

Olympiads School

Graph (II)

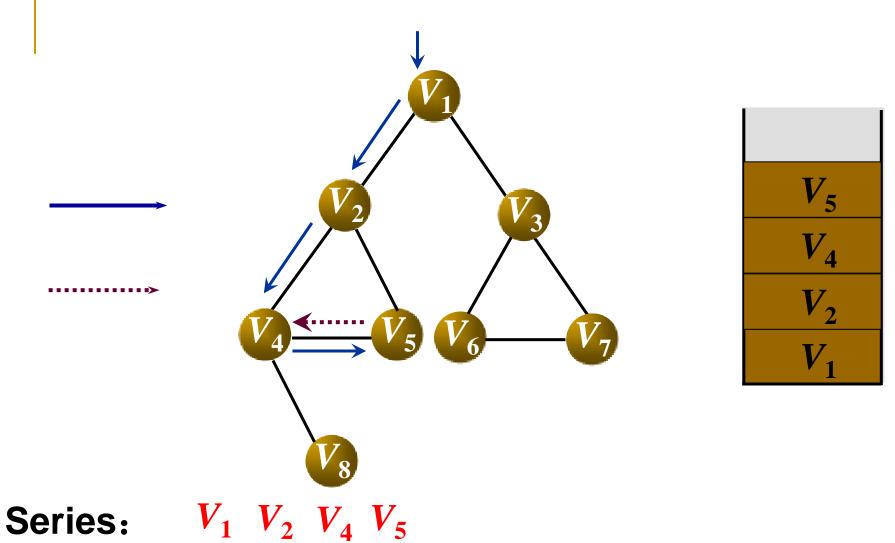
Bruce Nan

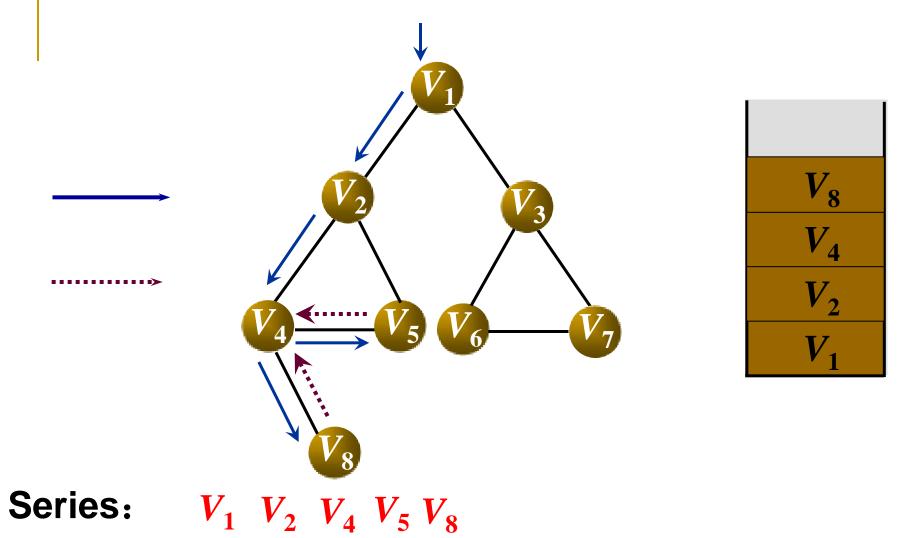
Depth Fist Search

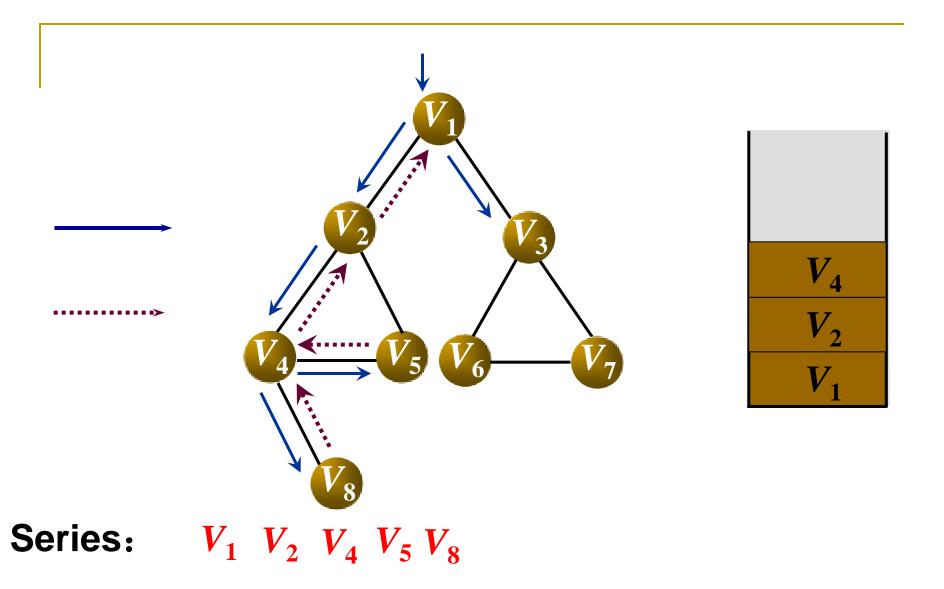
- Depth-first search uses the same idea as backtracking.
 - Exhaustively searching all possibilities by advancing if it is possible;
 - Backing up as soon as there is no unexplored possibility for further advancement.
- Recursive algorithms.

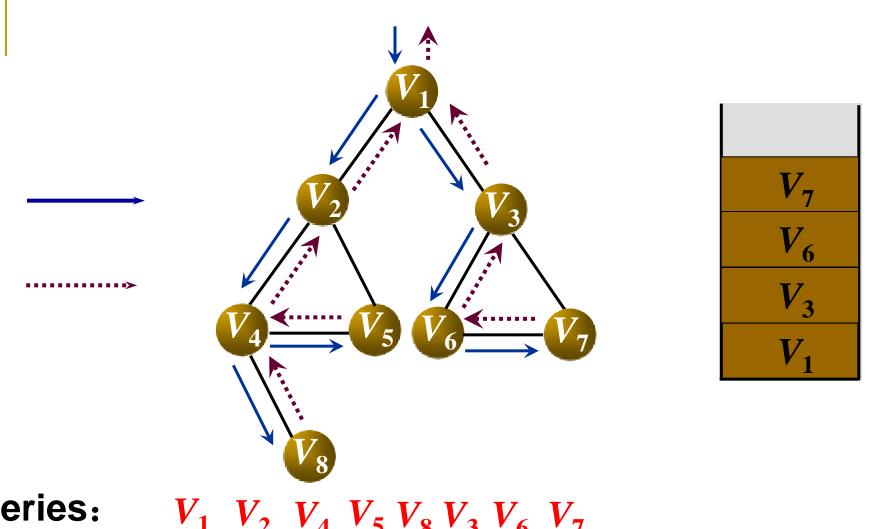
Depth First Search (DFS)

- A depth-first search (DFS) in an undirected graph G is like wandering in a labyrinth with a string and a can of red paint without getting lost.
- We start at vertex s, tying the end of our string to the point and painting s "visited". Next we label s as our current vertex called u.
- Now we travel along an arbitrary edge (u, v).
- If edge (u, v) leads us to an already visited vertex v we return to u.
- If vertex v is unvisited, we unroll our string and move to v, paint v "visited", set v as our current vertex, and repeat the previous steps.

