# MA5233 Computational Mathematics

Lecture 13: Multigrid

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### Multigrid algorithm

- ▶ Another algorithm for solving the Poisson equation.
- ► Key selling point: requires only  $\mathcal{O}(N)$  FLOP. As before, N denotes the matrix size.  $N = n^2$  for 2d Poisson on  $n \times n$  grid.
  - ▶ LU factorisation:  $\mathcal{O}(N^{3/2})$  FLOP in 2d.
  - ▶ Fourier transform:  $\mathcal{O}(N \log N)$  FLOP.
- ► Generality: somewhere between LU (solves any linear system) and fast Fourier transform (only solves very specific problem).

#### Outline

- ► Jacobi and Gauss-Seidel methods.
- Convergence analysis for Jacobi method.
- From Jacobi to multigrid.

### Jacobi iteration for Poisson equation

Assume we have an initial guess  $x^{(0)}$  for the linear system

$$(3+1)^2 \begin{pmatrix} 2 & -1 \\ -1 & 2 & -1 \\ & -1 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ b_3 \end{pmatrix}.$$

Solving each equation for the "diagonal" unknown yields

$$x_1^{(1)} = \frac{1}{2} \left( \frac{b_1}{(3+1)^2} + x_2^{(0)} \right),$$
  

$$x_2^{(1)} = \frac{1}{2} \left( \frac{b_2}{(3+1)^2} + x_1^{(0)} + x_3^{(0)} \right),$$
  

$$x_3^{(1)} = \frac{1}{2} \left( \frac{b_3}{(3+1)^2} + x_2^{(0)} \right).$$

The resulting  $x^{(1)}$  is not the exact solution in general, but we may hope that it is a better approximation to x than  $x^{(0)}$ .

Idea: iterate the map  $x^{(0)} \mapsto x^{(1)}$  until convergence.

See jacobi\_step() in 13\_multigrid.jl.

#### **Gauss-Seidel iteration**

Minor modification on Jacobi:

$$\begin{split} x_1^{(1)} &= \frac{1}{2} \left( \frac{b_1}{(3+1)^2} + x_2^{(0)} \right), \\ x_2^{(1)} &= \frac{1}{2} \left( \frac{b_2}{(3+1)^2} + x_1^{(1)} + x_3^{(0)} \right), \\ x_3^{(1)} &= \frac{1}{2} \left( \frac{b_3}{(3+1)^2} + x_2^{(1)} \right). \end{split}$$

See gauss\_seidel\_step() in 13\_multigrid.jl

Comparison with Jacobi iteration:

- ► Good: iteration can be done *in-place*:
  - ▶ Jacobi: read from  $x^{(0)}$ , write to  $x^{(1)}$ .
  - ► Gauss-Seidel: read to and write from single vector.
- Good: convergence is faster.
- ► Bad: loss of parallelism:
  - ▶ Jacobi: every entry of  $x^{(1)}$  can be computed independently.
  - ▶ Gauss-Seidel:  $x_k^{(1)}$  must be computed after  $x_\ell^{(1)}$  with  $\ell < k$ .

### Discussion of Jacobi-type methods

- ▶ Good: performing a single iteration is very fast if matrix is sparse.
- Bad: many iterations are needed to reach a reasonable accuracy.

See plot\_convergence() in 13\_multigrid.jl.

### Next steps

Goal: convergence estimate  $||x^{(k)} - x|| = \mathcal{O}(f(k))$ . Intermediate step: matrix formula for Jacobi iteration.

#### Remarks

- We will discuss only the Jacobi iteration in this lecture.
- Analysis of Gauss-Seidel is analogous but more complicated.

### Jacobi iteration, general definition

Let A be an invertible matrix with nonzero diagonal D. Jacobi iteration is defined as follows.

# **Algorithm 1** Jacobi iteration

- 1: **for** k = 1, 2, ... **do**
- 2:  $x^{(k)} = D^{-1} (b (A D) x^{(k-1)})$
- 3: end for

The next slide demonstrates that the general definition reduces to the concrete Jacobi iteration when applied to the Poisson matrix.

