

# Exercises Numerical Method I 2023/24

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## 3 Matrix Analysis

**Exercise 3.1.** Define formally the addition of matrices and multiplication by a scalar in  $\mathbb{R}^{n \times m}$ .

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**Exercise 3.2.** Let  $n \geq 1$ . Let us define the *identity matrix*  $I \in \mathbb{R}^{n \times n}$  given by  $I(i, i) = 1$  for all  $i = 1, \dots, n$ , and  $I(i, j) = 0$  if  $i \neq j$ . Show that,

- $AI = A$  for all  $A \in \mathbb{R}^{m \times n}$ .
  - $IB = B$  for all  $B \in \mathbb{R}^{n \times m}$ .
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**Exercise 3.3.** Show that the multiplication of matrices is associative, i.e., show that for all  $A \in \mathbb{R}^{n \times m}$ ,  $B \in \mathbb{R}^{m \times \ell}$  and  $C \in \mathbb{R}^{\ell \times k}$ :

$$(AB)C = A(BC).$$

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**Exercise 3.4.** Show that for all  $A \in \mathbb{R}^{n \times m}$ , and  $B, C \in \mathbb{R}^{m \times \ell}$ :

$$A(B + C) = AB + AC.$$

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**Exercise 3.5.** Show that following properties are satisfied:

1. The identity matrix  $I$  is a permutation matrix.
  2. Let  $P, Q \in \mathbb{R}^{n \times n}$  be permutation matrices. Then  $PQ \in \mathbb{R}^{n \times n}$  is a permutation matrix.
  3. Let  $P \in \mathbb{R}^{n \times n}$  be a permutation matrix. Then  $P'$  is a permutation matrix. Moreover,  $PP' = P'P = I$ ; that is,  $P$  is invertible and  $P' = P^{-1}$ .
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**Exercise 3.6.** Show that following properties are satisfied:

1. The identity matrix  $I$  is a diagonal matrix.
  2.  $D \in \mathbb{R}^{n \times n}$  is a diagonal matrix if and only if it is upper triangular and lower triangular.
  3. If  $L \in \mathbb{R}^{n \times n}$  is a lower triangular matrix,  $L'$  is an upper triangular matrix.
  4. If  $U \in \mathbb{R}^{n \times n}$  is an upper triangular matrix,  $U'$  is a lower triangular matrix.
  5. If  $L_1, L_2 \in \mathbb{R}^{n \times n}$  are lower triangular matrices, so are  $L_1 + L_2$  and  $L_1 L_2$ .
  6. If  $U_1, U_2 \in \mathbb{R}^{n \times n}$  are upper triangular matrices, so are  $U_1 + U_2$  and  $U_1 U_2$ .
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**Exercise 3.7.** Let  $A \in \mathbb{R}^{n \times n}$  and  $Z$  be the matrix satisfying

$$Z(i, j) = \begin{cases} 1 & \text{if } j = i + 1, \\ 0 & \text{otherwise.} \end{cases}$$

Compute  $ZZ'$ .

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**Exercise 3.8.** Let  $\|\cdot\|$  be a matrix norm defined in  $\mathbb{R}^{n \times n}$ . Show that the following properties are satisfied:

1.  $\|I\| \geq 1$ , for  $I$  the identity matrix.
2. If  $A \in \mathbb{R}^{n \times n}$  is a non-singular matrix,  $\|A^{-1}\| \geq \|A\|^{-1}$ .

3. If  $A \in \mathbb{R}^{n \times n}$ ,  $\|A^n\| \leq \|A\|^n$  for all  $n \in \mathbb{N}$ .

**Exercise 3.9.** Show that for all  $A \in \mathbb{R}^{n \times n}$ :

$$\|A\|_F = \sqrt{\text{tr}(AA')} = \sqrt{\text{tr}(A'A)}.$$

Remember that for any matrix  $B \in \mathbb{R}^{n \times n}$ , we have that  $\text{tr}(B) = \sum_{i=1}^n B(i, i)$ ; that is, the trace of a matrix is the sum of the terms in the diagonal.

**Exercise 3.10.** Show that for all  $A \in \mathbb{R}^{n \times n}$  and  $v \in \mathbb{R}^n$ :

$$\|Av\|_2 \leq \|A\|_F \|v\|_2.$$

Moreover, give an example in which  $\|\cdot\|_F$  is not the matrix norm induced by  $\|\cdot\|_2$ .

