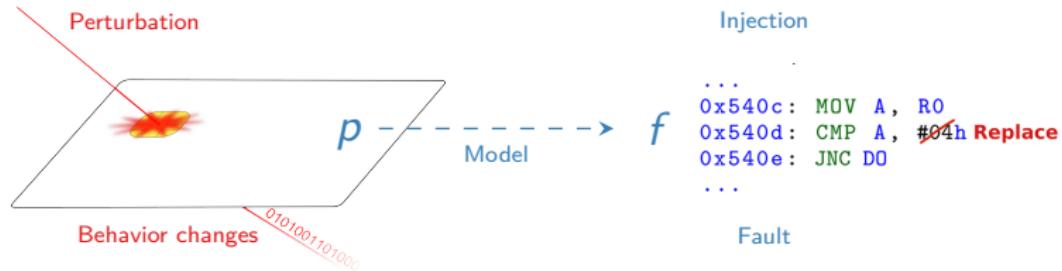


# Fault Model Inference in Practice

- (1) Laboratoire VERIMAG, Université de Grenoble-Alpes
- (2) CEA-LETI
- (3) SAFRAN IDENTITY AND SECURITY

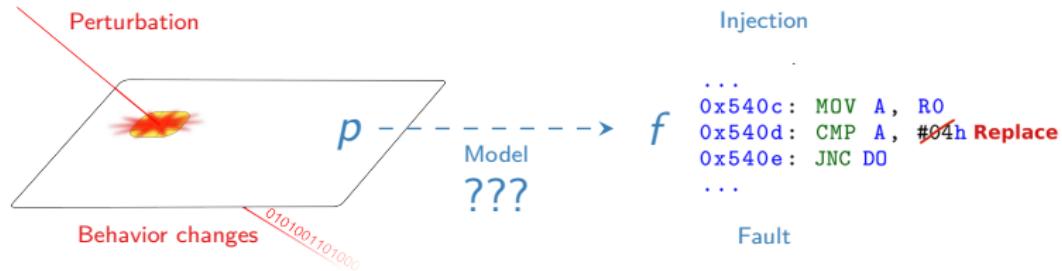
SERTIF Workshop, 2016-10-11

# Two spaces of parameters



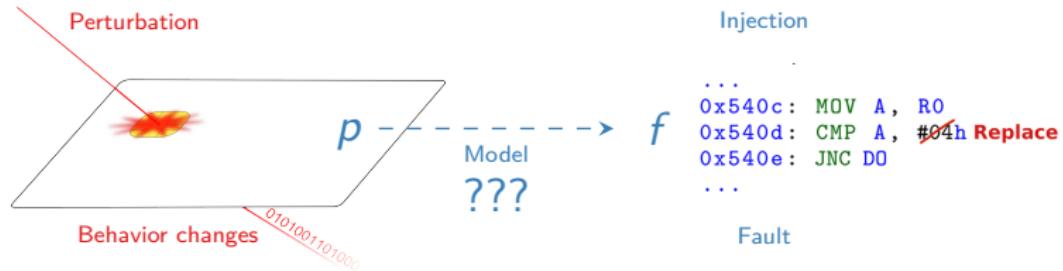
- ▶ Two spaces of parameters:
  - ▶ **parameter of the equipment**  $p \in \mathcal{P}$ :  
 $p \hat{=} (x = 12 \mu\text{m}, y = 24 \mu\text{m}, d = 3800 \text{ ns}, w = 850 \text{ ns})$
  - ▶ **effect on the code**  $f \in \mathcal{F}$ :  $f \hat{=} (i = 124, \text{store}([0x540d], 0))$

# Two spaces of parameters



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  - ▶ parameter of the equipment  $p \in \mathcal{P}$ :  
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  - ▶ effect on the code  $f \in \mathcal{F}$ :  $f \triangleq (i = 124, \text{store}([0x540d], 0))$
- ▶ How to model the effects of perturbation attack on code?