### **Example: Jacobi iteration for Poisson matrix**

Consider the matrix

$$A = (3+1)^2 \begin{pmatrix} 2 & -1 \\ -1 & 2 & -1 \\ & -1 & 2 \end{pmatrix}$$

The Jacobi iteration takes the form

$$\begin{split} x^{(1)} &= D^{-1} \left( b - (A - D) x^{(0)} \right) \\ &= \begin{pmatrix} 2 \\ 2 \\ 2 \end{pmatrix}^{-1} \left( \frac{1}{(3+1)^2} \begin{pmatrix} b_1 \\ b_2 \\ b_3 \end{pmatrix} - \begin{pmatrix} -1 \\ -1 \\ -1 \end{pmatrix} \begin{pmatrix} x_1^{(0)} \\ x_2^{(0)} \\ x_3^{(0)} \end{pmatrix} \right) \\ &= \begin{pmatrix} \frac{1}{2} \left( \frac{b_1}{(3+1)^2} + x_2^{(0)} \right) \\ \frac{1}{2} \left( \frac{b_2}{(3+1)^2} + x_1^{(0)} + x_3^{(0)} \right) \\ \frac{1}{2} \left( \frac{b_3}{(3+1)^2} + x_2^{(0)} \right) \end{pmatrix}. \end{split}$$

This is precisely the iteration we had before.

# Towards a convergence estimate for Jacobi iteration

Using the Jacobi iteration formula and b = Ax, we obtain

$$x^{(k)} - x = D^{-1} \left( b - (A - D) x^{(k-1)} \right) - x$$
$$= D^{-1} \left( Ax - Dx - (A - D) x^{(k-1)} \right)$$
$$= -D^{-1} \left( A - D \right) \left( x^{(k-1)} - x \right).$$

Applying this formula repeatedly yields

$$x^{(k)} - x = R^k (x^{(0)} - x)$$
 where  $R := -D^{-1} (A - D)$ .

Let us expand initial error in terms of eigenvectors  $u_{\ell}$  of R,

$$x^{(0)} - x = \sum_{\ell=1}^{N} c_{\ell} u_{\ell}.$$

Denoting the eigenvalue associated with  $u_{\ell}$  by  $\lambda_{\ell}$ , we obtain

$$x^{(k)} - x = R^k (x^{(0)} - x) = \sum_{\ell=1}^N c_\ell R^k u_\ell = \sum_{\ell=1}^N c_\ell \lambda_\ell^k u_\ell.$$

### Towards a convergence estimate for Jacobi iteration

From previous slide:

$$x^{(k)} - x = \sum_{\ell=1}^{N} c_{\ell} \lambda_{\ell}^{k} u_{\ell}.$$

Assume  $||u_{\ell}||=1$  and eigenvalues are sorted such that  $|\lambda_1|\leq\ldots\leq |\lambda_N|$ . Then,

$$||x^{(k)} - x|| \le \sum_{\ell=1}^{N} |c_{\ell}| |\lambda_{\ell}|^k \le \left(\sum_{\ell=1}^{N} |c_{\ell}|\right) |\lambda_{N}|^k.$$

#### Conclusion

Jacobi iterates  $x^{(k)}$  satisfy

$$||x^{(k)} - x|| \le C|\lambda_N|^k$$

where  $\lambda_N$  is the eigenvalue of largest absolute value of

$$R = -D^{-1}(A - D).$$

### Convergence rate of Jacobi iteration for Poisson equation

Let us consider

$$A = (N+1)^2 \begin{pmatrix} 2 & -1 & & \\ -1 & \ddots & \ddots & \\ & \ddots & \ddots & -1 \\ & & -1 & 2 \end{pmatrix} \in \mathbb{R}^{N \times N} \quad \implies \quad D = (N+1)^2 \begin{pmatrix} 2 & & & \\ & \ddots & & \\ & & \ddots & 2 \end{pmatrix}.$$

Let us introduce the following notation:

- $\lambda_{\ell} = (N+1)^2 \left(2 2\cos\left(\pi \frac{\ell}{N+1}\right)\right)$ : eigenvalues of A.
- $\hat{\lambda}_{\ell}$ : eigenvalues of  $R = -D^{-1}(A D)$ .

