

Birthday Gift - Editorial

Difficulty:

Medium - Hard

Prerequisites:

Strings

Hashing - ([Tutorial](#))

Manacher's Algorithm - ([Tutorial](#))

Suffix Array - ([Tutorial](#))

Binary Search - ([Tutorial](#))

Problem in Brief:

You are given a string S. You have to count the number of odd length substrings of S that are also palindromes.

Editorial:

There are 3 ways to solve this problem :-

- 1) Using Hashing
- 2) Using Manacher's Algorithm
- 3) Using Suffix Array

I will explain the Suffix Array algorithm as it is most intuitive out of all 3 ways.

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Consider an odd length palindromic substring of S. It must have its center at one of the indices of S. Instead of counting the substrings ad-hoc we will count the odd length palindromic substrings with its center at index i for all $1 \leq i \leq |S|$.

The count of odd length palindromic substrings with center at i is equal to precisely $(\text{MAX}(i) + 1) / 2$.

Where $\text{MAX}(i)$ = length of longest odd length palindromic substring with its center at i.

To find $\text{MAX}(i)$ we will build the Suffix Array on the concatenation of S and reverse of S.

We will then use binary search to find the max odd length substring using Longest Common Prefix from the Suffix Array DS.

We will then sum $(\text{MAX}(i) + 1) / 2$ for $1 \leq i \leq |S|$

And output the answer.

Time Complexity:

Building the suffix array is the costliest operation. Binary search will only take $\log(N)$ time per index. Hence the Time Complexity is

$O(N * \log^2(N))$

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