CS2109S: Introduction to AI and Machine Learning

Lecture 2: Solving Problems by Searching

25 August 2023

Admin: Assessments

- Midterm (35%)
 - Date/Time: Friday, 6 October (Week 7), 10:00 AM
 - Venue: Multipurpose Sports Hall (MPSH) (to be confirmed)
- Final (35%)
 - Date/Time: Saturday, 18 November, 8:00 PM Sunday, 19 November, 11:59 PM

Admin: Next Week Lecture

Next week (Friday, 1st September) is Polling Day

- There will be **no in-class lecture**
- Lecture will be recorded and published online
- You are expected to watch the lecture before the weekend is over

Admin: Tutorial

Tutorial starts next week!

- Tutorial allocations is available in Coursemology
- Tutorial swaps is allowed (see latest announcement)

Do not appeal through EduRec!

Recap

- What is AI?
- History lesson
- PEAS: Performance Measure, Environment, Actuators, Sensors
- Properties of the task environment
 - Fully observable, deterministic, episodic, static, discrete, single-agent
- Agents
 - Common: reflex, model-based, goal-based, utility-based, learning



Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Problem-Solving Agents

When the correct action to take is not immediately obvious, an agent may need to **plan-ahead**: to consider a *sequence* of actions that form a path to a goal state. Such an agent is called a **problem-solving agent**, and the computational process it undertakes is called **search**.



Want to go to Bucharest, how?

Credit: Passport & Plates

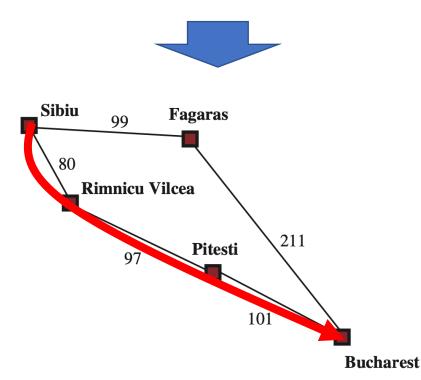
Problem-Solving Process

- Goal formulation
- Problem formulation
 - Create an abstract model of the relevant parts of the world
- Search
 - Simulates sequence of actions using the model, search until goal is reached
- Execution
 - Execute actions in the solution, one at a time

We assume that the task environment is **fully observable** and **deterministic**.

Thus, the solution is a fixed sequence of actions.

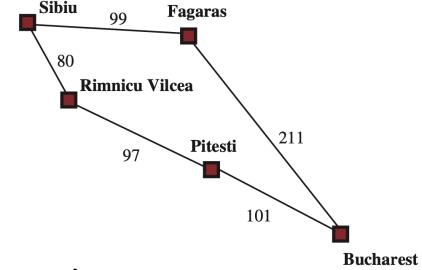




Problem Formulation

Formally, a search problem can be formulated as follows.

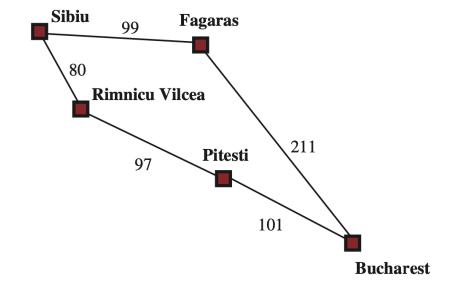
- States (state space).
- Initial state: the initial state of the agent
- Goal state(s)/test
- Actions: things that the agent can do
- Transition model: what each action does
- Action cost function: the cost of performing an action



A sequence of action form a path/trajectory. Solution is a path to a goal.

Problem Formulation: Romania

- States: {at Sibiu, at Fagaras, ..., at Bucharest}
- Initial state: at Sibiu
- Goal state(s)/test: at Bucharest
- Actions: go to neighboring city x
- Transition model: move to target city
- Action cost function: distance



Problem Formulation: Vacuum World

• States: see image

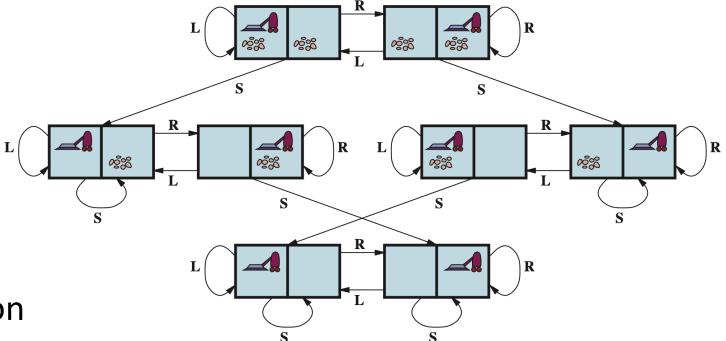
• Initial state: any state

• Goal state(s)/test: all clean

• Actions: suck, left, right

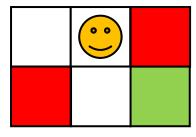
• Transition model: see image

• Action cost function: 1/action



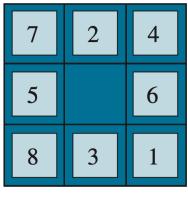
Problem Formulation: Grid World

- States: positions of agent
- **Initial state**: (0,1)
- Goal state(s)/test: (1,2)
- Actions: up, down, left, right
- Transition model: move if there is no wall, die on lava
- Action cost function: -1

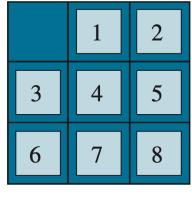


Problem Formulation: 8-Puzzle

- States: location of tiles
- Initial state: see image
- Goal state(s)/test: see image
- Actions: move blank left, right, up, down
- Transition model: move blank
- Action cost function: 1/move



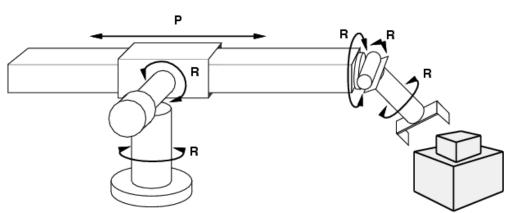




Goal State

Problem Formulation: Robot Assembly

- States: robot joints positions and angles
- Initial state: item not assembled
- Goal state(s)/test: item assembled
- Actions: see image
- Transition model: move/rotate joint
- Action cost function: rotation/displacement



Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Search Algorithms

A search algorithm takes in a search problem as input and returns a solution / failure. It is defined by the **order of node expansion**.

Evaluation criteria:

- Time complexity: number of nodes expanded
- Space complexity: maximum number of nodes in memory
- Completeness: does it return a solution if it exists?
- Optimality: does it always find the least-cost solution?

Measure: branching factor (b), depth (d), maximum depth (m)