Since  $D \propto I$ ,  $\lambda_{\ell}$  and  $\hat{\lambda}_{\ell}$  are related by  $\hat{\lambda}_{\ell} = -\frac{1}{2} \left( -\frac{\lambda_{\ell}}{(N+1)^2} - 2 \right)$ .

Inserting the known values for  $\lambda_{\ell}$  yields  $\hat{\lambda}_{\ell} = \cos\left(\pi \frac{\ell}{N+1}\right)$ .

Largest absolute value is achieved for  $\ell=1$  and  $\ell=\mathit{N}$  for which we have

$$\lambda_1 = -\lambda_n = \cos\left(\frac{\pi}{N+1}\right) = 1 - \frac{\pi^2}{(N+1)^2} + \mathcal{O}(N^{-4}).$$

See 13\_multigrid.jl.

### Summary of Jacobi convergence theory

Let  $\lambda_N$  be the eigenvalue of largest absolute value of  $R = -D^{-1}(A - D)$ .

Then, there exists a constant 
$$C \neq C(k)$$
 such that  $\|x^{(k)} - x\| \leq C |\lambda_N|^k$ .

For Poisson matrix

$$A = \begin{pmatrix} 2 & -1 \\ -1 & \ddots & \ddots \\ & \ddots & \ddots & -1 \\ & & -1 & 2 \end{pmatrix} \in \mathbb{R}^{N \times N},$$

we have

$$|\lambda_N| = 1 - \mathcal{O}(N^{-2}).$$

# Why Jacobi iteration for Poisson equation must be slow Consider

$$A = (N+1)^2 \begin{pmatrix} 2 & -1 \\ -1 & \ddots & \ddots \\ & \ddots & \ddots & -1 \\ & & -1 & 2 \end{pmatrix} \in \mathbb{R}^{N \times N}, \qquad b_i := \begin{cases} (N+1)^2 & \text{if } i = 1, \\ 0 & \text{otherwise.} \end{cases}$$

Solution x is given by  $x_i = 1 - \frac{i}{N+1} \neq 0$ .

Jacobi iterates  $x^{(k)}$  satisfy  $x_i^{(k)} = 0$  if i > k + 1.

Conclusion:  $x^{(k)}$  cannot be a good approximation to x for k < N.

See speed\_of\_propagation() in 13\_multigrid.jl.

### Why Krylov subspace methods for Poisson equation must be slow

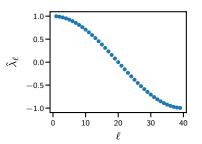
- Above argument also applies to unpreconditioned Krylov methods since similarly  $(A^k b)_i = 0$  if i > k + 1.
- ▶ A good preconditioner must introduce non-locality, i.e.  $P^{-1}$  should be a dense matrix. (ILU satisfies this requirement.)

### A closer look at Jacobi convergence theory

Recall: every Jacobi iteration multiplies error by R,

$$x^{(k)} - x = R^k (x^{(0)} - x).$$

Eigenvalues of R for Poisson matrix:  $\hat{\lambda}_{\ell} = \cos\left(\pi \frac{\ell}{N+1}\right)$ 

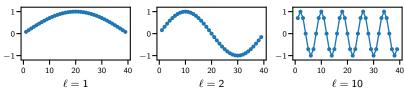


Only error components associated with  $\ell \approx 1$  and  $\ell \approx N$  converge slowly. Error components associated with  $\ell \approx \frac{N}{2}$  converge very fast.