**Exercise 3.11.** Let us consider  $(\mathbb{R}^n, \|\cdot\|_1)$ . The objective of this exercise is to show that its induced matrix norm is given by:

$$F(A) = \max_{j=1, \dots, n} \|A(:, j)\|_1.$$

1. Let  $A \in \mathbb{R}^{n \times n}$ . Show that  $\|Av\|_1 \leq F(A)$  for all  $v \in \mathbb{R}^n$  such that  $\|v\|_1 = 1$ .
2. Let  $A \in \mathbb{R}^{n \times n}$  and  $j_A$  be such that  $F(A) = \|A(:, j_A)\|_1$ . Show that  $\|Ae_{j_A}\|_1 = F(A)$ . Here  $e_{j_A}$  is the vector such that  $e_{j_A}(j_A) = 1$  and  $e_{j_A}(j) = 0$  if  $j \neq j_A$ .
3. Show that its induced matrix norm is given by  $F$ .

**Exercise 3.12.** Let us consider  $(\mathbb{R}^n, \|\cdot\|_\infty)$ . The objective of this exercise is to show that its induced matrix norm is given by:

$$F(A) = \max_{i=1, \dots, n} \|A(i, :)\|_1$$

1. Let  $A \in \mathbb{R}^{n \times n}$ . Show that  $\|Av\|_\infty \leq F(A)$  for all  $v \in \mathbb{R}^n$  such that  $\|v\|_\infty = 1$ .

2. Let  $A \in \mathbb{R}^{n \times n}$  and  $i_A$  be such that  $F(A) = \|A(i_A, :)\|_1$ . Let us define  $v$  as follows:  
 $v(j) = 1$  if  $A(i_A, j) \geq 0$  and  $v(j) = -1$  otherwise. Show that  $\|Av\|_\infty = F(A)$ .
  3. Show that its induced matrix norm is given by  $F$ .
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**Exercise 3.13.** Show that

$$F(A) = \sum_{i=1}^n \sum_{j=1}^n |A(i, j)|$$

is a matrix norm in  $\mathbb{R}^{n \times n}$ .

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**Exercise 3.14. Challenge.** Let  $f \in C^1(\mathbb{R}, \mathbb{R})$ . Then, show that for all  $x \in \mathbb{R}$  the absolute condition number of  $f$  on  $x$  is given by  $|f'(x)|$ . In addition, if  $f(x) \neq 0$ , show that the relative condition number of  $f$  on  $x$  is given by  $\frac{|xf'(x)|}{|f(x)|}$ .

## 4 Direct methods for the resolution of linear systems

**Exercise 4.1.** Let  $L \in \mathbb{R}^{n \times n}$  be a lower triangular matrix.

1. Show that  $\det(L) = L(1, 1) \cdots L(n, n)$ .
2. Show that  $L$  is non-singular if and only if  $L(i, i) \neq 0$  for all  $i \neq 0$ .

**Exercise 4.2.** Let  $L_1, L_2 \in \mathbb{R}^{n \times n}$  be lower triangular matrices. Show that:

$$(L_1 L_2)(i, i) = L_1(i, i) L_2(i, i), \quad i = 1, \dots, n.$$

**Exercise 4.3.** Let  $L \in \mathbb{R}^{n \times n}$  be a lower non-singular triangular matrices.

1. Prove that  $L^{-1}$  is also a lower triangular matrix. With that purpose, let us consider  $i, j \in \{1, \dots, n\}$  such that  $j > i$ . We have to show that  $L^{-1}(i, j) = 0$ . We recall that the inverse formula is given by:

$$L^{-1}(i, j) = (\det L)^{-1} (-1)^{i+j} \det(L([1:j-1, j+1:n], [1:i-1, i+1:n])).$$

In particular, we define  $\tilde{L} = L([1:j-1, j+1:n], [1:i-1, i+1:n])$ . Then,

- (a) Compute  $\tilde{L}(a, b)$  for  $a, b = 1, \dots, n$ .

**Hint:** you may want to consider separately the following four cases:

- $a \leq j-1, b \leq i-1$ ,
- $a \geq j, b \leq i-1$ ,
- $a \leq j-1, b \geq i$ ,
- $a \geq j, b \geq i$ .

- (b) Show that  $\tilde{L}$  is lower triangular.
- (c) Show that  $\tilde{L}(i, i) = 0$ .
- (d) Conclude that  $L^{-1}(i, j) = 0$ .

2. Show, with the Exercise 4.2, that  $(L^{-1})(i, i) = \frac{1}{L(i, i)}$ .

**Exercise 4.4.** State the analogous results of Exercises 4.1, 4.2 and 4.3 for upper triangular matrices. Prove them using that the adjoint matrix of an upper triangular matrix is a lower triangular matrix.

