

Algebra III

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Chapter 0: Review

Definition: Category

A category \mathcal{C} consists of the following data:

1. A class of objects, $\text{Obj}(\mathcal{C})$.
2. For any pair of objects $X, Y \in \text{Obj}(\mathcal{C})$, a set of morphisms $\text{Mor}_{\mathcal{C}}(X, Y)$, $\text{Hom}_{\mathcal{C}}(X, Y)$ or $\mathcal{C}(X, Y)$.
3. For any triple of objects $X, Y, Z \in \text{Obj}(\mathcal{C})$, a map

$$\begin{aligned} \text{Hom}_{\mathcal{C}}(Y, Z) \times \text{Hom}_{\mathcal{C}}(X, Y) &\rightarrow \text{Hom}_{\mathcal{C}}(X, Z) \\ (g, f) &\mapsto g \circ f \end{aligned}$$

called compositions subject to the following axioms:

1. Associativity: $f \circ (g \circ h) = (f \circ g) \circ h$ whenever this makes sense.
2. For every object $X \in \text{Obj}(\mathcal{C})$, there exists a morphism $\text{id}_X \in \text{Hom}_{\mathcal{C}}(X, X)$ such that

$$\text{id}_X \circ f = f \quad \text{and} \quad g \circ \text{id}_X = g, \quad \forall f \in \text{Hom}_{\mathcal{C}}(W, X), g \in \text{Hom}_{\mathcal{C}}(X, W)$$

Example 1

Let E be a set (or a class).

Define \mathcal{C} by taking $\text{Obj}(\mathcal{C}) = E$ and $\text{Hom}_{\mathcal{C}}(X, Y) = \begin{cases} \emptyset & \text{if } x \neq y \\ \{\text{id}_X\} & \text{if } x = y \end{cases}$.

Example 2

Let $\mathcal{C} = \text{Set}$ the category of all sets with set functions acting as morphisms.

Let $\mathcal{C} = \text{Grp}$ the category of all groups with group homomorphisms acting as morphisms.

Abelian Rings: Ab, Rings: Ring, Commutative Rings: CRing, Vector Spaces over F : Vect_F , Topological Spaces: Top, etc.

Example 3

Let G be a group (or more generally a monoid).

Define $\text{Obj}(\mathcal{C}) = \{*\}$, $\text{Hom}_{\mathcal{C}}(*, *) = G$ and

$$\text{Hom}_{\mathcal{C}}(*, *) \times \text{Hom}_{\mathcal{C}}(*, *) \rightarrow \text{Hom}_{\mathcal{C}}(*, *)$$

the group operator.

Example 4

Let (E, \leq) be a preordered set (i.e. reflexive and transitive).
Define \mathcal{C} by $\text{Obj}(\mathcal{C}) = E$,

$$\text{Hom}_{\mathcal{C}}(x, y) = \begin{cases} \emptyset & \text{if } x \not\leq y \\ \{f_{xy}\} & \text{if } x \leq y \end{cases}$$

Notation

If $f \in \text{Hom}_{\mathcal{C}}(X, Y)$ we write $X \xrightarrow{f} Y$ in \mathcal{C} .

Definition: Isomorphism

A morphism $f : X \rightarrow Y$ in \mathcal{C} is an isomorphism if $\exists g : Y \rightarrow X$ such that $g \circ f = \text{id}_X$ and $f \circ g = \text{id}_Y$.

Definition: Endomorphism

A morphism on X with $f : X \rightarrow X$.

Definition: Automorphism

An automorphism on X is just an isomorphism $f : X \xrightarrow{\sim} X$ from X to itself.
Note that $\text{Aut}_{\mathcal{C}}(X) \subseteq \text{End}_{\mathcal{C}}(X) = \text{Hom}_{\mathcal{C}}(X, X)$.

Remark:

The collection of all endomorphisms on X form a monoid.
The collection of all automorphisms on X forms a group called the automorphism group of X .

Example 1

Let $\mathcal{C} = \text{Set}$, $X = \{1, \dots, n\}$. Then $\text{Aut}_{\text{Set}}(\{1, \dots, n\}) = \text{Perm}(X) = S_n$.

Example 2

Let $\mathcal{C} = \text{Vect}_F$, $X = F^n$. Then $\text{Aut}_{\text{Vect}_F}(F^n) = \text{GL}_n(F)$.

Definition: Functors

Let \mathcal{C} and \mathcal{D} be categories.
A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ from \mathcal{C} to \mathcal{D} consists of the following data

1. For each object $X \in \text{Obj}(\mathcal{C})$, a chosen object $F(X) \in \text{Obj}(\mathcal{D})$.
2. For each pair of objects $X, Y \in \text{Obj}(\mathcal{C})$, a function

$$\begin{aligned} \text{Hom}_{\mathcal{C}}(X, Y) &\rightarrow \text{Hom}_{\mathcal{D}}(F(X), F(Y)) \\ f &\mapsto F(f) \end{aligned}$$

such that

1. For any two composable morphisms $X \xrightarrow{f} Y \xrightarrow{g} Z$ in \mathcal{C} , we have $F(g \circ f) = F(g) \circ F(f)$.
2. For each object $X \in \text{Obj}(\mathcal{C})$, $F(\text{id}_X) = \text{id}_{F(X)}$.

Example 1

For $\mathcal{D} := \mathcal{C}$, $\text{Id} : \mathcal{C} \rightarrow \mathcal{C}$, $X \mapsto X$, $f \mapsto f$.

Example 2: Forgetful Functors

$\mathcal{U} : \text{Grp} \rightarrow \text{Set}$ given as $(G, \cdot) \mapsto G$.

$\text{Ring} \rightarrow \text{Ab}$ given as $(R, +, \cdot) \mapsto (R, +)$.

Example 3: Tensors

Let R be a commutative ring, $M \in \text{Mod}_R$.

Then $\otimes_R M : \text{Mod}_R \rightarrow \text{Mod}_R$ and $\text{Hom}_R(M, -) : \text{Mod}_R \rightarrow \text{Mod}_R$.

Definition:

Let X be an object in a category \mathcal{C} and G a group.

An action of G on X is a group homomorphism $G \rightarrow \text{Aut}_{\mathcal{C}}(X)$.

Example 1

Let $\mathcal{C} = \text{Set}$.

A G -set is a set $X \in \text{Set}$ equipped with a group homomorphism

$$G \rightarrow \text{Perm}(X) = \text{Aut}_{\text{Set}}(X)$$

Exercise 1

A G -set is the same thing as a functor $G \rightarrow \text{Set}$, $* \mapsto X$, $\text{Hom}_{\mathcal{C}}(*, *) \rightarrow \text{Hom}_{\text{Set}}(X, X)$ ($G \rightarrow \text{Aut}_{\text{Set}}(X)$).

Definition: Adjunctions

Let \mathcal{C} and \mathcal{D} be categories and $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{D} \rightarrow \mathcal{C}$ be functors.

We say that F is left adjoint to G (and that G is right adjoint to F , and that we have a pair of adjoint functors) if for each object $X \in \text{Obj}(\mathcal{C})$ and $Y \in \text{Obj}(\mathcal{D})$, we have a bijection

$$\text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

which is “natural in X and Y ”:

For any $f : X \rightarrow X'$ in \mathcal{C} ,

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X'), Y) & \xrightarrow{\sim} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \\ - \circ F(f) \downarrow & & \downarrow - \circ f \\ \text{Hom}_{\mathcal{D}}(F(X), Y) & \xrightarrow{\sim} & \text{Hom}_{\mathcal{C}}(X, G(Y)) \end{array}$$

and for every $g : Y \rightarrow Y'$ in \mathcal{D}

$$\begin{array}{ccc} \mathrm{Hom}_{\mathcal{D}}(F(X), Y) & \xrightarrow{\sim} & \mathrm{Hom}_{\mathcal{C}}(X, G(Y)) \\ g \circ - \downarrow & & \downarrow G(g) \circ - \\ \mathrm{Hom}_{\mathcal{D}}(F(X), Y') & \xrightarrow{\sim} & \mathrm{Hom}_{\mathcal{C}}(X, G(Y')) \end{array}$$

We write

$$\begin{array}{c} \mathcal{C} \\ F \updownarrow G \\ \mathcal{D} \end{array}$$

Example 1

For $M \in \mathrm{Mod}_R$ we have

$$\begin{array}{c} \mathrm{Mod}_R \\ - \otimes_R M \updownarrow \mathrm{Hom}_R(M, -) \\ \mathrm{Mod}_R \end{array}$$

where

$$\begin{aligned} \mathrm{Hom}_R(M_1 \otimes M_2, N) &\cong \mathrm{Hom}_R(M_1, \mathrm{Hom}_R(M, \mathrm{Hom}_R(M_2, N))) \\ f &\mapsto (x \mapsto (y \mapsto f(x \otimes y))) \end{aligned}$$

Example 2

Let $R \xrightarrow{\phi} S$ be a ring homomorphism.

We can regard an S -module N as an R -module via

$$r \cdot x := \phi(r)x, \quad \forall r \in R, x \in N$$

This defines a functor $\mathrm{Mod}_S \rightarrow \mathrm{Mod}_R$ called a “restriction of scalars”, which has a left adjoint called “extension of scalars.”

$$\begin{array}{c} \mathrm{Mod}_R \\ S \otimes_R - \updownarrow \\ \mathrm{Mod}_S \end{array}$$

Recall

For commutative ring R , $\sim \mathrm{Mod}_R$.

e.g. $R = F$ a field, $\mathrm{Mod}_R \equiv \mathrm{Vect}_F$; $R = \mathbb{Z}$, $\mathrm{Mod}_R \equiv \mathrm{Ab}$.

Definition: R-Algebra

An R -algebra is an Abelian group $(A, +)$ that has both the structure of

1. an R -module and

2. a ring

which are compatible in that

$$r(ab) = (ra)b = a(rb), \quad \forall r \in R, a, b \in A$$

Example 1

The polynomial ring $R[x]$ is an R -algebra.

Example 2

The ring of $n \times n$ matrices $M_n(R)$ is an R -algebra.

Example 3

If $R \xrightarrow{\phi} S$ is a homomorphism of commutative rings, then S is an R -algebra via $r := \phi(r)a$, $\forall r \in R, a \in S$.

Example 4

$\mathbb{R} \hookrightarrow \mathbb{C}$. So \mathbb{C} is an \mathbb{R} -algebra.

$R \hookrightarrow R[x]$.

More generally, $R[x_1, x_2, \dots, x_n]$ is an R -algebra.

Commutative R-Algebras

An R -algebra is commutative if it is commutative as a ring.

$\text{CAlg}_R \subset \text{Alg}_R$.

Question: Why are polynomials important?

An algebraic perspective: they are the “free commutative algebras.”

Recall

For R a commutative ring, we have the notion of a free R -module – one that admits a basis.

Categorically, we have an adjunction.

$$\begin{array}{c} \text{Set} \\ f \updownarrow \mathcal{U} \\ \text{Mod}_R \end{array}$$

The left adjoint of the forgetful functor sends a set I to the free R -module with basis I .

$$F(I) = R^{(I)} = \bigoplus_{i \in I} R$$

The adjunction says that for any set I and R -module M ,

$$\begin{aligned} \text{Hom}_{\text{Mod}_R}(R^{(I)}, M) &\xrightarrow{\sim} \text{Hom}_{\text{Set}}(I, M) \\ \exists! R\text{-linear map } f: R^{(I)} &\rightarrow M \quad \leftarrow \{x_i\}_{i \in I} \\ e_i &\mapsto x_i \end{aligned}$$

Similarly, the forgetful functor $\mathcal{U} : \mathbf{CAlg}_R \rightarrow \mathbf{Set}$ has a left adjoint

$$\begin{array}{c} \mathbf{Set} \\ f \uparrow \downarrow \mathcal{U} \\ \mathbf{CAlg}_R \end{array}$$

which sends a set I to the “free commutative R -algebra on I .”

Explicitly, $F(I) = R[\{x_i\}_{i \in I}]$ the polynomial algebra with an indeterminate x_i for each $i \in I$.

Example 1

$$I = \{*\} \rightsquigarrow F(\{*\}) = R[x].$$

$$I = \{1, \dots, n\} \rightsquigarrow F(\{1, \dots, n\}) = R[x_1, \dots, x_n].$$

$$I = \mathbb{N} \rightsquigarrow F(\mathbb{N}) = R[x_1, x_2, \dots].$$

Adjunction

For any set I and commutative R -algebra $A \in \mathbf{CAlg}_R$, we have a bijection

$$\begin{aligned} \mathrm{Hom}_{\mathbf{CAlg}_R}(R[\{x_i\}_{i \in I}], A) &\cong \mathrm{Hom}_{\mathbf{Set}}(I, A) \\ \exists! R\text{-algebra homomorphism } R[\{x_i\}_{i \in I}] &\rightarrow A \leftarrow \{a_i\}_{i \in I} \\ &\quad x_i \mapsto a_i \end{aligned}$$

Example 1

Let A be a commutative R -algebra.

For any $a \in A$, there exists a unique R -algebra homomorphism $R[x] \rightarrow A$ which sends $X \mapsto a$.

Explicitly, $f(x) \mapsto f(a)$.

Corollary

Let $R \xrightarrow{\phi} S$ be a homomorphism of commutative rings.

For any $a \in S$, there is a unique ring $R[x] \xrightarrow{\bar{\phi}} S$ such that $\bar{\phi}|_R = \phi$ and $\bar{\phi}(X) = a$.

Example 1

Let $R \subseteq S$ be a subring.

For each $a \in S$, there is a unique ring homomorphism $R[x] \xrightarrow{\phi} S$ such that $\phi|_R = \mathrm{id}$ and $\phi'(X) = a$.

We call this the “evaluation at a .”

$$\begin{aligned} R[x] &\xrightarrow{\mathrm{ev}_a} S \\ f &\mapsto f(a) \end{aligned}$$

Definition: Subalgebra

Let A be a commutative R -algebra, and let $S \subset A$ be a subset.

The subalgebra of A generated by S , denoted $R[S]$, is the intersection of all subalgebras of A which contain S .

Explicitly,

$$R[S] = \{a \in A : \exists n \geq 1, s_1, \dots, s_n \in S, f \in R[x_1, \dots, x_n], a = f(s_1, \dots, s_n)\}$$

Example 1

Let $A = R[x]$. Then $A = R[x]$. That is, A is generated by $\{x\}$ as an algebra.

Similarly, $R[x_1, \dots, x_n]$ is generated as an algebra by $\{x_1, \dots, x_n\}$.

Example 2

If $R[x]/I$ with $I \subset R[x]$ an ideal, and $x := \overline{X} \in A$, then $A = R[x]$. That is, A is generated by $x = \overline{X}$ as an algebra.

More generally, if $I \subset R[x_1, \dots, x_n]$ an ideal, then $R[x_1, \dots, x_n]/I$ is generated by $\{\overline{x}_1, \dots, \overline{x}_n\}$.

Proposition

If $A \in \text{CAlg}_R$ is a finitely generated, commutative R -algebra, then $A \cong R[x_1, \dots, x_n]/I$ for some $n \geq 1$ and ideal $I \subset R[x_1, \dots, x_n]$.