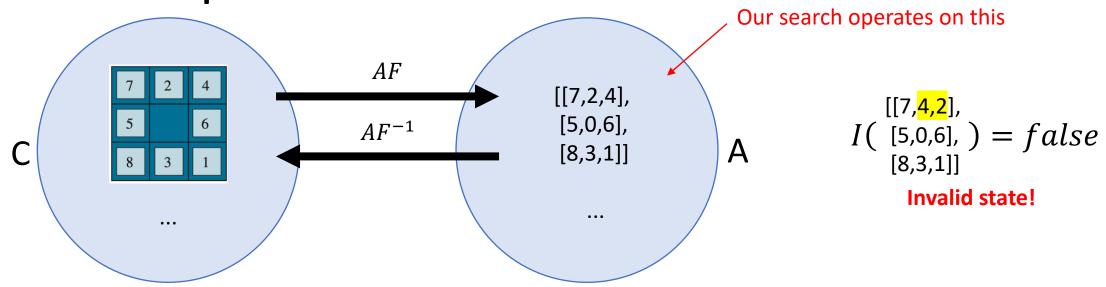
Tree Search

Define order of node expansion

```
create frontier
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```

Representation Invariant

An **abstraction function** AF maps a real-world (concrete) state $c \in C$ to an **abstract representation** $a \in A$.



Representation invariant I is a function such that I(a) = true for all legitimate representations $a \in A$. If I(a) = true, then $AF^{-1}(a) \in C$.

We need to make sure that our transition function satisfies the representation invariant i.e., produce valid states

Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Uninformed Search Algorithms

An uninformed search algorithm is given no clue about how close a state is to the goal(s).

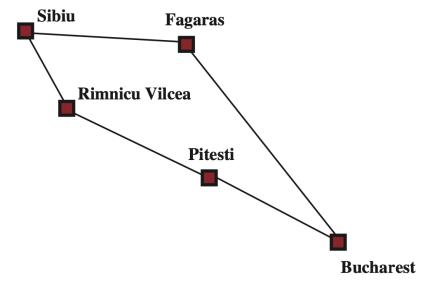
- Breadth-first search (BFS)
- Uniform-cost search
- Depth-first search (DFS)

Condition	Algorithm	Time Complexity
No Negative Weight Cycles	Bellman-Ford Algorithm	O(VE)
On Unweighted Graph (or equal weights)	BFS	O(V+E)
No Negative Weights	Dijkstra's Algorithm	$O((V+E)\log V)$
On Tree	BFS / DFS	O(V)
On DAG	Topological Sort	O(V+E)

CS2040S focus was on finding the shortest path (SSSP, APSP)

CS2109S: might not care about shortest path, or even the path

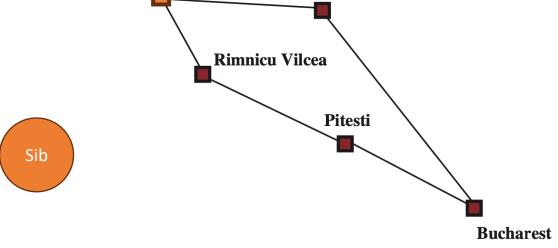
```
create frontier : queue
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```



```
create frontier : queue
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```

Queue:



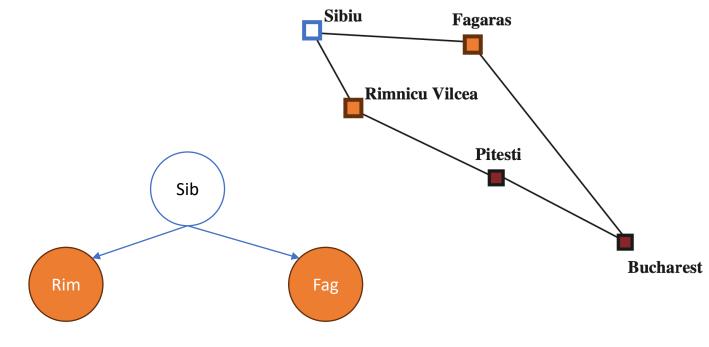


Fagaras

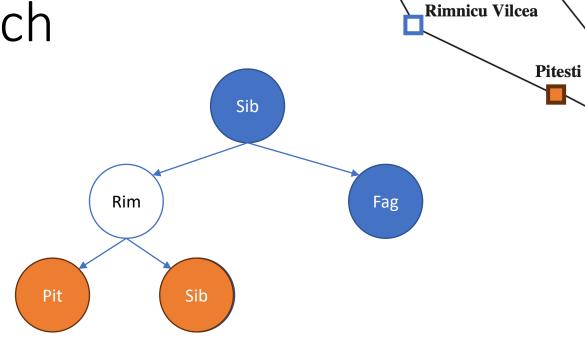
Sibiu

create frontier : queue insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure





create frontier : queue insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure



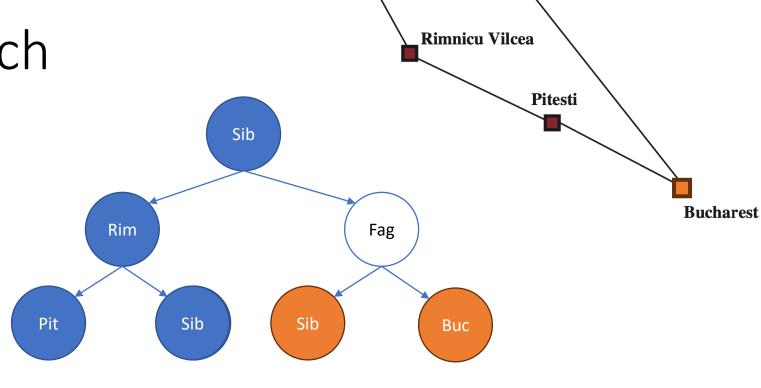
Sibiu

Fagaras

Bucharest



create frontier : queue insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure

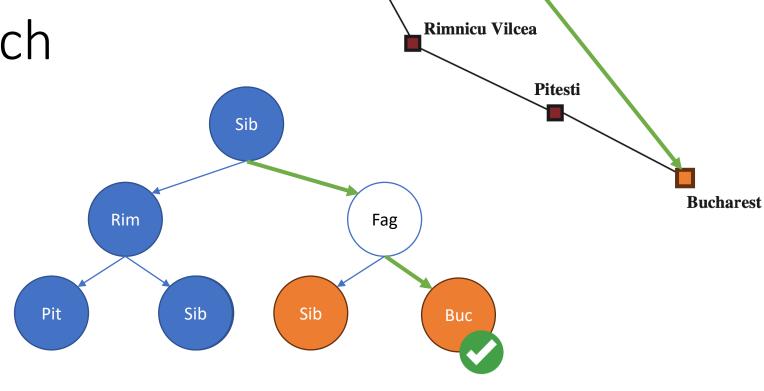


Sibiu

Fagaras



create frontier : queue insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure



Sibiu

Fagaras



create frontier : queue

insert initial state

while **frontier** is not empty:

state = frontier.pop()

for action in actions(state):

next state = transition(state, action)

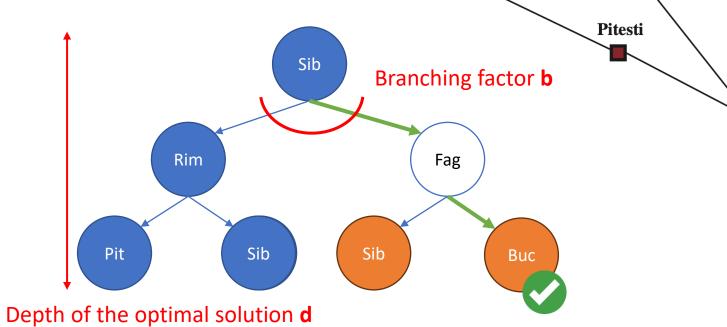
if next state is goal: return solution

frontier.add(next state)

return failure

Queue:





Sibiu

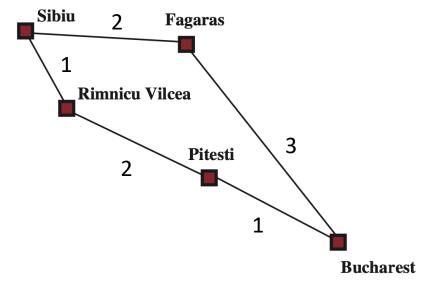
Fagaras

Bucharest

Rimnicu Vilcea

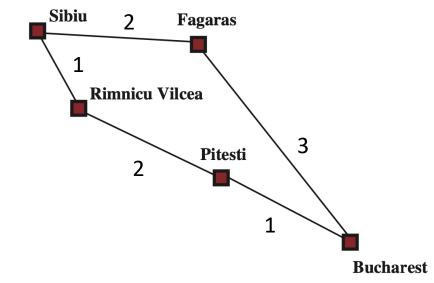
- Time complexity (# nodes expanded)?
 - $1 + b + b^2 + \dots + b^d = O(b^d)$
- Space complexity?
 - $O(b^d)$, worst case: expand the last child in a branch
- Complete? Yes, if B is finite
- Optimal? Yes, if uniform cost

