

Quadrature: Given  $f \in C([-1, 1])$ .  $I = \int_{-1}^1 f(x) dx = ?$

Distinct nodes  $x_0, \dots, x_n \in [-1, 1]$ ; weights  $w_0, \dots, w_n \in [-1, 1]$ .

$I \approx I_n = \sum_{k=0}^n w_k f(x_k)$ . Usually  $\{w_k\}$  come from poly. interp.

That is  $\int_{-1}^1 f(x) dx \approx \int_{-1}^1 \underset{\substack{\uparrow \\ \text{poly. interp.} \\ \text{of } f.}}{p_n(x)} dx = \sum_{k=0}^n w_k p(x_k) = \sum_{k=0}^n w_k f(x_k)$ .

### Three quadrature rules

- ① Newton-Cotes:  $\{x_j\}$  equispaced pts.  $\nexists |I - I_n|$  can diverge as  $n \rightarrow \infty$ . [Pölyo, 1933].
- ② Clenshaw-Curtis:  $\{x_j\} = \text{chebyshev pts.}$
- ③ Gauss quadrature:  $\{x_j\} = \text{Legendre pts, i.e., } \underset{\substack{\uparrow \\ \text{Legendre poly.}}}{P_{n+1/2}}(x_j) = 0$ .

Thm 19.1: For any quadrature rule from poly. interpolants,  
 $I = I_n$  if  $f \in P_n$

For Gauss,  $I = I_n$  if  $f \in P_{2n+1}$ .

Proof: If  $f \in P_{2n+1}$ , then  $f(x) = q_n(x)P_{n+1}(x) + r_n(x)$   $r_n, q_n \in P_n$ .

$$I = \int_{-1}^1 f(x) dx = \underbrace{\int_{-1}^1 q_n(x) P_{n+1}(x) dx}_{=0 \text{ by orthogonality}} + \int_{-1}^1 r_n(x) dx = \sum_{k=0}^n w_k r_n(x_k)$$

Also,  $f(x_j) = \underbrace{q_n(x_j) P_{n+1}(x_j)}_{=0} + r_n(x_j) = r_n(x_j)$ .

So,  $I = \int_1^1 f(x) dx = \sum_{k=0}^n w_k f(x_k)$ .

Implementation - Clenshaw - Curtis

$$\int_{-1}^1 T_k(x) dx = \begin{cases} 0 & \text{if } k \text{ is odd} \\ \frac{2}{1-k^2} & \text{if } k \text{ is even.} \end{cases}$$

Algorithm:

$$\int_{-1}^1 f(x) dx \approx \int_{-1}^1 \sum_{k=0}^n c_k T_k(x) dx = \sum_{k=0}^n c_k \int_{-1}^1 T_k(x) dx$$

get via DCT.

$$= \sum_{k=0}^{\lfloor n/2 \rfloor} c_{2k} \frac{2}{1-(2k)^2}$$

Clenshaw - Curtis can also be written as  $\int_{-1}^1 f(x) dx \approx \sum_{k=0}^n w_k f(x_k)$ .

Implementation - Gauss

[Golub - Welsch 1969]  $\{x_j\} = \text{Legendre pts: } P_{n+1}(x_j) = 0$

Idea: Find roots of  $P_{n+1}(x) = 0$  by eigenvalue problem.

$$P_0(x) = 1, \quad P_1(x) = x$$

$$(n+1)P_{n+1}(x) = (2n+1)xP_n(x) - nP_{n-1}(x), \quad n \geq 1.$$

$$\begin{bmatrix} 0 & 1 & & & \\ \frac{1}{2} & & & & \\ & 0 & \frac{n}{2} & & \\ & \frac{1}{2} & & & \\ & & 0 & \frac{n-1}{2} & \\ & & & & \ddots \\ & & & & & \frac{n}{2n+1} & 0 \end{bmatrix} \begin{bmatrix} P_0(x) \\ \vdots \\ P_n(x) \end{bmatrix} = \lambda \begin{bmatrix} P_0(x) \\ \vdots \\ P_n(x) \end{bmatrix}$$

Find eigenvalues of this!

Accuracy: How small is  $E_n(f) = I_n - I$ ?

Constructive:

$f = \text{Lipshitz cont.}$   $f(x) = \sum_{k=0}^{\infty} a_k T_k(x)$ .

So:  $E_n(f) = \sum_{k=0}^{\infty} a_k E_n(T_k) = \sum_{\substack{k=0 \\ k=\text{even}}}^{\infty} a_k E_n(T_k)$

CC:  $E_n(f) = \sum_{\substack{k=n+1 \\ k=\text{even}}}^{\infty} a_k E_n(T_k)$ , Gauss:  $E_n(f) = \sum_{\substack{k=2n+2 \\ k=\text{even}}}^{\infty} a_k E_n(T_k)$ .