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**Exercise 4.5.** Obtain the LUP decomposition of the following matrices:

1.  $A_1 = \begin{pmatrix} 1 & 1 & 0 & 4 \\ 2 & -1 & 5 & 0 \\ 5 & 2 & 1 & 2 \\ -3 & 0 & 2 & 6 \end{pmatrix},$

2.  $A_2 = \begin{pmatrix} 1 & 3 & 2 \\ 1 & 5 & 3 \\ 2 & 4 & -6 \end{pmatrix},$

3.  $A_3 = \begin{pmatrix} -5 & 2 & -1 \\ 1 & 0 & 3 \\ 3 & 1 & 6 \end{pmatrix}.$

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**Exercise 4.6.** Solve the following systems using the LUP decomposition obtained in Exercise 4.5:

1.

$$\begin{aligned} x(1) + x(2) + 4x(4) &= 5, \\ 2x(1) - x(2) + 5x(3) &= -6, \\ 5x(1) + 2x(2) + x(3) + 2x(4) &= 3, \\ -3x(1) + 2x(3) + 6x(4) &= 4. \end{aligned}$$

2.

$$\begin{aligned} x(1) + 3x(2) + 2x(3) &= 5, \\ x(1) + 5x(2) + 3x(3) &= 10, \\ 2x(1) + 4x(2) - 6x(3) &= -4. \end{aligned}$$

3.

$$-5x(1) + 2x(2) - x(3) = 2,$$

$$x(1) + 3x(3) = 2,$$

$$3x(1) + x(2) + 6x(3) = 2.$$

**Exercise 4.7. LU decomposition.** Let  $A \in \mathbb{R}^{n \times n}$  be a non-singular matrix. Then, we are going to show that there is a  $LU$  decomposition in which all the terms of the diagonal of  $L$  are 1 if and only if  $\det(A(1:i, 1:i)) \neq 0$  for all  $i = 1, \dots, n$ .

- Let us first show that if  $\det(A(1:k, 1:k)) = 0$  for some  $k$ , then  $A$  does not admit a  $LU$  decomposition in which all the terms of the diagonal of  $L$  are 1. We are going to prove this result by contradiction; that is, assuming that such  $L$  and  $U$  matrices exist. In this item we use the notation  $L_k = L(1:k, 1:k)$  and  $U_k = U(1:k, 1:k)$ :

(a) Using that  $A$  is non-singular, show that  $k < n$ .

(b) Show that:

$$L_k U_k = A(1:k, 1:k). \quad (\dagger)$$

(c) Show that  $\det(L_k) = 1$  and that  $\det(L) = 1$ .

(d) Using  $(\dagger)$ , show that  $\det(U_k) = 0$  and use this to obtain that  $\det(U) = 0$ .

(e) Using  $LU = A$ , show that  $\det(U) \neq 0$ .

(f) Obtain a contradiction.

- Let us now show that if  $\det(A(1:k, 1:k)) \neq 0$  for all  $k = 1, \dots, n$  there is a  $LU$  decomposition in which all the terms of the diagonal of  $L$  are 1. For that, we are going to adapt the proof of the LUP decomposition and show by induction that for all  $k = 1, \dots, n$  there are:

- $L_k \in \mathbb{R}^{n \times n}$  a lower triangular matrix such that  $L_k(i, i) = 1$  for all  $i = 1, \dots, k$ .
- $U_k \in \mathbb{R}^{n \times n}$  an upper triangular matrix such that  $U_k(i, i) \neq 0$  for all  $i = 1, \dots, k$ .

such that:

$$(L_k U_k)(:, 1:k) = A(:, 1:k). \quad (\ddagger)$$

With that purpose:

- (a) Obtain  $U_1$ . Ensure that  $U_1(1, 1) \neq 0$ .
- (b) Obtain  $L_1$  and prove that  $(\ddagger)$  is satisfied for  $k = 1$ .
- (c) From now on we are going to suppose that the result is satisfied for  $k$  and prove it for  $k + 1$  (here  $k = 1, \dots, n - 1$ ). Obtain  $U_{k+1}$  and prove that  $U_{k+1}(i, i) \neq 0$  for all  $i = 1, \dots, k$ .
- (d) Let  $\tilde{L}_{k+1} \in \mathbb{R}^{(k+1) \times (k+1)}$  be given by:

$$\tilde{L}_{k+1}(i, j) = \begin{cases} L_k(i, j) & j \leq k, \\ 0 & i = 1, \dots, k; \quad j = k + 1, \\ 1 & (i, j) = (k + 1, k + 1). \end{cases}$$

and  $\tilde{U}_{k+1} = U(1:k + 1, 1:k + 1)$ . Show that:

$$\tilde{L}_{k+1}\tilde{U}_{k+1} = A(1:k + 1, 1:k + 1).$$

- (e) Prove that  $U_{k+1}(k + 1, k + 1) \neq 0$ .
- (f) Obtain  $L_{k+1}$ . Show that  $L_{k+1}(i, i) = 1$  for all  $i = 1, \dots, k + 1$ .
- (g) Prove that  $(\ddagger)$  is satisfied for  $k + 1$ .

