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TeReSe, June 2024

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: $\mathsf{len}(\mathsf{nil}) \to \mathcal{O}$ $\mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))$

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 $\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,0,0])$

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
```

```
\begin{array}{c} & \mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & \mathsf{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \mathsf{s}(\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathcal{O},\mathsf{nil})))) & 1 + \mathsf{len}([0,0]) \end{array}
```

Introduction (TRS)

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```

```
\mathcal{R}_{len}: \frac{\mathsf{len}(\mathsf{nil})}{\mathsf{len}(\mathsf{cons}(x,y))} \to \mathcal{O}
```

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\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 2 + \operatorname{len}([0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([\ ]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\mathcal{O}))) & 3 \end{array}
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\begin{array}{ccc} & \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & \operatorname{len}([0,0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 1 + \operatorname{len}([0,0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{nil})))) & 2 + \operatorname{len}([0]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{nil})))) & 3 + \operatorname{len}([\ ]) \\ \to_{\mathcal{R}_{\mathit{len}}} & \operatorname{s}(\operatorname{s}(\operatorname{s}(\mathcal{O}))) & 3 \end{array}
```

Termination

 \mathcal{R} is terminating iff there is no infinite evaluation $t_0 \to_{\mathcal{R}} t_1 \to_{\mathcal{R}} \dots$

 $\mathcal{R}_{\mathit{len}}$: $\mathsf{len}(\mathsf{nil}) \to \mathcal{O}$ $\mathsf{len}(\mathsf{cons}(x,y)) \to \mathsf{s}(\mathsf{len}(y))$

```
\mathcal{R}_{len}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y)) \mathcal{B}_{com}: \operatorname{cons}(x,\operatorname{cons}(y,xs)) \to \operatorname{cons}(y,\operatorname{cons}(x,xs))
```

```
\mathcal{R}_{\mathit{len}}: \operatorname{len}(\operatorname{nil}) \to \mathcal{O} \operatorname{len}(\operatorname{cons}(x,y)) \to \operatorname{s}(\operatorname{len}(y))
```

 \mathcal{B}_{com} : $cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))$

$$[a,b]=[b,a]$$

```
\mathcal{R}_{len}: len(nil) \rightarrow \mathcal{O} len(cons(x,y)) \rightarrow s(len(y)) \mathcal{B}_{com}: cons(x,cons(y,xs)) \rightarrow cons(y,cons(x,xs))
```

$$[a,b] = [b,a]$$

$$\mathsf{len}(\mathsf{cons}(\mathcal{O},\mathsf{cons}(\mathsf{s}(\mathcal{O}),\mathsf{cons}(\mathcal{O},\mathsf{nil})))) \quad \mathsf{len}([0,1,0])$$

1 + len([0, 1])

Relative Termination of TRSs

 $s(len(cons(\mathcal{O}, cons(s(\mathcal{O}), nil))))$

```
\mathcal{R}_{len}: \qquad \qquad \operatorname{len}(\operatorname{nil}) \rightarrow \mathcal{O} \\ \operatorname{len}(\operatorname{cons}(x,y)) \rightarrow \operatorname{s}(\operatorname{len}(y))
\mathcal{B}_{com}: \qquad \operatorname{cons}(x,\operatorname{cons}(y,xs)) \rightarrow \operatorname{cons}(y,\operatorname{cons}(x,xs))
[a,b] = [b,a] \\ \operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(s(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) \quad \operatorname{len}([0,1,0]) \\ \rightarrow_{\mathcal{R}_{len}} \quad \operatorname{s}(\operatorname{len}(\operatorname{cons}(s(\mathcal{O}),\operatorname{cons}(\mathcal{O},\operatorname{nil})))) \quad 1 + \operatorname{len}([1,0]) \\ \rightarrow_{\mathcal{B}_{com}} \quad \operatorname{s}(\operatorname{len}(\operatorname{cons}(\mathcal{O},\operatorname{cons}(s(\mathcal{O}),\operatorname{nil})))) \quad 1 + \operatorname{len}([0,1]) \\ \rightarrow_{\mathcal{R}_{len}} \quad \operatorname{s}(\operatorname{s}(\operatorname{len}(\operatorname{cons}(s(\mathcal{O}),\operatorname{nil})))) \quad 2 + \operatorname{len}([1])
```

```
len(nil) \rightarrow \mathcal{O}
\mathcal{R}_{len}:
                                   len(cons(x, y)) \rightarrow s(len(y))
\mathcal{B}_{com}:
                           cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))
                                                  [a, b] = [b, a]
                        len(cons(\mathcal{O}, cons(s(\mathcal{O}), cons(\mathcal{O}, nil)))) len([0, 1, 0])
                       s(len(cons(s(\mathcal{O}), cons(\mathcal{O}, nil))))
                                                                                    1 + \text{len}([1, 0])
                       s(len(cons(\mathcal{O}, cons(s(\mathcal{O}), nil))))
                                                                                           1 + \text{len}([0, 1])
                       s(s(len(cons(s(\mathcal{O}), nil))))
                                                                                           2 + len([1])

ightarrow_{\mathcal{R}_{\mathit{len}}}
                       s(s(s(len(nil))))
         \rightarrow_{\mathcal{R}_{\mathit{len}}}
                                                                                           3 + \text{len}([])
```

Introduction (TRS)

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Relative Termination of TRSs

 $s(s(s(\mathcal{O})))$

```
\mathcal{R}_{len}:
                                                    len(nil) \rightarrow \mathcal{O}
                                      len(cons(x, y)) \rightarrow s(len(y))
\mathcal{B}_{com}:
                             cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))
                                                      [a, b] = [b, a]
                         len(cons(\mathcal{O}, cons(s(\mathcal{O}), cons(\mathcal{O}, nil)))) len([0, 1, 0])
                         s(len(cons(s(\mathcal{O}), cons(\mathcal{O}, nil))))
                                                                                          1 + \text{len}([1, 0])
                         s(len(cons(\mathcal{O}, cons(s(\mathcal{O}), nil))))
                                                                                                 1 + len([0,1])

ightarrow_{\mathcal{B}_{com}}
                         s(s(len(cons(s(\mathcal{O}), nil))))
                                                                                                 2 + len([1])

ightarrow_{\mathcal{R}_{\mathit{len}}}
          \to_{\mathcal{R}_{len}} \quad \mathsf{s}(\mathsf{s}(\mathsf{s}(\mathsf{len}(\mathsf{nil}))))
                                                                                                 3 + \text{len}([])
```

Introduction (TRS)

```
\mathcal{R}_{len}:
                                                        \mathsf{len}(\mathsf{nil}) \ 	o \ \mathcal{O}
                                         len(cons(x, y)) \rightarrow s(len(y))
```

 \mathcal{B}_{com} : $cons(x, cons(y, xs)) \rightarrow cons(y, cons(x, xs))$

```
len(cons(\mathcal{O}, cons(s(\mathcal{O}), cons(\mathcal{O}, nil)))) len([0, 1, 0])
               s(len(cons(s(\mathcal{O}), cons(\mathcal{O}, nil))))
                                                                                       1 + \text{len}([1, 0])
               s(len(cons(\mathcal{O}, cons(s(\mathcal{O}), nil))))
                                                                                          1 + \text{len}([0, 1])
\rightarrow_{\mathcal{B}_{com}}
              s(s(len(cons(s(\mathcal{O}), nil))))
                                                                                          2 + len([1])

ightarrow_{\mathcal{R}_{len}}
\rightarrow_{\mathcal{R}_{len}} s(s(s(len(nil))))
                                                                                          3 + \text{len}([])
              s(s(s(\mathcal{O})))
                                                                                          3
\rightarrow_{\mathcal{R}_{len}}
```

[a, b] = [b, a]

Relative Termination

 \mathcal{R}/\mathcal{B} is terminating iff there is no infinite evaluation

$$t_0 o_{\mathcal{R}} \circ o_{\mathcal{B}}^* \ t_1 o_{\mathcal{R}} \circ o_{\mathcal{B}}^* \dots$$

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Defined Symbols: len

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Defined Symbols: len , Constructor Symbols: cons, nil, s, $\mathcal O$

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$\operatorname{Sub}_D(r)$

 $\operatorname{Sub}_D(r) := \{t \mid t \text{ is a subterm of } r \text{ with defined root symbol}\}$

