

# Save Energy - Editorial

## Difficulty:

Medium

## Prerequisites:

Dynamic Programming - ([Tutorial](#))

Greedy Algorithm - ([Tutorial](#))

## Problem in Brief:

There are  $N$  towns in a line. Distance between town  $i$  and  $j$  is  $\text{abs}(j-i)$ . We want to go from town 0 to town  $N-1$ . From town  $i$ , we can jump to any town  $j > i$  and that costs  $(j-i)A[i] + (j^2 - i^2)A[i]^2$ , where  $A$  is given as input. We want to find the minimum cost needed.  $N \leq 10^5$ .

## Editorial:

There are 2 solutions. First solution using Greedy approach:

If you write cost function for path  $i \rightarrow j \rightarrow k$  and compare it with the cost function for path  $i \rightarrow k$  you can see that the former is better (costs less energy) if and only if either  $|A[j]| < |A[i]|$  or  $|A[j]| = |A[i]|$  with  $A[i] > 0$ .

From this observation it is quite easy to see that the optimal path would be greedy jumping from 0 to first  $i$  such that  $|A[i]| < |A[0]|$  or  $A[i] = -A[0]$  (if  $A[0] > 0$ ) or directly to  $n-1$  if there is no such  $i$  and then doing the same for  $i$  till we reach  $n-1$ . Jumping greedily to  $n-1$  solves the problem with  $O(N)$  complexity.

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Second solution using Dynamic Programming:

If we look at the nature of cost function it is easy to see that we can split each jump into consecutive steps, we can also solve it in more straightforward way — by computing a  $DP[i][j]$  which will tell us smallest cost to reach position  $i$  while using some  $A[x] = j$ . Now from position  $i$  we can either keep moving forward using same "old" value of  $A$ , or start using  $A[i]$ . Such solution works in  $O(N * 10^3)$  where  $10^3$  is count of distinct possible values of  $A[i]$ .

## **Time Complexity:**

The greedy algorithm passes the array once to find the answer. Hence the Time Complexity is

**$O(N)$**

## **Similar Problems:**

[First](#)

[Second](#)