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This lecture

- C-like C++
- Memory management in C/C++
 - Memory allocation: static, automatic and dynamic, Ch. 6.1.1,
 Ch. 12
 - Allocation and de-allocation, Ch. 12.1.2
 - 2-D and N-D Arrays, Ch. 3.5
 - Pass by value, by reference, by pointer, Ch. 6.2-6.2.4
 - Passing a function Ch. 10.3
 - Function pointers, Ch. 6.7



Java Memory Management

Java

- Memory is requested by the program
 - all Java objects are on the heap allocated with new
 - Local variables of primitive types and references are on the stack
- Memory is freed by the Java garbage collector
 - After the last reference is deleted
 - garbage collector must keep track (e.g., reference counting or mark and sweep (delete))
 - When the garbage collector executes



Java Garbage Collection

Mark and sweep algorithm

- Two pass algorithm
 - Mark all objects which are reachable
 - Sweep (delete) objects which are unmarked
- Heap is compacted
- Sun Java JDK 1.0 and 1.1

Generational collector

- Young ("Eden" and 2 survivor), tenured and permanent spaces
- Copying between spaces, i.e., compaction
- Garbage collections: Minor (young space only) and major
- Improves performance dramatically when most objects are short-lived, e.g., temporary objects
- Sun Java JDK 1.2 and later



Advantages and Disadvantages of Garbage Collection (GC)

Advantages

- Frees program designer from thinking about GC
 - increased productivity
- Avoids most memory bugs and leaks which can be extremely difficult to find and fix
- Security

Disadvantages

- Overhead
 - JVM needs to run a GC thread
 - JVM needs extra memory to store extra object information
- less control, soft real-time applications
- often leads to memory issues being overlooked (references not being released)



User Control of Garbage Collection in Java

Java has minimal user control

- Parameters to the Java VM
- Nulling of object references
- Pooling of objects
- Explicitly request a GC
- But still need to make sure that references are not kept too long (or forever)

Results

hard to predict and often counter-productive



Memory Management and Garbage Collection in C++

- Different mechanisms depending on type of memory
- Memory on the heap: No garbage collection!

Changes in C++11

- garbage collection interface is defined for compiler implementers
- optional vendor-specific implementations
- Uses notion of safe derivation of pointers but users remains in control through declaring, proposed options: relaxed (same as before), preferred, strict



Memory Allocation

Static Memory Allocation

 Allocated by the "linker" at the beginning of the program, and de-allocated when the program finishes

Automatic Memory Allocation

- Automatically allocated and de-allocated during the program execution
 - Examples include function arguments, function return values and local variables

Dynamic Memory Allocation

- Handled by explicit statements in the program. Allocation and de-allocation only by request
- No Garbage Collection



Memory Locations

- In the executable program code
 - allocated when program is loaded
 - global variables, static variables
- On the stack of the program
 - allocated automatically during execution
 - local variables, functions parameters, functions return values
- On the heap of the program
 - allocated as coded during execution
 - in C++: new new[] delete delete[]
 - in C: malloc calloc realloc free



new and delete Operator/Keyword

new type and new type []

- Allocates memory for objects, arrays, data types on the heap from the free store
- Returns nonzero pointer to a memory location on success
- May raise an exception std::bad alloc

delete pointer and delete [] pointer

- De-allocates a block of memory pointed to by a pointer and returns the memory to the free store
- Results unpredictable if used on not properly allocated memory



new and delete Example

- new and delete
 - used for object allocation and de-allocation
 - new is similar to Java
 - Error to call delete on variables not allocated with new
- Example: array allocation and de-allocation

```
const int max = 1000;
int numbers[max];
int *dynNumbers;
int arraySize = 0; cin >> arraySize;
if (!cin.fail()) {
   dynNumbers = new int[arraySize];
   delete[] dynNumbers;
}
delete[] numbers; | Illegal!
```

A closer Look at new and delete

3 steps during new

- allocate memory large enough to hold data of the requested type
 constructor in Java and C++
- construct the type
- return a pointer to the constructed type

2 steps during delete

- destruct the type
- return the memory to the free-store

finalize method in Java; destructor in C++



C-Style Memory (De-)Allocation

- No constructors and destructors in C
 - they are not called during free and malloc
- malloc return a pointer of type void *
 - memory is not initialized
 - Not type-safe
- malloc does not raise an exception on failure
 - Program needs to check for null pointer
- free behaves similar to delete
 - memory pointed to must have been allocated with malloc
 - multiple free calls will have unpredictable results
 - free with a null pointer should be Ok (ISO C)



