CSCI 2270 - CS 2: Data Structures





University of Colorado Boulder

Topics

Algorithm Analysis (Big O)

- Choosing the right algorithm can make a huge difference!
 - Government databases
 - Search engines
 - Air Traffic Control Games
 - Medical equipment
 - Computer/video games
- We can also time algorithms, but...

- We can count the number of operations that the algorithm performs.
 - Count the # of instructions it performs
- The # of instructions performed may vary based on:
 - Size of the input
 - The organization of the input
- The # of instructions can be written as a cost function on the input size.

Counting example

```
void printArray(int arr[], int size){
   for (int i = 0; i < size; i++){
      cout << arr[i] << endl;
   }
}</pre>
```

Operations performed on an array of length 10 = 32

- Let's not choose a particular input size (e.g. 10), but rather we will express a cost function for input of size n.
- Assume that the running time, t, of an algorithm is proportional to the number of operations.
- Express t as a function of n
 - Where t is the time required to process the data using some algorithm A
 - Denote a cost function as tA(n)
 - i.e. the running time of algorithm A, with input size n

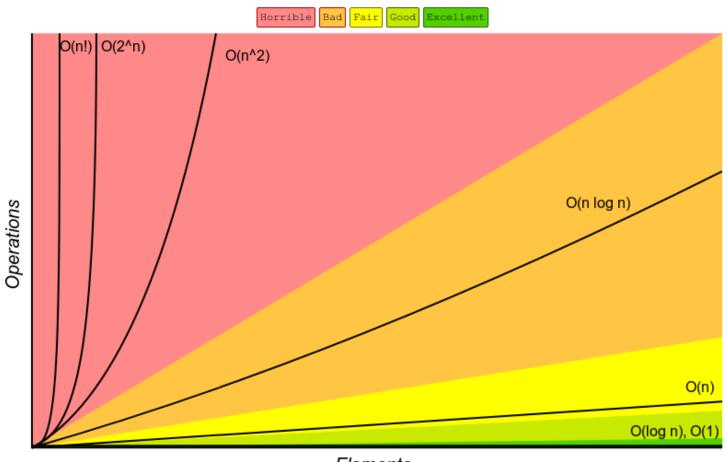
Go back to our example

```
void printArray(int arr[], int size){
   for (int i = 0; i < size; i++){
      cout << arr[i] << endl;
   }
}</pre>
```

Declare and initialize i: 1 Comparison, print array element, increment i: 3n Make comparison when i = n: 1

Consequently, our result is: t = 3n+2

- Big O Notation is the way of measuring the efficiency of an algorithm and how well it scales based on the size of the dataset.
- O notation approximates a cost function that allows us to estimate growth rate.
 - function f(n) is the algorithm's growth rate function
- Big O notation does not give a precise formulation of the cost function for a particular data size.



- O(1)/Constant Complexity
- O(n)/Linear Complexity
- O(log n)/Logarithmic Complexity
- O(n log n)
- O(n²)/Quadratic Complexity

- Arrays
 - Accessing an array index: O(1)
 - Insert item to beginning or middle: O(n)
 - Time taken will be proportional to the size of the list
 - Insert/Add item to end: O(1)
 - Assumes the array has space
 - Adding in middle: O(n)
 - Deletion from beginning or middle: O(n)
 - Shifting required after removing an element
 - Deleting from the end: O(1)

- Linked Lists
 - Insert node to beginning: O(1)
 - Same for deletion
 - Insert in middle: O(n) (linear time)
 - Must iterate over n elems to get to correct location
 - Same for deletion
 - Insert node to end: O(n)
 - Must iterate over n elements to get to end
 - Search: O(n)
 - Slower search times than arrays because we can't use indexes. Must traverse LL!
 - Using tail pointer, adding or removing from end will be O(1)

Arrays vs Linked Lists

- Arrays have faster search time
- Arrays utilize less memory per element
 - More memory needed per node in LL because additional storage needed for pointer
- Linked Lists have faster insertion/deletion time
- Linked Lists allow for dynamic size
 - Shrink and grow
- Linked Lists are more efficient memory-wise
 - Allocation & utilization

Questions

