Data Structure and Advanced Programming Homework 3

第一題

(a) Catalan Number with Recursive method

```
int catalanRecursive(int n)
{
    if (n <= 0)
        return 0;
    else if (n == 1)
        return 1;
    else
        return 2 * (2 * n - 1) * catalanRecursive(n - 1) / (n + 1);
}</pre>
```

(b) Knapsack Problem

```
input: w[ 1, ..., n], v[ 1, ..., n], V, B

define w = sum(w[i] * x[i]) where x[i] is in {0, 1}

define P(n, B) is the solution to the problem with n items, each with w[i] and v[i],

for w = 0 to B

P(0, w) <- 0

for i = 0 to n

for w = 0 to B

if w[i] > w

P(i, w) <- P(i - 1, w)

else

P(i, w) = max { P(i - 1, w), v[i] + P(i - 1, w - w[i])}

if P(n, B) > V

return true

else
return false
```

(c) Pseudocode for Augmenting Path

```
input: number of nodes n, edges m, an adjacency matrix adj[n][n]
     define Q \leftarrow vector, s \leftarrow s[0], t \leftarrow s[n]
     // below follows the functions in c++ library <vector>
     // in the queue/vector are Paths practiced in class
     Q.push_back(s)
     while(!Q.empty() && terminationNotFound)
          curPath <- Q.front()</pre>
          curNode <- Q.front().getTail()</pre>
          Q.erase(Q.begin)
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12
          isLoop <- false</pre>
13
          vector<Path> newPaths <- getNextWithResCap(curNode)</pre>
          for i <- 1 to curPath.size()</pre>
              for j <- 1 to curPath.size()</pre>
                   if curPath.node[i] == curPath.node[j]
                       isLoop <- true</pre>
          if isLoop
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              continue;
          else
              for i <- 0 to newPaths.size()</pre>
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24
                   if(newPaths[i].getTail() == t)
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                       return newPaths[i]
26
                   else
                       Q.push_back(newPaths[i])
28
     return -1
```

(d) Revised Maximum Flow Problem

(i) Multiple Sources and Terminations

新增兩個假想的起點 s' 與終點 t' 來連接多個起點 $\{s\}_{i=1}^3$ 與終點 $\{t\}_{i=1}^2$,並且定義 s' 與 $\{s\}_{i=1}^3$ 之間及 t' 與 $\{t\}_{i=1}^2$ 之間的邊的容量都是無限大,就可以應用原本的程式了。

(ii) Vertices with Capacities

把每個點 v_i 都分成兩個點 v_{in} 跟 v_{out} ,並且以有向邊從 v_{in} 指向 v_{out} ,定義這條邊的容量是原先指定點的容量,就可以應用原程式。