# Two spaces of parameters



- ▶ Two spaces of parameters:
  - ▶ **parameter of the equipment**  $p \in \mathcal{P}$ :  
 $p \triangleq (x = 12 \mu\text{m}, y = 24 \mu\text{m}, d = 3800 \text{ ns}, w = 850 \text{ ns})$
  - ▶ **effect on the code**  $f \in \mathcal{F}$ :  $f \triangleq (i = 124, \text{store}([0x540d], 0))$
- ▶ **How to model the effects of perturbation attack on code?**
- ▶ The model will depend on the equipment of attack, and the attacked device

# Fault as a relationship

- ▶ Fault:  $p \underset{f}{\rightsquigarrow} c$

$$(x = 12 \mu\text{m}, y = 24 \mu\text{m}, d = 3800 \text{ ns}) \underset{f_A}{\rightsquigarrow} (i = 124, \text{store(A, 0)})$$

- ▶ Fault model: set of faults

$$\{(x = 12, y = 24, d = 3000 + 200k)$$

$$\underset{f_A(k)}{\rightsquigarrow}$$

$$(i = 120 + k, \text{store(A, 0)}), k \in \mathbb{N}\}$$

- ▶ Probabilistic fault model to compute:

$$\Pr(F = f \mid p)$$

# Challenges

- ▶ The size of the space of parameters is too large  
Hundreds of years of attacks to cover the whole space!
  - ▶ Several faults can have the same effect:
    - ▶ Register corruption
    - ▶ Store instruction corruption
    - ▶ Memory corruption
- ⇒ **Black-box effect**

# Defeating the black-box effect

Lionel Rivière's PhD thesis: Fault model extraction

## Fault detection program

- ▶ Programs to disambiguate between possible faults.
- ▶ Get knowledge about the content of the black-box
- ▶ An example: EEPROM-RAM buffer copy
  - ▶ Executed from RAM
  - ▶ Sentinel RAM-RAM buffer copy

```
1 || ; main_loop:  
2 || 58: ldrb r5, [r0, #0] ; r5 <- @EEPROM  
3 || 5a: strb r5, [r2, #0] ; r5 -> @IO_EEPROM  
4 || 5c: ldrb r5, [r1, #0] ; r5 <- @RAM  
5 || 5e: strb r5, [r3, #0] ; r5 -> @IO_RAM  
6 || 60: add.w r0, r0, #1 ; @EEPROM += 1  
7 || 64: add.w r1, r1, #1 ; @RAM += 1  
8 || 68: add.w r2, r2, #1 ; @IO_EEPROM += 1  
9 || 6c: add.w r3, r3, #1 ; @IO_RAM += 1
```

However, obtained knowledge is partial

# Fault Model Inference Method

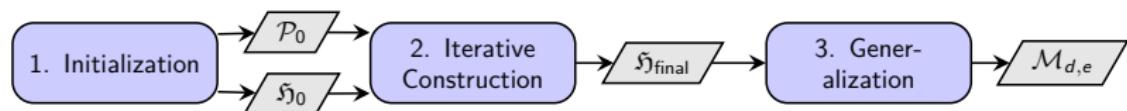
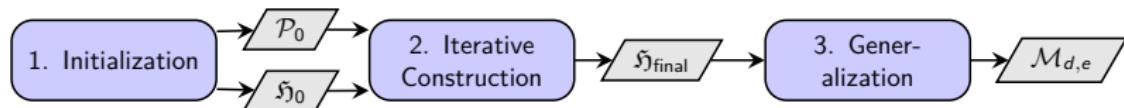