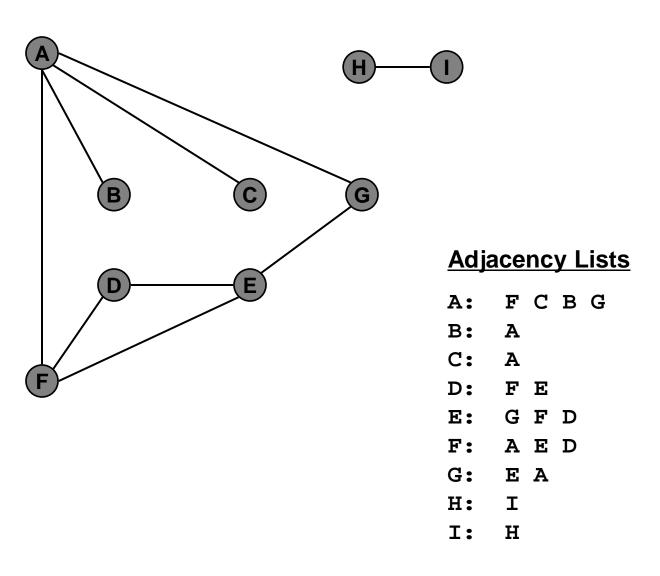


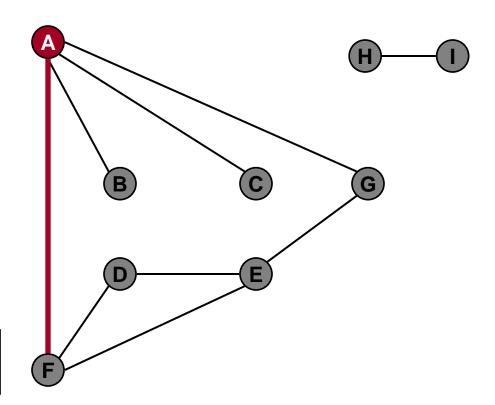






Series: V_1 V_2 V_4 V_5 V_8 V_3 V_6 V_7





F newly discovered

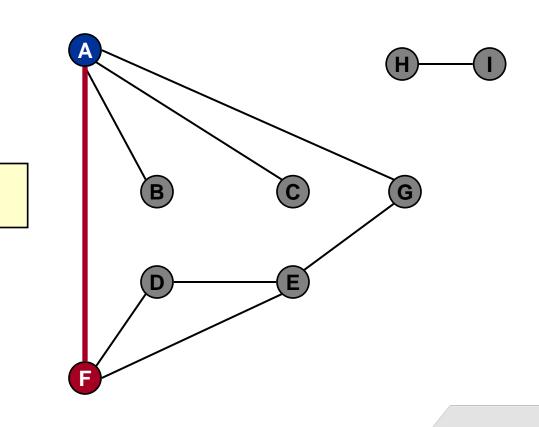
Undiscovered

Marked

Active

Finished

visit(A)
(A, F) (A, C) (A, B) (A, G)



Undiscovered

Marked

Active

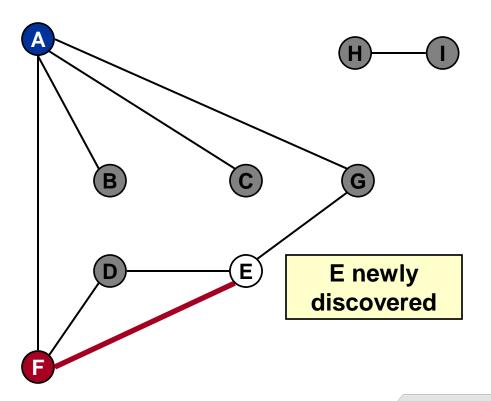
Finished

A already

marked

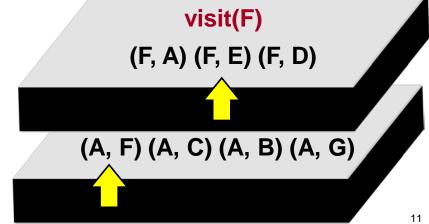
visit(F)
(F, A) (F, E) (F, D)

