# A Dependency Pair Framework for Relative Termination of Term Rewriting

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$$\mathcal{R}_{\mathit{len}}$$
:  $\operatorname{len}(\operatorname{nil}) \to \mathcal{O}$   $\operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))$ 

$$\mathcal{R}_{\mathit{len}}$$
:  $\mathsf{len}(\mathsf{nil}) \to \mathcal{O}$   $\mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))$ 

 $\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,0,0])$ 

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
```

```
\begin{array}{c} & \mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & \mathsf{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & 1 + \mathsf{len}([0,0]) \end{array}
```

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
```

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 \begin{array}{c} & \mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & \mathsf{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & 1 + \mathsf{len}([0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \mathsf{s}(\mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & 2 + \mathsf{len}([0]) \end{array}
```

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
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\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 2 + \operatorname{len}([0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([\ ]) \end{array}
```

```
\mathcal{R}_{len}: \frac{\mathsf{len}(\mathsf{nil})}{\mathsf{len}(\mathsf{cons}(x,y))} \to \mathcal{O}
```

```
\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,0,0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([0,0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 2 + \operatorname{len}([0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([\ ]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\mathcal{O}))) & 3 \end{array}
```

Introduction (TRS)

```
\mathcal{R}_{\mathit{len}}: \mathsf{len}(\mathsf{nil}) \to \mathcal{O} \mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))
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\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,0,0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([0,0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 2 + \operatorname{len}([0]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([\ ]) \\ \to_{\mathcal{R}_{len}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\mathcal{O}))) & 3 \end{array}
```

#### **Termination**

 $\mathcal{R}$  is terminating iff there is no infinite evaluation  $t_0 \to_{\mathcal{R}} t_1 \to_{\mathcal{R}} \dots$ 

$$\mathcal{R}_{\mathit{len}}$$
:  $\mathsf{len}(\mathsf{nil}) \to \mathcal{O}$   $\mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))$ 

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y)) \mathcal{B}_{com}: \operatorname{cons}(x,\operatorname{cons}(y,xs)) \to \operatorname{cons}(y,\operatorname{cons}(x,xs))
```

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
```

 $\mathcal{B}_{com}$ :  $cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))$ 

$$[a,b]=[b,a]$$

```
\mathcal{R}_{len}: len(nil) \rightarrow \mathcal{O}
len(cons(x,y)) \rightarrow s(len(y))
\mathcal{B}_{com}: cons(x,cons(y,xs)) \rightarrow cons(y,cons(x,xs))
[a,b] = [b,a]
len(cons(\mathcal{O},cons(s(\mathcal{O}),cons(\mathcal{O},nil)))) \quad len([0,1,0])
```

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O}
\operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
\mathcal{B}_{com}: \operatorname{cons}(x,\operatorname{cons}(y,xs)) \to \operatorname{cons}(y,\operatorname{cons}(x,xs))
[a,b] = [b,a]
\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) \operatorname{len}([0,1,0])
\to_{\mathcal{R}_{len}} \operatorname{s}(\operatorname{len}(\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) 1 + \operatorname{len}([1,0])
```

```
len(nil) \rightarrow \mathcal{O}
\mathcal{R}_{len}:
                                   len(cons(x, y)) \rightarrow s(len(y))
\mathcal{B}_{com}:
                           cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))
                                                  [a, b] = [b, a]
                        len(cons(\mathcal{O}, cons(s(\mathcal{O}), cons(\mathcal{O}, nil)))) len([0, 1, 0])
                       s(len(cons(s(\mathcal{O}), cons(\mathcal{O}, nil))))
                                                                                    1 + \text{len}([1, 0])
                       s(len(cons(\mathcal{O}, cons(s(\mathcal{O}), nil))))
                                                                                           1 + \text{len}([0, 1])
                       s(s(len(cons(s(\mathcal{O}), nil))))
                                                                                           2 + len([1])

ightarrow_{\mathcal{R}_{\mathit{len}}}
                       s(s(s(len(nil))))
         \rightarrow_{\mathcal{R}_{\mathit{len}}}
                                                                                           3 + \text{len}([])
```

3 + len([])

3

### Relative Termination of TRSs

 $\to_{\mathcal{R}_{len}} \quad \mathsf{s}(\mathsf{s}(\mathsf{s}(\mathsf{len}(\mathsf{nil}))))$ 

 $s(s(s(\mathcal{O})))$ 

```
\mathcal{R}_{len}: \frac{\mathsf{len}(\mathsf{nil}) \to \mathcal{O}}{\mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))}
\mathcal{B}_{com}: \frac{\mathsf{cons}(x,\mathsf{cons}(y,xs)) \to \mathsf{cons}(y,\mathsf{cons}(x,xs))}{\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(s(\mathcal{O}),\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,1,0])}
\xrightarrow{\mathcal{R}_{len}} \frac{\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O}),\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,1,0])}{\mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,1])}
\xrightarrow{\mathcal{R}_{len}} \frac{\mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathsf{s}(\mathcal{O}),\mathsf{nil})))) \quad \mathsf{len}([0,1])}{\mathsf{s}(\mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathsf{s}(\mathcal{O}),\mathsf{nil})))) \quad \mathsf{len}([0,1])}
```

Introduction (TRS)

3

### Relative Termination of TRSs

```
\mathcal{R}_{len}:
                                                        \mathsf{len}(\mathsf{nil}) \ 	o \ \mathcal{O}
                                         len(cons(x, y)) \rightarrow s(len(y))
```

 $\mathcal{B}_{com}$ :  $cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))$ 

$$\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,1,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([1,0]) \\ \to_{\mathcal{B}_{\mathit{com}}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{nil})))) & 1 + \operatorname{len}([0,1]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\operatorname{s}(\mathcal{O}),\operatorname{nil})))) & 2 + \operatorname{len}([1]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([1]) \end{array}$$

[a, b] = [b, a]

#### Relative Termination

 $\rightarrow_{\mathcal{R}_{len}}$ 

 $\mathcal{R}/\mathcal{B}$  is terminating iff there is no infinite evaluation

 $s(s(s(\mathcal{O})))$ 

$$t_0 \to_{\mathcal{R}} \circ \to_{\mathcal{B}}^* t_1 \to_{\mathcal{R}} \circ \to_{\mathcal{B}}^* \dots$$

$$\mathcal{R}_{\mathit{len}}$$
:  $|\mathsf{len}(\mathsf{nil}) \rightarrow \mathcal{O} \\ |\mathsf{len}(\mathsf{cons}(x,xs)) \rightarrow \mathsf{s}(\mathsf{len}(xs))$ 

```
\mathcal{R}_{len}: \begin{array}{ccc} & \mathsf{len}(\mathsf{nil}) & \to & \mathcal{O} \\ & \mathsf{len}(\mathsf{cons}(x,xs)) & \to & \mathsf{s}(\mathsf{len}(xs)) \end{array}
```

Defined Symbols: len

```
\mathcal{R}_{len}: len(nil) \rightarrow \mathcal{O}
len(cons(x, xs)) \rightarrow s(len(xs))
```

Defined Symbols: len , Constructor Symbols: cons, nil, s,  $\mathcal O$ 

```
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Defined Symbols: len , Constructor Symbols: cons, nil, s, O

### $\operatorname{Sub}_D(r)$

 $\operatorname{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$ 

$$\mathcal{R}_{len}$$
:  $\frac{\mathsf{len}(\mathsf{nil})}{\mathsf{len}(\mathsf{cons}(x,xs))} \to \frac{\mathcal{O}}{\mathsf{s}(\mathsf{len}(xs))}$ 

Defined Symbols: len, Constructor Symbols: cons, nil, s, O

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$$Sub_{D}(s(len(xs))) = \{len(xs)\}$$

$$\mathcal{R}_{len}$$
:  $\operatorname{len}(\operatorname{nil}) \to \mathcal{O}$   $\operatorname{len}(\operatorname{cons}(x, xs)) \to \operatorname{s}(\operatorname{len}(xs))$ 