Figure: Fault model inference

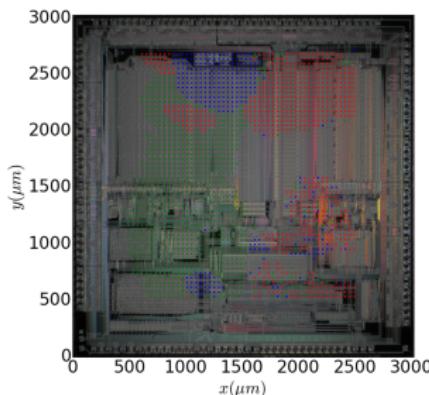
1. Initialization phase: parameter discovery to reduce the space of parameters
2. Iterative phase: physically attack several *ad-hoc* fault detection programs on the reduced space
3. Generalization phase: extend results to bigger set of parameters

# A Case Study



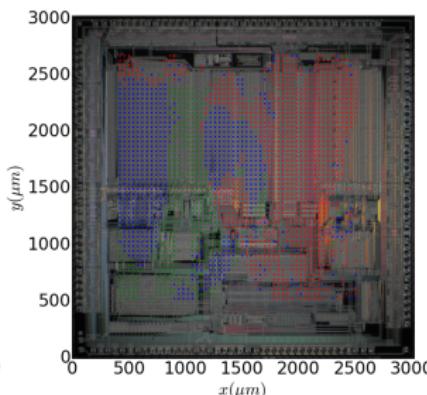
- ▶ “Card C”: “Unsecure” Cortex M-4 8 MHz
- ▶ Attacked with EM injector (100 µm copper loop with a 500 A current during 10 ns)
- ▶ The method in practice:
  1. Initialization phase: effect of the parameters of equipment
  2. Iterative phase: 3 successive programs
  3. Generalization phase

# Initialization phase: Effect of position and angle



:  $\theta = -90^\circ$

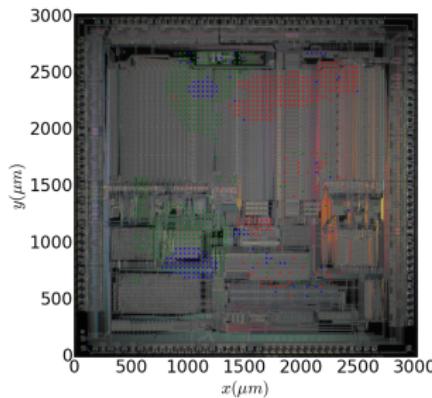
Choose one angle



:  $\theta = 0^\circ$

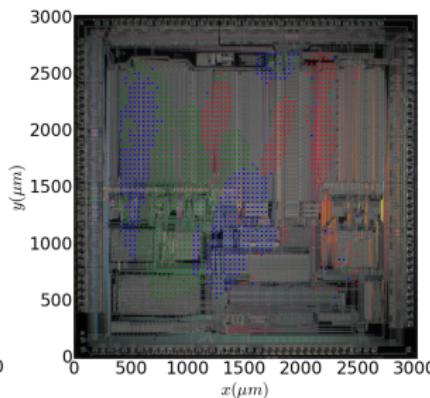
- ▶ EEPROM faults
- ▶ RAM/registers faults
- ▶ Mute

# Initialization phase: Effect of position and angle



:  $\theta = 90^\circ$

Choose one angle



:  $\theta = 180^\circ$

- ▶ EEPROM faults
- ▶ RAM/registers faults
- ▶ Mute

# Initialization phase: Effect of altitude

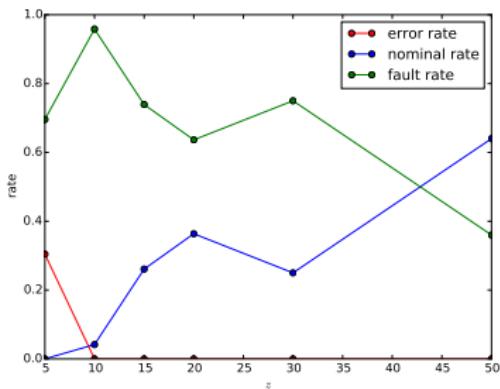
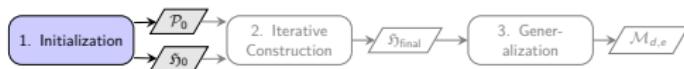


