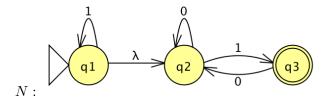
Due 9pm, Thursday, October 8, 2020

Overview. This is your second and final homework focusing on regular languages. The first question recalls the construction for converting an NFA to a DFA; the second has you construct a regular expression equivalent to an FA; the third and fourth have you design FAs equivalent to particular regular expressions, with the fourth requiring you to think beyond the general regex \rightarrow NFA construction; the last involves proving nonregularity using the pumping lemma (Section 1.4), which will be discussed in class in Monday and Wednesday. *Questions 1-4 are graded on correctness only; full credit for Question 5 depends on logical soundness/clarity.*

Question 0. List any consultations/collaborations for this assignment and elaborate to the extent helpful.



Question 1. For the NFA $N=(Q,\{0,1\},\delta,q_0,F)$ depicted above and relabeling $Q=\{1,2,3\}$, consider the equivalent DFA $M=(Q',\{0,1\},\delta',q'_0,F')$ constructed following the general procedure from class, i.e., with $Q'=\mathcal{P}(Q),\,\delta'(R,a)=E(\bigcup_{r\in R}\delta(r,a))$ for all $R\in Q'$ and $a\in\Sigma,\,q'_0=E(\{q_0\}),\,$ and $F'=\{R\in Q'\mid R\cap F\neq\emptyset\}$, where E(R) denotes the set of states ϵ -reachable from some $r\in R$.

- a) How many states are in Q' as defined above?
- b) What is the start state?
- c) How many of these states are reachable from the start state?
- d) What is the label of the state the DFA is in after digesting the characters 10010?

Question 2. Give a regular expression for the language recognized by the machines in Question 1.

Question 3. Example 1.56 in the text illustrates the general conversion from a regular expression to an NFA for the regular expression $(ab \cup a)^*$. As with the general conversion from NFAs to DFAs, this general conversion may result in more states than necessary. Submit a JFLAP file containing a two-state NFA that recognizes the language described by $(ab \cup a)^*$, and include a diagram in your pdf.

Question 4. Submit a JFLAP finite automaton that recognizes the language of $(0 \cup 1) \circ 0^*$, and include a diagram of it in your pdf.

Question 5. The following language over $\Sigma = \{1, \#\}$ contains #-separated lists of distinct unary values:

$$A = \{x_1 \# x_2 \# \dots \# x_k \mid k \ge 0, x_i = 1^* \text{ for all } i = 1 \dots k, x_i \ne x_j \text{ for all } i \ne j\}$$

Use the pumping lemma to show A is not regular. In other words, give a string $s \in A$ and argue that no matter how you partition s = xyz with |y| > 0, $|xy| \le p$, there exists some $i = 0, 2, 3, \ldots$ such that $xy^iz \notin A$.