Same example: in C

```
char *char1. *char40:
//Allocation of one char
char1 = (char*) malloc(1);
//Allocation of a string of 40 char
char40 = (char^*) malloc(40);
double *double1, *double50;
//Allocation of one double
double1 = (double*) malloc(sizeof(double));
double50 = (double*) malloc(50 * sizeof(double));
struct Node{
              int value:
              Node * left:
              Node * right;
};
Node *node1, *nodeN;
//Allocation of one node
node1 = (Node*) malloc(sizeof(Node));
//Allocation of N nodes
int N;
printf("How many nodes : ");
scanf ("%d ", &N);
nodeN = (Node*) malloc(N * sizeof(Node));
```

and in C++

```
char *char1, *char40;
//Allocation of one char
char1 = new char;
//Allocation of a string of 40 char
char40 = new char[40];
double *double1, *double50;
//Allocation of one double
double1 = new double:
double50 = new double(50);
struct Node{
              int value:
              Node * left:
              Node * right;
Node *node1, *nodeN;
//Allocation of one node
node1 = new Node;
//Allocation of N nodes
int N;
cout << "How many nodes : ";</pre>
cin >> N;
nodeN = new Node(N);
```

Array of Arrays: Definition and Initialization

- array holds garbage unless initialized
- inner braces are optional
- partial initialization is possible
- first dimension is optional
- "no" limit on the number of dimensions

Array of Arrays

- C/C++ does not have multidimensional arrays!
- BUT arrays of arrays
 - Example:
 - array of size 3 which holds arrays of size 4

a ₀₀	a ₀₁	a ₀₂	a ₀₃
a ₁₀	a ₁₁	a ₁₂	a ₁₃
a ₂₀	a ₂₁	a ₂₂	a ₂₃

a ₀	a ₀₀	a ₀₁	a ₀₂	a ₀₃
a ₁	a ₁₀	a ₁₁	a ₁₂	a ₁₃
a ₂	a ₂₀	a ₂₁	a ₂₂	a ₂₃

Element Access and Pointers

- Access by indices
- Pointer manipulation
 - Pointer to row (array)
 - A row is a onedimensional array, i.e., pointer to an array
 - Pointer to element
 - Pointer to first element

```
int numbers3D[3][4][5];
int element = numbers3D[1][2][0];
int numbers[3][4];
```

```
int (*row)[4] = &numbers[1];
int elementA = (*row)[3];
int* elementB = &numbers[0][2];
int* elementC = &numbers[0][0];
int* elementD = numbers[0];
```

Arrays cannot be assigned; they are not automatically copied



Semi-Dynamic Allocation of Array of Arrays

- Only the first dimension can be determined at run-time
 - Example: 2D array
 - number of rows at run-time
 - number of cols at compile

```
int numRows = 3;
int (*numbers)[4] = new int[numRows][4];
delete[] numbers;
```

Dynamic Allocation of Array of Arrays

Need to (de-)allocate all the arrays (arrays of arrays)

```
int **numbers = new int*[3];
for(int i=0; i<3; i++) {
  numbers[i] = new int[4];
}

for(int i=0; i<3; i++) {
  delete[] numbers[i];
}
delete[] numbers;</pre>
```

- row + 1 call to new
 - row + 1 separate memory location may be returned
- Contiguous memory layout may be important!



Memory Layout of Array of Arrays

Pointer manipulation (C-style)

```
int *tmp = new int[12];
int **numbers = new int*[3];
for(int r=0; r<numRows; r++) {
  numbers[r] = &tmp[r*numCols];
}</pre>
```

- Deallocation? Ugly!
- Instead:
 - Write a matrix, image etc. class with operators and accessors and use a one-dimensional array.
 - Use std::array (a class which wraps low-level arrays)
 - Use std::vector (similar to java.util.ArrayList)



Initialization by Value and by Pointer

- Variable is copied
 - Pointer is just another type

```
int *ptrNumA = \frac{1}{0 \times 0100}
ptrNumA = ptrNumB = 0 \times 0100
```

```
int numA = 1;
int numB = numA;
```

```
int *ptrNumA = &numA;
int *ptrNumB = ptrNumA;
```

- Pass by pointer is identical to pass by value
 - pointer is copied
 - object pointed to is not affected
 - similar effect than pass by reference as in Java
- Note: Arrays are passed by pointer even if we use array syntax!