**Exercise 4.8.** Let  $A \in \mathbb{R}^{n \times n}$  for  $n \geq 2$  be a matrix such that  $|A(i, i)| > \sum_{j \neq i} |A(i, j)|$  for all  $i = 1, \dots, n$  (a matrix with a *strictly dominant diagonal*).

1. Show that  $\det(A) \neq 0$ . For that:

- (a) Prove that the result holds for  $n = 2$ .
- (b) Suppose that the result is true for  $n$  and prove it for  $n + 1$ . For that, let  $B$  be given by:

$$B(i, j) = \begin{cases} A(i, 1) & i = 1, \dots, n + 1, \\ A(i, j) - \frac{A(1, j)}{A(1, 1)} A(i, 1) & \text{otherwise.} \end{cases}$$

- i. Prove that  $A(1, 1) \neq 0$  (which, in particular, implies that  $B$  is well-defined).
- ii. Prove that  $\det(A) = \det(B)$ .
- iii. Prove that  $B(1, j) = 0$  for  $j = 2, \dots, n + 1$ .



- iv. Prove that  $\det(B) = A(1, 1) \det(B(2:n+1, 2:n+1))$ .
  - v. Prove that  $B(2:n+1, 2:n+1)$  is a matrix of strictly dominant diagonal.
  - vi. Prove that  $\det(B(2:n+1, 2:n+1)) \neq 0$ .
  - vii. Prove that  $\det(A) \neq 0$ .
2. Prove that  $\det(A(1:k, 1:k)) \neq 0$  for all  $k = 1, \dots, n$ .
  3. Using Exercise 4.7, show that  $A$  admits a  $LU$  decomposition in which  $L(i, i) = 1$  for all  $i = 1, \dots, n$ .

**Exercise 4.9.** Let  $A \in \mathbb{R}^{n \times n}$  for  $n \geq 2$  be a matrix such that  $A(i, i) > \sum_{j \neq i} |A(i, j)|$  for all  $i = 1, \dots, n$  (a matrix with a *strictly dominant and positive diagonal*).

1. Show that  $\det(A) > 0$ . For that:
  - (a) Prove that the result holds for  $n = 2$ .
  - (b) Suppose that the result is true for  $n$  and prove it for  $n + 1$ . For that, let  $B$  be given by:
 
$$B(i, j) = \begin{cases} A(i, 1) & i = 1, \dots, n + 1, \\ A(i, j) - \frac{A(1, j)}{A(1, 1)} A(i, 1) & \text{otherwise.} \end{cases}$$
    - i. Show that  $A(1, 1) > 0$  (which, in particular, implies that  $B$  is well-defined).
    - ii. Prove that  $\det(A) = \det(B)$ .
    - iii. Prove that  $B(1, j) = 0$  for  $j = 2, \dots, n + 1$ .
    - iv. Prove that  $\det(B) = A(1, 1) \det(B(2:n+1, 2:n+1))$ .
    - v. Prove that  $B(2:n+1, 2:n+1)$  is a matrix of strictly dominant and positive diagonal.
    - vi. Prove that  $\det(B(2:n+1, 2:n+1)) > 0$ .
    - vii. Prove that  $\det(A) > 0$ .
2. Prove that  $\det(A(1:k, 1:k)) > 0$  for all  $k = 1, \dots, n$ .
3. Prove that if  $A$  is also symmetric, it admits a Cholesky decomposition.

**Exercise 4.10.** Show that the following matrices are positive definite and obtain their Cholesky decomposition:

1.  $\begin{pmatrix} 1 & 1 & 1 \\ 1 & 2 & 2 \\ 1 & 2 & 3 \end{pmatrix},$

2.  $\begin{pmatrix} 1 & 2 & 4 \\ 2 & 13 & 23 \\ 4 & 23 & 77 \end{pmatrix},$

3.  $\begin{pmatrix} 1 & -2 & -1 \\ -2 & 5 & 2 \\ -1 & 2 & 2 \end{pmatrix}.$

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**Exercise 4.11.** Solve these systems using the Cholesky decomposition obtained in Exercise 4.10:

1.

$$\begin{aligned}x(1) + x(2) + x(3) &= 3, \\x(1) + 2x(2) + 2x(3) &= 5, \\x(1) + 2x(2) + 3x(3) &= 6.\end{aligned}$$

2.

$$\begin{aligned}x(1) + 2x(2) + 4x(3) &= 1, \\2x(1) + 13x(2) + 23x(3) &= 2, \\4x(1) + 23x(2) + 7x(3) &= 4.\end{aligned}$$

3.