Figure: Influence of  $z$

Choose one  $z$

# Initialization phase: Effect of delay

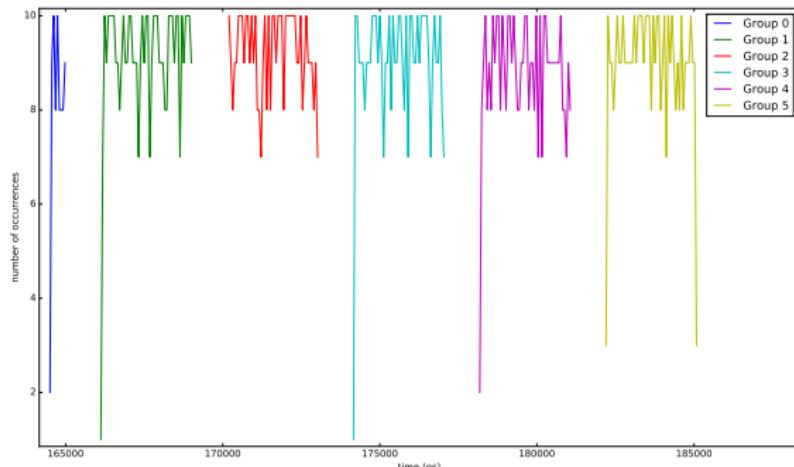
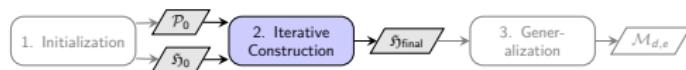


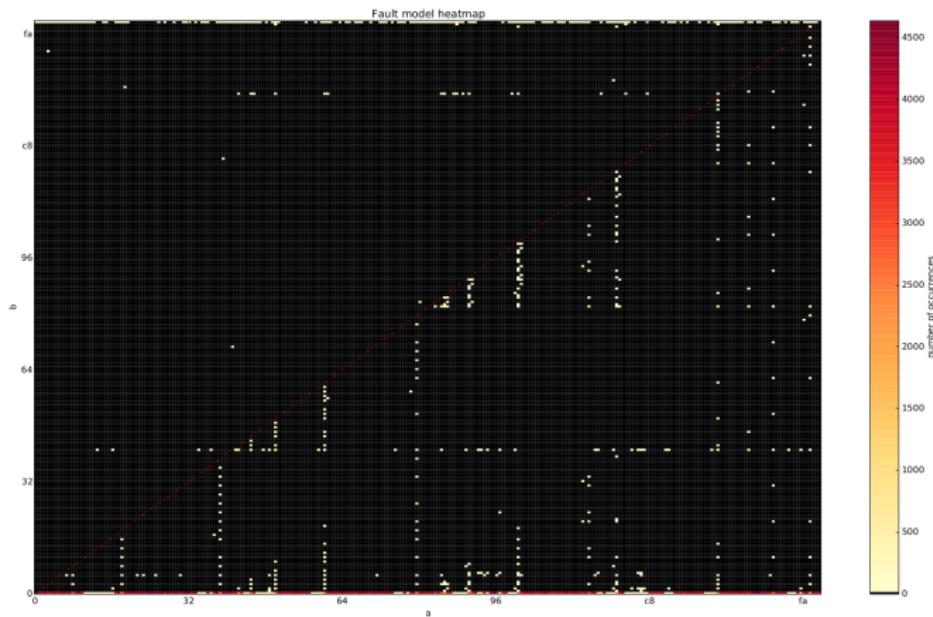
Figure: Fault count as a function of delay

