



Computer Networks

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Chapter 6.

Congestion Control and QoS

- Network Congestion
- Congestion Control in FR
- Traffic Management in ATM
- Internet QoS
- Resource Allocation and RSVP
- Differentiated Services



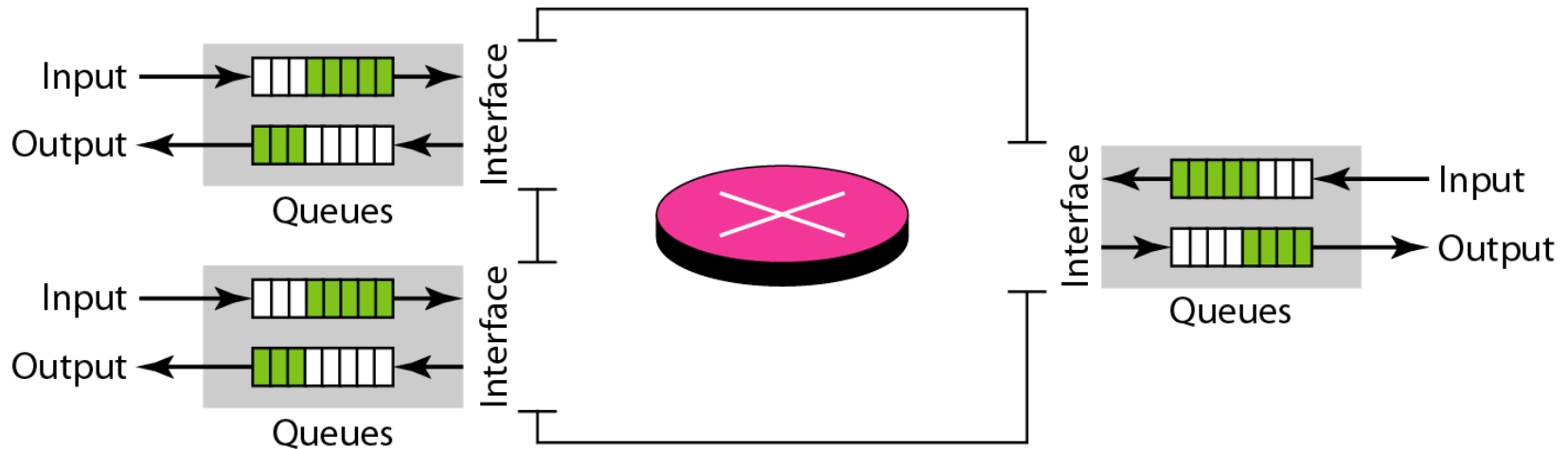
Network Congestion

- Congestion
 - Number of packets transmitted through the network approaches the packet handling capacity of the network
- One or more switches/routers becomes overloaded
 - Generally 80% utilization is critical
- Congestion control
 - Keep number of packets below level at which performance falls off dramatically



Queues at a Switch

- Switch **overloads** because receiving packets faster than it can forward

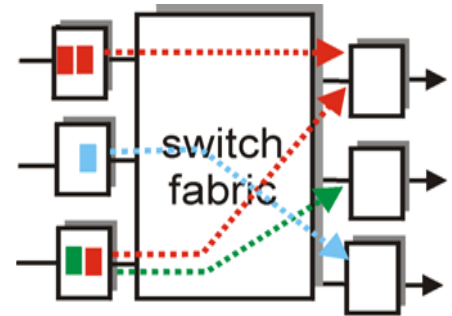




Causing Congestion

■ Congestion at switch

- Bursty traffic / poor topology
- Packet arrival rate **exceeds the outgoing link capacity**



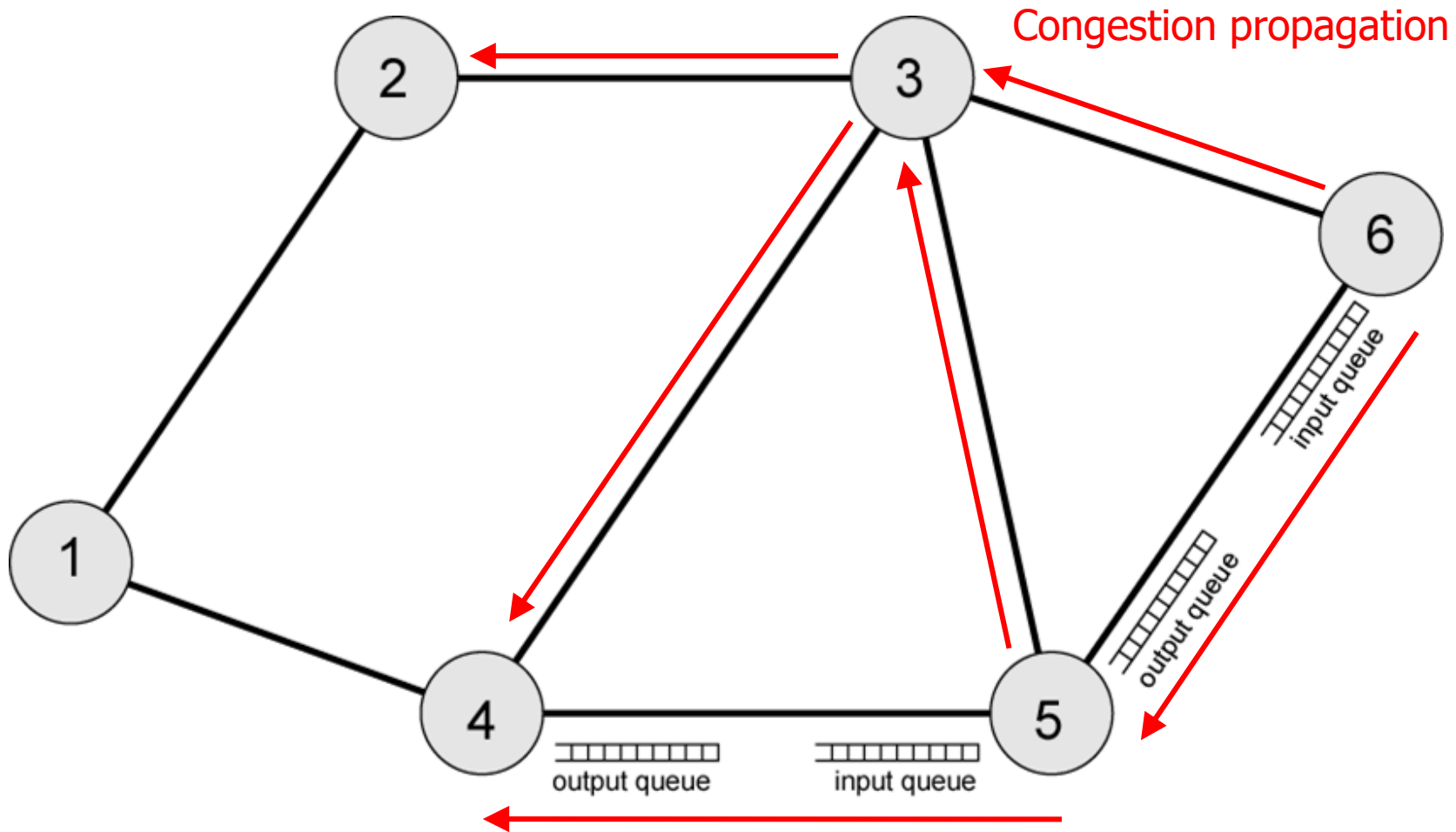
- Packet processing rate $<$ packet arriving rate
- **Insufficient memory** to store arriving packets

