} f_2 f_3 \\ f_1 \frac{\partial f_2}{\partial x_2} f_3 \\ f_1 f_2 \frac{\partial f_3}{\partial x_3} \end{bmatrix}.$$

These equations admit infinitely many solutions given by

$$x_{12}^\circ = (x_1^\circ, x_2^\circ, \bar{x}_3), \quad x_{13}^\circ = (x_1^\circ, \bar{x}_2, x_3^\circ), \quad x_{23}^\circ = (\bar{x}_1, x_2^\circ, x_3^\circ),$$

where  $\bar{x}_1$ ,  $\bar{x}_2$  and  $\bar{x}_3$  are arbitrary values.

ii) The Hessian matrix of  $f$  is

$$\nabla^2 f = \begin{bmatrix} \frac{\partial^2 f_1}{\partial x_1^2} f_2 f_3 & \frac{\partial f_1}{\partial x_1} \frac{\partial f_2}{\partial x_2} f_3 & \frac{\partial f_1}{\partial x_1} \frac{\partial f_3}{\partial x_3} f_2 \\ \frac{\partial f_1}{\partial x_1} \frac{\partial f_2}{\partial x_2} f_3 & \frac{\partial^2 f_2}{\partial x_2^2} f_1 f_3 & \frac{\partial f_2}{\partial x_2} \frac{\partial f_3}{\partial x_3} f_1 \\ \frac{\partial f_1}{\partial x_1} \frac{\partial f_3}{\partial x_3} f_2 & \frac{\partial f_2}{\partial x_2} \frac{\partial f_3}{\partial x_3} f_1 & \frac{\partial^2 f_3}{\partial x_3^2} f_1 f_2 \end{bmatrix}.$$

Hence

$$\nabla^2 f(x_{12}^\circ) = \begin{bmatrix} 0 & \alpha & 0 \\ \alpha & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix},$$

where

$$\alpha = \frac{\partial f_1}{\partial x_1}(x_1^\circ) \frac{\partial f_2}{\partial x_2}(x_2^\circ) f_3(\bar{x}_3).$$

The function  $\alpha$  is zero for  $\bar{x}_3 = x_3^\circ$  and it is non-zero otherwise. Hence,  $\nabla^2 f(x_{12}^\circ)$  is either identically zero or has trace zero, which means that it has a positive and a negative eigenvalue. In both cases, the points  $x_{12}^\circ$  cannot be local strict minimizers.

Similar considerations apply to  $x_{13}^\circ$  and  $x_{23}^\circ$ .

**Exercise 16** An alternative way to introduce Newton's method for the solution of a nonlinear equation is to consider the evaluation of the integral

$$f(x) = f(x_k) + \int_{x_k}^x \dot{f}(t) dt,$$

where  $\dot{f}$  denotes the derivative of the function  $f$ , by means of the so-called Newton-Cotes quadrature formula of order zero (the rectangular rule) yielding

$$f(x) \approx f(x_k) + (x - x_k)\dot{f}(x_k),$$

setting  $x = x_{k+1}$  and replacing the  $\approx$  sign with an  $=$  sign, thus yielding

$$f(x_{k+1}) = f(x_k) + (x_{k+1} - x_k)\dot{f}(x_k),$$

and setting  $f(x_{k+1}) = 0$ , thus obtaining the iteration

$$x_{k+1} = x_k - \frac{f(x_k)}{\dot{f}(x_k)}.$$

- a) Consider the evaluation of the integral by means of the Newton-Cotes quadrature formula of order one (the trapezoidal rule), that is

$$\int_{x_k}^x \dot{f}(t) dt \approx \frac{x - x_k}{2} (\dot{f}(x_k) + \dot{f}(x)).$$

- i) Determine a new iteration for the solution of the nonlinear equation  $f(x) = 0$ . Note that the obtained iteration, which is a modified Newton's iteration, is implicitly defined, that is  $x_{k+1}$  is a function of  $x_k$  and of  $\dot{f}(x_{k+1})$ .
- ii) An explicit iteration can be obtained replacing  $\dot{f}(x_{k+1})$  with  $\dot{f}(x^*)$ , where

$$x^* = x_k - \frac{f(x_k)}{\dot{f}(x_k)}.$$

Write the expression of the resulting modified Newton's iteration.

- b) Consider the problem of determining the square root of 2.
- i) Write Newton's iteration for the solution of this problem. Let  $x_0 = 1$  and apply three steps of Newton's iteration, that is compute the values  $x_1$ ,  $x_2$ , and  $x_3$  resulting from the application of Newton's iteration with the given initial point. Evaluate the absolute error  $e_k = |\sqrt{2} - x_k|$ .
- ii) Write the modified Newton's iteration for the solution of this problem. Let  $x_0 = 1$  and apply three steps of the modified Newton's iteration, that is compute the values  $x_1$ ,  $x_2$ , and  $x_3$  resulting from the application of the modified Newton's iteration with the given initial point. Evaluate the absolute error  $e_k = |\sqrt{2} - x_k|$ .
- iii) Compare the Newton's iteration and the modified Newton's iteration in terms of convergence speed and computational complexity.

#### Solution 16

- a) i) Consider the relation

$$f(x) = f(x_k) + \frac{x - x_k}{2} (\dot{f}(x) + \dot{f}(x_k)).$$

Setting  $x = x_{k+1}$  and  $f(x) = 0$  yields

$$0 = f(x_k) + \frac{x_{k+1} - x_k}{2} (\dot{f}(x_{k+1}) + \dot{f}(x_k)),$$

hence solving for  $x_{k+1}$  provides the iteration

$$x_{k+1} = x_k - 2 \frac{f(x_k)}{\dot{f}(x_{k+1}) + \dot{f}(x_k)}$$

- ii) The modified Newton's iteration is

$$x_{k+1} = x_k - 2 \frac{f(x_k)}{\dot{f}(x_k - f(x_k)/\dot{f}(x_k)) + \dot{f}(x_k)}.$$

b) To determine the square root of 2 consider the equation  $x^2 - 2 = 0$ .

i) Newton's iteration is given by

$$x_{k+1} = x_k - \frac{1}{2} \frac{x_k^2 - 2}{x_k}.$$

The sequence generated by Newton's iteration is

$$x_0 = 1, \quad x_1 = 1.5, \quad x_2 = 1.416666667, \quad x_3 = 1.414215686,$$

and this yields the sequence of the absolute error

$$e_0 = 0.414213562, \quad e_1 = 0.085786438, \quad e_2 = 0.002453105, \quad e_3 = 0.000002124.$$

ii) The modified Newton's iteration is

$$x_{k+1} = x_k - \frac{2(x_k^2 - 2)x_k}{3x_k^2 + 2}.$$

The sequence generated by the modified Newton's iteration is

$$x_0 = 1, \quad x_1 = 1.4, \quad x_2 = 1.414213198, \quad x_3 = 1.414213563,$$

and this yields the sequence of the absolute error

$$e_0 = 0.414213562, \quad e_1 = 0.014213562, \quad e_2 = 3.64 \times 10^{-7}, \quad e_3 = 1 \times 10^{-9}.$$

iii) The modified Newton's iteration is much faster (this is a general conclusion) and has similar complexity than the (classical) Newton's iteration.

**Exercise 17** Consider the function

$$f(x) = x_1^4 + x_1 x_2 + \frac{1}{2} x_2^2.$$

- Compute the stationary points of the function.
- Using second order sufficient conditions classify the stationary points determined in part a), that is say which is a local minimizer, or a local maximizer, or a saddle point.
- Sketch on the  $(x_1, x_2)$ -plane the level lines of the function  $f$ .
- Consider the point  $P_0 = (0, 0)$ .
  - Determine a direction  $d_0$  which is a descent direction for  $f$  at  $P_0$ .
  - Consider the problem of performing an exact line search along the direction  $d_0$  starting from  $P_0$ . Determine a solution to such a problem.

