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# System Programming

## (Chapter 2: Assembler)

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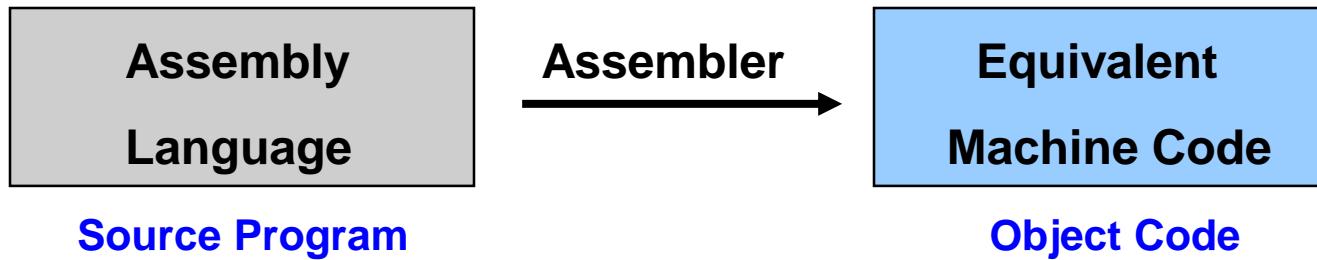
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# What is Assembler ?

## Fundamental Functions of an Assembler



- Translate mnemonic operation codes to their machine language equivalents.
- Assign machine addresses to symbol labels.

Other than the basic functions, designing of an assembler depends heavily upon the machine architecture (e.g., instruction format, addressing mode, etc.).

## Outline of Chapter 2

- Section 2.1 : Fundamental operations of an assembler using SIC.**
- Section 2.2 : Typical extensions via hardware considerations using SIC/XE.**
- Section 2.3 : Most commonly encountered machine-independent features.**
- Section 2.4 : Important design options of an assembler.**

# Information on the SIC Assembly Language

**Mnemonic instructions :** Section 1.3.1 and Appendix A.

Indexed addressing is indicated by adding the modifier “ ,**X** ” following the operand (line 160 of the program in next slide).

Lines beginning with “ . ” contain comments only.

**Assembler directives (pseudo-instructions) :** These statements are not translated into machine codes. Instead, they provide instructions to the assembler itself.

**START** : Specify name and starting address for the program.

**END** : Indicate the end of the source program and (optionally) specify the first executable instruction in the program.

**BYTE** : Generate character or hexadecimal constant, occupying as many bytes as needed to represent the constant.

**WORD** : Generate one-word integer constant.

**RESB** : Reserve the indicated number of bytes for a data area.

**RESW** : Reserve the indicated number of words for a data area.



# Example of a SIC Assembly Language Program

## Source statement

```

5      COPY      START    1000          COPY FILE FROM INPUT TO OUTPUT
10     FIRST     STL      RETADR        SAVE RETURN ADDRESS
15     CLOOP    JSUB     RDREC         READ INPUT RECORD
20     LDA      LENGTH        TEST FOR EOF (LENGTH = 0)
25     COMP     ZERO
30     JEQ      ENDFIL        EXIT IF EOF FOUND
35     JSUB     WRREC         WRITE OUTPUT RECORD
40     J       CLOOP
45     ENDFIL   LDA      EOF           LOOP
50     STA      BUFFER        INSERT END OF FILE MARKER
55     LDA      THREE         SET LENGTH = 3
60     STA      LENGTH        WRITE EOF
65     JSUB     WRREC         GET RETURN ADDRESS
70     LDL      RETADR        RETURN TO CALLER
75     RSUB
80     EOF      BYTE    C'EOF'
85     THREE    WORD    3
90     ZERO     WORD    0
95     RETADR   RESW    1
100    LENGTH   RESW    1
105    BUFFER   RESB    4096          LENGTH OF RECORD
110
115    .          .          4096-BYTE BUFFER AREA
120
125    RDREC    LDX     ZERO          SUBROUTINE TO READ RECORD INTO BUFFER
130    LDA      ZERO
135    RLOOP   TD      INPUT
140    JEQ      RLOOP
145    RD      INPUT
150    COMP     ZERO
155    JEQ      EXIT
160    STCH    BUFFER,X
165    TIX     MAXLEN
170    JLT      RLOOP
175    EXIT    STX     LENGTH
180    RSUB
185    INPUT   BYTE    X'F1'
190    MAXLEN WORD    4096          CLEAR LOOP COUNTER
195
200    .          .          CLEAR A TO ZERO
205    WRREC   LDX     ZERO
210    WLOOP   TD      OUTPUT
215    JEQ      WLOOP
220    LDCH    BUFFER,X
225    WD      OUTPUT
230    TIX     LENGTH
235    JLT      WLOOP
240    RSUB
245    OUTPUT  BYTE    X'05'
250
255    END      FIRST

```

The assembly code defines several subroutines:

- COPY**: Copies file from input to output. It saves the return address, reads the input record, tests for EOF (length 0), exits if EOF is found, writes the output record, loops, inserts the end-of-file marker, sets length to 3, writes EOF, gets the return address, and returns to the caller.
- RDREC**: Subroutine to read record into buffer. It clears the loop counter, clears register A to zero, tests the input device, loops until ready, reads character into register A, tests for end of record (X'00'), exits loop if EOR, stores character in buffer, loops unless max length has been reached, saves record length, and returns to the caller.
- WRREC**: Subroutine to write record from buffer. It clears the loop counter, tests the output device, loops until ready, gets character from buffer, writes character, loops until all characters have been written, returns to the caller, and codes for the output device.



# Assembly Language Program with Object Code

Line	Loc	Source statement			Object code
5	1000	COPY	START	1000	
10	1000	FIRST	STL	RETADR	141033
15	1003	CLOOP	JSUB	RDREC	482039
20	1006		LDA	LENGTH	001036
25	1009		COMP	ZERO	281030
30	100C		JEQ	ENDFIL	301015
35	100F		JSUB	WRREC	482061
40	1012		J	CLOOP	3C1003
45	1015	ENDFIL	LDA	EOF	00102A
50	1018		STA	BUFFER	0C1039
55	101B		LDA	THREE	00102D
60	101E		STA	LENGTH	0C1036
65	1021		JSUB	WRREC	482061
70	1024		LDL	RETADR	081033
75	1027		RSUB		4C0000
80	102A	EOF	BYTE	C'EOF'	454F46
85	102D	THREE	WORD	3	000003
90	1030	ZERO	WORD	0	000000
95	1033	RETADR	RESW	1	
100	1036	LENGTH	RESW	1	
105	1039	BUFFER	RESB	4096	
110		.	SUBROUTINE TO READ RECORD INTO BUFFER		
120		.			
125	2039	RDREC	LDX	ZERO	041030
130	203C		LDA	ZERO	001030
135	203F	RLOOP	TD	INPUT	E0205D
140	2042		JEQ	RLOOP	30203F
145	2045		RD	INPUT	D8205D
150	2048		COMP	ZERO	281030
155	204B		JEQ	EXIT	302057
160	204E		STCH	BUFFER, X	549039
165	2051		TIX	MAXLEN	2C205E
170	2054		JLT	RLOOP	38203F
175	2057	EXIT	STX	LENGTH	101036
180	205A		RSUB		4C0000
185	205D	INPUT	BYTE	X'F1'	F1
190	205E	MAXLEN	WORD	4096	001000
195		.	SUBROUTINE TO WRITE RECORD FROM BUFFER		
200		.			
210	2061	WRREC	LDX	ZERO	041030
215	2064	WLOOP	TD	OUTPUT	E02079
220	2067		JEQ	WLOOP	302064
225	206A		LDCH	BUFFER, X	509039
230	206D		WD	OUTPUT	DC2079
235	2070		TIX	LENGTH	2C1036
240	2073		JLT	WLOOP	382064
245	2076		RSUB		4C0000
250	2079	OUTPUT	BYTE	X'05'	05
255			END	FIRST	

Figure 2.2 Program from Fig. 2.1 with object code.



