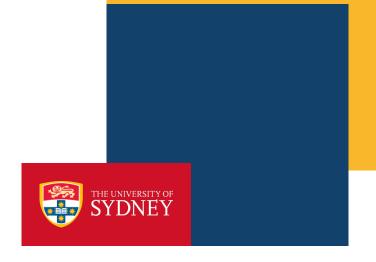
COMP9120

Week 6: Advanced SQL

Semester 1, 2025



Professor Athman Bouguettaya School of Computer Science

Warming up





Acknowledgement of Country

I would like to acknowledge the Traditional Owners of Australia and recognise their continuing connection to land, water and culture. We are currently on the land of the Gadigal People of the Eora nation and pay our respects to their Elders, past, present and emerging.





COMMONWEALTH OF AUSTRALIA

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New In-Person/Zoom Drop-in Helpdesk Sessions

- Due to the very low attendance of live ed sessions, I am replacing them with a combination of live and zoom drop-in helpdesk sessions. The new drop-in sessions started yesterday! The schedule can be found on Canvas @ https://canvas.sydney.edu.au/courses/63042/files/42180612? module item id=2681519
- Note the (online) zoom session links can be accessed following the Zoom link on canvas.

	Week	Date	Weekday	FROM	TO	Name	Туре	Room
	6	2-Apr-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	6	3-Apr-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	6	4-Apr-25	FRI	10am	12pm	Dipankar	in-person	J02.03.315.PNR Building.PNR Learning Studio 315
	6	2-Apr-25	WED	10am	11am	Abbey	online	
	6	2-Apr-25	WED	8pm	9pm	Abbey	online	
	6	3-Apr-25	THUR	3pm	4pm	Abbey	online	
	6	4-Apr-25	FRI	10am	11am	Abbey	online	
ſ	7	9-Apr-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	7	11-Apr-25	FRI	10am	12pm	Dipankar	in-person	J02.03.315.PNR Building.PNR Learning Studio 315
Ī	8	16-Apr-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	8	17-Apr-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
ſ	9	30-Apr-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	9	30-Apr-25	WED	8pm	9pm	Abbey	online	
	9	1-May-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	9	2-May-25	FRI	10am	12pm	Dipankar	in-person	F07.03.352.Carslaw Building.Carslaw Learning Studio 352
ſ	10	7-May-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	10	8-May-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	10	9-May-25	FRI	10am	12pm	Dipankar	in-person	F07.03.352.Carslaw Building.Carslaw Learning Studio 352
	10	7-May-25	WED	10am	11am	Abbey	online	
	10	7-May-25	WED	8pm	9pm	Abbey	online	
	10	8-May-25	THUR	3pm	4pm	Abbey	online	
	10	9-May-25	FRI	10am	11am	Abbey	online	
ſ	11	14-May-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	11	15-May-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	11	16-May-25	FRI	10am	12pm	Dipankar	in-person	F07.03.352.Carslaw Building.Carslaw Learning Studio 352
	11	14-May-25	WED	10am	11am	Abbey	online	
	11	14-May-25	WED	8pm	9pm	Abbey	online	
	11	15-May-25	THUR	3pm	4pm	Abbey	online	
	11	16-May-25	FRI	10am	11am	Abbey	online	
ſ	12	21-May-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	12	23-May-25	FRI	10am	12pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	13	28-May-25	WED	10am	12pm	Dipankar	in-person	J12.01.110-VR.The School of Information Technologies.SIT Computer Lab 110-VR
	13	29-May-25	THUR	3pm	5pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
	13	30-May-25	FRI	10am	12pm	Dipankar	in-person	A27.03.330.Edward Ford Building.Edward Ford Computer Lab 330/331
Ī								

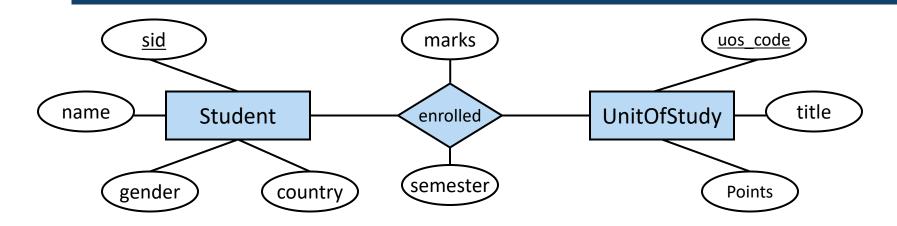


Call for a Volunteer for a "Lightning" Talk Next Week!