Choose one delay

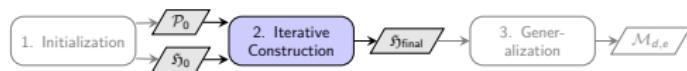
# Iterative Phase: Fault in EEPROM



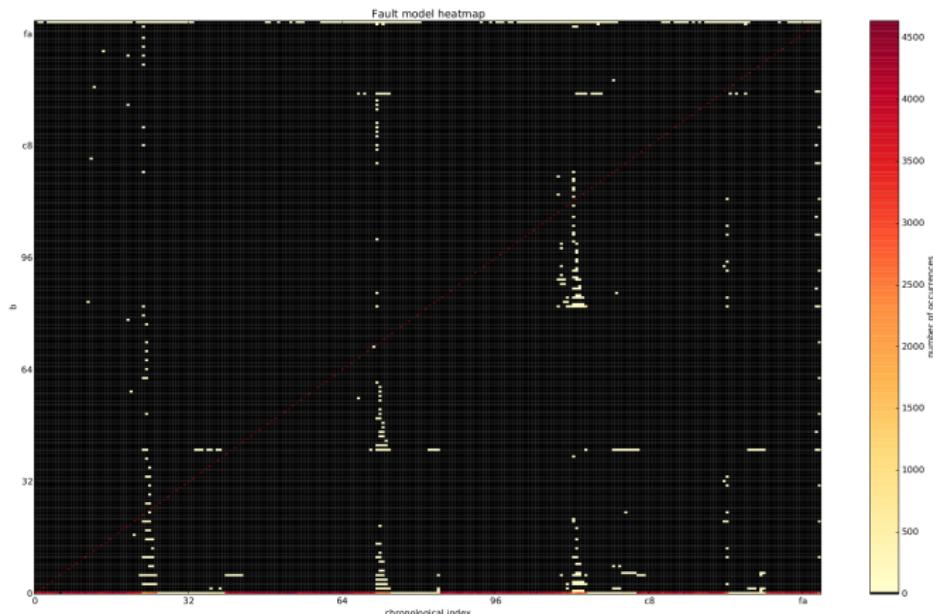
Previous knowledge: **None**



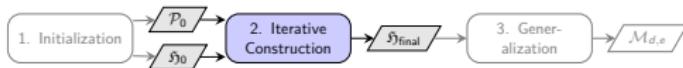
# Iterative Phase: Fault in EEPROM



Previous knowledge: **None**



# Iterative Phase: Results of Faults on EEPROM

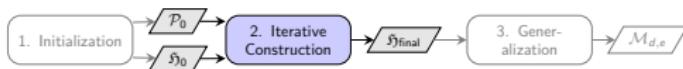


$p$  where only EEPROM reads are perturbed. Perturbations are:

- ▶ 16 consecutive bytes are faulted to 0x00 or 0xFF
- ▶ The first perturbed address has always a 16-bytes alignment

Goal: True for data EEPROM read. Check that on code!

# Iterative Phase: Effect on Code



Previous knowledge:

- ▶  $p$  where only EEPROM is faulted (no RAM or register faults)

Test:

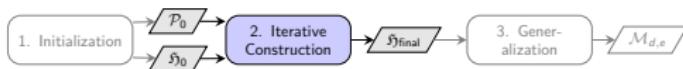
- ▶ Instruction 0x00: `movs r0, r0` is unchanged.

Program:

```
1 test_nop:  
2 ; initialization  
3 04: mov r0, I0 ; r0 <- @IO  
4 08: mov r4, I0_sentinel ; r4 <- @IO_sentinel  
5 0c: mov r1, #10 ; r1 <- 10  
6 10: mov r2, #20 ; r2 <- 20  
7 18: str r1, [r4]; r1 -> @IO_sentinel  
8 1c: str r2, [r4]; r2 -> @IO_sentinel  
9 20: movs r0, r0 ; NOP  
10 24: movs r0, r0 ; NOP  
11 ; [...]  
12 a0: movs r0, r0 ; NOP  
13 ; check in memory  
14 a4: str r1, [r0] ; r1 -> @IO  
15 a8: str r2, [r0+4] ; r2 -> @IO
```

Diagnostic: Success

# Iterative Phase: Offset Confirmation



Previous knowledge:

- ▶  $p$  where only EEPROM is faulted (no RAM or register faults)
- ▶ Instruction 0x00: `movs r0, r0`, r0 is unchanged.