$$\begin{aligned}x(1) - 2x(2) - x(3) &= 3, \\-2x(1) + 5x(2) + 2x(3) &= -7, \\-x(1) + 2x(2) + 2x(3) &= -3.\end{aligned}$$

## 5 Iterative methods for the resolution of linear systems

**Exercise 5.1.** Using Frobenius norm, show that we can obtain the solution of the following systems with the Jacobi method. In addition, compute the first two iterations, starting in  $x_0 = 0$ :

1.

$$\begin{aligned} 10x(1) + x(2) + x(3) &= 8, \\ x(1) - 5x(2) + x(3) &= 5, \\ 2x(1) - 2x(2) + 5x(3) &= -1. \end{aligned}$$

2.

$$\begin{aligned} 2x(1) + x(2) &= 1, \\ x(1) - 2x(2) &= 3. \end{aligned}$$

3.

$$\begin{aligned} 2x(1) - x(2) &= 1, \\ x(1) + 2x(2) &= 3. \end{aligned}$$

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**Exercise 5.2.** Let  $A \in \mathbb{R}^{n \times n}$  for  $n \geq 2$  be a matrix such that  $|A(i, i)| > \sum_{j \neq i} |A(i, j)|$  for all  $i = 1, \dots, n$  (a matrix with a *strictly dominant diagonal*).

1. Show that  $A(i, i) \neq 0$  for all  $i = 1, \dots, n$ .
2. Prove that if  $\|\cdot\|$  is the induced matrix norm in  $(\mathbb{R}^n, \|\cdot\|_\infty)$ , then:

$$\|D_A^{-1}(L_A + U_A)\| < 1. \quad (\dagger)$$

**Hint:** use the formula of the induced norm obtained in Exercise 3.12.

**Remark:** this result shows that the Jacobi method is convergent for matrices with a strictly dominant diagonal. Let us now use  $(\dagger)$  to obtain  $A$  is non-singular, a result also proved in Exercise 4.8 but with a different method.

3. Show that if  $Ax = 0$ , then

$$x = -D_A^{-1}(L_A + U_A)x. \quad (\dagger)$$

4. Prove, using  $(\dagger)$  and  $(\ddagger)$ , that if  $Ax = 0$ , then  $x = 0$ .

5. Explain why  $A$  is non-singular.

**Exercise 5.3. Gauss-Seidel iterative method.** In this exercise we are going to study the Gauss-Seidel iterative method. Let  $A \in \mathbb{R}^n$  a non-singular matrix with a diagonal with non-null entries. For that we consider a normed space  $(\mathbb{R}^n, \|\cdot\|)$  and we consider  $\mathbb{R}^{n \times n}$  with its induced matrix norm. The Gauss-Seidel iterative method is very similar to the Jacobi method, but, the iterative formula is given by:

$$x_{k+1} = D_A^{-1}(b - L_A x_{k+1} - U_A x_k) \quad (\dagger)$$

1. Write the equation satisfied by  $x_{k+1}(i)$  for all  $i = 1, \dots, n$ . In particular, you must show that these  $x_{k+1}(i)$  can be directly obtained with  $x_k(j)$  for  $j > i$  and with  $x_{k+1}(j)$  for  $j < i$ .
2. Show that if  $x_{k+1}$  satisfies  $(\dagger)$ , then it satisfies:

$$x_{k+1} = (L_A + D_A)^{-1}(b - U_A x_k).$$

3. In this item we are going to prove that if  $\bar{x} = A^{-1}b$  and if  $(x_k)_{k \geq 0}$  is a sequence satisfying  $(\dagger)$ :

$$\|x_k - \bar{x}\| \leq \|(L_A + D_A)^{-1}U_A\|^k \|x_0 - \bar{x}\| \quad (\ddagger)$$

For that:

- (a) Show that  $(\ddagger)$  is satisfied for  $k = 0$ .

- (b) Prove that  $\bar{x}$  satisfies:

$$\bar{x} = (L_A + D_A)^{-1}(b - U_A \bar{x})$$

- (c) Prove that:

$$x_{k+1} - \bar{x} = -(L_A + D_A)^{-1}U_A(x_k - \bar{x}).$$

(d) Prove that:

$$\|x_{k+1} - \bar{x}\| \leq \|(L_A + D_A)^{-1}U_A\| \|x_k - \bar{x}\|$$

(e) Conclude with an inductive argument.

4. State a sufficient condition for the convergence of the Gauss-Seidel method.
5. Obtain the first iteration with the Gauss-Seidel method of Item 1 of Exercise 5.1 starting in  $x_0 = 0$ .
6. **Challenge:** show that if  $A \in \mathbb{R}^{n \times n}$  for  $n \geq 2$  be a matrix such that  $|A(i, i)| > \sum_{j \neq i} |A(i, j)|$  for all  $i = 1, \dots, n$  (a matrix with a *strictly dominant diagonal*), then the Gauss-Seidel method is convergence.