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Dependency Pairs

If $f(\ell_1, \ldots, \ell_n) \to r$ is a rule and $g(r_1, \ldots, r_m) \in \operatorname{Sub}_D(r)$, then $f^{\#}(\ell_1,\ldots,\ell_n) \to g^{\#}(r_1,\ldots,r_m)$ is a dependency pair

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```
\mathcal{DP}(\mathcal{R}_{len}):
```

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: $|\mathsf{len}(\mathsf{nil})| o \mathcal{O} \ |\mathsf{len}(\mathsf{cons}(x,xs))| o \mathsf{s}(|\mathsf{len}(xs))$

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$$\mathcal{DP}(\mathcal{R}_{\mathit{len}})$$
: $\mathsf{len}^\#(\mathsf{cons}(x,xs)) o \mathsf{len}^\#(xs)$

Dependency Pairs Cont.

$$\begin{array}{c} \mathsf{len}(\mathsf{nil}) \to \mathcal{O} \\ \mathsf{len}(\mathsf{cons}(x,xs)) \to \mathsf{s}(\mathsf{len}(xs)) \end{array}$$

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Termination of $(\mathcal{D}, \mathcal{R})$

 $(\mathcal{D},\mathcal{R})$ is terminating iff there is no infinite evaluation

$$t_0 \rightarrow_{\mathcal{D}} \circ \rightarrow_{\mathcal{R}}^* t_1 \rightarrow_{\mathcal{D}} \circ \rightarrow_{\mathcal{R}}^* \dots$$

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Reminder: Relative Termination of \mathcal{R}/\mathcal{B}

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Theorem: Chain Criterion [Arts & Giesl 2000]

 \mathcal{R} is terminating iff $\mathcal{DP}(\mathcal{R})/\mathcal{R}$ is terminating

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 - Transform a "big" problem into simpler sub-problems

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 - if all $(\mathcal{D}_i, \mathcal{R}_i)$ are terminating, • Proc is sound: then $(\mathcal{D}, \mathcal{R})$ is terminating
 - if $(\mathcal{D}, \mathcal{R})$ is terminating, • *Proc* is complete: then all $(\mathcal{D}_i, \mathcal{R}_i)$ are terminating

Timeline



- 2000: DPs for termination [Arts & Giesl 2000, ...]
- 2006: Problem #106 of the RTA list of open problems:
 - "Can we use the dependency pair method to prove relative termination?"
- 2016: Properties of \mathcal{R}/\mathcal{B} that allow to analyze the DP problem $(\mathcal{DP}(\mathcal{R}), \mathcal{R} \cup \mathcal{B})$ [Iborra & Nishida & Vidal & Yamada 2016]
- 2023: Annotated Dependency Pairs for <u>Probabilistic</u> Rewriting [Kassing & Giesl 2023, Kassing & Dollase & Giesl 2024]
- 2024: Annotated Dependency Pairs for Relative Termination



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 $\mbox{Goal: DP approach better than } \mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B} \mbox{ (Termination of } \mathcal{R} \cup \mathcal{B})$

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Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$?

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Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

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Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$?

 \mathcal{R}_1 : $\mathbf{a} \to \mathbf{b}$

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Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 : $a \rightarrow b$ \mathcal{B}_1 :

 $\mathsf{b} \to \mathsf{a}$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

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Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

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 $\mathsf{b}\to\mathsf{a}$

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

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 $\mathsf{b}\to\mathsf{a}$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

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 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{\mathsf{f}} \to_{\mathcal{B}_2} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}) \to_{\mathcal{R}_2} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}}) \to_{\mathcal{B}_2} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f})) \to_{\mathcal{R}_2} \dots$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $a \rightarrow b$

 \mathcal{B}_1 :

 $b \rightarrow a$

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $f \rightarrow d(a, f)$

 $f \rightarrow_{\mathcal{B}_2} d(a, f) \rightarrow_{\mathcal{R}_2} d(b, f) \rightarrow_{\mathcal{B}_2} d(b, d(a, f)) \rightarrow_{\mathcal{R}_2} \dots$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $a \rightarrow b$

 \mathcal{B}_1 :

 $b \rightarrow a$

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $f \rightarrow d(a, f)$

 $f \rightarrow_{\mathcal{B}_2} d(a, f) \rightarrow_{\mathcal{R}_2} d(b, f) \rightarrow_{\mathcal{B}_2} d(b, d(a, f)) \rightarrow_{\mathcal{R}_2} \dots$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

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 \mathcal{R}_1 :

 $a \rightarrow b$

 \mathcal{B}_1 :

 $b \rightarrow a$

 $a \rightarrow_{\mathcal{R}_1} b \rightarrow_{\mathcal{B}_1} a \rightarrow_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $f \rightarrow d(a, f)$

 $f \rightarrow_{\mathcal{B}_2} d(a, f) \rightarrow_{\mathcal{R}_2} d(b, f) \rightarrow_{\mathcal{B}_2} d(b, d(a, f)) \rightarrow_{\mathcal{R}_2} \dots$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

$$\underline{f} \to_{\mathcal{B}_2} d(\underline{a}, f) \to_{\mathcal{R}_2} d(b, \underline{f}) \to_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_2/\mathcal{B}_2$ not terminating

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

 \mathcal{R}_1 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1\cup\mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1)=\varnothing)$

 \mathcal{R}_2 :

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 \mathcal{B}_2 :

 $f \to d(a,f)$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $\mathcal{R}_2/\mathcal{B}_2$ not terminating, but $\mathcal{DP}(\mathcal{R}_2)/\mathcal{R}_2 \cup \mathcal{B}_2$ terminating $(\mathcal{DP}(\mathcal{R}_2) = \varnothing)$

Goal: DP approach better than $\mathcal{DP}(\mathcal{R} \cup \mathcal{B})/\mathcal{R} \cup \mathcal{B}$ (Termination of $\mathcal{R} \cup \mathcal{B}$)

Sufficient to analyze $\mathcal{DP}(\mathcal{R})/\mathcal{R} \cup \mathcal{B}$?

$$\mathcal{R}_1$$
:

 $\mathsf{a} \to \mathsf{b}$

 \mathcal{B}_1 :

 $\mathsf{b}\to\mathsf{a}$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{R}_1/\mathcal{B}_1$ not terminating, but $\mathcal{DP}(\mathcal{R}_1)/\mathcal{R}_1 \cup \mathcal{B}_1$ terminating $(\mathcal{DP}(\mathcal{R}_1) = \varnothing)$

$$\mathcal{R}_2$$
:

 $\mathsf{a} o \mathsf{b}$

 \mathcal{B}_2 :

 $\mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_2/\mathcal{B}_2$ not terminating, but $\mathcal{DP}(\mathcal{R}_2)/\mathcal{R}_2 \cup \mathcal{B}_2$ terminating $(\mathcal{DP}(\mathcal{R}_2) = \varnothing)$

Domination

 ${\mathcal R}$ dominates ${\mathcal B} :\Leftrightarrow$ no defined symbol of ${\mathcal R}$ in a right-hand side of ${\mathcal B}$

 \mathcal{R}_3 : $\mathsf{a} \to \mathsf{b}$ \mathcal{B}_3 : $\mathsf{f}(x) \to \mathsf{c}(x,\mathsf{f}(x))$

 \mathcal{R}_3 : $\mathsf{a} \to \mathsf{b}$

$$\mathcal{B}_3$$
: $f(x) \to c(x, f(x))$

$$\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$$

 \mathcal{R}_3 : $\mathsf{a} \to \mathsf{b}$

 \mathcal{B}_3 : $f(x) \to c(x, f(x))$

 $\underline{\mathsf{f(a)}} \to_{\mathcal{B}_3} \underline{\mathsf{d(\underline{a},f(a))}} \to_{\mathcal{R}_3} \underline{\mathsf{d(b,\underline{f(a)})}} \to_{\mathcal{B}_3} \underline{\mathsf{d(b,d(\underline{a},f(a)))}} \to_{\mathcal{R}_2} \dots$

 \mathcal{R}_3 : $a \to b$

 \mathcal{B}_3 : $f(x) \to c(x, f(x))$

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(\underline{b}, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$

 \mathcal{R}_3 : $\mathbf{a} \to \mathbf{b}$ \mathcal{B}_3 : $\mathbf{f}(\mathbf{x}) \to \mathbf{c}(\mathbf{x}, \mathbf{f}(\mathbf{x}))$

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, \underline{d(\underline{a}, f(a))}) \to_{\mathcal{R}_2} \dots$

 \mathcal{R}_3 : $\mathbf{a} \to \mathbf{b}$ \mathcal{B}_3 : $\mathbf{f}(x) \to \mathbf{c}(x, \mathbf{f}(x))$

 $\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots$

 \mathcal{R}_3 : $\mathsf{a} o \mathsf{b}$ \mathcal{B}_3 : $f(x) \to c(x, f(x))$

$$\frac{\underline{\mathsf{f}(\mathsf{a})} \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}}, \mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b}, \underline{\mathsf{f}(\mathsf{a})}) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b}, \mathsf{d}(\underline{\mathsf{a}}, \mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3} \underbrace{\mathsf{not terminating}}$$

$$\mathcal{R}_3$$
: $\mathsf{a} \to \mathsf{b}$ \mathcal{B}_3 : $\mathsf{f}(x) \to \mathsf{c}(x,\mathsf{f}(x))$

$$\frac{\mathsf{f}(\mathsf{a}) \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}}(\underline{\mathsf{a}})) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

 \mathcal{R}_3 : $\mathsf{a} o \mathsf{b}$

$$\mathcal{B}_3$$
: $f(x) \to c(x, f(x))$

$$\frac{f(a) \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

Duplication

 \mathcal{B} is duplicating $:\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}: x$ occurs more often in r than in ℓ .

 \mathcal{R}_3 : $\mathsf{a} o \mathsf{b}$

$$\mathcal{B}_3$$
: $f(x) \to c(x, f(x))$

$$\frac{f(\underline{a})}{} \to_{\mathcal{B}_3} d(\underline{a}, f(\underline{a})) \to_{\mathcal{R}_3} d(\underline{b}, \underline{f(\underline{a})}) \to_{\mathcal{B}_3} d(\underline{b}, d(\underline{a}, f(\underline{a}))) \to_{\mathcal{R}_2} \dots$$

 $\mathcal{R}_3/\mathcal{B}_3$ not terminating, but $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3\cup\mathcal{B}_3$ terminating $(\mathcal{DP}(\mathcal{R}_3)=\varnothing)$

Duplication

 \mathcal{B} is duplicating $:\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}: x$ occurs more often in r than in ℓ .

$$\mathcal{R}_3$$
:

$$\mathsf{a}\to\mathsf{b}$$

$$\mathcal{B}_3$$
:

$$\mathcal{B}_3$$
: $f(x) \to c(x, f(x))$

$$\underline{f(a)} \rightarrow_{\mathcal{B}_3} d(\underline{a}, f(a)) \rightarrow_{\mathcal{R}_3} d(b, \underline{f(a)}) \rightarrow_{\mathcal{B}_3} d(b, \underline{d(\underline{a}, f(a))}) \rightarrow_{\mathcal{R}_2} \dots$$
 $\mathcal{R}_2/\mathcal{B}_2$ not terminating, but $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3$ terminating $(\mathcal{DP}(\mathcal{R}_3) = \mathcal{B}_3)$

$$\mathsf{nating}\;(\mathcal{DP}(\mathcal{R}_3)=0)$$

$\mathcal{R}_3/\mathcal{B}_3$ not terminating, but $\mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3$ terminating $(\mathcal{DP}(\mathcal{R}_3) = \varnothing)$

Duplication

 \mathcal{B} is duplicating : $\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V}$: x occurs more often in r than in ℓ .

DPs for Relative Termination [Iborra et al. 2016]

If \mathcal{R} dominates \mathcal{B} and \mathcal{B} is non-duplicating, then \mathcal{R}/\mathcal{B} is terminating iff $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$ is terminating

$$\mathcal{R}_3$$
: $a \to b$ \mathcal{B}_3 : $f(x) \to c(x, f(x))$

$$\frac{\underline{f(a)} \to_{\mathcal{B}_3} d(\underline{a}, f(a)) \to_{\mathcal{R}_3} d(b, \underline{f(a)}) \to_{\mathcal{B}_3} d(b, d(\underline{a}, f(a))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

Duplication

 $\mathcal{B} \text{ is duplicating } :\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V} \text{: } x \text{ occurs more often in } r \text{ than in } \ell.$

DPs for Relative Termination [Iborra et al. 2016]

If $\mathcal R$ dominates $\mathcal B$ and $\mathcal B$ is non-duplicating, then $\mathcal R/\mathcal B$ is terminating iff $\mathcal D\mathcal P(\mathcal R)/\mathcal R\cup\mathcal B$ is terminating

$$\mathcal{R}_{\mathit{len}}$$
: $\mathcal{B}_{\mathit{com}}$: $\mathsf{len}(\mathsf{nil}) \to \mathcal{O}$ $\mathsf{len}(\mathsf{cons}(x,xs)) \to \mathsf{s}(\mathsf{len}(xs))$ $\mathsf{cons}(x,\mathsf{cons}(y,xs)) \to \mathsf{cons}(y,\mathsf{cons}(x,xs))$

$$\mathcal{R}_3$$
: $a \to b$ \mathcal{B}_3 : $f(x) \to c(x, f(x))$

