

## Problem Set 8

Advanced Logic

30th October 2022

Due date: Friday, 4 November.

1. (50%) Prove the *Substitution Lemma for Terms*: for any signature  $\Sigma$ , any terms  $s$  and  $t$  of  $\Sigma$ , any variable  $v$ , any structure  $S$  for  $\Sigma$ , and any assignment  $g$  for  $S$ ,

$$\llbracket t[s/v] \rrbracket_S^g = \llbracket t \rrbracket_S^{g[v \mapsto \llbracket s \rrbracket_S^g]}$$

*Hint*: prove this by induction on the construction of terms (should it be  $t$  or  $s$ ?)

*Reminder*:

$$g[v \mapsto \llbracket s \rrbracket_S^g](u) = \begin{cases} \llbracket s \rrbracket_S^g & \text{if } u = v \\ gu & \text{otherwise} \end{cases}$$

2. (30%) Using the result from part 1, prove the *Substitution Lemma for Formulae*: for any signature  $\Sigma$ , any term  $s$  of  $\Sigma$ , any formula  $P$  of  $\Sigma$ , any variable  $v$ , any structure  $S$  for  $\Sigma$ , and any assignment  $g$  for  $S$ ,

$$S, g \models P[s/v] \text{ iff } S, g[v \mapsto \llbracket s \rrbracket_S^g] \models P$$

*Hint*: prove this by induction on the construction of formulae. In the induction step for quantifiers, you will need to separately consider the case of formulae that begin with  $\forall v$  (or  $\exists v$ ) and formulae that begin with  $\forall u$  (or  $\exists u$ ) for some other variable  $u$ . You may if you wish rely on the “Irrelevance Lemma” according to which if  $g(u) = h(u)$  for all  $u \in FV(Q)$ ,  $S, g \models Q$  iff  $S, h \models Q$ .

3. (10%) Using the result from part 2, prove the steps in the proof of the Soundness Theorem corresponding to the  $\forall$ Elim and  $\exists$ Intro rules. That is: show that if  $\Gamma \models \forall v P$  then  $\Gamma \models P[s/v]$  for every term  $s$ , and that if  $\Gamma \models P[s/v]$ ,  $\Gamma \models \exists v P$ .
4. (5%) Using the result from part 2, prove the step in the proof of the Soundness Theorem corresponding to the  $=$ Elim rule. That is: show that if  $\Gamma \models P[s/v]$  and  $\Gamma \models s = t$ , then  $\Gamma \models P[t/v]$ .
5. (5%) Using the result from part 2, prove the step in the proof of the Soundness Theorem corresponding to the  $\exists$ Elim rule. That is: show that if  $\Gamma \models \exists v P$  and  $\Gamma, P[u/v] \models Q$ , then  $\Gamma \models Q$ , provided that  $u$  is not free in  $\Gamma$ ,  $Q$ , or  $\exists v P$ .

Give examples to show that this can fail when (i)  $u$  is free in  $\Gamma$  though not in  $Q$  or  $\exists v P$ ; (ii)  $u$  is free in  $Q$  though not in  $\Gamma$  or  $\exists v P$ ; (iii)  $u$  is free in  $\exists v P$  though not in  $\Gamma$  or  $Q$ .

EXTRA CREDIT (2.5% each, up to a maximum of 10%) Prove the remaining facts about logical consequence required to complete the proof of the Soundness Theorem, namely:

1. If  $\Gamma \models P$  and  $\Gamma \models Q$  then  $\Gamma \models P \wedge Q$  ( $\wedge$ Intro)
2. If  $\Gamma \models P \wedge Q$  then  $\Gamma \models P$  and  $\Gamma \models Q$  ( $\wedge$ Elim1 and  $\wedge$ Elim2).
3. If  $\Gamma, P \models Q$  then  $\Gamma \models P \rightarrow Q$  ( $\rightarrow$ Intro).
4. If  $\Gamma \models P \rightarrow Q$  and  $\Gamma \models P$  then  $\Gamma \models Q$  ( $\rightarrow$ Elim).
5. If  $\Gamma, P \models Q$  and  $\Gamma, P \models \neg Q$  then  $\Gamma \models \neg P$  ( $\neg$ Intro).
6. If  $\Gamma \models \neg\neg P$  then  $\Gamma \models P$  (DNE).
7.  $\models t = t$  for every term  $t$  ( $=$ Intro).