## 6 Least Square Problem

**Exercise 6.1.** Minimize, using normal equations, the following functional:

$J(x) = \|Ax - b\|_2^2$ , for

1.

$$A = \begin{pmatrix} 1 & 1 & 1 \\ -1 & 1 & 1 \\ -1 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} 4 \\ 0 \\ 1 \\ 2 \end{pmatrix}$$

2.

$$A = \begin{pmatrix} -1 & 1 \\ 2 & 1 \\ 1 & -2 \end{pmatrix}, \quad b = \begin{pmatrix} 10 \\ 5 \\ 20 \end{pmatrix}$$

3.

$$A = \begin{pmatrix} 1 & 1 \\ -1 & 3 \\ 1 & 2 \end{pmatrix}, \quad b = \begin{pmatrix} -2 \\ 0 \\ 8 \end{pmatrix}.$$

In particular, give the vector  $x$  which minimizes  $J$ , and also provide the value  $J(x)$ .

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**Exercise 6.2. Optional exercise for which Matlab is required.** Minimize, using the QU decomposition with Householder reflections, the following functional  $J(x) = \|Ax - b\|_2^2$ :

1.

$$A = \begin{pmatrix} 1 & 1 & 1 \\ -1 & 1 & 1 \\ -1 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} 4 \\ 0 \\ 1 \\ 2 \end{pmatrix}$$

2.

$$A = \begin{pmatrix} -1 & 1 \\ 2 & 1 \\ 1 & -2 \end{pmatrix}, \quad b = \begin{pmatrix} 10 \\ 5 \\ 20 \end{pmatrix}$$

3.

$$A = \begin{pmatrix} 1 & 1 \\ -1 & 3 \\ 1 & 2 \end{pmatrix}, \quad b = \begin{pmatrix} -2 \\ 0 \\ 8 \end{pmatrix}.$$

In particular, give the vector  $x$  which minimizes  $J$ , and also provide the value  $J(x)$ .

**Exercise 6.3. Linear interpolation.** Let  $n \geq 2$ . Let us consider the problem of linear interpolation. In particular, let us consider the dataset  $\{(x_1, y_1), \dots, (x_n, y_n)\}$  for  $x_1, \dots, x_n$  different real numbers and  $y_1, \dots, y_n$  any real numbers. Let us consider the set of functions  $f_a(x) = a_0 + a_1x$ . We want to find the value  $a$  where the following minimum is obtained:

$$\min_{a \in \mathbb{R}^2} \sum_{i=1}^n |f_a(x_i) - y_i|^2 \quad (\dagger)$$

For that:

1. Show that the problem  $(\dagger)$  is equivalent to:

$$\min_{a \in \mathbb{R}^2} \left\| \begin{pmatrix} 1 & x_1 \\ \vdots & \vdots \\ 1 & x_n \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \end{pmatrix} - \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix} \right\|_2^2 \quad (\ddagger)$$

2. Solve  $(\ddagger)$  using normal equations (it suffices to find the values  $a_0$  and  $a_1$ ).

**Exercise 6.4. Uniqueness of QU decomposition.**

1. Let  $Q_1, Q_2 \in \mathbb{R}^{n \times n}$  be orthogonal matrices and  $U_1, U_2 \in \mathbb{R}^{n \times n}$  be upper triangular matrices satisfying

$$U_1(i, i) > 0, \quad U_2(i, i) > 0, \quad \forall i = 1, \dots, n$$

and such that:

$$Q_1 U_1 = Q_2 U_2.$$

Then,

(a) Show that:

$$Q'_2 Q_1 = U_2 U_1^{-1}$$

(b) Let  $Q \in \mathbb{R}^{n \times n}$  be an upper triangular orthogonal matrix. Show that  $Q$  is a diagonal matrix whose entries are 1 or -1. **Hint:** use that  $Q^{-1} = Q'$  to show that  $Q^{-1}$  is, at the same time, upper and lower triangular.

(c) Show that:

$$U_2 U_1^{-1} = I_n.$$

(d) Show that:

$$U_2 = U_1.$$

(e) Show that:

$$Q_1 = Q_2.$$

2. Let us consider  $A = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$  and  $U = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$ . Obtain all the orthogonal matrices  $Q$  such that  $QU = A$ .

3. **Challenge.** Let  $A \in \mathbb{R}^{n \times m}$  with  $n > m$  be a matrix with linearly independent columns. Using any method that provides the QU decomposition for non-singular square matrices in  $\mathbb{R}^{n \times m}$ , show that the QU decomposition is not unique even if we add the restriction that  $U(i, i) > 0$  for all  $i = 1, \dots, m$ .