# A Simple SIC Assembler

**Steps for translating source program to object code (not necessarily in the order given) :**

1. Convert mnemonic operation codes to their machine language equivalents – e.g., translate STL to 14 (line 10).
2. Convert symbolic operands to their equivalent machine addresses – e.g., translate RETADR to 1033 (line 10).
3. Build the machine instructions in the proper format.
4. Convert the data constants specified in the source program into their internal machine representations – e.g., translate EOF to 454F46 (line 80).
5. Write the object program and the assembly listing.

- **All of these functions except number 2 can easily be accomplished by sequential processing of the source program. However, the translation of addresses such as line 10 presents a problem (forward reference) – need 2 passes.**



# Object Program Format (from Slide 5)

```

H COPY 00100000107A
T 0010001E1410334820390010362810303010154820613C100300102A0C103900102D
T 00101E150C10364820610810334C0000454F46000003000000
T 0020391E041030001030E0205D30203FD8205D2810303020575490392C205E38203F
T 0020571C1010364C0000F1001000041030E02079302064509039DC20792C1036
T 002073073820644C000005
E 001000

```

## Header Record:

Col. 1  
Col. 2-7  
Col. 8-13  
Col. 14-19

**H**  
Program name  
Starting address of object program (hex)  
Length of object program in bytes (hex)

## Text Record:

Col. 1  
Col. 2-7  
Col. 8-9  
Col. 10-69

**T**  
Starting address for object code in this record (hex)  
Length of object code in this record in bytes (hex)  
Object code, represented in hexadecimal (2 columns per byte)

## End Record:

Col. 1  
Col. 2-7

**E**  
Address of first executable instruction in object program (hex)



# Two Passes of SIC Assembler

## Pass 1

1. Assign addresses to all statements in the program.
2. Save the values (addresses) assigned to all labels for use in Pass2.
3. Perform some processing of assembler directives. (This includes processing that affects address assignment, such as determining the length of data areas defined by BYTE, RESW, etc)

## Pass 2

1. Assemble instructions (translating operation codes and looking up addresses).
2. Generate data values defined by BYTE, WORD, etc.
3. Perform processing of assembler directives not done during Pass 1.
4. Write the object program and the assembly listing.



# Data Structures for Assembler (1)

## LOCCTR (Location Counter)

Variable that is used to help in the assignment of addresses.

Initialized to the beginning address specified in the START statement.

The length of the assembled instruction is added to LOCCTR.

Current value of LOCCTR gives the address to be associated with the label.

## OPTAB (Operation Code Table)

Contains mnemonic operation codes and their machine language equivalents.

Also contains information about instruction format and length.

In Pass 1, OPTAB is used to look up and validate operation codes.

In Pass 2, it is used to translate the operation codes to machine language.

OPTAB is usually organized as a hash table (key: mnemonic operation code).

OPTAB is a static table – entries are not normally added to or deleted from it.



# Data Structures for Assembler (2)

## SYMTAB (Symbol Table)

Store values (addresses) assigned to labels.

Contains the flags to indicate error conditions (e.g., labels defined in two different places, etc.).

May also contain other information about the data area or instruction labeled.

In Pass 1, labels are entered into SYMTAB as they are encountered in the source program, along with their assigned addresses (from LOCCTR).

In Pass 2, symbols used as operands are looked up in SYMTAB to obtain the addresses to be inserted in the assembled instructions.

SYMTAB is usually organized as a hash table for efficiency of insertion and retrieval.

Hashing function should be designed carefully to take into account the facts that entries are not rarely deleted from this table and programmers often select many labels that have similar characteristics (e.g., labels starting with the same characters - LOOP1, LOOP2, LOOPA, or with the same length – A, X, Y, Z).



# Assembler Algorithm

**Pass 1 of the assembler usually write an intermediate file that contains each source statement together with its assigned address, error indicators, etc.**

**This file is the input to the Pass 2.**

**Pointers into OPTAB and SYMTAB may be retained for each operation code and symbol used that allows us to avoid the need to repeat many of the table searching operations.**

**Slides 12 and 13 show the logic flow of the two passed of our assembler. (Text p. 53 and p. 54)**

**You are strongly urged to follow through the logic in these algorithms, applying them by hand to the program in Slide 4 to produce the object program of Slide 7.**



# Algorithm for Pass 1 of Assembler

**Pass 1:**

```

begin
    read first input line
    if OPCODE = 'START' then
        begin
            save #[OPERAND] as starting address
            initialize LOCCTR to starting address
            write line to intermediate file
            read next input line
        end {if START}
    else
        initialize LOCCTR to 0
    while OPCODE ≠ 'END' do
        begin
            if this is not a comment line then
                begin
                    if there is a symbol in the LABEL field then
                        begin
                            search SYMTAB for LABEL
                            if found then
                                set error flag (duplicate symbol)
                            else
                                insert (LABEL, LOCCTR) into SYMTAB
                        end {if symbol}
                    search OPTAB for OPCODE
                    if found then
                        add 3 {instruction length} to LOCCTR
                    else if OPCODE = 'WORD' then
                        add 3 to LOCCTR
                    else if OPCODE = 'RESW' then
                        add 3 * #[OPERAND] to LOCCTR
                    else if OPCODE = 'RESB' then
                        add #[OPERAND] to LOCCTR
                    else if OPCODE = 'BYTE' then
                        begin
                            find length of constant in bytes
                            add length to LOCCTR
                        end {if BYTE}
                    else
                        set error flag (invalid operation code)
                end {if not a comment}
                write line to intermediate file
                read next input line
            end {while not END}
            write last line to intermediate file
            save (LOCCTR - starting address) as program length
        end {Pass 1}
    
```



# Algorithm for Pass 2 of Assembler

**Pass 2:**

```

begin
    read first input line {from intermediate file}
    if OPCODE = 'START' then
        begin
            write listing line
            read next input line
        end {if START}
    write Header record to object program
    initialize first Text record
    while OPCODE ≠ 'END' do
        begin
            if this is not a comment line then
                begin
                    search OPTAB for OPCODE
                    if found then
                        begin
                            if there is a symbol in OPERAND field then
                                begin
                                    search SYMTAB for OPERAND
                                    if found then
                                        store symbol value as operand address
                                    else
                                        begin
                                            store 0 as operand address
                                            set error flag (undefined symbol)
                                        end
                                end {if symbol}
                            else
                                store 0 as operand address
                                assemble the object code instruction
                            end {if opcode found}
                        else if OPCODE = 'BYTE' or 'WORD' then
                            convert constant to object code
                        if object code will not fit into the current Text record then
                            begin
                                write Text record to object program
                                initialize new Text record
                            end
                            add object code to Text record
                        end {if not comment}
                    write listing line
                    read next input line
                end {while not END}
    write last Text record to object program
    write End record to object program
    write last listing line
end {Pass 2}