If you wish to volunteer for a 5 minute "lightning" talk in Week 7 (next week!) and have professional experience using databases and want to share it with the class, please send me an email or talk to me after class. Hurry up, first come first serve!



Running Example



Student					
<u>sid</u>	name	gender	country		
1001	Adam	М	AUS		
1002	Bob	М	ROK		
1003	Lily	F	AUS		
1004	Simon	М	GBR		
1005	Jesse	F	CHN		
1006	Adam	М	GER		

Enrolled				
<u>sid</u>	uos_code	semester	marks	
1001	COMP5138	2023-S2	72	
1002	COMP5702	2023-S2	85	
1003	COMP5138	2023-S2	67	
1006	COMP5318	2023-S2	94	
1003	ISYS3207	2023-S1	78	
1006	ISYS3207	2023-S2	40	

UnitOfStudy				
uos_code	title	points		
COMP5138	Relational DBMS	6		
COMP5318	Data Mining	6		
INFO6007	IT Project Management	6		
SOFT1002	Algorithms	12		
ISYS3207	IS Project	4		
COMP5702	Thesis	18		



Review: Set Operations

> Find id of all students who are enrolled in both 'COMP5138' and 'ISYS3207'.

SELECT id FROM Enrolled WHERE uos_code='COMP5138'
INTERSECT
SELECT id FROM Enrolled WHERE uos_code='ISYS3207'

> How about listing the **names** of students who are enrolled in both 'COMP5138' and 'ISYS3207'?

SELECT name FROM Student NATURAL JOIN Enrolled WHERE uos_code='COMP5138'

Is this correct?

INTERSECT

SELECT name FROM Student NATURAL JOIN Enrolled WHERE uos_code='ISYS3207'

Student			
<u>sid</u>	name	gender	country
1001	Adam	М	AUS
1002	Bob	М	ROK
1003	Lily	F	AUS
1004	Simon	М	GBR
1005	Jesse	F	CHN
1006	Adam	М	GER

Enrolled					
<u>sid</u>	uos_code	Semester	marks		
1001	COMP5138	2023-S2	72		
1002	COMP5702	2023-S2	85		
1003	COMP5138	2023-S2	67		
1006	COMP5318	2023-S2	94		
1003	ISYS3207	2023-S1	78		
1006	ISYS3207	2023-S2	40		





- Nested Queries
- > Aggregation and Grouping
- > NULL Values and Three-valued Logic





How about listing the names of students who enrolled in both 'COMP5138' and 'ISYS3207'?

```
SELECT name FROM Student NATURAL JOIN Enrolled WHERE uos_code='COMP5138'
INTERSECT
SELECT name FROM Student NATURAL JOIN Enrolled WHERE uos code='ISYS3207'
```



Correct SQL using nested queries

```
FROM Student
WHERE sid IN (
```

This subquery lists the **id** of all students who are enrolled in both 'COMP5138' and 'ISYS3207'

```
SELECT sid FROM Enrolled WHERE uos_code='COMP5138'
INTERSECT
SELECT sid FROM Enrolled WHERE uos_code='ISYS3207'
)
```



Set Comparison Operators in SQL

Format: Value v set-comparison R: v is a value in the outer query and R is the result of a subquery

v [NOT] IN R

- v IN R: tests whether v is in R: true ⇔ v ∈ R
- v **NOT IN** R: tests whether v is *not* in R: **true** ⇔ v ∉ R

> [NOT] EXISTS R

- **EXISTS** R: tests whether a set R is *not* empty: **true** \Leftrightarrow R \neq Ø
- **NOT EXISTS** R: tests whether a set R is empty: **true** \Leftrightarrow R = \emptyset

vop ALL R

- op can be <, ≤, >, ≥, =, ≠
- tests whether a predicate (i.e., **op**) is true for the whole set: **true** $\Leftrightarrow \forall t \in R : (v \text{ op } t)$

v op SOME R

- Same as v op ANY R
- tests whether a predicate (i.e., **op**) is true for at least one set element: **true** $\Leftrightarrow \exists t \in R : (v \text{ op } t)$