■ Effects caused at congested switch

- **Discard queued packets** to make room for new comings
- **Prevent additional packets** from entering the congested port (link-layer flow control)

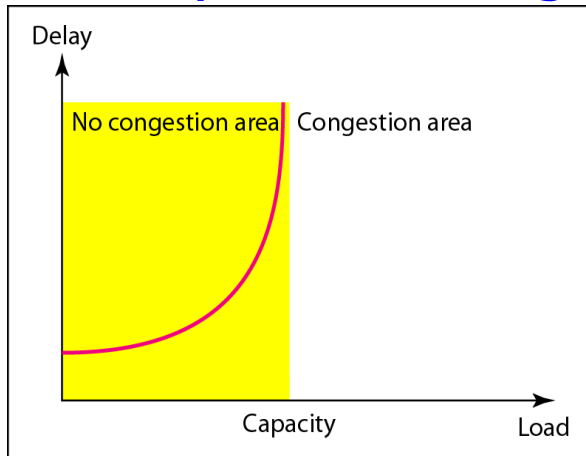


Interaction of Queues

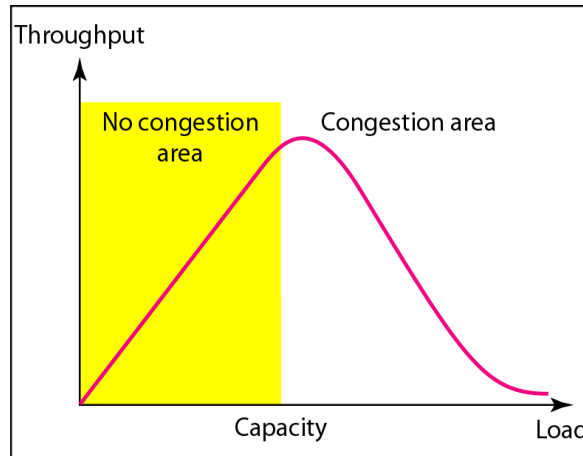


Network Utilization

■ Delay and Throughput vs. Network Load



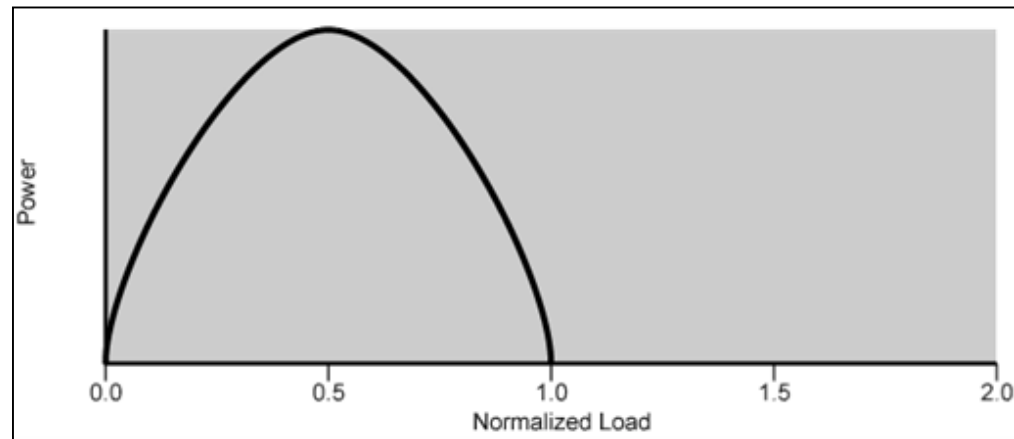
a. Delay as a function of load



b. Throughput as a function of load