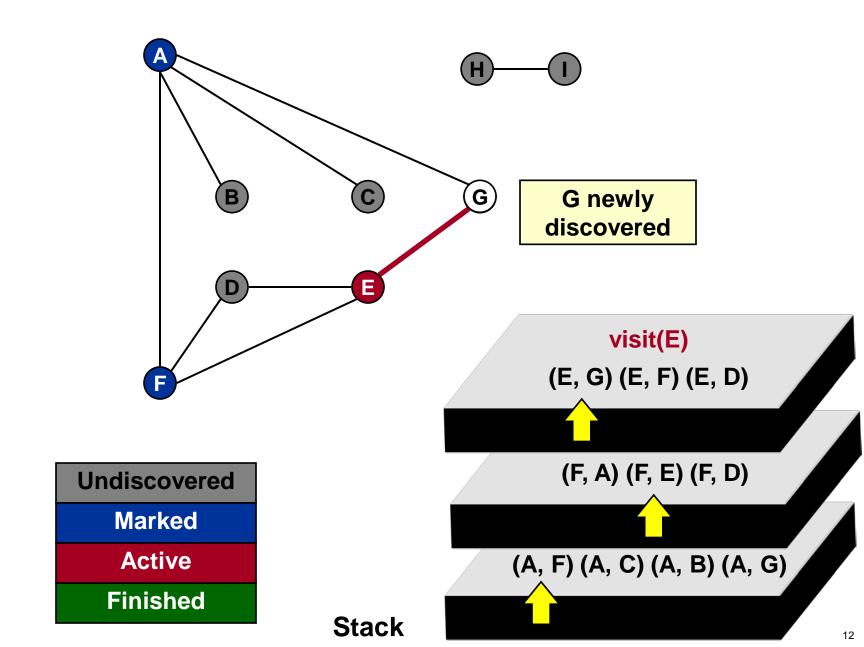
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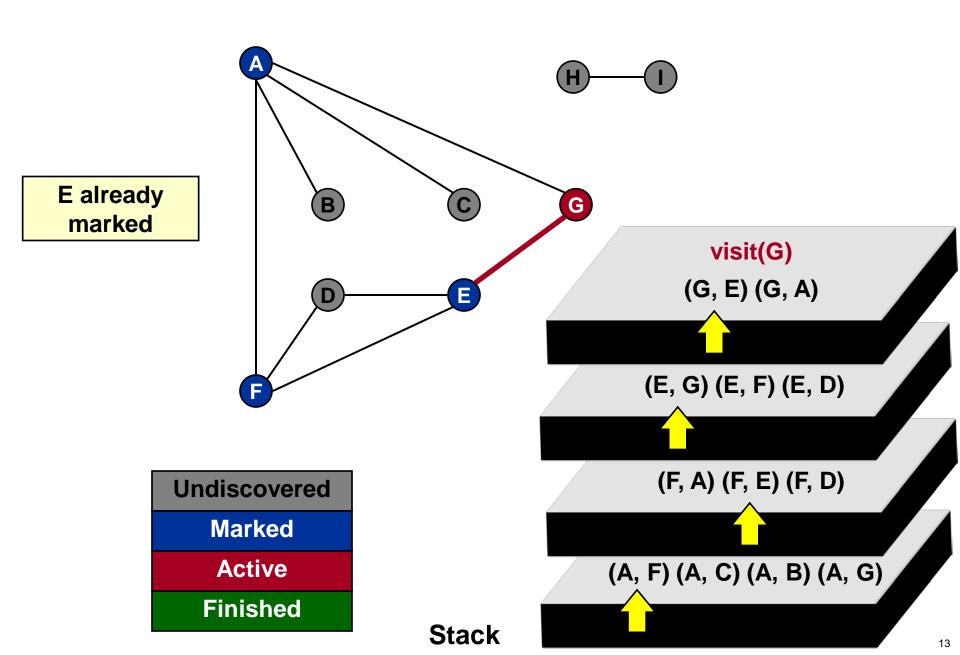


