

## Programming Assignment #2 Maximum Planar Subset (Dynamic Programming)

(a) Data structure (2%) How do you store the data of chords and/or other supporting information?

1. data of chord: 用一個一維vector去記錄每個vertices上chord的另一端的paired vertices

```
std::vector<int> chords(2*n); //先給各個vector一個已知的空間，之後才可以確定chords[a]=b不會出錯

while (fin >> a >> b){ //讀個line的前兩個element，第一個捨棄到junk中，把第二個存到data中
    if(a==0 && b==0){ //要兩個都用，不然有可能a=0只是它其中之一的vertices然後就停了
        break;
    }
    // data_b.push_back(a);
    // data_b.push_back(b); // data[0] will be the first data.

    ///cout<<"a: "<<a<<" b: "<<b<< endl;
    // chords.emplace_back(a, b);
    chords[a]=b;
    chords[b]=a;
    //cout<<"a="<<a<<" b="<<b<< endl;
    // data[1] will be the second data and so on.
}
```

2. other supporting information: 利用二維vector M 跟 S 去分別記錄maximum planar subset number 跟 the case of each M[i,j] is calculated from

```
std::vector<std::vector<int>> S;
std::vector<std::vector<int>> M;
```

(b) Algorithm (3%)

- Describe the algorithm of your program, including how your program finds the number of chords and the chords themselves of the maximum planar subset
- E.g. if you use dynamic programming, you should mention the recurrence equation and how to utilize your data structure to do calculation

1. Use Dynamic Programming with recurrence equation:

2. Construct table M[i,j] and S[i,j] with both size of  $2n \times 2n$

- M[i,j]: 用來存放某i,j區間下的 Maximum Planar Subset chord 數(即此區間最大互不交跨的chord的數量)
- S[i,j]: 用來存放各M[i,j] 是由recurrence equation中的哪個case計算出來的，以便在recurrence trace回去找各chord的時候有辦法回溯找subset

3. 設定base case: (M[i,j]對角線皆為0)

```
for(int i=0;i<=2*n-1;i++){
    M[i][i]=0; //initialize 對角線為0
}
```

4. 填滿M[i,j] 與 S[i,j]

從l=1 to l-2n (l: i, j間間格)

從i=0 to i=2n-1 (i: 範圍的左界)

j=i+l (j: 範圍的右界)

k=chords[j] (k 為j的chord 另一端的paired\_vertices)

利用k 的位置去判斷現在i,j為case 幾, M[i,j]要從哪個subset來

紀錄M[i,j] 為 recurrence equation case1 or 2 or 3

紀錄S[i,j] 為 計算M[i,j]時的case

```

for(int l=1;l<=2*n-1;l++){//左右i,j的兩兩間距
    //cout<<"ok12"<< endl;
    for(int i=0;i<=2*n-1-l;i++){

        int j=i+l;
        //cout<<"i="<<i<<" j="<<j<< endl;
        //cout<<"ok13 j="<<j<< endl;
        int k=chords[j]; //在j上的chord的另一端 (its chord pair)
        //cout<<"ok13 k="<<k<< endl;
        if(k<i || j<k){ //case1 <Debug3> 注意!!不可以這樣用!!if(!(i<=k<=j))
            //cout<<"ok13-1.1.1 i,j,k="<<i<<j<<k<< endl;
            M[i][j]=M[i][j-1];
            S[i][j]=1; //case1
            //cout<<"ok13-1.1.1="<<i<< endl;
        }else if(k==i){ //case2
            M[i][j]=M[i+1][j-1]+1; //+1就為kj (即ij上)這一條
            S[i][j]=2;
            //cout<<"ok13-1="<<i<< endl;
        }else{ //如果k在ij內 (k可=j, 但k!=i) //case3
            // M[i][j]=std::max(M[i][j-1],M[i][k-1]+1+M[k+1][j-1]); //M[i][k-1]+1+M[k+1][j-1] 左半(靠i那半)+kj這條+右半(靠j那半),
            // S[i][j]=3;
            if(M[i][j-1]>=M[i][k-1]+1+M[k+1][j-1]){ //如果兩種方法都一樣, 就選擇走M[i][j-1]這個方法的 (阿算出的最終chord 就會不同)
                M[i][j]=M[i][j-1]; //其實就是case1啦
                S[i][j]=1;
            }else{
                M[i][j]=M[i][k-1]+1+M[k+1][j-1];
                S[i][j]=4;
            }
            //cout<<"ok13-1="<<i<< endl;
        }
        //cout<<"ok14"<< endl;
    }
    //cout<<"ok15"<< endl;
}
}