```



# Example of a SIC/XE Program

Line	Source statement				
10	COPY FIRST	START	0		COPY FILE FROM INPUT TO OUTPUT
12		STL	RETADR		SAVE RETURN ADDRESS
13		LDB	#LENGTH		ESTABLISH BASE REGISTER
15	CLOOP	+JSUB	RDREC		READ INPUT RECORD
20		LDA	LENGTH		TEST FOR EOF (LENGTH = 0)
25		COMP	#0		EXIT IF EOF FOUND
30		JEQ	ENDFIL		WRITE OUTPUT RECORD
35		+JSUB	WRREC		LOOP
40		J	CLOOP		INSERT END OF FILE MARKER
45	ENDFIL	LDA	EOF		SET LENGTH = 3
50		STA	BUFFER		WRITE EOF
55		LDA	#3		RETURN TO CALLER
60		STA	LENGTH		LENGTH OF RECORD
65		+JSUB	WRREC		4096-BYTE BUFFER AREA
70		J	RETADR		
80	EOF	BYTE	C'EOF'		
95	RETADR	RESW	1		
100	LENGTH	RESW	1		
105	BUFFER	RESB	4096		
110		SUBROUTINE TO READ RECORD INTO BUFFER			
120					
125	RDREC	CLEAR	X		CLEAR LOOP COUNTER
130		CLEAR	A		CLEAR A TO ZERO
132		CLEAR	S		CLEAR S TO ZERO
133		+LDT	#4096		
135	RLOOP	TD	INPUT		TEST INPUT DEVICE
140		JEQ	RLOOP		LOOP UNTIL READY
145		RD	INPUT		READ CHARACTER INTO REGISTER A
150		COMPR	A,S		TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT		EXIT LOOP IF EOR
160		STCH	BUFFER,X		STORE CHARACTER IN BUFFER
165		TIXR	T		LOOP UNLESS MAX LENGTH
170		JLT	RLOOP		HAS BEEN REACHED
175	EXIT	STX	LENGTH		SAVE RECORD LENGTH
180		RSUB			RETURN TO CALLER
185	INPUT	BYTE	X'F1'		CODE FOR INPUT DEVICE
195		SUBROUTINE TO WRITE RECORD FROM BUFFER			
200					
205					
210	WRREC	CLEAR	X		CLEAR LOOP COUNTER
212		LDT	LENGTH		
215	WLOOP	TD	OUTPUT		TEST OUTPUT DEVICE
220		JEQ	WLOOP		LOOP UNTIL READY
225		LDCH	BUFFER,X		GET CHARACTER FROM BUFFER
230		WD	OUTPUT		WRITE CHARACTER
235		TIXR	T		LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP		HAVE BEEN WRITTEN
245		RSUB			RETURN TO CALLER
250	OUTPUT	BYTE	X'05'		CODE FOR OUTPUT DEVICE
255		END	FIRST		



# Example Program with Object Code

Line	Loc		Source statement		Object code
5	0000	COPY	START	0	
10	0000	FIRST	STL	RETADR	17202D
12	0003		LDB	#LENGTH	69202D
13			BASE	LENGTH	
15	0006	CLOOP	+JSUB	RDREC	4B101036
20	000A		LDA	LENGTH	032026
25	000D		COMP	#0	290000
30	0010		JEQ	ENDFIL	332007
35	0013		+JSUB	WRREC	4B10105D
40	0017		J	CLOOP	3F2FEC
45	001A	ENDFIL	LDA	EOF	032010
50	001D		STA	BUFFER	0F2016
55	0020		LDA	#3	010003
60	0023		STA	LENGTH	0F200D
65	0026		+JSUB	WRREC	4B10105D
70	002A		J	RETADR	3E2003
80	002D	EOF	BYTE	C'EOF'	454F46
95	0030	RETADR	RESW	1	
100	0033	LENGTH	RESW	1	
105	0036	BUFFER	RESB	4096	
110				SUBROUTINE TO READ RECORD INTO BUFFER	
120					
125	1036	RDREC	CLEAR	X	B410
130	1038		CLEAR	A	B400
132	103A		CLEAR	S	B440
133	103C		+LDT	#4096	75101000
135	1040	RLOOP	TD	INPUT	E32019
140	1043		JEQ	RLOOP	332FFA
145	1046		RD	INPUT	DB2013
150	1049		COMPR	A,S	A004
155	104B		JEQ	EXIT	332008
160	104E		STCH	BUFFER,X	57C003
165	1051		TIXR	T	B850
170	1053		JLT	RLOOP	3B2FEA
175	1056	EXIT	STX	LENGTH	134000
180	1059		RSUB		4F0000
185	105C	INPUT	BYTE	X'F1'	F1
195				SUBROUTINE TO WRITE RECORD FROM BUFFER	
200					
205					
210	105D	WRREC	CLEAR	X	B410
212	105F		LDT	LENGTH	774000
215	1062	WLOOP	TD	OUTPUT	E32011
220	1065		JEQ	WLOOP	332FFA
225	1068		LDCH	BUFFER,X	53C003
230	106B		WD	OUTPUT	DF2008
235	106E		TIXR	T	B850
240	1070		JLT	WLOOP	3B2FEF
245	1073		RSUB		4F0000
250	1076	OUTPUT	BYTE	X'05'	05
255			END	FIRST	



# Review of SIC/XE Addressing Modes

The addressing modes supported by the SIC/XE can be largely classified by two issues: how to calculate TA (Target Address) and how to interpret the calculated TA.

## How to calculate TA (Use b, p, x bits)

Direct addressing mode : TA = address  $\Rightarrow b = p = 0/1$

PC relative addressing mode : TA = (PC) + address  $\Rightarrow b = 0, p = 1$

Base relative addressing mode : TA = (B) + address  $\Rightarrow b = 1, p = 0$

## How to interpret TA (Use i, n bits)

Simple addressing mode : (TA)  $\Rightarrow i = 1, n = 1$

Indirect addressing mode : ((TA))  $\Rightarrow i = 0, n = 1$

Immediate addressing mode : TA  $\Rightarrow i = 1, n = 0$



# Features of SIC/XE Assembler

**Indirect addressing is indicated by adding prefix @ to the operand (line 70).**

**Immediate operands are denoted with the prefix # (line 25, 55, 133).**

**Instructions that refer to memory are normally assembled using either the PC relative or the base relative mode.**

**If the displacement required for both PC relative and base relative addressing are too large to fit into a 3-byte instruction, then the 4-byte extended format (format 4) must be used.**

**The extended instruction format is specified with the prefix + added to the operation code (line 15, 35, 65).**

**For speed-up, use register-to-register instructions, and immediate/indirect addressing as much as possible.**



# Assembling SIC/XE Instructions (1)

START statement now specifies a beginning address of 0 : re-locatable program => will be covered in the next section.

Translation of register-to-register instructions such as CLEAR (line 125) and COMPR (line 150) presents no new problems.

The conversion of register mnemonics to numbers can be done with a separate table; or SYMAB may be preloaded with the register names (A, X, etc) and their values (0, 1, etc).

Most of the register-to-memory instructions are assembled using either PC relative or base relative addressing.

If the resulting displacement is larger than the values in 12-bit field (For PC relative:  $-2048 \leq \text{disp} \leq 2047$ , for base relative:  $0 \leq \text{disp} \leq 4095$ ), the 4-byte format 4 (20 bit field) is used.

It is the programmer's responsibility that he/she should put + to represent format 4 instructions.

e.g., ) 15 0006 CLOOP +JSUB RDREC 4B101036



# Assembling SIC/XE Instructions (2)

## PC relative addressing

e.g. 1) 10 0000 FIRST STL RETADR 17202D

(PC) = 0003 (The PC is advanced *after* each instruction is fetched and *before* it is executed. Thus PC will contain the address of the *next* instruction)

Address of RETADR = 0030. Therefore, displacement =  $30 - 3 = 2D$ .

Simple addressing : i = 1, n = 1.

PC relative addressing : p = 1, b = 0, x = 0 (no indexing).

e.g. 2) 40 0017 J CLOOP 3F2FEC

(PC) = 001A

Address of CLOOP = 0006. Therefore displacement =  $6 - 1A = -14$ .

Negative values will be represented by 2's complement:  $-14 \Rightarrow FEC$ .



# Assembling SIC/XE Instructions (3)

**Base relative addressing :** similar to PC relative but differs in that the assembler does not know the content of the B register and it is the responsibility of the programmer to tell the assembler what the base register will contain during execution (via BASE assembler directive – line 13)

e.g. 1) **12 0003 LDB #LENGTH 69202D**  
**13 BASE LENGTH**

If we are to use register B for another purpose, we have to use the assembler directive NOBASE to inform the assembler that the contents of the base register can no longer be relied upon for addressing.

e.g. 2) **160 104E STCH BUFFER, X 57C003**

(B) = 0033 (the address of LENGTH)

Address of BUFFER = 0036. Therefore displacement =  $36 - 33 = 3$ .

x = 1, b = 1, p = 0, i = 1, n = 1 : Simple Base relative addressing.



# Assembling SIC/XE Instructions (4)

Choice between PC relative addressing and base relative addressing mode : arbitrary.

Try PC relative first, try base relative next. If it is not possible to represent the displacement with 12 bits, use format 4.

e.g. 1) 20 000A LDA LENGTH 032026 : PC relative

e.g. 2) 175 1056 EXIT STX LENGTH 134000 : Base relative

Immediate addressing

e.g. 1) 55 0020 LDA #3 010003

i = 1, other bits = 0.

e.g. 2) 133 103C +LDT #4096 75101000

4096 can't be represented with 12 bits. Use format 4.

e.g. 3) 12 0003 LDB #LENGTH 69202D

PC relative + immediate addressing.