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 $\operatorname{Sub}_{\mathcal{D}}(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$ 

$$\operatorname{Sub}_{\mathcal{D}}(\mathsf{s}(\mathsf{len}(xs))) = \{\mathsf{len}(xs)\}$$

#### Dependency Pairs

If  $f(\ell_1,\ldots,\ell_n)\to r$  is a rule and  $g(r_1,\ldots,r_m)\in \mathrm{Sub}_D(r)$ , then  $f^\#(\ell_1,\ldots,\ell_n)\to g^\#(r_1,\ldots,r_m)$  is a dependency pair

$$\mathcal{R}_{len}$$
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 $\operatorname{Sub}_{\mathcal{D}}(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$ 

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```
\mathcal{DP}(\mathcal{R}_{len}):
```

$$\mathcal{R}_{len}$$
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$$\mathcal{DP}(\mathcal{R}_{len})$$
:  $\operatorname{len}^{\#}(\operatorname{cons}(x, xs)) \rightarrow \operatorname{len}^{\#}(xs)$ 

## Dependency Pairs Cont.

$$\begin{array}{c} \mathsf{len}(\mathsf{nil}) \to \mathcal{O} \\ \mathsf{len}(\mathsf{cons}(x,xs)) \to \mathsf{s}(\mathsf{len}(xs)) \end{array}$$

$$\mathsf{len}^\#(\mathsf{cons}(x,xs)) \to \mathsf{len}^\#(xs)$$

### Termination of $(\mathcal{D}, \mathcal{R})$

 $(\mathcal{D},\mathcal{R})$  is terminating iff there is no infinite evaluation

$$t_0 \rightarrow_{\mathcal{D}} \circ \rightarrow_{\mathcal{R}}^* t_1 \rightarrow_{\mathcal{D}} \circ \rightarrow_{\mathcal{R}}^* \dots$$

$$\mathsf{len}(\mathsf{nil}) o \mathcal{O} \ \mathsf{len}(\mathsf{cons}(x,xs)) o \mathsf{s}(\mathsf{len}(xs))$$

$$\mathsf{len}^\#(\mathsf{cons}(x,xs)) \to \mathsf{len}^\#(xs)$$

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#### Reminder: Relative Termination of $\mathcal{R}/\mathcal{B}$

 $\mathcal{R}/\mathcal{B}$  is terminating iff there is no infinite evaluation

$$t_0 \rightarrow_{\mathcal{R}} \circ \rightarrow_{\mathcal{B}}^* t_1 \rightarrow_{\mathcal{R}} \circ \rightarrow_{\mathcal{B}}^* \dots$$

## Dependency Pairs Cont.

$$\operatorname{len}(\operatorname{nil}) \to \mathcal{O}$$
  
 $\operatorname{len}(\operatorname{cons}(x,xs)) \to \operatorname{s}(\operatorname{len}(xs))$ 

$$\mathsf{len}^\#(\mathsf{cons}(x,xs)) \to \mathsf{len}^\#(xs)$$

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#### Theorem: Chain Criterion [Arts & Giesl 2000]

 $\mathcal{R}$  is terminating iff  $\mathcal{DP}(\mathcal{R})/\mathcal{R}$  is terminating

- Key Idea:
  - Transform a "big" problem into simpler sub-problems

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  - $\bullet$  DP problems  $(\mathcal{D},\mathcal{R})$  with  $\mathcal{D}$  a set of DPs,  $\mathcal{R}$  a TRS

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- How do we start?:
  - (Chain Criterion) Use all rules and dependency pairs:  $(\mathcal{DP}(\mathcal{R}), \mathcal{R})$

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  - DP processors:  $Proc(\mathcal{D}, \mathcal{R}) = \{(\mathcal{D}_1, \mathcal{R}_1), \dots, (\mathcal{D}_k, \mathcal{R}_k)\}$

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  - Proc is sound: if all  $(\mathcal{D}_i, \mathcal{R}_i)$  are terminating, then  $(\mathcal{D}, \mathcal{R})$  is terminating

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  - Proc is sound: if all  $(\mathcal{D}_i, \mathcal{R}_i)$  are terminating, then  $(\mathcal{D}, \mathcal{R})$  is terminating
  - Proc is complete: if  $(\mathcal{D}, \mathcal{R})$  is terminating, then all  $(\mathcal{D}_i, \mathcal{R}_i)$  are terminating

### **Timeline**



- 2000: DP Framework for termination [Arts & Giesl 2000, ...]
- 2006: Problem #106 of the RTA list of open problems
  - "Can we use the dependency pair method to prove relative termination?"
- 2016: Properties of  $\mathcal{R}/\mathcal{B}$  that allow to analyze the DP problem  $(\mathcal{DP}(\mathcal{R}), \mathcal{R} \cup \mathcal{B})$  [Iborra & Nishida & Vidal & Yamada 2016]
- 2023: Annotated Dependency Pairs for <u>Probabilistic</u> Rewriting [Kassing & Giesl 2023, Kassing & Dollase & Giesl 2024]
- 2024: Annotated Dependency Pairs for Relative Termination



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Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$ ?

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Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

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 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$ 

 $\mathcal{R}_2$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_2$ :

 $f \to d(a,f)$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$ 

 $\mathcal{R}_2$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_2$ :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$ 

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $a \rightarrow b$ 

 $\mathcal{B}_1$ :

 $b \rightarrow a$ 

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$ 

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a, f) \rightarrow_{\mathcal{R}_2} d(b, f) \rightarrow_{\mathcal{B}_2} d(b, d(a, f)) \rightarrow_{\mathcal{R}_2} \dots$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$ 

 $\mathcal{R}_2$ :

 $\mathsf{a} \to \mathsf{b}$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(\underline{b}, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $a \rightarrow b$ 

 $\mathcal{B}_1$ :

 $b \rightarrow a$ 

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$ 

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a, f) \rightarrow_{\mathcal{R}_2} d(b, f) \rightarrow_{\mathcal{B}_2} d(b, d(a, f)) \rightarrow_{\mathcal{R}_2} \dots$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$ 

 $\mathcal{R}_2$ :

 $\mathsf{a} \to \mathsf{b}$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $a \rightarrow b$ 

 $\mathcal{B}_1$ :

 $\mathsf{b} \to \mathsf{a}$ 

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$ 

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

$$\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_2/\mathcal{B}_2$  not terminating

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

 $\mathcal{R}_1$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_1$ :

 $\mathsf{b}\to\mathsf{a}$ 

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$ 

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$ 

 $\mathcal{R}_2$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_2$ :

 $f \to d(a,f)$ 

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$ 

 $\mathcal{R}_2/\mathcal{B}_2$  not terminating, but  $\mathcal{DP}(\mathcal{R}_2)/\mathcal{R}_2 \cup \mathcal{B}_2$  terminating  $(\mathcal{DP}(\mathcal{R}_2) = \varnothing)$ 

Goal: DP approach better than  $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$  (Termination of  $\mathcal{R} \cup \mathcal{B}$ )

Sufficient to analyze  $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$ ?

$$\mathcal{R}_1$$
:

 $a \rightarrow b$ 

 $\mathcal{B}_1$ :

 $b \rightarrow a$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$  not terminating, but  $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$  terminating  $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$ 

$$\mathcal{R}_2$$
:

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_2/\mathcal{B}_2$  not terminating, but  $\mathcal{DP}(\mathcal{R}_2)/\mathcal{R}_2 \cup \mathcal{B}_2$  terminating  $(\mathcal{DP}(\mathcal{R}_2) = \varnothing)$ 

#### **Domination**

 $\mathcal{R}$  dominates  $\mathcal{B}:\Leftrightarrow$  no defined symbol of  $\mathcal{R}$  in a right-hand side of  $\mathcal{B}$ 

 $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$  $\mathcal{R}_3$ :  $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{R}_3$ :  $\mathsf{a} \to \mathsf{b}$ 

 $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$ 

 $\mathcal{R}_3$ :  $\mathbf{a} \to \mathbf{b}$   $\mathcal{B}_3$ :  $\mathbf{f}(\mathbf{x}) \to \mathbf{c}(\mathbf{x}, \mathbf{f}(\mathbf{x}))$ 

 $\underline{f(a)} \to_{\mathcal{B}_3} \underline{d(\underline{a},f(a))} \to_{\mathcal{R}_3} \underline{d(b,\underline{f(a)})} \to_{\mathcal{B}_3} \underline{d(b,d(\underline{a},f(a)))} \to_{\mathcal{R}_2} \dots$ 

 $\mathcal{R}_3$ :  $\mathbf{a} \to \mathbf{b}$   $\mathcal{B}_3$ :  $\mathbf{f}(x) \to \mathbf{c}(x, \mathbf{f}(x))$ 

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(\underline{b}, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$ 

 $\mathcal{R}_3$ :  $a \to b$   $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, \underline{d(\underline{a}, f(a))}) \to_{\mathcal{R}_2} \dots$ 

$$\mathcal{R}_3$$
:  $\mathbf{a} \to \mathbf{b}$   $\mathcal{B}_3$ :  $\mathbf{f}(x) \to \mathbf{c}(x, \mathbf{f}(x))$ 

$$\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_3$ :  $\mathsf{a} \to \mathsf{b}$ 

 $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

$$\frac{\mathsf{f}(\mathsf{a})}{\mathsf{b}_3} \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}}(\underline{\mathsf{a}})) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_3/\mathcal{B}_3$  not terminating

$$\mathcal{R}_3$$
:  $a \to b$   $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

$$\underline{\mathsf{f}(\mathsf{a})} \mathop{\rightarrow}_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \mathop{\rightarrow}_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}(\mathsf{a})}) \mathop{\rightarrow}_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \mathop{\rightarrow}_{\mathcal{R}_2} \ldots$$

 $\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$ 

 $\mathcal{R}_3$ :

 $\mathsf{a} o \mathsf{b}$ 

$$\mathcal{B}_3$$
:  $f(x) \to c(x, f(x))$ 

$$\underline{\mathsf{f}(\mathsf{a})} \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}(\mathsf{a})}) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_3/\mathcal{B}_3$  not terminating, but  $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3\cup\mathcal{B}_3$  terminating  $(\mathcal{DP}(\mathcal{R}_3)=\varnothing)$ 

#### **Duplication**

 $\mathcal{B}$  is duplicating  $:\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}: x$  occurs more often in r than in  $\ell$ .

 $\mathcal{R}_3$ :

 $\mathsf{a} o \mathsf{b}$ 

$$\mathcal{B}_3$$
:  $f(x) \to c(x, f(x))$ 

$$\underline{\mathsf{f}(\mathsf{a})} \mathop{\rightarrow_{\mathcal{B}_3}} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \mathop{\rightarrow_{\mathcal{R}_3}} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}(\mathsf{a})}) \mathop{\rightarrow_{\mathcal{B}_3}} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \mathop{\rightarrow_{\mathcal{R}_2}} \ldots$$

 $\mathcal{R}_3/\mathcal{B}_3$  not terminating, but  $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3\cup\mathcal{B}_3$  terminating  $(\mathcal{DP}(\mathcal{R}_3)=\varnothing)$ 

#### **Duplication**

 $\mathcal{B}$  is duplicating  $:\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}: x$  occurs more often in r than in  $\ell$ .

 $\mathcal{R}_3$ :

 $\mathsf{a} o \mathsf{b}$ 

$$\mathcal{B}_3$$
:  $f(x) \to c(x, f(x))$ 

 $f(a) \rightarrow_{\mathcal{B}_3} d(\underline{a}, f(a)) \rightarrow_{\mathcal{R}_3} d(b, f(a)) \rightarrow_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \rightarrow_{\mathcal{R}_2} \dots$  $\mathcal{R}_3/\mathcal{B}_3$  not terminating, but  $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3$  terminating  $(\mathcal{DP}(\mathcal{R}_3) = \varnothing)$ 

#### **Duplication**

 $\mathcal{B}$  is duplicating : $\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}$ : x occurs more often in r than in  $\ell$ .

