



## PulsarCast

## Scaling Pub-Sub over the distributed web

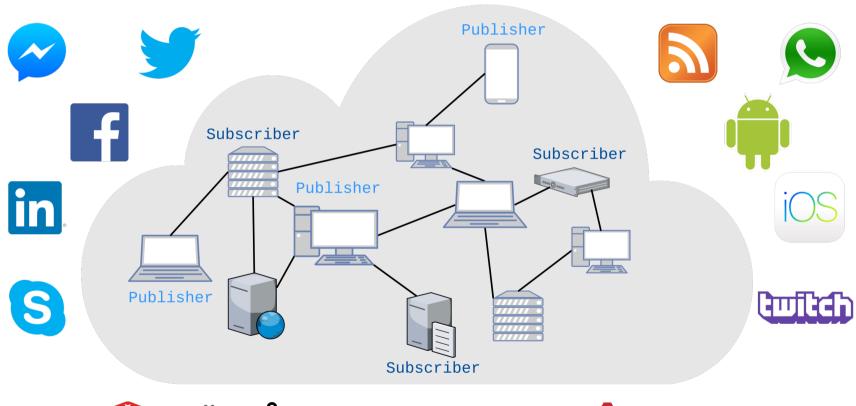
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#### Motivation

#### World Wide Web















## Pub-Sub Paradigm

- Communication paradigm providing full decoupling in:
  - Time
  - Space
  - Synchronisation





#### **Problems**

Lack of a pub-sub system that:

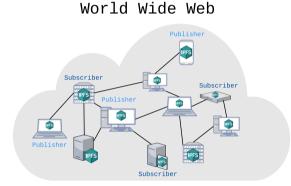
- Scales for the web
- Provides reliability:
  - Delivery guarantees
  - Data persistence





## Objectives

- Decentralised pub-sub architecture using IPFS<sup>1</sup>:
  - Highly scalable
  - Reliable
  - Assures persistence



- Develop a system that meets these requirements
- Evaluate its performance on multiple environments

[1] - https://ipfs.io/





#### Related Work

#### **Pub-Sub Systems**

# Subscribers Message Event e=<1,2,3,1,2>

#### Web Technologies







## **Design Dimensions**

- Subscription model
- Network architecture
- Overlay structure
- Subscription management
- Event dissemination





## **Subscription Model**

- Topic based
  - E.g. Redis<sup>2</sup>, Scribe[1], Tera[7], Poldercast[10]
- Type based
  - E.g. Hermes[4]
- Content based
  - E.g. Gryphon[8], Siena[9], Meghdoot[2]

[2] - https://redis.io/





## **Topic Based**

- A simple notion of group
- Members of the group receive every message
- Can build a complex hierarchy
- Lacks expressiveness

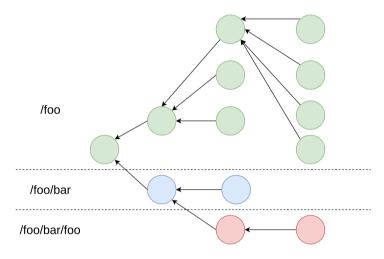


Fig.1: Events on a topic hierarchy

 Note: Type based can be seen as a special case of the topic based model fit for type based languages





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#### **Content Based**

- Filter events based on their content
- Support for rich and flexible subscriptions
- Requires complex filtering and message forwarding

```
Message

{
   exchange: "Euronext Lisboa",
   company: "CTT",
   order: "buy",
   number: "100",
   price: "5.55",
}

   Subscription

{
   exchange: "Euronext Lisboa",
   order: "buy",
   number: ">50",
   price: "<10",
}</pre>
```

Fig.2: JSON<sup>3</sup> subscription example

[3] - https://www.json.org/

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#### **Network Architecture**

- Centralised:
  - E.g. Redis, Kafka<sup>4</sup>, Gryphon
- Decentralised:
  - E.g. Scribe, Meghdoot, Poldercast

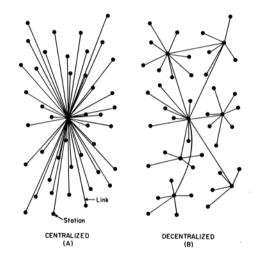


Fig.3: Network architecture overview

[4] - http://kafka.apache.org/documentation/#design





#### Centralised

- Focus:
  - Reliability (with replication)
  - Consistency
- Lacks:
  - Scalability
  - Data Throughput

