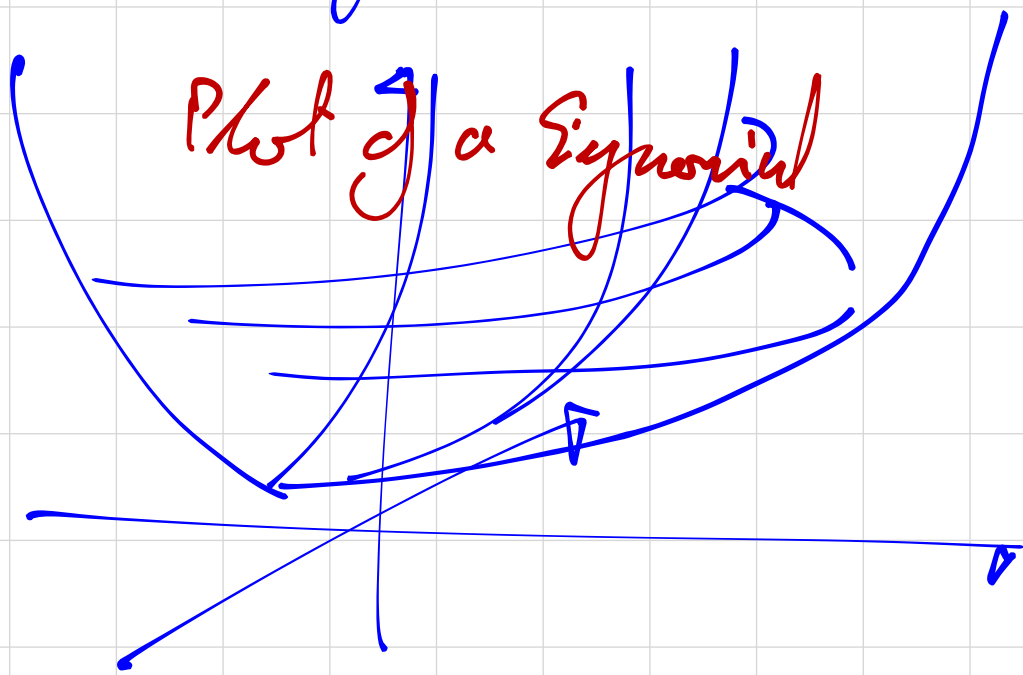


Thoughts on presentation:

- * Remind the back on Rose Kunkle present ☐
- * Get on Onuma Ven before begin! ☐
- * Look into L.'s funny slide techs ☐
- * Do the proof in recorded + written ☐
+ speed up
- * Images ☐



Example 1: Plot

Something on
the Coef
[Need on image!]

- * IPM idea
- * Do the presentation in two parts.
27 Oct
3 Nov
1 Nov 45 min
20 min - 5 min interger
20 min - 5 min ..
20 min - 15 min low
20 min last steps
Total 10 is a Recog.
- * Table should be

READING GROUP TALK: SIGNOMIAL OPTIMIZATION

TABLE OF CONTENTS / OVERVIEW

① PRELIMINARIES

Def: AM/GM-EXPONENTIALS

Def: Relation Entropy \square better Notation

\square Lemma AM/GM-EXPO's are characterized by Relation Entropy relation on Coef's
 \square Complete Proof linear relation on Expo's

\square Example

\square PICTURE (plot the example)

② The cone of Sums of AM/GM-Exponentials (SAGE)

Def: $C_{SAGE} := \sum \dots \square$ Picture Ann
Def: $SAGE := \sum$
 $C_{SAGE} := ?$

\square Proposition: Characterizing C_{SAGE} in terms of relation entropy

③ SP (Signomial programming)

Unconstrained SP

$\hookrightarrow f_{SAGE}$
 $\hookrightarrow f_{SAGE}^{(p)} \leftarrow$ Hierarchy

Constrained SP

$\hookrightarrow f_{SAGE}^{(p,q)}$

\square Picture/Gif of sequence of AM-GM-E Approx

④ \square Summary of the pipeline.

Definitions

SIGNOMIAL: $g(x) := \sum_{j=1}^l c_j \exp(\alpha^{(j)T} x)$

when $c_j \in \mathbb{R}$, $\alpha^{(j)} \in \mathbb{R}^n$ are parameters, $j \in [l]$

SIGNOMIAL OPTIMIZATION

$$\begin{array}{ll} \min & g(x) \\ \text{s.t.} & g_i(x) \geq 0 \quad \forall i \in [m] \\ & h_k(x) = 0 \quad \forall k \in [d] \end{array}$$

$$g, g_1, \dots, g_m, h_1, \dots, h_d \in \text{SIGNOMIALS}$$

SPECIAL CASES

AM/GM-EXPONENTIALS

- ① SIGNOMIAL g
- ② $g(x) \geq 0 \quad \forall x$
- ③ c_j 's > 0 EXCEPT POSSIBLY for one Denoter

$$\hookrightarrow g(x) = b \exp(\alpha^T x) + \sum_{j=1}^l c_j \exp(\alpha^{(j)T} x)$$
$$b \in \mathbb{R}, c \in \mathbb{R}_{++}^l; \alpha, \alpha^{(1)}, \dots, \alpha^{(l)} \in \mathbb{R}^n$$

