

A Cut Into Classical Planning

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ProSeminar KI

Outline

Introduction

- Definition

- Representation

Finding a plan

- SATPlan

- Graphplan

 - General

 - Forward Expansion (FE)

 - Backward Search (BS)

Conclusion

The Human wants the machine (agent) to imitate her.

Examples:

- ▶ Humans make coffee by following a step-by-step plan: Get a cup, press the big START button on the coffee machine, add some sugar, ...
- ▶ Humans travel on holiday by following a step-by-step plan: Buy a train ticket to Berlin *by July 17*, get on the train to Berlin *at 8 am*, check in at the hotel *by 12 am*, ...

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State-Transition System

State-Transition System

$$\Sigma=(S,A,E,\gamma)$$

- ▶ S is a set of states
- ▶ A is a set of actions
- ▶ E is a set of Events
- ▶ γ is a state-transition function

Restrictions

Classical planning deals with planning problems under certain restrictions:

Restrictions

- ▶ A1: S is a finite set of states
- ▶ A2: Σ is fully observable
- ▶ A3: Σ is deterministic
- ▶ A4: The planner handles only restricted goals

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- ▶ A5: Σ is static
- ▶ A6: Sequential plans
- ▶ A7: Implicit time
- ▶ A8: Offline planning

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Planning Problem

Planning Problem

$$P = (\Sigma, s_0, g)$$

- ▶ Σ is the planning domain
- ▶ s_0 is the initial state
- ▶ g is the goal state or set of goal states

Problem: Is there a sequence of actions that will lead from the initial state to the goal state?

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PDDL

PDDL

- ▶ Planning Domain Description Language
- ▶ Describes the parts of the state-transition system
- ▶ Provides a generalized description for planning Problems
- ▶ Main input language at International Planning Competitions

States

states

- ▶ Conjunction of fluents
- ▶ Fluents: ground, functionless atoms
- ▶ Unique names imply distinct objects

Example 1

$$at(r_1, l_1) \wedge loaded(r_1, c_1)$$

Robot r_1 is at location l_1 and loaded with cargo c_1

States

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Action schemas

Action Schemas

- ▶ List of object names
- ▶ Preconditions
- ▶ Effects

Action move

Action($move(r_1, l_1, l_2)$)

Precondition : $at(r_1, l_1) \wedge robot(r_1) \wedge location(l_1) \wedge location(l_2)$

Effect : $\neg at(r_1, l_1) \wedge at(r_1, l_2)$

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Example

Example: Cargo-transport

$Init(at(r_1, l_1) \wedge at(c_1, l_1) \wedge robot(r_1) \wedge cargo(c_1) \wedge location(l_1) \wedge location(l_2))$
 $Goal(at(c_1, l_2))$

Example

Action(load(r, c, l)

Precon : $at(r, l) \wedge at(c, l) \wedge robot(r) \wedge cargo(c) \wedge loaction(l)$

Effect : $\neg at(c, l) \wedge loaded(r, c)$

Action(unload(r, c, l)

Precon : $at(r, l) \wedge loaded(r, c) \wedge robot(r) \wedge cargo(c) \wedge location(l)$

Effect : $\neg loaded(r, c) \wedge at(c, l)$

Action(move(r, from, to)

Precon : $at(r, from) \wedge robot(r) \wedge location(from) \wedge location(to)$

Effect : $\neg at(r, from) \wedge at(r, to)$

Example

- ▶ Actions affect predicates

Solution-Plan

load(r_1, c_1, l_1)

move(r_1, l_1, l_2)

unload(r_1, c_1, l_2)

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SAT-Problem

SAT-Problem

Given a formula φ .

Does there exist a model that satisfies φ ?

steps to follow

- ▶ Translate the classical planning problem into a satisfiability problem
- ▶ Determine if there exists a plan by solving the satisfiability problem with a satisfiability decision procedure.
- ▶ Extract the plan from the assignments determined by the satisfiability decision procedure

states

- ▶ still ground, functionless atoms
- ▶ additionally negations are allowed

example

$$at(r_1, l_1) \wedge loaded(r_1, c_1)$$

solutions with second location

$$\mu_1 = (at(r_1, l_1) \leftarrow true, loaded(r_1, c_1) \leftarrow true, at(r_1, l_2) \leftarrow false)$$

$$\mu_2 = (at(r_1, l_1) \leftarrow true, loaded(r_1, c_1) \leftarrow true, at(r_1, l_2) \leftarrow true)$$

state-transitions

- ▶ Add a timestep i to every predicate
- ▶ An action itself is a predicate now
- ▶ The action implies preconditions and effects

Example

$$\text{move}(r_1, l_1, l_2, s_1) \Rightarrow \\ (at(r_1, l_1, s_1) \wedge \neg at(r_1, l_2, s_1) \wedge \neg at(r_1, l_1, s_2) \wedge at(r_1, l_2, s_2))$$

