

# Digital Logic Design: Homework 1

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Unless otherwise specified, all numbers are decimal.

## 1.4

a)

$$\begin{aligned} &1457.11 \\ &= 1457 + 0.11 \\ &= (16 \cdot 91 + 1) + 0.11 \\ &= (16 \cdot (16 \cdot 5 + 11) + 1) + 0.11 \\ &= (16 \cdot (16 \cdot 5 + 11) + 1) + \left(\frac{1}{16} \cdot 1.76\right) \\ &= (16 \cdot (16 \cdot 5 + 11) + 1) + \left(\frac{1}{16} \cdot \left(1 + \frac{1}{16} \cdot 12.16\right)\right) \\ &\text{Thus } 1457.11_{10} = 5B1.1C_{16}. \end{aligned}$$

## 1.7

b)

$-14 = 110010_2$ ,  $-32 = 100000_2$  in 2's complement representation.

$$\begin{array}{r} 110010 \\ + 100000 \\ \hline 1010010 \end{array}$$

The result, after discarding the carry from the sign bit, is positive. It is clear that there is an overflow.

c)

$-25 = 100111_2$ ,  $18 = 010010_2$  in 2's complement representation.

$$\begin{array}{r} 100111 \\ + 010010 \\ \hline 111001 \end{array}$$

There is no carry from the sign bit and there is no overflow.

Using 1's complement representation,  $-25 = 100110_2$ ,  $18 = 010010_2$ .

$$\begin{array}{r} 100110 \\ + 010010 \\ \hline 111000 \end{array}$$

There is no carry from the sign bit and there is no overflow.

## 1.8

If 2's complement is used, the system's range is  $-2^7$  to  $+(2^7 - 1)$ , or  $-128$  to  $+127$ . If 1's complement is used, the system's range is  $-(2^7 - 1)$  to  $+(2^7 - 1)$ , or  $-127$  to  $+127$ .

## 1.10

a)

$$1305.375$$

$$= 1305 + 0.375$$

$$= (16 \cdot 81 + 9) + 0.375$$

$$= (16 \cdot (16 \cdot 5 + 1) + 9) + 0.375$$

$$= (16 \cdot (16 \cdot 5 + 1) + 9) + \frac{1}{16} \cdot 6$$

$$\text{Thus } 1305.375_{10} = 519.6_{16}.$$

Conversion to decimal is easy by converting each hexadecimal digit to its binary representation:

$$0101\ 0001\ 1001\ .\ 0110$$

$$\text{That is, } 519.6_{16} = 010100011001.0110_2.$$

d)

$$1644.875$$

$$= 1644 + 0.875$$

$$= (16 \cdot 102 + 12) + 0.875$$

$$= (16 \cdot (16 \cdot 6 + 6) + 12) + 0.875$$

$$= (16 \cdot (16 \cdot 6 + 6) + 12) + \frac{1}{16} \cdot 14$$

$$\text{Thus } 1644.875_{10} = 66C.E_{16}.$$

Conversion to decimal is easy by converting each hexadecimal digit to its binary representation:

$$0110\ 0110\ 1100\ .\ 1110$$

$$\text{That is, } 66C.E_{16} = 011001101100.1110_2.$$

## 1.19

b)

Shown below is division with binary numbers:

$$\begin{array}{r} 11011 \\ 1110 \overline{) 110000001} \\ \underline{1110} \phantom{00000000} \\ 10100 \phantom{00000000} \\ \underline{1110} \phantom{00000000} \\ 11000 \phantom{00000000} \\ \underline{1110} \phantom{00000000} \\ 10101 \phantom{00000000} \\ \underline{1110} \phantom{00000000} \\ 111 \phantom{00000000} \end{array}$$

$$14 \cdot 27 + 7 = 385,$$

$$1110_2 \cdot 11011_2 + 111_2 = 110000001_2,$$

The answer of  $11011_2$  is correct, with remainder  $111_2$ .

## 1.37

a)

$$\begin{array}{r} 01001 \\ -11010 \\ \hline \end{array}$$

We interpret this as 1's complement:

$$\begin{array}{r} 01001 \\ +00101 \\ \hline 01110 \end{array}$$

And as 2's complement:

$$\begin{array}{r} 01001 \\ +00110 \\ \hline 01111 \end{array}$$

No carry occurs on the sign bits, so there is no overflow.

c)

$$\begin{array}{r} 10110 \\ -01101 \\ \hline \end{array}$$

We interpret this as 1's complement:

$$\begin{array}{r} 10110 \\ +10010 \\ \hline 101000 \\ \hline 01001 \end{array}$$

And as 2's complement:

$$\begin{array}{r} 10110 \\ +10011 \\ \hline 101001 \\ \hline 01001 \end{array}$$

There is overflow with both methods.

e)

$$\begin{array}{r} 11100 \\ -10101 \\ \hline \end{array}$$

We interpret this as 1's complement:

$$\begin{array}{r} 11100 \\ +01010 \\ \hline 100110 \\ \hline 00111 \end{array}$$

And as 2's complement:

$$\begin{array}{r} 11100 \\ +01011 \\ \hline 100111 \\ \hline 00111 \end{array}$$

There is no overflow.

## 2.3

a)

$X'Y'Z + (X'Y'Z)' = 1X'Y'Z + 1(X'Y'Z)' = 1$ , by the uniting theorem of boolean addition.

f)

$(A + BC) + (DE + F)(A + BC)' = (A + BC) + (DE + F)$ , by the elimination theorem of boolean addition.