# Assembling SIC/XE Instructions (5)

## Indirect addressing

e.g.) **70 002A J @RETADDR 3E2003**

PC relative + indirect addressing.

(PC) = 002D

Address of RETADDR = 30. Therefore, displacement =  $30 - 2D = 3$ .



# Program Relocation

It is often desirable to have more than one program at a time sharing the memory and other resources of the machine.  
(Multi-programming) => The actual starting address of the program is not known until load time.

e.g.) 5      1000    COPY START 1000 <= should be loaded at 1000  
              55      101B                  LDA      THREE                  00102D

- If this program is loaded at 2000 and memory address 102D should be modified.
- Since the assembler does not know the actual location where the program will be loaded, it identifies those parts of the object program and let the loader know.
- An object program that contains the information necessary to perform this modification is called a *re-locatable* program.

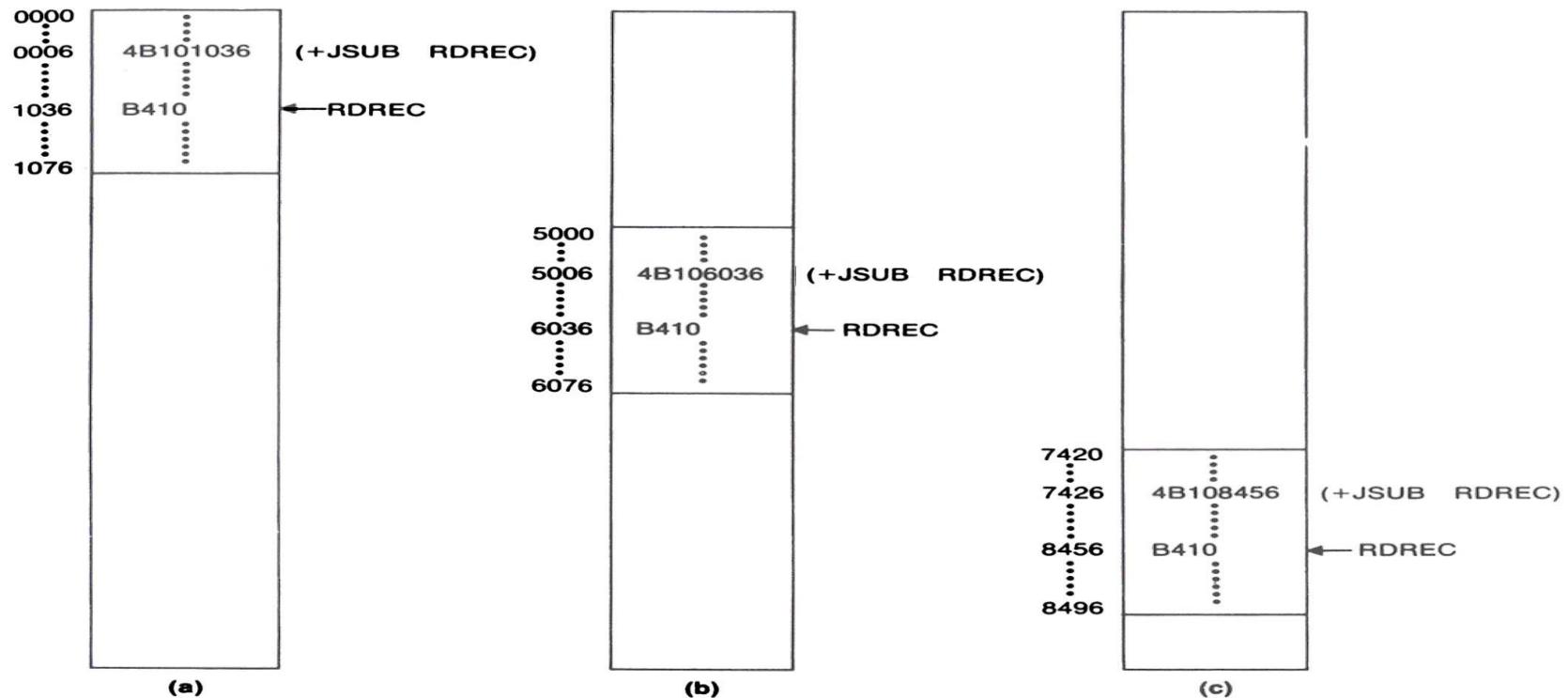
# Example of a Program Relocation

Fig 2.5 and Fig 2.6

15 CLOOP +JSUB RDREC 4B101036

...

125 RDREC CLEAR X B410



# Modification Record (1)

Relocation can be solved in the following way:

When the assembler generates the object code for the JSUB instruction we are considering, it will insert the address of RDREC *relative to the start of the program* (starting address should be 0 – START 0 instruction).

The assembler also produce a command for the loader, instructing it to add the beginning address of the program to the address field in the JSUB instruction at load time. => **Modification record**.

**Modification Record:**

Col. 1      M

Col. 2-7      Starting location of the address field to be modified, relative to the beginning of the program (hex).

Col. 8-9      Length of the address field to be modified, in half bytes (hex).



# Modification Record (2)

The length is stored in **half-bytes** because the address field to be modified may not occupy an integral number of bytes. (e.g., 20 bit address -> 5 half bytes)

The starting location is the **location of the byte** containing the leftmost bits of the address field to be modified.

If the length contains an odd number of half bytes, it is assumed to begin in the middle of the first byte at the starting location.

```
HCOPY 00000001077
T0000001D17202D69202D4B1010360320262900003320074B10105D3F2FEC032010
T00001D130F20160100030F200D4B10105D3E2003454F46
T0010361DB410B400B44075101000E32019332FFADB2013A00433200857C003B850
T0010531D3B2FEA1340004F0000F1B410774000E32011332FFA53C003DF2008B850
T001070073B2FEF4F00005
M00000705
M00001405
M00002705
E000000
```



# Rules for Modification

The only parts of the program that require modification at load time are those that specify direct (as opposed to relative) addresses.

For this SIC/XE program, the only such direct addresses are found in extended format (4 bytes) instructions.

This is an advantage of relative addressing.

# Program with Additional Features (Fig 2.9)

Line	Source statement				
5	COPY	START	0		COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR		SAVE RETURN ADDRESS
13		LDB	#LENGTH		ESTABLISH BASE REGISTER
14		BASE	LENGTH		
15	CLOOP	+JSUB	RDREC		READ INPUT RECORD
20		LDA	LENGTH		TEST FOR EOF (LENGTH = 0)
25		COMP	#0		
30		JEQ	ENDFIL		EXIT IF EOF FOUND
35		+JSUB	WRREC		WRITE OUTPUT RECORD
40		J	CLOOP		LOOP
45	ENDFIL	LDA	=C'EOF'		INSERT END OF FILE MARKER
50		STA	BUFFER		
55		LDA	#3		SET LENGTH = 3
60		STA	LENGTH		
65		+JSUB	WRREC		WRITE EOF
70		J	@RETADR		RETURN TO CALLER
93		LTORG			
95	RETADR	RESW	1		
100	LENGTH	RESW	1		LENGTH OF RECORD
105	BUFFER	RESB	4096		4096-BYTE BUFFER AREA
106	BUFEND	EQU			
107	MAXLEN	EQU	BUFEND-BUFFER		MAXIMUM RECORD LENGTH
110					
115		SUBROUTINE TO READ RECORD INTO BUFFER			
120					
125	RDREC	CLEAR	X		CLEAR LOOP COUNTER
130		CLEAR	A		CLEAR A TO ZERO
132		CLEAR	S		CLEAR S TO ZERO
133		+LDT	#MAXLEN		
135	RLOOP	TD	INPUT		TEST INPUT DEVICE
140		JEQ	RLOOP		LOOP UNTIL READY
145		RD	INPUT		READ CHARACTER INTO REGISTER A
150		COMPR	A,S		TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT		EXIT LOOP IF EOR
160		STCH	BUFFER,X		STORE CHARACTER IN BUFFER
165		TIXR	T		LOOP UNTIL MAX LENGTH
170		JLT	RLOOP		HAS BEEN REACHED
175	EXIT	STX	LENGTH		SAVE RECORD LENGTH
180		RSUB			RETURN TO CALLER
185	INPUT	BYTE	X'F1'		CODE FOR INPUT DEVICE
195					
200		SUBROUTINE TO WRITE RECORD FROM BUFFER			
205					
210	WRREC	CLEAR	X		CLEAR LOOP COUNTER
212		LDT	LENGTH		
215	WLOOP	TD	=X'05'		TEST OUTPUT DEVICE
220		JEQ	WLOOP		LOOP UNTIL READY
225		LDCH	BUFFER,X		GET CHARACTER FROM BUFFER
230		WD	=X'05'		WRITE CHARACTER
235		TIXR	T		LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP		HAVE BEEN WRITTEN
245		RSUB			RETURN TO CALLER
255		END	FIRST		