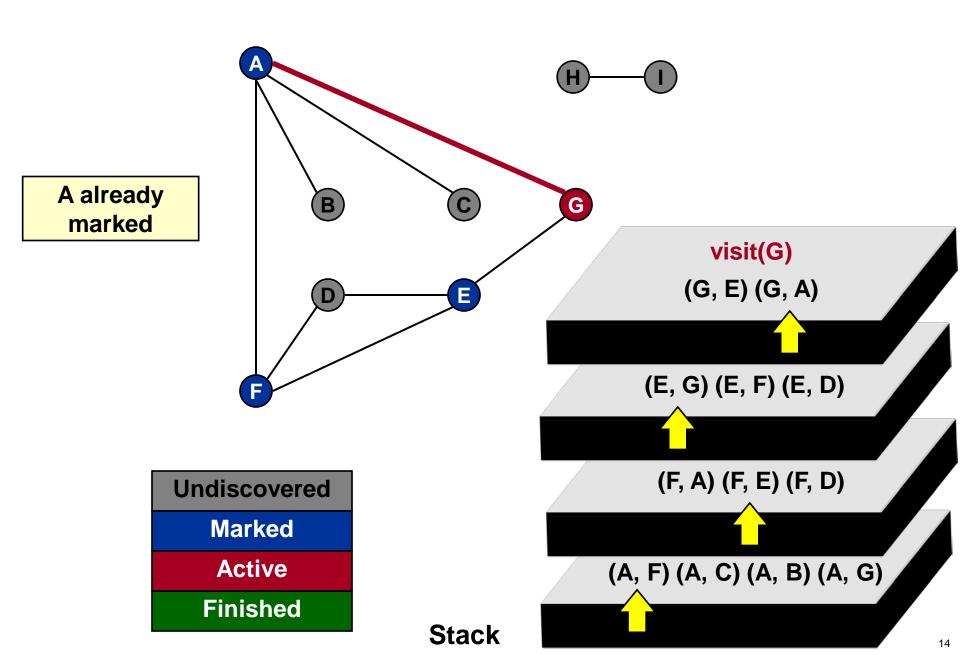
Stack

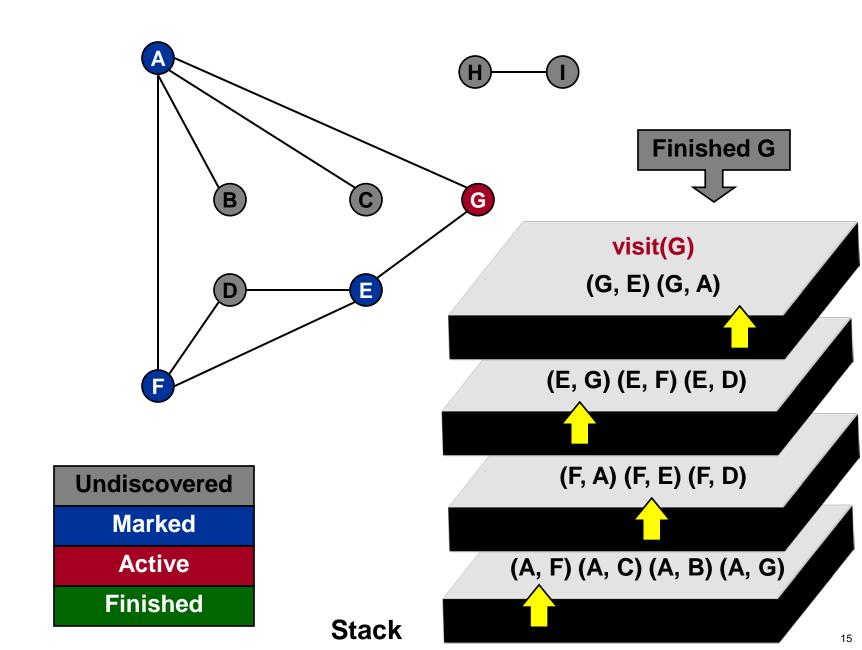


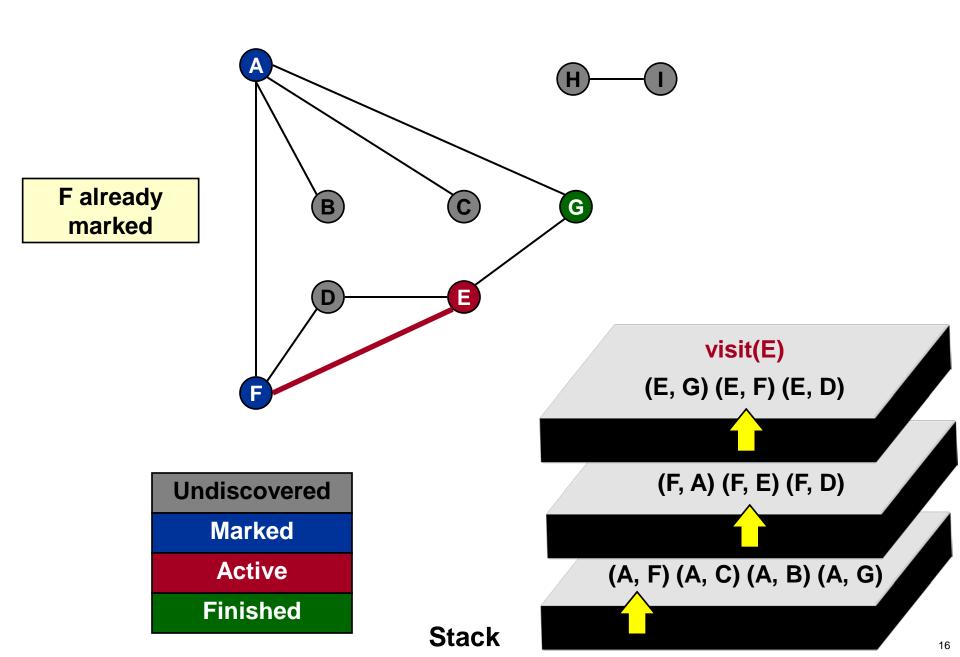


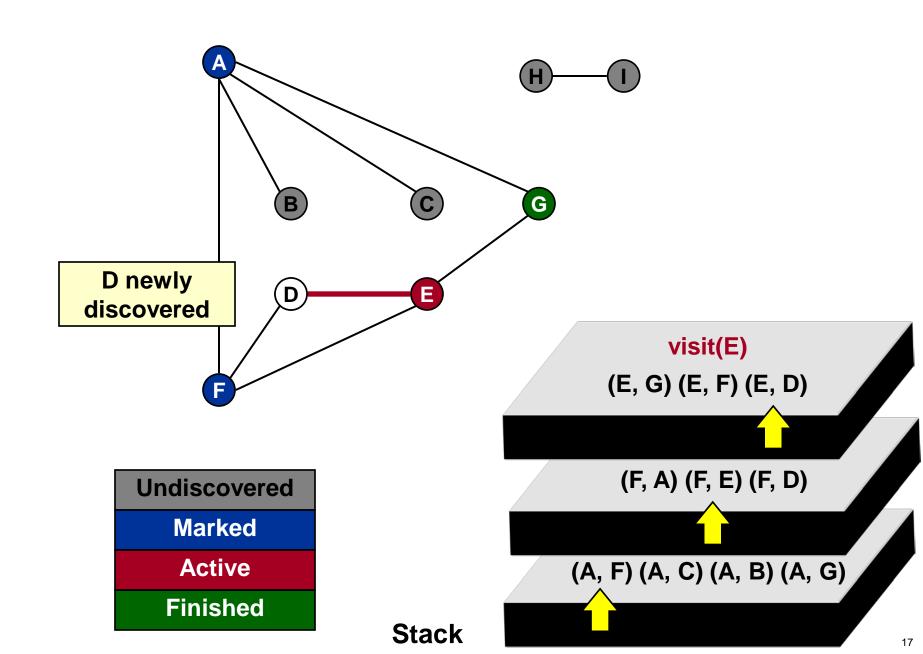


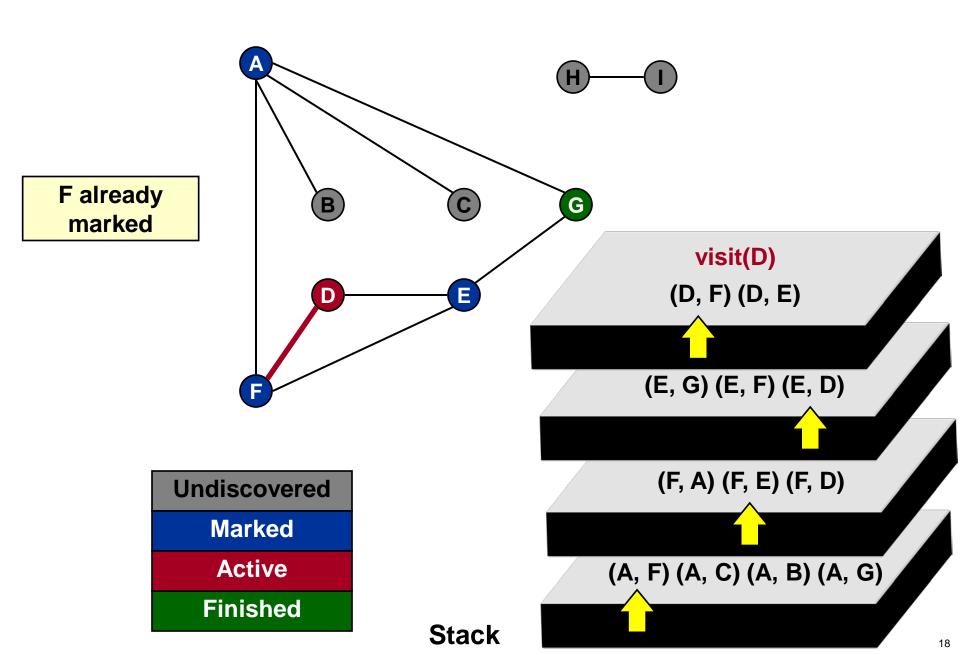


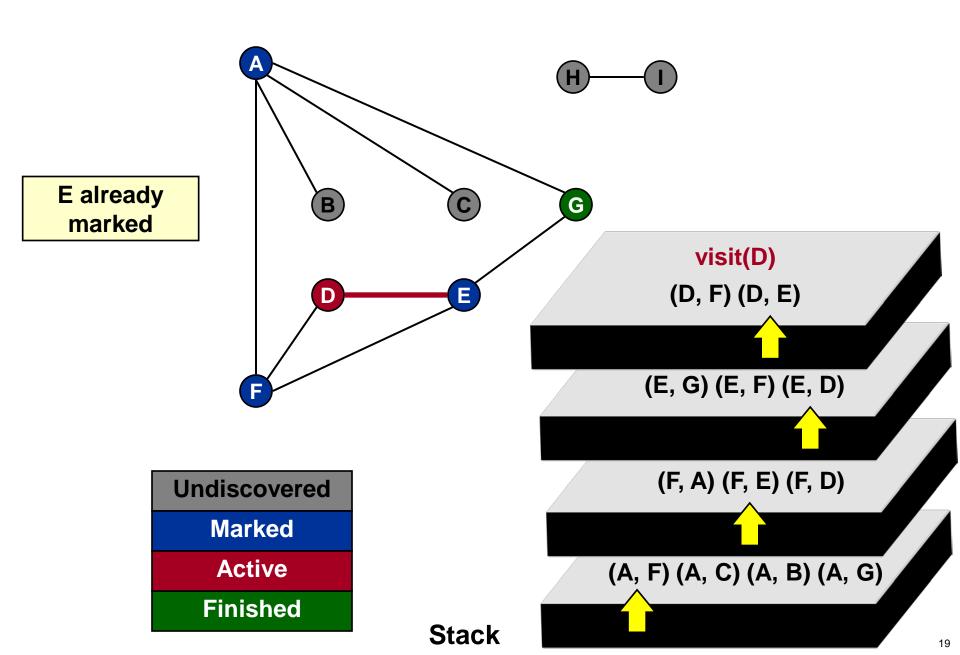


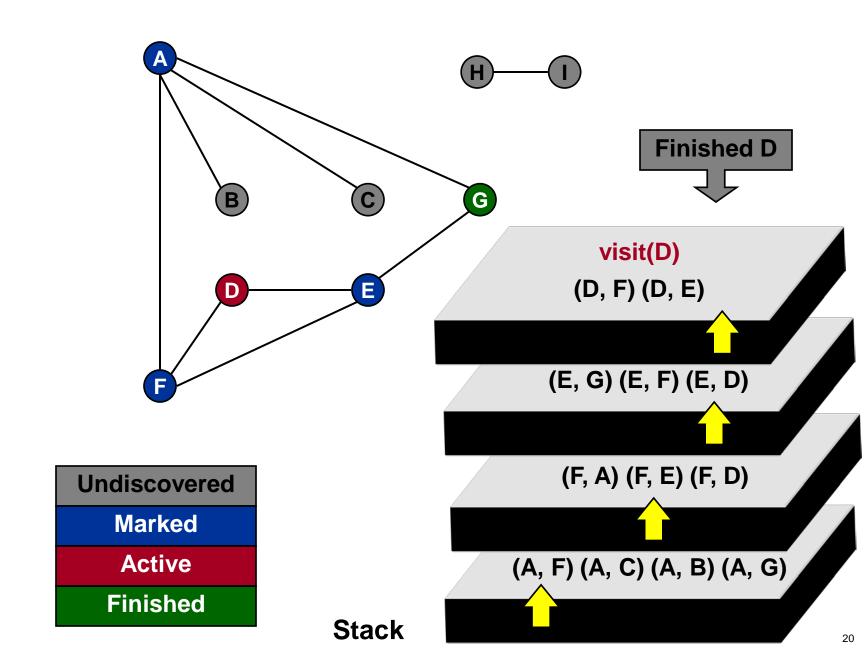


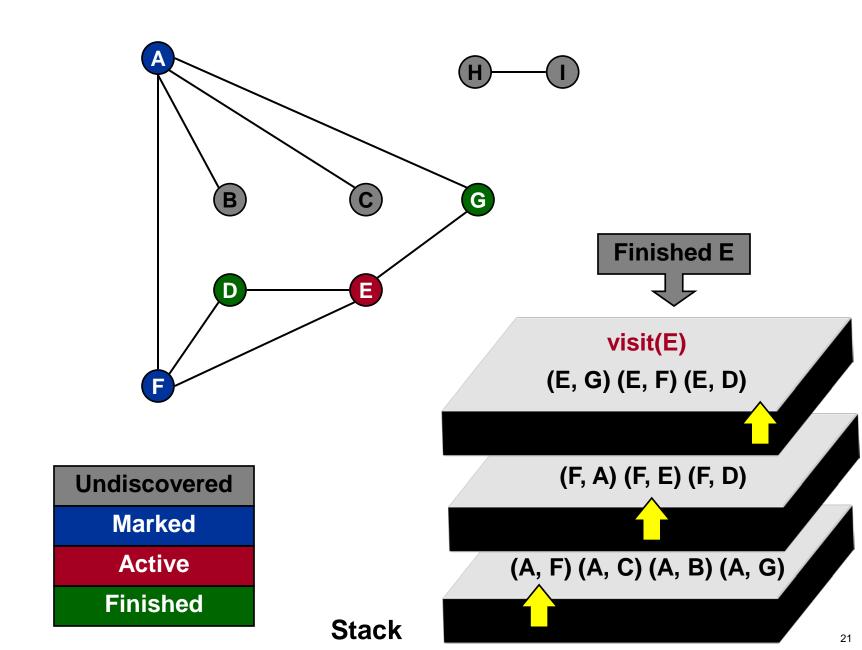


