

# Computational Homework8 Report

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## 1 Introduction to Gauss Quadrature

### 1.1 Fundamental concepts of Gauss Quadrature

Suppose  $f(x)$  is a polynomial of degree at most  $2n - 1$ . We seek quadrature weights  $\{A_i\}$  and nodes  $\{x_i\}$  such that

$$\int_a^b \omega(x) f(x) dx = \sum_{i=1}^n A_i f(x_i),$$

where  $\omega(x)$  is a given weight function.

Let  $\{S_k(x)\}_{k=0}^\infty$  denote a sequence of orthogonal polynomials with respect to the inner product

$$\langle S_k, S_m \rangle = \int_a^b \omega(x) S_k(x) S_m(x) dx.$$

By the division algorithm for polynomials: Suppose  $p(x), s(x) \in \mathcal{P}(\mathbf{F})$  with  $s(x) \neq 0$ . Then there exist unique polynomials  $q(x), r(x) \in \mathcal{P}(\mathbf{F})$  such that

$$p(x) = s(x)q(x) + r(x), \quad \deg r < \deg s.$$

For our purposes, write

$$f(x) = Q_{n-1}(x)S_n(x) + R_{n-1}(x),$$

where  $\deg Q_{n-1} \leq n - 1$ ,  $\deg R_{n-1} < n$ .

Plugging into the integral, and using the orthogonality:

$$\int_a^b \omega(x) f(x) dx = \int_a^b \omega(x) Q_{n-1}(x) S_n(x) dx + \int_a^b \omega(x) R_{n-1}(x) dx.$$

Since  $R_{n-1}(x)$  is a polynomial of degree at most  $n - 1$  and  $S_n(x)$  is orthogonal to all polynomials of lower degree,

$$\int_a^b \omega(x) Q_{n-1}(x) S_n(x) dx = 0,$$

thus

$$\int_a^b \omega(x) f(x) dx = \int_a^b \omega(x) R_{n-1}(x) dx.$$

The quadrature rule must then be exact for all polynomials of degree at most  $2n - 1$ , hence,

$$\sum_{i=1}^n A_i f(x_i) = \sum_{i=1}^n A_i (Q_{n-1}(x_i) S_n(x_i) + R_{n-1}(x_i)).$$

If we choose  $x_i$  to be the roots of  $S_n(x)$ , i.e.,  $S_n(x_i) = 0$  for  $i = 1, \dots, n$ , then

$$\sum_{i=1}^n A_i Q_{n-1}(x_i) S_n(x_i) = 0,$$

so

$$\int_a^b \omega(x) R_{n-1}(x) dx = \sum_{i=1}^n A_i R_{n-1}(x_i).$$

Because  $R_{n-1}$  is an arbitrary degree  $n - 1$  polynomial, we can construct the weights  $A_i$  by requiring the quadrature to be exact for the Lagrange basis polynomials  $l_j(x)$  for  $j = 1, 2, \dots, n$ :

Let

$$l_j(x) = \prod_{\substack{i=1 \\ i \neq j}}^n \frac{x - x_i}{x_j - x_i}, \quad l_j(x_i) = \delta_{ij}.$$

Then,

$$A_j = \int_a^b \omega(x) l_j(x) dx.$$

Moreover, since  $x_j$  is a root of  $S_n(x)$ , we have

$$S_n(x) = \prod_{i=1}^n (x - x_i), \quad S'_n(x_j) = \prod_{\substack{i=1 \\ i \neq j}}^n (x_j - x_i).$$

Hence, the Lagrange polynomial can be written as

$$l_j(x) = \frac{S_n(x)}{(x - x_j) S'_n(x_j)}.$$

Thus, the integral becomes:

$$\int_a^b \omega(x) f(x) dx = \sum_{i=1}^n A_i f(x_i),$$

where

$$A_j = \int_a^b \omega(x) l_j(x) dx = \int_a^b \omega(x) \frac{S_n(x)}{(x - x_j) S'_n(x_j)} dx.$$

To transform the interval  $[a, b]$  to  $[-1, 1]$ , let

$$x = \frac{b-a}{2}\eta + \frac{b+a}{2}, \quad \eta \in [-1, 1].$$

Then,

$$\int_a^b \omega(x) f(x) dx = \frac{b-a}{2} \int_{-1}^1 \Omega(\eta) F(\eta) d\eta,$$

where  $\Omega(\eta) = \omega(x(\eta))$ ,  $F(\eta) = f(x(\eta))$ .

Finally, the Gauss quadrature rule becomes:

$$\int_a^b \omega(x) f(x) dx \approx \frac{b-a}{2} \sum_{i=1}^n a_i F(\eta_i),$$

where  $\eta_i$  are the roots of the orthogonal polynomial  $S_n$  with respect to the inner product

$$\langle S_k, S_m \rangle = \int_a^b \omega(x) S_k(x) S_m(x) dx.$$

## 2 Numerical Integration: Quadrature Formulas, Nodes, and Weights

We now illustrate the application and effectiveness of Gauss quadrature rules with two examples. The first uses Gauss-Legendre quadrature (weight  $\omega(x) = 1$ ), and the second uses Gauss-Chebyshev quadrature (weight  $\omega(x) = \frac{1}{\sqrt{1-x^2}}$ ).