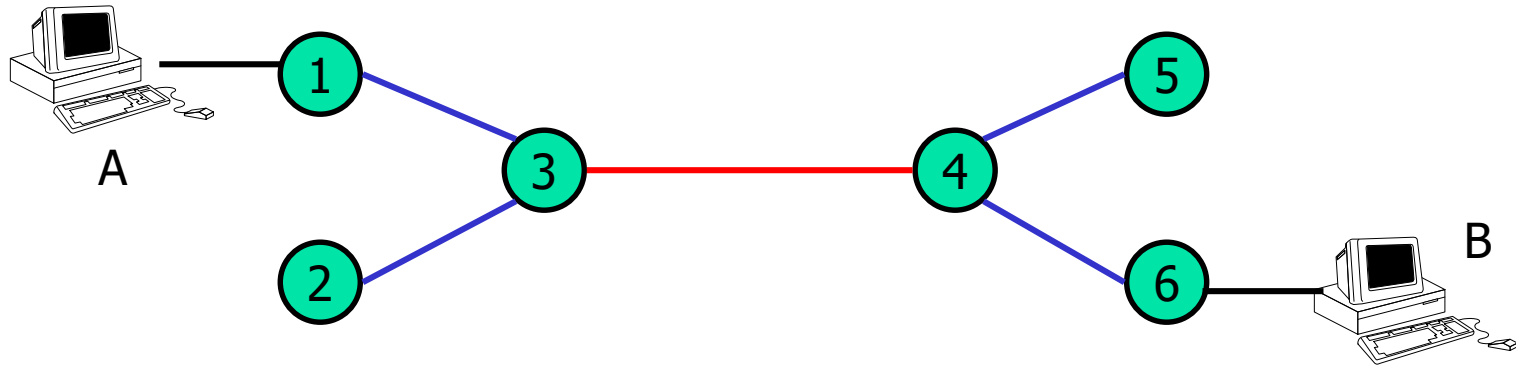
■ Communication Power

$$Power = \frac{Throughput}{Delay}$$





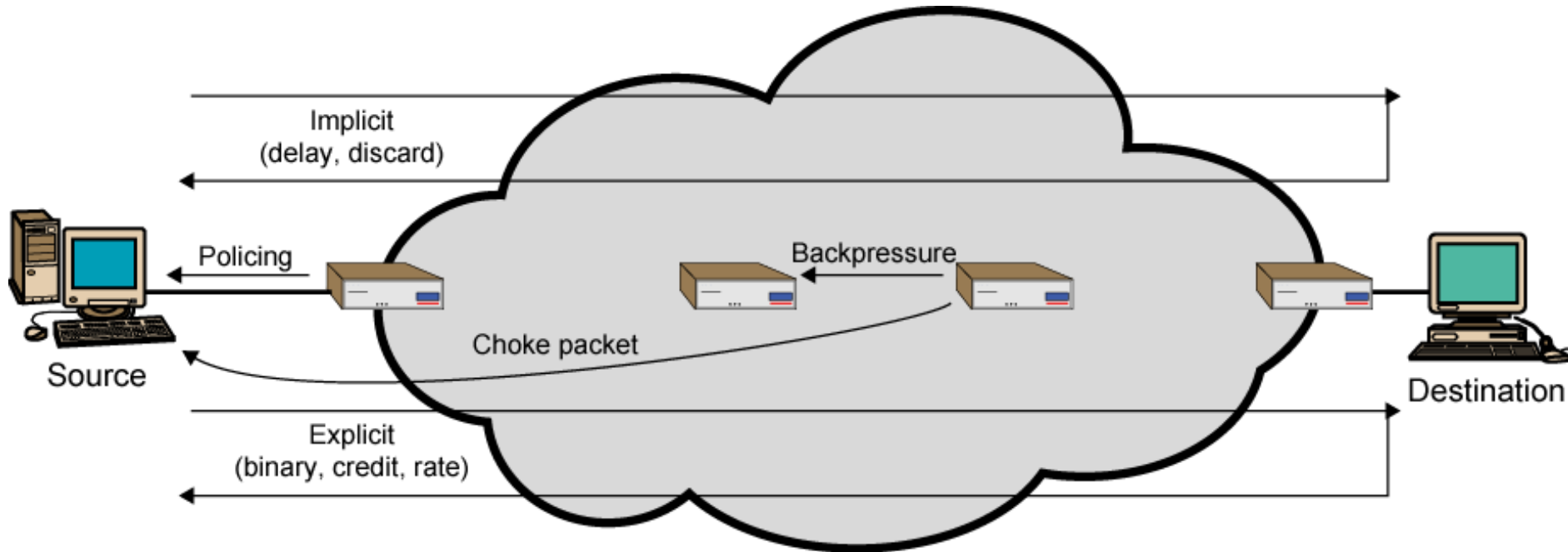
Bottleneck Effect



- Assume all the links have similar capacity, and run in full for both direction
- Then switches 3 and 4 will be **in congestion**



Mechanisms for Congestion Control

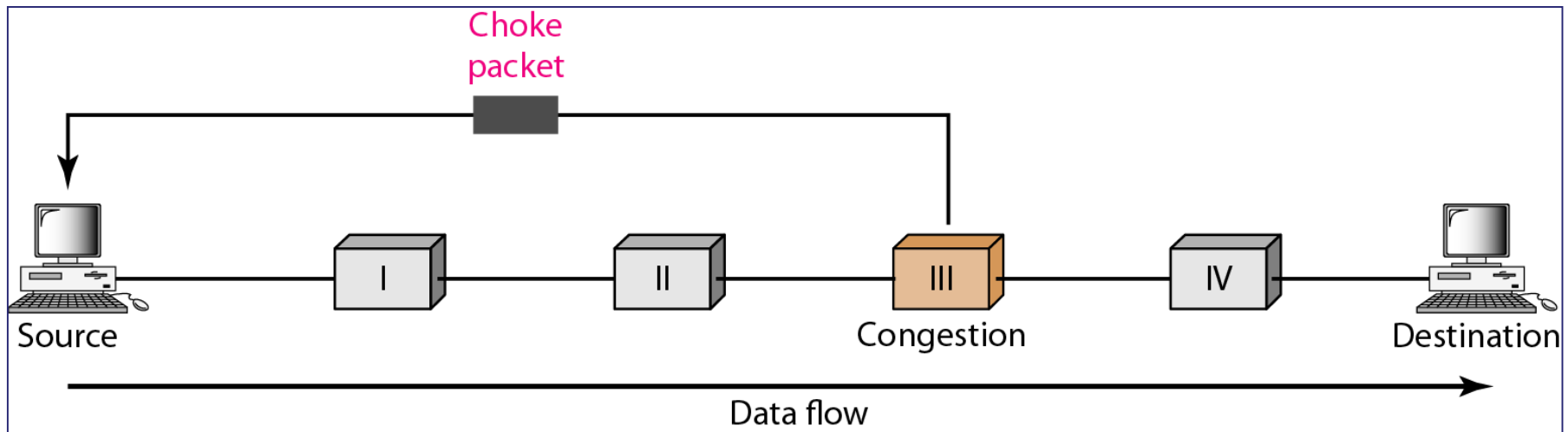


- Choke Packet
- Backpressure
- Warning bit
- Congestion window
- Random early discard
- Traffic shaping