Let us take a closer look at the eigenfunctions for  $\ell \approx 1$ .

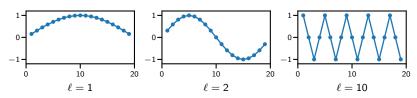
### A closer look at Jacobi convergence theory

Eigenvectors of R are given by  $(u_\ell)_i = \sin(\pi \frac{\ell i}{n+1})$ .



#### Observations:

- $\blacktriangleright$   $\ell$  indicates the *frequency* (number of oscillations) of  $u_{\ell}$ .
- ▶  $u_\ell$  with small  $\ell$  are *smooth*: only few grid points are necessary to capture the rough shape of the curve.



### The multigrid idea

- Use Jacobi to reduce high-frequency errors.
   We have seen that Jacobi is good at reducing intermediate-frequency errors.
   We will see shortly how to use Jacobi to reduce high-frequency errors.
- Solve for low-frequency errors on coarser grid ("coarse grid correction").This switching between grids is why the method is called multigrid.

### Benefits of switching to coarser grid

- Obvious: problem becomes smaller and hence more manageable.
- Less obvious: switching between grids allows us to reuse Jacobi smoothing on coarser levels.

### Topics for the following slides

- ► How to use Jacobi to reduce high-frequency errors?
- How to perform coarse grid correction?

# Jacobi smoothing for high-frequency errors

Observation in plot on slide 13:

High-frequency errors oscillate:  $Ru_{\ell} \approx (-1) u_{\ell}$  for  $\ell \approx N$ .

#### Consequences:

▶ If we replace Jacobi iteration  $x^{(k)} = D^{-1} (b - (A - D) x^{(k-1)})$  with

$$x^{(k)} = (1 - \theta) x^{(k-1)} + \theta D^{-1} (b - (A - D) x^{(k-1)})$$

for some  $\theta \in (0,1]$ , then high-frequency errors approximately cancel.

▶ More precisely, the error recursion formula becomes

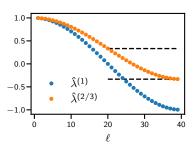
$$x^{(k)} - x = \underbrace{\left( (1 - \theta) I - \theta D^{-1} (A - D) \right)}_{R_{\theta}} (x^{(k-1)} - x),$$

and the eigenvalues of  $R_{\theta}$  are

$$\hat{\lambda}_{\ell}^{(\theta)} = (1-\theta) + \theta \, \hat{\lambda}_{\ell} = (1-\theta) + \theta \, \cos\big(\pi \tfrac{\ell}{N+1}\big).$$

# Jacobi smoothing for high-frequency errors (continued)

▶ For  $\theta = \frac{2}{3}$ , all error components for  $\ell \geq \frac{N}{2}$  are multiplied by  $\frac{1}{3}$  in every Jacobi iteration.



# Jacobi smoothing for high-frequency errors (conclusion)

Relaxed Jacobi iteration

$$x^{(k)} = \frac{1}{3} x^{(k-1)} + \frac{2}{3} D^{-1} (b - (A - D) x^{(k-1)})$$

efficiently reduces high-frequency error.

Method is called *relaxed* because it takes only  $\frac{2}{3}$  of the step proposed by Jacobi.

After a few steps of relaxed Jacobi, we obtain  $\tilde{x}^{(0)}$  such that

$$\tilde{x}^{(0)} - x = \underbrace{\sum_{\ell=1}^{N/2} c_{\ell} \, u_{\ell}}_{A} + \underbrace{\sum_{\ell=N/2+1}^{N} c_{\ell} \, u_{\ell}}_{B}$$

where B is small and A is smooth.

Next topic: coarse grid correction for eliminating the smooth error A.

### Coarse grid correction

Assume N = 2N' + 1.