#### DPs for Relative Termination [Iborra et al. 2016]

If  $\mathcal{R}$  dominates  $\mathcal{B}$  and  $\mathcal{B}$  is non-duplicating, then  $\mathcal{R}/\mathcal{B}$  is terminating iff  $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$  is terminating

$$\mathcal{R}_3$$
:  $a \to b$   $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

$$\frac{\mathsf{f}(\mathsf{a}) \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}}(\underline{\mathsf{a}})) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

#### **Duplication**

 $\mathcal{B}$  is duplicating : $\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}$ : x occurs more often in r than in  $\ell$ .

#### DPs for Relative Termination [Iborra et al. 2016]

If  $\mathcal{R}$  dominates  $\mathcal{B}$  and  $\mathcal{B}$  is non-duplicating, then  $\mathcal{R}/\mathcal{B}$  is terminating iff  $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$  is terminating

$$\mathcal{R}_{len}$$
:  $\mathcal{B}_{com}$ :
$$|\operatorname{len}(\operatorname{nil}) \to \mathcal{O}$$

$$|\operatorname{len}(\cos(x, xs)) \to \operatorname{s}(\operatorname{len}(xs)) \qquad \cos(x, \cos(y, xs)) \to \cos(y, \cos(x, xs))$$

$$\mathcal{R}_3$$
:  $a \to b$   $\mathcal{B}_3$ :  $f(x) \to c(x, f(x))$ 

$$\frac{f(\underline{\mathsf{a}}) \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(b,\underline{f(\underline{\mathsf{a}})}) \to_{\mathcal{B}_3} \mathsf{d}(b,\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

#### **Duplication**

 $\mathcal{B}$  is duplicating : $\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}$ : x occurs more often in r than in  $\ell$ .

#### DPs for Relative Termination [Iborra et al. 2016]

If  $\mathcal R$  dominates  $\mathcal B$  and  $\mathcal B$  is non-duplicating, then  $\mathcal R/\mathcal B$  is terminating iff  $\mathcal D\mathcal P(\mathcal R)/\mathcal R\cup\mathcal B$  is terminating

$$\mathcal{R}_{len}$$
:  $\mathcal{B}_{com}$ :  $len(nil) \to \mathcal{O}$   $len(cons(x,xs)) \to s(len(xs))$   $cons(x,cons(y,xs)) \to cons(y,cons(x,xs))$ 

 $\mathcal{R}_{\mathit{len}}/\mathcal{B}_{\mathit{com}}$  terminates  $\Leftrightarrow \mathcal{DP}(\mathcal{R}_{\mathit{len}})/\mathcal{R}_{\mathit{len}} \cup \mathcal{B}_{\mathit{com}}$  terminates

# Annotated Dependency Pairs

$$\mathcal{R}_2$$
:  $b o a$ 

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_2$ :  $\mathsf{a} \to \mathsf{b}$ 

 $\mathcal{B}_2$ :

 $\mathsf{b} \to \mathsf{a}$ 

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{egin{array}{l} a^\# o b^\# \ a o b \end{array}
ight\}$ 

 $\label{eq:AB2} \mathcal{A}(\mathcal{B}_2) \text{:} \qquad \left\{ \begin{array}{l} b^\# \to a^\# \\ b \to a \end{array} \right\}$ 

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $b \rightarrow a$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{egin{array}{l} \mathsf{a}^\# o \mathsf{b}^\# \ \mathsf{a} o \mathsf{b} \end{array}
ight\}$ 

 $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ egin{array}{l} b^\# 
ightarrow a^\# \ b 
ightarrow a \end{array} 
ight\}$ 

 $a^{\#}$ 

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $b \rightarrow a$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ egin{align*} \mathbf{a}^{\#} & \rightarrow \mathbf{b}^{\#} \\ \mathbf{a} & \rightarrow \mathbf{b} \end{array} 
ight\}$ 

 $\begin{array}{ccc} \mathcal{A}(\mathcal{B}_2) \text{:} & \left\{ \begin{array}{c} b^\# \to a^\# \\ b \to a \end{array} \right\} \end{array}$ 

$$\mathsf{a}^\# o_{\mathcal{A}(\mathcal{R}_1)}^{(\#)} \mathsf{b}^\#$$

 $\mathcal{R}_2$ :

 $a \rightarrow b$ 

 $\mathcal{B}_2$ :

 $b \rightarrow a$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

$$\mathcal{A}(\mathcal{R}_2)$$
:  $\left\{ egin{align*} \mathbf{a}^\# & \mathbf{b}^\# \ \mathbf{a} & \mathbf{b} \end{array} 
ight\}$ 

 $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ egin{array}{l} \mathsf{b}^\# o \mathsf{a}^\# \ \mathsf{b} o \mathsf{a} \end{array} 
ight\}$ 

$$\left\{\begin{array}{c} \mathsf{b}^{\#} \to \mathsf{a}^{\#} \\ \mathsf{b} \to \mathsf{a} \end{array}\right\}$$

$$\mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} \mathsf{a}^\#$$

 $\mathcal{R}_2$ :  $a \to b$ 

 $\mathcal{B}_2$ :

 $\mathsf{b}\to\mathsf{a}$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

$$\mathcal{A}(\mathcal{R}_2): \left\{ \begin{array}{l} \mathbf{a}^{\#} \to \mathbf{b}^{\#} \\ \mathbf{a} \to \mathbf{b} \end{array} \right\}$$

$$\mathcal{A}(\mathcal{B}_2)$$
:  $\left\{egin{array}{l} \mathsf{b}^\# o \mathsf{a}^\# \ \mathsf{b} o \mathsf{a} \end{array}
ight\}$ 

$$\mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} \mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

 $\mathcal{R}_2$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_2$ :

 $\mathsf{b}\to\mathsf{a}$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

$$\frac{\mathcal{A}(\mathcal{R}_2):}{\left\{ \begin{array}{l} a^{\#} \rightarrow b^{\#} \\ a \rightarrow b \end{array} \right\} }$$

$$egin{aligned} \mathcal{A}(\mathcal{B}_2) \colon & \left\{egin{aligned} \mathsf{b}^\# & o \mathsf{a}^\# \ \mathsf{b} & o \mathsf{a} \end{aligned}
ight\} \end{aligned}$$

$$a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} b^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

$$\mathsf{a} \to_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}$$

 $\mathcal{R}_2$ :

 $\mathsf{a}\to\mathsf{b}$ 

 $\mathcal{B}_2$ :

 $b \rightarrow a$ 

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

$$\mathcal{A}(\mathcal{R}_2)$$
:  $\left\{ egin{align*} \mathbf{a}^\# & \mathbf{b}^\# \ \mathbf{a} & \mathbf{b} \end{array} 
ight\}$ 

$$\mathcal{A}(\mathcal{B}_2)$$
:  $\left\{ egin{array}{ll} \mathsf{b}^\# o \mathsf{a}^\# \ \mathsf{b} o \mathsf{a} \end{array} 
ight\}$ 

$$a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} b^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

$$\mathsf{a}\to_{\mathcal{A}(\mathcal{R}_1)}\mathsf{b}$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \ o_{\mathcal{P}}^{(\#)} \circ ( o_{\mathcal{P}} \cup o_{\mathcal{S}})^* \ t_2 \ o_{\mathcal{P}}^{(\#)} \circ ( o_{\mathcal{P}} \cup o_{\mathcal{S}})^* \ \dots$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 $\mathcal{R}_2$ :  $\mathsf{a} o \mathsf{b}$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ : a o b  $\mathcal{B}_2$ : f o d(a,f)

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathsf{a} o \mathsf{b}$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$ 

$$\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^-) \\ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

 $f^{\#}$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^-) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

$$\mathsf{f}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^\#,\mathsf{f})$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^-) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

$$\mathsf{f}^{\#} \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^{\#},\mathsf{f}) \to^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \mathsf{d}(\mathsf{b},\mathsf{f})$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^-,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

f#

#### Relative $(\mathcal{P},\mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

 $\mathcal{R}_2$ :  $\mathsf{a} o \mathsf{b}$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$ 

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\boldsymbol{\mathsf{f}^{\#}} \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a},\mathsf{f}^{\#})$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^-,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

 $\mathsf{f}^\# \to^{(\#)}_{^{\mathcal{A}(\mathcal{R}_2)}} \mathsf{d}(\mathsf{a},\mathsf{f}^\#) \to_{\mathcal{A}(\mathcal{R}_2)} \mathsf{d}(\mathsf{b},\mathsf{f}^\#)$ 

#### Relative $(\mathcal{P},\mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 $\mathcal{R}_2$ :  $\mathsf{a} o \mathsf{b}$ 

 $\mathcal{B}_2$ :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$ 

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{egin{array}{l} \mathsf{a}^\# o \mathsf{b} \ \mathsf{a} o \mathsf{b} \end{array}
ight\}$ 

 $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^{-},\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

$$f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(\mathsf{a}, f^{\#}) \rightarrow_{\mathcal{A}(\mathcal{R}_2)} d(\mathsf{b}, f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \ldots$$

#### Relative $(\mathcal{P},\mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathsf{a} o \mathsf{b}$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$ 

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$ 

 $f^{\#}$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a^{\#}, f^{\#})$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $f \rightarrow d(a, f)$  $a \rightarrow b$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$  $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ 

 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a^{\#}, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, f^{\#})$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $\mathcal{B}_2$ :  $a \rightarrow b$  $f \rightarrow d(a, f)$ 

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$ 

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$  $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a^{\#}, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, d(a^{\#}, f^{\#}))$ 

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 $\mathcal{R}_2$ :  $a \rightarrow b$   $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$ :  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ 

 $\mathcal{A}(\mathcal{B}_2)$ :  $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ 

$$f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^{\#}, f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \mathsf{d}(\mathsf{b}, f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{b}, \mathsf{d}(\mathsf{a}^{\#}, f^{\#})) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \dots$$

#### Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$  is terminating iff there is no infinite evaluation

$$t_1 \, \, o_{\mathcal{P}}^{(\#)} \circ ( o_{\mathcal{P}} \cup o_{\mathcal{S}})^* \, \, t_2 \, \, \, o_{\mathcal{P}}^{(\#)} \circ ( o_{\mathcal{P}} \cup o_{\mathcal{S}})^* \, \ldots$$

 $\mathcal{R}_2$ :  $a \rightarrow b$   $\mathcal{B}_2$ :

 $f \rightarrow d(a, f)$ 

$$\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$$

$$\mathcal{A}(\mathcal{R}_2)$$
:  $\left\{ egin{align*} \mathsf{a}^\# &\to \mathsf{b} \ \mathsf{a} &\to \mathsf{b} \end{array} 
ight\}$ 

$$egin{aligned} \mathcal{A}(\mathcal{B}_2) \colon & \left\{ egin{aligned} \mathsf{f}^\# & \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \ \mathsf{f} & \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{aligned} 
ight\} \end{aligned}$$

$$f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^{\#}, \mathsf{f}^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \mathsf{d}(\mathsf{b}, \mathsf{f}^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{b}, \mathsf{d}(\mathsf{a}^{\#}, \mathsf{f}^{\#})) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \dots$$

#### Chain Criterion

For  $\mathcal{B}$  non-duplicating:  $\mathcal{R}/\mathcal{B}$  is terminating iff  $(\mathcal{A}(\mathcal{R}), \mathcal{A}(\mathcal{B}))$  is terminating

24/[4,3]

$$24/[4,3] = (24/4)/3$$

$$24/[4,3] = (24/4)/3 = 2$$

$$24/[4,3] = (24/4)/3 = 2$$

