

# Exercises

- Infer the type (if possible) of  $\lambda f x \rightarrow x (f x)$ .

	$\frac{}{[H = B]}$
$] \frac{}{[I = D]}$	
	$f : B, x : D \vdash f : H$
$f : B, x : D \vdash x : I$	
$\frac{}{[F = B]}$	$\frac{}{[H = I \rightarrow G]}$
$f : B, x : D \vdash x : F$	$f : B, x : D \vdash f x : G$
	$\frac{}{[F = G \rightarrow E]}$
$f : B, x : D \vdash x (f x) : E$	
$\frac{}{[C = D \rightarrow E]}$	
$f : B \vdash \lambda x \rightarrow x (f x) : C$	
$\frac{}{[A = B \rightarrow C]}$	
$\vdash \lambda f x \rightarrow x (f x) : A$	

We collect the equations:

$A = B \rightarrow C$
$C = D \rightarrow E$
$F = G \rightarrow E$
$F = B$
$H = I \rightarrow G$
$H = B$

$$I = D$$

Now apply the type inference algorithm (TODO).

- Prove that the following are isomorphisms:

- $\text{Prod } (B \rightarrow A) (C \rightarrow A) \sim \text{Sum } B \ C \rightarrow A$

```
q2a-to : {A B C : Set} -> Prod (B -> A) (C -> A) -> Sum B  
C -> A
```

```
q2a-to (pair f g) (left b) = f b
```

```
q2a-to (pair f g) (right c) = g c
```

```
q2a-from : {A B C : Set} -> (Sum B C -> A) -> Prod (B -> A  
) (C -> A)
```

```
q2a-from fg = pair (\ b -> fg (left b)) (\ c -> fg (right  
c))
```

To check that these are mutually inverse, let  $f : B \rightarrow A$  and  $g : C \rightarrow A$  be arbitrary functions. We need  $\text{q2a-from } (\text{q2a-to } (\text{pair } f \ g)) = \text{pair } f \ g$ . We check this by computation:

```
q2a-from (q2a-to (pair f g)) = pair (\ b -> (q2a-to pair f  
g) (left b)) (\ c -> (q2a-to pair f g) (right c))  
                                = pair (\ b -> f b) (\ c -> g c  
)  
                                = pair f g
```

We also need to check the opposite order, so suppose  $fg : \text{Sum } B \ C \rightarrow A$  is an arbitrary function. We need  $q2a\text{-to } (q2a\text{-from } fg) = fg$ . Now  $fg$  itself is a function, so to check for equality, let  $bc : \text{Sum } B \ C$  be an arbitrary element. We need to check that  $q2a\text{-to } (q2a\text{-from } fg) \ bc = fg \ bc$ . To check this, we consider both cases for  $bc$ :

```
q2a-to (q2a-from fg) (left b) = q2a-to (pair (\ b -> fg (left b)) (\ c -> fg (right c))) (left b)
                                     = (\ b -> fg (left b))
                                     = fg (left b)
```

```
q2a-to (q2a-from fg) (right c) = q2a-to (pair (\ b -> fg (left b)) (\ c -> fg (right c))) (right c)
                                     = (\ c -> fg (right c))
                                     = fg (right c)
```

- $C \rightarrow \text{Prod } A \ B \sim \text{Prod } (C \rightarrow A) \ (C \rightarrow B)$

```
q2b-to : {A B C : Set} -> (C -> Prod A B) -> Prod (C -> A)
(C -> B)
```

```
q2b-to fg = pair (\ c -> proj1 (fg c)) (\ c -> proj2 (fg c))
```

```
q2b-from : {A B C : Set} -> Prod (C -> A) (C -> B) -> C ->
Prod A B
```

```
q2b-from (pair f g) = \ c -> pair (f c) (g c)
```

To check that these are inverse in one direction, let  $fg : C \rightarrow \text{Prod } A \ B$  be an arbitrary function. We need that  $q2b\text{-from } (q2b\text{-to } fg) = fg$ . Since  $fg$  is a function, we check this for every input  $c : C$ , so we need that  $q2b\text{-from } (q2b\text{-to } fg) \ c = fg \ c$ .

```

q2b-from (q2b-to fg) c = q2b-from (pair (\ c -> proj1 (fg
c)) (\ c -> proj2 (fg c))) c
                        = (\ c -> pair ((\ c -> proj1 (fg c
)) c) ((\ c -> proj2 (fg c)) c)) c
                        = pair ((\ c -> proj1 (fg c)) c) ((
\ c -> proj2 (fg c)) c)
                        = pair (proj1 (fg c)) (proj2 (fg c)
)
                        = fg c