Formulas

Formulas

- ▶ Initial state: $\bigwedge_{f \in s_0} f_0 \wedge \bigwedge_{f \notin s_0} \neg f_0$
- ▶ Goal state: $\bigwedge_{f \in g^+} f_n \wedge \bigwedge_{f \in g^-} \neg f_n$
- ▶ Actions: $a_i \Rightarrow \left(\bigwedge_{p \in \text{precond}(a)} p_i \wedge \bigwedge_{e \in \text{effects}(a)} e_{i+1} \right)$

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Formulas

Formulas

- Explanatory frame axioms:

$$\left(\neg f_i \wedge f_{i+1} \Rightarrow \left(\bigvee_{a \in A \mid f_i \in \text{effects}^+(a)} a_i \right) \right) \wedge$$
$$\left(f_i \wedge \neg f_{i+1} \Rightarrow \left(\bigvee_{a \in A \mid f_i \in \text{effects}^-(a)} a_i \right) \right)$$

- Complete exclusion axioms: $\neg a_i \vee \neg b_i$

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- Complete exclusion axioms: $\neg a_i \vee \neg b_i$

Example

Example: Cargo-transport

Init-Goal

- ▶ (*init*) $at(r_1, l_1, 0) \wedge at(c_1, l_1, 0) \wedge \neg at(r_1, l_2, 0) \wedge \dots$
- ▶ (*goal*) $at(c_1, l_2, 3) \wedge \neg at(r_1, l_1, 2) \wedge \dots$

Actions

- ▶ *move1* $move(r_1, l_1, l_2, 0) \Rightarrow$
 $(at(r_1, l_1, 0) \wedge at(r_1, l_2, 1) \wedge \neg at(r_1, l_1, 1))$
- ▶ *move2* $move(r_1, l_2, l_1, 0) \Rightarrow$
 $(at(r_1, l_2, 0) \wedge at(r_1, l_1, 1) \wedge \neg at(r_1, l_2, 1))$
- ▶ ...

Example

Example: Cargo-transport

Explanatory frame axioms

- ▶ $\neg at(r_1, l_1, 0) \wedge at(r_1, l_1, 1) \Rightarrow move(r_1, l_2, l_1, 0)$
- ▶ $\neg at(r_1, l_2, 0) \wedge at(r_1, l_2, 1) \Rightarrow move(r_1, l_1, l_2, 0)$
- ▶ $at(r_1, l_1, 0) \wedge \neg at(r_1, l_1, 1) \Rightarrow move(r_1, l_1, l_2, 0)$
- ▶ $at(r_1, l_2, 0) \wedge \neg at(r_1, l_2, 1) \Rightarrow move(r_1, l_2, l_1, 0)$
- ▶ ...

Example

Example: Cargo-transport

Complete exclusion axioms

- ▶ $\neg \text{move}(r_1, l_1, l_2, 0) \vee \neg \text{move}(r_1, l_2, l_1, 0)$
- ▶ ...

Solution

- ▶ $\text{load}(r_1, c_1, l_1, 0), \text{move}(r_1, l_1, l_2, 1), \text{unload}(r_1, c_1, l_2, 2)$

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The Planning Graph is Graphplan's main data structure.

- ▶ Layered graph in which one layer corresponds to one time step in the plan
- ▶ Layer i consists of one set of actions that are applicable at layer i and one set of literals that *could* be true at layer i .
- ▶ For every positive and negative literal p , we add the persistence action α_p with precondition p and effect p .

Forward Graph Expansion vs. Backward Search

Mutex Links

Literals

Actions

Goals

Backtracking

Forward Graph Expansion vs. Backward Search

Forward Graph Expansion

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The planning graph is expanded up to the point that it might be possible to achieve the goal.

Notation:

Literals: $at(r, l_1) \longrightarrow Arl_1,$

$at(c_1, l_2) \longrightarrow Ac_1l_2,$

$loaded(r, c_1) \longrightarrow Lrc_1,$

etc...

Actions: $move(r, l_1, l_2) \longrightarrow Mrl_1l_2,$

$unload(r, c_1, l_1) \longrightarrow Urc_1l_1,$

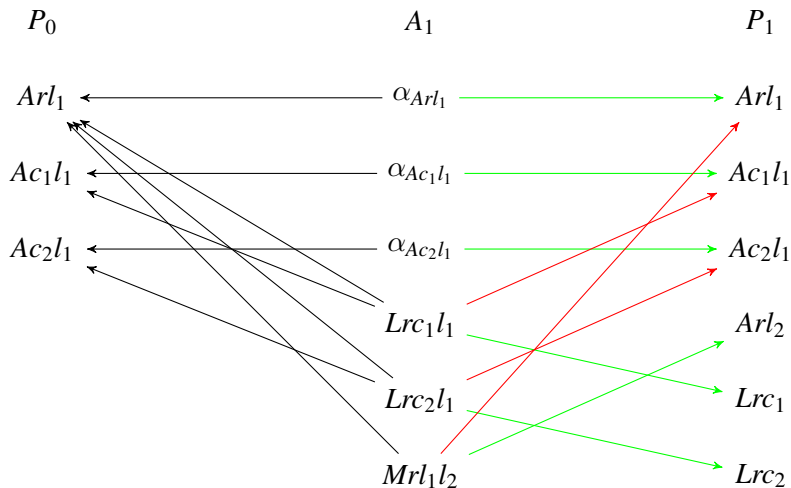
etc...