# Example Program with Object Code (Fig 2.10)

Line	Loc		Source statement	Object code
5	0000	COPY	START O	
10	0000	FIRST	STL RETADR	17202D
13	0003		LDB #LENGTH	69202D
14			BASE LENGTH	
15	0006	CLOOP	+JSUB RDREC	4B101036
20	000A		LDA LENGTH	032026
25	000D		COMP #0	290000
30	0010		JEQ ENDFIL	332007
35	0013		+JSUB WRREC	4B10105D
40	0017		J CLOOP	3F2FEC
45	001A	ENDFIL	LDA =C'EOF'	032010
50	001D		STA BUFFER	0F2016
55	0020		LDA #3	010003
60	0023		STA LENGTH	0F200D
65	0026		+JSUB WRREC	4B10105D
70	002A		J @RETADR	3E2003
93	002D	*	LTORG =C'EOF'	454F46
95	0030	RETADR	RESW 1	
100	0033	LENGTH	RESW 1	
105	0036	BUFFER	RESB 4096	
106	1036	BUFEND	EQU *	
107	1000	MAXLEN	EQU BUFEND-BUFFER	
110				
115			SUBROUTINE TO READ RECORD INTO BUFFER	
120				
125	1036	RDREC	CLEAR X	B410
130	1038		CLEAR A	B400
132	103A		CLEAR S	B440
133	103C		+LDT #MAXLEN	75101000
135	1040	RLOOP	TD INPUT	E32019
140	1043		JEQ RLOOP	332FFA
145	1046		RD INPUT	DB2013
150	1049		COMPR A,S	A004
155	104B		JEQ EXIT	332008
160	104E		STCH BUFFER,X	57C003
165	1051		TIXR T	B850
170	1053		JLT RLOOP	3B2FEA
175	1056	EXIT	STX LENGTH	134000
180	1059		RSUB	4F0000
185	105C	INPUT	BYTE X'F1	F1
195				
200			SUBROUTINE TO WRITE RECORD FROM BUFFER	
205				
210	105D	WRREC	CLEAR X	B410
212	105F		LDT LENGTH	774000
215	1062	WLOOP	TD =X'05'	E32011
220	1065		JEQ WLOOP	332FFA
225	1068		LDCH BUFFER,X	53C003
230	106B		WD =X'05'	DF2008
235	106E		TIXR T	B850
240	1070		JLT WLOOP	3B2FEF
245	1073		RSUB	4F0000
255			END FIRST	
	1076		=X'05'	05



# Literals (1) (강의제외 참조용)

## Using Constants (Review)

Immediate addressing

e.g. 1) LDA #3

BYTE directive : Generate *character* or *hexadecimal* constant.

e.g. 1) LABEL1 BYTE C'EOF'

e.g. 2) LABEL2 BYTE X'05'

WORD directive: Generate *one-word integer* constant.

e.g. 1) LABEL1 WORD 3

- **Literal : a constant operand used as a part of the instruction. Identified with the prefix =, followed by a specification of the literal value (same notation as in the BYTE statement).**

e.g. 1) 45	001A	ENDFILL	LDA	=C'EOF'	032010
------------	------	---------	-----	---------	--------

e.g. 2) 215	1062	WLOOP	TD	=X'05'	E32011
-------------	------	-------	----	--------	--------



# Literals (2) (강의제외 참조용)

## Difference between an immediate operand and a literal.

Immediate addressing : The operand value is assembled as part of the machine instruction.

Literal : The assembler generates the specified value as a constant at some other memory location. The *address* of this generated constant is used as the target address for the machine instruction. => same as the constant definition with a label attached to it (Compare line 45 and 215 in Fig 2.10 and Fig 2.6)

- **Where are the literals placed ? => *literal pools* (all of the literal operands used in a program are gathered together into one or more literal pools)**
- **Location of the literal pools**
  - At the end of the program (following the END statement). – Check Fig 2.10
  - After the LTORG assembler directive. Create a literal pool that contains all of the literal operands used since the previous LTORG (or the beginning of the program). – Check line 93 in Fig 2.10.



# Literals (3) (강의제외 참조용)

## Why LTORG ?

Avoid placing literal operand too far away from the instruction so that we can use PC relative addressing mode.

Avoid having to use extended format instructions when referring to the literals.

### ■ Most assemblers recognize duplicate literals :

- Comparison of the character strings defining them. (line 215 and 230 in Fig 2.10)
- Additional saving is possible if we look at the generated data value (e.g., =C'EOF' and =X'454F46' specify identical operand values). If the benefits realized are not enough to justify the additional complexity, store both literals.

### ■ Literals whose value depends upon their location in the program (often denoted by the symbol \*).

- e.g.) BASE \*  
              LDB      =\*
- Should be careful when detecting duplicate literals. If this literal appears on line 13 in Fig 10, it will specify an operand with value 0003. If the same literal appears on line 55, it will specify an operand with value 0020.



## Handling Literal Operands in the Assembler(강의제외 참조용)

### Basic Data Structure : a literal table called LITTAB.

Contains the literal name, the operand value and length, and the address assigned to the operand when it is placed in a literal pool.

LITTAB is often organized as a hash table, using the literal name or value as the key.

### Steps for handling literal operands.

#### Pass 1 :

Search LITTAB for the specified literal name.

If the literal is already there, no action is needed; If not, the literal is added to LITTAB (leaving the address unassigned).

When pass 1 encounters a LTORG or the end of the program, assign addresses to the literals currently in the table. The location counter is updated.

#### Pass 2 :

Search LITTAB for each literal operand encountered and insert the values at the appropriate places in the object program.

If a literal value represents an address, the assembler also generates the appropriate Modification Record.



# EQU (Equate) Assembler Directive (1)

**EQU (Equate) allows programmer to define symbols and specify their values.**

Syntax : *symbol*    EQU    *value*

The symbol is also entered into the SYMTAB.

The value may be given as a constant or as any expression involving constants and previously defined symbols.

**One common use of EQU is to establish symbolic names that can be used for improved readability. (Compare with #define in C)  
(e.g.)**

+LDT #4096 can be written as

```
MAXLEN EQU 4096  
        +LDT #MAXLEN
```



## EQU (Equate) Assembler Directive (2)

**Another common use of EQU is in defining names for registers if the assembler expects register numbers instead of names.**

(e.g.)

RMO 0, 1 can be written as RMO A, X if

A	EQU	0
X	EQU	1
L	EQU	2

..... are defined.

**Case for a machine with general purpose registers (e.g., R0, R1, R2, ...)**

(e.g.)

BASE	EQU	R1
COUNT	EQU	R2
INDEX	EQU	R3

# Restrictions Common to EQU

In the case of EQU,

All symbols used on the right-hand side of the statement must have been defined previously in the program.

(e.g.) The sequence

ALPHA RESW 1

BETA EQU ALPHA --- is valid, however

BETA EQU ALPH

ALPHA RESW 1 --- is not valid.



## ORG (Origin) Assembler Directive (1) (강의제외 참조용)

**ORG (Origin) is used to indirectly assign values to symbols.**

Syntax : ORG value

The assembler resets its location counter (LOCCTR) to the specified value =>  
This will affect the values of all labels defined until the next ORG.