$$\frac{\mathsf{f}(\mathsf{a}) \to_{\mathcal{B}_3} \mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a})) \to_{\mathcal{R}_3} \mathsf{d}(\mathsf{b},\underline{\mathsf{f}}(\underline{\mathsf{a}})) \to_{\mathcal{B}_3} \mathsf{d}(\mathsf{b},\mathsf{d}(\underline{\mathsf{a}},\mathsf{f}(\mathsf{a}))) \to_{\mathcal{R}_2} \dots}{\mathcal{R}_3/\mathcal{B}_3 \text{ not terminating, but } \mathcal{DP}(\mathcal{R}_3)/\mathcal{R}_3 \cup \mathcal{B}_3 \text{ terminating } (\mathcal{DP}(\mathcal{R}_3) = \varnothing)$$

Duplication

 $\mathcal{B} \text{ is duplicating } :\Leftrightarrow \exists \ell \to r \in \mathcal{B}, x \in \mathcal{V} \text{: } x \text{ occurs more often in } r \text{ than in } \ell.$

DPs for Relative Termination [Iborra et al. 2016]

If \mathcal{R} dominates \mathcal{B} and \mathcal{B} is non-duplicating, then \mathcal{R}/\mathcal{B} is terminating iff $\mathcal{DP}(\mathcal{R})/\mathcal{R}\cup\mathcal{B}$ is terminating

$$\mathcal{R}_{len}$$
: \mathcal{B}_{com} : $len(nil) \to \mathcal{O}$ $len(cons(x,xs)) \to s(len(xs))$ $cons(x,cons(y,xs)) \to cons(y,cons(x,xs))$

 $\mathcal{R}_{\mathit{len}}/\mathcal{B}_{\mathit{com}}$ terminates $\Leftrightarrow \mathcal{DP}(\mathcal{R}_{\mathit{len}})/\mathcal{R}_{\mathit{len}} \cup \mathcal{B}_{\mathit{com}}$ terminates

Annotated Dependency Pairs

$$\mathcal{R}_2$$
: $a \to b$ \mathcal{B}_2 : $b \to a$

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 \mathcal{R}_2 : $\mathsf{a}\to\mathsf{b}$ \mathcal{B}_2 :

 $b \rightarrow a$

 $\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^\# \rightarrow a^\#$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathsf{a}} \to_{\mathcal{R}_1} \underline{\mathsf{b}} \to_{\mathcal{B}_1} \underline{\mathsf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^\# \rightarrow a^\#$

 $a^{\#}$

 \mathcal{R}_2 :

 $\mathsf{a}\to\mathsf{b}$

 \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^\# \rightarrow a^\#$

$$\mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}^\#$$

 \mathcal{R}_2 : $\mathsf{a}\to\mathsf{b}$ \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^{\#} \rightarrow a^{\#}$

$$\mathsf{a}^\# o_{\mathcal{A}(\mathcal{R}_1)}^{(\#)} \mathsf{b}^\# o_{\mathcal{A}(\mathcal{B}_1)}^{(\#)} \mathsf{a}^\#$$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^{\#} \rightarrow a^{\#}$

$$\mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} \mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^{\#} \rightarrow a^{\#}$

$$\mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} \mathsf{a}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

$$\mathsf{a} \to_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}$$

 \mathcal{R}_2 :

 $a \rightarrow b$

 \mathcal{B}_2 :

 $b \rightarrow a$

$$\underline{\mathbf{a}} \to_{\mathcal{R}_1} \underline{\mathbf{b}} \to_{\mathcal{B}_1} \underline{\mathbf{a}} \to_{\mathcal{R}_1} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$:

 $a^\# \rightarrow b^\#$

 $\mathcal{A}(\mathcal{B}_2)$: $b^{\#} \rightarrow a^{\#}$

$$a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} b^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_1)} a^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{R}_1)} \dots$$

$$\mathsf{a} o_{\mathcal{A}(\mathcal{R}_1)} \mathsf{b}$$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

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 \mathcal{R}_2 : a o b \mathcal{B}_2 : f o d(a,f)

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 $\mathcal{R}_2\hbox{:}\qquad \quad \mathsf{a}\to\mathsf{b}\qquad \qquad \mathcal{B}_2\hbox{:}\qquad \quad \mathsf{f}\to\mathsf{d}(\mathsf{a},\mathsf{f})$

$$\underline{f} \to_{\mathcal{B}_2} d(\underline{a},f) \to_{\mathcal{R}_2} d(b,\underline{f}) \to_{\mathcal{B}_2} d(b,d(\underline{a},f)) \to_{\mathcal{R}_2} \dots$$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 \mathcal{R}_2 : \mathcal{B}_2 : $a \rightarrow b$ $f \rightarrow d(a, f)$

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ $\mathcal{A}(\mathcal{B}_2)$: $\left\{ egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$

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 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 \mathcal{R}_2 : $\mathsf{a} o \mathsf{b}$ \mathcal{B}_2 : $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $f^{\#}$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

 \mathcal{R}_2 : $\mathsf{a} \to \mathsf{b}$ \mathcal{B}_2 : $\mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(a^{\#}, f^{\#})$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 \mathcal{R}_2 : \mathcal{B}_2 : $a \rightarrow b$ $f \rightarrow d(a, f)$

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$

 $\mathcal{A}(\mathcal{B}_2)$: $\left\{ egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$ $\mathcal{A}(\mathcal{R}_2)$: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$

 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a^{\#}, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, f^{\#})$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 \mathcal{R}_2 : \mathcal{B}_2 : $a \rightarrow b$ $f \rightarrow d(a, f)$

 $f \rightarrow_{\mathcal{B}_2} d(a,f) \rightarrow_{\mathcal{R}_2} d(b,f) \rightarrow_{\mathcal{B}_2} d(b,d(a,f)) \rightarrow_{\mathcal{R}_2} \dots$

 $\mathcal{A}(\mathcal{R}_2)$: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$ $\mathcal{A}(\mathcal{B}_2)$: $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$

 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a^{\#}, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, f^{\#}) \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(b, d(a^{\#}, f^{\#}))$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 \mathcal{R}_2 : $a \rightarrow b$ \mathcal{B}_2 :

 $f \rightarrow d(a, f)$

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

 $\mathcal{A}(\mathcal{R}_2)$: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$

 $\mathcal{A}(\mathcal{B}_2)$: $\left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array} \right\}$

 $f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(a^{\#}, f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(b, f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(b, d(a^{\#}, f^{\#})) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \dots$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

 \mathcal{R}_2 : $a \rightarrow b$ \mathcal{B}_2 :

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 $\mathcal{A}(\mathcal{R}_2)$: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$

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 $f^{\#}$

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 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 \mathcal{R}_2 : $a \to b$ \mathcal{B}_2 : $f \to d(a,f)$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $\mathsf{f}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^\#,\mathsf{f})$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* t_2 \rightarrow_{\mathcal{P}}^{(\#)} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \dots$$

 \mathcal{R}_2 : $a \rightarrow b$ \mathcal{B}_2 :

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$$\mathcal{A}(\mathcal{R}_2)$$
: $\left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\}$

$$\mathcal{A}(\mathcal{B}_2)$$
: $\left\{egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^-,\mathsf{f}^-) \ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array}
ight\}$

$$\mathsf{f}^{\#} \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a}^{\#},\mathsf{f}) \to^{(\#)}_{\mathcal{A}(\mathcal{R}_2)} \mathsf{d}(\mathsf{b},\mathsf{f})$$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

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f#

Relative $(\mathcal{P},\mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

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 \mathcal{R}_2 : $\mathsf{a} o \mathsf{b}$ \mathcal{B}_2 : $\mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $\mathsf{f}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{d}(\mathsf{a},\mathsf{f}^\#)$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

$$t_1 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ t_2 \rightarrow^{(\#)}_{\mathcal{P}} \circ (\rightarrow_{\mathcal{P}} \cup \rightarrow_{\mathcal{S}})^* \ \dots$$

 \mathcal{R}_2 : \mathcal{B}_2 : $a \rightarrow b$ $f \rightarrow d(a, f)$

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 $f^{\#} \rightarrow^{(\#)}_{A(\mathcal{B}_2)} d(a, f^{\#}) \rightarrow_{\mathcal{A}(\mathcal{R}_2)} d(b, f^{\#})$

Relative $(\mathcal{P},\mathcal{S})$ -Chain

 $(\mathcal{P},\mathcal{S})$ is terminating iff there is no infinite evaluation

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 \mathcal{R}_2 : $\mathsf{a} \to \mathsf{b}$ \mathcal{B}_2 : $\mathsf{f} \to \mathsf{d}(\mathsf{a},\mathsf{f})$

 $\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$

 $f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} d(a,f^{\#}) \rightarrow_{\mathcal{A}(\mathcal{R}_2)} d(b,f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \dots$

Relative $(\mathcal{P}, \mathcal{S})$ -Chain

 $(\mathcal{P}, \mathcal{S})$ is terminating iff there is no infinite evaluation

 \mathcal{R}_2 : $a \rightarrow b$ \mathcal{B}_2 :

 $f \rightarrow d(a, f)$

$$\underline{f} \rightarrow_{\mathcal{B}_2} d(\underline{a}, f) \rightarrow_{\mathcal{R}_2} d(b, \underline{f}) \rightarrow_{\mathcal{B}_2} d(b, d(\underline{a}, f)) \rightarrow_{\mathcal{R}_2} \dots$$

$$\mathcal{A}(\mathcal{R}_2)$$
: $\left\{ egin{array}{l} \mathsf{a}^\# o \mathsf{b} \ \mathsf{a} o \mathsf{b} \end{array} \right\}$

$$\mathcal{A}(\mathcal{B}_2)$$
: $\left\{egin{array}{l} \mathsf{f}^\# o \mathsf{d}(\mathsf{a}^\#,\mathsf{f}^\#) \ \mathsf{f} o \mathsf{d}(\mathsf{a},\mathsf{f}) \end{array}
ight\}$

$$f^{\#} \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_{2})} d(a,f^{\#}) \rightarrow_{\mathcal{A}(\mathcal{R}_{2})} d(b,f^{\#}) \rightarrow^{(\#)}_{\mathcal{A}(\mathcal{B}_{2})} \dots$$

Chain Criterion

For \mathcal{B} non-duplicating: \mathcal{R}/\mathcal{B} is terminating iff $(\mathcal{A}_1(\mathcal{R}), \mathcal{A}_2(\mathcal{B}))$ is terminating

- Our objects we work with:
 - \bullet (Relative) ADP Problems $(\mathcal{P},\mathcal{S})$ with \mathcal{P} and \mathcal{S} set of ADPs

- Our objects we work with:
 - ullet (Relative) ADP Problems $(\mathcal{P},\mathcal{S})$ with \mathcal{P} and \mathcal{S} set of ADPs
- How do we start?:
 - (Chain Criterion) Use all main ADPs with 1-annotation $\mathcal{A}_1(\mathcal{R})$ and all base ADPs with 2-annotations $\mathcal{A}_2(\mathcal{B})$

- Our objects we work with:
 - (Relative) ADP Problems $(\mathcal{P}, \mathcal{S})$ with \mathcal{P} and \mathcal{S} set of ADPs
- How do we start?:
 - (Chain Criterion) Use all main ADPs with 1-annotation $A_1(\mathcal{R})$ and all base ADPs with 2-annotations $A_2(\mathcal{B})$
- How do we create smaller problems?:
 - DP Processors: $Proc(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P}_1, \mathcal{S}_1), \dots, (\mathcal{P}_k, \mathcal{S}_k)\}$

- Our objects we work with:
 - \bullet (Relative) ADP Problems $(\mathcal{P},\mathcal{S})$ with \mathcal{P} and \mathcal{S} set of ADPs
- How do we start?:
 - (Chain Criterion) Use all main ADPs with 1-annotation $\mathcal{A}_1(\mathcal{R})$ and all base ADPs with 2-annotations $\mathcal{A}_2(\mathcal{B})$
- How do we create smaller problems?:
 - DP Processors: $Proc(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P}_1, \mathcal{S}_1), \dots, (\mathcal{P}_k, \mathcal{S}_k)\}$
 - Proc is sound: if all $(\mathcal{P}_i, \mathcal{S}_i)$ are terminating, then $(\mathcal{P}, \mathcal{S})$ is terminating

- Our objects we work with:
 - (Relative) ADP Problems $(\mathcal{P}, \mathcal{S})$ with \mathcal{P} and \mathcal{S} set of ADPs
- How do we start?:
 - (Chain Criterion) Use all main ADPs with 1-annotation $A_1(\mathcal{R})$ and all base ADPs with 2-annotations $A_2(\mathcal{B})$
- How do we create smaller problems?:
 - DP Processors: $Proc(\mathcal{P}, \mathcal{S}) = \{(\mathcal{P}_1, \mathcal{S}_1), \dots, (\mathcal{P}_k, \mathcal{S}_k)\}$
 - if all $(\mathcal{P}_i, \mathcal{S}_i)$ are terminating, • Proc is sound: then $(\mathcal{P}, \mathcal{S})$ is terminating
 - if $(\mathcal{P}, \mathcal{S})$ is terminating, • *Proc* is complete: then all $(\mathcal{P}_i, \mathcal{S}_i)$ are terminating

24/[4, 3]

$$24/[4,3] = (24/4)/3$$

$$24/[4,3] = (24/4)/3 = 2$$

$$24/[4,3] = (24/4)/3 = 2$$