Example: Nested Queries

Find the id of the students with the highest mark

```
FROM Enrolled
WHERE marks >= ALL ( SELECT marks
FROM Enrolled )
```

> Find name of the students who did *not* enroll in 2023-S2 semester.

```
FROM Student
WHERE sid NOT IN ( SELECT sid
FROM Enrolled
WHERE semester = '2023-S2')
```





- > A view is a *virtual* table
 - Defined through a SQL query and used as a table in other queries
 - Normally evaluated on each use
 - Provides an abstraction
 - Provides extra security
 - Convenient way of encapsulating repetitive and complex queries as tables

CREATE VIEW student_enrollment **AS**

SELECT sid, name, title, semester

FROM student NATURAL JOIN Enrolled NATURAL JOIN unitofstudy

SELECT *

FROM student_enrollment

ORDER BY name;





> Find the name of male students who are enrolled in units that have the lowest credit point

Solution: [Without view and step-by-step nested query]

(SELECT points FROM UnitOfStudy)));





> Find male students' names who are enrolled in units that have the lowest credit point

Solution: [Using View]

CREATE VIEW MaleStudents AS

SELECT sid, name

FROM Student

WHERE gender = 'M';

CREATE VIEW LowestCreditPointUnit AS

SELECT uos_code

FROM UnitOfStudy

WHERE points <= ALL (SELECT points

FROM UnitOfStudy);

SELECT name

FROM MaleStudents

WHERE sid IN (SELECT sid

FROM Enrolled

WHERE uos_code IN (SELECT uos_code

FROM LowestCreditPointUnit));



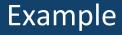


- Nested Queries
- Aggregation and Grouping
- > NULL Values and Three-valued Logic





- > Besides retrieving data, SQL supports several aggregation operations
 - COUNT, SUM, AVG, MAX, MIN (also called aggregate functions)
 - Except **COUNT**, all aggregations apply to a *single attribute*
 - These operations apply to *duplicates*, unless **DISTINCT** is specified





- How many courses are there?
 - SELECT COUNT(*) FROM unitofstudy
- Find the highest mark for 'COMP5138'?
 - SELECT MAX(marks) FROM Enrolled
 WHERE uos code = 'COMP5138'
- Find the average mark of 'COMP5138'?
 - SELECT AVG(marks) FROM EnrolledWHERE uos_code = 'COMP5138'
- How many students are enrolled?
 - SELECT COUNT(DISTINCT sid) FROM Enrolled

Enrolled				
<u>sid</u>	uos_code	Semester	marks	
1001	COMP5138	2020-S2	72	
1002	COMP5702	2020-S2	85	
1003	COMP5138	2020-S2	67	
1006	COMP5318	2020-S2	94	
1003	ISYS3207	2020-S1	78	
1006	ISYS3207	2020-S2	40	

UnitOfStudy				
uos_code	title	points		
COMP5138	Relational DBMS	6		
COMP5318	Data Mining	6		
INFO6007	IT Project Management	6		
SOFT1002	Algorithms	12		
ISYS3207	IS Project	4		
COMP5702	MIT Research Project	18		





Find the id of the students who have the highest mark in COMP5138

SELECT sid

FROM enrolled

WHERE uos_code = 'COMP5138' AND

marks = (SELECT MAX(marks)

FROM enrolled

WHERE uos_code = 'COMP5138');

	Enrolled				
sid	uos_code	Semester	marks		
1001	COMP5138	2020-S2	72		
1002	COMP5702	2020-S2	85		
1003	COMP5138	2020-S2	67		
1006	COMP5318	2020-S2	94		
1003	ISYS3207	2020-S1	78		
1006	ISYS3207	2020-S2	40		





> Find the names of the students who have the highest mark in 'Relational DBMS' course