### 2.1 Example 1: Gauss-Legendre Quadrature

We approximate

$$I_1 = \int_0^\pi e^{3x} \cos(x) dx$$

using Gauss-Legendre quadrature with  $n = 2, 3, 4$  nodes. The true value is computed exactly as

$$I_1 = -3717.7943.$$

```
def If(x):
    return 1/10*(np.sin(x)+3*np.cos(x))*np.exp(3*x)

I_f = If(np.pi) - If(0)
```

The quadrature approximations and errors are summarized below:

$n$	True Value	Numerical Value	Absolute Error	Relative Error
2	-3717.7943	-2082.9894	1634.8050	0.4397
3	-3717.7943	-3504.0510	213.7433	0.0575
4	-3717.7943	-3711.2726	6.5217	0.00175

```
def f(x):
    return np.exp(3*x)*np.cos(x)

class GaussQuadrature:
    def __init__(self,f,n,a,b,weight="1"):
        self.f=f
        self.n=n
        self.a=a
        self.b=b
        self.weight=weight

    def cal_quad(self):
        if self.weight=="1":
            nodes,weights=roots_legendre(self.n)
            xn=(self.b-self.a)/2*nodes+(self.b+self.a)/2
            return (self.b-self.a)/2*sum(self.f(xn)*weights)

GQ = GaussQuadrature(f, 2, 0, np.pi)
#GQ = GaussQuadrature(f, 3, 0, np.pi)
#GQ = GaussQuadrature(f, 4, 0, np.pi)
Q_f = GQ.cal_quad()
I_f = If(np.pi) - If(0)
abs_err_f = abs(Q_f - I_f)
rel_err_f = abs_err_f / abs(I_f)
print("【Gauss-Legendre on [0, pi]】")
print("True value      :", I_f)
print("Numerical value :", Q_f)
print("Absolute error  :", abs_err_f)
print("Relative error  :", rel_err_f)
print()
```

As the number of nodes  $n$  increases, the Gauss-Legendre quadrature's accuracy improves dramatically:

- With  $n = 2$ , the error is large, because the function  $e^{3x} \cos(x)$  is highly non-polynomial and varies rapidly.
- At  $n = 3$ , accuracy already improves by more than an order of magnitude.
- By  $n = 4$ , the quadrature nearly matches the true value (relative error  $< 0.2\%$ ).

## 2.2 Example 2: Gauss-Chebyshev Quadrature for Weighted Integral

For the weighted integral:

$$I_2 = \int_{-1}^1 \frac{1}{\sqrt{1-x^2}\sqrt{16-x^2}} dx$$

**Reduction to Elliptic Integral (Sketch)** To compute

$$I_2 = \int_{-1}^1 \frac{1}{\sqrt{1-x^2}\sqrt{16-x^2}} dx,$$

make the substitution  $x = \sin \theta$ . The bounds change to  $\theta \in [-\pi/2, \pi/2]$ , and the integral simplifies:

$$I_2 = \int_{-\pi/2}^{\pi/2} \frac{1}{\sqrt{16-\sin^2 \theta}} d\theta = 2 \int_0^{\pi/2} \frac{1}{\sqrt{16-\sin^2 \theta}} d\theta$$

Using a standard formula for elliptic integrals:

$$\int_0^{\pi/2} \frac{1}{\sqrt{a^2 - b^2 \sin^2 \theta}} d\theta = \frac{1}{a} K\left(\frac{b}{a}\right)$$

Setting  $a = 4$ ,  $b = 1$ , we get:

$$I_2 = \frac{1}{2} K\left(\frac{1}{4}\right)$$

where  $K(k)$  is the complete elliptic integral of the first kind.

```
k = 1/4
m = k ** 2
K_k = ellipk(m)
I_wg_true = 0.5 * K_k
```

we apply both Gauss-Legendre (which does not take the singular weight into account) and Gauss-Chebyshev (matching the weight):

**Gauss-Legendre Quadrature ( $n = 2$ ) for the weighted function:**

True value	0.798121
Numerical value	0.618853
Absolute error	0.179268
Relative error	0.2246

```
def wg(x):
    return 1/(np.sqrt(1-x**2)*np.sqrt(16-x**2))

GQ = GaussQuadrature(wg, 2, -1, 1)
Q_wg = GQ.cal_quad()
abs_err_wg = abs(Q_wg - I_wg_true)
rel_err_wg = abs_err_wg / abs(I_wg_true)
print("【Gauss-Legendre on [-1, 1] for wg(x)】")
print("True value      :", I_wg_true)
print("Numerical value :", Q_wg)
print("Absolute error  :", abs_err_wg)
print("Relative error   :", rel_err_wg)
```

### Gauss-Chebyshev Quadrature ( $n = 2$ ):

True value	0.798121
Numerical value	0.797965
Absolute error	0.000156
Relative error	0.00020

```
def g(x):
    return 1/np.sqrt(16-x**2)

GQ = GaussQuadrature(g, 2, -1, 1, "cheby")
Q_g = GQ.cal_quad()
abs_err_g = abs(Q_g - I_wg_true)
rel_err_g = abs_err_g / abs(I_wg_true)
print("【Gauss-Chebyshev with weight on [-1, 1]】")
print("True value      :", I_wg_true)
print("Numerical value :", Q_g)
print("Absolute error  :", abs_err_g)
print("Relative error   :", rel_err_g)
```

```
print()
```

**Observations:**

- For the second case, Gauss-Chebyshev quadrature almost exactly matches the true value even with  $n = 2$ , demonstrating how matching the weight to the quadrature rule dramatically improves accuracy.
- Applying Gauss-Legendre quadrature to weighted integrals (with a singular weight) gives large error, further showing the importance of adapting the quadrature rule to the specific weight function of the integral.