RELATIVE ENTROPY

$$D: \mathbb{R}^n \times \mathbb{R}^n \ni (v, \lambda) \mapsto \sum_{j=1}^n v_j \log\left(\frac{v_j}{\lambda_j}\right) \in \mathbb{R} \cup \{\pm\infty\}$$

WEIGHTED AM/GM INEQUALITIES

$$\frac{w_1 x_1 + \dots + w_n x_n}{w} \geq \sqrt[w]{x_1^{w_1} \dots x_n^{w_n}}$$

$$\frac{1}{w} \sum_{j=1}^n w_j x_j \geq \sqrt[w]{\prod_{j=1}^n x_j^{w_j}}$$

Where $w_1, w_2, \dots, w_n > 0$ and $w = \sum_{j=1}^n w_j$
 $x_1, x_2, \dots, x_n \geq 0$

Proof: ✓

$$\ln\left(\frac{1}{w} \sum_{j=1}^n w_j x_j\right) \geq \sum_{j=1}^n \frac{w_j}{w} \ln x_j \quad \because \text{Jensen} \Rightarrow \ln \text{ is strictly concave}$$
$$= \ln\left(\sqrt[w]{\prod_{j=1}^n x_j^{w_j}}\right)$$

The result follows from apply. exp on both sides which is valid since \ln is strictly inc. □

CONVEX CONJUGATE ^{ZR}: X a Topo. space, X^* Topo. dual
 $\langle, \rangle: X, X^* \rightarrow \mathbb{R}$ duality for
 $f^*: X^* \ni x^* \mapsto \sup_{x \in X} \langle x^*, x \rangle - f$

Recall:

SLATER'S CONDITION: ^{ZK} (Sufficient condition for Strong Duality in Convex optimization problems)

Let $A \in \mathbb{R}^{n \times k}$, $b \in \mathbb{R}^k$, $f_i(x)$ convex $\forall i \in [m] \cup \{0\}$

$$\begin{array}{ll} \min & f_0(x) \\ \text{s.t.} & f_i(x) \leq 0 \quad \forall i \in [m] \\ & Ax = b \end{array}$$

Satisfies SLATER'S condition if

$$\exists x^* \in \bigcap_{i=0}^m \text{Dom}(f_i) \text{ s.t. } f_i(x^*) < 0 \quad \forall i \in [m] \\ Ax^* = b$$

Strictly feasible
Pr

If SLATER'S condition is satisfied then
STRONG Duality holds, i.e.,

LEMMA 2

Let $Q \in \mathbb{R}^{n \times \ell}$, $c \in \mathbb{R}_+^\ell$, $b \in \mathbb{R}$, $\alpha \in \mathbb{R}^n$

$$\textcircled{1} \quad g(x) := b \exp(\alpha^T x) + \sum_{j=1}^{\ell} c_j \exp(\alpha^{(j)T} x) \geq 0$$

\Leftrightarrow

$$\textcircled{2} \quad \forall v \in \mathbb{R}_+^\ell \text{ s.t. } D(v, c) - b \leq 0 \quad \text{(1) Cref's}$$

$$[\alpha^{(1)}, \dots, \alpha^{(\ell)}]v = \sum_{j=1}^{\ell} v_j \alpha^{(j)} = (\mathbb{1}^T v) \alpha \quad \text{(2) Cref's}$$

Proof: (via AM/GM)
 " \Leftarrow "

Suppose you are given a v satisfying $\textcircled{2}$

$$\frac{1}{2} \sum_{j=1}^{\ell} \left(\frac{v_j / \mathbb{1}^T v}{v_j / \mathbb{1}^T v} \right) c_j \exp(\alpha^{(j)T} x) \geq \left(\prod_{j=1}^{\ell} \left(\frac{c_j \exp(\alpha^{(j)T} x)}{v_j / \mathbb{1}^T v} \right)^{v_j / \mathbb{1}^T v} \right)^{\frac{1}{2}}$$