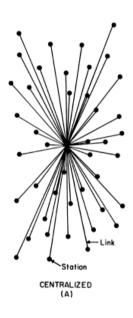


Fig.4: Centralised architecture example





#### Decentralised

- Focus:
  - Scalability
  - Data Throughput
- Lacks:
  - Reliability
  - Consistency
  - Persistence



Fig.5: Decentralised architecture example





## **Network Overlay**

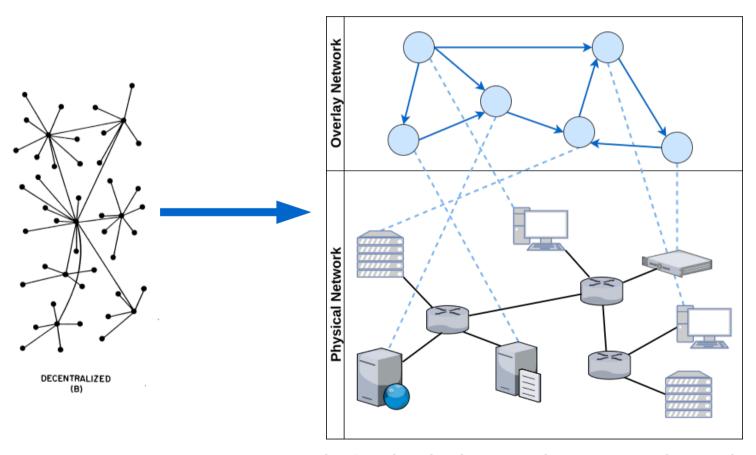


Fig.6: Physical network vs network overlay

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## **Unstructured Overlay**

- Each peer connects to a subset of nodes
- No clear structure or hierarchy
- Usage of **gossip based membership** protocols (e.g. Cyclon[11]) to help preserve:
  - Network diameter
  - Average degree





## Structured Overlay

- Peers organise according to a specific structure (e.g. ring, tree, multi-dimensional space).
- Common approach is to use a Distributed Hash Table
  - E.g. Chord[5], Kademlia[3], CAN[6]





## Structured Overlay - DHT

- Peer identifiers evenly spread across key space
- Ensures the content is evenly distributed
- Queries usually solved in logarithmic time

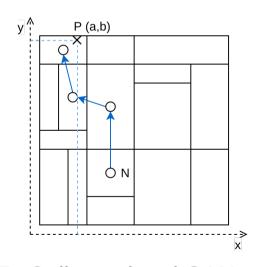


Fig.7: 2 dimensional CAN routing

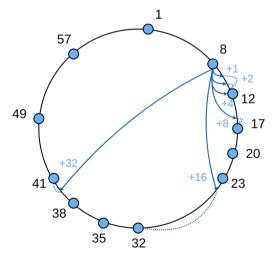


Fig.8: Chord ring





- Network mesh with state stored at each node
  - E.g. Siena
- Using rendezvous nodes and dissemination trees
  - E.g. Scribe
- Clustering of similar subscriptions
  - E.g. using multiple overlays, like Poldercast or Tera
- Usage of per attribute DHTs or multidimensional DHTs (for content based subscritpions)
  - E.g. Meghdoot, Mercury





#### **Event Dissemination**

Solutions heavily depend on subscription management

Matching intrinsically happens in the network

(distributed)

