07 Basic Concepts of DBMS

- Database Schema
- Tables, Rows, and Columns
- DBMS Keys
- Normalization (1NF, 2NF, 3NF, BCNF)
- Indexes

01 Database Schema

- A schema is the structure of a database, defining how data is organized into tables
- It includes the tables, fields, data types, and relationships between tables
- The schema acts as a blueprint for the database, guiding how data is stored, organized, and accessed
- For example, in an e-commerce database, you might have a schema with tables for Customers, Orders, and Products

02 Tables, Rows, and Columns

Tables

- The fundamental building blocks of a database
- A table is a collection of related data entries organized in rows and columns
- Example: A Users table might have columns like UserID, Username, Email, and Password

Rows

- Also known as records or tuples, rows represent a single entry in a table
- Example: A single row in the Users table could represent one user's data, such
 as 1, "johndoe", "john@example.com", "password123"

Columns

- Also known as fields or attributes, columns define the data type and characteristics of the data stored in each row
- **Example:** The Email column in the Users table might be defined as a string that stores the user's email addresses

03 DBMS Keys

1. Primary Key

• **Definition:** A primary key is a unique identifier for each record in a table. It ensures that no two rows have the same value in this column or set of columns.

Properties:

- Uniqueness: Each value in the primary key column(s) must be unique across the table.
- Not Null: A primary key cannot contain NULL values.
- A table can have only one primary key.
- **Example:** In a Students table, the StudentID column could be the primary key, uniquely identifying each student.

2. Foreign Key

 Definition: A foreign key is a column or a set of columns in one table that refers to the primary key in another table. It is used to establish a relationship between the two tables.

Properties:

- The foreign key value must match an existing primary key value in the referenced table, ensuring referential integrity.
- A table can have multiple foreign keys.
- **Example:** In an Orders table, the CustomerID column could be a foreign key that references the CustomerID primary key in a Customers table, linking each order to a customer.

3. Candidate Key

 Definition: A candidate key is a column, or a set of columns, that can uniquely identify any record in a table. Each table can have multiple candidate keys.

Properties:

- Uniqueness: Each candidate key must have unique values across the table.
- Not Null: A candidate key cannot contain NULL values.
- One of the candidate keys is selected as the primary key, while others can be considered as alternate keys.
- **Example:** In a Students table, both StudentID and Email could be candidate keys because they can uniquely identify each student.

4. Super Key

• **Definition:** A super key is any combination of columns that uniquely identifies a row in a table. It includes candidate keys as well as any superset of candidate keys.

Properties:

- Uniqueness: A super key can uniquely identify each row in a table.
- Can be composed of one or more columns.
- Super keys include all candidate keys, and any candidate key is a minimal super key.
- Example: In a Students table, both StudentID and the combination of StudentID +
 Email can be super keys.

5. Alternate Key

- **Definition:** An alternate key is any candidate key that is not selected as the primary key. It is an alternative option for unique identification of records in the table.
- Properties:
 - Like candidate keys, alternate keys must be unique and not null.
- **Example:** If StudentID is chosen as the primary key in the Students table, Email would be an alternate key.

6. Composite Key

- **Definition:** A composite key is a primary key that consists of two or more columns that together uniquely identify a row in a table.
- Properties:
 - The combination of the columns in the composite key must be unique.
 - None of the individual columns can uniquely identify the row by themselves.
- **Example:** In a CourseRegistrations table, a combination of StudentID and CourseID could be used as a composite key to uniquely identify each registration.

7. Unique Key

- **Definition:** A unique key is a constraint that ensures all values in a column or a set of columns are unique across the table, similar to a primary key but can allow one NULL value.
- Properties:
 - Uniqueness: Ensures all values are unique.
 - Allows NULL values, but only one NULL per column.
 - A table can have multiple unique keys.
- **Example:** In a Users table, the Username column might be set as a unique key to prevent duplicate usernames.

8. Surrogate Key

 Definition: A surrogate key is an artificial key, usually an auto-incremented number, that is added to a table to act as a primary key.

Properties:

- Surrogate keys have no business meaning and are typically generated by the database system.
- Often used when no natural primary key exists or when natural keys are too complex.
- **Example:** In a Books table, an auto-incremented BookID might be used as a surrogate key, even if each book has a unique ISBN.

9. Natural Key

 Definition: A natural key is a key that has a logical relationship to the data it represents, often derived from the data itself.