- 抑制分组
- 反压
- 警告位
- 拥塞窗口
- 随机早期丢弃
- 流量整形

(1) Choke Packet

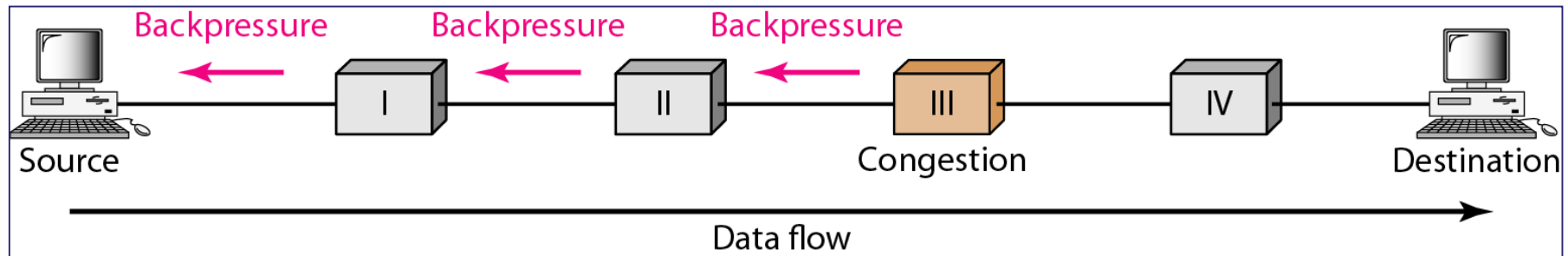
- Control packet
 - Generated at **congested node**
 - Sent to source node
- **Source quench:** using ICMP to notify source
 - From router or destination, sent for every discarded packet



(2) Backpressure

■ Hop-by-Hop Choke Packets

- Propagation time $>$ transmission time (long distance or high speed link)
- **Choke packets** from router to source are not effective
- Require each hop to reduce its transmission



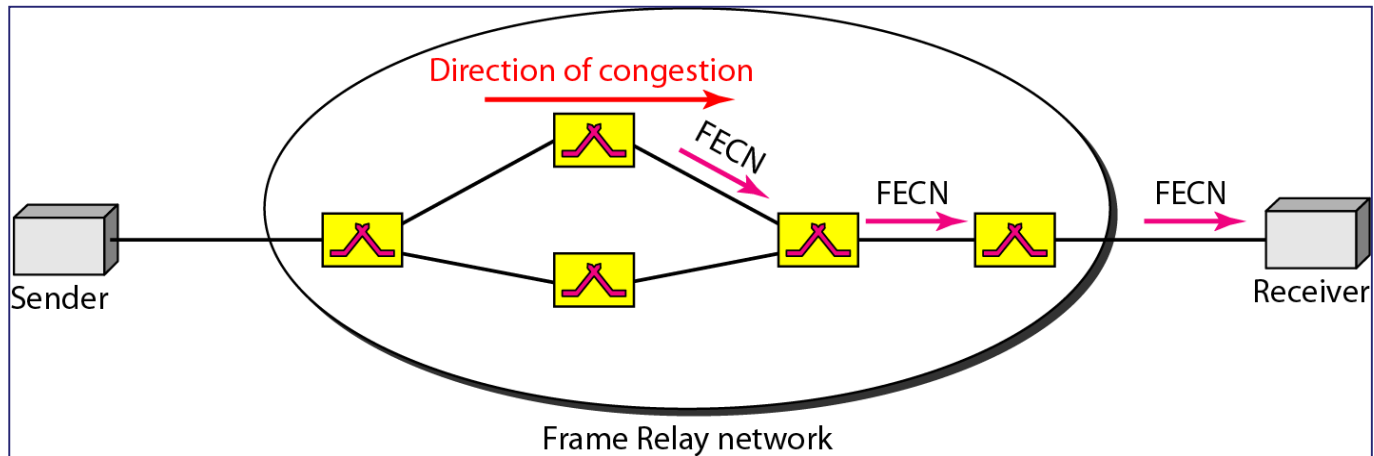
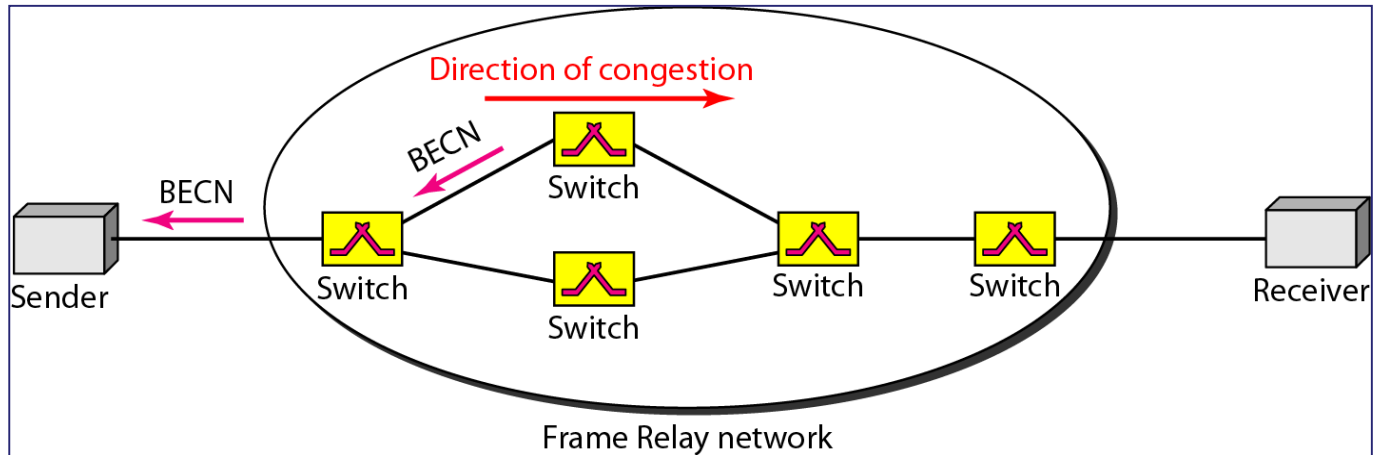