```
create frontier: priority queue (path cost)
insert initial state
while frontier is not empty:
  state = frontier.pop()
  if state is goal: return solution
  for action in actions(state):
    next state = transition(state, action)
    frontier.add(next state)
return failure
```

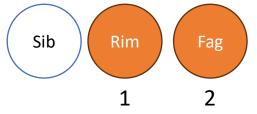


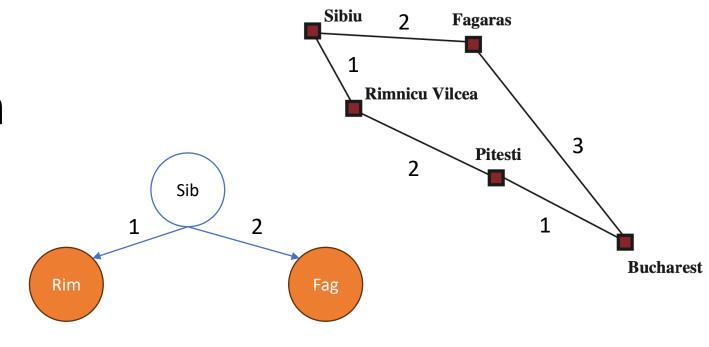
```
create frontier: priority queue (path cost)
insert initial state
while frontier is not empty:
  state = frontier.pop()
  if state is goal: return solution
  for action in actions(state):
    next state = transition(state, action)
    frontier.add(next state)
return failure
```



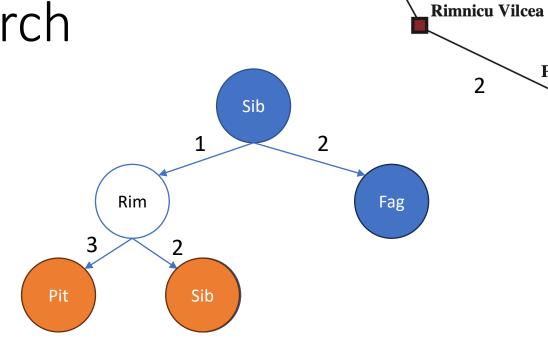


create frontier: priority queue (path cost) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure





create frontier: priority queue (path cost) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure



Sibiu

Fagaras

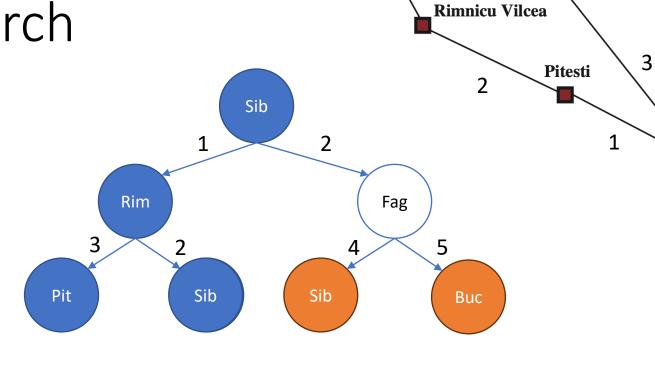
Pitesti

3

Bucharest



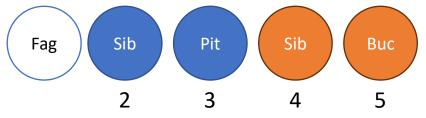
create frontier: priority queue (path cost) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure



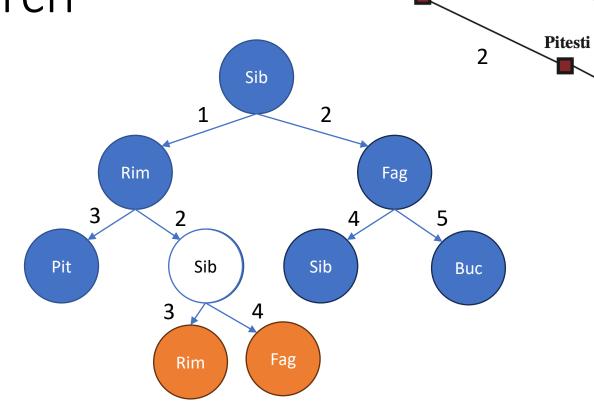
Sibiu

Fagaras

Bucharest



create frontier: priority queue (path cost) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure



Sibiu

Fagaras

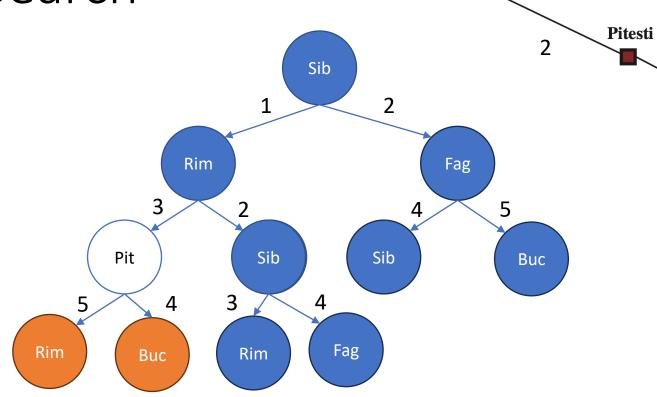
3

Bucharest

Rimnicu Vilcea



create frontier: priority queue (path cost) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure



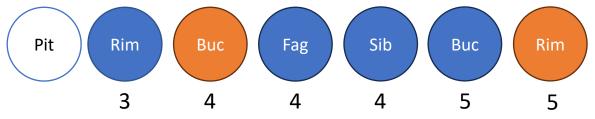
Sibiu

Fagaras

3

Bucharest

Rimnicu Vilcea



create frontier: priority queue (path cost)

insert initial state

while **frontier** is not empty:

state = frontier.pop()

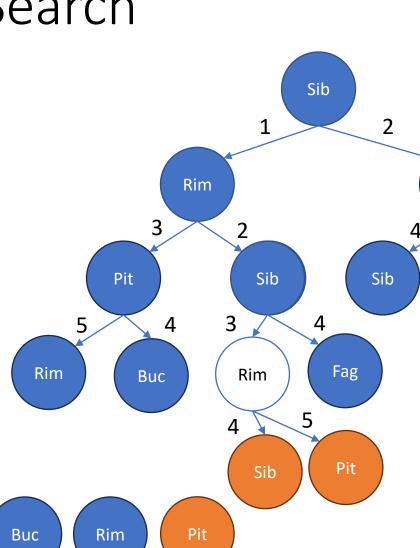
if state is goal: return solution

for action in actions(state):

next state = transition(state, action)

frontier.add(next state)

return failure



Sibiu

Fag

Fagaras

Pitesti

3

Bucharest

Rimnicu Vilcea

Buc

Priority Queue:

Rim Buc Fag Sib Sib Buc Rim Pit
4 4 4 4 5 5 5

Uniform-cost Search

create frontier: priority queue (path cost)

insert initial state

while **frontier** is not empty:

state = frontier.pop()

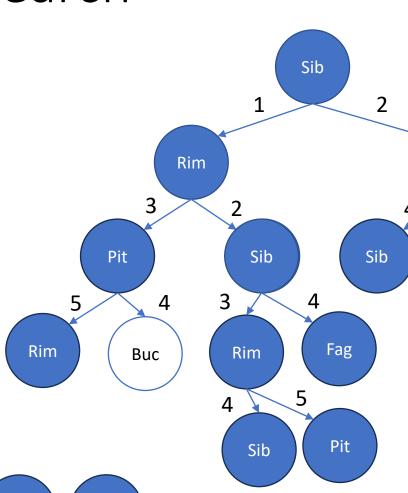
if state is goal: return solution

for action in actions(state):

next state = transition(state, action)

frontier.add(next state)

return failure



Sibiu

Fag

Fagaras

Pitesti

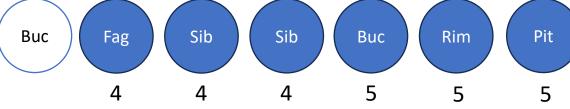
3

Bucharest

Rimnicu Vilcea

Buc

Priority Queue:



Uniform-cost Search

create frontier: priority queue (path cost)

insert initial state

while **frontier** is not empty:

state = frontier.pop()

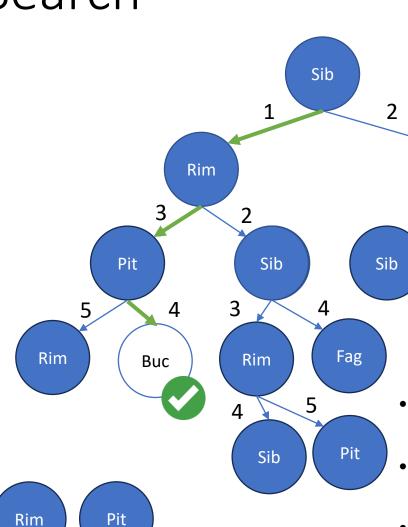
if state is goal: return solution

for action in actions(state):

next state = transition(state, action)

frontier.add(next state)

return failure



Time complexity (# nodes expanded)?

Fagaras

Pitesti

3

Bucharest

Rimnicu Vilcea

- $O(b^{C^*/\epsilon})$, C^* cost of optimal solution
- **Space complexity?**

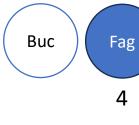
Buc

- $O(b^{C^*/\epsilon})$
- **Complete?** Yes, if step cost $\geq \epsilon$
- **Optimal?** Yes

Sibiu

Fag













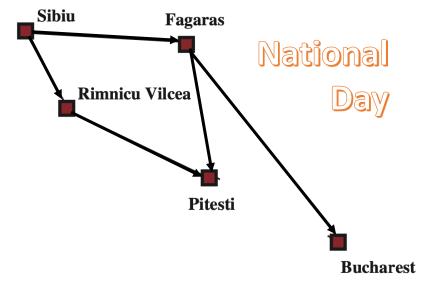
Sib







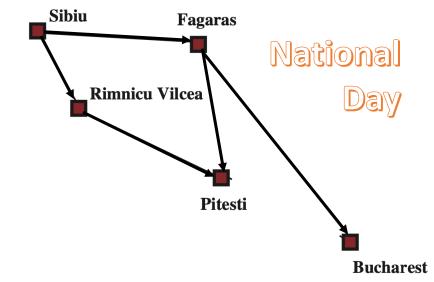
```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```



```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```

Stack:

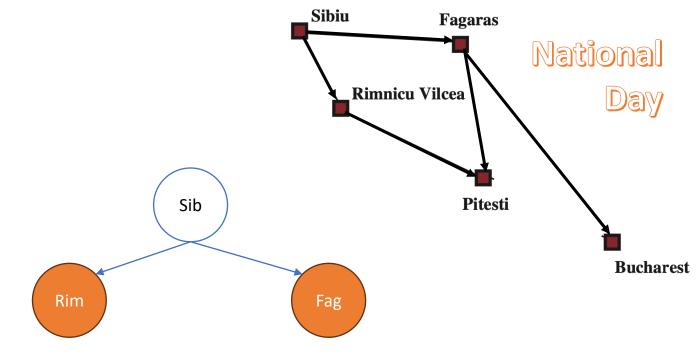


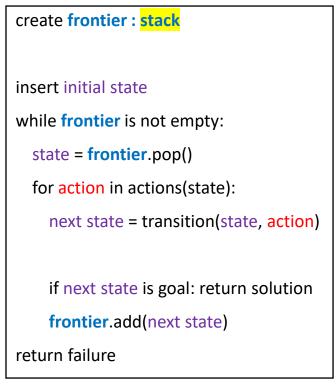


```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```

Stack:

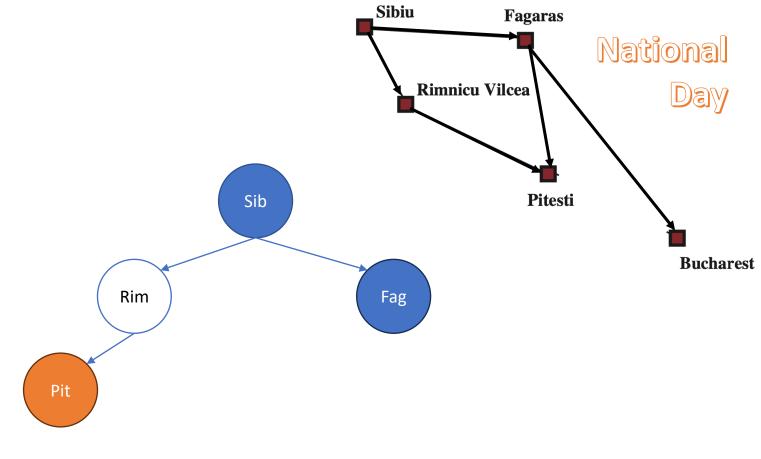




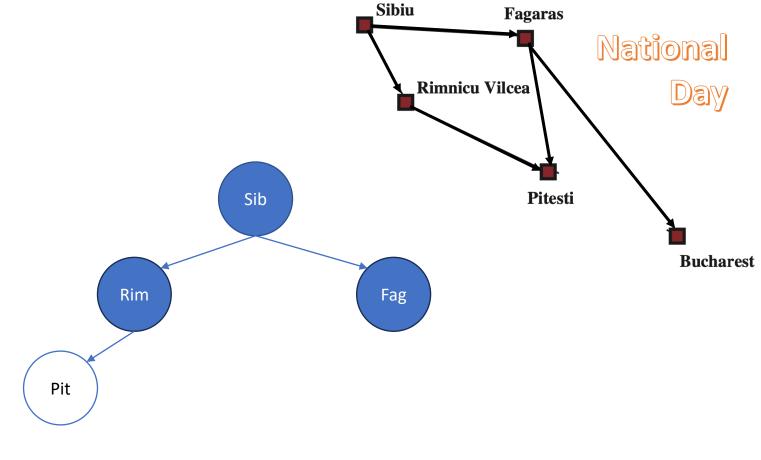








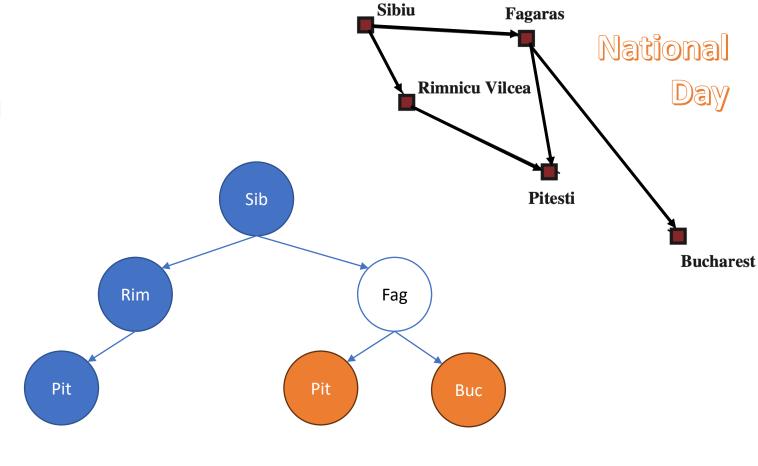
```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```



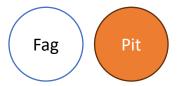
Stack:



create frontier: stack insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure



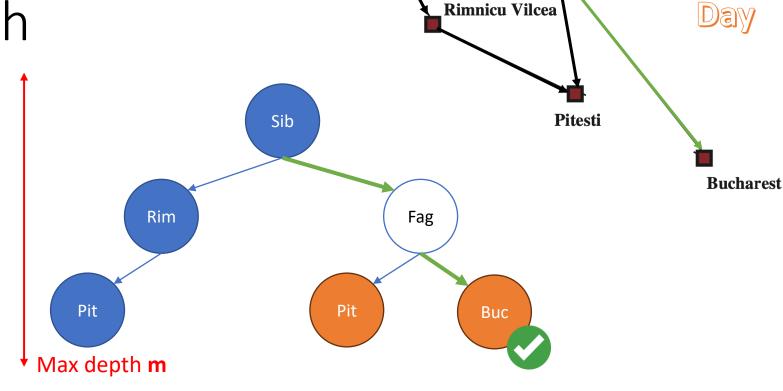
Stack:



create frontier: stack insert initial state while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state is goal: return solution frontier.add(next state) return failure

Stack:





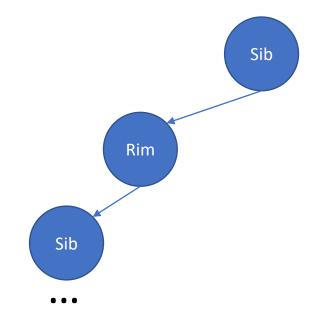
Sibiu

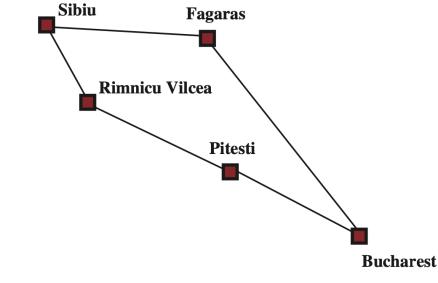
Fagaras

National

- Time complexity (# nodes expanded)?
 - $O(b^m)$
- Space complexity?
 - 0(bm)
- Complete?

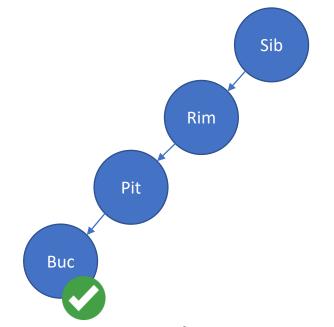
```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```