The value is a constant or an expression involving constants and previously defined symbols.

**Consider the case where we are defining a symbol table with the following structure:**

	Symbol	Value	Flags
STAB (100 entries)	6 byte symbol	1 word	2 bytes

.....

STAB      RESB      1100



## ORG (Origin) Assembler Directive (2) (강의제외 참조용)

**Two ways to efficiently access the symbol table (STAB).**

**Method 1:** Use indexed addressing (X register) with EQU directive.

(e.g.) Define

SYMBOL	EQU	STAB
VALUE	EQU	STAB+6
FLAGS	EQU	STAB+9

And access each field (e.g., VALUE field) using “**LDA VALUE, X**”, where X will contains the offset of the desired entry from the beginning of the table.

**Method 2:** Use ORG directive.

(e.g.) Define

STAB	RESB	1100	
	ORG	STAB	: Reset LOCCTR to STAB
SYMBOL	RESB	6	: 1 <sup>st</sup> 6 bytes are reserved for SYMBOL
VALUE	RESW	1	: 2 <sup>nd</sup> 3 bytes are reserved for VALUE
FLAGS	RESB	2	: 3 <sup>rd</sup> 2 bytes are reserved for FLAGS
	ORG	STAB+1100	: Set LOCCTR back to its previous value : STAB+1100 can be omitted (simply ORG)



# Expressions (1)

## Definitions

**Expression** : Terms with operators such as +, -, \*, /.

**Terms** : Constants, user-defined symbols, or special term (e.g., \* - value of current LOCCTR)

## Examples from Fig 2.10

(e.g. 1) 106            BUFEND            EQU        \*

(e.g. 2) 107            MAXLEN            EQU        BUFEND-BUFFER

**The values of terms and expressions are either relative or absolute.**

## Terms

Constant : absolute term.

Symbols on instructions and data areas : relative terms.

Symbol whose value is given by EQU : either absolute term or relative term.

# Expressions (2)

## Expressions

An expression that contains only absolute terms : absolute expression.

- (e.g.) 10-3, DATA-10 (where DATA is declared as a constant).

Absolute expressions may also contain relative terms if the relative terms occur in pairs and the terms in each such pair have opposite signs.

- (e.g.) 10 + BUFEND - BUFFER

A relative term is the one in which all of the relative terms except one can be paired as described above; the remaining unpaired relative term must have a positive sign.

- (e.g.) BUFEND – BUFFER + LENGTH

No relative terms may enter into a multiplication or division operation.

Expressions that do not meet the conditions given above generate errors.

- (e.g.) BUFEND + BUFFER, 100 – BUFFER, or 3 \* BUFFER

A relative term represents some location within the program. When relative terms are paired with opposite signs, the dependency on the program starting address is canceled out -> absolute value.



## Expressions (3)

To determine the type of an expression, we must keep track of the types for all symbols.

For this purpose, we need a flag in the symbol table to indicate the type of value (absolute or relative) in addition to the value itself.

Symbol	Type	Value
RETADR	R	0030
BUFFER	R	0036
BUFEND	R	1036
MAXLEN	A	1000

With this information, the assembler can easily determine the type of each expression and generate Modification records for relative values.



## Program Blocks

So far, all assembled programs were treated as a unit (single block of object code) although they contained subroutines, data areas, etc. => Machine instructions and data appeared in the same order as they were written in the source program.

Some assemblers allow the generated machine instructions and data to appear in the object program in a different order by creating several independent parts in the program.

**Program Blocks** : Segments of code that are rearranged within a single object program unit.

Syntax : USE      *Block Name*

Blocks with same names are treated as same blocks. If the name is omitted, it is the unnamed block.

Example of program blocks : Fig. 2.11 and 2.12 (line 92, 103, 123, 183) shows 3 blocks, *Unnamed*, *CDATA*, *CBLKS*.



# An Example with Multiple Program Blocks

Line	Source statement				
5	COPY	START	0		COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR		SAVE RETURN ADDRESS
15	CLOOP	JSUB	RDREC		READ INPUT RECORD
20		LDA	LENGTH		TEST FOR EOF (LENGTH = 0)
25		COMP	#0		
30		JEQ	ENDFIL		EXIT IF EOF FOUND
35		JSUB	WRREC		WRITE OUTPUT RECORD
40		J	CLOOP		LOOP
45	ENDFIL	LDA	=C'EOF'		INSERT END OF FILE MARKER
50		STA	BUFFER		
55		LDA	#3		SET LENGTH = 3
60		STA	LENGTH		
65		JSUB	WRREC		WRITE EOF
70		J	@RETADR		RETURN TO CALLER
92		USE	CDATA		
95	RETADR	RESW	1		
100	LENGTH	RESW	1		LENGTH OF RECORD
103		USE	CBLKS		
105	BUFFER	RESB	4096		4096-BYTE BUFFER AREA
106	BUFEND	EQU	*		FIRST LOCATION AFTER BUFFER
107	MAXLEN	EQU	BUFEND-BUFFER		MAXIMUM RECORD LENGTH
110		SUBROUTINE TO READ RECORD INTO BUFFER			
120					
123		USE			
125	RDREC	CLEAR	X		CLEAR LOOP COUNTER
130		CLEAR	A		CLEAR A TO ZERO
132		CLEAR	S		CLEAR S TO ZERO
133		+LDT	#MAXLEN		
135	RLOOP	TD	INPUT		TEST INPUT DEVICE
140		JEQ	RLOOP		LOOP UNTIL READY
145		RD	INPUT		READ CHARACTER INTO REGISTER A
150		COMPR	A,S		TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT		EXIT LOOP IF EOR
160		STCH	BUFFER,X		STORE CHARACTER IN BUFFER
165		TIXR	T		LOOP UNLESS MAX LENGTH
170		JLT	RLOOP		HAS BEEN REACHED
175	EXIT	STX	LENGTH		SAVE RECORD LENGTH
180		RSUB			RETURN TO CALLER
183	INPUT	USE	CDATA		
185		BYTE	X'F1'		CODE FOR INPUT DEVICE
195		SUBROUTINE TO WRITE RECORD FROM BUFFER			
200					
205		USE			
208	WRREC	CLEAR	X		CLEAR LOOP COUNTER
210		LDI	LENGTH		
212	WLOOP	TD	=X'05'		TEST OUTPUT DEVICE
215		JEQ	WLOOP		LOOP UNTIL READY
220		LDCH	BUFFER,X		GET CHARACTER FROM BUFFER
225		WD	=X'05'		WRITE CHARACTER
230		TIXR	T		LOOP UNTIL ALL CHARACTERS
235		JLT	WLOOP		HAVE BEEN WRITTEN
240		RSUB	USE		RETURN TO CALLER
245			CDATA		
252		LTOORG	FIRST		
253	END				



# An Example with Object Code

Line	Loc/Block		Source statement		Object code
5	0000	0	COPY	START	0
10	0000	0	FIRST	STL	RETADR
15	0003	0	CLOOP	JSUB	4B2021
20	0006	0		LDA	RDREC
25	0009	0		COMP	032060
30	000C	0		#0	290000
35	000F	0		JEQ	ENDFIL
40	0012	0		JSUB	332006
45	0015	0	ENDFIL	J	WRREC
50	0018	0		LDA	4B203B
55	001B	0		=C'EOF'	3F2FEE
60	001E	0		STA	032055
65	0021	0		LDA	0F2056
70	0024	0		STA	#3
92	0000	1		JSUB	010003
95	0000	1		J	LENGTH
100	0003	1	RETADR	LDA	0F2048
103	0000	2	LENGTH	STA	WRREC
105	0000	2	BUFFER	JSUB	4B2029
106	1000	2	BUFEND	J	GRETADR
107	1000	2	MAXLEN	USE	3E203F
110				RESW	CDATA
115				1	
120				1	
123	0027	0		RESW	
125	0027	0	RDREC	USE	
130	0029	0		CLEAR	B410
132	002B	0		CLEAR	B400
133	002D	0		CLEAR	B440
135	0031	0	RLOOP	+LDT	#MAXLEN
140	0034	0		TD	75101000
145	0037	0		JEQ	E32038
150	003A	0		RD	332FFA
155	003C	0		COMPR	DB2032
160	003F	0		A,S	A004
165	0042	0		JEQ	332008
170	0044	0		STCH	57A02F
175	0047	0	EXIT	TIXR	B850
180	004A	0		JLT	3B2FEA
183	0006	1		STX	13201F
185	0006	1	INPUT	RSUB	4F0000
195				USE	
200				BYTE	CDATA
205				X'F1'	F1
208	004D	0			
210	004D	0	WRREC		
212	004F	0		USE	
215	0052	0	WLOOP	CLEAR	B410
220	0055	0		LDT	LENGTH
225	0058	0		TD	772017
230	005B	0		=X'05'	E3201B
235	005E	0		JEQ	332FFA
240	0060	0		LDCH	53A016
245	0063	0		BUFFER,X	DF2012
252	0007	1		WD	3B2FEE
253	0007	1		=X'05'	4F0000
	000A	1		LTORG	
				=C'EOF'	454F46
				=X'05'	05
255			END	FIRST	



## Pass 1 for a Program with Program Blocks

### Pass 1

A separate location counter (LOCCTR) for each program block is maintained (Fig 2.12).