```
\mathcal{R}_{divL}:
                                   minus (x, \mathcal{O}) \to x
                  (a)
                  (b) minus (s(x), s(y)) \rightarrow minus (x, y)
                                   div (\mathcal{O}, s(y)) \to \mathcal{O}
                        \mathsf{div} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\mathsf{div}(\mathsf{minus} \ (x,y),\mathsf{s}(y)))
                  (d)
                            divL (x, nil) \rightarrow x
                  (e)
                  (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4$$

```
\mathcal{R}_{divL}:
                              minus (x, \mathcal{O}) \to x
                (a)
               (b) minus (s(x), s(y)) \rightarrow minus (x, y)
               (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
               (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
               (e) \operatorname{divL}(x, \operatorname{nil}) \to x
               (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{divL}:
                               minus (x, \mathcal{O}) \to x
                (b) minus (s(x), s(y)) \rightarrow minus (x, y)
                (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
                (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                (e) \operatorname{divL}(x, \operatorname{nil}) \to x
                (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                         \begin{array}{lll} (a) & \text{minus } (x,\mathcal{O}) \rightarrow x \\ (b) & \text{minus } (\mathsf{s}(x),\mathsf{s}(y)) \rightarrow \mathsf{minus } \ (x,y) \\ (c) & \text{div } (\mathcal{O},\mathsf{s}(y)) \rightarrow \mathcal{O} \end{array} 
                         (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                         (e) \operatorname{divL}(x, \operatorname{nil}) \to x
                         (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                        (g)
                                    \operatorname{divL} (x, \operatorname{cons}(y, xs)) \to \operatorname{divL} (x, \operatorname{switch} (y, xs))
                        (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                                       switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                       \begin{array}{lll} (a) & \text{minus } (x,\mathcal{O}) \rightarrow x \\ (b) & \text{minus } (\mathsf{s}(x),\mathsf{s}(y)) \rightarrow \mathsf{minus } \ (x,y) \\ (c) & \text{div } (\mathcal{O},\mathsf{s}(y)) \rightarrow \mathcal{O} \end{array} 
                      (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                                      divL (x, nil) \rightarrow x
                      (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                      (g) divL (x, cons(y, xs)) \rightarrow divL (x, switch (y, xs))
                      (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                                  switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

$$\mathcal{R}_{\mathsf{divL}} : \\ (a) & \mathsf{minus} \ (x,\mathcal{O}) \to x \\ (b) & \mathsf{minus} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{minus} \ (x,y) \\ (c) & \mathsf{div} \ (\mathcal{O},\mathsf{s}(y)) \to \mathcal{O} \\ (d) & \mathsf{div} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\mathsf{div}(\mathsf{minus} \ (x,y),\mathsf{s}(y))) \\ (e) & \mathsf{divL} \ (x,\mathsf{nil}) \to x \\ (f) & \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL}(\mathsf{div} \ (x,y),xs) \\ \\ \mathcal{B}_{\mathsf{com}} : & & (\mathbf{g}) & \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL} \ (x,\mathsf{switch} \ (y,xs)) \\ (h) & \mathsf{switch} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{cons}(y,\mathsf{switch} \ (x,xs)) \\ (i) & & \mathsf{switch} \ (x,xs) \to \mathsf{cons}(x,xs) \\ \end{cases}$$

$$24/[4,3] \rightarrow_{\mathcal{B}_{com}} 24/[\hat{4},3]$$

Example: Division

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

$$\mathcal{R}_{\mathsf{divL}}:$$

$$(a) \qquad \mathsf{minus} \ (x,\mathcal{O}) \to x$$

$$(b) \qquad \mathsf{minus} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{minus} \ (x,y)$$

$$(c) \qquad \mathsf{div} \ (\mathcal{O},\mathsf{s}(y)) \to \mathcal{O}$$

$$(d) \qquad \mathsf{div} \ (\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\mathsf{div}(\mathsf{minus} \ (x,y),\mathsf{s}(y)))$$

$$(e) \qquad \mathsf{divL} \ (x,\mathsf{nil}) \to x$$

$$(f) \qquad \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL}(\mathsf{div} \ (x,y),xs)$$

$$\mathcal{B}_{\mathsf{com}}:$$

$$(g) \qquad \mathsf{divL} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{divL} \ (x,\mathsf{switch} \ (y,xs))$$

$$(h) \qquad \mathsf{switch} \ (x,\mathsf{cons}(y,xs)) \to \mathsf{cons}(y,\mathsf{switch} \ (x,xs))$$

$$(i) \qquad \mathsf{switch} \ (x,xs) \to \mathsf{cons}(x,xs)$$

Example: Division

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{R}_{\mathsf{divL}}:
                (a) minus (x, \mathcal{O}) \to x

(b) minus (s(x), s(y)) \to minus (x, y)
                (c) div (\mathcal{O}, s(y)) \to \mathcal{O}
                (d) div (s(x), s(y)) \rightarrow s(div(minus (x, y), s(y)))
                            divL (x, nil) \rightarrow x
                (f) divL (x, cons(y, xs)) \rightarrow divL(div (x, y), xs)
\mathcal{B}_{com}:
                (g) divL (x, cons(y, xs)) \rightarrow divL (x, switch (y, xs))
                (h) switch (x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                        switch (x, xs) \rightarrow cons(x, xs)
```