(3) Warning Bit

- Special bits set in the packet header by switches
 - Alerts end systems of increasing congestion
 - End systems take steps to reduce offered load
- Backwards
 - Congestion avoidance in opposite direction to congested packet
 - Assume congestion will burst up quickly
- Forwards
 - Congestion avoidance in same direction as congested packet
 - Assume congestion will cumulate slowly

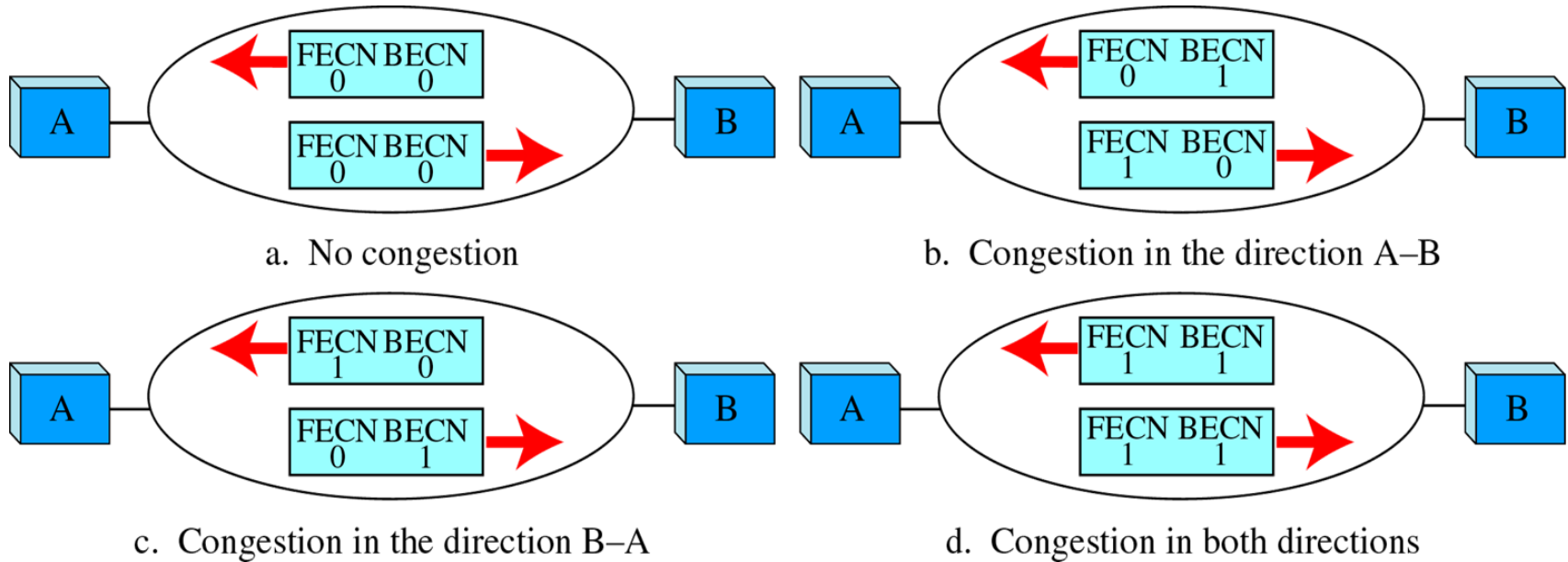


BECN and FECN in FR





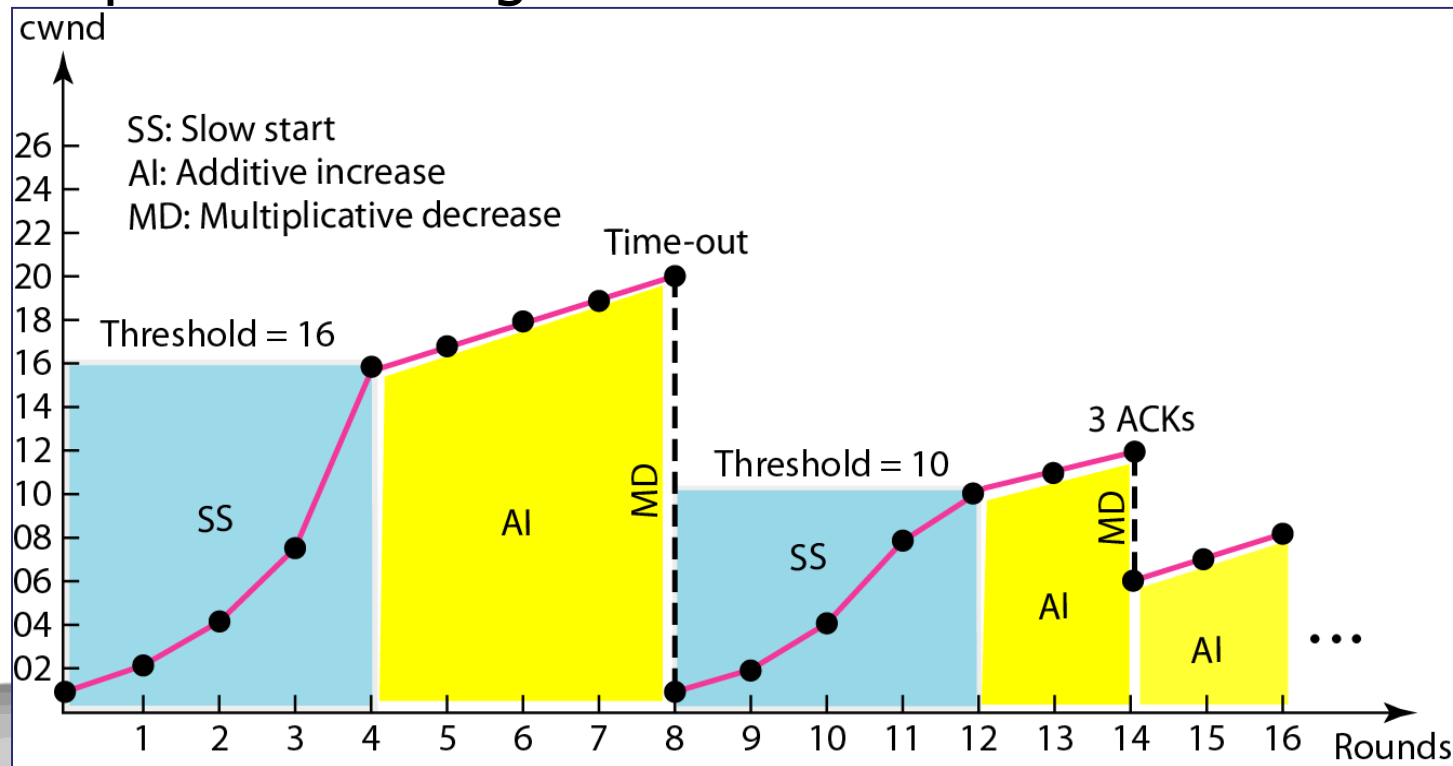
Four Cases of Congestion



(4) Congestion Window

■ Control congestion at hosts

- Packet timeout as a signal of network congestion
- **Dynamic send window management** (as in TCP) to hold the packet sending