Weighted AM/GM
 [Recall]

$$= \prod_{j=1}^{\ell} \left(\frac{c_j \exp(\frac{v_j \alpha^{(j)T} x}{\mathbb{1}^T v})}{v_j / \mathbb{1}^T v} \right) = \prod_{j=1}^{\ell} \left(\frac{c_j}{v_j / \mathbb{1}^T v} \right)^{v_j / \mathbb{1}^T v} \exp(\alpha^T x)$$

$$\sum_j v_j \alpha^{(j)T} = \mathbb{1}^T v \alpha^T \quad (2.2)$$

Where

$$\prod_{j=1}^{\ell} \left(\frac{c_j}{v_j / \mathbb{1}^T v} \right)^{v_j / \mathbb{1}^T v} \geq b$$

$$\prod_{j=1}^l \left(\frac{c_j}{v_j / \mathbb{1}^T v} \right)^{v_j / \mathbb{1}^T v} = \exp \{ -D(\frac{v}{\mathbb{1}^T v}, c) \}$$

∴ Def of Rel. Ent.
[Recall]

$$\exp \{ -D(\frac{v}{\mathbb{1}^T v}, c) \} \geq -\xi [D(\frac{v}{\mathbb{1}^T v}, c) + \log(\xi) - 1] \quad \forall \xi > 0$$

∴ Definition of Convex Conjugate: \hookrightarrow

$$f^*: X^* \ni x^* \mapsto \sup_{\xi \geq 0} \langle x^*, \xi \rangle - f(\xi) \quad (1)$$

Claim: $\exp(-x)$ is the Conv-Conj of $x \log x$.

$$\exp(-p) = \sup_{\xi \geq 0} (1-p) \cdot \xi - \xi \log \xi$$

[Recall]

$$= -\xi \left(D(\frac{v}{\mathbb{1}^T v}, c) + \log(\xi) - \log(e) \right)$$

$$= -\xi \left(\sum_{j=1}^l \frac{v_j}{\mathbb{1}^T v} \log \left(\frac{v_j}{c_j \mathbb{1}^T v} \right) + \log \left(\frac{\xi}{e} \right) \right)$$

$$= -\sum_{j=1}^l \xi \frac{v_j}{\mathbb{1}^T v} \log \left(\frac{\xi v_j}{\mathbb{1}^T v} \cdot \frac{1}{c_j e} \right)$$

$$= -D\left(\frac{\xi v}{\mathbb{1}^T v}, ec\right)$$

∴ Def. Rel. Ent.

$$\geq -D(v, ec) \quad \text{∴ Choose } \xi := \mathbb{1}^T v$$

$$\geq -b \quad \text{∴ (2.1)}$$

□

" \Rightarrow "

$$g(x) \geq 0$$

\Leftrightarrow

$$g(x) \exp(-\alpha^T x) \geq 0$$

$\Delta \Rightarrow$

$$\circ \circ \exp(-\alpha^T x) > 0$$

$$\sum_{j=1}^l c_j \exp((\alpha^{(j)})^T x - \alpha^T x) \geq -b$$

$$\circ \circ \text{Def } g(x)$$

Consider the problem

$$P := \inf_{x \in \mathbb{R}^n} c^T \epsilon$$

$$\text{s.t. } \exp(\alpha^{(j)T} x - \alpha^T x) = \epsilon_j \quad \forall j \in [l]$$

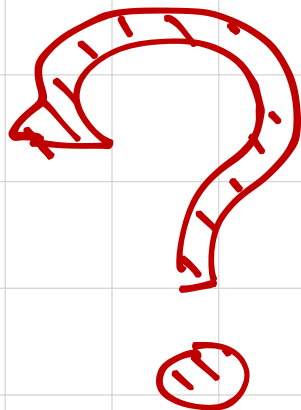
We want $P \geq -b$

$$P = \inf_{x \in \mathbb{R}^n, \epsilon = \mathbb{R}^l} c^T \epsilon$$

$$\text{s.t. } \exp(\alpha^{(j)T} x - \alpha^T x) \leq \epsilon_j \quad \forall j \in [l]$$

$$\circ \circ \begin{aligned} &\leq \because \text{Def } \inf \\ &\geq \because c, \epsilon \geq 0 \end{aligned}$$

$$\sup_{\lambda \in \mathbb{R}_+^l} \inf_{x \in \mathbb{R}^n, \epsilon = \mathbb{R}^l} (1 - \lambda)^T (\epsilon_j - \exp(\dots))_{j=1}^l$$



$$\begin{array}{l} \text{S.t.} \\ \text{S.F.} \end{array} \quad - \sum_{j=1}^l v_j \log\left(\frac{v_j}{e c_j}\right)$$

$$[\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(l)}] v = (\pi^T v) \alpha$$

\Leftrightarrow

$$\begin{array}{l} \text{S.t.} \\ \text{S.F.} \end{array} \quad -D(v, eC)$$

$$[\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(l)}] v = (\pi^T v) \alpha$$

Immer SLATER'S CONDITION

[RECALL]