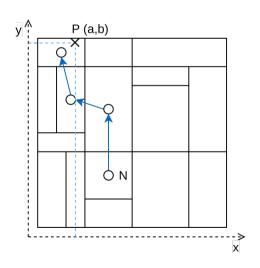


Fig.7: 2 dimensional CAN routing

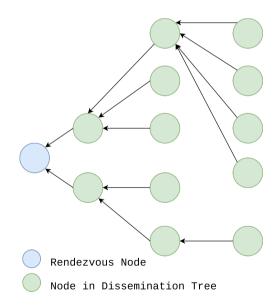


Fig.9: Dissemination tree for a topic







	Subscription Model	Architecture	Overlay Structure	Subscription Management	Event Dissemination
Gryphon	Content	(C) broker hierarchy	N/A	Each broker responsible for a subscription scheme	Tree hierarchy
Siena	Content	(C) broker mesh	N/A	Keep state at each node for subscription routes	Flood with cached state
Jedi	Content	(C) broker hierarchy	N/A	Keep state at each node for subscription routes	Tree hierarchy
Bayeux	Topic	Decentralised	Tapestry DHT	Rendezvous node	Multicast tree
Scribe	Topic	Decentralised	Pastry DHT	Rendezvous node	Multicast tree
Meghdoot	Content	Decentralised	CAN DHT	Point in CAN DHT	CAN routing
Hermes	Туре	Decentralised	Pastry DHT	Rendezvous node	Multicast tree
Tera	Topic	Decentralised	Gossip based overlay	Unstructured overlay per topic	Random walks & flooding
Mercury	Content	Decentralised	Ring based DHTs	Overlay per attribute in schema	Route through ring overlays
Sub-2-Sub	Content	Decentralised	Ring based DHT & gossip overlay	Clustering of similar subscriptions	Gossip & ring overlay routing
Poldercast	Topic	Decentralised	Ring based DHT, Vicinity & Cyclon	Ring overlay per topic	Ring overlay routing

(C) – Centralised

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	Relay Free Routing	Delivery Guarantees	Fault Tolerance
Gryphon	N/A	yes	Best effort
Siena	N/A	yes	Best effort
Jedi	N/A	yes	Best effort
Bayeux	No	yes	Best effort, no subscription persistence
Scribe	No	yes	Best effort, no subscription persistence
Meghdoot	No	yes	Replicated subscriptions
Hermes	No	yes	Best effort
Tera	No	no	Best effort
Mercury	No	yes	Best effort
Sub-2-Sub	No	no	Best effort
Poldercast	Yes	Yes (every publisher is a subscriber)	High resilience to churn, no subscription persistence

**Delivery Guarantees** – Event delivery guarantees under normal network conditions **Fault Tolerance** – Mechanisms to guarantee successful event delivery and subscription matching under heavy churn

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## Web Technologies

- The Javascript<sup>6</sup> ecosystem
- New network protocols that facilitate P2P communication
- New P2P applications that leverage all of this
  - E.g. IPFS

[5] - https://www.ecma-international.org/publications/standards/Ecma-262.htm





## The Javascript Ecosystem

 Javascript is one of the main programming languages for the web<sup>6</sup>



- Thanks to NodeJS<sup>7</sup>, now runs in servers also
- Its package manager, **NPM**<sup>8</sup>, is one of the largest package registries in the world



 NPM powers a UNIX-like culture of small modules that work well together

<sup>[6] -</sup> https://insights.stackoverflow.com/survey/2017

<sup>[7] -</sup> https://nodejs.org

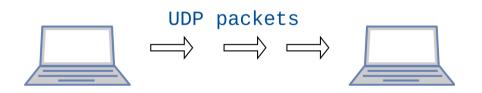
<sup>[8] -</sup> https://www.npmjs.com/





#### **New Protocols**

- **WebRTC**<sup>9</sup> made possible full-duplex communication between browsers without an intermediary.
- QUIC<sup>10</sup> and uTP<sup>11</sup> provided alternatives to TCP, bringing reliability and order delivery over UDP





<sup>[9] -</sup> https://www.w3.org/TR/webrtc/

<sup>[10] -</sup> https://datatracker.ietf.org/wg/quic/about/

<sup>[11] -</sup> http://www.bittorrent.org/beps/bep\_0029.html





#### **IPFS**

- A P2P hypermedia protocol designed to create a persistent, content-addressable network for the web
- Uses a modular approach through libp2p<sup>12</sup> to solve common challenges of P2P applications

• At its core, IPFS uses a Merkle DAG13



[12] - https://libp2p.io

[13] - https://github.com/ipld/specs/tree/master/ipld





#### Merkle DAG

- A graph structure used to store and represent linked data.
- Each node can be linked to based on the hash of its content.
- Referred to as IPLD in the IPFS ecosystem.
- Offers immutability.
- IPNS built to offer mutability





## **Proposed Solution**

- A topic based, pub-sub module for the IPFS ecosystem that provides:
  - Data persistence
  - Delivery guarantees
  - High scalability





