

MA5233 Computational Mathematics

Lecture 6: Sparse Matrices

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Sparse Matrices

Example

```
julia> n = 100_000
      A = Tridiagonal(
          fill(-1.0,n-1),
          fill( 2.0,n),
          fill(-1.0,n-1)
      )
      b = rand(100_000)
      @time A \ b;
0.011402 seconds (...)
```

```
julia> n = 10_000
      A = rand(n,n)
      b = rand(n)
      @time A \ b;
8.071135 seconds (...)
```

First matrix is 10x bigger, yet $A \setminus b$ is roughly 1000x faster.
How is this possible?

Sparse Matrices

Outline

- ▶ LU factorisation of triangular matrices
- ▶ Poisson equation in two dimensions
- ▶ General sparse matrix formats
- ▶ Sparse matrices in Julia
- ▶ Eigenvalues and -vectors of two-dimensional Laplacian.

Sparse Matrices

Theorem

LU factorisation of tridiagonal matrix is tridiagonal.

Proof sketch.

$$\begin{pmatrix} x & x & \\ x & x & x \\ & x & x \end{pmatrix} = \begin{pmatrix} 1 & & \\ x & 1 & \\ & & 1 \end{pmatrix} \begin{pmatrix} x & x & \\ & x & x \\ & x & x \end{pmatrix}$$
$$= \begin{pmatrix} 1 & & \\ x & 1 & \\ & x & 1 \end{pmatrix} \begin{pmatrix} x & x & \\ & x & x \\ & & x \end{pmatrix}$$

Corollary

LU factorisation of tridiagonal matrix takes $\mathcal{O}(n)$ FLOP instead of $\mathcal{O}(n^3)$.

Corollary of corollary

One-dimensional Poisson equation can be solved very efficiently!

Sparse Matrices

Discretising the two-dimensional Poisson equation

Functions:

- ▶ Introduce mesh $\Omega_n \times \Omega_n$ where

$$\Omega_n := \left\{ x_k := \frac{k}{n+1} \mid k = 0, \dots, n+1 \right\}.$$

- ▶ Replace function $f(x, y)$, with vector of point-values $f(x_{k_1}, x_{k_2})$.
- ▶ Use *lexicographical ordering* to arrange these point values as vector:

$$f_{k_1+n(k_2-1)} := f(x_{k_1}, x_{k_2}).$$

Example for lexicographical ordering

1	5	9	13
2	6	10	14
3	7	11	15
4	8	12	16

Sparse Matrices

Discretising the two-dimensional Poisson equation

Derivatives:

$$\frac{\partial^2 u}{\partial x^2} \longrightarrow (n+1)^2 (u_{i+1} - 2u_i + u_{i-1})$$

$$\frac{\partial^2 u}{\partial y^2} \longrightarrow (n+1)^2 (u_{i+n} - 2u_i + u_{i-n})$$

Laplacian now becomes

$$\Delta_n(i_1 + n(i_2 - 1), j_1 + n(j_2 - 1)) = \dots$$

$$= (n+1)^2 \begin{cases} -4 & \text{if } |i_1 - j_1| = 0 \text{ and } |i_2 - j_1| = 0, \\ 1 & \text{if } |i_1 - j_1| = 0 \text{ and } |i_2 - j_1| = 1, \\ 1 & \text{if } |i_1 - j_1| = 1 \text{ and } |i_2 - j_1| = 0, \\ 0 & \text{otherwise} \end{cases}$$

Sparse Matrices

Two-dimensional Laplacian matrix

$$\Delta_n \propto \begin{pmatrix} \begin{bmatrix} -4 & 1 & & & \\ 1 & \ddots & \ddots & & \\ & \ddots & \ddots & 1 & \\ & & 1 & -4 & \end{bmatrix} & \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & 1 \end{bmatrix} & & & \\ \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & 1 \end{bmatrix} & \begin{matrix} \ddots & & & \\ & \ddots & & \\ & & \ddots & \end{matrix} & \begin{matrix} \ddots & & \\ & \ddots & \\ & & \ddots \end{matrix} & & \\ & \begin{matrix} \ddots & & \\ & \ddots & \\ & & \ddots \end{matrix} & \begin{matrix} \ddots & & \\ & \ddots & \\ & & \ddots \end{matrix} & \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & 1 \end{bmatrix} & \\ & & \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & 1 \end{bmatrix} & \begin{bmatrix} -4 & 1 & & & \\ 1 & \ddots & \ddots & & \\ & \ddots & \ddots & 1 & \\ & & 1 & -4 & \end{bmatrix} \end{pmatrix}$$

Sparse Matrices

Observation

Two-dimensional Laplacian is no longer tridiagonal.