$$24/[4,3] \to_{\mathcal{B}_{com}} 24/[\hat{4},3] \to_{\mathcal{B}_{com}} 24/[3,\hat{4}] \to_{\mathcal{B}_{com}} 24/[3,4]$$

Example: Division

$$24/[4,3] = (24/4)/3 = 2 = (24/3)/4 = 24/[3,4]$$

```
\mathcal{A}_1(\mathcal{R}_{\text{divL}}):
                               (a)
                                                     minus^{\#}(x,\mathcal{O}) \to x
                               (b) minus\#(s(x), s(y)) \rightarrow minus\#(x, y)
                               (c)
                                                      \operatorname{div}^{\#}(\mathcal{O}, s(v)) \to \mathcal{O}
                              (d1) \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}(\operatorname{minus}^{\#}(x, y), s(y)))
                              (d2) \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
                               (e)
                                                        divL^{\#}(x, nil) \rightarrow x
                              (f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
                              (f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
\mathcal{A}_2(\mathcal{B}_{com}):
                              (g)
                                         \operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs)) \to \operatorname{divL}^{\#}(x,\operatorname{switch}^{\#}(y,xs))
                              (h) switch \#(x, cons(y, xs)) \rightarrow cons(y, switch \#(x, xs))
                               (i)
                                                     switch^{\#}(x,xs) \rightarrow cons(x,xs)
```

$$\mathcal{P}_2$$
: $\left\{ egin{aligned} \mathbf{a}^\# & \to \mathbf{b} \\ \mathbf{a} & \to \mathbf{b} \end{aligned}
ight\}$

$$\mathcal{S}_2$$
:
$$\left\{ \begin{array}{l} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{array} \right\}$$

$$\mathcal{P}_2 \colon \quad \left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\} \qquad \qquad \mathcal{S}_2 \colon \left\{ \begin{array}{l} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{array} \right\}$$

$$\mathit{Proc}_{\mathsf{DG}}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

$$\mathcal{P}_2: \qquad \left\{ \begin{array}{l} \mathsf{a}^\# \to \mathsf{b} \\ \mathsf{a} \to \mathsf{b} \end{array} \right\} \qquad \qquad \mathcal{S}_2: \quad \left\{ \begin{array}{l} \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) \\ \mathsf{f} \to \mathsf{d}(\mathsf{a}, \mathsf{f}) \end{array} \right\}$$

