

LAB_11 ICS 202

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Task 1: find Longest non-overlapping suffix that is also a prefix:

Declare Prefix and LPS, then run one loop to iterate through all characters, each time we compare their prefix and suffix, if they are equal, assign LPS string to the longest common.

Code:

Output:

```
Enter a pattern to search for: ABABCABAB
Longest non-overlapping suffix that is also a prefix is: ABAB its length is: 4
Enter a pattern to search for: ABCDE
No non-overlapping suffix that is also a prefix.
```

Task 2: brute force string matching algorithm:

First we check if the pattern is already exist anywhere, if yes, then we iterate through all characters, each time we start we the next, if the next characters equals to the pattern, print, else skip to next.

Code:

```
// task 2
lusage
public static void overlappingPatternSearch(String T, String P){
    //check if pattern exist anywhere
    if (T.indexOf(P) == -1){
        System.out.println("Pattern not found.");
        return;
    }

    for (int i = 0; i < T.length()-P.length()+1; i++) {
        // check if pattern equals to the next strings that has the same length
        if (P.equals(T.substring(i,P.length()+i))){
            //print
            System.out.println(T);
            System.out.println(" ".repeat(i)+P);
            System.out.println(" ".repeat(i)+i);
        }
    }
}</pre>
```

Output:

```
Enter a text string T: ABABABCDABABK
Enter a pattern string P: ABAB
ABABABCDABABK
ABAB

0
ABABABCDABABK
ABAB
2
ABABABCDABABK
ABAB
2
ABABABCDABABK
ABAB
8
```

```
Enter a text string T: THIS IS KFUPM
Enter a pattern string P: YES
Pattern not found.
```

Task 3.1: generate all proper overlapping suffixes and prefixes of the string:

First we declare a prefix and a current string for the suffix, we have two loop, each time the first one run, it stores the current string from the first to the current index, then we run a second loop to add every individual character to prefix and suffix, then compare them if they are equal or not, if they are print the prefix and its length.

Code:

Output:

```
Substring: A

Substring: AB

Proper prefix: A, Proper suffix: B

Substring: ABC

Proper prefix: A, Proper suffix: BC

Substring: ABCA

Proper prefix: AB, Proper suffix: BC

Substring: ABCA

Proper prefix: AB, Proper suffix: A*1

Proper prefix: AB, Proper suffix: CA

Proper prefix: AB, Proper suffix: CA

Proper prefix: ABC, Proper suffix: A*1

Proper prefix: ABC, Proper suffix: AA

Proper prefix: ABC, Proper suffix: AA

Proper prefix: ABC, Proper suffix: ABCA

Substring: ABCABB

Proper prefix: ABC, Proper suffix: BCA

Substring: ABCABB

Proper prefix: ABC, Proper suffix: AB *2

Proper prefix: ABC, Proper suffix: CABB

Proper prefix: ABC, Proper suffix: BCABB

Substring: ABCABBC

Proper prefix: ABC, Proper suffix: CABB

Proper prefix: ABC, Proper suffix: CABB

Proper prefix: ABC, Proper suffix: CABC

Proper prefix: ABC, Proper suffix: BCABC

Proper prefix: ABC, Proper suffix: CABC

Proper prefix: ABC, Proper suffix: CABC

Proper prefix: ABC, Proper suffix: BCABC

Proper prefix: ABC, Proper suffix: CABC

Proper prefix: ABCA, Proper suffix: CABC

Proper prefix: ABCADA

Proper pre
```

Task 3.2: Manually, find the nextArray (lps array):

(a) ABCDE

(-)	<u>(w) 12 622</u>							
j	Pattern [0 j-1]	Proper prefixes	Proper Suffixes	next[j]				
	••) •]							
0	-	null	null	-1				
1	Α	λ	-	0				
2	AB	λ, Α	В	0				
3	ABC	λ, Α, ΑΒ	C, BC	0				
4	ABCD	λ, A, AB, ABC	D, CD, BCD	0				
5	ABCDE	λ, A, AB, ABC, ABCD	E, DE, CDE, BCDE	0				

The next array is:

(b) AAAAA

j	Pattern [0j-1]	Proper prefixes	Proper Suffixes	next[j]
0	-	null	null	-1
1	Α	λ	-	0
2	AA	λ,Α	Α	1
3	AAA	λ, Α,ΑΑ	A,AA	2
4	AAAA	λ, Α,ΑΑ,ΑΑΑ	A,AA,AAA	3
5	AAAAA	λ, Α, ΑΑ, ΑΑΑ, ΑΑΑΑ	A, AA, AAA,AAAA	4

The next array is:

(c) ABABAMK

j	Pattern [0j-1]	Proper prefixes	Proper Suffixes	next[j]
0	-	null	null	-1
1	Α	λ	-	0
2	AB	λ, Α	В	0
3	ABA	λ, Α, ΑΒ	A,BA	1
4	ABAB	λ, A, AB, ABA	B, AB, BAB	2
5	ABABA	λ, A, AB, ABA, ABAB	A, BA, ABA, BABA	3
6	ABABAM	λ, A, AB, ABA, ABAB, ABABA	M, AM, BAM, ABAM, BABAM	0
7	ABABAMK	λ, A, AB, ABA, ABAB, ABABAM	K, MK, AMK, BAMK, ABAMK, BABAMK	0

Task 4: Knuth-Morris-Pratt (KMP) Implementation:

We declare M and N, which are the lengths of the pattern and the text, then we Preprocess the pattern (calculate lps[] array), so lps[] will hold the longest prefix suffix values for pattern, we find lps by calling another method called "computeLPSArray(String pattern)" which basically returns a Next[j] pattern, Compare characters at text.charAt(i) and pattern.charAt(j), f they match, increment both i and j, If j reaches the end of the pattern (j == M), a match is found. Add the starting index to the indexes string and update j using the LPS array. Return the indexes string or a message if no match is found.

Code:

Output:

```
Enter a text : ABABCABABCABABCABABCABABKKKABABCABAB
Enter a pattern to search for: ABABCABAB
Pattern found at these text starting indexes: 0 7 12 17 29
```