```

5. 利用recursive 從M[0,11] 一路跟著S[i][j] 紀錄的case, 根據各case 的 recurrence equation 往回回溯, 若為case2 與 case4(case3的第二個情形), 則i,j為其中之一選擇的chord, 並存入 maximumPlanarSubset\_chord中 直到 j<=i 中止

```

void trace_maximumPlanarSubset_chord(int i,int j){
    //cout<<"now i = "<<i<<"now j = "<<j<< endl;
    //中止條件: trace回去了如果j<=i 左下角含對角線就停止這個subtree, return 0;
    if(j<=i){
        return;
    }

    if(S[i][j]==1){ //case1 //ij上面沒線, 來自M[i][j]=M[i][j-1];
        trace_maximumPlanarSubset_chord(i,j-1);
    }else if(S[i][j]==2){ //case2 //ij上面有線, 來自M[i][j]=M[i+1][j-1]+1;
        //cout<<"push i = "<<i<<"push j = "<<j<< endl;
        trace_maximumPlanarSubset_chord(i+1,j-1);

        maximumPlanarSubset_chord.push_back({i,j}); //放在回call後面, 所以出來的chord還是會照原順序
    }else if(S[i][j]==4){ //case4 //ij上面有線, 來自M[i][j]=M[i][k-1]+1+M[k+1][j-1];
        int k=chords[j];
        //cout<<"push i = "<<i<<"push j = "<<j<< endl;
        trace_maximumPlanarSubset_chord(i,k-1);
        trace_maximumPlanarSubset_chord(k+1,j-1);
        maximumPlanarSubset_chord.push_back({k,j}); //<Debug4> 這個case是kj在連線! //放在回call後面, 所以出來的chord還是會照原順序
    }
    return;
}
}

```

(c) Time Complexity Analysis (3%)

- In terms of number of chords  $n$
- Analysis based on your implementation code

Ans:

因

$$T(\text{maximumPlanarSubset}()) = \theta(n^2) \text{ \& } T(\text{trace\_maximumPlanarSubset\_chord}()) = \theta(n^2)$$