The LOCCTR for a block is initialized to 0 when the block is first begun (Line 92 in Fig 2.12).

The current LOCCTR value is saved when switching to another block, and the saved value is restored when resuming a previous block (Line 123 in Fig 2.12).

The address of each label is relative to the start of each block.

The latest LOCCTR value for each block indicates the length of that block.

At the end of Pass 1, the assembler constructs a table as shown below :

Block Name	Block Number	Address	Length
(default)	0	0000	0066
CDATA	1	0066	000B
CBLKS	2	0071	1000



## Pass 2 for a Program with Program Blocks

### Pass 2

The assembler needs the address for each symbol relative to the start of the object program (not the start of each block).

The assembler simply adds the location of the symbol, relative to the start of each block, to the assigned block starting address.

(e.g. 1)

Line	Loc	Block	Source Statement	Object Code
20	0006	0	LDA LENGTH	032060
.....				
100	0003	1	LENGH RESW 1	

- Relative address of LENGTH relative to block 1 (CDATA) = 0003
- The starting address for CDATA is 0066.
- Target address of LENGTH =  $0003 + 0066 = 0069$ .
- Use PC-relative addressing (PC = 009).
- Displacement =  $0069 - 0009 = 0060$ .



## Advantages of using Program Blocks

### Reduce the addressing problems.

No need to use extended format (line 15, 35, 65).

No need to use base relative addressing (no LDB and BASE statements).

### Reduce the problems of placing literals.

Include a LTORG statement in the CDATA block to be sure that the literals are placed ahead of any large data areas.

### The use of program blocks is one way of satisfying both machine consideration and human factors :

Machine consideration: Part of the object program appear in memory in a particular order.

Human factors: Improve readability by putting data areas close to the source statements that reference them -> Need long jump and extended addressing.



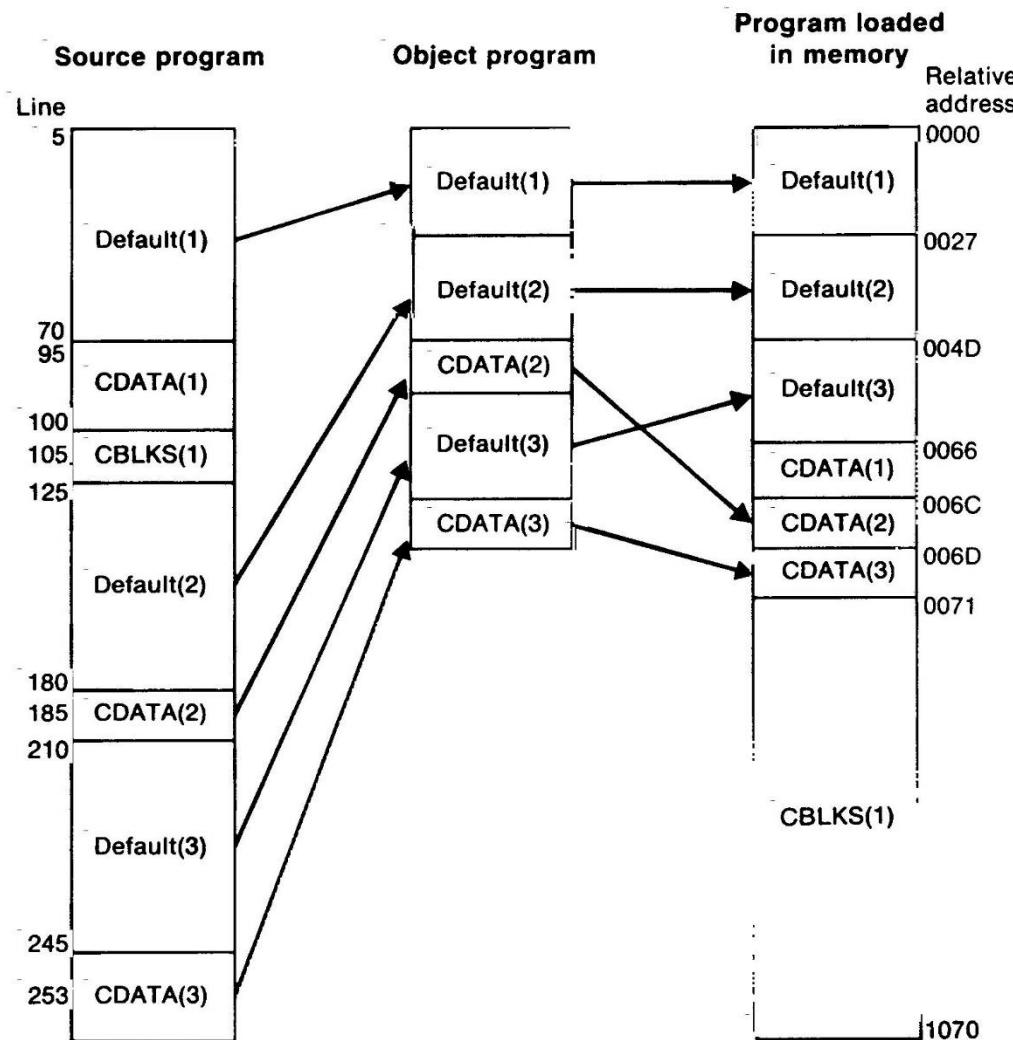
## Object Program for Fig 2.11

**It is not necessary to physically rearrange the generated code in the object program to place the pieces of each program block together.**

The assembler simply writes the object code as it is generated during pass 2 and insert the proper load address in each Text record => The loader will rearrange them.

Line 5-70 Line 125-180 Line 185 Line 210-245 LTORG constants	<pre> HCOPY 000000001071 T0000001E1720634B20210320602900003320064B203B3F2FEE0320550F2056010003 T00001E090F20484B20293E203F T0000271DB410B400B44075101000E32038332FFADB2032A00433200857A02FB850 T00004093B2FEA13201F4F0000 T00006C01F1 T00004D19B410772017E3201B332FFA53A016DF2012B8503B2FEE4F0000 T00006D04454F4605 E000000 </pre>
--	--

# Assembly and Loading Process



# Control Sections

A part of the program that can be assembled, loaded and relocated independently of the others.

Different control sections are most often used for subroutines or other logical subdivisions of a program.

Syntax : *Section Name*    CSECT

The assembler establishes a separate location counter (beginning at 0) for each control section.

The instructions in one control section may need to refer to instructions or data located in another section.

**EXTDEF** (External Definition): for external symbols that are defined in this control section and may be used by other sections.

**EXTREF** (External Reference): for symbols that are used in this control section and are defined elsewhere.

Control section names do not need to be named in an EXTDEF statement.

Same symbols can be used in different control sections (line 107 and line 190).



