Models of Computation Assessed Coursework I

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This document is a skeleton which provides macros to help typeset your solution to the coursework.

Rules can be typeset like this:

$$\begin{array}{ll} \text{E.VAR} \, \frac{x \in \text{dom}(s)}{\langle x, s, h \rangle \, \Downarrow_e \, \langle s(x), s, h \rangle} & \text{E.NUM} \, \frac{}{\langle n, s, h \rangle \, \Downarrow_e \, \langle n, s, h \rangle} \\ \\ \text{E.ADD} \, \frac{\langle E_1, s, h \rangle \, \Downarrow_e \, \langle n_1, s', h' \rangle \quad \langle E_2, s', h' \rangle \, \Downarrow_e \, \langle n_2, s'', h'' \rangle \quad n_3 = n_1 \pm n_2}{\langle E_1 + E_2, s, h \rangle \, \Downarrow_e \, \langle n_3, s'', h'' \rangle} \\ \\ \text{E.NEW} \, \frac{a \notin \text{dom}(h) \quad (a+1) \notin \text{dom}(h)}{\langle \text{newpair}, s, h \rangle \, \Downarrow_e \, \langle \ulcorner a \urcorner, s, h [a \mapsto 0] [a+1 \mapsto 0] \rangle} \\ \\ \text{E.FST} \, \frac{\langle E, s, h \rangle \, \Downarrow_e \, \langle \ulcorner a \urcorner, s', h' \rangle \quad a \in \text{dom}(h')}{\langle E.\text{fst}, s, h \rangle \, \Downarrow_e \, \langle h'(a), s', h' \rangle} \\ \\ \text{E.SND} \, \frac{\langle E, s, h \rangle \, \Downarrow_e \, \langle \ulcorner a \urcorner, s', h' \rangle \quad a + 1 \in \text{dom}(h')}{\langle E.\text{snd}, s, h \rangle \, \Downarrow_e \, \langle h'(a+1), s', h' \rangle} \end{array}$$

Derivations can be typeset as in Figure 1 Types:

$$\Gamma; s; h \vdash \mathsf{well} ext{-typed}$$

$$\Gamma \vdash E : \tau$$

$$h \Vdash v : \tau$$

Suppose that $\Gamma; s; h \vdash \mathsf{well-typed}$. We wish to show that

$$\Gamma \vdash E : \tau \implies h \Vdash v : \tau.$$

Figure 1: A derivation