This will complicate both the data structures and algorithms.

Common data structures for sparse matrices

- ▶ Coordinate list
- ▶ Compressed sparse column (CSC)

Sparse Matrices

Coordinate list format

Three vectors i, j, v of length nnz (number of nonzeros) such that

```
A = zeros(n,n)
for k = 1:nnz
    A[i[k],j[k]] = v[k]
end
```

Example

$$A = \begin{pmatrix} 0.2 & & & & & \\ 0.6 & & & & & \\ & 0.6 & 1.0 & & & \end{pmatrix} \longrightarrow \begin{aligned} i &= (1 & 2 & 3 & 2 & 3)^T \\ j &= (1 & 1 & 2 & 3 & 3)^T \\ v &= (0.2 & 0.6 & 0.6 & 0.7 & 1.0)^T \end{aligned}$$

Properties

- ▶ Convenient to assemble sparse matrix.
- ▶ A bit wasteful since j contains many repeated entries.

Sparse Matrices

Compressed sparse column (CSC) format

Three vectors p, i, v with

$$\text{length}(p) == n+1, \quad \text{length}(i) == \text{length}(v) == \text{nnz}$$

such that

```
A = zeros(n,n)
for j = 1:n
    for k = p[j]:p[j+1]-1
        A[i[k],j] = v[k]
    end
end
```

CSC is the format most commonly used in practice.

Example

$$A = \begin{pmatrix} 0.2 & & & & \\ 0.6 & & 0.7 & & \\ & 0.6 & 1.0 & & \end{pmatrix} \longrightarrow \begin{array}{l} p = (\quad 1 \quad 3 \quad 4 \quad 6 \quad)^T \\ i = (\quad 1 \quad 2 \quad 3 \quad 2 \quad 3 \quad)^T \\ v = (\quad 0.2 \quad 0.6 \quad 0.6 \quad 0.7 \quad 1.0 \quad)^T \end{array}$$

Sparse Matrices

Sparse matrices in Julia

- ▶ Sparse matrix tools are provided by the `SparseArrays` package.
Type using `SparseArrays` before calling any of the functions listed below.
- ▶ Assemble a sparse matrix: `sparse(i,j,v)`
See section on coordinate lists above for meaning of `i,j,v`.
- ▶ Extract `i,j,v` from sparse matrix `A`: `i,j,v = findnz(A)`
- ▶ Sparse identity matrix: `sparse(I, (n,n))`
- ▶ Sparse matrix of zeros: `spzeros(n,n)`
- ▶ Convert to full matrix: `Matrix(A)`

Sparse Matrices

Assembling the 2d Laplacian matrix, the tedious way

See `laplacian_2d_tedious()` in `6_sparse_matrices.jl`.

Assembling the 2d Laplacian matrix, the clever way

2d Laplacian operator is given by $\Delta = \frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}$.

Let us imitate this in the discrete case: $\Delta_n = \Delta_n^{(2,1)} + \Delta_n^{(2,2)}$ where

$$\begin{aligned}\Delta_n^{(2,1)}(i_1 + n(i_2 - 1), j_1 + n(j_2 - 1)) &= \dots \\ &= (n+1)^2 \begin{cases} -2 & \text{if } |i_1 - j_1| = 0 \text{ and } |i_2 - j_1| = 0, \\ 1 & \text{if } |i_1 - j_1| = 1 \text{ and } |i_2 - j_1| = 0, \\ 0 & \text{otherwise,} \end{cases}\end{aligned}$$

$$\begin{aligned}\Delta_n^{(2,2)}(i_1 + n(i_2 - 1), j_1 + n(j_2 - 1)) &= \dots \\ &= (n+1)^2 \begin{cases} -2 & \text{if } |i_1 - j_1| = 0 \text{ and } |i_2 - j_1| = 0, \\ 1 & \text{if } |i_1 - j_1| = 0 \text{ and } |i_2 - j_1| = 1, \\ 0 & \text{otherwise.} \end{cases}\end{aligned}$$