SELECT name **FROM** student **WHERE** sid **IN** (

SELECT sid **FROM** enrolled **NATURAL JOIN** unitofstudy

WHERE title = 'Relational DBMS' and marks = (SELECT MAX(marks)

FROM enrolled NATURAL JOIN unitofstudy WHERE title = 'Relational DBMS'))

Student					
<u>sid</u>	name	gender	country		
1001	Adam	М	AUS		
1002	Bob	М	ROK		
1003	Lily	F	AUS		
1004	Simon	М	GBR		
1005	Jesse	F	CHN		
1006	Adam	М	GER		

Enrolled				
<u>sid</u>	uos_code	Semester	marks	
1001	COMP5138	2023-S2	72	
1002	COMP5702	2023-S2	85	
1003	COMP5138	2023-S2	67	
1006	COMP5318	2023-S2	94	
1003	ISYS3207	2023-S1	78	
1006	ISYS3207	2023-S2	40	

UnitOfStudy				
uos_code	title	points		
COMP5138	Relational DBMS	6		
COMP5318	Data Mining	6		
INFO6007	IT Project Management	6		
SOFT1002	Algorithms	12		
ISYS3207	IS Project	4		
COMP5702	Thesis	18		





- Instead of aggregating *all* (qualifying) tuples into a single value, sometimes we want to apply aggregation to each of several *groups* of tuples.
- > Example: Find the total sales amount of each company

Sales Table

id	company	amount
1	IBM	5500
2	DELL	4500
3	IBM	6500

SELECT company, SUM(amount)
FROM Sales

company	amount
IBM	16500
DELL	16500
IBM	16500

SELECT company, **SUM**(amount) **FROM** Sales **GROUP BY** company

company	amount		
IBM	12000		
DELL	4500		





Group By Queries

In SQL, we can "partition" a relation into groups according to the values of one or more attributes:

SELECT target-list

FROM relation-list

WHERE qualification

GROUP BY grouping-list

HAVING group-qualification

SELECT company, **SUM**(amount)

FROM Sales

GROUP BY company

- A group is a set of tuples that have the same value for all attributes in the grouping-list.
 - Example: 'IBM' for the company attribute in the Sales table.
- > NOTE: Attributes in the **SELECT / HAVING** clause must be either in *aggregate* functions or from the *grouping-list*
 - Intuitively, each result tuple corresponds to a group, and these attributes must have a single value per group.





SELECT company, **SUM**(amount) **FROM** Sales **GROUP BY** company

Sales Table

id	compa	ny	amour	nt			company	amount
1	IBM		5500				IBM	5500
2	DELL		4500			_	IBM	6500
3	IBM		6500				DELL	4500
		СО	mpany	a	mount			
		IBN	1	120	000			
		DEL		450	00			



Filtering Groups with HAVING Clause

List courses and their average marks

```
FROM enrolled

GROUP BY uos code
```

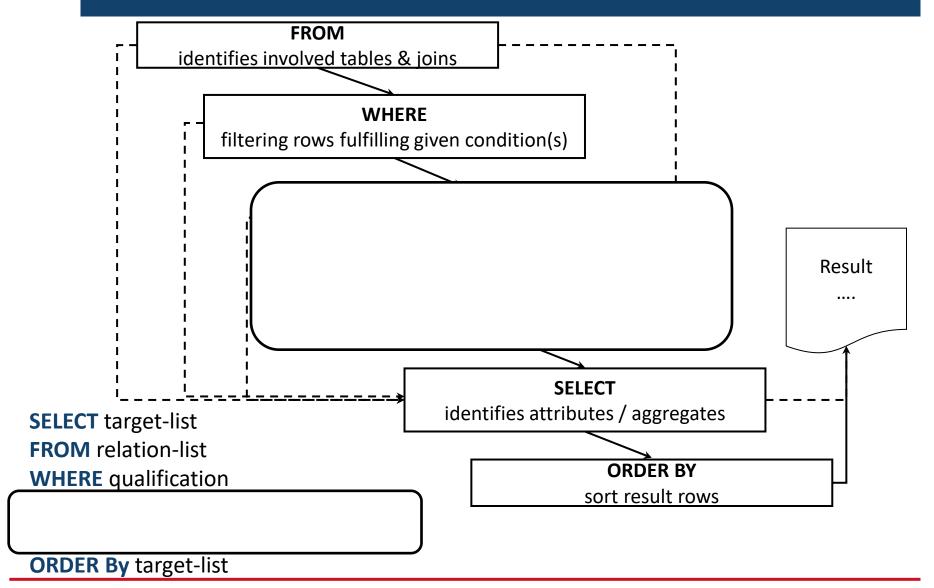
- > HAVING clause can further filter groups to fulfill a predicate
 - Example: Find course code and average mark of courses that have average marks > 60

```
SELECT uos_code, AVG(mark)
FROM enrolled
GROUP BY uos_code
HAVING AVG(mark) > 60
```