- Time complexity (# nodes expanded)?
 - $O(b^m)$
- Space complexity?
 - O(bm)
- Complete? No, when depth is infinite or can go back and forth (loops)
- Optimal?

```
create frontier: stack
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```



- Time complexity (# nodes expanded)?
 - $O(b^m)$
- Space complexity?
 - 0(bm)
- Complete? No, when depth is infinite or can go back and forth (loops)

Sibiu

Fagaras

Pitesti

Rimnicu Vilcea

Can be shorter!

Bucharest

Optimal? No

How do we handle infinite depth?

Outline

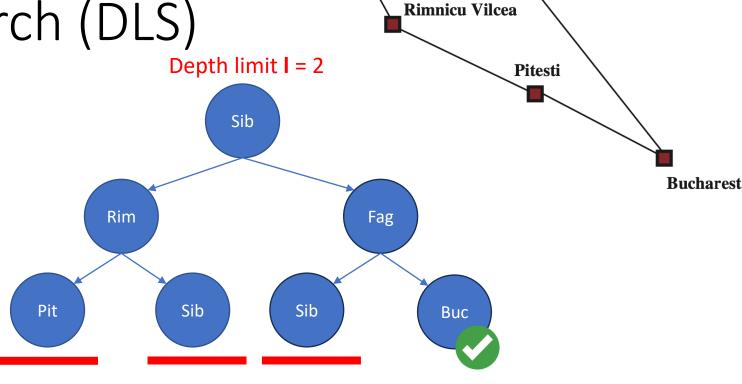
- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)

Variants of uninformed search algorithms

- Depth-limited search
- Iterative deepening search
- Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Depth-limited Search (DLS)

- Limit the search depth
- Backtrack once the depth limit is reached



Sibiu

Fagaras

How do we know the depth of the solution?

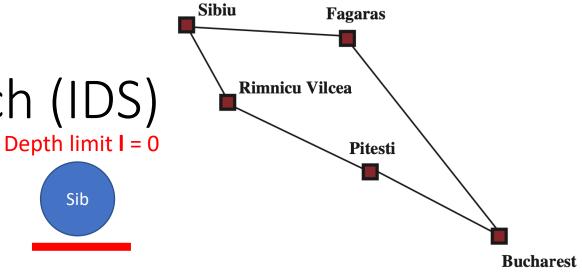
We don't know:

• Time complexity (# nodes expanded)?

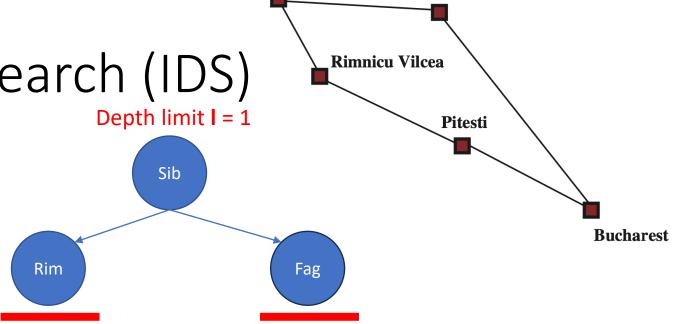
•
$$b^0 + b^1 + \dots + b^l = O(b^l)$$

- Space complexity?
 - 0(bl)
- Complete? No
- Optimal? No

- Do depth-limited search with max depth 0 ... N
- Return solution if found
- Increase the depth otherwise



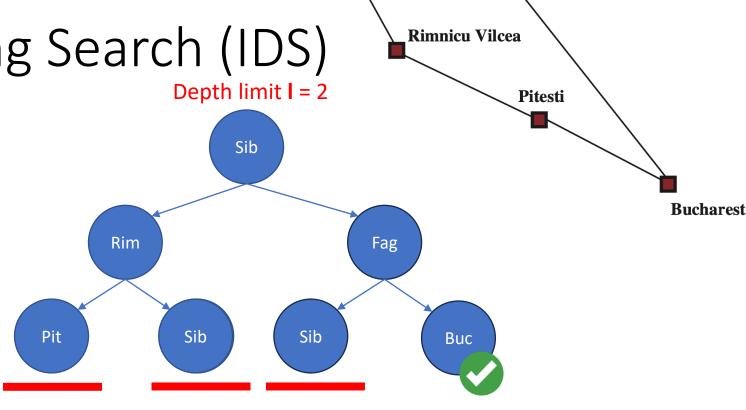
- Do depth-limited search with max depth 0 ... N
- Return solution if found
- Increase the depth otherwise



Sibiu

Fagaras

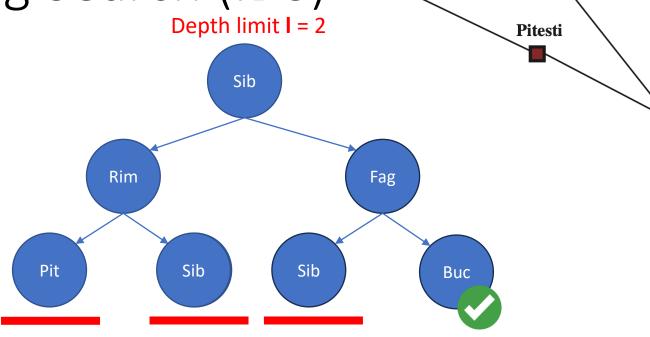
- Do depth-limited search with max depth 0 ... N
- Return solution if found
- Increase the depth otherwise



Sibiu

Fagaras

- Do depth-limited search with max depth 0 ... N
- Return solution if found
- Increase the depth otherwise



Sibiu

Fagaras

Bucharest

Rimnicu Vilcea

Time complexity (# nodes expanded)?

Overhead! •
$$b^0 + (b^0 + b^1) + \dots + (b^0 + \dots + b^d)$$

• $= (d+1)b^0 + db^1 + (d-1)b^2 + \dots + 2b^{d-1} + b^d = O(b^d)$

- Space complexity?
 - 0(bd)
- Complete? Yes
- **Optimal?** Yes, if uniform step cost

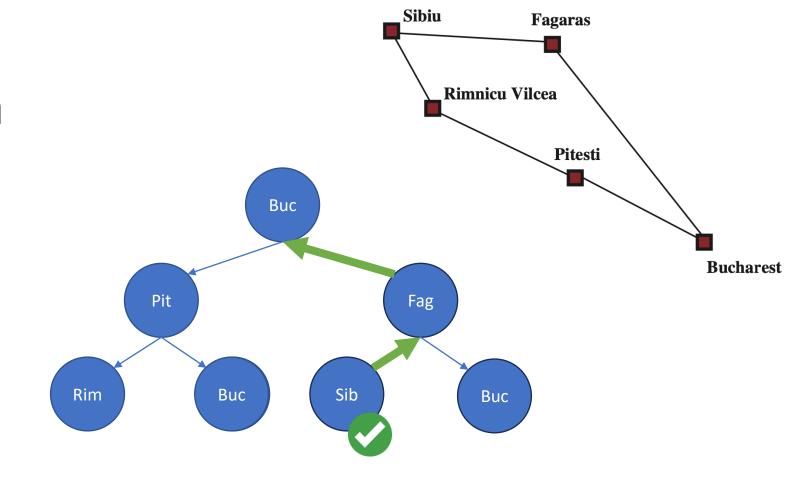
For b=10, d=5,

- N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111
- $N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$

Overhead = (123,456 - 111,111)/111,111 = 11%

Backward Search

Search from the goal



Bidirectional Search

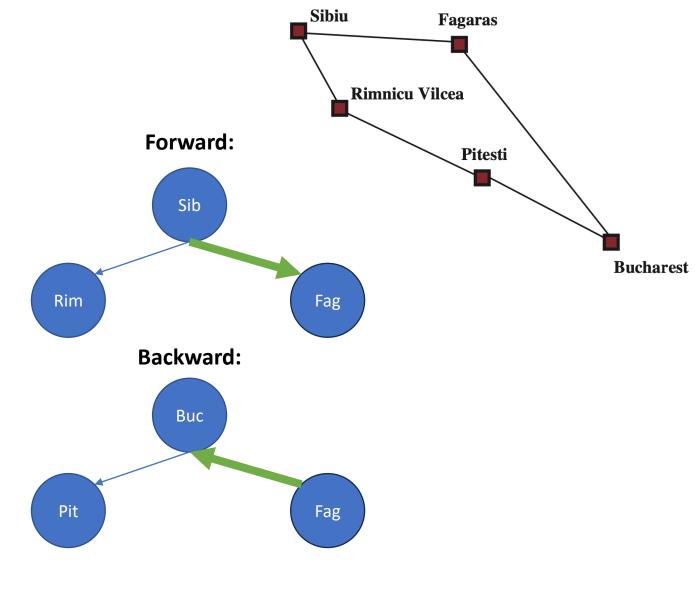
Combine search:

- Forward (from the start)
- Backward (from the goal)

Stop when two searches meet

Intuition:

$$2 \times O(b^{d/2}) < O(b^d)$$



Issues?

- Operators need to be reversible
- Many goal states

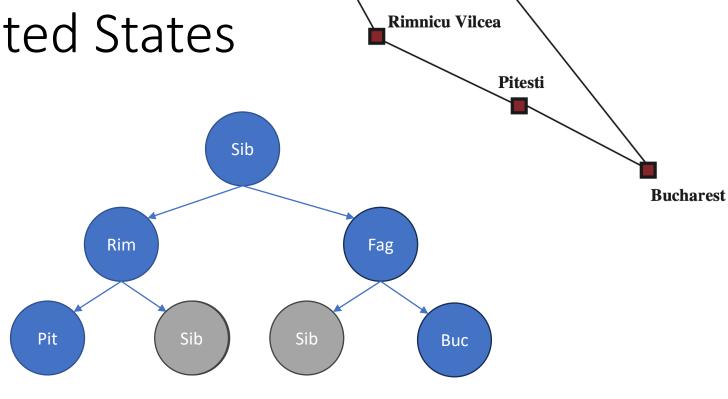
How to efficiently check if a node appears in the other search tree?

Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Dealing with Repeated States

- Remember states that are already visited
- Don't visit again



Sibiu

Fagaras

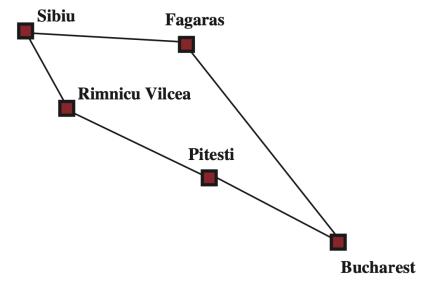
(use hashtable/dictionary!)

Tree Search Keep track of visited nodes Graph Search

```
create frontier
insert initial state
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
    if next state is goal: return solution
    frontier.add(next state)
return failure
```

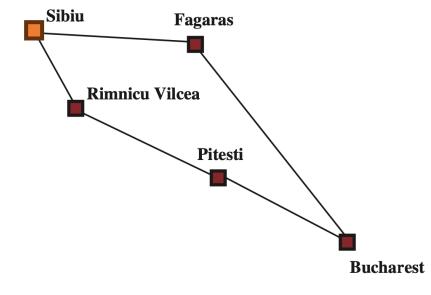
```
create frontier
create visited
insert initial state to queue and visited
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
     if next state in visited: continue
    if next state is goal: return solution
    frontier.add(next state)
    visited.add(next state)
return failure
```