```
\mathcal{R}_{divL}:
                              minus (x, \mathcal{O}) \to x
               (a)
               (b) minus (s(x), s(y)) \rightarrow minus (x, y)
                             div (\mathcal{O}, s(y)) \to \mathcal{O}
                    div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
               (d)
                        divL (x, nil) \rightarrow x
               (e)
               (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4$$

```
\mathcal{R}_{divL}:
                              minus (x, \mathcal{O}) \to x
                (a)
               (b) minus (s(x), s(y)) \rightarrow minus (x, y)
               (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
               (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
               (e) \operatorname{divL}(x, \operatorname{nil}) \to x
               (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{divL}:
                               minus (x, \mathcal{O}) \to x
                (b) minus (s(x), s(y)) \rightarrow minus (x, y)
                (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
                (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                (e) \operatorname{divL}(x, \operatorname{nil}) \to x
                (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                  (a) minus (x, \mathcal{O}) \to x

(b) minus (s(x), s(y)) \to minus (x, y)
                  (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
                  (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                  (e) \operatorname{divL}(x, \operatorname{nil}) \to x
                  (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                  (g)
                          \operatorname{divL} (x, \operatorname{cons}(y, xs)) \to \operatorname{divL} (x, \operatorname{switch} (y, xs))
                  (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                             switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                      \begin{array}{lll} (a) & \text{minus } (x,\mathcal{O}) \to x \\ (b) & \text{minus } (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{minus } \ (x,y) \\ (c) & \text{div } (\mathcal{O},\mathsf{s}(y)) \to \mathcal{O} \end{array} 
                      (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                                     divL (x, nil) \rightarrow x
                      (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                     (g) divL (x, cons(y, xs)) \rightarrow divL (x, switch (y, xs))
                     (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                                 switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

$$\mathcal{R}_{\mathsf{divL}}:$$

$$(a) \qquad \mathsf{minus} \ (x,\mathcal{O}) \to x$$

$$(b) \qquad \mathsf{minus} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{minus} \ (x,y)$$

$$(c) \qquad \mathsf{div} \ (\mathcal{O},\mathsf{s}(y)) \to \mathcal{O}$$

$$(d) \qquad \mathsf{div} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\mathsf{div}(\mathsf{minus} \ (x,y),\mathsf{s}(y)))$$

$$(e) \qquad \mathsf{divL} \ (x,\mathsf{nil}) \to x$$

$$(f) \qquad \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL}(\mathsf{div} \ (x,y),xs)$$

$$\mathcal{B}_{\mathsf{com}}:$$

$$(g) \qquad \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL} \ (x,\mathsf{switch} \ (y,xs))$$

$$(h) \quad \mathsf{switch} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{cons}(y,\mathsf{switch} \ (x,xs))$$

$$(i) \qquad \mathsf{switch} \ (x,xs) \to \mathsf{cons}(x,xs)$$

$$24/[4,3] \rightarrow_{\mathcal{B}_{com}} 24/[\hat{4},3]$$

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

$$\mathcal{R}_{\mathsf{divL}} : \\ (a) & \mathsf{minus} \ (x, \mathcal{O}) \to x \\ (b) & \mathsf{minus} \ (\mathsf{s}(x), \mathsf{s}(y)) \to \mathsf{minus} \ (x, y) \\ (c) & \mathsf{div} \ (\mathcal{O}, \mathsf{s}(y)) \to \mathcal{O} \\ (d) & \mathsf{div} \ (\mathsf{s}(x), \mathsf{s}(y)) \to \mathsf{s}(\mathsf{div}(\mathsf{minus} \ (x, y), \mathsf{s}(y))) \\ \\ (e) & \mathsf{divL} \ (x, \mathsf{nil}) \to x \\ (f) & \mathsf{divL} \ (x, \mathsf{cons}(y, xs)) \to \mathsf{divL}(\mathsf{div} \ (x, y), xs) \\ \\ \mathcal{B}_{\mathsf{com}} : & (g) & \mathsf{divL} \ (x, \mathsf{cons}(y, xs)) \to \mathsf{divL} \ (x, \mathsf{switch} \ (y, xs)) \\ (h) & \mathsf{switch} \ (x, \mathsf{cons}(y, xs)) \to \mathsf{cons}(y, \mathsf{switch} \ (x, xs)) \\ (i) & & \mathsf{switch} \ (x, xs) \to \mathsf{cons}(x, xs) \\ \end{cases}$$

$$24/[4,3] \rightarrow_{\mathcal{B}_{com}} 24/[\hat{4},3] \rightarrow_{\mathcal{B}_{com}} 24/[3,\hat{4}]$$

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                (a) minus (x, \mathcal{O}) \to x

(b) minus (s(x), s(y)) \to minus (x, y)
                (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
                (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                            divL (x, nil) \rightarrow x
                (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                (g) divL (x, cons(y, xs)) \rightarrow divL (x, switch (y, xs))
                (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                        switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] \to_{\mathcal{B}_{com}} 24/[\hat{4},3] \to_{\mathcal{B}_{com}} 24/[3,\hat{4}] \to_{\mathcal{B}_{com}} 24/[3,4]$$

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{A}_1(\mathcal{R}_{\text{divL}}):
                               (a)
                                                     minus^{\#}(x,\mathcal{O}) \to x
                               (b) minus\#(s(x), s(y)) \rightarrow minus\#(x, y)
                               (c)
                                                      \operatorname{div}^{\#}(\mathcal{O}, s(v)) \to \mathcal{O}
                              (d1) \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}(\operatorname{minus}^{\#}(x, y), s(y)))
                              (d2) \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
                               (e)
                                                        divL^{\#}(x, nil) \rightarrow x
                              (f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
                              (f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
\mathcal{A}_2(\mathcal{B}_{com}):
                              (g)
                                         \operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs)) \to \operatorname{divL}^{\#}(x,\operatorname{switch}^{\#}(y,xs))
                              (h) switch \#(x, cons(y, xs)) \rightarrow cons(y, switch \#(x, xs))
                               (i)
                                                     switch^{\#}(x,xs) \rightarrow cons(x,xs)
```

