

# Step 2

Modes of Composition in functional Scala programming

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Lambdajam Online 2020

## Step 2: Validating an Order Message

- We will assume that the system shouldn't fully trust incoming orders and will validate their Customer and Skus
  - Could be a B2B (business-to-business) or franchised integration, or simply an team boundary in a large organisation
- Customers and Skus are managed by different data stores, abstracted by interfaces `CustomerLookup[F]` and `SkuLookup[F]` respectively
  - These would likely be databases or key-value stores in production, but we'll use in-memory implementations for testing
- Some Skus can only be sold to customers in particular regions. Each customer has a region associated with them. We'll load this info from the lookups as well.

# Zooming in on the problem

```
object OrderProcessor {  
  
  def resolveOrderMsg[F[_]: Sync: Parallel: SkuLookup: CustomerLookup](msg: OrderMsg): F[CustomerOrder] =  
    ???  
  
}
```

# Zooming in on the problem

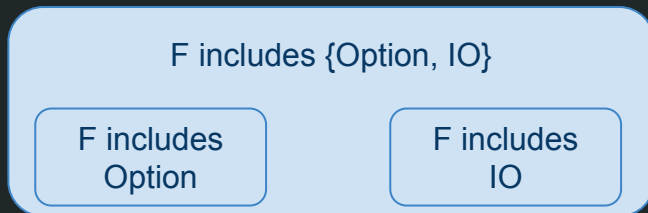
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```

F, the effect type, has four typeclass constraints this time

We'll go through them one by one

# Modes of Composition

- *An effectful value can incorporate several different effects*
  - *Tagless Final style and/or Monad Transformers*
- If part of our program includes an Option effect, and another includes an IO effect, how can we compose them together?
  - The effects present in our program are a *Set of Constraints*. We can add to the set when needed.



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```
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```

Sync is a key typeclass from Cats-effect. Its ability is self-described as to “suspend the execution of side-effects” in the F context.

It implies capabilities similar to Monix’s Task, or the “IO monad” from Cats Effect or Haskell

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```

The Parallel typeclass from Cats library allows parts of a task or computation to be run in parallel with each other

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```
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```

SkuLookup is a domain capability trait defined by the application (rather than from a library)

It provides a single capability: looking up a Sku record from a Sku code string



# Zooming in on the problem

```
object OrderProcessor {  
  
  def resolveOrderMsg[F[_]: Sync: Parallel: SkuLookup: CustomerLookup](msg: OrderMsg): F[CustomerOrder] =  
    ???  
  
}
```

CustomerLookup is another domain capability  
trait: looking up a Customer record from a  
Customer ID string

# Zooming in on the problem

```
object OrderProcessor {  
  
  def resolveOrderMsg[F[_]: Sync: Parallel: SkuLookup: CustomerLookup](msg: OrderMsg): F[CustomerOrder] =  
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```

OrderMsg is the record we parsed from JSON in the previous exercise

# Zooming in on the problem

```
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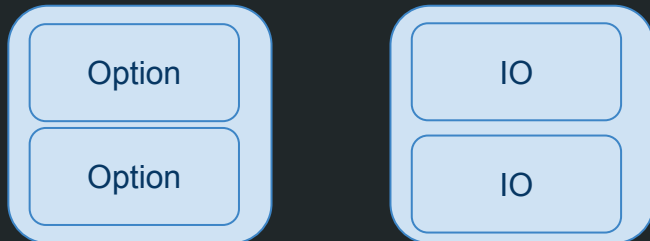
CustomerOrder is a validated domain object  
with invariants enforced by the type system

# Customer and Sku Lookups

- Both Customer and Sku Lookups are effectful queries
  - There is one Customer lookup, but potentially many Sku lookups
- None depends on any other, so they should be done in parallel
- How can we compose multiple effectful queries in parallel?

# Modes of Composition

- *Independent effectful values can be run together concurrently*
  - *Applicative (or Semigroupal) operators*
- Multiple independent effectful values can be composed in parallel into a larger value of the same type using `Semigroupal.tupledN` operators
  - Operations can't depend upon each other's outputs and happen in an undefined order
- The “meaning” of the composition still depends upon the type of effect present in the value



# Modes of Composition

- *An effectful value depends upon the result of another, or must happen after*
  - *Monadic operators*
- **FlatMap**, aka **Bind**, composes two effectful values into a larger value of the same type
  - The second operations happens after the first and may depend upon the value it yields
  - Flatmap operator is written symbolically as `>>=`
- The “meaning” of the composition depends upon the type of effect present in the value



# Parallel Composition

- When we have a fixed number of effectful values that we wish to compose in parallel
  - Put them in a tuple
  - Invoke `parTupled` on the tuple
    - The components are all run concurrently and a tuple of the results is returned
  - Common variant: `parMapN(f)`
    - The components are all run concurrently, and the function `f` is invoked on the results (ie simply maps over the result)

# Sku Lookups: Parallel Traversal

- Lets look at Sku lookups. We have a variable number of items to work with
  - `case class OrderMsg(customerId: String, skuQuantities: NonEmptyChain[(String, Int)])`
  - `trait SkuLookup[F[_]] { def resolveSku(s: String): F[Either[String, Sku]] }`
- We can invoke `parTraverse` on a `NonEmptyChain` to run an effectful function over each item, returning a `NonEmptyChain` of the results



# SKU Lookups: Getting Rid of the Either

- `SkuLookup` returns us an `Either[String, Sku]`
  - How can we get at the `String` and deal with the `Either`?
- We can take advantage of the error facility built-in to the `Sync` effect
  - `Sync[F]` extends `MonadError[F, Throwable]`
  - The same combinator we used in previous problem, `errorValueFromEither`, can lift the left side of the `Either` into an error effect
- To feed the result of `parTraverse` over `resolveSku` into `errorValueFromEither`, we need to compose them sequentially (ie using monads)
  - Use **`flatMap`** or its alias `>>=` ie `<expression1>.>>=(result => <expression2>)`

# PosInt: A look at Refinement Types

- Note `SkuQuantity.quantity` is a `PosInt`
  - A positive int is an `Int` certified by the type system to be  $> 0$
  - Cannot order zero or a negative quantity of something
- `PosInt` is a *refinement type* provided by the Refined library
  - Refined types have a base type (`Int` in this case) refined with extra constraints (positiveness in this case)
- How do we obtain a `PosInt`?
  - `PosInt` provides a `PosInt.fromF[F](<an int>): F[PosInt]` that will raise an error if the passed `Int` isn't positive

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# Have a go

- **resolveOrderMsg** has been broken out into building block parts
  - The final composition of the parts is already filled in
- Try to implement the building blocks to implement **resolveOrderMsg**
- Run the unit tests to check your work
  - `step2/test`

# Advice:

- Scala's Type Inference is limited
  - Methods that take the effect type  $F$  as a type parameter will likely need to have  $F$  specified
  - If not, likely to see compiler error messages about diverging or ambiguous implicits