```
create frontier : queue
create visited
insert initial state to queue and visited
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
     if next state in visited: continue
    if next state is goal: return solution
    frontier.add(next state)
    visited.add(next state)
return failure
```



create frontier : queue create visited insert initial state to queue and visited while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state in visited: continue if next state is goal: return solution frontier.add(next state) visited.add(next state) return failure



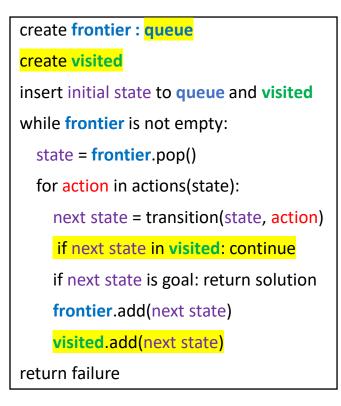


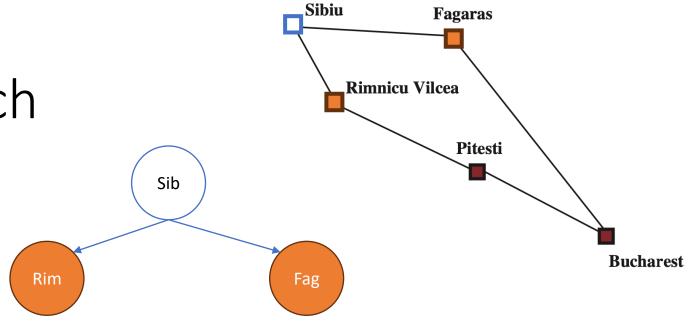
Queue:

Sib

Visited:

Sib





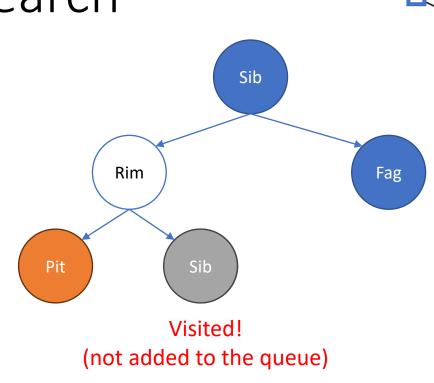
Queue:



Visited:







Queue:







Sibiu

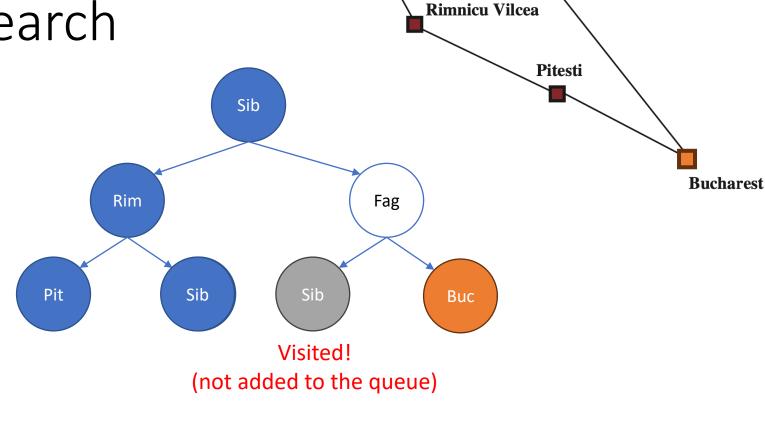
Fagaras

Pitesti

Bucharest

Rimnicu Vilcea

create frontier : queue create visited insert initial state to queue and visited while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state in visited: continue if next state is goal: return solution frontier.add(next state) visited.add(next state) return failure



Sibiu

Fagaras

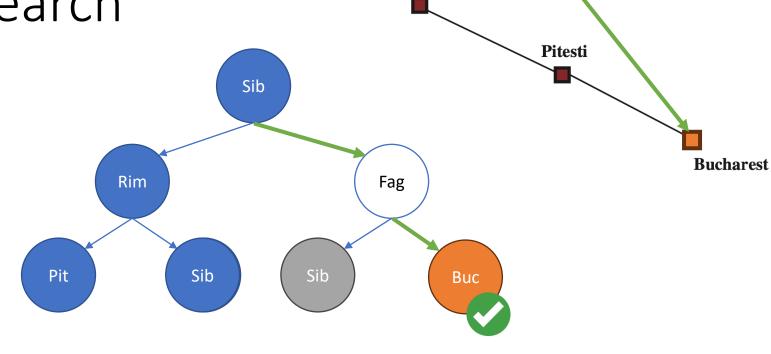
Queue:



Visited:



create frontier : queue create visited insert initial state to queue and visited while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state in visited: continue if next state is goal: return solution frontier.add(next state) visited.add(next state) return failure



Sibiu

Fagaras

Rimnicu Vilcea

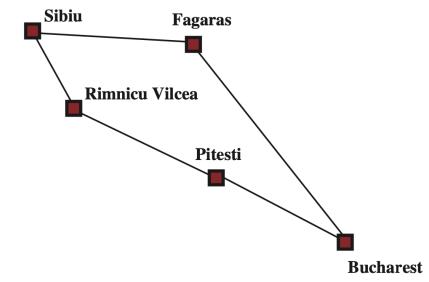
Queue:



Visited:

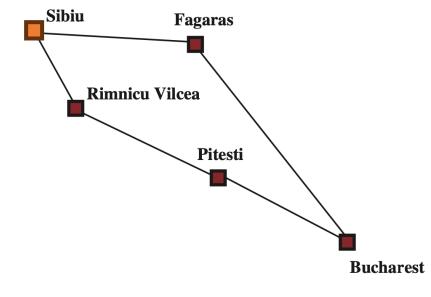