**Solution 17**

- The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} 4x_1^3 + x_2 \\ x_1 + x_2 \end{bmatrix}.$$

Replacing the second equation in the first yields  $x_1(4x_1^2 - 1) = 0$ . Hence, the stationary points are

$$P_1 = (0, 0), \quad P_2 = \left(\frac{1}{2}, -\frac{1}{2}\right), \quad P_3 = \left(-\frac{1}{2}, \frac{1}{2}\right).$$

- b) The Hessian matrix of the function  $f$  is

$$\nabla^2 f(x) = \begin{bmatrix} 12x_1^2 & 1 \\ 1 & 1 \end{bmatrix}.$$

Note that

$$\nabla^2 f(P_1) = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}$$

is indefinite and

$$\nabla^2 f(P_2) = \nabla^2 f(P_3) = \begin{bmatrix} 3 & 1 \\ 1 & 1 \end{bmatrix}$$

is positive definite. Hence  $P_2$  and  $P_3$  are local minimizers, and  $P_1$  is a saddle point.

- c) The level lines of  $f$  can be sketched using the following considerations.

- Around the minimizers the level lines are closed.
- The value of  $f$  at the saddle point  $P_1$  is 0.
- The value of  $f$  at the local minimizers  $P_2$  and  $P_3$  is  $-1/16$ . There is a closed level line which goes through the saddle point and encircles both local minimizers.

A sketch of the level lines is in the figure below.

- d) Note that  $\nabla f(P_0) = 0$ , hence for any direction  $d$  the scalar product  $\nabla' f d$  is zero, *i.e.* it is not possible to use first order sufficient conditions to establish if a direction is a descent direction.

- i) Let, for example,  $d_0 = [1, -1]'$  and consider the restriction of the function  $f$  along  $d_0$ , with initial point  $P_0$ , namely

$$f(P_0 + \alpha d_0) = \alpha^2 \left( -\frac{1}{2} + \alpha^2 \right).$$

For any  $\alpha > 0$  and sufficiently small (namely  $\alpha \in (0, 1/\sqrt{2})$ )

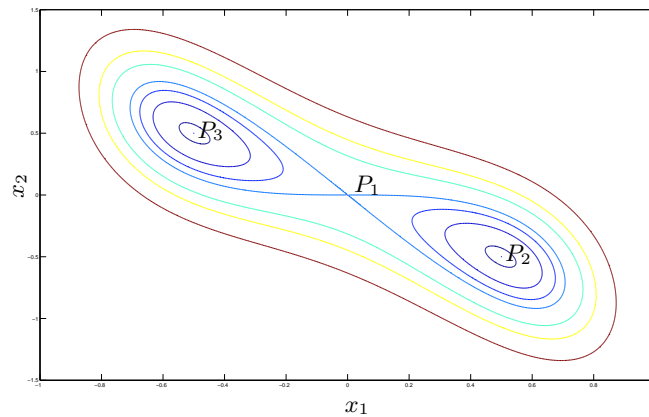
$$f(P_0) > f(P_0 + \alpha d_0),$$

hence  $d_0$  is a descent direction for  $f$  at  $P_0$ .

- ii) To solve an exact line search problem along  $d_0$  at  $P_0$  one has to find the global minimizer, if it exists, of  $f(P_0 + \alpha d_0)$ . Note that the function

$$f(P_0 + \alpha d_0) = \alpha^2(-1/2 + \alpha^2)$$

is radially unbounded (and bounded from below), hence possesses a global minimizer, which is a stationary point. The stationary points of this function are  $\alpha = 0$  (local maximizer) and  $\alpha = \pm 1/2$  (local minimizer). Hence, an exact line search along  $d_0$ , starting at  $P_0$ , gives either the point  $P_2$  or the point  $P_3$ .



**Exercise 18** Consider the function

$$f(x_1, x_2) = \sin(x_1^2 + x_2^2).$$

- Sketch on the  $(x_1, x_2)$ -plane the level lines of the function  $f$ .
- Compute the stationary points of the function.
- Explain why second order sufficient conditions of optimality are inadequate to classify some of the stationary points of the functions.
- Consider the change of variable

$$x_1 = \rho \cos \theta, \quad x_2 = \rho \sin \theta,$$

with  $\rho \geq 0$  and  $\theta \in (-\pi, \pi]$ .

- Rewrite the function  $f$  in the new variables. Note that the function depends only upon the variable  $\rho$ .
- Compute the stationary points of the function  $f$  as a function of  $\rho$  and classify these stationary points.
- Exploiting the results in part d.ii) classify the stationary points of the function  $f$ .

**Solution 18**

- Note that the function is constant on any circle centered at the origin, *i.e.* on any set of the form  $x_1^2 + x_2^2 = R^2$ . A sketch of the level lines is therefore as in the figure below.
- The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} 2x_1 \cos(x_1^2 + x_2^2) \\ 2x_2 \cos(x_1^2 + x_2^2) \end{bmatrix}.$$

Hence, the point  $(0, 0)$  is a stationary point and all points such that

$$x_1^2 + x_2^2 = \frac{\pi}{2} + k\pi,$$

with  $k$  integer, are stationary points.

- The Hessian matrix of the function  $f$  is

$$\nabla^2 f(x) = 2 \cos(x_1^2 + x_2^2) I - 4 \sin(x_1^2 + x_2^2) \begin{bmatrix} x_1^2 & x_1 x_2 \\ x_1 x_2 & x_2^2 \end{bmatrix}.$$

Note that

$$\nabla^2 f(0) = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}$$

is positive definite, hence the point  $(0, 0)$  is a local minimizer. To classify the stationary points such that  $x_1^2 + x_2^2 = \frac{\pi}{2} + k\pi$ , note that at such points  $P_k$

$$\nabla^2 f(P_k) = \mp 4 \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \begin{bmatrix} x_1 & x_2 \end{bmatrix},$$

hence  $\nabla^2 f(P_k)$  is singular, and this does not enable the use of second order sufficient conditions of optimality (which require the Hessian to be non-singular).

- The function  $f$  in the new variables is given by

$$f(\rho, \theta) = \sin \rho^2,$$

hence it is a function of  $\rho$  only.

ii) The stationary points of the function  $\sin \rho^2$  are all points such that

$$\frac{df}{d\rho} = 2\rho \cos \rho^2 = 0.$$

These are given by

$$\rho = 0 \quad \rho^2 = \frac{\pi}{2} + k\pi,$$

with  $k$  any non-negative integer.

iii) Note that

$$\frac{d^2 f}{d\rho^2} = 2 \cos \rho^2 - 4\rho^2 \sin \rho^2,$$

hence the point  $\rho = 0$  is a local minimizer, the points

$$\rho^2 = \frac{\pi}{2} + 2k\pi,$$

with  $k$  any non-negative integer, are local maximizers, and the points

$$\rho^2 = \frac{\pi}{2} + (2k+1)\pi,$$

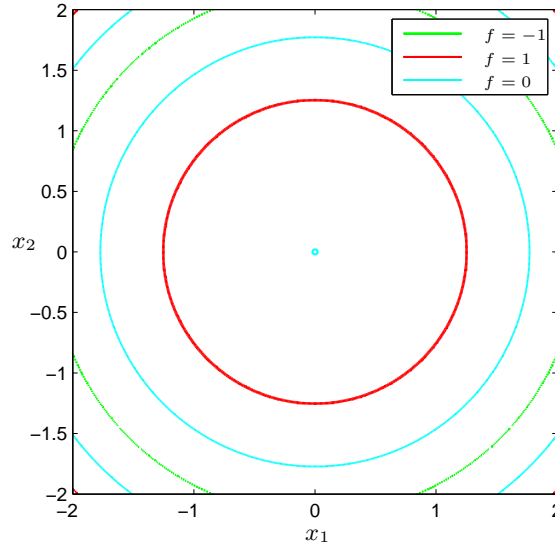
with  $k$  any non-negative integer, are local minimizers. This implies that the point  $(x_1, x_2) = (0, 0)$  is a local strict minimizer, the points  $(x_1, x_2)$  such that

$$x_1^2 + x_2^2 = \frac{\pi}{2} + 2k\pi,$$

with  $k$  any non-negative integer, are local non-strict maximizers, and the points  $(x_1, x_2)$  such that

$$x_1^2 + x_2^2 = \frac{\pi}{2} + (2k+1)\pi,$$

with  $k$  any non-negative integer, are local non-strict minimizers.



