MA5233 Computational Mathematics

Lecture 4: QR Factorisation

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2019/2020

QR factorisation

- An alternative way to solve linear systems.
- ► A basic building block in many more advanced algorithms.

Outline

- QR factorisation: definition and applications.
- ► Algorithms:
 - Gram-Schmidt orthogonalisation
 - Householder reflection
 - Givens rotation
- ▶ qr() in Julia.
- Complexity and stability.

QR factorisation

Any $A \in \mathbb{K}^{m \times n}$ can be written as A = QR, where

- ▶ $Q \in \mathbb{K}^{m \times m}$ is orthogonal $(Q^H Q = I)$, and
- ▶ $R \in \mathbb{K}^{m \times n}$ is upper triangular (R(i,j) = 0 for i > j).

Applications

▶ Solve linear systems (assuming m = n):

$$QRx = b \iff x = R^{-1}Q^Hb.$$

▶ Determine orthogonal basis (see next slide).

Key property of QR factorisation

Q factor provides an orthogonal basis for range of A.

More precisely:

$$A[:,1] = Q[:,1] R[1,1],$$

$$A[:,2] = Q[:,1] R[1,2] + Q[:,2] R[2,2],$$

$$A[:,3] = Q[:,1] R[1,3] + Q[:,2] R[2,3] + Q[:,3] R[3,3]$$
...

Conclusion: span A[:, 1 : k] = span Q[:, 1 : k].

Illustration of QR decomposition

Case m > n:

Case m < n:

Red part of Q factor is redundant since it gets multiplied by zeros in R. QR factorisation without red part is known as "thin" QR.

Algorithms for computing the QR factorisation

- ► Gram-Schmidt orthogonalisation
- ► Householder reflections
- Givens rotations

All of these have strengths and weaknesses.

We next go through each of them.

Gram-Schmidt theorem

Let $q_1, \ldots, q_k \in \mathbb{K}^n$ be orthonormal and $a_{k+1} \in \mathbb{K}^n$ be linearly independent of the q_k . Then,

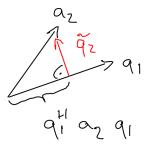
$$ilde{q}_{k+1} := \mathsf{a}_{k+1} - \sum_{\ell=1}^k \mathsf{q}_\ell^H \mathsf{a}_{k+1} \, \mathsf{q}_\ell$$

is orthogonal to q_1, \ldots, q_k .

Proof.

$$q_{\ell}^{H} \, \tilde{q}_{k+1} = q_{\ell}^{H} \, a_{k+1} - q_{\ell}^{H} q_{\ell} \, q_{\ell}^{H} a_{k+1} = 0.$$

Gram-Schmidt theorem, pictorial version



QR factorisation via Gram-Schmidt theorem

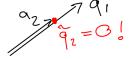
Rewrite Gram-Schmidt formula

$$ilde{q}_{k+1} = a_{k+1} - \sum_{\ell=1}^k q_\ell^H a_{k+1} \, q_\ell$$

in the form

$$a_{k+1} = \underbrace{\|q_{k+1}\|_2}_{R[k+1,k+1]} q_{k+1} + \sum_{\ell=1}^k \underbrace{q_{\ell}^H a_{k+1}}_{R[\ell,k+1]} q_{\ell}.$$

Breaking the Gram-Schmidt algorithm



Classical Gram-Schmidt

$$ilde{q}_{k+1}^{(0)} := a_{k+1}, \qquad ilde{q}_{k+1}^{(\ell+1)} := ilde{q}_{k+1}^{\ell} - q_{\ell}^{H} extbf{a}_{k+1} \, q_{\ell}$$

Modified Gram-Schmidt

$$\tilde{q}_{k+1}^{(0)} := a_{k+1}, \qquad \tilde{q}_{k+1}^{(\ell+1)} := \tilde{q}_{k+1}^{(\ell)} - q_{\ell}^{H} \tilde{q}_{k+1}^{(\ell)} q_{\ell}$$

Modified Gram-Schmidt is more numerically stable.

