

Explanation of the Shortest Remaining Time First (SRTF) scheduling algorithm implementation:

1. Process Class:

```
```java
class Processss {
 int id;
 int arrivalTime;
 int burstTime;
 int remainingTime;

 Processss(int id, int arrivalTime, int burstTime) {
 this.id = id;
 this.arrivalTime = arrivalTime;
 this.burstTime = burstTime;
 this.remainingTime = burstTime;
 }
}
```
```

This class represents a process with an ID, arrival time, burst time, and remaining time. The remaining time is initially set to the burst time and decreases as the process is executed.

2. Main Method and Process Generation:

```
```java
public static void main(String[] args) {
 List<Processss> processsses = new ArrayList<>();
 Random rand = new Random(0); // For reproducibility

 for (int i = 0; i < 79; i++) {
 int arrivalTime = rand.nextInt(11); // Random arrival time between
0 and 10
 int burstTime = rand.nextInt(10) + 1; // Random burst time between
1 and 10
 processsses.add(new Processss(i, arrivalTime, burstTime));
 }

 srtfScheduling(processsses);
}
```
```

The main method generates 79 processes with random arrival times (0-10) and burst times (1-10). It then calls the SRTF scheduling method.

3. SRTF Scheduling Method:

```
```java
public static void srtfScheduling(List<Process> processes) {
 processes.sort(Comparator.comparingInt(p -> p.arrivalTime));
 // ... (initialization of variables)

 while (complete != n) {
 // Find process with shortest remaining time
 for (int j = 0; j < n; j++) {
 if ((processes.get(j).arrivalTime <= currentTime) &&
 (processes.get(j).remainingTime < minRemainingTime) &&
 processes.get(j).remainingTime > 0) {
 minRemainingTime = processes.get(j).remainingTime;
 shortest = j;
 check = true;
 }
 }

 // ... (process execution and time updates)
 }

 // ... (calculate turnaround times and print results)
}
```
```

This is the core of the SRTF algorithm:

- It first sorts processes by arrival time.
- In each time unit, it finds the process with the shortest remaining time among the arrived processes.
- It executes this process for one time unit, updates remaining time, and checks for completion.
- This continues until all processes are complete.

4. Transient Event Handling:

```
```java
if (currentTime == transientEventTime && !transientEventOccurred) {
 processes.get(shortest).remainingTime += 3;
 System.out.println("Transient event: Process P" +
 (processes.get(shortest).id + 1) +
 " remaining time adjusted by +3 units at time " +
 currentTime);
 transientEventOccurred = true;
}
```
```

This section simulates a transient event at a specific time (5 in this case), increasing the remaining time of the currently executing process by 3 units.

5. Results Calculation and Printing:

```
```java
System.out.println("Process\tBurst Time\tArrival Time\tWaiting
Time\tTurnaround Time");
for (int i = 0; i < n; i++) {
 System.out.println("P" + (processses.get(i).id + 1) + "\t\t\t" +
processses.get(i).burstTime +
"\t\t\t" + processses.get(i).arrivalTime + "\t\t\t"
+ waitingTimes[i] +
"\t\t\t\t\t" + turnaroundTimes[i]);
}
```
```

```
double avgWaitingTime = Arrays.stream(waitingTimes).average().orElse(0.0);
double avgTurnaroundTime =
Arrays.stream(turnaroundTimes).average().orElse(0.0);

System.out.printf("\nAverage waiting time = %.2f\n", avgWaitingTime);
System.out.printf("Average turnaround time = %.2f\n", avgTurnaroundTime);
```
```

This section prints the results for each process and calculates average waiting and turnaround times.

## Key Features of this Implementation:

1. Preemptive Scheduling: The algorithm can switch processes at any time if a shorter job arrives.
2. Dynamic Priority: The priority (shortest remaining time) is recalculated at every time unit.
3. Transient Event Simulation: Demonstrates how unexpected events can affect scheduling.
4. Performance Metrics: Calculates waiting time and turnaround time for each process and their averages.

This implementation provides a realistic simulation of the SRTF scheduling algorithm, incorporating the complexity of varying arrival times and a transient event. It demonstrates how SRTF prioritizes processes with the shortest remaining time, potentially leading to better average waiting times compared to non-preemptive algorithms.