**Exercise 6.5. QU decomposition and Gram-Schmidt.** Let  $A \in \mathbb{R}^{n \times m}$  with  $n \geq m$  a matrix with linearly independent columns. In this exercise we are going to use the Gram-Schmidt orthonormalization algorithm to obtain a QU decomposition of  $A$ .

1. Let us consider the vectors given by:

$$\begin{cases} u_1 = A(:, 1), & v_1 = \frac{u_1}{\|u_1\|_2}, \\ u_2 = A(:, 2) - (A(:, 2) \cdot v_1)v_1, & v_2 = \frac{u_2}{\|u_2\|_2}, \\ \vdots & \vdots \\ u_m = A(:, m) - \sum_{j=1}^{m-1} (A(:, m) \cdot v_j)v_j, & v_m = \frac{u_m}{\|u_m\|_2}. \end{cases} \quad (*)$$

Prove that  $u_i \neq 0$  for all  $i = 1, \dots, m$ .



2. Show, by induction on  $m$ , that  $\{v_1, \dots, v_m\}$  is an orthonormal set in  $\mathbb{R}^n$ .

3. Show that:

$$A(:, j) = \sum_{k=1}^j (A(:, j) \cdot v_k) v_k \quad (\dagger)$$

**Hint:** show that  $\|v_j\|_2 = A(:, j) \cdot v_j$  by multiplying the expression in  $(*)$  by  $v_j$ . Alternatively, you may use properties of the orthonormal set proved in the Lineal Algebra course.

4. Provide an algorithm to obtain some vectors  $\{v_{m+1}, \dots, v_n\}$  so that  $\{v_1, \dots, v_n\}$  is an orthonormal basis of  $\mathbb{R}^n$ . **Hint:** you may use Gram-Schmidt and vectors from the canonical basis.

5. Let us consider  $Q$  the matrix defined by  $Q(i, j) = v_j(i)$ ; that is, the matrix whose columns are given by the elements of  $\{v_1, \dots, v_n\}$ . Show that  $Q$  is orthogonal.

6. Prove that for any matrix  $B_1 \in \mathbb{R}^{n \times m}$  and  $B_2 \in \mathbb{R}^{m \times \ell}$ , we have that:

$$(B_1 B_2)(:, j) = \sum_{k=1}^m B_2(k, j) B_1(:, k). \quad (\ddagger)$$

7. Let us consider  $U$  the matrix given by

$$U(i, j) = \begin{cases} A(:, j) \cdot v_i & i \leq j, \\ 0 & i > j \end{cases}$$

Prove, using  $(\dagger)$  and  $(\ddagger)$ , that  $QU = A$ .

**Exercise 6.6. Challenge.** QU decomposition has been obtained with reflections in a plane, but it can also be obtained with rotations, as rotation matrices are orthogonal matrix. Develop an analogue theory using rotations instead of reflections, and provide several examples.

## 7 Time complexity of Numerical Methods

**Exercise 7.1.** Let  $a_n, b_n$  be essentially positive sequences and  $c_n$  be any sequence. Prove that:

1.  $c_n = \mathcal{O}(|c_n|)$ .
2.  $\mathcal{O}(a_n) = \mathcal{O}(\lambda a_n), \quad \forall \lambda > 0$ .
3. If  $c_n = \mathcal{O}(b_n)$  and  $b_n = \mathcal{O}(a_n)$ , then  $c_n = \mathcal{O}(a_n)$ .

**Exercise 7.2.** Prove that the  $\Theta$  notation is an equivalence relation in the set of essentially positive sequences.

**Exercise 7.3.** Let  $a_n, b_n$  be essentially positive sequences. Prove that:

1.  $\Theta(a_n)\Theta(b_n) = \Theta(a_n b_n)$
2.  $\Theta(a_n) + \Theta(b_n) = \Theta(a_n + b_n)$

**Exercise 7.4.** Are the following statements true or false? Why?

1. If  $a_n = \mathcal{O}(n)$ , then  $\lim_{n \rightarrow \infty} \frac{a_n}{n^2} = 0$ .
2. Let  $a_n, b_n$  be essentially positive sequences. Then, either  $a_n = \mathcal{O}(b_n)$  or  $b_n = \mathcal{O}(a_n)$ .
3. Let  $a_n$  be an essentially positive sequence. Then,  $2a_n = \Theta(a_n)$ .
4. Let  $a_n, b_n$  be essentially positive sequences. If  $\lim_{n \rightarrow \infty} \frac{a_n}{b_n} = 1$ , then  $b_n = \Theta(a_n)$ .
5. Let  $a_n, b_n$  be essentially positive sequences. If  $b_n = \mathcal{O}(a_n)$ , then  $e^{b_n} = \mathcal{O}(e^{a_n})$ .

