
MODULE *ORSet*

EXTENDS *Naturals, Sequences, Bags, TLC, SEC*

CONSTANTS

Data the set of data

Instance $\triangleq [d : \text{Data}, r : \text{Replica}, k : \text{Nat}]$

VARIABLES

sSet, *sSet*[*r*]: set of active *Instance*(*s*) maintained by *r* \in *Replica*

seq, *seq*[*r*]: local sequence number at replica *r* \in *Replica*

incoming, *incoming*[*r*]: incoming messages at replica *r* \in *Replica*

msg,

messageSet

vars $\triangleq \langle sSet, seq, incoming, msg, messageSet, SECvars \rangle$

Msg $\triangleq [r : \text{Replica}, S : \text{SUBSET } Instance, seq : \text{Nat}, update : \text{SUBSET } Update]$

Network $\triangleq \text{INSTANCE } ReliableNetwork$

TypeOK \triangleq

$\wedge sSet \in [Replica \rightarrow \text{SUBSET } Instance]$

$\wedge seq \in [Replica \rightarrow \text{Nat}]$

Init \triangleq

$\wedge sSet = [r \in Replica \mapsto \{\}]$

$\wedge seq = [r \in Replica \mapsto 0]$

$\wedge Network!RInit$

$\wedge SECInit$

Add(*d*, *r*) \triangleq

$\wedge seq' = [seq \text{ EXCEPT } ![r] = @ + 1]$

$\wedge sSet' = [sSet \text{ EXCEPT } ![r] = @ \cup \{[d \mapsto d, r \mapsto r, k \mapsto seq'[r]]\}]$

$\wedge SECUpdate(r, seq[r])$

$\wedge \text{UNCHANGED } \langle incoming, msg, messageSet \rangle$

Remove(*d*, *r*) \triangleq

$\wedge \text{LET } D \triangleq \{ins \in sSet[r] : ins.d = d\}$

$\text{IN } sSet' = [sSet \text{ EXCEPT } ![r] = @ \setminus D]$

$\wedge seq' = [seq \text{ EXCEPT } ![r] = @ + 1]$

$\wedge SECUpdate(r, seq[r])$

$\wedge \text{UNCHANGED } \langle incoming, msg, messageSet \rangle$

Broadcast(*s*, *m*) \triangleq

$[r \in Replica \mapsto \text{IF } s = r \text{ THEN } incoming[s]$

$\text{ELSE } incoming[r] \oplus SetToBag(\{m\})]$

$$\begin{aligned}
Send(r) &\triangleq \\
&\wedge Network!RBroadcast(r, [r \mapsto r, S \mapsto sSet[r], seq \mapsto seq[r], \\
&\quad update \mapsto OpUpdate(r)]) \\
&\wedge SECSend(r) \\
&\wedge UNCHANGED \langle sSet, seq \rangle
\end{aligned}$$

$$\begin{aligned}
Receive(r) &\triangleq \\
&\wedge Network!RDeliver(r) \\
&\wedge SECDeliver(r, msg'[r]) \\
&\wedge sSet' = [sSet \text{ EXCEPT } ![r] = @ \cup msg'[r].S] \\
&\wedge UNCHANGED \langle seq \rangle
\end{aligned}$$

$$\begin{aligned}
Next &\triangleq \\
&\vee \exists r \in Replica : \exists a \in Data : \\
&\quad Add(a, r) \vee Remove(a, r) \\
&\vee \exists r \in Replica : \\
&\quad Send(r) \vee Receive(r)
\end{aligned}$$

$$Spec \triangleq Init \wedge \Box [Next]_{vars}$$

$$Read(r) \triangleq \{ins.d : ins \in sSet[r]\}$$

QC: Quiescent Consistency

$$Quiescence \triangleq \forall r \in Replica : incoming[r] = \langle \rangle$$

$$Convergence \triangleq \forall r, s \in Replica : Read(r) = Read(s)$$

$$QC \triangleq Quiescence \Rightarrow Convergence$$

SEC: Strong Eventual Consistency

$$SEC \triangleq \forall r1, r2 \in Replica : SameUpdate(r1, r2) \Rightarrow Read(r1) = Read(r2)$$

\ * Modification History
\ * Last modified Thu May 16 10:34:44 CST 2019 by zfwang
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