- 1. Compute residual  $r := b A\tilde{x}^{(0)}$ .
- 2. Approximate residual on coarser grid:

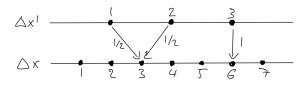
$$r' \in \mathbb{R}^{N'}, \qquad r'_i = \frac{r_{2i-1}}{4} + \frac{r_{2i}}{2} + \frac{r_{2i+1}}{4}.$$

3. Solve on coarse grid:  $\Delta x' := A'^{-1}r'$ .

### Coarse grid correction (continued)

4. Interpolate coarse grid correction  $\Delta x'$  to fine grid:

$$\Delta x \in \mathbb{R}^N, \qquad \Delta x_{2i} := \Delta x_i', \qquad \Delta x_{2i+1} := \frac{\Delta x_i' + \Delta x_{i+1}'}{2}$$



5. Update  $\tilde{x}^{(1)} := \tilde{x}^{(0)} + \Delta x$ .

See twogrid\_step() in 13\_multigrid.jl.

# Idea of coarse grid correction

Smoothness of  $\tilde{x}^{(0)}-x$  implies smoothness of  $r=-A(\tilde{x}^{(0)}-x)$ . This smoothness ensures that coarse problem is "close" to fine problem. Hence,  $\Delta x \approx A^{-1} \, r = x - \tilde{x}$  and  $\tilde{x} + \Delta x \approx x$ .

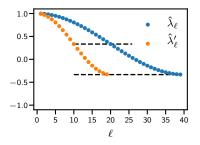
### Solving the coarse grid problem

Step 3 of coarse grid correction:  $\Delta x' = A'^{-1} r'$ .

How to solve this linear system?

### Key observation

Intermediate frequencies on fine grid  $\rightarrow$  high frequencies on coarse grid.



Idea: apply smoothing  $\rightarrow$  coarsening  $\rightarrow$  smoothing  $\rightarrow$  ... recursively. See multigrid\_step() in 13\_multigrid.jl.

### Multigrid convergence theory

The algorithm implemented in twogrid\_step() is as follows:

- 1. Compute  $\tilde{x}^{(0)}$  from  $x^{(0)}$  using one relaxed Jacobi step.
- 2. Compute and apply coarse grid correction,  $\tilde{x}^{(1)} = \tilde{x}^{(0)} + \Delta x$
- 3. Compute  $x^{(1)}$  from  $\tilde{x}^{(1)}$  using one relaxed Jacobi step.

The error recursion formula for relaxed Jacobi from slide 16 yields

$$\tilde{x}^{(0)} - x = R_{2/3}(x^{(0)} - x)$$
 and  $x^{(1)} - x = R_{2/3}(\tilde{x}^{(1)} - x)$ .

The following slides derive a similar formula for the coarse grid correction.

### Multigrid convergence theory (continued)

Observations:

▶ Approximation step can be written as  $r' = \frac{1}{2} P^T r$  where

$$\frac{1}{2}P^{T} = \begin{pmatrix} \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & & & \\ & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & & & \\ & & & \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & & & & & \\ & & & & \ddots & \ddots & \ddots \end{pmatrix}.$$

Interpolation step can be written as  $\Delta x = P \Delta x'$  where

$$P = \begin{pmatrix} \frac{1}{2} & & \\ \frac{1}{2} & \frac{1}{2} & \\ & \frac{1}{2} & \frac{1}{2} & \\ & & \frac{1}{2} & \frac{1}{2} & \\ & & & \frac{1}{2} & \ddots \end{pmatrix}.$$

### Multigrid convergence theory (continued)

Coarse grid correction takes the form  $\tilde{x}^{(1)} = \tilde{x}^{(0)} + \Delta x$  where

$$\Delta x = \frac{1}{2} P A'^{-1} P^T r = -\frac{1}{2} P A'^{-1} P^T A (\tilde{x}^{(0)} - x).$$

Error recursion becomes

$$\tilde{x}^{(1)} - x = \tilde{x}^{(0)} - x + \Delta x 
= \left(I - \frac{1}{2} P A'^{-1} P^T A\right) (\tilde{x}^{(0)} - x).$$