#### Dependency Graph Processor

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ 

$$\mathcal{S}_2$$
: 
$$\left\{ \begin{array}{l} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{array} \right\}$$

$$\begin{array}{ll} \mathcal{P}_2 \colon & \left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\} & \mathcal{S}_2 \colon \left\{ \begin{array}{l} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{array} \right\} \end{array}$$

$$\mathit{Proc}_{\mathit{DG}}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

$$\mathit{Proc}_{\mathit{DG}}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \emptyset)$ 

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 | | \( \) (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

#### $(\mathcal{P}, \mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$ 

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$ 

$$S_2$$
:  $\left\{ f^\# \to d(a^\#, f^\#) \right\}$ 

$$\mathit{Proc}_{\mathit{DG}}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph:

#### $(\mathcal{P},\mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$ 

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$ 

$$\mathcal{S}_2$$
:  $\left\{ f^\# \to \mathsf{d}(\mathsf{a}^\#, f^\#) \right\}$   $f \to \mathsf{d}(\mathsf{a}, \mathsf{f})$ 

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

#### $(\mathcal{P}, \mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$ 

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ \\ (\text{sound \& complete}) \}$$

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{\mathbf{a}^\# \to \mathbf{b}}$$
 
$$\boxed{\mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#, \mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ \\ (\text{sound \& complete}) \}$$

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{\mathbf{a}^\# \to \mathbf{b}}$$
 
$$\boxed{\mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#, \mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^{\#}\sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^{\#}\sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ \\ (\text{sound \& complete}) \}$$

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{\mathtt{a}^\# \to \mathtt{b}} \longleftarrow \boxed{\mathtt{f}^\# \to \mathtt{d}(\mathtt{a}^\#, \mathtt{f}^\#)}$$

- $\bullet$  directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ \\ (\text{sound \& complete}) \}$$