# Example of Control Sections (Fig 2.15)

Line	Source statement			
5	COPY	START	0	COPY FILE FROM INPUT TO OUTPUT
6		EXTDEF	BUFFER, BUFEND, LENGTH	
7		EXTREF	RDREC, WRREC	
10	FIRST	STL	RETADR	SAVE RETURN ADDRESS
15	CLOOP	+JSUB	RDREC	READ INPUT RECORD
20		LDA	LENGTH	TEST FOR EOF (LENGTH = 0)
25		COMP	#0	
30		JEQ	ENDFIL	EXIT IF EOF FOUND
35		+JSUB	WRREC	WRITE OUTPUT RECORD
40		J	CLOOP	LOOP
45	ENDFIL	LDA	=C'EOF'	INSERT END OF FILE MARKER
50		STA	BUFFER	
55		LDA	#3	SET LENGTH = 3
60		STA	LENGTH	
65		+JSUB	WRREC	WRITE EOF
70		J	@RETADR	RETURN TO CALLER
95	RETADR	RESW	1	
100	LENGTH	RESW	1	LENGTH OF RECORD
103		LTORG		
105	BUFFER	RESB	4096	4096-BYTE BUFFER AREA
106	BUFEND	EQU	*	
107	MAXLEN	EQU	BUFEND-BUFFER	
109	RDREC	CSECT		
110	.			
115	.			SUBROUTINE TO READ RECORD INTO BUFFER
120	.			
122		EXTREF	BUFFER, LENGTH, BUFEND	
125		CLEAR	X	CLEAR LOOP COUNTER
130		CLEAR	A	CLEAR A TO ZERO
132		CLEAR	S	CLEAR S TO ZERO
133		LDT	MAXLEN	
135	RLOOP	TD	INPUT	TEST INPUT DEVICE
140		JEQ	RLOOP	LOOP UNTIL READY
145		RD	INPUT	READ CHARACTER INTO REGISTER A
150		COMPR	A, S	TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT	EXIT LOOP IF EOR
160		+STCH	BUFFER, X	STORE CHARACTER IN BUFFER
165		TIXR	T	LOOP UNLESS MAX LENGTH
170		JLT	RLOOP	HAS BEEN REACHED
175	EXIT	+STX	LENGTH	SAVE RECORD LENGTH
180	INPUT	RSUB	X'F1'	RETURN TO CALLER
185	MAXLEN	BYTE		CODE FOR INPUT DEVICE
190		WORD	BUFEND-BUFFER	
193	WRREC	CSECT		
195	.			
200	.			SUBROUTINE TO WRITE RECORD FROM BUFFER
205	.			
207		EXTREF	LENGTH, BUFFER	
210		CLEAR	X	CLEAR LOOP COUNTER
212		+LDT	LENGTH	
215	WLOOP	TD	=X'05'	TEST OUTPUT DEVICE
220		JEQ	WLOOP	LOOP UNTIL READY
225		+LDCH	BUFFER, X	GET CHARACTER FROM BUFFER
230		WD	=X'05'	WRITE CHARACTER
235		TIXR	T	LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP	HAVE BEEN WRITTEN
245		RSUB		RETURN TO CALLER
255		END	FIRST	



# Fig 2.15 with Object Code (Fig 2.16)

Line	Loc	Source statement			Object code
5	0000	COPY	START	0	
6			EXTDEF	BUFFER, BUFEND, LENGTH	
7			EXTREF	RDREC, WRREC	
10	0000	FIRST	STL	RETADR	172027
15	0003	CLOOP	+JSUB	RDREC	4B1000000
20	0007		LDA	LENGTH	032023
25	000A		COMP	#0	290000
30	000D		JEQ	ENDFIL	332007
35	0010		+JSUB	WRREC	4B1000000
40	0014		J	CLOOP	3F2FEC
45	0017	ENDFIL	LDA	=C'EOF'	032016
50	001A		STA	BUFFER	0F2016
55	001D		LDA	#3	010003
60	0020		STA	LENGTH	0F200A
65	0023		+JSUB	WRREC	4B1000000
70	0027		J	@RETADR	3E2000
95	002A	RETADR	RESW	1	
100	002D	LENGTH	RESW	1	
103			LTORG		
105	0030	*	=C'EOF'		454F46
106	0033	BUFFER	RESB	4096	
107	1033	BUFEND	EQU	*	
	1000	MAXLEN	EQU	BUFEND-BUFFER	
109	0000	RDREC	CSECT		
110			SUBROUTINE TO READ RECORD INTO BUFFER		
115			EXTREF	BUFFER, LENGTH, BUFEND	
120			CLEAR	X	B410
122			CLEAR	A	B400
125	0000		CLEAR	S	B440
130	0002		LDT	MAXLEN	77201F
132	0004		TD	INPUT	E3201B
133	0006		JEQ	RLOOP	332FFA
135	0009	RLOOP	RD	INPUT	DB2015
140	000C		COMPR	A,S	A004
145	000F		JEQ	EXIT	332009
150	0012		+STCH	BUFFER, X	579000000
155	0014		TIXR	T	B850
160	0017		JLT	RLOOP	3B2FE9
165	001B		RSUB	LENGTH	131000000
170	001D		BYTE	X'F1'	4F0000
175	0020	EXIT	WORD	BUFEND-BUFFER	F1
180	0024				000000
185	0027	INPUT			
190	0028	MAXLEN			
193	0000	WRREC	CSECT		
195			SUBROUTINE TO WRITE RECORD FROM BUFFER		
200			EXTREF	LENGTH, BUFFER	
205			CLEAR	X	B410
207			+LDT	LENGTH	771000000
210	0000		TD	=X'05'	E32012
212	0002		JEQ	WLOOP	332FFA
215	0006	WLOOP	+LDCH	BUFFER, X	539000000
220	0009		WD	=X'05'	DF2008
225	000C		TIXR	T	B850
230	0010		JLT	WLOOP	3B2FEE
235	0013		RSUB	END	4F0000
240	0015			FIRST	
245	0018			=X'05'	05
255	001B				



# Handling External References by the Assembler

Since the instructions in one control section may need to refer to instructions or data located in another section, the assembler generates information so that the linker can perform the required linking.

(e.g. 1)

15 0003 CLOOP +JSUB RDREC 4B1 <u>00000</u>
--

RDREC is named in the EXTREF statement.

Since the assembler has no idea where the control section containing RDREC will be loaded, it generates an extended format instruction with an address of zero.

(e.g. 2)

160 0017 +STCH BUFFER, X 579 <u>00000</u>
---

(e.g. 3) – Check the difference between line 107 and line 190.

190 0028 MAXLEN WORD BUFEND-BUFFER <u>000000</u>
--



# Define Record and Refer Record

## Define Record:

Col. 1	D
Col. 2-7	Name of external symbol defined in this control section.
Col. 8-13	Relative address of symbol within this control section (hex).
Col. 14-73	Repeat information in Col. 2-13 for other external symbols.

## Refer Record:

Col. 1	R
Col. 2-7	Name of external symbol referred to in this control section.
Col. 8-73	Name of other external reference symbols.

## Modification Record (revised):

Col. 1	M
Col. 2-7	Starting address of the field to be modified, relative to the beginning of the <u>control section</u> (hex).
Col. 8-9	Length of the field to be modified, in half bytes (hex).
Col. 10	Modification flag (+ or -).
Col. 11-16	External symbol whose value is to be added to or subtracted from the indicated field.



# Object Program for Fig 2.15

```

HCOPY 000000001033
DBUFFER000033BUFEND001033LENGTH00002D
RRDREC WRREC
T0000001D1720274B100000320232900003320074B1000003F2FEC0320160F2016
T00001D0D0100030F200A4B1000003E2000
T00003003454F46
M00000405+RDREC
M00001105+WRREC
M00002405+WRREC
E000000

HRDREC 000000000028
RBUFFERLENGTHBUFEND
T0000001DB410B400B44077201FE3201B332FFADB2015A00433200957900000B850
T00001D0E3B2FE913100004F0000F1000000
M00001805+BUFFER
M00002105+LENGTH
M00002806+BUFEND
M00002806-BUFFER
E

HWRREC 00000000001C
RLENGTHBUFFER
T0000001CB41077100000E32012332FFA53900000DF2008B8503B2FEE4F000C05
M00000305+LENGTH
M00000D05+BUFFER
E

```