**Exercise 19** A nonlinear least-squares problem is an unconstrained optimization problem of the form

$$\min_x \frac{1}{2} \sum_{i=1}^m r_i^2(x),$$

where  $x \in \mathbb{R}^n$ . The functions  $r_1, r_2, \dots, r_m$  are called residuals and the objective function can be rewritten as  $\frac{1}{2}r'(x)r(x)$ , with

$$r(x) = \begin{bmatrix} r_1(x) \\ \vdots \\ r_m(x) \end{bmatrix}.$$

- a) Write Newton's iteration for the solution of the considered least-square problem.
- b) Gauss-Newton's iteration for the solution of the considered least-square problem is given by

$$x_{(k+1)} = x_{(k)} - [J'(x_{(k)})J(x_{(k)})]^{-1}J'(x_{(k)})r(x_{(k)}),$$

where

$$J(x) = \begin{bmatrix} \frac{\partial r_1}{\partial x_1} & \cdots & \frac{\partial r_1}{\partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial r_m}{\partial x_1} & \cdots & \frac{\partial r_m}{\partial x_n} \end{bmatrix}$$

and  $x_{(k)} = [x_{k,1}, \dots, x_{k,n}]'$ . Discuss the differences between Newton's iteration and Gauss-Newton's iteration.

(Hint: consider the *difference* between the Hessian of  $\frac{1}{2}r'(x)r(x)$  and the matrix  $J'(x)J(x)$ .) Discuss under what conditions the Gauss-Newton direction

$$d_{GN} = -[J'(x)J(x)]^{-1}J'(x)r(x)$$

is a descent direction.

- c) Assume  $m = 2$ ,  $x = (x_1, x_2)$  and

$$r_1(x) = x_1 + x_2 - x_1x_2 + 2, \quad r_2(x) = x_1 - e^{x_2}.$$

- i) Sketch on the  $(x_1, x_2)$ -plane the set of points  $r_1(x) = 0$  and  $r_2(x) = 0$ , hence argue that the considered least-square problem has two (global) solutions. Find an approximation of these global solutions using graphical considerations.
- ii) Write explicitly Gauss-Newton's iteration for the considered problem.
- iii) Compute three iterations of Gauss-Newton's methods from the initial conditions  $(0, 0)$ . Evaluate the residuals at  $(0, 0)$  and at the last iteration.
- iv) Comment on the convergence speed and complexity of Gauss-Newton's method.

### Solution 19

- a) Newton's method for the minimization of the function

$$f(x) = \frac{1}{2}r'(x)r(x)$$

is described by the iteration

$$x_{k+1} = x_k - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k),$$

where

$$\nabla f = \left[ \frac{\partial r}{\partial x} \right]' r = J'(x)r$$

(with  $J = \frac{\partial r}{\partial x}$ , as defined above) and

$$\nabla^2 f = \left[ \frac{\partial r}{\partial x} \right]' \left[ \frac{\partial r}{\partial x} \right] + \sum_{i=1}^m r_i \nabla^2 r_i.$$



- b) The difference between Newton's method and Gauss-Newton's method is in the matrix that it is inverted. In Newton's method this is the Hessian of the function to be minimized, in Gauss-Newton's method this is one term of the Hessian, which can be computed using only first derivatives. Gauss-Newton's direction is a descent direction if

$$\nabla' f d_{GN} = -r'(x)J(x)[J'(x)J(x)]^{-1}J'(x)r(x) < 0,$$

which holds at all points in which  $J(x)$  is full rank and  $r(x) \neq 0$ .

- c) i) The sets  $r_1(x) = r_2(x) = 0$  are displayed in the figure below. These sets have two points of intersection, hence the least square problem has only two solutions, which are both global minimizers of the function  $\frac{1}{2}r'r$ . From the graph one sees that the minimizers are approximately given by the points  $(0.1, -2.3)$  and  $(5.4, 1.7)$ .
- ii) Note that

$$J(x) = \begin{bmatrix} 1 - x_2 & 1 - x_1 \\ 1 & -e^{x_2} \end{bmatrix}$$

and

$$d_{GN}(x) = -\frac{1}{(x_1 - 1) + e^{x_2}(x_2 - 1)} \begin{bmatrix} -x_2 e^{x_2} + x_1 x_2 e^{x_2} - 2e x_2 x_1 - e^{x_2} + x_1^2 - x_1 \\ -e^{x_2} + x_2 e^{x_2} - 2 - x_2 \end{bmatrix}.$$

Hence, Gauss-Newton iteration can be written as

$$x_{k+1} = x_k + d_{GN}(x_k).$$

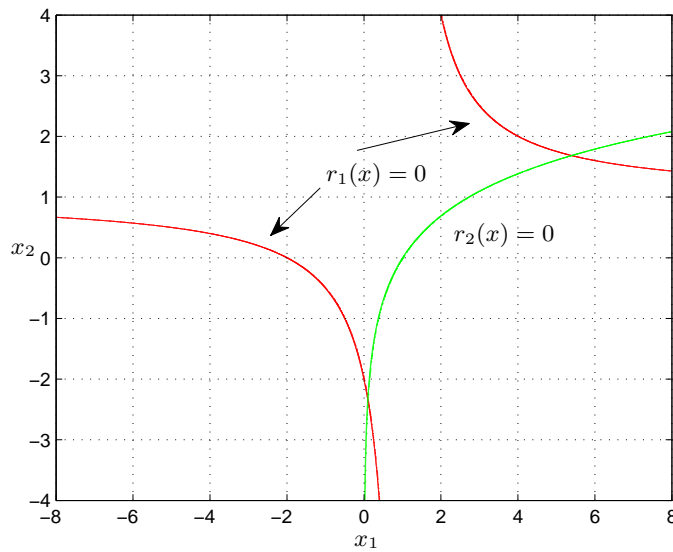
- iii) Let  $x_{(0)} = (0, 0)$ . The residuals at  $(0, 0)$  are  $r_1(0) = 2$  and  $r_2(0) = -1$ . The first three elements of the sequence generated by Gauss-Newton's iteration are

$$x_{(1)} = (-0.5, -1.2), \quad x_{(2)} = (0.0974, -2.1345), \quad x_{(3)} = (0.096347, -2.319927),$$

and the value of the residuals after three iterations are

$$r_1(x_{(3)}) = -0.000061740, \quad r_2(x_{(3)}) = -0.001933.$$

- iv) It is worth noting the fast convergence rate despite the fact that the iteration does not use second derivatives and a line search parameter.



**Exercise 20** Consider the function

$$f(x_1, x_2) = \frac{1}{2}x_1^2 \left( \frac{1}{6}x_1^2 + 1 \right) + x_2 \arctan x_2 - \frac{1}{2} \ln(x_2^2 + 1).$$

- a) Compute the unique stationary point of the function.
- b) Using second order sufficient conditions of optimality show that the stationary point determined in part a) is a local minimizer. Show, in addition, that the function is convex. Finally, show that the local minimizer is a global minimizer.  
(Hint: convexity of a function  $f$  is implied by the condition  $\nabla^2 f(x) > 0$  for all  $x$ .)