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{ \mathsf{a}^\# \to \mathsf{b}} \qquad \boxed{ \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) } \qquad \boxed{ }$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\mathcal{P}_2$$
:  $\left\{ egin{aligned} \mathsf{a}^\# & \to \mathsf{b} \\ \mathsf{a} & \to \mathsf{b} \end{aligned} 
ight\}$   $\mathcal{S}_2$ :  $\left\{ egin{aligned} \mathsf{f}^\# & \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) \\ \mathsf{f} & \to \mathsf{d}(\mathsf{a}, \mathsf{f}) \end{aligned} 
ight\}$ 

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ \\ | \ \mathcal{Q} \in \mathtt{SCC}^{(\mathcal{P},\mathcal{S})}_{\mathcal{P}} \ \}$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$ 

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ 

$$\mathcal{S}_2$$
:  $\left\{ egin{array}{l} f^\# 
ightarrow d(a^\#, f^\#) \\ f 
ightarrow d(a, f) \end{array} \right\}$ 

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC_{\mathcal{P}}^{(\mathcal{P}, \mathcal{S})} \}$$
(sound & complete)

$$Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing) = \{(\mathcal{S}_2, \mathcal{P}_2)\}$$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{ \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) }$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$ 

$$\mathcal{S}_2$$
:  $\left\{ f^\# \to d(a^\#, f^\#) \atop f \to d(a, f) \right\}$ 

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC_{\mathcal{P}}^{(\mathcal{P}, \mathcal{S})} \}$$
(sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$ 

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph: 
$$\boxed{\mathbf{f}^\# \to \mathsf{d}(\mathbf{a}^\#, \mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

 $S_2$ :  $\begin{cases} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{cases}$ 

$$\mathcal{P}_2$$
:  $\left\{egin{array}{l} \mathsf{a}^\# o \mathsf{b} \ \mathsf{a} o \mathsf{b} \end{array}
ight\}$ 

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC_{\mathcal{P}}^{(\mathcal{P}, \mathcal{S})} \}$$
(sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$ 

- $\bullet$  directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\begin{array}{ll} \mathcal{P}_2 \colon & \left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\} & \mathcal{S}_2 \colon \left\{ \begin{array}{l} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{array} \right\} \end{array}$$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, (\mathcal{S} \cap \mathcal{Q}) \cup \flat ((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC_{\mathcal{P}}^{(\mathcal{P}, \mathcal{S})} \}$$
(sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$ 

$$(\mathcal{A}(\mathcal{P}_2),\mathcal{S}_2)$$
-Dependency Graph: 
$$\boxed{\mathtt{a}^\# \to \mathtt{b}} \qquad \boxed{\mathtt{f}^\# \to \mathtt{d}(\mathtt{a}^\#,\mathtt{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$ 

$$S_2$$
:  $\left\{ f^\# \to d(a^\#, f^\#) \atop f \to d(a, f) \right\}$ 

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in \mathtt{SCC}^{(\mathcal{P},\mathcal{S})}_{\mathcal{P}} \cup \mathtt{Lasso} \}$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$ 

- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

$$\mathcal{P}_2$$
:  $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$ 

$$\mathcal{S}_2$$
:  $\left\{ f^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) \right\}$   $f \to \mathsf{d}(\mathsf{a}, \mathsf{f})$ 

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in \mathtt{SCC}^{(\mathcal{P},\mathcal{S})}_{\mathcal{P}} \cup \mathtt{Lasso} \}$$
 (sound & complete)

$$Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2) = \{(\mathcal{P}_2, \mathcal{S}_2)\}$$

$$(\mathcal{A}(\mathcal{P}_2), \mathcal{S}_2)$$
-Dependency Graph:
$$\mathbf{a}^\# \to \mathbf{b} \qquad \qquad \mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#, \mathbf{f}^\#)$$

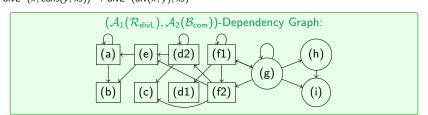
- ullet directed graph whose nodes are the ADPs from  $\mathcal{P} \cup \mathcal{S}$
- there is an arc from  $s \to t$  to  $v \to w$  iff  $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$  for substitutions  $\sigma_1, \sigma_2$ .

```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                          div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                        \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
(d1)
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                           (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                     \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                                      switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                           (i)
                            \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
 (e)
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```

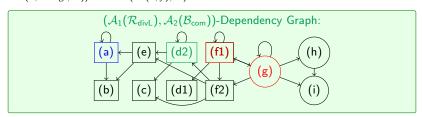
```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
                minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
 (b)
                          \operatorname{div}^{\#}(\mathcal{O}, \mathsf{s}(v)) \to \mathcal{O}
 (c)
                                                                                                                                          \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
(d1)
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                             (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                                        switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                             (i)
 (e)
                            divL^{\#}(x, nil) \rightarrow x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```

```
(\mathcal{A}_1(\mathcal{R}_{\text{divL}}), \mathcal{A}_2(\mathcal{B}_{\text{com}}))-Dependency Graph:
```

```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                          div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                     divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
                                                                                                                         (g)
(d1)
                     \operatorname{\mathsf{div}}^\#(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{\mathsf{div}}(\min \mathsf{us}^\#(x,y),\mathsf{s}(y)))
                                                                                                                         (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                                   switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                         (i)
 (e)
                            \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```

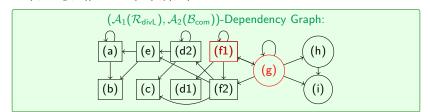


```
(a)
                        minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                         div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                 divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}(\operatorname{minus}^{\#}(x, y), s(y)))
                                                                                                                     (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                    \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                              switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                     (i)
 (e)
                           \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



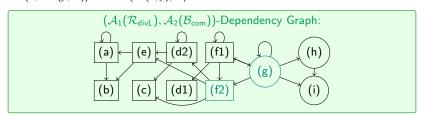
SCC:  $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f2)\}$ 

```
(a)
                       minus^{\#}(x,\mathcal{O}) \to x
 (b)
              minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                        div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                (g)
                                                                                                                            divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}(\operatorname{minus}^{\#}(x, y), s(y)))
                                                                                                                (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                   \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                         switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                 (i)
 (e)
                          \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



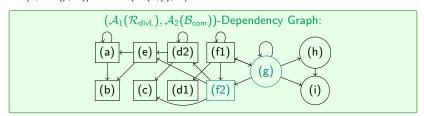
SCC:  $\{(a)\}$ ,  $\{(d2)\}$ , and  $\{(g), (f2)\}$ Lasso:  $\{(g), (f1)\}$  and  $\{(g), (f2)\}$ 

```
(a)
                       minus^{\#}(x,\mathcal{O}) \to x
 (b)
              minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                        div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                (g)
                                                                                                                            divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}(\operatorname{minus}^{\#}(x, y), s(y)))
                                                                                                                (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                   \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                         switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                 (i)
 (e)
                          \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



SCC:  $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f2)\}$ Lasso:  $\{(g), (f1)\}\$  and  $\{(g), (f2)\}\$ 

```
(a)
                     minus^{\#}(x,\mathcal{O}) \to x
             minus^{\#}(s(x), s(y)) \rightarrow minus (x, y)
 (b)
                      div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                    divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                  \operatorname{div}^{\#}(s(x), s(y)) \to s(\operatorname{div}(\min s(x, y), s(y)))
                                                                                                         (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                  \operatorname{div}^{\#}(s(x), s(y)) \to s(\operatorname{div} (\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                         (i)
 (e)
                        \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}(x, y), xs)
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```



SCC:  $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f2)\}$ Lasso:  $\{(g), (f1)\}\$  and  $\{(g), (f2)\}\$ 

 $(f2) \ \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad (g) \ \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$ 

$$(f2) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad (g) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$$

Find natural polynomial interpretation Pol

$$(f2) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad (g) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$$

Find natural polynomial interpretation Pol such that  $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \geq_{Pol}$  and

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) \qquad (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
    Find natural polynomial interpretation Pol such that \flat(\mathcal{P} \cup \mathcal{S}) \subseteq \triangleright_{Pol} and
         Pol(\text{divL}^{\#}(x, \text{cons}(y, xs))) \geq Pol(\text{divL}^{\#}(x, \text{switch}(y, xs))) + Pol(\text{switch}^{\#}(y, xs))
                                Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) > Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs))
```