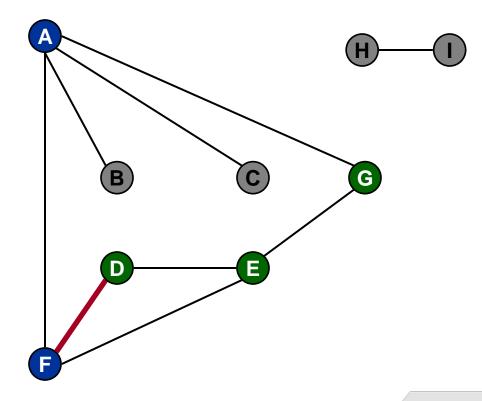












Undiscovered

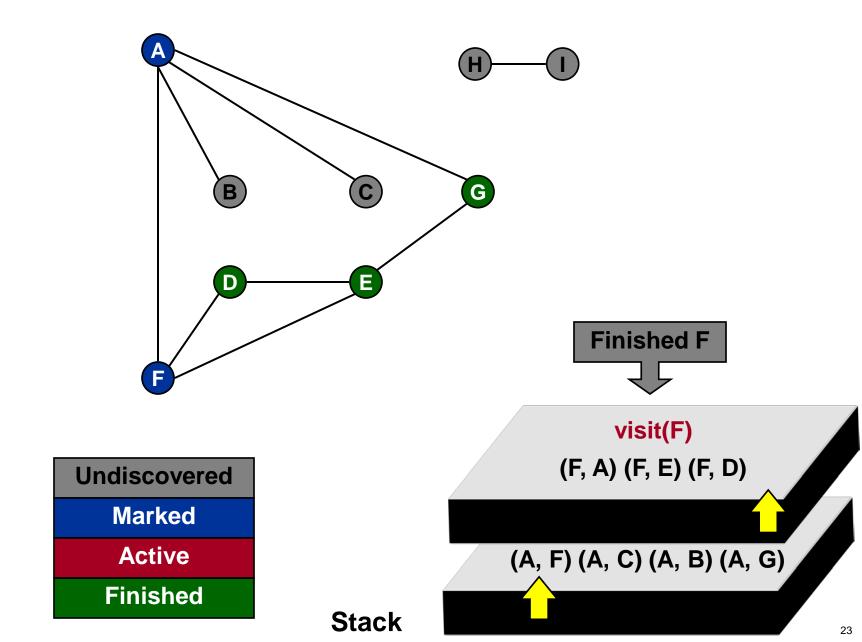
D already

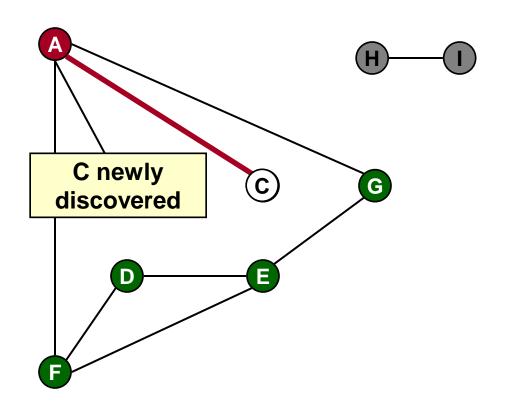
marked

Marked Active

Finished

visit(F)
(F, A) (F, E) (F, D)
(A, F) (A, C) (A, B) (A, G)

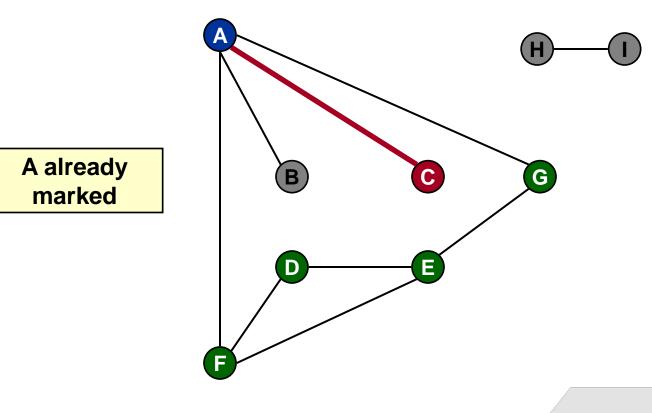




Stack



