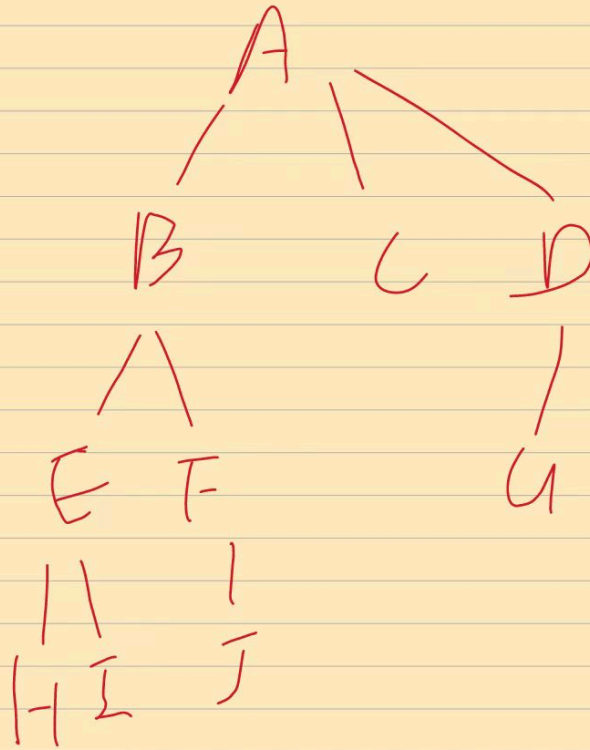


作业六

1

(1)



括号表示法

(A (B (E (H) (I)) (F (J))) (C) (D (G))))

(2)

深度: 3

高度: 4

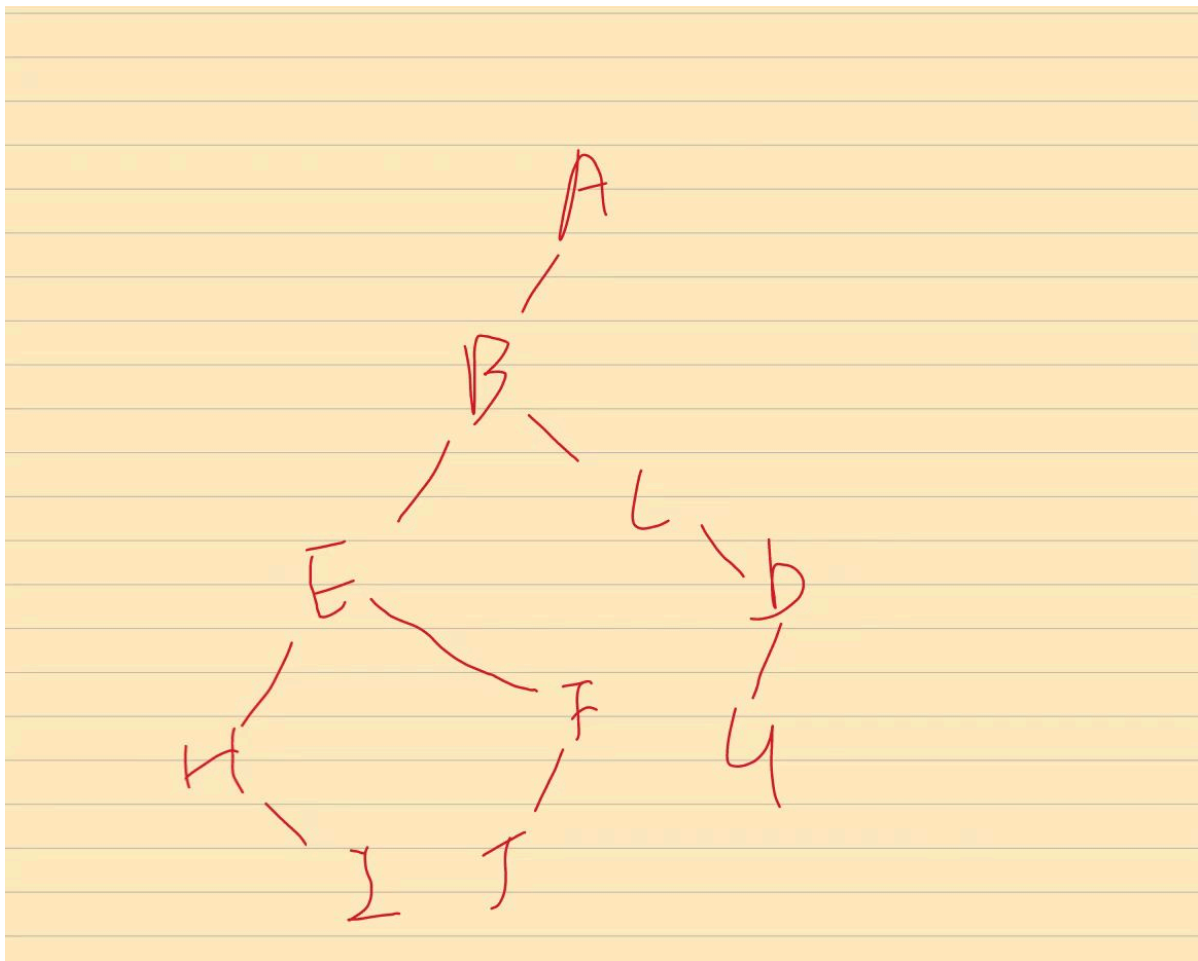
```
#include<iostream>
int result=0;
void dfs(int path,Tree<T>*&root){
    if(root==NULL){
```

```

        if(path>result){
            result=path;
        }
        return ;
    }
    for(int i=0;i<root->children.size();i++){
        dfs(path+1,root->children[i]);
    }
    return ;
}
int main(){
    dfs(0,root);
    cout<<result<<endl;
}

```

(3)



```

BinaryTreeNode *treeToBinaryTree(TreeNode *root)
{
    if(!root){
        return NULL;
    }
    BinaryTreeNode *binarytreeroot = new BinaryTreeNode(root->value);
    if(root->children.size()){
        binarytreeroot->left = treeToBinaryTree(root->children[0]);
        BinaryTreeNode *current =binarytreeroot->left;
        for (int i = 1; i < root->children.size();i++){
            current->right = treeToBinaryTree(root->children[i]);

```

```

        current = current->right;
    }
}
return binarytreeroot;
}

```

2

(1)

$$\begin{aligned}
 & \left[\frac{1 - 3^{L-1}}{-2} + 1, \frac{1 - 3^L}{-2} \right] \\
 &= \left[\frac{3^{L-1} + 1}{2}, \frac{3^L - 1}{2} \right]
 \end{aligned}$$

(2)

$$\frac{N - 1}{k}$$

(3)

$$N * k + i$$

3

每次合并时，若两个树的高度不一样，则总是将矮的树根指向较高的树的跟，因此合并后的树的高度不变，即 $\max(h_1, h_2)$

但是，当两个树的高度相同，合并后的树 = $h_1 + 1$

则当我们的树的高度为 h 时，至少需要进行 2^h 个节点的合并

因此

$$\begin{aligned}
 2^{h_{\max}} &\leq N \\
 h_{\max} &\leq \log_2 N
 \end{aligned}$$

则得证