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$

$$\mathcal{P}_2$$
: $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$

$$S_2: \begin{cases} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{cases}$$

$$\mathit{Proc}_{\mathsf{DG}}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph:

$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$

$$\mathcal{P}_2$$
: $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$

$$S_2$$
: $\left\{ f^\# \to d(a^\#, f^\#) \atop f \to d(a, f) \right\}$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$(\mathcal{P}, \mathcal{S})$ -Dependency Graph

ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$

$$\mathcal{P}_2$$
: $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$

$$S_2$$
: $\left\{ f^\# \to d(a^\#, f^\#) \atop f \to d(a, f) \right\}$

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\mathcal{P}_2$$
: $\left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$

$$S_2$$
: $\left\{ f^\# \to d(a^\#, f^\#) \atop f \to d(a, f) \right\}$

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph:
$$\boxed{ \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) }$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^{\#}\sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^{\#}\sigma_2$ for substitutions σ_1, σ_2 .

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$Proc_{DG}(\mathcal{P},\mathcal{S}) = \{$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph:
$$\boxed{ \mathsf{f}^\# \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) }$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\begin{array}{ccc} \mathcal{P}_2 \colon & \left\{ \begin{smallmatrix} a^\# \to b \\ a \to b \end{smallmatrix} \right\} & & \mathcal{S}_2 \colon \left\{ \begin{smallmatrix} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{smallmatrix} \right\} \end{array}$$

$$\mathit{Proc}_{\mathit{DG}}(\mathcal{P},\mathcal{S}) = \{ & | \mathcal{Q} \in \mathtt{SCC} \}$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing)$

```
(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)-Dependency Graph:
a^{\#} \rightarrow b f^{\#} \rightarrow d(a^{\#}, f^{\#})
```

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\begin{array}{ccc} \mathcal{P}_2 \colon & \left\{ \begin{smallmatrix} a^\# \to b \\ a \to b \end{smallmatrix} \right\} & & \mathcal{S}_2 \colon \left\{ \begin{smallmatrix} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{smallmatrix} \right\} \end{array}$$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, (\mathcal{S} \cap \mathcal{Q}) \cup b((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC \}$$
(sound & complete)

$$Proc_{DG}(\mathcal{P}_2 \cup \mathcal{S}_2, \varnothing) = \{(\mathcal{S}_2, \mathcal{P}_2)\}$$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)\text{-Dependency Graph:}$$

$$\boxed{\mathbf{f}^\# \to \mathsf{d}(\mathbf{a}^\#, \mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{\flat(\mathcal{P} \cup \mathcal{S})}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\mathcal{P}_2$$
: $\left\{ egin{aligned} \mathsf{a}^\# & \to \mathsf{b} \\ \mathsf{a} & \to \mathsf{b} \end{aligned}
ight\}$ \mathcal{S}_2 : $\left\{ egin{aligned} \mathsf{f}^\# & \to \mathsf{d}(\mathsf{a}^\#, \mathsf{f}^\#) \\ \mathsf{f} & \to \mathsf{d}(\mathsf{a}, \mathsf{f}) \end{aligned} \right\}$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \forall ((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC \\ \text{(sound \& complete)}$$

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$

$$(\mathcal{A}(\mathcal{P}_2) \cup \mathcal{S}_2, \varnothing)$$
-Dependency Graph:
$$\boxed{\mathbf{a}^\# \to \mathbf{b}} \boxed{\mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#, \mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\begin{array}{ll} \mathcal{P}_2 \colon & \left\{ \begin{smallmatrix} a^\# \to b \\ a \to b \end{smallmatrix} \right\} & \mathcal{S}_2 \colon \left\{ \begin{smallmatrix} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{smallmatrix} \right\} \end{array}$$

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \flat ((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$Proc_{DG}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, (\mathcal{S} \cap \mathcal{Q}) \cup b((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in SCC \}$$
(sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$

$$(\mathcal{A}(\mathcal{P}_2),\mathcal{S}_2)\text{-Dependency Graph:}$$

$$\boxed{\mathbf{a}^\# \to \mathbf{b}} \qquad \qquad \boxed{\mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#,\mathbf{f}^\#)}$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\mathcal{P}_2: \qquad \left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$$

$$S_2: \begin{cases} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{cases}$$

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \forall ((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in \texttt{SCC} \cup \texttt{Lasso} \}$$
 (sound & complete)

 $Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2)$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

$$\mathcal{P}_2: \qquad \left\{ \begin{array}{l} a^\# \to b \\ a \to b \end{array} \right\}$$

$$S_2: \begin{cases} f^\# \to d(a^\#, f^\#) \\ f \to d(a, f) \end{cases}$$

$$\textit{Proc}_{\textit{DG}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \cap \mathcal{Q}, \ (\mathcal{S} \cap \mathcal{Q}) \cup \forall ((\mathcal{P} \cup \mathcal{S}) \setminus \mathcal{Q})) \mid \mathcal{Q} \in \texttt{SCC} \cup \texttt{Lasso} \}$$
 (sound & complete)

$$Proc_{DG}(\mathcal{P}_2, \mathcal{S}_2) = \{(\mathcal{P}_2, \mathcal{S}_2)\}$$

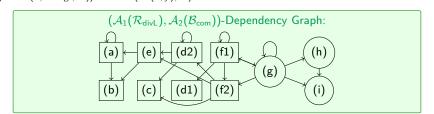
$$(\mathcal{A}(\mathcal{P}_2), \mathcal{S}_2)$$
-Dependency Graph:
$$\mathbf{f}^\# \to \mathbf{d}(\mathbf{a}^\#, \mathbf{f}^\#)$$

- ullet directed graph whose nodes are the ADPs from $\mathcal{P} \cup \mathcal{S}$
- there is an arc from $s \to t$ to $v \to w$ iff $t^\# \sigma_1 \to_{b(\mathcal{P}_1 \cup S)}^* v^\# \sigma_2$ for substitutions σ_1, σ_2 .

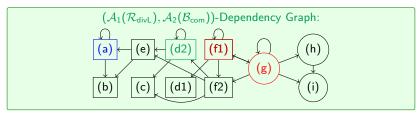
```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                          div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                        \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
(d1)
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                           (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                     \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                                      switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                           (i)
                            \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
 (e)
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```

```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
 (b)
                          \operatorname{div}^{\#}(\mathcal{O}, \mathsf{s}(v)) \to \mathcal{O}
 (c)
                                                                                                                                        \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
(d1)
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                           (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                                      switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                            (i)
 (e)
                            divL^{\#}(x, nil) \rightarrow x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
                                                              (\mathcal{A}_1(\mathcal{R}_{\text{divL}}), \mathcal{A}_2(\mathcal{B}_{\text{com}}))-Dependency Graph:
```

```
(a)
                         minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                          div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                      divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                     \operatorname{\mathsf{div}}^\#(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{\mathsf{div}}(\min \mathsf{us}^\#(x,y),\mathsf{s}(y)))
                                                                                                                         (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                     \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                                    switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                          (i)
 (e)
                            \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```

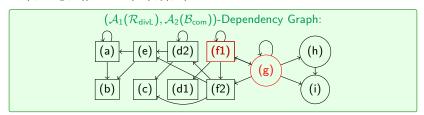


```
(a)
                        minus^{\#}(x,\mathcal{O}) \to x
 (b)
               minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                         div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                                    divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                        (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup><math>(x, xs))
                    \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}^{\#}(\operatorname{minus}(x,y),\mathsf{s}(y)))
(d2)
                                                                                                                                                  switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                        (i)
 (e)
                            \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



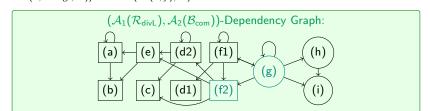
SCC: $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f1)\}$