- c) Consider the problem of minimizing the function  $f$  using Newton's method.

- i) Write Newton's iteration for the minimization of the function  $f$ .
- ii) Perform 4 steps of Newton's iteration with starting point

$$(x_1, x_2) = (1, 2).$$

- d) Consider the function

$$f_2(x_2) = x_2 \arctan x_2 - \frac{1}{2} \ln(x_2^2 + 1).$$

- i) Using the iteration derived in part c.i) write Newton's iteration for the minimization of the function  $f_2$ .
- ii) Write the Newton's iteration in part d.i) in the form

$$x_2(k+1) = \psi(x_2(k)).$$

Write explicitly the function  $\psi$ .

- iii) Plot on the same graph the functions  $x_2$  and  $\psi(x_2)$ . Exploiting the graph explain why Newton's iteration for the minimization of  $f_2$  converges for initial conditions sufficiently close to zero, and diverges otherwise.  
(Hint: use the graph to *execute* Newton's iteration graphically.)
- e) Exploiting the results in part d) and the fact that the function  $f$  is the sum of two functions of one variable each, determine (qualitatively) for which initial points the Newton's iteration for the minimization of  $f$  converges to the minimizer.

**Solution 20**

- a) The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} \\ \frac{\partial f_2}{\partial x_2} \end{bmatrix} = \begin{bmatrix} \frac{1}{3}x_1(x_1^2 + 3) \\ \arctan x_2 \end{bmatrix}.$$

Hence, the point  $(0, 0)$  is the unique stationary point.

- b) The Hessian matrix of the function  $f$  is

$$\nabla^2 f(x) = \begin{bmatrix} x_1^2 + 1 & 0 \\ 0 & \frac{1}{1+x_2^2} \end{bmatrix}.$$

Note that

$$\nabla^2 f(0) = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

is positive definite, hence the point  $(0, 0)$  is a local minimizer. In addition,  $\nabla^2 f > 0$  for all  $(x_1, x_2)$ , hence the function is convex. For convex function, a stationary point is a global minimizer, hence  $(0, 0)$  is a global minimizer.

- c) i) Newton's iteration, considering that the function  $f$  is the sum of a function of  $x_1$  and of a function of  $x_2$ , gives two *decoupled* equations, namely

$$x_1^{k+1} = \frac{2}{3} \frac{x_1^3}{1+x_1^2} \quad x_2^{k+1} = x_2 - (1+x_2^2) \arctan x_2.$$

- ii) The first five elements of the sequences  $\{x_1^k\}$  and  $\{x_2^k\}$  are

$$x_1^0 = 1, \quad x_1^1 = 1/3, \quad x_1^2 = 1/45, \quad x_1^3 = 1/136755, \quad x_1^4 = 1/3836373661058445 \approx 0,$$

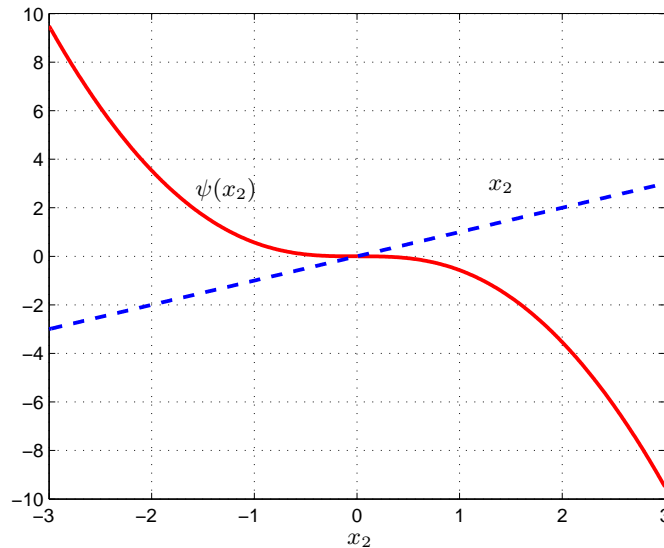
and

$$x_2^0 = 2, \quad x_2^1 = -3.5357, \quad x_2^2 = 13.95095909, \quad x_2^3 = -279.3440667, \quad x_2^4 = 122016.9990.$$

- d) i) The iteration is the same as the " $x_2$ " iteration in part c.i).  
 ii) The function  $\psi$  is given by

$$\psi(x_2) = x_2 - (1+x_2^2) \arctan(x_2).$$

- iii) The graphs of the considered functions are displayed in the figure below. One can use the graph to show how Newton's iteration works. In fact, pick a point  $x_2^k$  on the  $x_2$ -axis, and *lift* it on the graph of the function  $\psi$ . Then *move* the point horizontally on the graph of the function  $x_2$ , and then vertically on the  $x_2$ -axis. This is the point  $x_2^{k+1}$ . Iterating the procedure one can construct the sequence  $\{x_2^k\}$ . Using this approach, one concludes that if  $x_2^0$  is sufficiently close to zero the iteration yields a sequence converging to  $x_2 = 0$ . If  $|x_2^0|$  is large, then the sequence diverges.  
 e) As shown in part c.i), Newton's iteration is composed of two decoupled iterations. The iteration for  $x_1$  yields a globally converging sequence, whereas the iteration for  $x_2$  converges only for  $|x_2^0|$  sufficiently small (to be precise, for  $|x_2^0| < 1.39\dots$ ). Hence, for all initial points  $(x_1^0, x_2^0)$  such that  $|x_2^0| < 1.39\dots$ , the iteration yields a sequence converging to the global minimizer.



**Exercise 21** The company  $XYZ$  has invested £20000 to develop a new product. The product can be manufactured for £2 per unit. The company then performs a marketing research. The conclusion of the research is that if the company spends £ $a$  on advertising then it can sell the product at price £ $p$  per unit and it will sell  $2000 + 4\sqrt{a} - 20p$  units.

- a) Compute the revenue for sales as a function of  $a$  and  $p$ .
- b) Compute the overall costs associated to the production and commercialization of the product, that is the development cost plus the production cost and the advertising cost, as a function of  $a$  and  $p$ .
- c) Compute the company's profit as a function of  $a$  and  $p$ .
- d) The company wishes to select  $a$  and  $p$  to maximize the profit. Pose this problem as an unconstrained optimization problem (disregard the non-negativity conditions on  $a$  and  $p$ ).
- e) Compute the unique stationary point of the profit. Using second order sufficient conditions of optimality show that the stationary point is a local maximizer.
- f) Assume that the company is forced to fix the sale price of the product to  $p = \tilde{p}$ , with  $\tilde{p} > 2$ .
  - i) Determine the optimal advertising cost as a function of  $\tilde{p}$ .
  - ii) Determine the optimal profit as a function of  $\tilde{p}$ .
  - iii) Plot the optimal profit as a function of the fixed price  $\tilde{p}$  and show that as  $\tilde{p}$  increases the profit becomes negative.