Suppose  $f = \text{analytic}$  in Bernstein ellipse with  $\rho > 1$ , with  $|f(x)| \leq M$ .  
Then,  $|a_k| \leq 2M \rho^{-k}$ .

CC:  $|E_n(f)| \leq \sum_{\substack{k=n+1 \\ k=\text{even}}}^{\infty} |a_k| |E_n(T_k)|$   
 $\leq \sum_{\substack{k=n+1 \\ k=\text{even}}}^{\infty} \frac{2M \rho^{-k}}{\cancel{\rho^{-k}}} |E_n(T_k)|$   
 $\leq 8M \frac{\rho^{-n-1}}{1 - \rho^{-2}}.$

Gauss:  $|E_n(f)| \leq \frac{8M \rho^{-2n-2}}{1 - \rho^{-2}}.$

## Quadrature $\longleftrightarrow$ rational functions

$$I = \int_{-1}^1 f(x) dx, \quad I_n = \sum_{k=0}^n w_k f(x_k).$$

Claim: Every quadrature can be thought of as:

- ① Approximate  $f(x)$  by rational interpolant.
- ② Integrate rational interpolant exactly.

$$\text{Let } \int_{-1}^1 f(x) dx \approx \sum_{k=0}^n w_k f(x_k).$$

Take  $r(x) = \frac{p(x)}{q(x)}$  where  $p$  and  $q$  are of degree  $\leq n$ .

so that  $r(x_j) = f(x_j)$ .

### LAGRANGE FORM:

Lagrange func.  $l(x) = (x-x_0) \dots (x-x_n)$

Lagrange weights  $V_k = \frac{1}{\prod_{\substack{i=0 \\ i \neq k}}^n (x_k - x_i)}$

$$q(x) = l(x) \sum_{k=0}^n \frac{w_k}{x - x_k} q(x_k) \quad \left\{ \begin{array}{l} \text{rhs is a poly. of deg } \leq n. \\ \text{rhs interpolants } q(x) \text{ at } x_k. \end{array} \right.$$

Therefore, we know that  $q(x_j) r(x_j) = p(x_j)$ .

$$\text{Therefore } p(x) = l(x) \sum_{k=0}^n \frac{V_k}{x - x_k} q(x_k) r(x_k).$$

Hence,

$$r(x) = \frac{p(x)}{q(x)} = \frac{\sum_{k=0}^n \frac{\mu_k r(x_k)}{x - x_k}}{\sum_{k=0}^n \frac{\mu_k}{x - x_k}} \quad \mu_k = v_k q(x_k).$$

Therefore,

$$\begin{aligned} \int_{-1}^1 f(x) dx &\approx \int_{-1}^1 r(x) dx = \int_{-1}^1 \frac{\sum_{k=0}^n \frac{\mu_k r(x_k)}{x - x_k}}{\sum_{k=0}^n \frac{\mu_k}{x - x_k}} dx \\ &= \sum_{k=0}^n \left( \underbrace{\int_{-1}^1 \frac{\frac{\mu_k}{x - x_k}}{\sum_{k=0}^n \frac{\mu_k}{x - x_k}} dx}_{= w_k} \right) r(x_k) \end{aligned}$$

\* If you want to understand a quadrature rule, then first understand what rational interpolant it is building.

Example: Monte Carlo:

- ① Sample  $f(x)$  at  $x_0, \dots, x_n$  (random locations)
- ② Take  $\int_{\Omega} f(x) dx = |\Omega| \frac{1}{n+1} \sum_{k=0}^n f(x_k)$ .  
 $\uparrow$   
 vol. of  $\Omega$ .

Equivalent to:

① Approximate  $f(x)$  by a constant (using LS).

② Integrate constant exactly.

① Solve  $\begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}^T c = \begin{bmatrix} f(x_0) \\ \vdots \\ f(x_n) \end{bmatrix}$  in least squares sense.

$A$   $\underline{b}$

$$A^T A = n+1, \quad A^T b = \sum_{k=0}^n f(x_k), \quad A^T A c = A^T b$$
$$\text{so } c = \frac{1}{n+1} \sum_{k=0}^n f(x_k).$$

② Integrate const. exactly:

$$\int_{\Omega} f(\underline{x}) d\underline{x} \approx \int_{\Omega} c d\underline{x} = |\Omega| \cdot \frac{1}{n+1} \sum_{k=0}^n f(x_k)$$

Understand  $f(\underline{x}) \approx c$  by least squares...

High-order Monte Carlo in 1D

①

$$\begin{bmatrix} 1 & x_0 & x_0^2 \\ \vdots & \vdots & \vdots \\ 1 & x_N & x_N^2 \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ c_2 \end{bmatrix} = \begin{bmatrix} f(x_0) \\ \vdots \\ f(x_N) \end{bmatrix}$$

② Integrate  $c_0 + c_1x + c_2x^2$  exactly.