```

Now for the opposite order, let  $f : C \rightarrow A$  and  $g : C \rightarrow B$ . We need that  $q2b\text{-to } (q2b\text{-from } (pair \ f \ g)) = pair \ f \ g$ . To check that these two elements are equal, considering that it is a pair, we need to check that the coordinates are equal, that is, that  $proj1 \ (q2b\text{-to } (q2b\text{-from } (pair \ f \ g))) = proj1 \ (pair \ f \ g)$  and  $proj2 \ (q2b\text{-to } (q2b\text{-from } (pair \ f \ g))) = proj2 \ (pair \ f \ g)$ . In other words, we need to check that  $proj1 \ (q2b\text{-to } (q2b\text{-from } (pair \ f \ g))) = f$  and  $proj2 \ (q2b\text{-to } (q2b\text{-from } (pair \ f \ g))) = g$ . Now both of these coordinates are functions, respectively with types  $C \rightarrow A$  and  $C \rightarrow B$ . So to check that the coordinates are equal, we need to check that for every input, they have the same output. So let  $c : C$  be an arbitrary element. We check, by computation, that  $proj1 \ (q2b\text{-to } (q2b\text{-from } (pair \ f \ g)))$

$c = f\ c$ , and that  $\text{proj2}\ (\text{q2b-to}\ (\text{q2b-from}\ (\text{pair}\ f\ g)))\ c = g\ c$

```
proj1 (q2b-to (q2b-from (pair f g))) c = proj1 (pair (\ c
-> proj1 ((q2b-from (pair f g)) c)) (\ c -> proj2 ((q2b-fr
om (pair f g)) c)))
                                           = (\ c -> proj1 ((q
2b-from (pair f g)) c)) c
                                           = proj1 ((q2b-from
(pair f g)) c)
                                           = proj1 ((\ c -> (p
air (f c) (g c))) c)
                                           = proj1 (pair (f c)
(g c))
                                           = f c
```

```
proj2 (q2b-to (q2b-from (pair f g))) c = proj2 (pair (\ c
-> proj2 ((q2b-from (pair f g)) c)) (\ c -> proj2 ((q2b-fr
om (pair f g)) c)))
                                           = (\ c -> proj2 ((q
2b-from (pair f g)) c)) c
                                           = proj2 ((q2b-from
(pair f g)) c)
                                           = proj2 ((\ c -> (p
air (f c) (g c))) c)
                                           = proj2 (pair (f c)
(g c))
                                           = g c
```

- **(Advanced):** Extend the algorithm of type inference to STLC with **Sum** and **Prod**.
- Define the *multiplication* operation, as *repeated addition* (defined in the lecture).

*Hint:*

$$m * 1 = m$$

$$m * (1 + n) = m * 1 + m * n = m + (m * n)$$

```
mult : Nat -> Nat -> Nat
mult n 0 = 0
mult n (suc k) = add n (mult n k)
```

- Define the function **max : Nat -> Nat -> Nat** which returns the maximum of two **Nat**s

```
max : Nat -> Nat -> Nat
max zero k          = k
max (suc n) zero    = suc n
max (suc n) (suc k) = suc (max n k)
```

- Define the function **maxl : List Nat -> Nat** which returns the maximum element of a list, or **zero** if the list is empty.

```
maxl : List Nat -> Nat
maxl nil = zero
maxl (cons n ns) = max n (maxl ns)
```

- Define the function **zip : {A B : Set} -> List A -> List B**

-> `List (Pair A B)` which converts two lists into a list of pairs of elements. If one of the lists is longer than the other the extra elements are ignored.

```
zip : {A B : Set} -> List A -> List B -> List (Pair A B)
zip (cons a as) (cons b bs) = cons (pair a b) (zip as bs)
zip _ _ = nil
```

- Define the function `intl : {A : Set} -> List A -> List A -> List A` which *interleaves* two lists `xxxx...` and `yyyy....` into `xyxyxy...`. If one of the lists is longer than the other the extra elements are ignored.

```
intl : {A : Set} -> List A -> List A -> List A
intl (cons a as) (cons b bs) = cons a (cons b (intl as bs)
)
intl _ _ = nil
```