**Solution 21**

- a) The revenue for sales is given by

$$\text{revenue} = p(2000 + 4\sqrt{a} - 20p).$$

- b) The costs are

$$\begin{aligned} \text{production cost} &= 2(2000 + 4\sqrt{a} - 20p), \\ \text{development cost} &= 20000, \\ \text{advertising cost} &= a. \end{aligned}$$

Hence

$$\text{total cost} = 24000 + 8\sqrt{a} - 40p + a.$$

- c) The profit is given by

$$\text{profit} = p(2000 + 4\sqrt{a} - 20p) - (24000 + 8\sqrt{a} - 40p + a).$$

- d) The optimization problem is

$$\max_{a,p} p(2000 + 4\sqrt{a} - 20p) - (24000 + 8\sqrt{a} - 40p + a).$$

- e) The stationary points of the profit are the solutions of the equations

$$0 = \frac{\partial \text{profit}}{\partial a} = 2\frac{p}{\sqrt{a}} - \frac{4}{\sqrt{a}} - 1, \quad 0 = \frac{\partial \text{profit}}{\partial p} = 2\frac{p}{\sqrt{a}} - \frac{4}{\sqrt{a}} - 1.$$

The only solution is

$$a^* = \frac{60025}{4} = 15006.25, \quad p^* = \frac{253}{4} = 63.25.$$

The Hessian of the profit at the stationary point is

$$H(a^*, p^*) = - \begin{bmatrix} \frac{2}{60025} & -\frac{4}{245} \\ -\frac{4}{245} & 40 \end{bmatrix},$$

which is negative definite, hence the point  $(a^*, p^*)$  is a local maximizer.

f) The profit for fixed price is

$$\text{profit fix price} = \tilde{p}(2000 + 4\sqrt{a} - 20\tilde{p}) - (24000 + 8\sqrt{a} - 40\tilde{p} + a).$$

i) The optimal advertising cost  $\tilde{a}^*$  is given by the solution of the equation

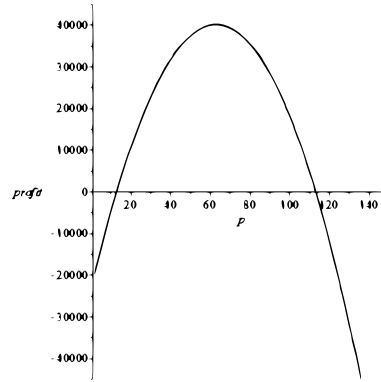
$$0 = \frac{\partial \text{profit fix price}}{\partial a},$$

which gives  $\tilde{a}^* = 4(\tilde{p} - 2)^2$ .

ii) The resulting optimal profit is

$$\text{profit fix price}^* = 2024\tilde{p} - 16\tilde{p}^2 - 23984.$$

iii) The optimal profit as a function of the fixed price  $\tilde{p}$  is displayed in the graph below. Note that as  $\tilde{p}$  increases the optimal profit becomes negative (because of the term  $-16\tilde{p}^2$ ).



**Exercise 22** Consider the function

$$f(x) = (x_1 - 2)^4 + (x_1 - 2)^2 x_2^2 + (x_2 + 1)^2.$$

- Compute the unique stationary point  $x_*$  of the function  $f$ .
- Using second order sufficient conditions of optimality show that the stationary point determined in part a) is a local minimizer. Hence, show that  $f$  is radially unbounded and that the stationary point determined in part a) is the global minimizer of  $f$ .
- Write the modified Newton's iteration for the minimization of the function  $f$  given by

$$x_{k+1} = x_k - [\nabla^2 f(x_*)]^{-1} \nabla f(x_k).$$

- Run five steps of the modified Newton's iteration in part c) from the starting point  $(1.5, 0)$ .
- Run four steps of the modified Newton's iteration in part c) from the starting point  $(1, 0)$ .
- Show that the research directions generated by the modified Newton's iteration in part c) are descent directions satisfying the condition of angle. Explain why the iteration is not globally convergent.

**Solution 22**

- The stationary points of the function  $f$  are computed by solving the equations

$$0 = \nabla f = \begin{bmatrix} 2(x_1 - 2)(2(x_1 - 2)^2 + x_2^2) \\ 2x_1^2 x_2 - 8x_1 x_2 + 10x_2 + 2 \end{bmatrix}.$$

As a result, the point  $x_* = (2, -1)$  is the unique stationary point.

- b) The Hessian matrix of the function  $f$  is

$$\nabla^2 f(x) = \begin{bmatrix} 12(x_1 - 2)^2 + 2x_2^2 & 4(x_1 - 2)x_2 \\ 4(x_1 - 2)x_2 & 2(x_1 - 2)^2 + 2 \end{bmatrix},$$

hence

$$\nabla^2 f(x_*) = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}.$$

Since  $\nabla^2 f(x_*) > 0$ ,  $x_*$  is a minimizer of  $f$ . Note now that

$$0 \leq (x_1 - 2)^4 + (x_2 + 1)^2 \leq f,$$

and the function  $(x_1 - 2)^4 + (x_2 + 1)^2$  takes non-negative values and it is radially unbounded (it is the sum of two squares, one involving  $x_1$  and one involving  $x_2$ ). Hence,  $f$  is radially unbounded, and since  $f(x_*) = 0$ ,  $x_*$  is the global minimizer of  $f$ .

- c) The modified Newton's iteration is given by

$$x_{k+1} = x_k - \frac{1}{2} \nabla f(x_k) = \begin{bmatrix} x_{k,1} - (x_{k,1} - 2)(2(x_{k,1} - 2)^2 + x_{k,2}^2) \\ -4x_{k,2} - x_{k,2}x_{k,1}^2 + 4x_{k,2}x_{k,1} - 1 \end{bmatrix}.$$

- d) The points generated by the modified Newton's iteration from the starting point  $x_0 = (3/2, 0)$  are

$$x_1 = (1.75, -1), \quad x_2 = (2.03125, -0.9375), \quad x_3 = (2.003723145, -0.9990844731),$$

$$x_4 = (2.000006711, -0.9999861507), \quad x_5 = (2.000000000, -1.000000000).$$

- e) The points generated by the modified Newton's iteration from the starting point  $x_0 = (1, 0)$  are

$$x_1 = (3, -1), \quad x_2 = (0, 0), \quad x_3 = (16, -1), \quad x_4 = (-5486, 195).$$

- f) The research direction used in the modified Newton's iteration is  $-1/2 \nabla f(x_k)$ , which is nothing else than the direction of the anti-gradient, hence it is a descent direction satisfying the condition of angle. The reason why the method is not globally convergent is that the line search parameter is *fixed* to  $\alpha = 1/2$ , and this may not yield a descent algorithm at each step.

**Exercise 23** Consider the function

$$f(x) = \frac{1}{2}x_1^2 + \frac{m}{2}x_2^2,$$

with  $m > 0$ . The function has a global minimizer at  $x_* = 0$ .

- a) Show that the gradient algorithm with exact line search for the function  $f$  can be written as

$$x_{k+1} = x_k - \frac{x_{k,1}^2 + m^2 x_{k,2}^2}{x_{k,1}^2 + m^3 x_{k,2}^2} \begin{bmatrix} x_{k,1} \\ m x_{k,2} \end{bmatrix}$$

- b) Let  $m = 9$  and  $x_0 = [9, 1]'$ . Show that the sequence of points generated by the gradient algorithm is given by

$$x_k = \begin{bmatrix} 9 \\ (-1)^k \end{bmatrix} (0.8)^k.$$

(Hint: assume that for the given values of  $m$  and  $x_0$  the quantity

$$\frac{x_1^2 + m^2 x_2^2}{x_1^2 + m^3 x_2^2}$$

remains constant for all iterations of the algorithm.)

- c) Compute the speed of convergence of the sequence generated by the algorithm and in particular show that

$$\frac{\|x_{k+1} - x_\star\|}{\|x_k - x_\star\|} = \text{constant}$$

for every  $k$ , where  $\|v\| = \sqrt{v'v}$ .