**Exercise 7.5.** Let  $A \in \mathbb{R}^{n \times n}$ . Let us consider as elementary step the following ones:

- Consulting an entry of  $A$ .
- Obtain the maximum of two elements.
- Additions.
- Taking the absolute value.

Show that the computational complexity of obtaining the matrix norm induced by the supreme norm is  $\Theta(n^2)$ . For that:

1. Using the equivalence in Exercise 3.12, show that it can be done in  $\mathcal{O}(n^2)$  steps.
2. Get a contradiction if for any algorithm that gets the norm of any matrix does not consult all the entries of all the matrices.
3. Conclude.

**Exercise 7.6.** Let  $L, U \in \mathbb{R}^{n \times n}$  be lower and upper non-singular triangular matrices respectively such that  $L(i, i) = 1$  for all  $i = 1, \dots, n$ , and  $b \in \mathbb{R}^n$ . Let us consider as elementary steps consulting a term, additions, multiplications and divisions of real numbers.

1. Show that solving the systems  $LUx = b$  requires  $\mathcal{O}(n^2)$  elementary steps. For that:
  - (a) Show that solving the systems  $Ly = b$  requires  $\mathcal{O}(n^2)$  elementary steps.
  - (b) Show that solving the systems  $Ux = y$  requires  $\mathcal{O}(n^2)$  elementary steps.
  - (c) Conclude.
2. Let us now provide a lower bound for a number of steps required for an algorithm to solve  $LUx = b$ . For that:
  - (a) Compute the number entries of  $L$  that are not given by the restrictions.
  - (b) Let us consider  $L = I$ ,  $U = I$  and  $b = (1, \dots, 1)'$ . Obtain the solution of  $LUx = b$ .
  - (c) Let us suppose that we replace an entry of  $L$  that is not given by the restrictions by 1. Obtain the solution of  $LUx = b$ .

- (d) Get a contradiction if an algorithm does not consult all the entries of  $L$  that are not obtained by the restrictions.
  - (e) Get a lower bound for a number of steps required for an algorithm to solve  $LUx = b$ .
3. Show that solving the systems  $LUx = b$  requires  $\Theta(n^2)$  elementary steps.

**Exercise 7.7.** Let  $A \in \mathbb{R}^{n \times n}$  be a non-singular matrix. Let us consider as elementary steps additions, subtractions, multiplications, divisions, computing the absolute value and taking the maximum of two numbers. Let us consider that multiplying by a permutation matrix is  $\mathcal{O}(n^2)$ , as explained throughout the course.

1. Show that there is  $C > 0$  such that the obtention of  $P_1$ ,  $P_1A$ ,  $U_1$ , and  $L_1$ , and  $U_1$  requires at most  $Cn^2$  elementary steps.
2. Show that there is  $C > 0$  independent of  $n$  or  $k$  such that the obtention of  $U_{k+1}$ ,  $P_{k+1}$ ,  $P_{k+1}A$  and  $L_{k+1}$  is at most  $Cn^2$ . For that, we show that we already know:  $U_k$ ,  $P_k$ ,  $P_kA$  and  $L_k$ . To do this:
  - (a) Get a bound for the number of steps necessary to compute  $\tilde{U}_{k+1}$ .
  - (b) Get a bound for the number of steps necessary to compute  $\tilde{P}_{k+1}$ .
  - (c) Get a bound for the number of steps necessary to compute  $P_{k+1}$ ,  $P_{k+1}A$  and  $U_{k+1}$ .
  - (d) Get a bound for the number of steps necessary to compute  $L_{k+1}$ .
  - (e) Conclude.
3. Conclude that the LUP decomposition requires at most  $\mathcal{O}(n^3)$  elementary steps.

**Exercise 7.8.** Let  $A \in \mathbb{R}^{n \times n}$  be a non-singular matrix.

1. Let us suppose that we compute the determinant of  $A$  with the Leibniz formula seen in your Linear Algebra course; that is, with

$$\det(A) = \sum_{\sigma \in S_n} \text{sgn}(\sigma) \prod_{i=1}^n a_{\sigma(i),i}.$$

How many multiplications are required if we are applying blindly that formula?

2. Let us suppose that obtaining the  $LUP$  decomposition of  $A$  requires  $\mathcal{O}(n^3)$  steps (as shown in Exercise 7.7). Provide an algorithm that computes the determinant of  $A$  in  $\mathcal{O}(n^3)$  steps (with an upper bound in the number of steps that belongs to  $\mathcal{O}(n^3)$ ). The elementary steps are the same ones as in Exercise 7.7.