



THE UNIVERSITY OF  
**SYDNEY**

CONFIDENTIAL EXAM PAPER

This paper is not to be removed from the exam venue.

Computer Science

EXAMINATION

Semester 1- Practice, 2020

COMP2123/9123 Data Structures and Algorithms

**EXAM WRITING TIME:** 2 hours

**READING TIME:** 10 minutes

**EXAM CONDITIONS:**

This is a CLOSED book examination.

All submitted work must be **done individually** without consulting someone else's solutions, or external resources like the Internet or a textbook, in accordance with the University's "Academic Dishonesty and Plagiarism" policies.

**MATERIALS PERMITTED IN THE EXAM VENUE:**

One A4 sheet of handwritten and/or typed notes double sided

**MATERIALS TO BE SUPPLIED TO STUDENTS:**

Nothing

**INSTRUCTIONS TO STUDENTS:**

Type your answers in your text editor and submit via Canvas. Handwritten responses will **not** be accepted.

Start by typing your student ID at the top of the first page of your submission. Do **not** type your name.

Submit only your answers to the questions. Do **not** copy the questions.

**For examiner use only:**

Problem	1	2	3	4	Total
Marks					
Out of	10	10	20	20	60

**Problem 1.**

a) Analyze the time complexity of this algorithm.

[3 marks]

```

1: def COMPUTE(A)
2:   result  $\leftarrow$  0
3:   for i = 0; i < n; i ++ do
4:     if A[i] > i then
5:       result  $\leftarrow$  result + A[i]
6:   return result

```

b) Solve the following recursion:

[3 marks]

$$T(n) = \begin{cases} T(n/2) + O(1) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

c) We are planning a board games event and we're using one of the shelves in my office to store the games. Unfortunately the shelf only has a certain amount of space  $S$ , so we need to carefully pick which games we want to bring. Every game takes some space  $s_i$  and has a fun factor  $f_i$  that indicates how much fun it is to play that game (for  $1 \leq i \leq n$ ). [4 marks]

We want to maximize the amount of fun we'll have, so we want to maximize the sum of the fun factors of the games we pick (i.e.,  $\max \sum_{\text{picked game } i} f_i$ ), while

making sure that the games fit on my shelf, so the sum of the space the games we pick take should be at most  $S$  (i.e.,  $\sum_{\text{picked game } i} s_i \leq S$ ). For simplicity, you can

assume that all  $f_i$ ,  $s_i$ , and  $S$  are all distinct positive integers.

The strategy of PICKLARGEST is to always pick the game with the highest fun factor until my shelf is full: it sorts the games by their fun factor  $f_i$  in decreasing order and adds a game when its required space is less than the remaining space on the shelf.

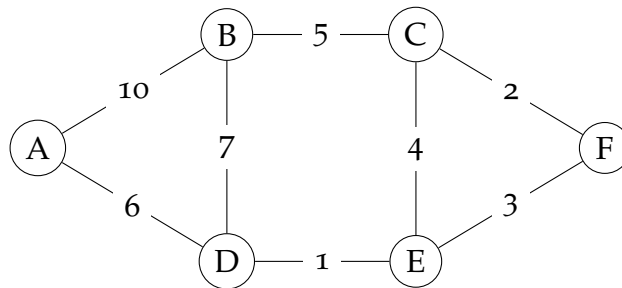
```

1: def PICKLARGEST(all  $f_i$  and  $s_i$ ,  $S$ )
2:   currentSpace  $\leftarrow$  0
3:   currentFun  $\leftarrow$  0
4:   Sort games by  $f_i$  and renumber such that  $f_1 \geq f_2 \geq \dots \geq f_n$ 
5:   for  $i \leftarrow 1$ ;  $i \leq n$ ;  $i++$  do
6:     if currentSpace +  $s_i \leq S$  then
7:       currentSpace  $\leftarrow$  currentSpace +  $s_i$            ▷ Pick the ith game
8:       currentFun  $\leftarrow$  currentFun +  $f_i$ 
9:   return currentFun

```

Show that PICKLARGEST doesn't always return the correct solution by giving a counterexample.

**Problem 2.** Consider the following edge weighted undirected graph  $G$ :



**Your task is to:**

- a) Compute a minimum spanning tree  $T$  of  $G$ . List the edges in  $T$ . [7 marks]
- b) Indicate the order in which the edges are added by Kruskal's algorithm. [3 marks]

(You do **not** have to explain your answer.)

**Problem 3.** Consider the *Dynamic Matrix* ADT for representing an matrix  $A = \{a_{i,j}\}_{i,j=1}^n$  that supports the following operations:

- **CREATE()**: creates a  $1 \times 1$  matrix where  $a_{1,1} = 0$ .
- **SET/GET( $i, j$ )**: set or get the value of the entry  $a_{i,j}$ .
- **INCREASE-SIZE**: If the current size of the matrix is  $n \times n$ , increase it to  $n + 1 \times n + 1$  such that the new entries are set of 0. In other words,  $A$  becomes  $A'$  such that  $a'_{i,j} = a_{i,j}$  if  $1 \leq i, j \leq n$ , and  $a'_{i,j} = 0$  otherwise.

**Your task** is to come up with a data structure implementation for the Dynamic Matrix ADT that uses  $O(n^2)$  space, where  $n$  is the size of the matrix, and **CREATE**, **SET**, **GET** take  $O(1)$  and **INCREASE-SIZE** takes  $O(n)$  time. Remember to:

- Describe your data structure implementation in plain English. [8 marks]
- Prove the correctness of your data structure. [7 marks]
- Analyze the time and space complexity of your data structure. [5 marks]

**Problem 4.** Let  $G$  be a connected undirected graph on  $n$  vertices. We say that two distinct spanning trees  $T$  and  $S$  of  $G$  are *one swap away* from each other if  $|T \cap S| = n - 2$ ; that is,  $T$  and  $S$  differ in only one edge.

For two distinct spanning trees  $T$  and  $S$  we say that  $R_1, R_2, \dots, R_k$  form a *swapping sequence* from  $T$  to  $S$  if:

1.  $R_1 = T$ ,
2.  $R_k = S$ , and
3. for any  $1 \leq i < k$ , the trees  $R_i$  and  $R_{i+1}$  are one swap away from each other

**Your task** is to design a polynomial time algorithm that given  $G$  and two spanning trees  $T$  and  $S$  of  $G$ , constructs a minimum length swapping sequence. Remember to:

- a) Describe your algorithm in plain English. [8 marks]
- b) Prove the correctness of your algorithm. [7 marks]
- c) Analyze the time complexity of your algorithm. [5 marks]