

Lab #3. ROP & Challenges

Prof. Jaeseung Choi

Dept. of Computer Science and Engineering

Sogang University

General Information

■ Check "Lab #3" in *Assignment* tab of *Cyber Campus*

- Skeleton code (Lab3.tgz) is attached in the post
- Deadline: **11/15 Friday 23:59**
- Submission will be accepted in that post, too
- Late submission due: **11/17 Sunday 23:59 (-20% penalty)**
- Delay penalty is applied uniformly (**not problem by problem**)

■ Please read the instructions in this slide carefully

- This slide is a step-by-step tutorial for the lab
- It also contains important submission guidelines
 - If you do not follow the guidelines, **you will get penalty**

Remind: Cheating Policy

- **Cheating (code copy) is strictly forbidden in this course**
 - Read the orientation slide once more
- **Don't ask for solutions in the online community**
 - TA will regularly monitor the communities
- **Sharing your code with others is as bad as copying**
 - Your cooperation is needed to manage this course successfully
- **You must submit a report as well**
 - More instructions are provided at the end of this slide

Skeleton Code Structure

- Copy Lab3.tgz into CSPRO server and decompress it
 - You must connect to csproN.sogang.ac.kr (N = 2, 3, or 7)
- Skeleton code has similar structure to the previous lab
 - 3-1/ ... 3-4/ : Problems that you have to solve
 - 3-5/ : *Bonus problem* for practice (**not included in grading**)
 - But this one can be important when preparing the lab exam
 - **check.py**, **config**: Files for self-grading
- This slide will provide a guide on writing ROP exploit
 - It also provides a detailed tutorial for solving 3-1

```
jschoi@cspro2:~$ tar -xzf Lab3.tgz
jschoi@cspro2:~$ ls Lab3
3-1  3-2  3-3  3-4  3-5  check.py  config
```

Reading serect.txt

- In the lecture slide, we talked about `execve()` function
 - But there are other variants of like `execv()`, `execl()`, ...
- In this lab, you must run the following code* with ROP
 - There are other ways to read `secret.txt`, but **don't use them**
 - Ex) Using `system()` instead, or spawning a shell with `execv()`
 - They may not work and you can even **get 0 point** in such cases
 - For instance, `system()` does not work properly with SUID

```
// You can run "cat secret.txt" with execv() as follow.  
char *argv[3];  
argv[0] = "/bin/cat";  
argv[1] = "secret.txt";  
argv[2] = NULL;  
execv(argv[0], argv);
```

Example: Problem 3-1

- Target program (twice.c / twice.bin) is given

```
void run_cat(char *filepath) {  
    char *argv[3];  
    argv[0] = "/bin/cat";  
    argv[1] = filepath;  
    argv[2] = NULL;  
    execv(argv[0], argv);  
}
```

Your goal is to execute this function with "secret.txt"

...

```
void vuln(void) {  
    char buf[20];  
    printf("Input your message in stack buffer: ");  
    read(0, buf, 64);  
}
```

You can see that BOF occurs here

Finding ROP Gadgets

- In principle, you must disassemble all the addresses in the code section, which contains assembly instructions
- Pwntools offers `ROP()` API that does this automatically
 - `print(rop.rdi)`: Print gadgets that can affect `%rdi` register
- Tip: You can use `p64()` function to write concise code
 - FYI, `u64()` function performs conversion in opposite direction

```
p = process("./twice.bin")
rop = ROP("./twice.bin")
# You can print the gadget information as follow.
print(rop.rdi)

# The following two lines have the same meaning.
rdi_gadget = b"\xb3\x12\x40\x00\x00\x00\x00\x00"
rdi_gadget = p64(0x4012b3) # More concise
```