## 2.5

a)

$$\begin{aligned}
& (A+B)(C+B)(D'+B)(ACD'+E) \\
&= (A(C+B)+B(C+B))(D'+B)(ACD'+E) \\
&= (AC+AB+BC+B^2)(D'+B)(ACD'+E) \\
&= (AC+B)(D'+B)(ACD'+E) \\
&= (AC+B)(D'(ACD'+E)+B(ACD'+E)) \\
&= (AC+B)(AC(D')^2+D'E+ABCD'+BE) \\
&= (AC+B)(ACD'+BE+D'E) \\
&= AC(ACD'+BE+D'E)+B(ACD'+BE+D'E) \\
&= A^2C^2D'+ACBE+ACD'E+BACD'+B^2E+BD'E) \\
&= ACD'+ACBE+ACD'E+BACD'+BE+BD'E) \\
&= ACD'+BE
\end{aligned}$$

b)

$$\begin{aligned}
& (A'+B+C')(A'+C'+D)(B'+D') \\
&= (A'(B'+D')+B(B'+D')+C'(B'+D'))(A'+C'+D) \\
&= (A'B'+A'D'+BB'+BD'+C'B'+C'D')(A'+C'+D) \\
&= (A'B'+A'D'+BD'+B'C'+C'D')(A'+C'+D) \\
&= A'(A'B'+A'D'+BD'+B'C'+C'D')+C'(A'B'+A'D'+BD'+B'C'+C'D') \\
&= A'A'B'+A'A'D'+A'BD'+A'B'C'+A'C'D'+C'A'B'+C'A'D'+C'BD'+C'B'C'+C'C'D'+DA'B'+DA'D'+DBD'+DB'C'+DC'D' \\
&= A'B'+A'D'+A'BD'+A'B'C'+A'C'D'+A'B'C'+A'C'D'+BC'D'+B'C'+C'D'+A'B'D+0+0+B'C'D+0 \\
&= A'B'+A'D'+B'C'+C'D'
\end{aligned}$$

## 2.6

a)

$$\begin{aligned}
& AB+C'D' \\
&= (AB+C')(AB+D') \\
&= (A+C')(B+C')(A+D')(B+D')
\end{aligned}$$

f)

$$\begin{aligned}
& A+BC+DE \\
&= (A+B)(A+C)+DE \\
&= ((A+B)(A+C)+D)((A+B)(A+C)+E) \\
&= ((A+B+D)(A+C+D))((A+B+E)(A+C+E)) \\
&= (A+B+D)(A+C+D)(A+B+E)(A+C+E)
\end{aligned}$$

## 2.8

a)

$$[(AB)' + C'D]' = AB(C'D)' = AB(C + D') = ABC + ABD'$$

## 2.9

a)

$$\begin{aligned} F &= ((A + (A + B)')' + (A + B)')(A + (A + B)')' \\ &= ((A'(A + B)) + (A'B'))(A + A'B')' \\ &= ((A'(A + B)) + (A'B'))(A' + (A'B')') \\ &= ((A'(A + B)) + (A'B'))(A' + A + B) \\ &= ((A'(A + B)) + (A'B'))(1) \\ &= (A'A + A'B + (A'B')) \\ &= 0 + A'B + A'B' \\ &= A' \end{aligned}$$

b)

$$\begin{aligned} G &= (((R + S + T)'P((R + S)'T))'T)' \\ &= (R + S + T)'P((R + S)'T) + T' \\ &= R'S'T'PR'S'T + T' \\ &= PR'S'TT' + T' \\ &= 0 + T' \\ &= T' \end{aligned}$$

## 2.12

a)

$(X + Y'Z) + (X + Y'Z)' = 1(X + Y'Z) + 1(X + Y'Z)' = 1$ , by the unification theorem of boolean addition.

f)

$$\begin{aligned} &(V' + U + W)[(W + X) + Y + UZ'] + [(W + X) + UZ' + Y] \\ &= (V' + U + W)[(W + X) + UZ' + Y] + [(W + X) + UZ' + Y] \\ &= V' + U + W, \text{ by the absorption theorem of boolean addition.} \end{aligned}$$

## 2.15

a)

$$\begin{aligned}
 f(A, B, C, D) &= [A + (BCD)'][(AD)' + B(C' + A)] \\
 dual(f) &= [A(B + C + D)'] + [(A + D)'(B + C'A)] \\
 complement(f) &= [A'(B' + C' + D')'] + [(A' + D')'(B' + CA')] \\
 &= A'BCD + AD(B' + CA')
 \end{aligned}$$

## 2.16

a)

$$\begin{aligned}
 f(A, B, C, D) &= [A + (BCD)'][(AD)' + B(C' + A)] \\
 dual(f) &= [A(B + C + D)'] + [(A + D)'(B + C'A)]
 \end{aligned}$$

## 2.21

A	B	C	F	G	H
1	1	1	1	<b>X</b>	1
1	1	0	0	<b>0</b>	0
1	0	1	0	<b>1</b>	1
1	0	0	0	<b>0</b>	0
0	1	1	1	<b>X</b>	1
0	1	0	0	<b>1</b>	1
0	0	1	1	<b>X</b>	1
0	0	0	0	<b>0</b>	0

X represents an unspecified value in the above table.

## 2.22

a)

$$\begin{aligned}
 &A'B' + A'CD + A'DE' \\
 &= A'(B' + CD + DE') \\
 &= A'((B' + C)(B' + D) + DE') \\
 &= A'((B' + C + DE')(B' + D + DE')) \\
 &= A'(B' + C + DE')(B' + D) \\
 &= A'(B' + C + D)(B' + C + E')(B' + D) \\
 &= A'(B' + D)(B' + C + E')
 \end{aligned}$$

## 2.24

a)

$$[(XY')' + (X' + Y)'Z] = (X' + Y) + (XY')Z = X' + Y + XY'Z$$