**Solution 23**

- a) Note that

$$\nabla f = \begin{bmatrix} x_1 \\ mx_2 \end{bmatrix},$$

hence the gradient algorithm is described by the iteration

$$x_{k+1,1} = x_{k,1} - \alpha x_{k,1}, \quad x_{k+1,2} = x_{k,2} - \alpha m x_{k,2}.$$

Replacing  $x_{k+1}$  in  $f$  yields

$$f(x_{k+1}) = \frac{1}{2} (x_{k,1}^2 + m x_{k,2}^2) - \alpha (x_{k,2}^2 + m^2 x_{k,2}^2) + \frac{1}{2} (x_{k,2}^2 + m^3 x_{k,2}^2) \alpha^2.$$

To obtain the exact linear search parameter one has to compute the stationary point of  $f(x_{k+1})$  as a function of  $\alpha$  (since  $f(x_{k+1})$  is convex in  $\alpha$ ), that is

$$\alpha_\star = \frac{x_{k,1}^2 + m^2 x_{k,2}^2}{x_{k,2}^2 + m^3 x_{k,2}^2}.$$

As a result, the gradient algorithm with exact line search is given by

$$x_{k+1} = x_k - \alpha_\star \nabla f(x_k),$$

as given in the question.

- b) As indicated in the question, for the considered initial condition and value of  $m$  the value of  $\alpha_\star$  is constant, namely

$$\alpha_\star = \frac{x_{0,1}^2 + m^2 x_{0,2}^2}{x_{0,1}^2 + m^3 x_{0,2}^2} = 1/5.$$

As a result, the gradient iteration is given by

$$x_{k+1,1} = \frac{4}{5} x_{k,1}, \quad x_{k+1,2} = -\frac{4}{5} x_{k,2}.$$

This yields

$$x_{k,1} = x_{0,1} \left(\frac{4}{5}\right)^k = 9 \left(\frac{4}{5}\right)^k, \quad x_{k,2} = x_{0,2} \left(-\frac{4}{5}\right)^k = (-1)^k \left(\frac{4}{5}\right)^k.$$

- c) Note that  $x_\star = 0$ , hence

$$\|x_{k+1}\|^2 = \left(9 \left(\frac{4}{5}\right)^{k+1}\right)^2 + \left((-1)^{k+1} \left(\frac{4}{5}\right)^{k+1}\right)^2 = 82 \left(\frac{4}{5}\right)^{2(k+1)},$$

$$\|x_k\|^2 = 82 \left(\frac{4}{5}\right)^{2k},$$

thus

$$\frac{\|x_{k+1} - x_\star\|}{\|x_k - x_\star\|} = \frac{4}{5}.$$

The sequence thus converges with linear speed of convergence.

**Exercise 24** Consider the problem of computing the average of four numbers,  $a_1$ ,  $a_2$ ,  $a_3$  and  $a_4$ . This problem can be posed as an unconstrained optimization problem as follows

$$\min_x f(x),$$

with

$$f(x) = (x - a_1)^2 + (x - a_2)^2 + (x - a_3)^2 + (x - a_4)^2.$$

- a) Compute the unique stationary point  $x_*$  of the function  $f$  and show that  $x_*$  is indeed the average of  $a_1$ ,  $a_2$ ,  $a_3$  and  $a_4$ .
- b) Using second order sufficient conditions of optimality show that the stationary point determined in part a) is a local minimizer. Hence, show that  $f$  is radially unbounded and that the stationary point determined in part a) is the global minimizer of  $f$ .
- c) Assume  $a_1 = 1$ ,  $a_2 = 2$ ,  $a_3 = 3$  and  $a_4 = -6$ .
  - i) Write the gradient method for the minimization of the function  $f$  and determine the exact line search parameter  $\alpha^*$ .
  - ii) Consider the gradient method with line search parameter  $\alpha = \gamma \alpha^*$ , with  $\gamma \in [0, 3]$ . Determine for which values of  $\gamma$  the iteration yields a converging sequence and, for these values of  $\gamma$ , determine the speed of convergence of the sequence.

**Solution 24** Note that

$$\begin{aligned} f(x) &= (x - a_1)^2 + (x - a_2)^2 + (x - a_3)^2 + (x - a_4)^2 \\ &= x^2 - 2a_1 + a_1^2 + x^2 - 2a_2 + a_2^2 + x^2 - 2a_3 + a_3^2 + x^2 - 2a_4 + a_4^2 \\ &= 4x^2 - 2(a_1 + a_2 + a_3 + a_4)x + a_1^2 + a_2^2 + a_3^2 + a_4^2. \end{aligned}$$

- a) The first order necessary condition of optimality is

$$0 = \nabla f = 8x - 2(a_1 + a_2 + a_3 + a_4),$$

which yields

$$x^* = \frac{a_1 + a_2 + a_3 + a_4}{4}.$$

Clearly,  $x^*$  is the average of  $a_1$ ,  $a_2$ ,  $a_3$ ,  $a_4$ .

- b) The second order sufficient condition of optimality is

$$\nabla^2 f = 8 > 0,$$

hence  $x^*$  is a local minimizer. Note that  $f$  is strictly convex and this implies that  $x^*$  is global minimizer.

- c) Note that  $f(x) = 4x^2 + c^2$ , with  $c^2 = a_1^2 + a_2^2 + a_3^2 + a_4^2$ .

- i) The gradient is

$$\nabla f = 8x$$

and the gradient algorithm gives

$$x_{k+1} = x_k - 8\alpha_k x_k = (1 - 8\alpha_k)x_k.$$

Note that

$$f(x_{k+1}) = 4(1 - 16\alpha_k + 64\alpha_k^2)x_k^2 + c^2$$

and

$$f(x_k) = 4x_k^2 + c^2$$

yields

$$f(x_{k+1}) - f(x_k) = 4(64\alpha_k^2 - 16\alpha_k)x_k^2.$$



To find the exact line search parameter solve

$$0 = \frac{\partial[f(x_{k+1}) - f(x_k)]}{\partial\alpha} = 4(128\alpha_k - 16)x_k^2,$$

obtaining

$$\alpha^* = \frac{1}{8}.$$

ii) Note now that

$$x_{k+1} = x_k - 8\gamma\alpha^*x_k = (1 - \gamma)x_k.$$

To have convergence we need

$$|1 - \gamma| < 1.$$

Hence,  $\gamma \in (0, 2)$ . For  $\gamma = 0$  or  $\gamma = 2$ ,  $|x_{k+1}| = |x_k|$ , and the sequence does not converge. For  $\gamma \in (2, 3]$   $|x_{k+1}| > |x_k|$ , hence the sequence diverges. For  $\gamma \in (0, 2)$  the speed of convergence is linear, since

$$\frac{|x_{k+1} - x^*|}{|x_k - x^*|} = \frac{|x_{k+1}|}{|x_k|} = |1 - \gamma|.$$

**Exercise 25** Consider the problem of minimizing the function

$$\min_x f(x),$$

with

$$f(x) = \frac{1}{2}x_1^2 - x_1x_2 + \frac{1}{4}x_2^4 - \frac{1}{3}x_2^3.$$

- Compute the stationary points of the function  $f$ .
- Using second order sufficient conditions of optimality classify the stationary points determined in part a). Hence, determine the global minimizer of  $f$ .
- Consider the problem of minimizing the function using the so-called gradient method with extrapolation, that is the method defined by the iteration

$$x_{k+1} = x_k - \alpha_k \nabla f(x_k) + \beta_k(x_k - x_{k-1}),$$

with  $\alpha_k > 0$  and  $\beta_k \in [0, 1)$ , for all  $k \geq 0$ , and  $x_{-1} = x_0$ . Let  $x_0 = (1, 1)$ .