Attaching GDB to Process

■ Assume that you wrote the exploit code below

- It uses ROP gadget to change the value of `%rdi` into `0x4142`

■ Let's use `gdb` to check if this works as expected

- Previously, we launched `gdb` and started a process from there
- This time, let's run the script and attach to the *running process*

```
p = process("./twice.bin")

# You can use this line to pause the script for a while.
input("Attach GDB now and press enter to continue: ")
...
print(p.recvuntil(b"stack buffer: "))
rdi_gadget = p64(0x4012b3)
p.send(b"a" * 0x28 + rdi_gadget + p64(0x4142))

input("Done, but let me wait for a while...")
```


Attaching GDB to Process

- You must open **two terminals** and switch between them
 - When launching `gdb`, specify the **process id (pid)** to attach

Step 1. Start the exploit script (1st terminal)

```
jschoi@csp2:~Lab3/3-1$ ./exploit-twice.py  
[+] Starting local process './twice.bin': pid 6936  
Attach GDB now and press enter to continue:
```

Step 2. Attach and set breakpoints (2nd terminal)

```
jschoi@csp2:~Lab3/3-1$ gdb -q ./twice.bin 6936  
Reading symbols from ./twice.bin...  
...  
(gdb) b * 0x4011f0  
Breakpoint 1 at 0x4011f0  
(gdb) c  
Continuing.
```

Attaching GDB to Process

- You must open **two terminals** and switch between them
 - In the 2nd terminal, you can use the **gdb** commands to debug

Step 3. Resume the exploit script (1st terminal)

```
jschoi@cspro2:~Lab3/3-1$ ./exploit-twice.py  
[+] Starting local process './twice.bin': pid 6936  
Attach GDB now and press enter to continue: Let's go!  
b'Input your message in global buffer: '  
b'Input your message in stack buffer: '  
Done, but let me wait for a while...
```

You must type in something like this

Step 4. Now the breakpoint is hit (2nd terminal)

```
...  
Breakpoint 1, 0x00000000004011f0 in vuln ()  
(gdb) x/2xg $rsp  
0x7ffe17aa3248: 0x00000000004012b3    0x00000000000004142
```

Demonstration

Obtaining Function Offset

- For problem 3-2, you will have to obtain the offset of a function within the `libc` library
 - Recall that you need this information to figure out the address of `execv()` function, using memory disclosure
- You can do this easily by using the `pwntools` API
 - Then you don't have to hard-code constants in your script

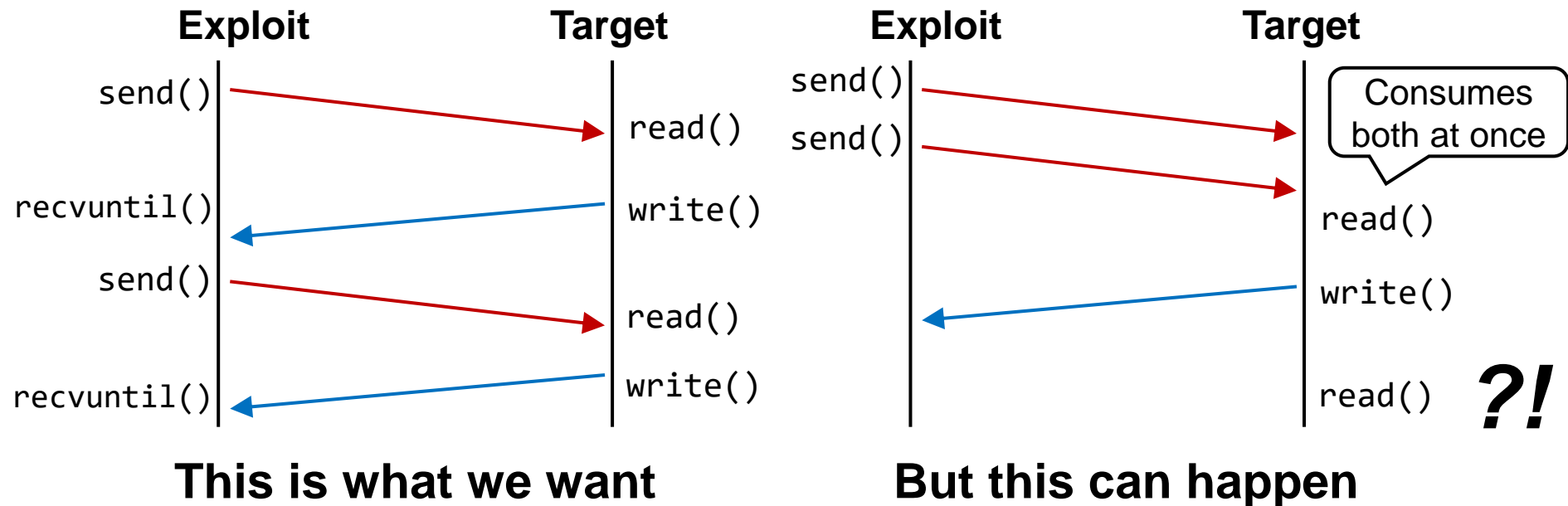
```
# You can investigate the offset of libc functions as follow.  
libc = ELF("/lib/x86_64-linux-gnu/libc.so.6")  
read_offset = libc.symbols['read']  
execv_offset = libc.symbols['execv']  
print("Offset of read() within library: 0x%x" % read_offset)  
print("Offset of execv() within library: 0x%x" % execv_offset)
```