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
    Find natural polynomial interpretation Pol such that \flat(\mathcal{P} \cup \mathcal{S}) \subseteq \triangleright_{Pol} and
         Pol(\text{divL}^{\#}(x, \text{cons}(y, xs))) \geq Pol(\text{divL}^{\#}(x, \text{switch}(y, xs))) + Pol(\text{switch}^{\#}(y, xs))
                                Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) > Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs))
```

$$Proc_{RP}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>})) \}$$
 (sound & complete)

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
     Find natural polynomial interpretation Pol such that \flat(\mathcal{P} \cup \mathcal{S}) \subseteq \triangleright_{Pol} and
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                                Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) > Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs))
                                       Proc_{RP}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>})) \}
                                                                    (sound & complete)
                                                       Proc_{RP}(\{(f2)\},\ldots)
```

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

Find natural polynomial interpretation *Pol* such that  $\flat(\mathcal{P} \cup \mathcal{S}) \subset \triangleright_{Pol}$  and

$$Pol(\operatorname{divL}^{\#}(x, \cos(y, xs))) \geq Pol(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs))) + Pol(\operatorname{switch}^{\#}(y, xs))$$

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$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>})) \} \\ \text{(sound \& complete)} \end{array}$$

$$Proc_{RP}(\{(f2)\},\ldots)$$

```
\operatorname{divL}_{Pol}^{\#}(x, xs) = xs \operatorname{switch}_{Pol}^{\#}(x, xs) = 0
cons_{Pol}(x, xs) = xs + 1 switch_{Pol}(x, xs) = xs + 1
```

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

Find natural polynomial interpretation *Pol* such that  $\flat(\mathcal{P} \cup \mathcal{S}) \subset \triangleright_{Pol}$  and

$$\begin{array}{ccc} & \textit{Pol}(\mathsf{divL}^\#(x,\mathsf{cons}(y,xs))) & \geq & \textit{Pol}(\mathsf{divL}^\#(x,\mathsf{switch}(y,xs))) + \textit{Pol}(\mathsf{switch}^\#(y,xs)) \\ & \times s + 1 & \geq & \times s + 1 \end{array}$$

$$\begin{array}{ccc} & Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) & > & Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)) \\ & & \times s + 1 & \geq & \times s \end{array}$$

$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{(\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>}))\} \\ \text{(sound \& complete)} \end{array}$$

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```
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```

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

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$$\begin{array}{ccc} & \textit{Pol}(\mathsf{divL}^\#(x,\mathsf{cons}(y,xs))) & \geq & \textit{Pol}(\mathsf{divL}^\#(x,\mathsf{switch}(y,xs))) + \textit{Pol}(\mathsf{switch}^\#(y,xs)) \\ & \times s + 1 & \geq & \times s + 1 \end{array}$$

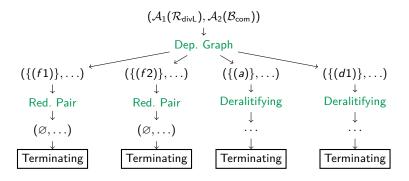
$$\begin{array}{ccc} & Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) & > & Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)) \\ & & \times s + 1 & \geq & \times s \end{array}$$

$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>})) \} \\ \text{(sound \& complete)} \end{array}$$

$$Proc_{RP}(\{(f2)\},\ldots)=\{(\varnothing,\ldots)\}$$

```
\operatorname{divL}_{Pol}^{\#}(x, xs) = xs \operatorname{switch}_{Pol}^{\#}(x, xs) = 0
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```

### Final Relative Termination Proof



⇒ Relative termination is proved automatically!

## Implementation and Experiments

Fully implemented in AProVE

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### **Relative rewriting** (130 benchmarks):

		new AProVE	NaTT	old AProVE	$T_TT_2$	MultumNonMulta
ſ	YES	91 (32)	68 (10)	48 (5)	39 (3)	0 (0)
	NO	13 (0)	5 (0)	13 (0)	7 (0)	13 (0)

### Implementation and Experiments

### Fully implemented in AProVE

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	new AProVE	NaTT	old AProVE	T <sub>T</sub> T <sub>2</sub>	MultumNonMulta
YES	91 (32)	68 (10)	48 (5)	39 (3)	0 (0)
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#### **Relative string rewriting** (403 benchmarks):

ſ		MultumNonMulta	Matchbox	AProVE	ADPs
Ī	YES	274	274	209	71

# Implementation and Experiments

### Fully implemented in AProVE

### Relative rewriting (130 benchmarks):

	new AProVE	NaTT	old AProVE	T <sub>T</sub> T <sub>2</sub>	MultumNonMulta
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ſ		MultumNonMulta	Matchbox	AProVE	ADPs
Ī	YES	274	274	209	71

### **Equational rewriting** (76 benchmarks):

		AProVE	MU-TERM	ADPs	
Ì	YES	66	64	36	l

• First DP framework specifically for relative termination

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- Annotated Dependency Pairs:

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  - o Reduction Pair Processor
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- Adapted the core processors from DP framework:
  - o Dependency Graph Processor o Usable Terms Processor
  - Reduction Pair Processor
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- Fully implemented in AProVE.
- Future Work:
  - Further Processors to (dis)-prove relative termination
  - Analyze further possibilities to use ADPs



$$\mathcal{R}_2$$
:  $\mathsf{a}(x) \to \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} \to \mathsf{a}(\mathsf{f})$ 

$$\mathcal{R}_2$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{a}(\mathsf{f})$ 

$$\underline{f} \to_{\mathcal{B}_2} \underline{a(f)} \to_{\mathcal{R}_2} b(\underline{f}) \to_{\mathcal{B}_2} b(\underline{a(f)}) \to_{\mathcal{R}_2} \dots$$

$$\mathcal{R}_2$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{a}(\mathsf{f})$   $\underbrace{\mathsf{f}} o \mathsf{b}_2 \underbrace{\mathsf{a}(\mathsf{f})} o \mathsf{e}_2 \underbrace{\mathsf{b}(\underline{\mathsf{f}})} o \mathsf{e}_2 \underbrace{\mathsf{b}(\underline{\mathsf{a}}(\mathsf{f}))} o \mathsf{e}_2 \dots$   $\mathcal{A}(\mathcal{R}_2)$ :  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{A}(\mathcal{B}_2)$ :  $\mathsf{f} o \mathsf{a}^\#(\mathsf{f}^\#)$ 

f#

$$\mathcal{R}_2$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{a}(\mathsf{f})$  
$$\underbrace{\mathsf{f} o_{\mathcal{B}_2} \ \mathsf{\underline{a}}(\mathsf{f}) o_{\mathcal{R}_2} \ \mathsf{b}(\underline{\mathsf{f}}) o_{\mathcal{B}_2} \ \mathsf{b}(\underline{\mathsf{a}}(\mathsf{f})) o_{\mathcal{R}_2} \dots}_{\mathcal{A}(\mathcal{R}_2)$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{A}(\mathcal{B}_2)$ :  $\mathsf{f} o \mathsf{a}^\#(\mathsf{f}^\#)$ 

$$\mathsf{f}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{a}^\#(\mathsf{f}^\#)$$

$$\mathcal{R}_2$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{a}(\mathsf{f})$   $\underbrace{\mathsf{f}} o_{\mathcal{B}_2} \underbrace{\mathsf{a}(\mathsf{f})} o_{\mathcal{R}_2} \mathsf{b}(\underline{\mathsf{f}}) o_{\mathcal{B}_2} \mathsf{b}(\underline{\mathsf{a}}(\mathsf{f})) o_{\mathcal{R}_2} \dots$ 

$$\mathcal{A}(\mathcal{R}_2)$$
:  $\mathsf{a}(x) \to \mathsf{b}(x)$   $\mathcal{A}(\mathcal{B}_2)$ :  $\mathsf{f} \to \mathsf{a}^\#(\mathsf{f}^\#)$ 

$$f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} a^{\#}(f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} b(f^{\#})$$

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:  $\mathsf{a}(x) \to \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} \to \mathsf{a}(\mathsf{f})$ 

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:  $\mathsf{a}(x) \to \mathsf{b}(x)$   $\mathcal{A}(\mathcal{B}_2)$ :  $\mathsf{f} \to \mathsf{a}^\#(\mathsf{f}^\#)$ 

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$$a(x) \rightarrow b(x)$$

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$$a(x) \rightarrow b(x)$$
  $a(x) \rightarrow b(x,x)$ 

$$\mathcal{R}_2$$
:  $\mathsf{a}(x) o \mathsf{b}(x)$   $\mathcal{B}_2$ :  $\mathsf{f} o \mathsf{a}(\mathsf{f})$ 

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$$a(x) \rightarrow b(x)$$
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$$a(x) \rightarrow b(x)$$
  $a(x) \rightarrow b(x,x)$   $a(x) \rightarrow b(x,x)$ 

#### Chain Criterion

For  $\mathcal B$  non-duplicating:  $\mathcal R/\mathcal B$  is terminating iff  $(\mathcal A(\mathcal R),\mathcal A(\mathcal B))$  is terminating