```
create frontier : stack
create visited
insert initial state to queue and visited
while frontier is not empty:
  state = frontier.pop()
  for action in actions(state):
    next state = transition(state, action)
     if next state in visited: continue
    if next state is goal: return solution
    frontier.add(next state)
    visited.add(next state)
return failure
```



create frontier: stack create visited insert initial state to queue and visited while **frontier** is not empty: state = frontier.pop() for action in actions(state): next state = transition(state, action) if next state in visited: continue if next state is goal: return solution frontier.add(next state) visited.add(next state) return failure



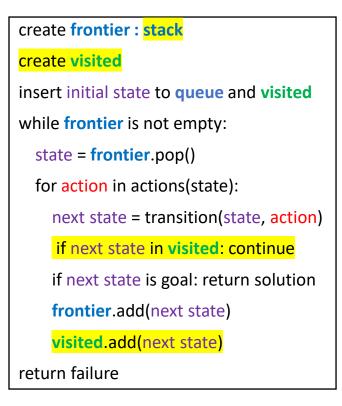


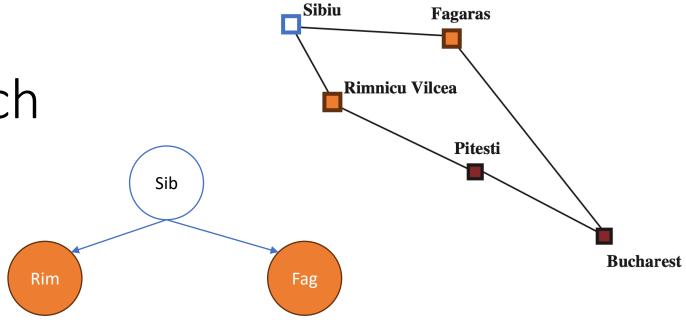
Stack:

Sib

Visited:

Sib



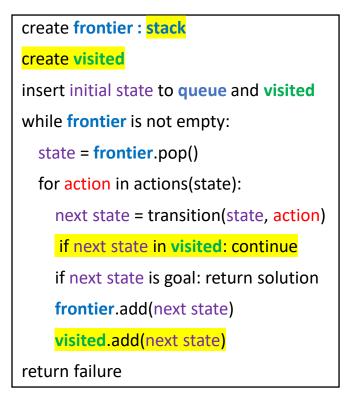


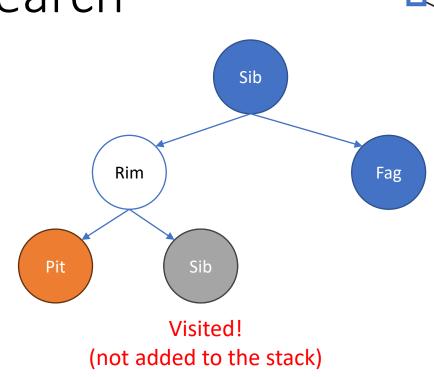
Stack:



Visited:







Stack:



Visited:



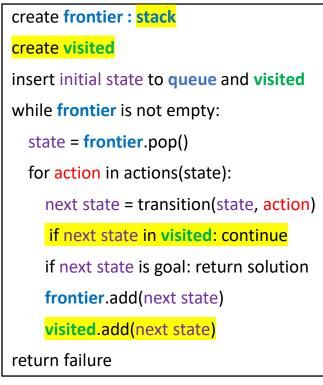
Sibiu

Fagaras

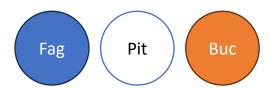
Pitesti

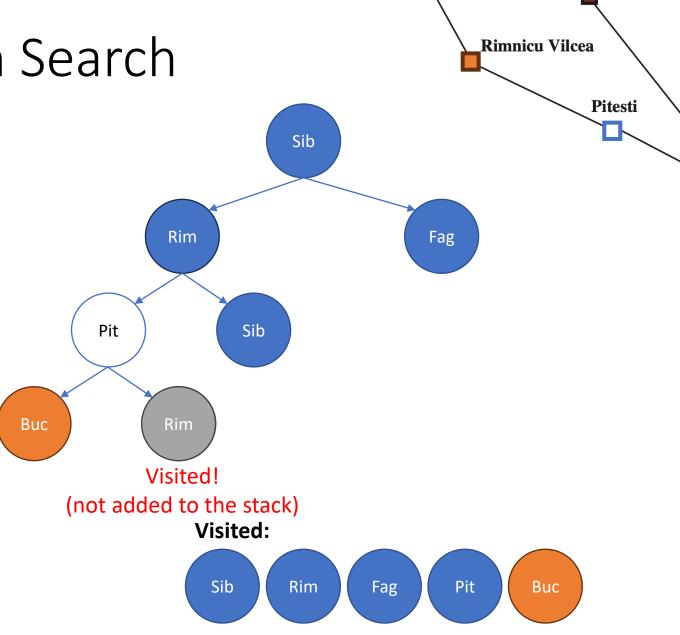
Bucharest

Rimnicu Vilcea



Stack:

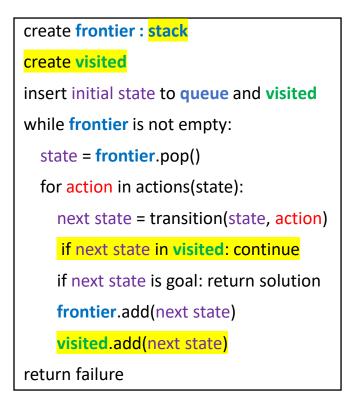


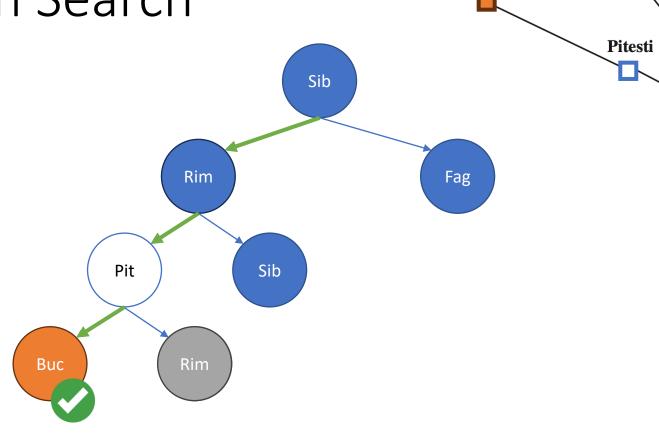


Sibiu

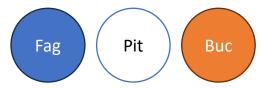
Fagaras

Bucharest





Stack:



Visited:



Sibiu

Fagaras

Bucharest

Rimnicu Vilcea

Choosing a Search Strategy

Depends on the problem:

- Number of goal states
- Distribution of goal states in search tree
- Finite/infinite branching factor/depth
- Repeated states
- Need for optimality?
- Need to know if there is no solution?

Summary: Uninformed Search Algorithms

Search algorithms

- Breadth-first search queue, explore layer by layer
- Uniform-cost search priority queue (path cost)
- Depth-first search **stack**, go deep first then backtrack

Variants:

- Depth limited search limit max depth of the search
- Iterative deepening search try DLS with depth limit 0, ..., N
- Bidirectional search search from the **start** and the **goal**, meet in the **middle**

Dealing with repeated states

Graph search – don't visit nodes that are already visited

Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

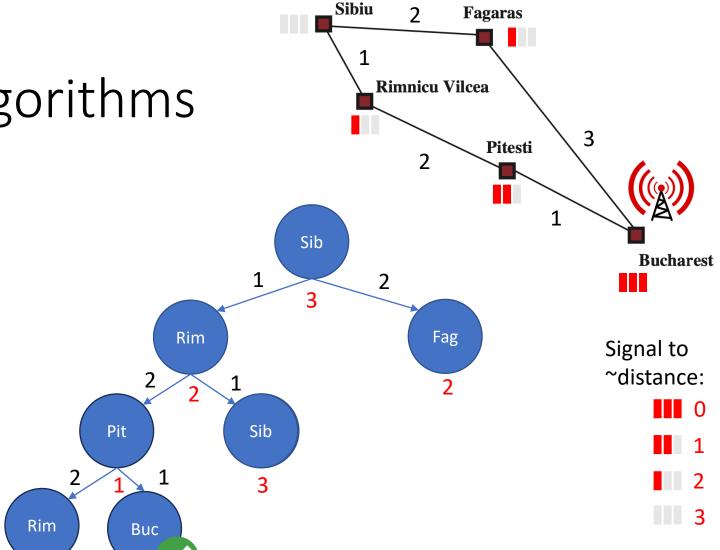
Informed Search Algorithms

Uninformed search:

Search blindly

Informed search:

Use domain information to guide the search

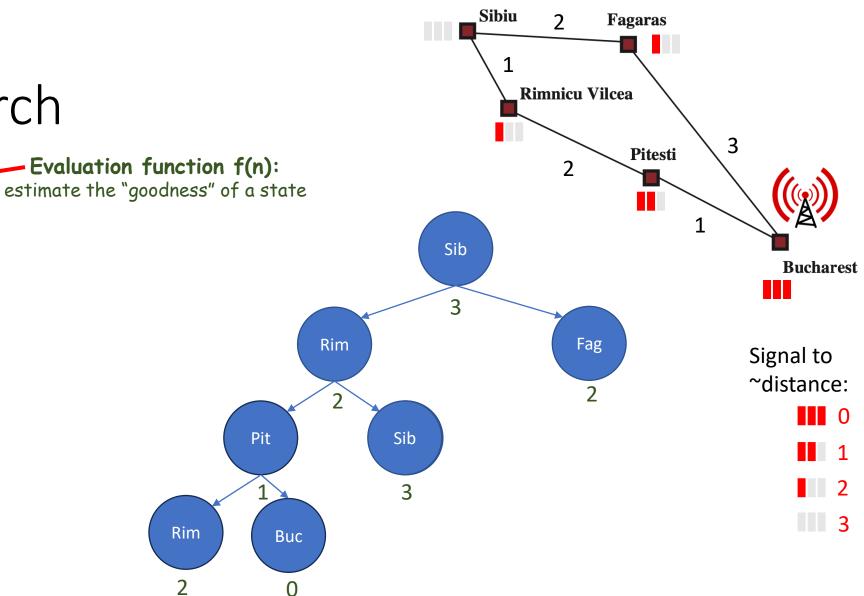


Best-first Search

create frontier : priority queue f(n) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure

Special cases:

- Greedy best-first search
- A* search



Greedy Best-first Search

create frontier: priority queue f(n)

insert initial state

while **frontier** is not empty:

state = frontier.pop()

if state is goal: return solution

for action in actions(state):

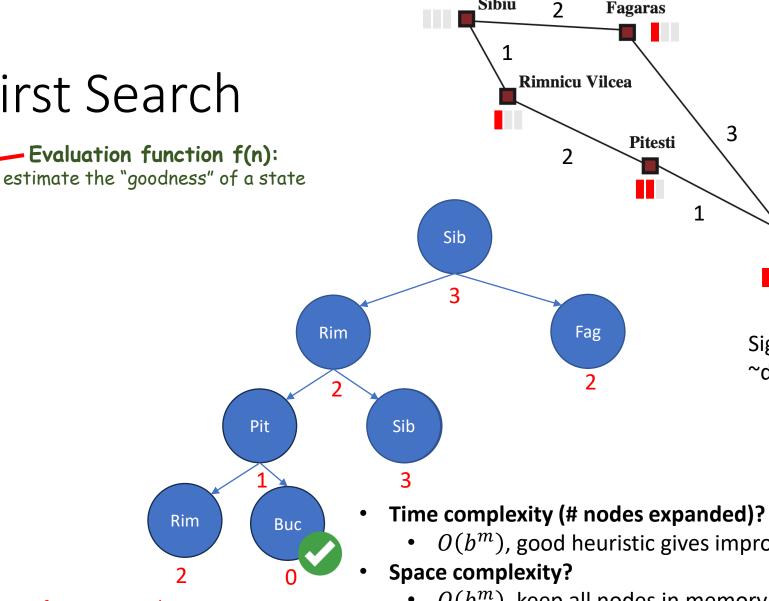
next state = transition(state, action)

frontier.add(next state)

return failure

f(n) = h(n)

Heuristic: estimated cost from n to goal



Sibiu

 $O(b^m)$, good heuristic gives improvement

Bucharest

Signal to

~distance:

0

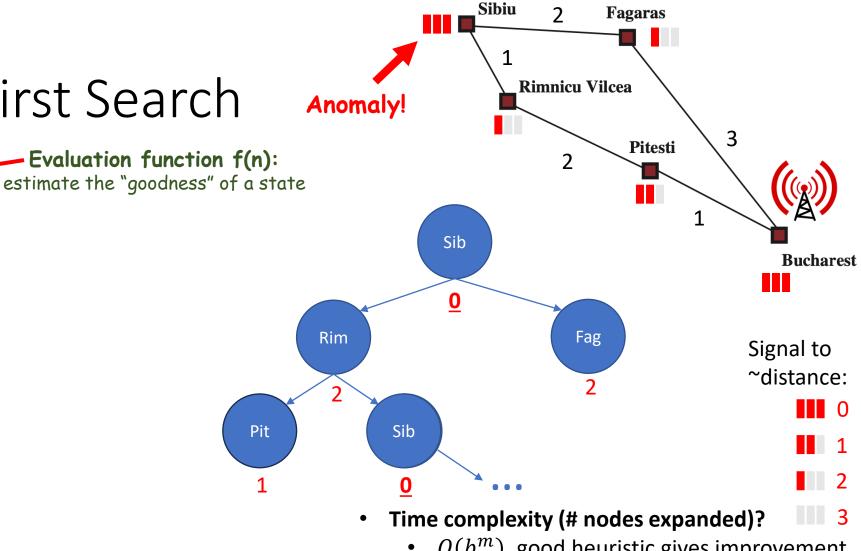
- $O(b^m)$, keep all nodes in memory
- **Complete?**

Greedy Best-first Search

create frontier: priority queue f(n) insert initial state while **frontier** is not empty: state = frontier.pop() if state is goal: return solution for action in actions(state): next state = transition(state, action) frontier.add(next state) return failure

f(n) = h(n)Heuristic: estimated cost from n to goal

Doesn't consider the cost so far!



- $O(b^m)$, good heuristic gives improvement
- **Space complexity?**
 - $O(b^m)$, keep all nodes in memory
- **Complete? No**
- **Optimal? No**

A* Search

create frontier: priority queue f(n)

insert initial state

while **frontier** is not empty:

state = frontier.pop()

if state is goal: return solution

for action in actions(state):

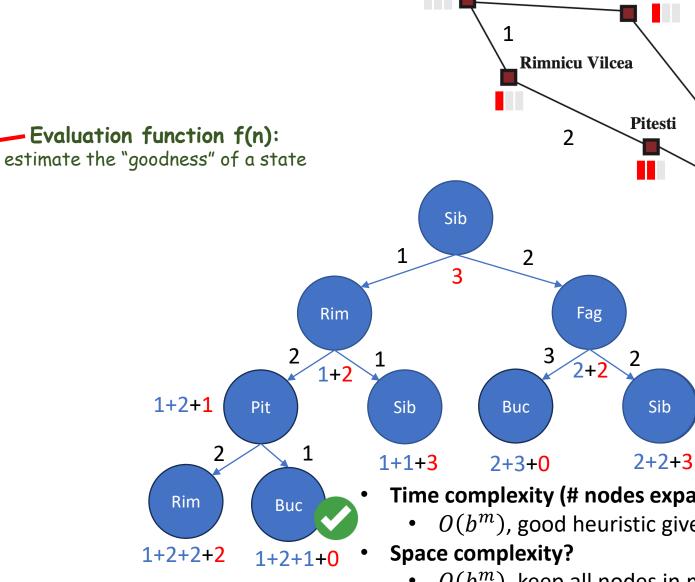
next state = transition(state, action)

frontier.add(next state)

return failure

$$f(n) = g(n) + h(n)$$

Cost so far to reach n Heuristic: estimated cost from n to goal



Sibiu

- Time complexity (# nodes expanded)?
 - $O(b^m)$, good heuristic gives improvement

Fagaras

Bucharest

Signal to

~distance:

0

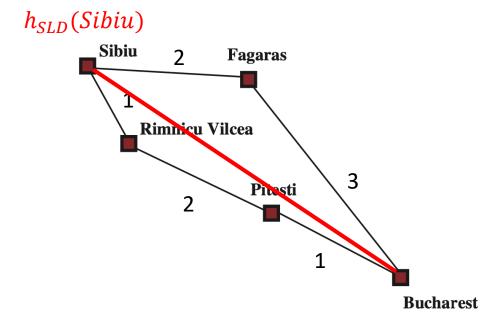
- $O(b^m)$, keep all nodes in memory
- **Complete?** Yes
- Optimal? Yes*

Admissible Heuristics

A heuristic h(n) is admissible if for every node n, h(n) \leq h*(n), where h*(n) is the **true cost** to reach the goal state from n.