(5) Random Early Discard

- Control congestion at routers (switches)
 - Combined with congestion window at hosts
- Internet (TCP) global synchronization problem
 - Traffic burst fills queues so packets lost, TCP connections enter slow start
 - Traffic drops so network under utilized, connections leave slow start at same time causing burst again
- Handle the problem – RED
 - Router randomly discards packets before buffer becomes completely full



The RED Algorithm

■ Compute average queue length

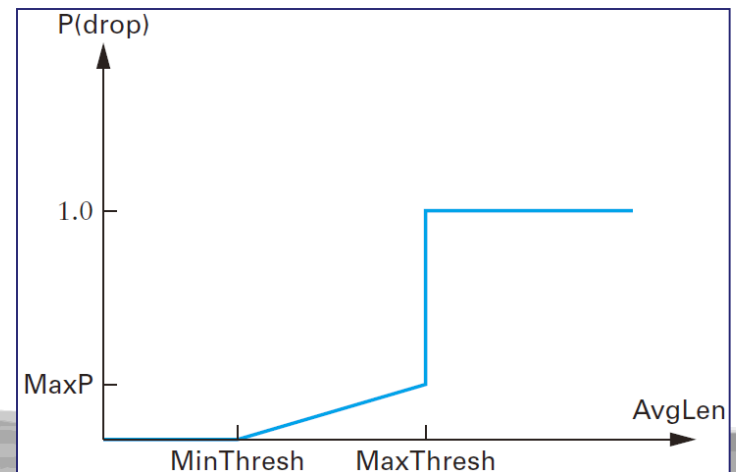
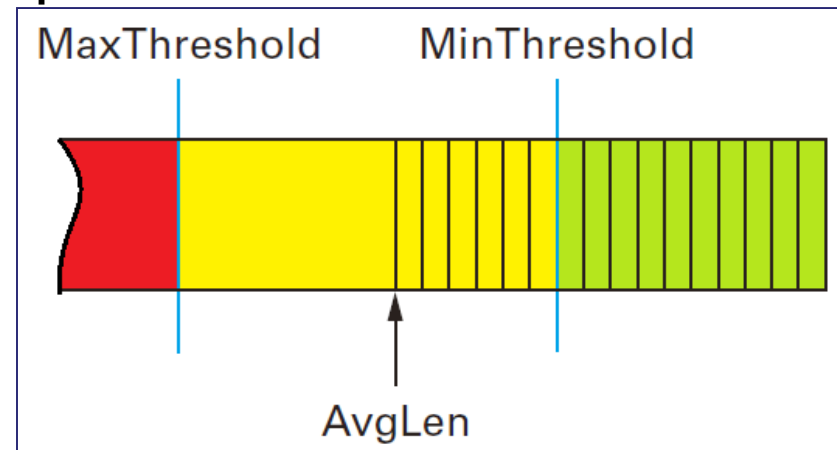
$$\text{avgLen} = (1-\omega) \times \text{avgLen} + \omega \times \text{sampleLen}$$