Hints

- In 3-1, be careful in handling newline (`\n`) and null (`\0`)
 - Recall that `pwntools`' `sendline()` implicitly appends `'\n'`
- In 3-2, you must leak the addresses of `libc` functions
 - Try to disclose the library addresses stored in **GOT**
 - A function's **GOT** entry is filled in when it's *called for the first time*
- In 3-3, you will have to exploit a **format string bug** to disclose the memory content of an arbitrary address
- In 3-4, you must exploit a **use-after-free** vulnerability
 - First, examine the behavior of memory allocator, by writing a simple program with `malloc()` and `free()` sequence
 - In other words, think about how to make the allocator return the freed block that you want

Caution: Reliability of Exploit

- In this lab, your script has to be especially careful in interacting with the target program carefully
 - Make sure that you send and receive message **step by step**
 - If not, your exploit code may not work reliably (if it doesn't work during the actual grading for this reason, I will deduct point)

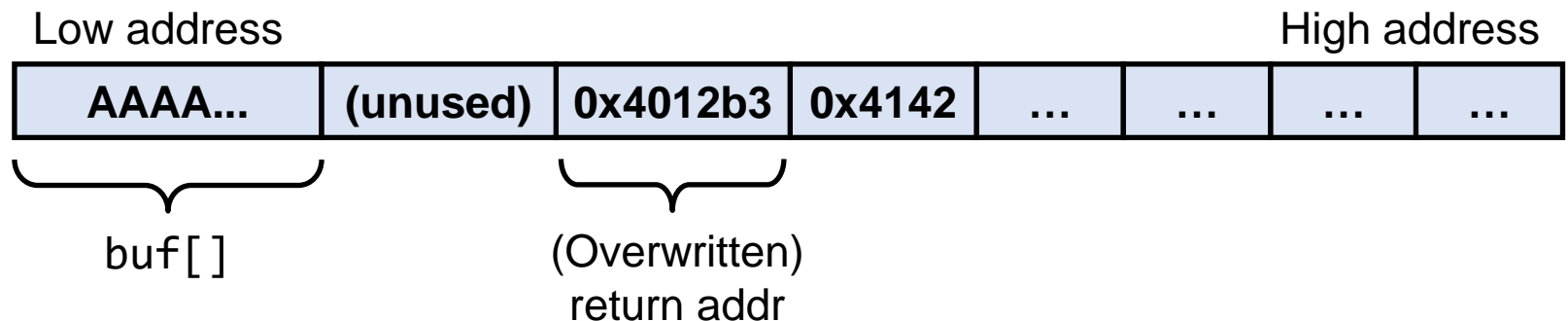


Report Guideline

- **Write report for 3-2 and 3-3 (not required for 3-1 and 3-4)**
 - The role of report is to prove that you solved them on your own
 - **If you couldn't solve a problem, don't have to write its report**
 - Report will not give you point; it is only used to deduct point
- **This time, I will provide a template for each problem**
 - Make sure that your report contains the requested content
- **If you used ChatGPT to write your exploit code, clearly describe it in your report (review the orientation slide)**