visit(A)
(A, F) (A, C) (A, B) (A, G)

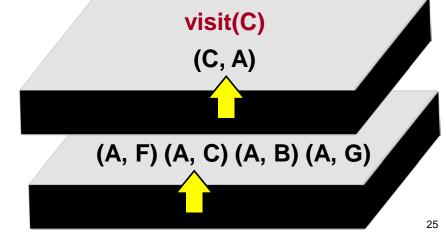


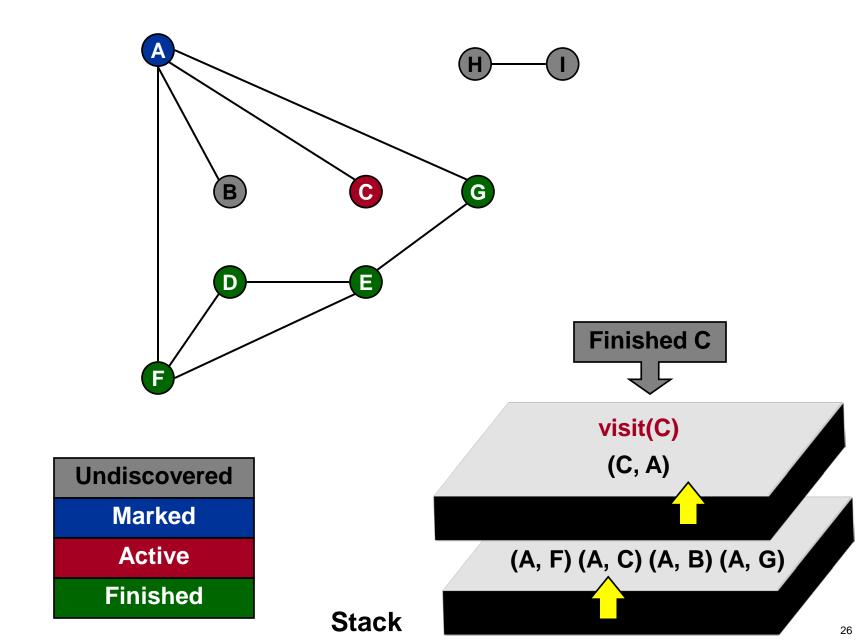
Undiscovered

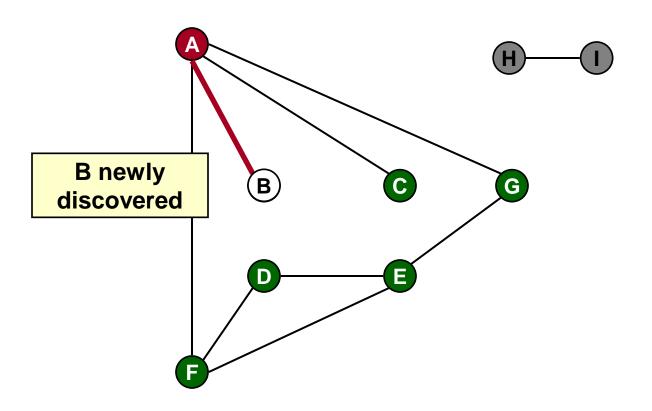
Marked

Active

Finished

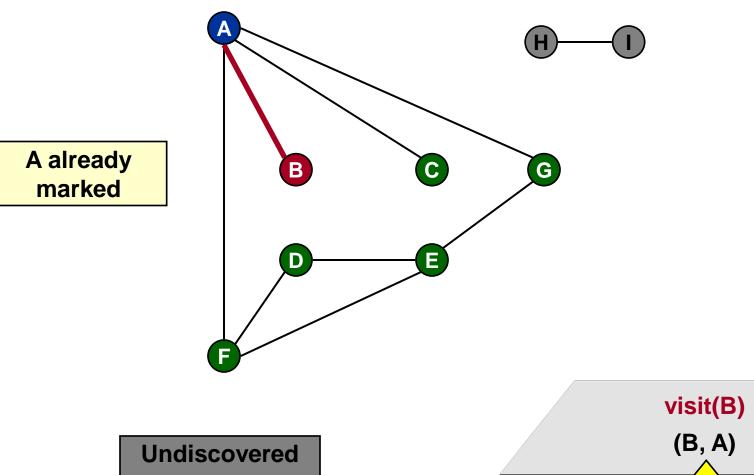








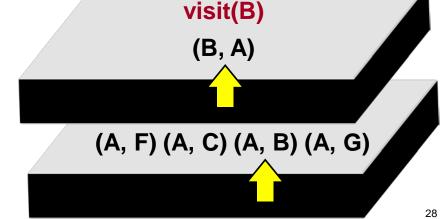
visit(A)
(A, F) (A, C) (A, B) (A, G)