$$\Rightarrow T(n) = \theta(n^2)$$

```

class MaxPlanarSubset {
public:
    MaxPlanarSubset(int n, std::vector<int> chords) //c++ class 傳array要這樣const int* chords
    : n(n), chords(chords) {}

    void maximumPlanarSubset() {
        //Debug2-0> 不行這樣! 過了initialization 後, 要用assign!
        //Debug2-1> 要用assign!
        M.assign(2 * n, std::vector<int>(2 * n, 0)); //Debug2-2> 0-2n-1共有2n個數!! std::vector<std::vector<uint16_t>> M;
        S.assign(2 * n, std::vector<int>(2 * n, 0));

        for (int i=0; i<=2*n-1; i++) {
            M[i][i]=0; //initialize 對角線為0
        }

        for (int i=1; i<=2*n-1; i++) { //左右i,j 的兩端間距
            for (int j=i+1; j<=2*n-1; j++) {
                int k=i+1;
                int k=chords[j]; //在j上的chord的另一端 (its chord pair)

                if (k<1 || j<k) { //case1 <Debug3> 注意!! 不可以這樣用!! if (i<=k<=j)
                    M[i][j]=M[i][j]-1;
                    S[i][j]=1; //case1
                    //cout<<"ok13-1.1.1="<<i<<"<<endl;
                } else if (k==i) { //case2
                    M[i][j]=M[i+1][j]-1+1; //+1就為k (即i)上這一條
                    S[i][j]=2;
                    //cout<<"ok13-1.1="<<i<<"<<endl;
                } else { //如果k在i,j內 (k可=j, 但k=i-1) //case3
                    if (M[i][j-1]>M[i][k-1]+1+M[k+1][j]-1) { //如果兩種方法都一樣, 就選擇這M[i][j]-1這個方法的 (阿算出的最終chord 就會不同)
                        M[i][j]=M[i][j]-1; //其實就是case1囉
                        S[i][j]=1;
                    } else {
                        M[i][j]=M[i][k-1]+1+M[k+1][j]-1;
                        S[i][j]=4;
                    }
                    //cout<<"ok13-1="<<i<<"<<endl;
                }
                //cout<<"ok14"<<"<<endl;
            }
            //cout<<"ok15"<<"<<endl;
        }
        return;
    }

    std::vector<std::vector<int>> get_maximumPlanarSubset_chord() {
        trace_maximumPlanarSubset_chord(0, 2*n-1);
        return maximumPlanarSubset_chord;
    }

private:
    int n;
    std::vector<int> chords;
    std::vector<std::vector<int>> maximumPlanarSubset_chord; //uint16_t
    std::vector<std::vector<int>> S;
    std::vector<std::vector<int>> M;

    void trace_maximumPlanarSubset_chord(int i, int j) {
        //cout<<"now i = "<<i<<"<<"now j = "<<j<<"<<endl;
        //停止條件: trace回去到了如果j<=i 左下三角形對角線就停止這個subtree, return 0;
        if (j<=i) {
            return;
        }

        if (S[i][j]==1) { //case1 //i,j上面沒有線, 來自M[i][j]=M[i][j]-1;
            trace_maximumPlanarSubset_chord(i, j-1);
        } else if (S[i][j]==2) { //case2 //i,j上面有線, 來自M[i][j]=M[i+1][j]-1+1;
            //cout<<"push i = "<<i<<"<<"push j = "<<j<<"<<endl;
            trace_maximumPlanarSubset_chord(i+1, j-1);
        } else {
            maximumPlanarSubset_chord.push_back(i, j); //放在回call後面, 所以出來的chord還是會照順序
        }

        if (S[i][j]==4) { //case4 //i,j上面有線, 來自M[i][j]=M[i][k-1]+1+M[k+1][j]-1;
            int k=chords[j];
            //cout<<"push i = "<<i<<"<<"push j = "<<j<<"<<endl;
            trace_maximumPlanarSubset_chord(i, k-1);
            trace_maximumPlanarSubset_chord(k+1, j-1);
            maximumPlanarSubset_chord.push_back(k, j); //Debug4> 這個case是k,j在連線!! 放在回call後面, 所以出來的chord還是會照順序
        }
        return;
    }
};
    
```

Case 1 Case 2 Case 3 Case 4

$$\begin{aligned}
 T(N) &= \left(\frac{1}{3} + \frac{1}{6}\right) T(N-1) + \frac{1}{3} \times T(N-2) + \frac{1}{6} T\left(\frac{N}{2}\right) + \theta(1) \\
 &= \left(\frac{1}{3}\right) T(N-1) + \frac{1}{3} \times T(N-2) + \frac{1}{3} T\left(\frac{N}{2}\right) + \theta(1) \\
 &= \frac{1}{3} T_1 + \frac{1}{3} T_2 + \frac{1}{6} T_3 = \left(\frac{1}{3} + \frac{1}{3} + \frac{1}{6}\right) \theta(N) = \theta(N) \quad \text{✗} \\
 \begin{cases} T_1 = T(N-1) + \theta(1) \doteq \theta(N) \\ T_2 = T(N-2) + \theta(1) \doteq \theta(N) \\ T_3 = 2 T\left(\frac{N}{2}\right) + \theta(1) \doteq \theta(N^{\log_2 N}) = \theta(N) \end{cases}
 \end{aligned}$$

(d) README (2%)

- Clear instruction of how to compile and execute your program
- You may refer to PA#1 README

(also written in file README)

This is README file for Algorithm PA#2

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Date: 2024/11/3

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SYNOPSIS:

./bin/mps <input\_file\_name> <output\_file\_name>

ex: ./bin/mps inputs/12.in outputs/12.out

This program supports finding the maximum planar subset of a  $0 \sim 2n-1$  vertices

=====

DIRECTORY:

bin/        executable binary

doc/        reports

inputs/    input data (all chords)

outputs/   output result (maximum planar subset chords)

src/        source C++ codes

utility/   checker

=====

HOW TO COMPILE:

Under the root directory of this PA, simply type

make

=====

HOW TO RUN:

cd PA2/

./bin/mps <input\_file\_name> <output\_file\_name>

For example,

under <student\_id>\_pa1

./bin/mps inputs/12.in outputs/12.out

=====

OTHER NOTICE: