

# ***Candyland***

## **Difficulty:**

Hard

## **Prerequisites:**

Divide and Conquer([Theory](#))

## **Problem in Brief:**

Given a tree where each node has a value associated with it. For each node, find the value that appears maximum number of times in its subtree.

## **Editorial:**

To solve this problem, we will make use of the divide and conquer strategy known as smaller to larger.

Consider the example problem:

Given a tree, every vertex has color. Query is how many vertices in subtree of vertex  $v$  are colored with color  $c$ ?

First, we have to calculate the size of the subtree of every vertex. It can be done with simple dfs:

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```
int sz[maxn];  
void getsz(int v, int p){  
    sz[v] = 1; // every vertex has itself in its  
    subtree  
    for(auto u : g[v])  
        if(u != p){  
            getsz(u, v);  
            sz[v] += sz[u]; // add size of child u to  
            its parent(v)  
        }  
}
```

Now we have the size of the subtree of vertex  $v$  in  $sz[v]$ .

```
map<int, int> *cnt[maxn];  
void dfs(int v, int p){  
    int mx = -1, bigChild = -1;  
    for(auto u : g[v])  
        if(u != p){  
            dfs(u, v);  
            if(sz[u] > mx)  
                mx = sz[u], bigChild = u;  
        }  
    if(bigChild != -1)
```

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```
    cnt[v] = cnt[bigChild];
else
    cnt[v] = new map<int, int> ();
    (*cnt[v])[ col[v] ] ++;
    for(auto u : g[v])
        if(u != p && u != bigChild){
            for(auto x : *cnt[u])
                (*cnt[v])[x.first] += x.second;
        }
    //now (*cnt[v])[c] is the number of vertices in subtree of vertex v
    that has color c. You can answer the queries easily.

}
```

The above code is the use of smaller to larger technique. Each color in any node can travel upward upto  $\log(N)$  times as when it moves to a larger map the size becomes twice at the very least.

We can easily modify this code to solve our problem.

## Time Complexity:

We have seen that each element at any node may move up at most  $\log(N)$  times and it takes  $\log(N)$  time for each insertion as we use a map. Hence the Time Complexity is

**$O(N * \log^2(N))$**

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