- Argue that the first step of the gradient method with extrapolation coincides with the first step of the gradient method.
- Run one iteration of the gradient method with extrapolation and determine the point  $x_1$ . Note that  $x_1$  is a function of  $\alpha_0$  hence write a condition on  $\alpha_0$  such that the algorithm is a descent algorithm, that is  $f(x_1) < f(x_0)$ . Explain why  $\beta_0$  does not appear in the descent condition  $f(x_1) < f(x_0)$ .
- Pick  $\alpha_k = 1/2$  for all  $k$ . Run one more iteration of the gradient method with extrapolation (using as initial condition the point  $x_1$  determined in part c.ii), that is compute the point  $x_2$ . Determine a condition on  $\beta_1$  yielding a descent algorithm. Explain why  $\beta_1 = 0$  is a feasible selection of  $\beta_1$  and argue that it is not the best selection.

**Solution 25**

- The first order necessary condition of optimality is

$$0 = \nabla f = \begin{bmatrix} x_1 - x_2 \\ -x_1 + x_2^3 - x_2^2 \end{bmatrix},$$

which gives the equations

$$\begin{aligned} x_1 &= x_2, \\ x_2(x_2^2 - x_2 - 1) &= 0. \end{aligned}$$

The stationary points are therefore  $P_1 = (0, 0)$ ,  $P_2 = \left(\frac{1+\sqrt{5}}{2}, \frac{1+\sqrt{5}}{2}\right)$  and  $P_3 = \left(\frac{1-\sqrt{5}}{2}, \frac{1-\sqrt{5}}{2}\right)$ .

- b) The Hessian matrix of the function  $f$  is

$$\nabla^2 f(x) = \begin{bmatrix} 1 & -1 \\ -1 & 3x_2^2 - 2x_2 \end{bmatrix}.$$

Evaluating the matrix at the stationary points yields

$$\nabla^2 f(P_1) = \begin{bmatrix} 1 & -1 \\ -1 & 0 \end{bmatrix},$$

which is indefinite,

$$\nabla^2 f(P_2) = \begin{bmatrix} 1 & -1 \\ -1 & \frac{7+\sqrt{5}}{2} \end{bmatrix},$$

which is positive definite, and

$$\nabla^2 f(P_3) = \begin{bmatrix} 1 & -1 \\ -1 & \frac{7-\sqrt{5}}{2} \end{bmatrix},$$

which is also positive definite. Hence,  $P_1$  is a saddle point, whereas  $P_2$  and  $P_3$  are two local minimizers. Computing the function at this points yields  $f(P_2) = -1.0075$  and  $f(P_3) = -0.0758$ . Since  $f$  is radially unbounded, *i.e.*

$$\lim_{|x| \rightarrow +\infty} f(x) = +\infty,$$

the point  $P_2$  is the global minimizer of  $f$ .

- c) i) In the first step of the gradient method with extrapolation the term multiplied by  $\beta_0$  is zero (because of the way the algorithm is initialized), hence the iteration coincides with the gradient iteration.  
ii) Running one iteration of the algorithm from the indicated initial conditions yields the point

$$x_1 = \begin{bmatrix} 1 \\ 1 + \alpha_0 \end{bmatrix}.$$

Note that, consistently with the answer to c.i), the parameter  $\beta_0$  does not contribute to the point  $x_1$ . To check the descent condition note that

$$f(x_1) - f(x_0) = \frac{1}{12}\alpha_0(3\alpha_0^3 + 8\alpha_0^2 + 6\alpha_0 - 12),$$

hence  $\alpha_0$  should be selected such that

$$\alpha_0(3\alpha_0^3 + 8\alpha_0^2 + 6\alpha_0 - 12) < 0$$

which is the case for  $\alpha_0$  strictly positive and sufficiently small (approximately smaller than 0.82).

- iii) Using  $\alpha_k = 1/2$  and running one more iteration of the algorithm yields

$$x_2 = \begin{bmatrix} 5/4 \\ 23/16 + 1/2\beta_1 \end{bmatrix}.$$

The descent condition is now

$$f(x_2) - f(x_1) = -0.0788 - 0.1729\beta_1 + 0.415\beta_1^2 + 0.138\beta_1^3 + 0.015625\beta_1^4 < 0,$$

which shows that  $\beta_1$  should be non-negative and smaller than approximately 0.6. The selection  $\beta_1 = 0$  gives a descent condition because for such value of  $\beta_1$  one has essentially the gradient iteration, for which the descent condition holds for the given selection of  $\alpha_1$ . However,  $\beta_1 = 0$  is not *optimal* since one could have a greater decrease selecting a strictly positive value of  $\beta_1$ . The optimal selection for this particular case is approximately  $\beta_1 = 0.2$ .

**Exercise 26** The proximal method is a descent method in which the problem

$$\min_x f(x)$$

is replaced by the sequence of modified problems

$$\min_x \left( f(x) + \frac{1}{2\gamma_k} \|x - x_k\|^2 \right),$$

where  $x_k$  is the current estimate of the solution of the problem and  $x_{k+1}$  is the solution of the modified minimization problem, and  $\gamma_k > 0$ .

Consider the quadratic function

$$f(x) = \frac{1}{2} x' Q x + c' x + d,$$

with  $Q = Q' > 0$ . Recall that the function has a global minimizer at  $x^* = -Q^{-1}c$ .

- Write the optimization problem used in the proximal method and state under what conditions the problem has a unique solution.
- Solve explicitly the optimization problem arising from the proximal method, that is determine  $x_{k+1}$  as a function of  $x_k$ . In particular, write the relation between  $x_{k+1}$  and  $x_k$  in the form

$$x_{k+1} = Ax_k + b,$$

in which  $A$  is a matrix and  $b$  is a vector. (Note that  $A$  and  $b$  are functions of  $k$ .) Write explicitly the matrix  $A$  and the vector  $b$  as a function of  $Q$ ,  $c$  and  $\gamma_k$ .

- Determine the fixed point  $\bar{x}$  of the equation  $x_{k+1} = Ax_k + b$ , that is the point  $\bar{x}$  such that

$$\bar{x} = A\bar{x} + b,$$

and show that the point is the global minimizer of the quadratic function.

- Show that the iteration arising from the proximal method is globally convergent for all  $\gamma_k > 0$ . This can be achieved using the following steps.
  - Show that  $A = (\gamma_k Q + I)^{-1}$  and that  $A'A < I$ .
  - Write the iteration arising from the proximal method in the form  $x_{k+1} - x^* = A(x_k - x^*) + \tilde{b}$  and show that  $\tilde{b} = 0$ .
  - Exploit the results in parts d.i) and d.ii) to demonstrate the global convergence claim. Discuss also the effect of the parameter  $\gamma_k$  on the speed of convergence of the algorithm.

**Solution 26**

- The optimization problem used in the proximal method is

$$\begin{aligned} x_{k+1} &= \operatorname{argmin}_x \left( \frac{1}{2} x' Q x + c' x + d + \frac{1}{2\gamma_k} (x - x_k)' (x - x_k) \right) \\ &= \operatorname{argmin}_x \frac{1}{2} x' \left( Q + \frac{1}{\gamma_k} I \right) x + \left( c' - \frac{1}{\gamma_k} x_k' \right) x + d + \frac{1}{2\gamma_k} x_k' x_k. \end{aligned}$$

Note that the function to be minimized is again a quadratic function. Hence, it has a unique minimizer if and only if the matrix  $\left( Q + \frac{1}{\gamma_k} I \right)$  is positive definite (which is always the case since  $Q > 0$  and  $\gamma_k > 0$ ).