```
(f2) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

$$(f2)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \\ \qquad (g)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs)) \\$$

Find Com-monotonic and Com-invariant reduction pair  $(\succsim,\succ)$ 

#### Reduction Pair

- $\bullet$   $\succsim$  is reflexive, transitive, and closed under contexts and substitutions,
- ullet is a well-founded order and closed under substitutions
- $\bullet \succsim \circ \succ \circ \succsim \subseteq \succ.$

$$(f2)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \\ \qquad (g)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs)) \\$$

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#### Reduction Pair

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#### Com-monotonic

If  $s_1 \succ s_2$ , then  $\text{Com}_2(s_1,t) \succ \text{Com}_2(s_2,t)$  and  $\text{Com}_2(t,s_1) \succ \text{Com}_2(t,s_2)$ 

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#### Reduction Pair

- ullet is reflexive, transitive, and closed under contexts and substitutions,
- $\bullet \ \, \succsim \circ \, \succ \circ \, \succsim \subseteq \, \succ.$

#### Com-monotonic

If  $s_1 \succ s_2$ , then  $\text{Com}_2(s_1,t) \succ \text{Com}_2(s_2,t)$  and  $\text{Com}_2(t,s_1) \succ \text{Com}_2(t,s_2)$ 

#### Com-invariant

Let  $\sim = \succsim \cap \precsim$ , then

- $Com_2(s_1, s_2) \sim Com_2(s_2, s_1)$
- $\mathsf{Com}_2(s_1, \mathsf{Com}_2(s_2, s_3)) \sim \mathsf{Com}_2(\mathsf{Com}_2(s_1, s_2), s_3)$

```
 (f2) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(\mathsf{div}(x,y),x\mathsf{s}) \qquad \qquad (g) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,x\mathsf{s}))
```

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim \text{ and } \ell^{\#} \succsim \operatorname{ann}(r) \text{ for all } \ell \to r \in \mathcal{P} \cup \mathcal{S}$
- $\ell^\# \succ \operatorname{ann}(r)$  for all  $\ell \to r \in \mathcal{P}_{\succ}$

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim$  and  $\ell^\# \succsim \operatorname{ann}(r)$  for all  $\ell \to r \in \mathcal{P} \cup \mathcal{S}$
- $\ell^{\#} \succ \operatorname{ann}(r)$  for all  $\ell \rightarrow r \in \mathcal{P}_{\succ}$

```
\ell^{\#} \qquad \succeq \qquad \operatorname{ann}(r) \\ \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \qquad \succeq \qquad \operatorname{Com}_{2}(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs)), \operatorname{switch}^{\#}(y, xs))
```

```
(f2) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) \qquad \qquad (g) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim \text{ and } \ell^{\#} \succsim \operatorname{ann}(r) \text{ for all } \ell \to r \in \mathcal{P} \cup \mathcal{S}$
- $\ell^{\#} \succ \operatorname{ann}(r)$  for all  $\ell \rightarrow r \in \mathcal{P}_{\succ}$

$$\ell^{\#} \succeq \operatorname{ann}(r)$$

$$\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \succeq \operatorname{Com}_{2}(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs)), \operatorname{switch}^{\#}(y, xs))$$

$$\begin{aligned} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) &= \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \} \\ & (\text{sound \& complete}) \end{aligned}$$

```
(f2) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) \qquad \qquad (g) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

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$$\textit{Proc}_{\textit{RP}}(\{(f2)\}, \ldots)$$

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim$  and  $\ell^\# \succsim \operatorname{ann}(r)$  for all  $\ell \to r \in \mathcal{P} \cup \mathcal{S}$
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$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \} \\ \text{(sound \& complete)} \end{array}$$

$$Proc_{RP}(\{(f2)\},\ldots)$$

```
\begin{array}{llll} \operatorname{Com}_{2Pol}(x,y) & = & x+y & \operatorname{switch}_{Pol}^{\#}(x,xs) & = & 0 \\ \operatorname{cons}_{Pol}(x,xs) & = & xs+1 & \operatorname{switch}_{Pol}(x,xs) & = & xs+1 \\ \operatorname{divL}_{Pol}^{\#}(x,xs) & = & xs & \dots \end{array}
```

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim$  and  $\ell^\# \succsim \operatorname{ann}(r)$  for all  $\ell \to r \in \mathcal{P} \cup \mathcal{S}$
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```
\begin{array}{ccc} \ell^{\#} & \succeq & \operatorname{ann}(r) \\ \operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs)) & \succeq & \operatorname{Com}_2(\operatorname{divL}^{\#}(x,\operatorname{switch}(y,xs)),\operatorname{switch}^{\#}(y,xs)) \\ \operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs))) & \geq & \operatorname{\textit{Pol}}(\operatorname{Com}_2(\operatorname{divL}^{\#}(x,\operatorname{switch}(y,xs)),\operatorname{switch}^{\#}(y,xs))) \end{array}
```

$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \} \\ \text{(sound \& complete)} \end{array}$$

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```
\operatorname{Com}_{2Pol}(x,y) = x+y \quad \operatorname{switch}_{Pol}^{\#}(x,xs) = 0

\operatorname{cons}_{Pol}(x,xs) = xs+1 \quad \operatorname{switch}_{Pol}(x,xs) = xs+1

\operatorname{divL}_{Pol}^{\#}(x,xs) = xs \quad \dots
```

```
 (f2) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(\mathsf{div}(x,y),x\mathsf{s}) \qquad \qquad (g) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,x\mathsf{s}))
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```

$$Proc_{RP}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \}$$
(sound & complete)

$$Proc_{RP}(\{(f2)\},\ldots)$$

```
\operatorname{\mathsf{Com}}_{2Pol}(x,y) = x+y \quad \operatorname{\mathsf{switch}}_{Pol}^\#(x,xs) = 0 \\ \operatorname{\mathsf{cons}}_{Pol}(x,xs) = xs+1 \quad \operatorname{\mathsf{switch}}_{Pol}(x,xs) = xs+1 \\ \operatorname{\mathsf{divL}}_{Pol}^\#(x,xs) = xs \quad \dots
```

```
 (f2) \ \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad \qquad (g) \ \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))
```

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```
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```

$$Proc_{RP}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \}$$
 (sound & complete)

$$Proc_{RP}(\{(f2)\},\ldots)=\{(\varnothing,\ldots)\}$$

```
Com_2_{Pol}(x, y) = x + y switch_{Pol}^{\#}(x, xs) = 0

cons_{Pol}(x, xs) = xs + 1 switch_{Pol}(x, xs) = xs + 1

divL_{Pol}^{\#}(x, xs) = xs ...
```