Section 2: Characterizing the Cone of Sums of AM/GM Exponentials

Exponents (fixed)

$$(\alpha, \alpha^{(1)}, \dots, \alpha^{(L)})$$

Coefficients ("Variables")

$$(b, c_1, \dots, c_L)$$

$$g(x) := b \exp(\alpha^T x) + \sum_{j=1}^L c_j \exp(\alpha^{(j)T} x) \geq 0$$



$$\forall v \in \mathbb{R}_+^L$$

$$\text{s.t. } Av := \begin{bmatrix} \alpha^{(1)} & \alpha^{(2)} & \dots & \alpha^{(L)} \\ | & | & & | \\ | & | & & | \\ | & | & & | \end{bmatrix} v = (\mathbf{1}^T v) \alpha$$

$$D(v, c) - b \leq 0$$

} Linear Constant on v

} Entropy Const. S on v and w.r.t. c, b

What are we gonna do with this local?

max

b

$$\text{s.t. } \sum_{j=1}^L c_j \exp(\alpha^{(j)T} x) - b \geq 0$$

$g(x)$ with

$$\alpha = 0$$

$$b \leq 0$$



$$\text{s.t. } \forall v \in \mathbb{R}_+^L$$

$$\text{s.t. } Av := \begin{bmatrix} \alpha^{(1)} & \alpha^{(2)} & \dots & \alpha^{(L)} \\ | & | & & | \\ | & | & & | \\ | & | & & | \end{bmatrix} v = 0$$

$$D(v, c) - b \leq 0$$

Definition

For a finite set of vectors $M \subseteq \mathbb{R}^n$

$$\text{SAGE}(M) := \left\{ f = \sum_{i=1}^l f_i : \left| \begin{array}{l} f_i \text{ is AM/GM Expo} \\ \text{using Exponent. in } M \end{array} \right. \right\}$$

So if $M = \{\alpha^{(j)}\}_{j=1}^l$ then

$$\text{SAGE}(\{\alpha^{(j)}\}_{j=1}^l) = \text{the paper gives a description with only } l \text{ terms..}$$

SAGE $\not\subset$ all Nonnegative Polynomial

Example

$$(\exp(x_1) - \exp(x_2) - \exp(x_3))^2 \geq 0$$

is not SAGE because ... WHY

$$\sum_{i=1}^3 \exp\left(2e_i^T \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}\right) + \exp([0, 1, 1]x)$$

$$- \exp([1, 1, 0]x) - \exp([0, 1, 1]x)$$

$$C_{\text{SAGE}}(\{\alpha^{(j)}\}_{j=2}^L) := \left\{ c \in \mathbb{R}^L \mid \sum_{j=2}^L c_j \exp(\alpha^{(j)T} x) \in \text{SAGE}(\{\alpha^{(j)}\}_{j=2}^L) \right\}$$

"The SET of all coefficients that correspond to AM/GM-Expo's in exponent $\{\alpha^{(j)}\}_{j=2}^L$."

Proposition 6 \Downarrow

$$C_{\text{SAGE}}(\{\alpha^{(j)}\}_{j=2}^L) := \text{Def}$$

$$= \text{Characterization 1}$$

$$= \text{Characterization 2}$$

} Ugly.

$$C_{\text{SAGE}}^*(\{\alpha^{(j)}\}_{j=2}^L) = \{ \dots \text{Relative Entropy} \dots \}$$

Point 3

Unconstrained Signomial Programming

$$f_* = \inf_{x \in \mathbb{R}^n} f(x)$$

\Leftrightarrow

$$f_* = \sup_{\gamma \in \mathbb{R}} \gamma$$

s.t. $f(x) - \gamma \in \text{"Nonnegative Signomials"}$

NON-Negativity
Characteristics

Assume w.l.o.g. $f(x) = \sum_{j=2}^l c_j \exp(\alpha^{(j)T} x)$

with $c_1 \in \mathbb{R}, \dots ?$

Relax:

$$f_{\text{SAGE}} = \sup_{\gamma \in \mathbb{R}} \gamma$$

s.t. $f(x) - \gamma \in \text{SAGE}(\{\alpha^{(j)}\}_{j=2}^l)$

Prop 6.

$$f_{\text{SAGE}} = \sup_{\gamma \in \mathbb{R}} \gamma$$

s.t. $(c_1 - \gamma, c_2, \dots, c_l) \in \text{SAGE}(\{\alpha^{(j)}\}_{j=2}^l)$

(Make steps explicit)

EXTEND

$$f_{\text{SAGE}}^{(P)} = \sup_{\gamma \in \mathbb{R}} \gamma$$

s.t. $\left(\sum_{j=2}^l \exp(\alpha^{(j)T} x) \right)^P [f(x) - \gamma] \in \text{SAGE}(\{\alpha^{(j)}\}_{j=2}^l)$

Question:

Is Symon's/optimized NP-hard?
↳ what is the proof.

Notes: Make decisions : is a different choice.