- The solution of the optimization problem is

$$x^* = x_{k+1} = - \left( Q + \frac{1}{\gamma_k} I \right)^{-1} \left( c - \frac{1}{\gamma_k} x_k \right).$$

This can be rewritten as

$$x_{k+1} = \frac{1}{\gamma_k} \left( Q + \frac{1}{\gamma_k} I \right)^{-1} x_k - \left( Q + \frac{1}{\gamma_k} I \right)^{-1} c.$$

Hence,  $A = \frac{1}{\gamma_k} \left( Q + \frac{1}{\gamma_k} I \right)^{-1}$  and  $b = - \left( Q + \frac{1}{\gamma_k} I \right)^{-1} c$ .

c) The fixed point  $\bar{x}$  of the equation above is such that

$$\bar{x} = \frac{1}{\gamma_k} \left( Q + \frac{1}{\gamma_k} I \right)^{-1} \bar{x} - \left( Q + \frac{1}{\gamma_k} I \right)^{-1} c.$$

Multiplying on left by  $\left( Q + \frac{1}{\gamma_k} I \right)$  yields

$$\left( Q + \frac{1}{\gamma_k} I \right) \bar{x} = \frac{1}{\gamma_k} \bar{x} - c,$$

which gives

$$\bar{x} = -Q^{-1}c,$$

*i.e.* the global minimizer of the quadratic function.

d) i) Trivially

$$A = \frac{1}{\gamma_k} \left( Q + \frac{1}{\gamma_k} I \right)^{-1} = (\gamma_k)^{-1} \left( Q + \frac{1}{\gamma_k} I \right)^{-1} = (\gamma_k Q + I)^{-1}.$$

Observe now that  $A'A < I$ . Multiplying both sides with the matrix  $(A'A)^{-1}$  (which exists because  $A'A$  is positive definite) yields

$$I < (\gamma_k Q + I)(\gamma_k Q + I)',$$

hence

$$\gamma_k^2 Q'Q + \gamma_k(Q' + Q) > 0,$$

which holds by positivity of  $Q$  and  $\gamma_k$ .

ii) We add and subtract  $x^*$  and  $Ax^*$  to the equation in part b) obtaining

$$x_{k+1} - x^* = A(x_k - x^*) + b - x^* + Ax^*.$$

Defining  $\tilde{b} = b - x^* + Ax^*$ , it remains to prove that  $\tilde{b} = 0$ . Multiplying the expression of  $\tilde{b}$  on the left-hand side by  $\left( Q + \frac{1}{\gamma_k} I \right)$  yields

$$\begin{aligned} \left( Q + \frac{1}{\gamma_k} I \right) \tilde{b} &= \left( Q + \frac{1}{\gamma_k} I \right) \left[ \frac{1}{\gamma_k} \left( Q + \frac{1}{\gamma_k} I \right)^{-1} x^* - \left( Q + \frac{1}{\gamma_k} I \right)^{-1} c - x^* \right] \\ &= -c - \left( Q + \frac{1}{\gamma_k} I \right) x^* + \frac{1}{\gamma_k} I x^* = -c - Qx^*. \end{aligned}$$

The claim is proved substituting the minimizer  $x^* = -Q^{-1}c$  in the last equation.

iii) The equation

$$x_{k+1} - x^* = A(x_k - x^*)$$

is a linear difference equation in which all the eigenvalues of the dynamic matrix  $A$  have modulus strictly smaller than one. Hence, the state  $x_k - x^*$  converges globally to zero, *i.e.*  $x_k$  converges to the optimal solution  $x^*$ . The greater the value of  $\gamma_k$ , the smaller the modulus of the eigenvalues of  $A$ . Thus, increasing  $\gamma_k$  corresponds to a faster convergence of the algorithm.

**Exercise 27** The Levenberg-Marquardt algorithm is a modification of Newton's method for the solution of nonlinear equations. In the case of the scalar equation

$$f(x) = 0,$$

with  $x \in \mathbb{R}$  and  $f$  differentiable, the Levenberg-Marquardt algorithm can be written as (note that  $f'$  denotes the first derivative of  $f$  with respect to  $x$ )

$$x_{k+1} = x_k - 2 \frac{f(x_k)}{f'(x_k) + f'(\bar{x})}, \quad \bar{x} = x_k - \frac{f(x_k)}{f'(x_k)}.$$

Consider now the problem of minimizing the function

$$q(x) = \frac{x^4}{4} + \frac{4}{3}x^3 - 10x$$

(note that the global minimizer of the function  $q$  is  $x = 1.365230013$ ).

- Re-cast the considered minimization problem as the problem of finding the solution of a scalar equation.
- Write Newton's iteration for the solution of the equation determined in part a).
- Run four iterations of the Newton's iteration in part b) with  $x_0 = 3$  and evaluate the first four values of the sequence of the relative errors

$$RE_{k+1}^N = \frac{x_{k+1} - x^*}{(x_k - x^*)^3},$$

that is evaluate  $RE_1^N$ ,  $RE_2^N$ ,  $RE_3^N$  and  $RE_4^N$ . Hence argue that Newton's method does not have speed of convergence of order three. Explain why this is not un-expected.

- Write now the Levenberg-Marquardt algorithm for the solution of the equation determined in part a).  
(Hint: write the algorithm as two equations, that is do not substitute  $\bar{x}$  into the first equation of the algorithm.)
- Run four iterations of the Levenberg-Marquardt algorithm in part d) with  $x_0 = 3$  and evaluate the first four values of the sequence of the relative errors

$$RE_{k+1}^{LM} = \frac{x_{k+1} - x^*}{(x_k - x^*)^3}.$$

Hence argue that the Levenberg-Marquardt algorithm is faster than Newton's algorithm. (It is well-known that the Levenberg-Marquardt algorithm has, under similar assumptions o those required by Newton's method, speed of convergence of order three.)

### Solution 27

- The minimization problem can be re-cast as the problem of finding the stationary points of the function  $q$ , that is as the problem of solving the scalar equation

$$k(x) = x^3 + 4x^2 - 10 = 0.$$

As noted above, this equation has a solution at  $x = 1.365230013$ , which is actually the only solution. Note also that the second derivative of  $f$  at  $x = 1.365230013$  is positive, hence the point is a local minimizer.

- Newton's iteration for the solution of the equation  $k(x) = 0$  is

$$x_{k+1} = x_k - \frac{x_k^3 + 4x_k^2 - 10}{x_k(3x_k + 8)} = \frac{2x_k^3 + 4x_k^2 + 10}{x_k(3x_k + 8)}.$$

c) Setting  $x_0 = 3$  yields

$$\begin{array}{ll} x_1 = 1.960784314, & RE_1 = 0.1363174326, \\ x_2 = 1.486238507, & RE_2 = 0.5728643018, \\ x_3 = 1.371823522, & RE_3 = 3.721080242, \\ x_4 = 1.365251224, & RE_4 = 73.99652652. \end{array}$$

Since the sequence of the relative errors diverges the speed of convergence of the method is not of order three (note the *cube* in the denominator of the definition of the relative error). This is not un-expected, since under the given conditions one can only claim quadratic speed of convergence of Newton's method.

d) The Levenberg-Marquardt iteration is given by the two equations

$$x_{k+1} = x_k - 2 \frac{x_k^3 + 4x_k^2 - 10}{(x_k(3x_k + 8)) + (\bar{x}(3\bar{x} + 8))}, \quad \bar{x} = x_k - \frac{x_k^3 + 4x_k^2 - 10}{x_k(3x_k + 8)}.$$

e) Setting  $x_0 = 3$  yields

$$\begin{array}{ll} x_1 = 1.644853060, & RE_1 = 0.06400339279, \\ x_2 = 1.369582249, & RE_2 = 0.1990643754, \\ x_3 = 1.365230035, & RE_3 = 0.2668611603, \\ x_4 = 1.365230013, & RE_4 \approx 0. \end{array}$$

Since the sequence of the relative errors converges to zero, the speed of convergence of the Levenberg-Marquardt iteration is at least of order three, definitely faster than Newton's iteration.