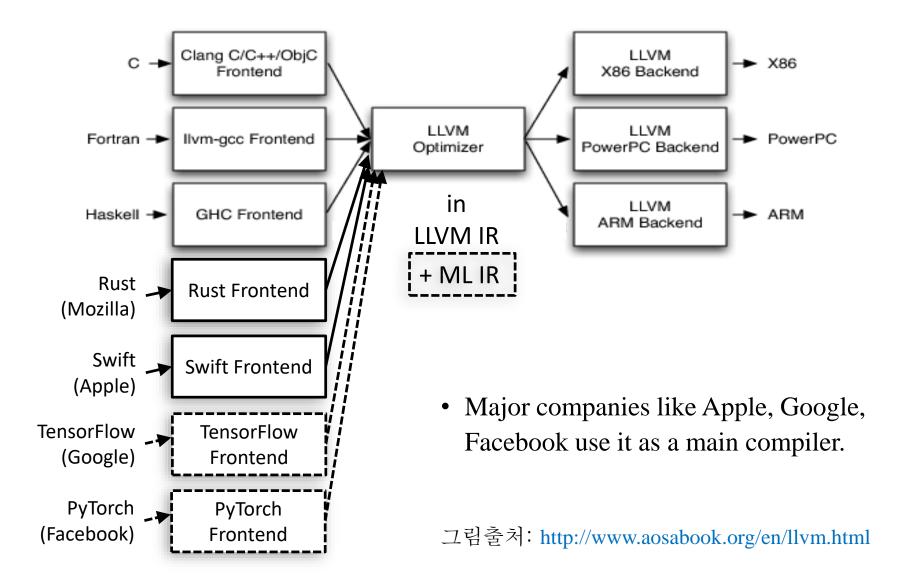
LLVM IR

Chung-Kil Hur

What is LLVM?



A running example

```
LLVM IR
int f(int x, int y) {
 int z;
 if (x == y) {
  z = x + y;
 } else {
  z = x - y;
 return z + z + z;
```

LLVM IR: Control Flow Graph (CFG)

C LLVM IR

```
int f(int x, int y)
int f(int x, int y) {
                                           br (x \neq y)
 int z;
 if (x == y) {
                          use-def chains
  z = x + y;
 } else {
  z = x - y;
 return z + z + z;
                                               ret z + z + z
```

LLVM IR: Static Single Assignment (SSA)

C

LLVM IR

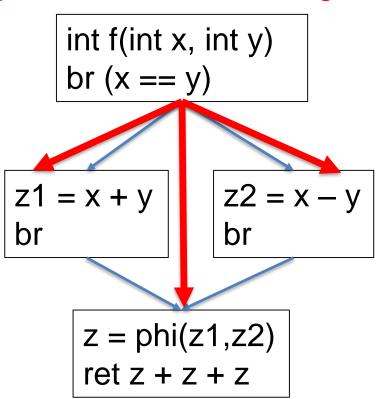
```
int f(int x, int y) {
 int z;
 if (x == y) {
  z = x + y;
 } else {
  z = x - y;
 return z + z + z;
```

```
int f(int x, int y)
    br(x == y)
  = x + y
                z^2 = x - y
br
     z = phi(z1,z2)
     ret z + z + z
```

LLVM IR: Domination Tree

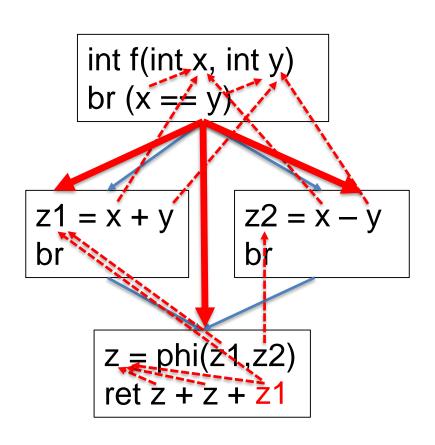
A dominates B if A is always visited before visiting B





LLVM IR: Checking use-after-def

We can check whether a variable is used after it is defined using use-def and dom tree.



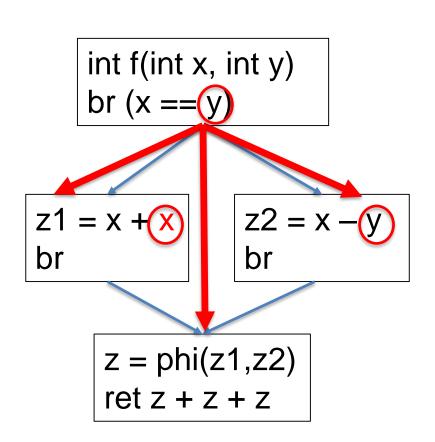
An example optimization

Goal

- replace y by x when x == y

Algorithm

 replace all uses of y that are dominated by the left successor of x == y



LLVM Syntax

```
int f(int x, int y)
br (x == y)
```

$$z1 = x + y$$
 br

$$z2 = x - y$$
 br

$$z = phi(z1,z2)$$

ret $z + z + z$

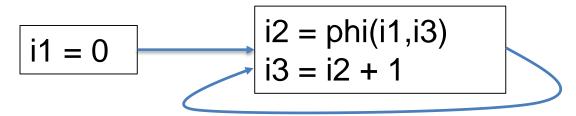
```
define i32 @f(i32 %x, i32 %y) {
entry:
  %0 = icmp eq i32 %x, %y
  br i1 %0, label %=btrue, label %bfalse
btrue:
  %z1 = add nsw i32 %x, %y
  br label %end
bfalse:
  %z2 = sub nsw i32 %x, %y
  br label %end
end:
  %z = phi i32 [ %z1, %btrue ],
               [ %z2, %bfalse ]
  %1 = add nsw i32 %z, %z
  %2 = add nsw i32 %1, %z
  ret i32 %2
```

Registers

Registers

- Syntax: %0, %x, %gil, ...
- Can use an unbounded number of registers (ie, infinitely many)
- Each register has a type: i1, i32, i64, f32, f64, i32*, ...
- Define: write a value (of the right type) into a register
- Use: read the stored value from it
- Q: Can a register be updated multiple times?

A: Yes, in the presence of loop (ie, cycle in CFG)



Q: What's difference between LLVM registers and C variables?

A: C variables reside in memory (ie, have addresses, sharable) but registers do not (ie, have no addresses, unsharable)

Memory: Stack vs. Heap

```
// Stack Allocation
                                          define i32 @f() {
int f() {
                                             %xptr = alloca i32
   int x;
                                             store i32 42, i32* %xptr
   x = 42; //*(&x) = 42
                                             call void @g(i32* %xptr)
   g(&x);
                                             %0 = load i32, i32* %xptr
   return x; // *(&x)
                                             ret i32 %0
// Heap Allocation
int f() {
                                          define i32 @f() {
  int* xptr;
                                             %xptr = call i32* @malloc(i64 4)
  xptr = malloc(4);
                                             store i32 42, i32* %xptr
  *xptr = 42;
                                             call void @g(i32* %xptr)
  g(xptr);
                                             %0 = load i32, i32* %xptr
  return *xptr;
                                             ret i32 %0
```