NOTE: Predicates in the HAVING clause are applied after the formation of groups,
 whereas predicates in the WHERE clause are applied before forming groups

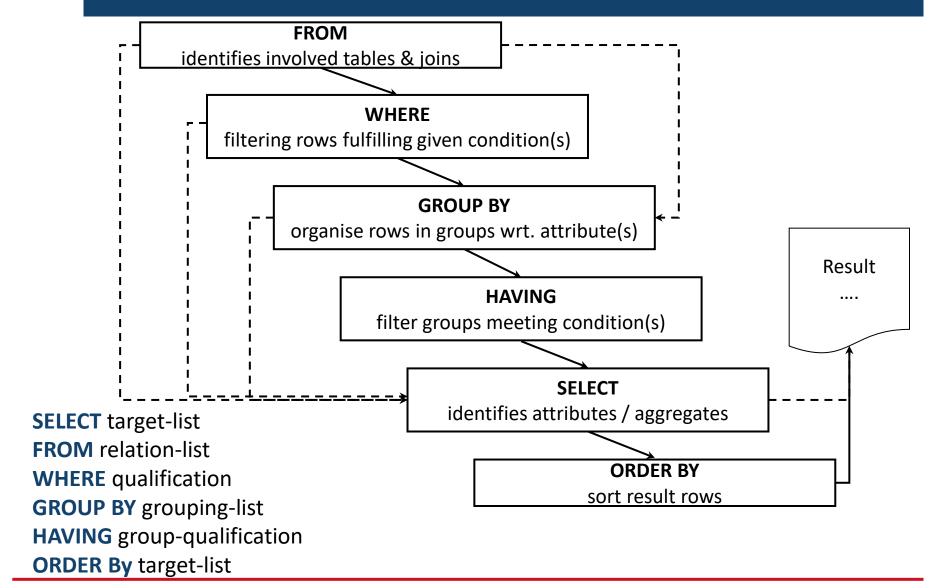


Query-Clause Evaluation Order





Query-Clause Evaluation Order





Evaluation Example

> Find the maximum marks of 6-credit point Units of Study with at least 2 students

SELECT uos_code, MAX(marks)
FROM Enrolled NATURAL JOIN UnitOfStudy
WHERE points = 6
GROUP BY uos_code
HAVING COUNT(*) >= 2

1. Enrolled and UnitOfStudy are joined

FROM Enrolled NATURAL JOIN UnitOfStudy

2. Tuples that fail the WHERE condition are discarded

WHERE credit_points = 6

4	uos_code character (8)	sid integer	semester character varying	marks integer	title character varying (30)	points integer
1	COMP5138	1001	2023-S2	72	Relational DBMS	6
2	COMP5702	1002	2023-S2	85	Thesis	18
3	COMP5138	1003	2023-S2	67	Relational DBMS	6
4	COMP5318	1006	2023-S2	94	Data Mining	6
- 5	ISYS3207	1003	2023-S2	78	IS Project	4
6	ISYS3207	1006	2023-S1	40	IS Project	4



Evaluation Example (cont'd)

3. Remaining tuples are partitioned into groups by the value of attributes in the grouping-list (uos_code).

GROUP BY uos_code

4. Groups which fail the **HAVING** condition are discarded.

HAVING COUNT(*) >= 2

uos_code character (8)	sid integer	semester character varying	marks integer	title character varying (30)	points integer
COMP5138	1001	2023-S2	72	Relational DBMS	6
COMP5138	1003	2023-S2	67	Relational DBMS	6
				'	