Properties:

- Has business meaning and is not artificially created.
- Could be a candidate key or even a primary key if it uniquely identifies records.
- **Example:** In a Employees table, an EmployeeID derived from an employee's badge number might be a natural key

04 Normalization (1NF, 2NF, 3NF, BCNF)

Normalization

 The process of organizing data in a database to reduce redundancy and improve data integrity

1NF (First Normal Form)

- Ensures that each column contains atomic values (indivisible) and that each column in a table contains unique data
- **Example:** A table with repeated columns for phone numbers would not be in 1NF; instead, each phone number should be in a separate row

2NF (Second Normal Form)

- Achieved when a table is in 1NF, and all non-primary key attributes are fully functionally dependent on the primary key
- **Example:** In a table with a composite primary key, all non-key columns should depend on the entire key, not just part of it

3NF (Third Normal Form)

- Achieved when a table is in 2NF, and all the attributes are functionally dependent only on the primary key
- **Example:** If a table contains a column City that depends on ZipCode, which in turn depends on the CustomerID, moving City and ZipCode to a new table

linked by CustomerID brings the table into 3NF

BCNF (Boyce-Codd Normal Form)

- A stricter version of 3NF where every determinant (a column or set of columns on which another column is fully dependent) must be a candidate key
- **Example:** If a table has more than one candidate key, all of them must be a determinant for the table to be in BCNF

1. First Normal Form (1NF)

1NF is the simplest form of normalization and requires that:

- Each table has a primary key
- Each column contains atomic (indivisible) values
- Each column contains only one value per row (no repeating groups or arrays)

Example:

Consider a table storing student information and their enrolled courses:

| StudentID | StudentName | Courses |
|-----------|-------------|--------------------|
| 1 | John | Math, Science |
| 2 | Jane | English, History |
| 3 | Bob | Math, English, Art |

Problem: The Courses column contains multiple values (a list of courses), violating the atomic value rule.

1NF Solution: Split the courses into individual rows:

| StudentID | StudentName | Course |
|-----------|-------------|---------|
| 1 | John | Math |
| 1 | John | Science |
| 2 | Jane | English |
| 2 | Jane | History |
| 3 | Bob | Math |
| 3 | Bob | English |
| 3 | Bob | Art |

Now, each column contains atomic values, and the table is in 1NF.

2. Second Normal Form (2NF)

2NF is achieved when:

- The table is in 1NF
- All non-key attributes are fully functionally dependent on the primary key

Example:

Consider a CourseRegistrations table in 1NF with a composite primary key:

| StudentID | CourseID | StudentName | CourseName | Instructor |
|-----------|----------|-------------|------------|------------|
| 1 | 101 | John | Math | Dr. Smith |
| 2 | 102 | Jane | English | Dr. Brown |
| 3 | 101 | Bob | Math | Dr. Smith |

Problem: The StudentName depends only on StudentID, and CourseName and Instructor depend only on CourseID, not on the entire composite primary key (StudentID + CourseID).

2NF Solution: Split the table into two tables:

Students Table:

| StudentID | StudentName |
|-----------|-------------|
| 1 | John |
| 2 | Jane |
| 3 | Bob |

Courses Table:

| CourseID | CourseName | Instructor |
|----------|------------|------------|
| 101 | Math | Dr. Smith |
| 102 | English | Dr. Brown |

CourseRegistrations Table:

| StudentID | CourseID |
|-----------|----------|
| 1 | 101 |
| 2 | 102 |
| 3 | 101 |

Now, the CourseRegistrations table is in 2NF because all non-key attributes are fully dependent on the entire primary key.

3. Third Normal Form (3NF)

3NF is achieved when:

- The table is in 2NF
- There are no transitive dependencies (i.e., non-key attributes do not depend on other non-key attributes)

Example:

Consider a Students table in 2NF:

| StudentID | StudentName | AdvisorID | AdvisorName |
|-----------|-------------|-----------|-------------|
| 1 | John | 201 | Dr. Allen |
| 2 | Jane | 202 | Dr. Baker |
| 3 | Bob | 201 | Dr. Allen |

Problem: AdvisorName depends on AdvisorID, which is not a primary key but is a non-key attribute. This creates a transitive dependency.

3NF Solution: Split the table into two tables:

Students Table:

| StudentID | StudentName | AdvisorID |
|-----------|-------------|-----------|
| 1 | John | 201 |
| 2 | Jane | 202 |
| 3 | Bob | 201 |

Advisors Table:

| AdvisorID | AdvisorName |
|-----------|-------------|
| 201 | Dr. Allen |
| 202 | Dr. Baker |

Now, the Students table is in 3NF because there are no transitive dependencies.