Variable Initialization by Reference

- Reference to the variable is created
 - Same as in Java
 - Reference is an alias to a variable
 - "just another name"

```
int numA = 1;
int &numB = numA; // int * const numB = &numA;

numB = 3 + numB; //(*numB) = 3 + (*numB);
numB++; // (*numB)++;
```

 In general: Variable initialization is the same as argument passing into methods and functions: "Pass by Reference"

```
void myFunction( int (&arg)[4] );
...
int numbers[4];
myFunction( numbers );
```

Passing Arrays

 Low-level arrays can be passed by pointer to the first element, by pointer or by reference

```
int a[]{0,1,2,3,4,5,6,7};
int sumOfP1( int *ptr);
int sumOfP2( int ptr[]);
int sumOfP3( int ptr[8]);
```

Pass by pointer to first element! Size unknown!

```
for (auto a: ptr) {}
```

```
int sumOfPArray( int (*ptr)[8]);
```

```
int sumOfRefArray( int (&ref)[8]);
```

Pointer (top) or reference (bottom) to an int array of size 8.

```
for ( auto a : ref ) {}
```

Passing Arrays

 Low-level arrays can be passed by pointer to the first element, by pointer or by reference

```
int a[]{0,1,2,3,4,5,6,7};
int sumOfP1( int *ptr);
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Pass by pointer to first element! Size unknown!

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```
int sumOfPArray( int (*ptr)[8]);
```

```
int sumOfRefArray( int (&ref)[8]);
```

Pointer (top) or reference (bottom) to an int array of size 8.

```
for ( auto a : ref ) {}
```



Function Pointers

- STL, GUIs, etc. expect to pass a callback function.
 - We can use objects with operator overloading (functors).
 - But in simple cases a function is enough
 - Example: lessThan

Using Function Pointers

- Function pointer types have to match
 - arguments and return type (unlike function overloading)
- Function pointers have awkward syntax
 - We can use simplified notations, or, we can use typedefs
 - Example: lessThan (continued)

Calling Function through Pointers

- explicitly dereferenced
- implicitly dereferenced

```
optional
```

```
// Function declaration
bool lessThan(const Point2D&, const Point2D&);
// no typedef
Point2D ptA, ptB;
lessThan(ptA, ptB); // direct call of function, no ptr
bool (*ptr) (const Point2D&, const Point2D&) = &lessThan;
(*ptr)(ptA,ptB);
ptr(ptA,ptB);
// using a typedef C++11 notation
using pt compare=bool (*) (const Point2D&, const Point2D&);
pt compare c = &lessThan;
(*c) (ptA, ptB);
c(ptA,ptB);
```

References with auto

- auto type deduction will use the base type
 - reference is not part of base type
 - aside: top-level const are also not part of the base type

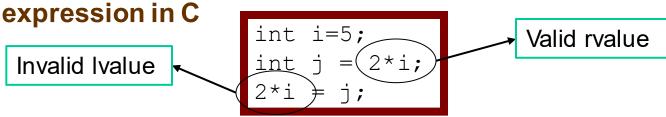
```
int i, &iRef = i, *iPtr = &i;
auto j = iRef; // j is not a reference
auto k = iPtr; // k is a pointer
auto &jRef = j; // jRef is a reference
```

- But there is also decltype which works differently
 - Need to discuss LValue and RValue



LValue and Rvalue Expression

Origin of term comes simply from the place of an



- A non-rigorous C++ definition
 - LValue expressions refer to a memory location
 - yields an object or function
 - uses the memory location
 - RValue are the rest (non-LValue expressions).

Some Operators require LValues

- Assignment operators need LValue as left operand (the classic)
- Address-of-operator require a LValue operand
- Dereference and subscripting yields an LValue (you can assign to it).
- Increment and decrement need a LValue operand and the prefix version yields a LValue

```
int foo();
int i=5;
int *j = &i;
int array[5];
array[3] = 3;
++i *= 3; // i = (5+1)* 3 = 18
```

References with decltype

<u>decitype</u> deduces the type of the specified expression.

auto deduces types based on values being assigned to the variable.

decltype deduction will evaluate expressions

- brackets can be used to define references
- using an expression that yields a LValue (can appear on the lhs of an assignment) will make it a reference
- top-level const and reference are used to deduce type

```
int i = 5, &iRef = i, *iPtr = &i;//iRef=5, *iPtr=5
decltype(iRef) jRef = i; // reference jRef = 5
decltype(i) iVal=4; // iVal=4 is not a reference
decltype((i)) kRef = i; // reference kRef = 5
decltype(*iPtr) lRef = i; // reference lRef = 5
```



Auto with const and References C++14

- As we have seen auto uses the underlying type, e.g.,
 const int or int& become int. This is the same as normal initialization rules (templates).
- auto and decltype infer types differently
- In C++14 we can use decltype type inference with auto

```
int i = 5, &iRef = i; //iRef = 5
auto j = iRef; // j is not a reference = 5
auto &jRef = i; // jRef is a reference = 5
decltype(auto) jRef2 = iRef; // jRef2 = 5 is also a reference
decltype(auto) jRef3 = (i); // jRef3 = 5 is also a reference
```

Next Lectures

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- Object-oriented design
 - Class relationships: association, aggregation, generalization and inheritance
 - Pointer attributes
 - Copy construction and assignment
 - Polymorphism: Virtual functions, abstract classes and dynamic cast
 - Exceptions Basics
 - Inline functions, static members, constexpr