COMD5318	1006	2022 62	0.4	Data Mining	6
COIVIT 33 18	1000	2023-32	74	Data Willing	0

5. ONE result tuple is generated per group

SELECT uos_code, **MAX**(mark)

uos_code	MAX(mark)		
COMP5138	72		





> Find the uos code in which students have the lowest average mark

Solution:

SELECT uos_code

FROM enrolled

GROUP BY uos_code

HAVING AVG(marks) <= **ALL (SELECT AVG** (marks)

FROM enrolled

GROUP BY uos code);

FROM enrolled

GROUP BY uos_code

HAVING AVG(marks) = (SELECT MIN(AVG (marks)))

FROM enrolled

GROUP BY uos_code);

NOTE: aggregate function calls cannot be nested

Short break

Let us have some fun again...







- Nested Queries
- > Aggregation and Grouping
- > NULL Values and Three-valued Logic





- It is possible for tuples to have a null value, denoted by NULL, for some of their attributes
 - NULL signifies that a value <u>does not exist</u> or is <u>not applicable</u> -- it does <u>not mean "0" or "blank"!</u>
 - Integral part of SQL to handle missing / unknown information
- > The predicate IS NULL or IS NOT NULL can be used to check for null values
 - e.g. Find students who don't have a mark for an assessment yet.

SELECT sid

FROM enrolled

WHERE marks IS NULL





- > The <u>result</u> of any <u>arithmetic</u> expression involving NULL is NULL
 - e.g. 5 + NULL returns NULL
- Any comparison with NULL returns unknown
 - e.g. 5 < NULL or NULL <> NULL or NULL = NULL
- > Result of WHERE clause predicate is treated as false if it evaluates to unknown
 - e.g: SELECT sid FROM enrolled WHERE marks < 50 ignores all students without a mark





> Three-valued logic for Boolean operations

- OR:
 - (unknown OR true) = true
 - (unknown OR false) = unknown
 - (unknown OR unknown) = unknown
- AND:
 - (true AND unknown) = unknown
 - (false AND unknown) = false
 - (unknown AND unknown) = unknown
- NOT:
 - (NOT unknown) = unknown

> It is equivalent to the following



Unexpected Behavior of NULL Values

SELECT *
FROM enrolled
WHERE marks < 25 OR marks >= 25

The students whose marks are unknown will not be returned!

SELECT *
FROM enrolled
WHERE marks < 25 OR marks >= 25 OR marks IS NULL

This SQL now lists all students!



NULL Values and Aggregation

- Aggregate functions except COUNT(*) ignore NULL values on the aggregated attributes
- > Result of an aggregate function is NULL if all rows have null values.
- > Examples:
 - Minimum mark of all courses

SELECT MIN(marks)

-- ignores tuples with nulls on mark

FROM enrolled

Number of all courses

SELECT COUNT(*)

-- counts *all* tuples

FROM enrolled

- Number of all courses with a mark

SELECT COUNT(marks)

-- ignores tuples with nulls on mark

FROM enrolled



You should now be able to...

> ...formulate even complex SQL Queries

- Nested queries
- Aggregate functions
- Grouping and Having conditions
- Handling NULL values

Let us menti again!

Q/A session







- > Kifer/Bernstein/Lewis (2nd edition)
 - Chapter 5
- > Ramakrishnan/Gehrke (3rd edition the 'Cow' book)
 - Chapter 5
- > Ullman/Widom (3rd edition)
 - Sections 6.3 and 6.4
- Silberschatz/Korth/Sudarshan (5th edition 'sailing boat')
 - Sections 3.1-3.6
- > Elmasri/Navathe (5th edition)
 - Sections 8.4 and 8.5.1





- Database Application Development
 - SQL in Client Code
 - Call-level Database APIs
 - Server-Side Application Development with Stored Procedures

> Readings:

- Kifer/Bernstein/Lewis book, Chapter 8
- Ramakrishnan/Gehrke (Cow book), Chapter 6
- Ullman/Widom, Chapter 9