4. Boyce-Codd Normal Form (BCNF)

BCNF is a stricter version of 3NF. A table is in BCNF if:

- It is in 3NF
- For every functional dependency (A → B), A should be a super key

Example:

Consider a ClassAssignments table:

| ClassID | TeacherID | TeacherName |
|---------|-----------|-------------|
| 1 | 101 | Mr. Smith |
| 2 | 102 | Ms. Johnson |
| 1 | 103 | Ms. Brown |

Problem: ClassID and TeacherID both determine TeacherName, but neither ClassID nor TeacherID is a super key by themselves.

BCNF Solution: Create separate tables:

Classes Table:

| ClassID | TeacherID |
|---------|-----------|
| 1 | 101 |
| 2 | 102 |
| 1 | 103 |

Teachers Table:

| TeacherID | TeacherName |
|-----------|-------------|
| 101 | Mr. Smith |
| 102 | Ms. Johnson |
| 103 | Ms. Brown |

Now, each determinant is a super key, and the table is in BCNF.

05 Indexes

- Indexes are special database objects that improve the speed of data retrieval operations on a database table
- By creating an index on one or more columns, you can make queries that filter or sort data by those columns much faster
- Types:

- Primary Index: Automatically created on the primary key of a table
- Secondary Index: Created on columns other than the primary key to speed up query performance
- **Example:** An index on the Email column in a Users table would make searching for users by email faster

1. Creating Indexes

Basic Index

- A basic index is created on a single column of a table
- It speeds up queries that filter or sort by that column

Example

• Imagine a Users table with a LastName column. To speed up searches based on LastName, you can create an index:

```
CREATE INDEX idx_lastname
ON Users (LastName);
```

• The idx_lastname index is created on the LastName column of the Users table. Queries that search for users by LastName will be faster because the database can use this index to guickly locate the rows

Composite Index

- A composite index is created on multiple columns
- It improves performance on queries that filter or sort based on the combination of these columns

Example

Suppose you frequently run queries that filter by both LastName and FirstName. You
can create a composite index on both columns:

```
CREATE INDEX idx_name
ON Users (LastName, FirstName);
```

• The idx_name index is created on both LastName and FirstName columns. This index will be useful for queries that involve both columns, such as searching for users by both their last and first names

2. Using Indexes in Queries

- Indexes are automatically used by the DBMS when they are appropriate
- You don't need to specify that you want to use an index; the DBMS handles this

Example

Consider the Users table and the following query:

```
SELECT *
FROM Users
WHERE LastName = 'Smith';
```

 If an index on LastName exists, the DBMS will use it to quickly find all users with the last name 'Smith' rather than scanning the entire table

3. Unique Index

- A unique index ensures that the values in the indexed column(s) are unique across the table
- This is often used to enforce uniqueness on a column that is meant to have unique values, such as email addresses or usernames

Example

 To ensure that each email address in the Users table is unique, you can create a unique index:

```
CREATE UNIQUE INDEX idx_email
ON Users (Email);
```

The idx_email index ensures that all values in the Email column are unique. If you
try to insert a duplicate email address, the DBMS will reject it

4. Full-Text Index

- A full-text index is used for searching text within a column
- It is particularly useful for searching large amounts of text data

Example

• If the Articles table has a Content column with large text fields, you can create a full-text index to enable efficient searches:

```
CREATE FULLTEXT INDEX idx_content
ON Articles (Content);
```

Explanation: The idx_content full-text index allows you to perform full-text searches
within the Content column, making it faster to search for keywords or phrases

5. Dropping Indexes

 Indexes can be dropped if they are no longer needed, which can help improve performance for data modification operations

Example

• To remove the idx_lastname index from the Users table:

```
DROP INDEX idx_lastname
ON Users;
```

 The idx_lastname index is removed from the Users table, which might improve performance for inserts, updates, or deletes but may slow down queries that used this index.

6. Index Maintenance

- Indexes need to be maintained to ensure optimal performance
- This involves periodically rebuilding or reorganizing indexes to optimize their efficiency, especially after significant changes to the data

Example

To rebuild an index in MySQL, you can use:

```
ALTER TABLE Users

DROP INDEX idx_lastname,

ADD INDEX idx_lastname (LastName);
```

This command drops the existing idx_lastname index and then recreates it.
 Rebuilding indexes can help improve performance if they become fragmented