#### Architecture

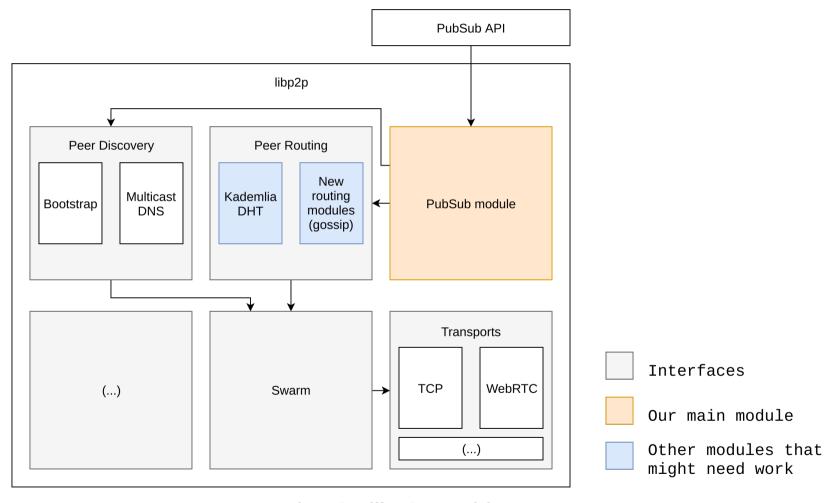


Fig.13: libp2p architecture

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## Libp2p Kademlia DHT

- put: insert a value with a given key.
- get: get a value of a given key.
- findPeer: find the peer with the given peer ID.
- getClosestPeers: find the k closest peers to a given key.
- provide: let the network know that this peer can also distribute a given key.
- findProviders: find providers for a given key.





#### Distributed Data Structures

#### **Topic Descriptor**

#### **Event Descriptor**

```
topic: {
   name: <topic-name>, // Name of the topic
   link: <merkle-link> // Link to the topic of this message
},
   publisher: <publisher node ID>
   parent: <merkle-link to previous event>,
   metadata: <json-object>, // Timestamp and other relevant info
   payload: <json-object>, // The actual message content
}
```

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#### Distributed Data Structures

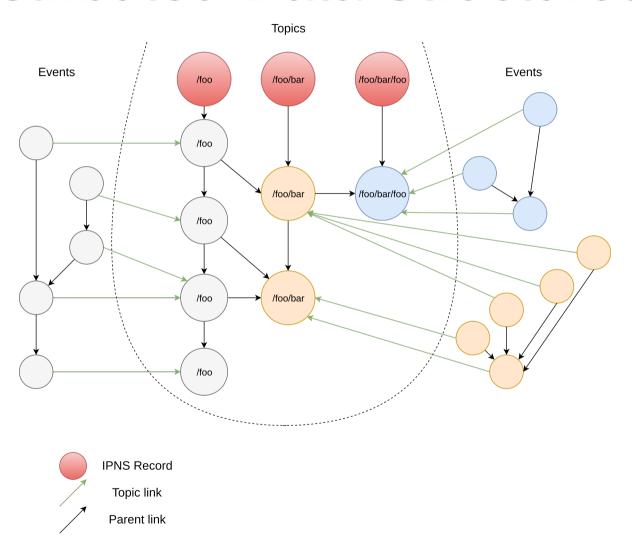


Fig.14: Overall data structures architecture





- Using the Kademlia DHT, build a tree structure when creating a new subscription
  - Node checks if it is part of the dissemination tree
  - If not:
    - Issue a query for the closest nodes to the topic ID
    - Select the closest
    - Issue a special command for it to join the tree
    - Receiving node registers initiator as its child in the tree and repeats the process





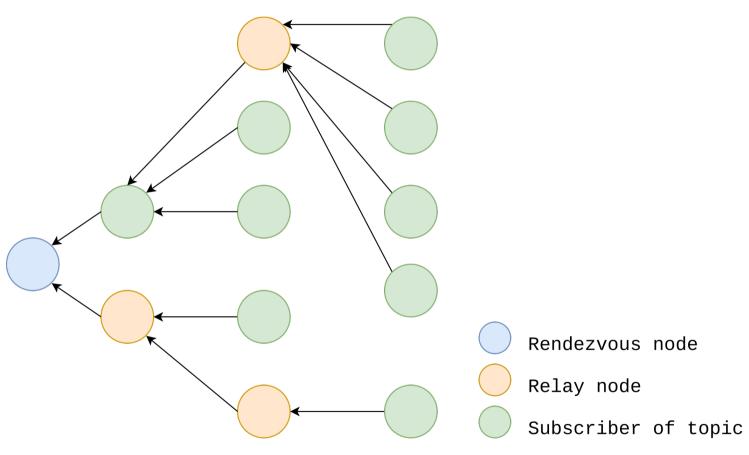


Fig.15: Dissemination tree for a topic in Pulsarcast

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 Upon successfully subscribing to a topic a peer will request for all the leaf events from the event stream