Test:

- ▶ Only aligned blocks of 16 consecutive addresses are affected.

Program:

```
1 test_align:  
2 ; initialization  
3 ; [...]  
4 20: movs r0, r0 ; NOP  
5 24: movs r0, r0 ; NOP  
6 ; [...]  
7 78: movs r0, r0 ; NOP  
8 7c: adds r1, #1 ; r1 <- r1 + 1  
9 80: adds r1, #1 , r1 <- r1 + 1  
10 84: movs r0, r0 ; NOP  
11 ; [...]  
12 a0: movs r0, r0 ; NOP  
13 ; check in memory  
14 ; [...]
```

Diagnostic: Success



# Generalization Phase: Final Extracted Model



| Parameter             | Effect                             | Probability |
|-----------------------|------------------------------------|-------------|
| $(d = d_0 + k\delta)$ | 16 bytes: $(a_d \rightarrow 0x00)$ | 16%         |
| $(d = d_0 + k\delta)$ | 16 bytes: $(a_d \rightarrow 0xFF)$ | 0.3%        |

# Laser Fault Model

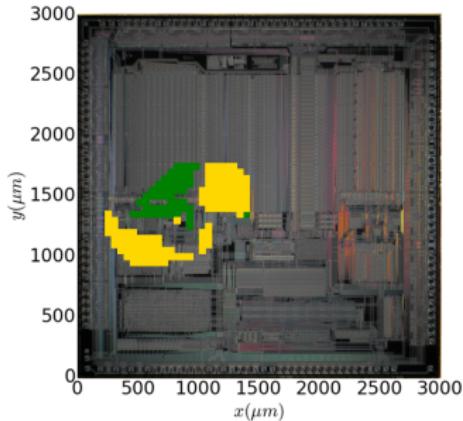


Figure: Cartography: 0xFF, 0x00

| Parameter                               | Effect                             | Probability |
|---|------------------------------------|-------------|
| $(x = x_0, y = y_0, d = d_0 + k\delta)$ | 16 bytes: $(a_d \rightarrow 0x00)$ | 21%         |
| $(x = x_1, y = y_1, d = d_0 + k\delta)$ | 16 bytes: $(a_d \rightarrow 0xFF)$ | 69%         |

# Conclusion on Fault Model Inference

- ▶ 4 inferred models
  - ▶ On 3 cards (2 Cortex-M, 1 proprietary CISC)
  - ▶ 2 with laser, 2 with EM
- ▶ High sensibility to equipment parameters
- ▶ New probabilistic aspect
- ▶ Fault Detection Programs in sequence to defeat the black-box effect
- ▶ Find model at the device level to reuse with various applications
- ▶ *ad-hoc* method... fault detection program database?

| Fault   | Pr   |
|---|------|
| $a \rightarrow 0 \mid a \neq 0$                                       | 4.8% |
| $a \rightarrow b \mid a \neq 0 \wedge d(a, b) \leq 1\%$               | 1.8% |
| $a \rightarrow b \mid a \neq 0 \wedge 1\% < d(a, b) \leq 20\%$        | 1.6% |
| $a \rightarrow b \mid a \neq 0 \wedge b \neq 0 \wedge d(a, b) > 20\%$ | 1.3% |
| $(a \rightarrow 0, a' \rightarrow 0) >   (a, a') \neq 0$              | 0.5% |

Table: Card A, EM

| Fault   | Pr    |
|---|-------|
| Bitreset of 1 byte: $a \rightarrow 0 \mid a \neq 0$ | 4.32% |
| Bitreset of 2 bytes                                 | 2.93% |
| Bitreset of 3 bytes                                 | 3.13% |
| Bitreset of 4 bytes                                 | 2.98% |
| Bitreset of 5 bytes                                 | 6.56% |
| Bitreset of 6 bytes                                 | 2.48% |

Table: Card B, laser