

# Model

## Parameters

### 卡车相关参数

- $\alpha$  : fixed cost per day per truck
- $\beta$  : transportation cost per package per unit distance
- $s$  : a sequence of index of different capacities
- $C_s$  : capacity of *sth* type of truck
- $L$  : max number of legs allowed to be traveled by a truck
- $D$  : max distance allowed to be traveled by a truck
- *Speed* : average speed of trucks, if necessary it can be truck specific
- *DrivingTimePerDay* : driving time per day allowed for trucks
- $b_{ij}^\tau = 1$  if  $\tau$  contains arc(i,j)
- $L_{so}$  : the number of trucks available starting from origin o with capacity of  $C_s$

### 节点相关参数

- $q^p$  : quantity of pickup and delivery demand  $p$
- $l_{i,j}$  : distance of arc( $i, j$ )

### Auxiliary graph $G'(V', A')$

- $V'$  : for each  $u \in V$ , associate T vertices:  $u_1, u_2, \dots, u_T$
- $A_T = \{(u_t, u_{(t+1) \bmod T}) | u \in V, t \in \{1, 2, \dots, T\}\}$
- $\tilde{A} = \{(u_t, (w_{(t+l(u,w)) \bmod T})) | (u, w) \in A, t \in \{1, 2, \dots, T\}\}$
- $A' = A_T \cup \tilde{A}$
- cost:  $A_T = 0$        $\tilde{A} = l(u, w)$

## Decision variables

- $z_\tau$  : the number of trucks choose to run in cycle  $\tau$
- $x_{i,j}^p$  : a split of demand  $q^p$  shipped on arc  $(i, j) \in \tilde{A} \cup A_T$

# Sets

- $V$ : set of nodes
- $A$ : set of arcs
- $P$ : set of demand O-D pairs
- $S$ : set of index of different capacities of trucks

# Indices

- $i, j$ : index of nodes
- $(i, j)$ : index of arcs
- $p$ : index of O-D pairs

# Const

$$b_i^p = \begin{cases} q^p & i = o(p) \\ -q^p & i = d(p) \\ 0 & \text{otherwise} \end{cases}$$

**Minimize**

$$\sum_{\tau \in \theta} \sum_{(i,j) \in \tilde{A} \cup A_T} \frac{\alpha_{ij} b_{ij}^{\tau} z_{\tau}}{\text{Speed} * \text{DrivingTimePerDay}} + \sum_{s \in S} \sum_{o \in O} \sum_{\tau \in \theta_{so}} \gamma_{so} z_{\tau} + \sum_{p \in P} \sum_{(i,j) \in \tilde{A} \cup A_T} \beta_{ij} x_{ij}^p$$

**Subject to:**

$$\sum_{(i,j) \in \delta^+(i)} x_{ij}^p - \sum_{(j,i) \in \delta^-(i)} x_{ji}^p = b_i^p \quad \forall p \in P, i \in V' \quad (1)$$

$$\sum_{p \in P} x_{ij}^p \leq \sum_{s \in S} \sum_{\tau \in \theta_s} C_s b_{ij}^{\tau} z_{\tau} \quad \forall (i, j) \in \tilde{A} \quad (2)$$

$$\sum_{\tau \in \theta_{so}} z_{\tau} + q_{so} = L_{so} \quad \forall s \in S, \forall o \in O \quad (3)$$

$$x_{ij}^p \geq 0 \quad \forall p \in P, (i, j) \in \tilde{A} \cup A_T \quad (4)$$

$$z_{\tau} \in Z \quad \forall \tau \in \theta \quad (5)$$

$$q_{so} \geq 0 \quad \forall s \in S, o \in O \quad (6)$$

$$\sum_{\tau \in \theta} b_{ij}^{\tau} z_{\tau} \leq 1 \quad \forall (i, j) \in \tilde{A}$$