Example

Classical Gram-Schmidt:

$$\begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 0 & \varepsilon & 0 \\ 0 & 0 & \varepsilon \end{pmatrix} \longrightarrow \begin{pmatrix} 1 & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} \\ 1 & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} \\ 0 & 1 & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-2} \\ 0 & 0 & 1 \end{pmatrix}$$

Modified Gram-Schmidt:

$$\begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 0 & \varepsilon & 0 \\ 0 & 0 & \varepsilon \end{pmatrix} \longrightarrow \begin{pmatrix} 1 & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} \\ 1 & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} & \varepsilon_{\mathsf{mach}} \, \varepsilon^{-1} \\ 0 & 1 & \varepsilon_{\mathsf{mach}}^2 \, \varepsilon^{-2} \\ 0 & 0 & 1 \end{pmatrix}$$

Householder reflection theorem

Assume $u \in \mathbb{K}^n$ such that ||u|| = 1. Then,

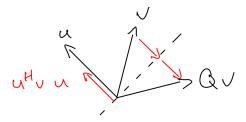
$$Q := I - 2 uu^H$$

is an orthogonal matrix, and $Q^2 = I$.

Proof.

$$Q^{H} Q = Q^{2} = (I - 2 uu^{H}) (I - 2 uu^{H}) = I - 4uu^{H} + 4u (u^{H}u) u^{H} = I.$$

Householder reflection theorem, pictorial version



QR factorisation via Householder reflections

Idea: choose u such that $Qa_1=\pm \alpha e_1$. Then,

$$Q\begin{pmatrix} x & x & x \\ x & x & x \\ x & x & x \end{pmatrix} = \begin{pmatrix} x & x & x \\ 0 & x & x \\ 0 & x & x \end{pmatrix}.$$

How to find $Q = I - 2uu^H$?

Determining Householder *u*

- ▶ Reflections preserve length; hence $\alpha = ||a_1||$.
- ▶ $a_1 u u^H a_1 = \pm ||a_1|| e_1$ implies $u \propto a_1 \mp ||a_1|| e_1$.
- ► Choose $Qa_1 = -\operatorname{sign}(a_{11}) \|a_1\| e_1$ to avoid cancellation in u. a_{11} denotes first component of a_1 .



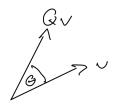
Givens rotation

$$Q:=egin{pmatrix} \cos(heta) & -\sin(heta) \ \sin(heta) & \cos(heta) \end{pmatrix}$$
 represents a rotation by $heta.$

Example

$$\begin{bmatrix} \cos 90^{\circ} & \sin 90^{\circ} \\ -\sin 90^{\circ} & \cos 90^{\circ} \end{bmatrix} \begin{bmatrix} \alpha_{1} \\ \alpha_{2} \end{bmatrix} = \underbrace{5 \quad 90^{\circ}}$$

Serious example



QR factorisation via Givens rotation

Similar idea as in Householder case:

$$\begin{pmatrix} x & x & x \\ x & x & x \\ x & x & x \end{pmatrix} \rightarrow \begin{pmatrix} x & x & x \\ 0 & x & x \\ x & x & x \end{pmatrix} \rightarrow \begin{pmatrix} x & x & x \\ 0 & x & x \\ 0 & x & x \end{pmatrix} \rightarrow \begin{pmatrix} x & x & x \\ 0 & x & x \\ 0 & 0 & x \end{pmatrix}$$

QR factorisation in Julia

F = qr(A) computes "factorisation object".

- Access factors through F.Q, F.R, F.p (vector) and F.P (matrix).
- ► F.Q * F.U == A[:,F.p] == A * F.P.
- x = F\b computes solution.

Computational cost of QR factorisation

 $\mathcal{O}(mn \min\{m,n\})$ floating-point operations.

Stability of Householder QR factorisation

Q,R computed by Householder QR satisfy

$$QR = A + \delta A$$
 where $\frac{\|\delta A\|}{\|A\|} = \mathcal{O}(\varepsilon_{\mathsf{mach}}).$

Hence, Householder QR factorisation is backwards stable.

References and further reading

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