Inserting the relaxed Jacobi error recursion yields

$$x^{(1)} - x = R_{2/3} \left( I - \frac{1}{2} P A'^{-1} P^T A \right) R_{2/3} \left( x^{(0)} - x \right).$$

### Multigrid convergence theory (continued)

Last formula on previous slide implies that convergence rate of twogrid\_step() is given by eigenvalue of largest absolute value of

$$R_{2g} = R_{2/3} \left( I - \frac{1}{2} P A'^{-1} P^T A \right) R_{2/3}.$$

#### Observations:

- ▶ We already know the eigenvalues and eigenvectors of  $R_{2/3}$ , A and A'.
- ▶ It can be shown that with  $(S_N)_{\ell k} := \sin(\pi \frac{\ell k}{N+1})$  we have

$$\left(S_{2N'+1} \, P \, S_{N'}\right)_{\ell,\ell'} \neq 0 \quad \iff \quad \ell = \ell' \text{ or } \ell = 2N'+1-\ell'.$$

▶ The nonzero entries of  $S_{2N'+1} P S_{N'}$  can be computed explicitly.

Using these results, lengthy calculations will reveal that  $R_{2g}$  has only two distinct eigenvalues, namely  $\frac{1}{9}$  and 0.

Disclaimer: I did not actually do these calculations, and it is possible that the above result is not 100% accurate. However, the  $\frac{1}{9}$  convergence rate is confirmed numerically in 13\_multigrid.jl.

### Multigrid convergence theory (final remarks)

Key result from previous slide:

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Convergence rate of multigrid is independent of grid size N!
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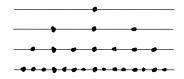
This is what sets multigrid apart from Jacobi and Krylov type methods.

#### Further remarks:

- ► Convergence depends on dimension. Convergence is slower in 2d.
- ▶ twogrid\_step() performs two relaxed Jacobi steps and hence reduces the high-frequency error by  $\frac{1}{9}$ . Heuristically, the result on the previous slide says that the coarse grid correction leads to the same reduction in the low-frequency errors.
- twogrid\_step() is not a practical algorithm for more complicated problems because we solve the coarse grid problem explicitly. However, it is reasonable to assume and confirmed numerically that multigrid\_step() performs only slightly worse.

### Computational cost of multigrid

▶ Multigrid employs a stack of grids of sizes  $N_{\ell} := 2^{\ell} - 1$ .



- For each grid point in any grid in the multigrid stack of grids, we perform  $\mathcal{O}(1)$  operations.
- ► Total number of grids points (on any level):

$$\sum_{\ell=1}^{L} (2^{\ell} - 1) = \frac{2^{L+1} - 1}{2 - 1} - L = \mathcal{O}(2^{L}) = \mathcal{O}(N_{L}).$$

#### Conclusion:

FLOP count of multigrid is proportional to number of grid points on finest level!

### Summary of multigrid method

Each iteration requires  $\mathcal{O}(N)$  FLOP and reduces error by factor  $\rho$  independent of N.

### Why multigrid for Poisson equation can be fast

Recall finite speed of propagation problem for Jacobi-type and Krylov subspace methods.

The coarse grid correction step eliminates this problem for multigrid algorithms: Jacobi iteration on coarse grid propagates information much faster than Jacobi step on fine grid.

#### Geometric and algebraic multigrid

The algorithm presented here is known as *geometric* multigrid since it exploits the geometric interpretation of the linear system that we try to solve.

There is a generalisation called *algebraic* multigrid which tries to form the stack of grids by only looking at the sparsity pattern of A. Generally, algebraic multigrid is easier to use (you only have to provide A and no further information) but performance is worse.

### References and further reading

Most of this lecture is based on Chapter 6 of the following book:

J. W. Demmel. *Applied Numerical Linear Algebra*. Society for Industrial and Applied Mathematics (1997), doi:10.1137/1.9781611971446