This will allow him to rebuild the graph by requesting

the parent identifiers

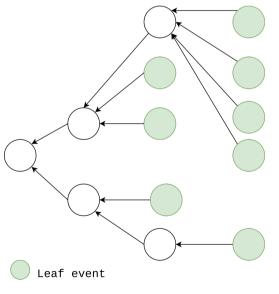


Fig.16: Leaf events in a event stream





## **Event Dissemination**

- When a new event is created by a node:
  - Issue a Kademlia put command to persist the event in the network
  - If node is part of the tree send it through its child and parent links
  - Else, send it to the rendezvous node which will disseminate it through the tree
- At each node, the event will be persisted with a systemwide probability p





## Quality of Service

- Replicate rendezvous nodes using the DHT getClosestNodes and provide methods. Synchronise them using a simple gossip overlay
- Missing events in a stream can be solved by querying parent links
- The persistence of data across the network allows load to be evenly distributed while giving us fault tolerant capabilities
- Note: The IPNS bottleneck can be considered a matter of ownership





#### **Evaluation**

- Using either PeerSim<sup>14</sup> or IPTB<sup>15</sup>
- Using synthetic datasets of both events and subscriptions:
  - Of different sizes
  - With different distributions
- Using our system, IPFS current pub-sub and a baseline

<sup>[14] -</sup> http://peersim.sourceforge.net/

<sup>[15] -</sup> https://github.com/whyrusleeping/iptb





#### **Evaluation - Metrics**

- Ratio of messages sent by each node (correlated with the CPU, memory and bandwidth usages).
- Ratio of throughput speedup vs disk storage used at each node.
- Ratio between latency reduction vs disk storage used at each node.
- Ratio of subscriptions covered under heavy network churn.
- Ratio of subscriptions covered after a severe network partition and its recovery.
- State of event streams at each node under heavy network churn.
- State of event streams at each node after a network partition and its recovery.





#### Conclusion

- Pub-sub systems are vital for distributed applications
- Lack of a system that:
  - Scales for the web
  - Focuses persistence and reliability
- Overview of relevant pub-sub systems and web technology
- Proposed solution of a decentralised pub-sub system for the web using IPFS





#### Thank you for your presence!

## Questions?

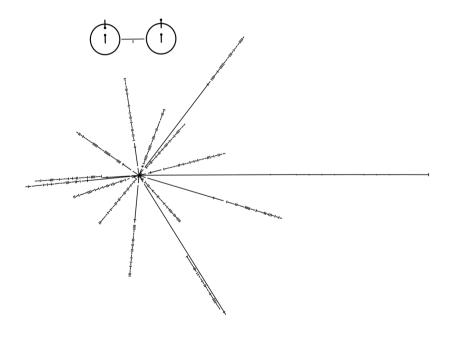


Fig.17: Pulsar map used for the pioneer plaque





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#### Merkle DAG

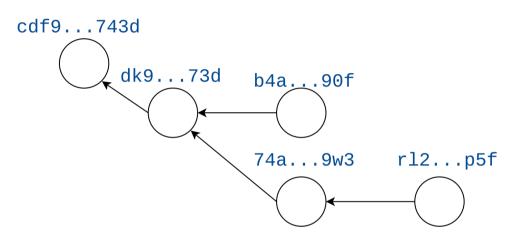


Fig.10: Example of a Merkle DAG

```
{
  "content": {
    "key": "value"
  },
    "merkle-link": "cdf9...743d"
}
Hash(...)

dk9...73d
```

Fig.11: Merkle Node creation

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### **IPNS**

A way to offer mutability.

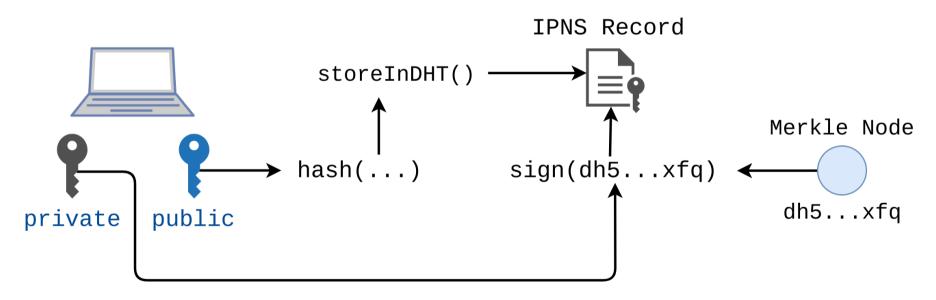


Fig.12: IPNS record overview