Calculate average queue size avgLen

if $\text{avgLen} < \text{TH}_{\min}$
queue packet

else if $\text{TH}_{\min} \leq \text{avgLen} < \text{TH}_{\max}$
calculate probability p
with probability p discard packet
else with probability $1-p$ queue packet

else if $\text{avg} \geq \text{TH}_{\max}$
discard packet



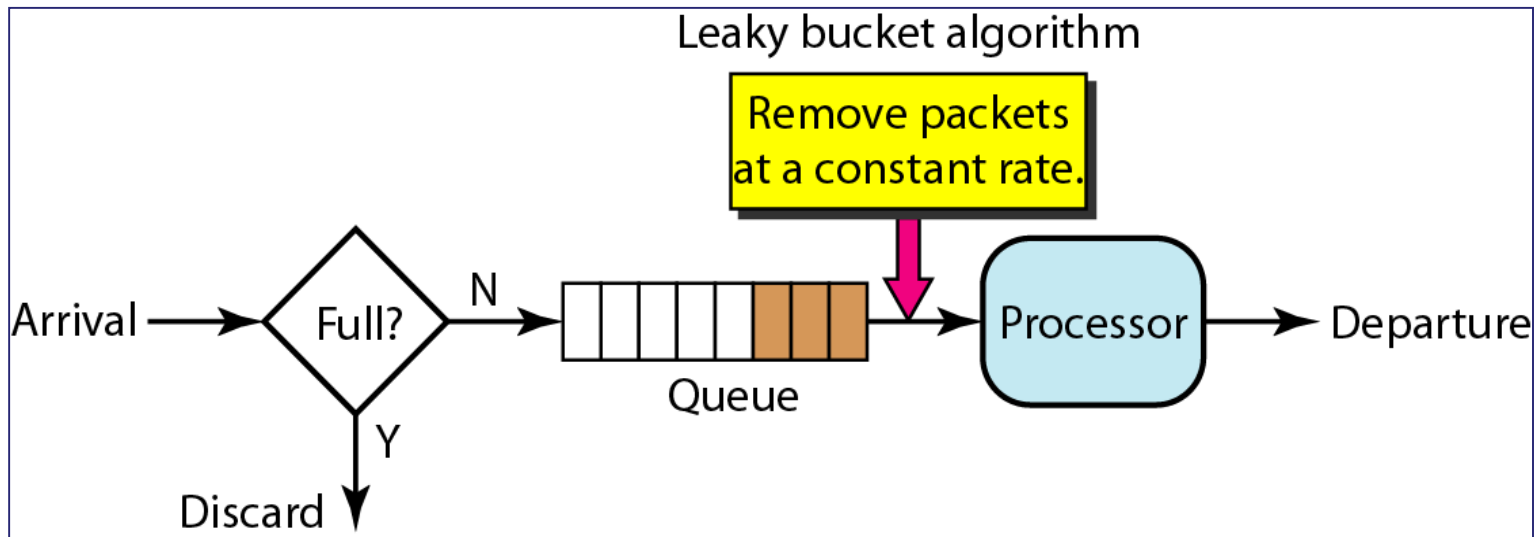


(6) Traffic Shaping

- Shape the traffic (**packet flow**) before it enters the network
 - **Control the rate** at which packets are sent
 - At connection set-up, host and end switch negotiate a traffic pattern (shape)
- **Two traffic shaping algorithms**
 - Leaky Bucket
 - Token Bucket

Leaky Bucket

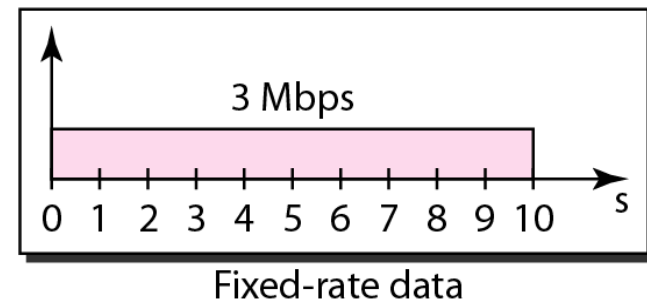
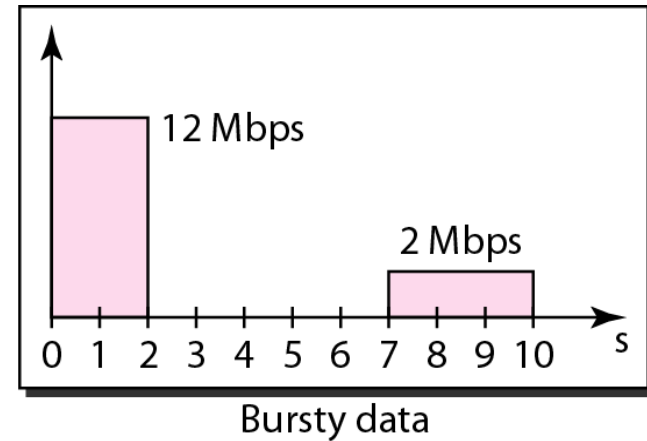
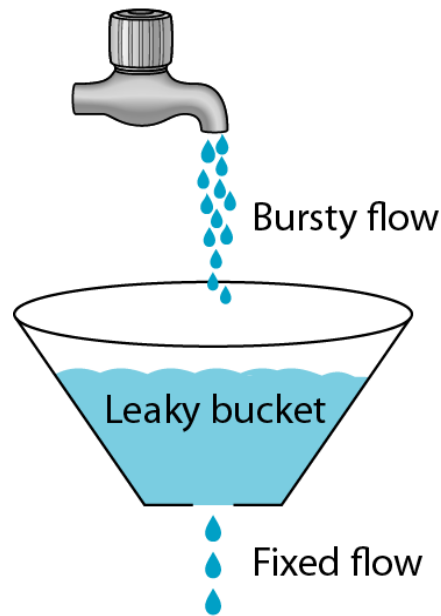
- Shape bursty traffic into **fixed-rate traffic** by averaging the data rate
- May drop the packets if the bucket is full