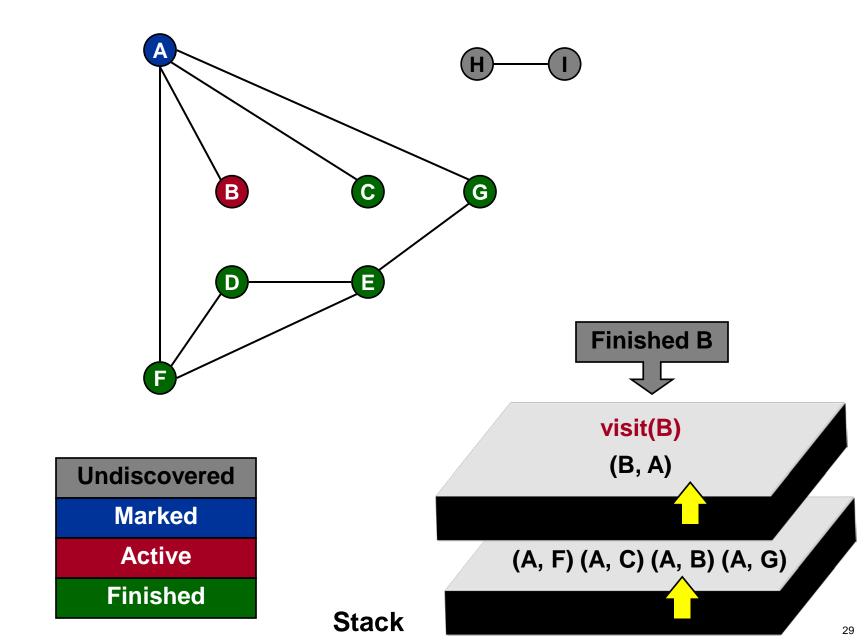
Marked

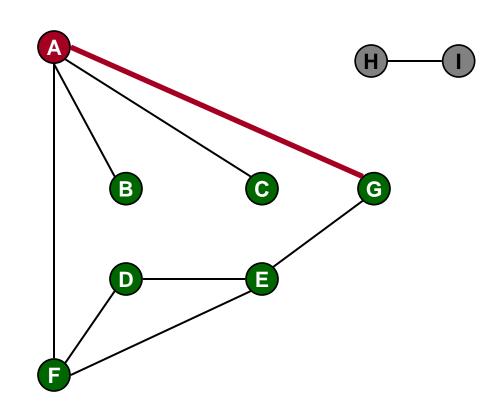
Active

Finished



Stack



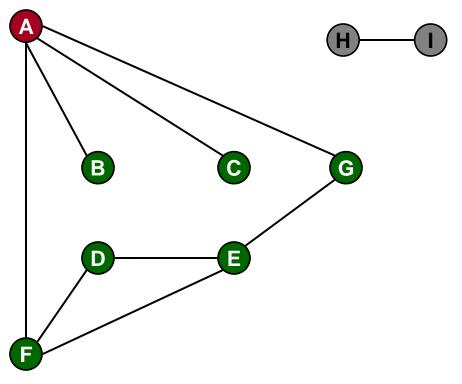


Undiscovered Marked **Active Finished**

G already

finished

visit(A) (A, F) (A, C) (A, B) (A, G) 30

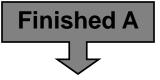


Undiscovered

Marked

Active

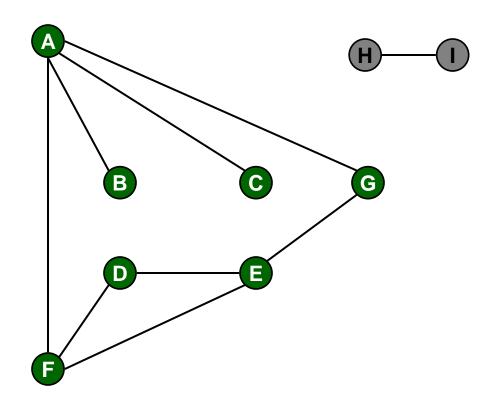
Finished



visit(A)

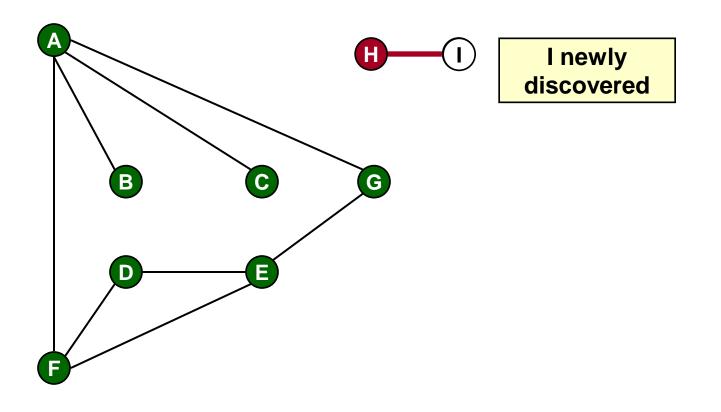
(A, F) (A, C) (A, B) (A, G)