Sparse Matrices

$$\Delta_n^{(2,1)} \propto \begin{pmatrix} \begin{bmatrix} -2 & 1 & & \\ 1 & \ddots & \ddots & \\ & \ddots & \ddots & 1 \\ & & 1 & -2 \end{bmatrix} & & \\ & \ddots & \\ & & \begin{bmatrix} -2 & 1 & & \\ 1 & \ddots & \ddots & \\ & \ddots & \ddots & 1 \\ & & 1 & -2 \end{bmatrix} \end{pmatrix}$$

Sparse Matrices

$$\Delta_n^{(2,2)} \propto \begin{pmatrix} \begin{bmatrix} -2 & & \\ & \ddots & \\ & & -2 \end{bmatrix} & \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix} & & \\ \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix} & \begin{matrix} \cdot & & \\ & \ddots & \\ & & \cdot \end{matrix} & \begin{matrix} \cdot & & \\ & \ddots & \\ & & \cdot \end{matrix} & \\ & \begin{matrix} \cdot & & \\ & \ddots & \\ & & \cdot \end{matrix} & \begin{matrix} \cdot & & \\ & \ddots & \\ & & \cdot \end{matrix} & \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix} & \\ & & \begin{bmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{bmatrix} & \begin{bmatrix} -2 & & \\ & \ddots & \\ & & -2 \end{bmatrix} \end{pmatrix}$$

Sparse Matrices

Kronecker product of matrices

$$A \otimes B := \begin{pmatrix} A[1, 1] B & \cdots & A[1, n] B \\ \vdots & \ddots & \vdots \\ A[n, 1] B & \cdots & A[n, n] B \end{pmatrix}$$

Hence, the 2d Laplacian $\Delta_n^{(2)}$ can be expressed in terms of 1d $\Delta_n^{(1)}$ as

$$\Delta_n^{(2)} = \Delta_n^{(1)} \otimes I + I \otimes \Delta_n^{(1)}.$$

See `laplacian_2d_clever()` in `6_sparse_matrices.jl` for how to do this in Julia.

General rule

If A, B are “one-dimensional” operators, then $A \otimes B$ is the “two-dimensional” operator which applies A in one dimension and B in the other dimension.

Be careful about which operator applies to which dimension.

Sparse Matrices

Kronecker product of matrices

$$a \otimes b := \begin{pmatrix} a[1] b \\ \vdots \\ a[n] b \end{pmatrix}$$

Theorem

$$(A \otimes B)(a \otimes b) = (Aa) \otimes (Bb)$$

Proof. Straightforward but tedious computations.

Sparse Matrices

Eigenvalues and -vectors of 2d Laplacian

Assume λ_ℓ, u_ℓ are eigenpairs of $\Delta_n^{(1)}$.

Then, eigenpairs of $\Delta_n^{(2)}$ are given by

$$\lambda_{\ell_1, \ell_2} := \lambda_{\ell_1} + \lambda_{\ell_2}, \quad u_{\ell_1, \ell_2} := u_{\ell_1} \otimes u_{\ell_2}.$$

Proof.

$$\begin{aligned} \Delta_n^{(2)} u_{\ell_1, \ell_2} &= \left(\Delta_n^{(1)} \otimes I + I \otimes \Delta_n^{(1)} \right) (u_{\ell_1} \otimes u_{\ell_2}) \\ &= (\Delta_n^{(1)} u_{\ell_1}) \otimes u_{\ell_2} + u_{\ell_1} \otimes (\Delta_n^{(1)} u_{\ell_2}) \\ &= \lambda_{\ell_1} u_{\ell_1} \otimes u_{\ell_2} + \lambda_{\ell_2} u_{\ell_1} \otimes u_{\ell_2} \\ &= (\lambda_{\ell_1} + \lambda_{\ell_2}) u_{\ell_1} \otimes u_{\ell_2} \\ &= (\lambda_{\ell_1} + \lambda_{\ell_2}) u_{\ell_1, \ell_2}. \end{aligned}$$