

Intro to Algorithms: Homework #10

Due on April 28, 2021

Prof. Zaki

Jared Gridley

Lab Results:

Results:

g50:

```
cut edges ('0', '103', 57), ('0', '1073', 14), ('0', '1083', 39), ('0', '1158', 3), ('0',  
'1187', 13), ('0', '1219', 6), ('0', '1262', 55), ('0', '1318', 4), ('0', '1336', 29),  
( '0', '1343', 18), ('0', '1363', 19), ('0', '1417', 21), ('0', '1433', 18), ('0', '1468',  
26), ('0', '1484', 25), ('0', '1495', 31), ('0', '150', 9), ('0', '1512', 38), ('0',  
'160', 50), ('0', '1639', 8), ('0', '1688', 9), ('0', '1697', 11), ('0', '173', 7), ('0',  
'1754', 16), ('0', '1778', 11), ('0', '1783', 1), ('0', '1799', 24), ('0', '189', 23),  
( '0', '1901', 7), ('0', '1964', 33), ('0', '197', 40), ('0', '2021', 7), ('0', '2049',  
3), ('0', '210', 3), ('0', '2139', 6), ('0', '214', 15), ('0', '2183', 2), ('0', '221',  
17), ('0', '2249', 31), ('0', '2399', 18), ('0', '2413', 15), ('0', '2473', 14), ('0',  
'2509', 6), ('0', '276', 21), ('0', '329', 3), ('0', '384', 26), ('0', '404', 4), ('0',  
'451', 7), ('0', '520', 19), ('0', '529', 2), ('0', '543', 16), ('0', '597', 15), ('0',  
'6', 50), ('0', '605', 23), ('0', '647', 4), ('0', '67', 16), ('0', '676', 2), ('0',  
'68', 2), ('0', '695', 8), ('0', '715', 40), ('0', '741', 8), ('0', '762', 13), ('0',  
'818', 10), ('0', '821', 35), ('0', '86', 58), ('0', '891', 1), ('0', '956', 33), ('0',  
'966', 6)
```

maxflow value 1224

g5:

```
flow ('0', '24', 59), ('0', '25', 7), ('0', '31', 44), ('0', '32', 3)
```

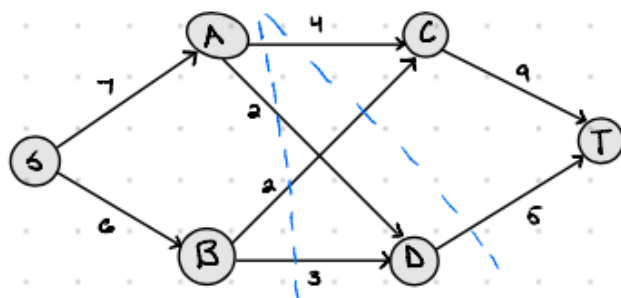
maxflow value 113

g10:

```
flow ('0', '103', 2), ('0', '126', 5), ('0', '136', 12), ('0', '138', 15), ('0', '158',  
43), ('0', '159', 35), ('0', '184', 2), ('0', '199', 55), ('0', '20', 63), ('0', '205',  
1), ('0', '218', 3), ('0', '244', 16), ('0', '249', 23), ('0', '266', 8), ('0', '276',  
1), ('0', '289', 28), ('0', '290', 27), ('0', '297', 19), ('0', '305', 48), ('0', '323',  
23), ('0', '343', 8), ('0', '360', 32), ('0', '368', 22), ('0', '381', 10), ('0', '387',  
2), ('0', '406', 62), ('0', '415', 71), ('0', '416', 4), ('0', '429', 11), ('0', '43',  
4), ('0', '435', 22), ('0', '442', 16), ('0', '444', 28), ('0', '56', 12), ('0', '58',  
3), ('0', '60', 19), ('0', '68', 9), ('0', '8', 1), ('0', '88', 14)
```

maxflow value 779

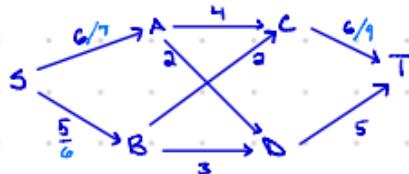
7.17) Network:



a) Max Flow and min cut

1) Find Shortest path

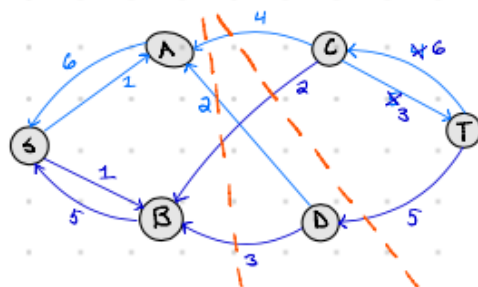
2) Increase flow along the path



Max Flow: 11 units

min cut: Diagonal through $A \rightarrow C$, $B \rightarrow C$, $D \rightarrow T$

b) Residual Graph G_f (with edge capacities). Mark vertices reachable from S and T.



All nodes are reachable from T.
Nodes A and B are reachable from S.

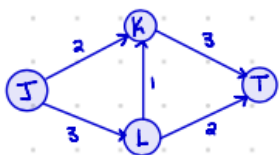
c) Bottleneck Edges

Edge $A \rightarrow C$
 $B \rightarrow C$

($B \rightarrow D$, $D \rightarrow T$) also a bottleneck, but both are required to increase max flow

Figure 1: Page 1

d) Simple Example without Bottlenecks



e) Algorithm to Identify Bottlenecks

Bottlenecks (G):
 $F(e) = 0 \quad \forall e \in G$
 while \exists a path from s to t in G^p
 P : Shortest path from $s \rightarrow t$ (BFS)
 $c(P)$: capacity of P
augment the flow along P
 recompute G
 recompute G^p

Edmunds-Karp
 Algorithm
 $O(|V| + |E|)$

+

Run BFS on (G^p, s) , but only on residual edges (no back edges)

A = nodes reachable from s

$O(|V| + |E|)$

Run BFS on (G^p, t) , but only with residual edges

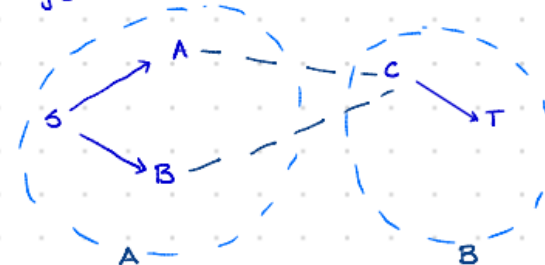
B = nodes reachable from t

For each edge $(u, v) \in G$:

if $(u \in A \text{ and } v \in B)$:

(u, v) is a bottleneck

Break the graph into only the residual edges:



Any edges that cross between A and B in the original graph are bottlenecks.

Total time
 $O(|V| + |E|^2)$

Figure 2: Page 2

7.21) Give efficient Algorithm to Find critical edges.

critical edge \rightarrow edge that exhausts its capacity

\hookrightarrow Edmond-Karp Algorithm

$\hookrightarrow G^f \rightarrow$ traverse from t , find all edges without a back edge \Rightarrow Critical

Critical Edges (G, s):

Edmond-Karp
 $O((V+E)^2)$

```

f(e) = 0  $\forall e \in G$ 
while  $\exists$  a path from  $s$  to  $t$  in  $G^f$ 
  P: Shortest path from  $s \rightarrow t$  (BFS)
  c(P): capacity of P
  augment the flow along P
  recompute  $G$ 
  recompute  $G^f$ 

```

Run modified BFS on G^f , starting from T :

critical = []

Q = [T] // Priority Queue

prev = T

```

while Q not empty:
  u = eject Q
  backedge = False
  for all  $v \in N(u)$ :
    if not visited:
      add v to Q, set visited flag
    if  $v == prev$ :
      Add (u,v) to critical
  prev = u

```

return critical

BFS
 $O(V)$

Total Time:
 $O((V+E)^2)$

Figure 3: Page 3