```
(a)
                       minus^{\#}(x,\mathcal{O}) \to x
 (b)
              minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                        div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                   (g)
                                                                                                                               divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                   (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                    \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                             switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                    (i)
 (e)
                           \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



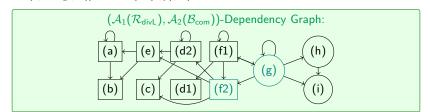
SCC: $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f1)\}$ Lasso: $\{(g), (f1)\}\$ and $\{(g), (f2)\}\$

```
(a)
                       minus^{\#}(x,\mathcal{O}) \to x
 (b)
              minus^{\#}(s(x), s(y)) \rightarrow minus^{\#}(x, y)
                        div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                    (g)
                                                                                                                               divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                    \operatorname{div}^{\#}(\mathsf{s}(x),\mathsf{s}(y)) \to \mathsf{s}(\operatorname{div}(\operatorname{minus}^{\#}(x,y),\mathsf{s}(y)))
                                                                                                                    (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch<sup>#</sup>(x, xs))
                    \operatorname{div}^{\#}(s(x), s(y)) \rightarrow s(\operatorname{div}^{\#}(\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                             switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                                    (i)
 (e)
                           \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}^{\#}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



SCC: $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f1)\}$ Lasso: $\{(g), (f1)\}\$ and $\{(g), (f2)\}\$

```
(a)
                     minus^{\#}(x,\mathcal{O}) \to x
 (b)
             minus^{\#}(s(x), s(y)) \rightarrow minus(x, y)
                      div^{\#}(\mathcal{O}, s(v)) \rightarrow \mathcal{O}
 (c)
                                                                                                                    divL^{\#}(x, cons(y, xs)) \rightarrow divL^{\#}(x, switch^{\#}(y, xs))
(d1)
                  \operatorname{div}^{\#}(s(x), s(y)) \to s(\operatorname{div}(\min s(x, y), s(y)))
                                                                                                         (h) switch<sup>#</sup>(x, cons(y, xs)) \rightarrow cons(y, switch (x, xs))
                  \operatorname{div}^{\#}(s(x), s(y)) \to s(\operatorname{div} (\operatorname{minus}(x, y), s(y)))
(d2)
                                                                                                                                switch^{\#}(x, xs) \rightarrow cons(x, xs)
                                                                                                         (i)
 (e)
                        \operatorname{divL}^{\#}(x,\operatorname{nil}) \to x
(f1) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}(\operatorname{div}(x, y), xs)
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs)
```



SCC: $\{(a)\}, \{(d2)\}, \text{ and } \{(g), (f1)\}$ Lasso: $\{(g), (f1)\}\$ and $\{(g), (f2)\}\$

Reduction Pair Processor (sound & complete)

 $(f2) \ \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad (g) \ \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$

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Find natural polynomial interpretation Pol

$$(f2) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(\mathsf{div}(x,y),x\mathsf{s}) \qquad (g) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,x\mathsf{s}))$$

```
\begin{array}{lll} \operatorname{divL}^{\#}_{Pol}(x,xs) & = & xs & \operatorname{switch}^{\#}_{Pol}(x,xs) & = & 0 \\ \operatorname{cons}_{Pol}(x,xs) & = & xs+1 & \operatorname{switch}_{Pol}(x,xs) & = & xs+1 \end{array}
```

$$(f2) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad (g) \; \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$$

such that $\flat(\mathcal{P} \cup \mathcal{S}) \subset >_{Pol}$ and

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

```
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```

such that $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \triangleright_{Pol}$ and

$$Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) \geq Pol(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs))) + Pol(\operatorname{switch}^{\#}(y, xs))$$

$$Pol(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs))) > Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs))$$

```
(f2) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) (g) \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

$$\begin{array}{lll} \operatorname{\sf divL}^\#_{Pol}(x,xs) &=& xs & \operatorname{\sf switch}^\#_{Pol}(x,xs) &=& 0 \\ \operatorname{\sf cons}_{Pol}(x,xs) &=& xs+1 & \operatorname{\sf switch}_{Pol}(x,xs) &=& xs+1 \\ & \dots & & & & & & & & & & & & & & \end{array}$$

such that $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \triangleright_{Pol}$ and

$$Pol(\operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs))) \geq Pol(\operatorname{divL}^{\#}(x,\operatorname{switch}(y,xs))) + Pol(\operatorname{switch}^{\#}(y,xs))$$

$$Pol(\operatorname{divL}^{\#}(x, \cos(y, xs))) > Pol(\operatorname{divL}^{\#}(\operatorname{div}(x, y), xs))$$

$$\label{eq:proc_RP} \begin{split} \textit{Proc}_\textit{RP}(\mathcal{P},\mathcal{S}) &= \{ (\mathcal{P} \setminus \mathcal{P}_{>}, (\mathcal{S} \setminus \mathcal{P}_{>}) \cup \flat(\mathcal{P}_{>})) \} \\ & \text{(sound \& complete)} \end{split}$$

$$Proc_{RP}(\{(f2)\},\ldots)$$

Reduction Pair Processor (sound & complete)

$$(f2) \ \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \rightarrow \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) \qquad (g) \ \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \rightarrow \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))$$

Find natural polynomial interpretation Pol

$$\operatorname{divL}_{Pol}^{\#}(x,xs) = xs \quad \operatorname{switch}_{Pol}^{\#}(x,xs) = 0$$

$$\operatorname{cons}_{Pol}(x,xs) = xs+1 \quad \operatorname{switch}_{Pol}(x,xs) = xs+1$$

$$\cdots$$

such that $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \geq_{\textit{Pol}}$ and

$$\frac{\operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)))}{xs + 1} \geq \frac{\operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs))) + \operatorname{\textit{Pol}}(\operatorname{switch}^{\#}(y, xs))}{xs + 1}$$

$$\frac{\text{Pol}(\text{divL}^{\#}(x, \text{cons}(y, xs)))}{xs + 1} > \frac{\text{Pol}(\text{divL}^{\#}(\text{div}(x, y), xs))}{xs}$$

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$$Proc_{RP}(\{(f2)\},\ldots)$$

Reduction Pair Processor (sound & complete)

$$(\mathit{f2})\;\mathsf{divL}^{\#}(x,\mathsf{cons}(y,\mathit{xs}))\to\mathsf{divL}^{\#}(\mathsf{div}(x,y),\mathit{xs}) \qquad (\mathit{g})\;\mathsf{divL}^{\#}(x,\mathsf{cons}(y,\mathit{xs}))\to\mathsf{divL}^{\#}(x,\mathsf{switch}^{\#}(y,\mathit{xs}))$$

Find natural polynomial interpretation Pol

such that $\flat(\mathcal{P} \cup \mathcal{S}) \subset >_{Pol}$ and

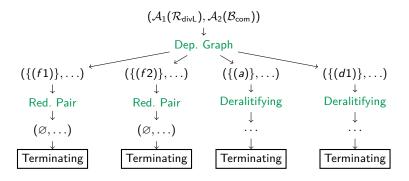
$$\frac{\operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)))}{xs + 1} \geq \frac{\operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs))) + \operatorname{\textit{Pol}}(\operatorname{switch}^{\#}(y, xs))}{xs + 1}$$

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$$Proc_{RP}(\{(f2)\},\ldots)=\{(\varnothing,\ldots)\}$$

Final Relative Termination Proof



⇒ Relative termination is proved automatically!

Fully implemented in AProVE

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Relative rewriting (130 benchmarks):

		new AProVE	NaTT	old AProVE	T_TT_2	MultumNonMulta
ľ	YES	91 (32)	68 (10)	48 (5)	39 (3)	0 (0)
	NO	13 (0)	5 (0)	13 (0)	7 (0)	13 (0)

Fully implemented in AProVE

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Relative string rewriting (403 benchmarks):

ſ		MultumNonMulta	Matchbox	AProVE	ADPs
Ī	YES	274	274	209	71

Fully implemented in AProVE

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	new AProVE	NaTT	old AProVE	T _T T ₂	MultumNonMulta
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Relative string rewriting (403 benchmarks):

		MultumNonMulta	Matchbox	AProVE	ADPs
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Equational rewriting (76 benchmarks):

		AProVE	MU-TERM	ADPs	
Ì	YES	66	64	36	l

• First DP framework specifically for relative termination

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- Annotated Dependency Pairs:

$$\left\{ \begin{array}{ll} \mathsf{divL}^\#(x,\mathsf{cons}(y,xs)) & \to \mathsf{divL}(\mathsf{div}^\#(x,y),xs) \\ \mathsf{divL}(x,\mathsf{cons}(y,xs)) & \to \mathsf{divL}(\mathsf{div}(x,y),xs) \end{array} \right\}$$

- First DP framework specifically for relative termination
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- Adapted the core processors from DP framework:
 - Dependency Graph Processor
 - Usable Terms Processor
 - Reduction Pair Processor
- o Deralitifying Processor

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- First DP framework specifically for relative termination
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- Adapted the core processors from DP framework:
 - o Dependency Graph Processor o Usable Terms Processor
 - Reduction Pair Processor
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- Fully implemented in AProVE.
- Future Work:
 - Further Processors to (dis)-prove relative termination
 - Analyze further possibilities to use ADPs



$$\mathcal{R}_2$$
: $\mathsf{a}(\mathsf{x}) o \mathsf{b}(\mathsf{x})$ \mathcal{B}_2 : $\mathsf{f} o \mathsf{a}(\mathsf{f})$

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$$\underline{f} \to_{\mathcal{B}_2} \underline{a(f)} \to_{\mathcal{R}_2} b(\underline{f}) \to_{\mathcal{B}_2} b(\underline{a(f)}) \to_{\mathcal{R}_2} \dots$$

f#

$$\mathcal{A}(\mathcal{R}_2)$$
: $\mathsf{a}(x) \to \mathsf{b}(x)$ $\mathcal{A}(\mathcal{B}_2)$: $\mathsf{f} \to \mathsf{a}^\#(\mathsf{f}^\#)$

$$\mathsf{f}^\# \to^{(\#)}_{\mathcal{A}(\mathcal{B}_2)} \mathsf{a}^\#(\mathsf{f}^\#)$$

$$\mathcal{R}_2$$
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$$a(x) \rightarrow b(x)$$

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 $a(x) \rightarrow b(x,x)$

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$$a(x) \rightarrow b(x)$$
 $a(x) \rightarrow b(x,x)$ $a(x) \rightarrow b(x,x)$

Chain Criterion

For $\mathcal B$ non-duplicating: $\mathcal R/\mathcal B$ is terminating iff $(\mathcal A_1(\mathcal R),\mathcal A_2(\mathcal B))$ is terminating