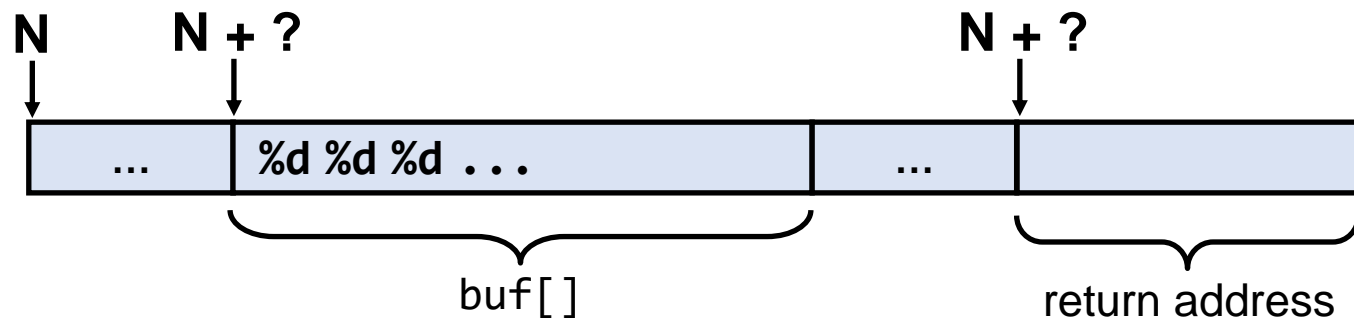
Report Template for 3-2

- Draw the state of stack frame after your input overflows the buffer (see the example below)
 - Draw it just as I did in the lecture slide of the ROP chapter
 - Explain the **meaning (role) of each memory** block in the figure
 - Ex) If it's a gadget address, explain **what that gadget does**
 - Ex) If it's an address of a function, explain **what arguments you are trying to pass**, and **why you are doing that**
 - If you are passing a pointer (memory address) as a function argument, explain what is stored in that address



Report Template for 3-3

- Draw the state of `main()`'s stack frame immediately before `printf()` is called
 - Clearly indicate the positions of `buf[]` and `saved return address` in the stack frame
 - `N` must be the value of `%rsp` at address `0x4011e9`
 - You will be entering format specifiers as input ("`%d%d%d...`"); so explain `which stack position is consumed` by each specifier
 - Justify `why you repeat` each format specifier `for certain number of times`



Problem Information

- There are four problems you have to solve (25 pt. each)
 - Problem 3-1: `twice.bin` (★)
 - Problem 3-2: `substr.bin` (★★★)
 - Problem 3-3: `fsb.bin` (★★☆)
 - Problem 3-4: `item.bin` (★★☆)
- You'll get the point for each problem if the exploit works
 - No partial point for non-working exploit
- If the report does not clearly explain how you analyzed and solved the problem, you will **lose points**
 - You can write the report in Korean or English
 - Due to the limited time, I will randomly select a problem to grade

Submission Guideline

■ You should submit four exploit scripts and report

- Problem 3-1: `exploit-twice.py`
- Problem 3-2: `exploit-substr.py`
- Problem 3-3: `exploit-fsb.py`
- Problem 3-4: `exploit-item.py`
- **Don't forget the report: `report.pdf`**
- 3-5 is a bonus problem, so you don't have to submit it

■ Submission format

- Upload these files directly to *Cyber Campus* (**do not zip them**)
- **Do not change the file name** (e.g., adding any prefix or suffix)
- If your submission format is wrong, you will get **-20% penalty**