Stack

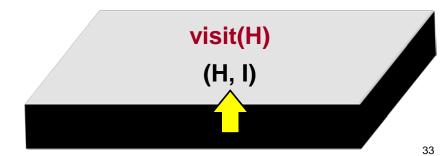


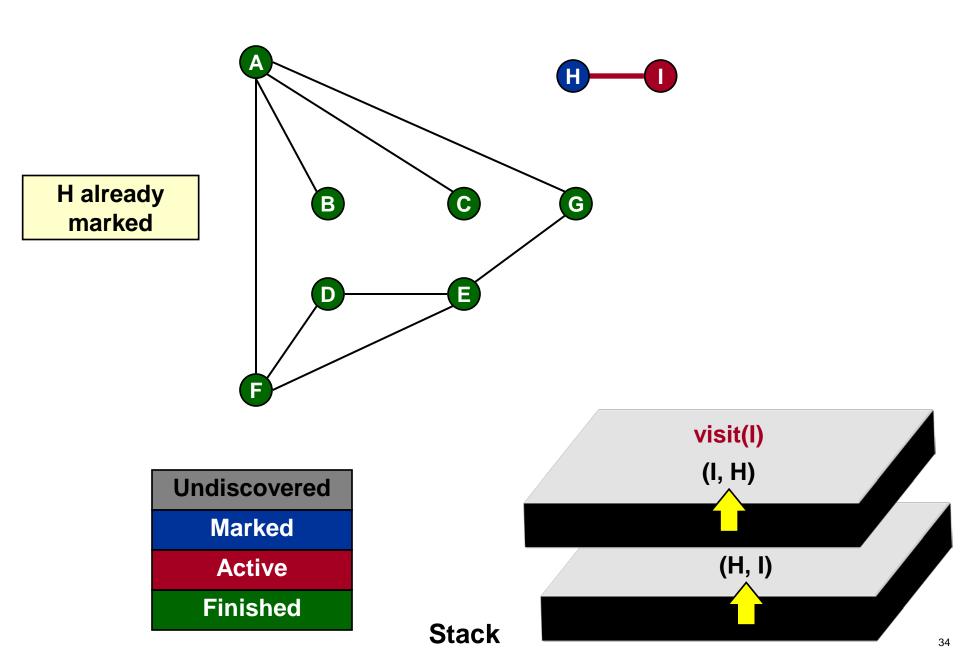


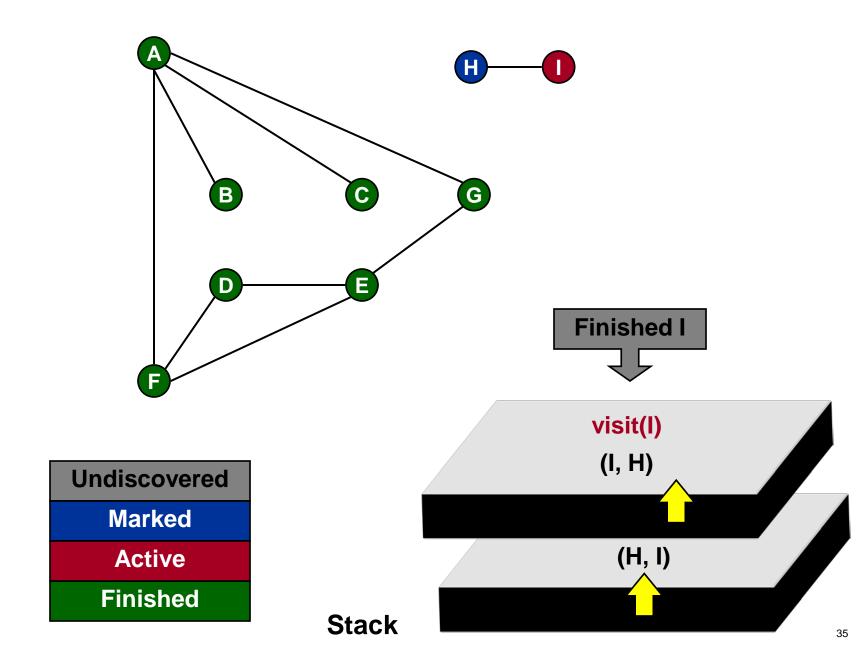
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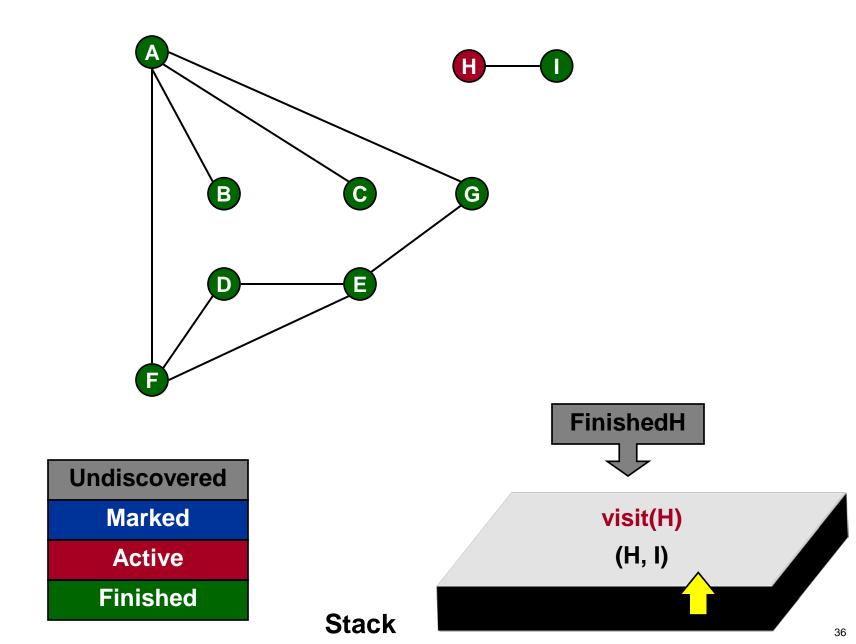


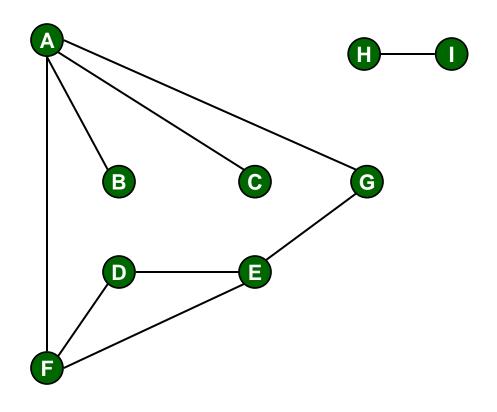














DFS Implementation

BFS: Queue

DFS: Stack

 The beauty of implementing dfs recursively is that recursion eliminates the need to keep an explicit stack

Depth-First Search Algorithm

 The beauty of implementing dfs recursively is that recursion eliminates the need to keep an explicit stack

```
Algorithm DFS(v); Input: A vertex v in a graph
Output: A labeling of the edges as "discovery" edges and "backedges"
for each edge e incident on v do
    if edge e is unexplored then let w be the other endpoint of e
    if vertex w is unexplored then label e as a discovery edge
        recursively call DFS(w)
    else label e as a backedge
```

DFS—Using Adjacent Matrix

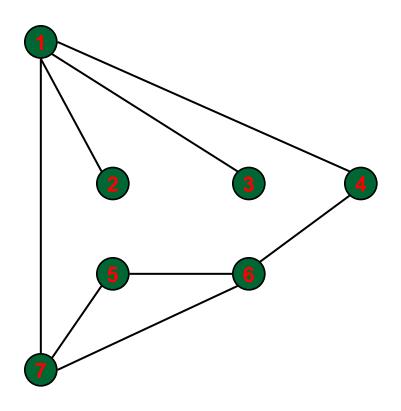
```
void DFSTraverse (int v)
{
    cout<<vertex[v]; visited [v]=1;
    for (j=0; j<vertexNum; j++)
        if (arc[v][j]==1 && visited[j]==0)
            DFSTraverse ( j );
}</pre>
```

Thinking problem

Finding cycle in an undirected graph

Given an undirected graph, how to find whether there is a cycle? Design your algorithm and evaluate its time complexity.

Practice



Input:

7 //number of vertices

8 //number of edges

1 2 //edge 1—2

13

46

56

57

67

Output:

Yes //there is cycle

Thank You