```
(f2)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \\ \qquad (g)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs)) \\
```

Find Com-monotonic and Com-invariant reduction pair (\succsim,\succ)

Reduction Pair

- \bullet \succsim is reflexive, transitive, and closed under contexts and substitutions,
- > is a well-founded order and closed under substitutions
- $\bullet \ \, \succsim \circ \, \succ \circ \, \succsim \subseteq \, \succ.$

$$(f2)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \\ \qquad (g)\; \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs)) \\$$

Find Com-monotonic and Com-invariant reduction pair (\succsim,\succ) such that

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- ullet \succ is a well-founded order and closed under substitutions
- $\bullet \ \succsim \circ \succ \circ \succsim \subseteq \succ.$

Com-monotonic

If $s_1 \succ s_2$, then $\text{Com}_2(s_1,t) \succ \text{Com}_2(s_2,t)$ and $\text{Com}_2(t,s_1) \succ \text{Com}_2(t,s_2)$

$$(f2) \ \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(\mathsf{div}(x,y),xs) \qquad \qquad (g) \ \mathsf{divL}(x,\mathsf{cons}(y,xs)) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,xs))$$

Find Com-monotonic and Com-invariant reduction pair (\succsim,\succ) such that

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- ≿∘≻∘≿ ⊆ ≻.

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Com-invariant

Let $\sim = \succsim \cap \precsim$, then

- $Com_2(s_1, s_2) \sim Com_2(s_2, s_1)$
- $Com_2(s_1, Com_2(s_2, s_3)) \sim Com_2(Com_2(s_1, s_2), s_3)$

```
 (f2) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(\mathsf{div}(x,y),x\mathsf{s}) \qquad \qquad (g) \ \mathsf{divL}(x,\mathsf{cons}(y,x\mathsf{s})) \to \mathsf{divL}^\#(x,\mathsf{switch}^\#(y,x\mathsf{s}))
```

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim \text{ and } \ell^{\#} \succsim \operatorname{ann}(r) \text{ for all } \ell \to r \in \mathcal{P} \cup \mathcal{S}$
- $\ell^{\#} \succ \operatorname{ann}(r)$ for all $\ell \to r \in \mathcal{P}_{\succ}$

```
(f2) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(\operatorname{div}(x, y), xs) \qquad \qquad (g) \operatorname{divL}(x, \operatorname{cons}(y, xs)) \to \operatorname{divL}^{\#}(x, \operatorname{switch}^{\#}(y, xs))
```

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```
\ell^{\#} \qquad \succeq \qquad \operatorname{ann}(r) \\ \operatorname{divL}^{\#}(x, \operatorname{cons}(y, xs)) \qquad \succeq \qquad \operatorname{Com}_{2}(\operatorname{divL}^{\#}(x, \operatorname{switch}(y, xs)), \operatorname{switch}^{\#}(y, xs))
```

```
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$$Proc_{RP}(\mathcal{P}, \mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \}$$
 (sound & complete)

```
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```

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$$\begin{aligned} \textit{Proc}_{\textit{RP}}(\mathcal{P}, \mathcal{S}) &= \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \} \\ & (\text{sound \& complete}) \end{aligned}$$
$$\textit{Proc}_{\textit{RP}}(\{(f2)\}, \ldots)$$

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$$Proc_{RP}(\{(f2)\},\ldots)$$

```
\operatorname{Com}_{2Pol}(x,y) = x+y \quad \operatorname{switch}_{Pol}^{\#}(x,xs) = 0

\operatorname{cons}_{Pol}(x,xs) = xs+1 \quad \operatorname{switch}_{Pol}(x,xs) = xs+1

\operatorname{divL}_{Pol}^{\#}(x,xs) = xs \quad \dots
```

Find Com-monotonic and Com-invariant reduction pair (\succeq, \succ) such that

- $\flat(\mathcal{P} \cup \mathcal{S}) \subseteq \succsim$ and $\ell^\# \succsim \operatorname{ann}(r)$ for all $\ell \to r \in \mathcal{P} \cup \mathcal{S}$
- $\ell^{\#} \succ \operatorname{ann}(r)$ for all $\ell \to r \in \mathcal{P}_{\succ}$

```
\begin{array}{cccc} \ell^{\#} & \succeq & \operatorname{ann}(r) \\ \operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs)) & \succeq & \operatorname{Com}_2(\operatorname{divL}^{\#}(x,\operatorname{switch}(y,xs)),\operatorname{switch}^{\#}(y,xs)) \\ \operatorname{\textit{Pol}}(\operatorname{divL}^{\#}(x,\operatorname{cons}(y,xs))) & \geq & \operatorname{\textit{Pol}}(\operatorname{Com}_2(\operatorname{divL}^{\#}(x,\operatorname{switch}(y,xs)),\operatorname{switch}^{\#}(y,xs))) \end{array}
```

$$\begin{array}{l} \textit{Proc}_{\textit{RP}}(\mathcal{P},\mathcal{S}) = \{ (\mathcal{P} \setminus \mathcal{P}_{\succ}, (\mathcal{S} \setminus \mathcal{P}_{\succ}) \cup \flat(\mathcal{P}_{\succ})) \} \\ \text{(sound \& complete)} \end{array}$$

 $Proc_{RP}(\{(f2)\},\ldots)$

```
\operatorname{Com}_{2\operatorname{Pol}}(x,y) = x+y \quad \operatorname{switch}_{\operatorname{Pol}}^\#(x,xs) = 0

\operatorname{cons}_{\operatorname{Pol}}(x,xs) = xs+1 \quad \operatorname{switch}_{\operatorname{Pol}}(x,xs) = xs+1

\operatorname{divL}_{\operatorname{Pol}}^\#(x,xs) = xs \quad \dots
```

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$$Proc_{RP}(\{(f2)\},\ldots)=\{(\varnothing,\ldots)\}$$

```
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