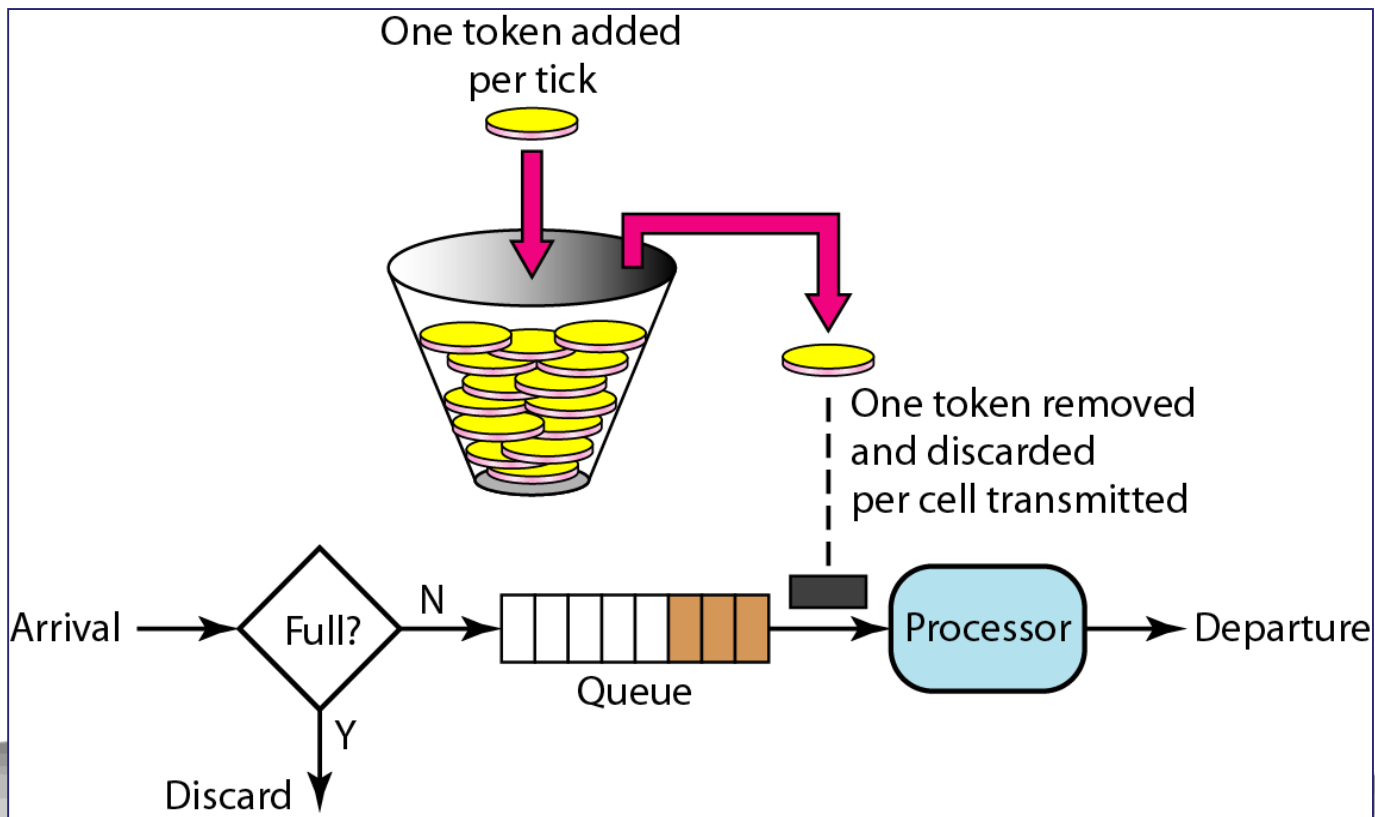
Leaky Bucket

- Do nothing when input is idle
- Packet output rate is **fixed**



Token Bucket

- Use token to control the output traffic, allowing **vary output rate**
- Token generation rate is fixed, may drop token (**not packet**) when bucket full



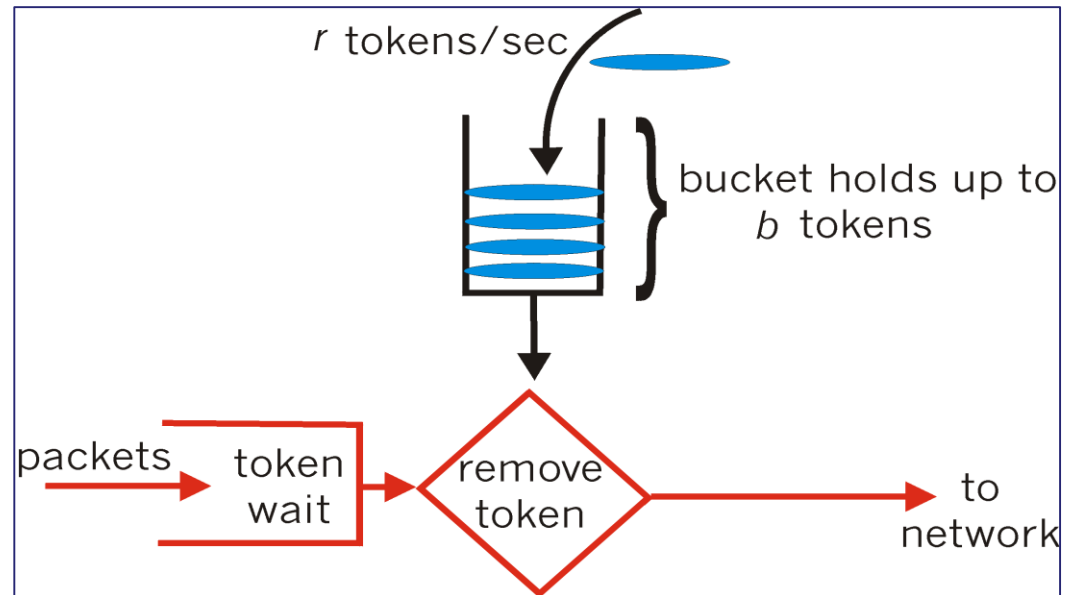


Token Bucket

- Token bucket is more powerful in traffic shaping

3 metrics defined

- Average traffic rate
- Burst traffic rate
- Maximum burst size





Summary

- Mechanisms for Network Congestion Control
 - Choke packet
 - Backpressure
 - Warning bit
 - Congestion window
 - Random early discard
 - Traffic shaping



Homework

- 第7章: R19, P18, P20