An admissible heuristic **never over- estimates** the cost to reach the goal, i.e., it is a **conservative estimate**.

Theorem: if h(n) is admissible, A* using tree search is optimal



Example: $h_{SLD}(n)$ never overestimates

the actual road distance

Admissible Heuristics: Optimality

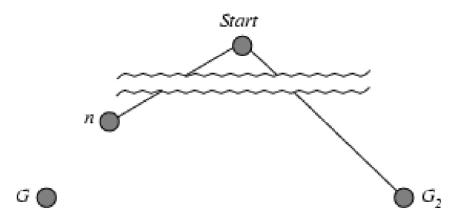
Suppose some suboptimal goal G_2 has been generated and is in the fringe. Let n be an unexpanded node in the fringe such that n is on a shortest path to an optimal goal G.

$$f(G) = g(G) + h(G) = g(G) + 0$$

 $f(G_2) = g(G_2) + h(G_2) = g(G_2) + 0$
 $g(G_2) > g(G)$ since G_2 is suboptimal
 $f(G_2) > f(G)$

$$h(n) \le h^*(n)$$
 since h is admissible $f(n) = g(n) + h(n) \le g(n) + h^*(n) = f(G)$

$$f(n) \le f(G) < f(G_2)$$
, G_2 will never be expanded



Consistent Heuristics

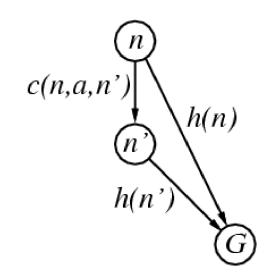
A heuristic h(n) is consistent if for every node n, every successor n' of n generated by any action a, $h(n) \le c(n,a,n') + h(n')$

If h is consistent, we have

$$f(n') = g(n') + h(n')$$

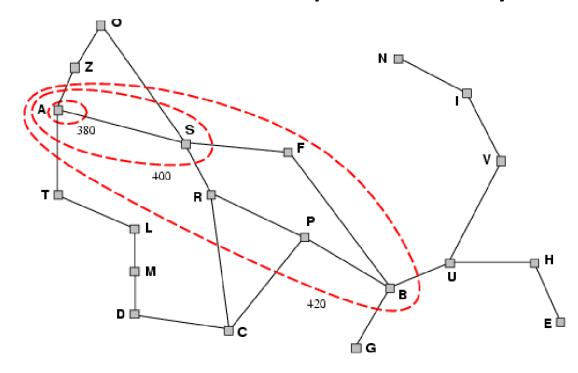
= $g(n) + c(n,a,n') + h(n')$
 $\geq g(n) + h(n) = f(n)$

i.e., f(n) is **non-decreasing** along any path



Theorem: If h(n) is consistent, A* using graph search is optimal

Consistent Heuristics: Optimality



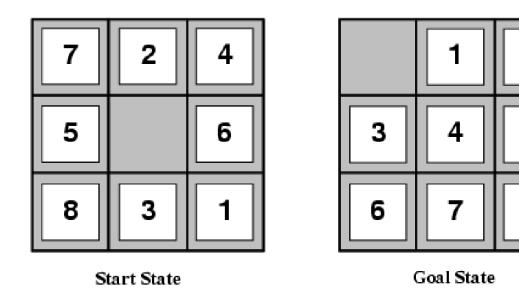
A * expands nodes in order of increasing f-cost

Gradually adds "f-contours" of nodes

Contour i has all nodes with $f = f^{(i)}$, where $f^{(i)} < f^{(i+1)}$

Dominance

If $h_2(n) \ge h_1(n)$ for all n (both admissible) then h_2 dominates h_1 . h_2 is better for search.



Heuristics

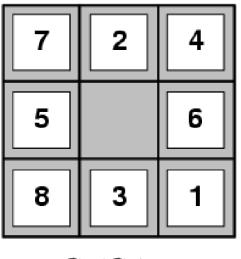
- h_1 number of misplaces tiles
- h₂ total Manhattan distance

 h_2 dominates h_1

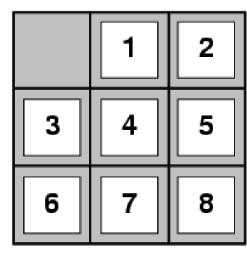
If each tile is at most one distance away from the goal, then $h_2 = h_1$, otherwise $h_2 > h_1$

"Inventing" Admissible Heuristics

A problem with **fewer restrictions** on the actions is called a <u>relaxed</u> <u>problem</u>. The cost of an <u>optimal solution</u> to a relaxed problem is an **admissible heuristic** for the original problem.







Goal State

Original:

A tile can only move to adjacent blank

Relaxations:

- Each tile can move anywhere
 - h_1 number of misplaces tiles
- Each tile can move to any adjacent square
 - h₂ total Manhattan distance

Outline

- Problem-solving agents
- Search algorithms
- Uninformed search algorithms
 - Breadth-first Search (BFS)
 - Uniform-cost search
 - Depth-first Search (DFS)
- Variants of uninformed search algorithms
 - Depth-limited search
 - Iterative deepening search
 - Bidirectional search
- Dealing with repeated states
- Informed search algorithms
 - Greedy best-first search
 - A* search
 - Heuristics
- Variants of A*

Variants of A*

- Iterative Deepening A* (IDA*)
 - Use iterative deepening search
 - Cutoff using f-cost [f(n) = g(n) + h(n)] instead of depth
- Simplified Memory-bounded A* (SMA*)
 - Drop the nodes with worst f-cost if memory is full

Summary: Informed Search Algorithms

- Informed search: guide search with domain information
- Best-first search
 - Greedy best-first search
 - f(n) = h(n), heuristic estimate of cost from n to goal
 - A* search
 - f(n) = g(n) + h(n), cost so far + heuristic
- Heuristics
 - Admissible: $h(n) \le h^*(n)$
 - Consistent: $h(n) \le c(n,a,n') + h(n')$
 - Dominant: if $h1(n) \le h2(n)$, h2 dominant
- Creating admissible heuristic: true cost of the relaxed problem
- A* variants: IDA*, SMA*
 - Idea: prune to save memory

Summary

- Problem-solving agents
- Uninformed search algorithms
 - Breadth-first Search (BFS) layer by layer
 - Uniform-cost search Djikstra
 - Depth-first Search (DFS) go deep first
- Variants of uninformed search algorithms
 - Depth-limited search set max depth
 - Iterative deepening search try DLS with depth limit 0, ..., N
 - Bidirectional search combine forward and backward search
- Dealing with repeated states visit state only once
- Informed search algorithms
 - Greedy best-first search f(n) = h(n)
 - A^* search f(n) = g(n) + h(n)
 - Heuristics: admissibility, consistency, dominance
- Variants of A*: IDA*, SMA*

Coming Up Next Week

- Local search
 - Hill climbing
 - Simulated annealing
 - Beam search
 - Genetic algorithms
- Adversarial search
 - Games vs search problems
 - Minimax
 - Alpha-beta pruning

To Do

- Lecture Training 2
 - +100 Free EXP
 - +50 Early bird bonus
- Problem Set 0
 - Due Saturday, 26th August (Tomorrow)
- Tutorial Swaps
 - Due Sunday, 27th August