Cargo Transport (CT) problem:

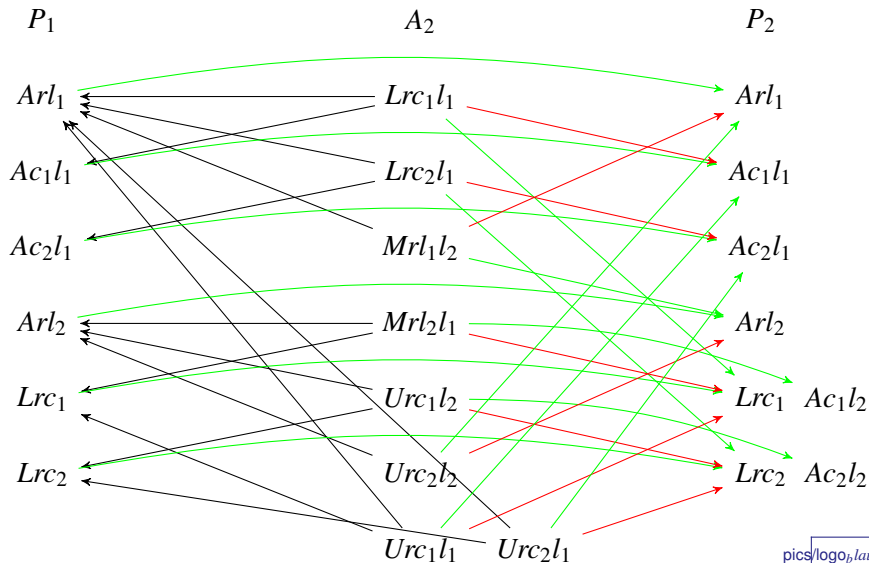
$Init(Arl_1, Ac_1l_1, Ac_2l_1)$

$Goal(Ac_1l_2, Ac_2l_2)$

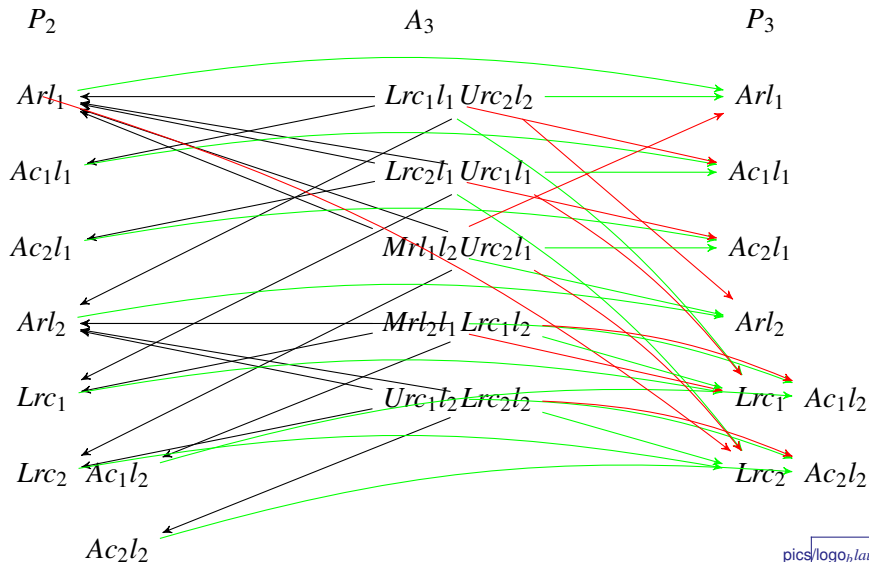
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Mutex Links indicate whether two literals or actions are incompatible with each other.

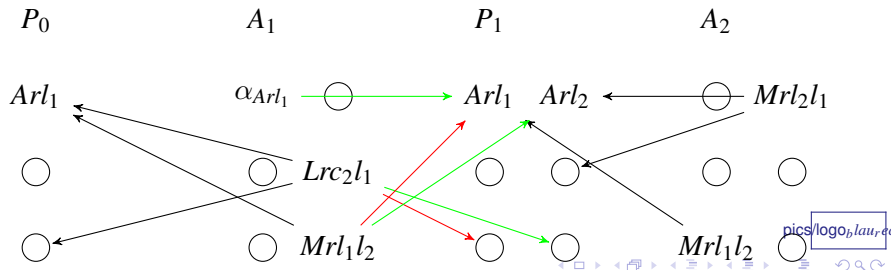
Mutex Links for Actions:

Two actions a_1 and a_2 are mutex in P_i if and only if (iff) one of the following conditions holds:

- ▶ *Interference*: action a_1 deletes a positive effect or a precondition of action a_2 or vice versa (*Dependence*)
- ▶ *Competing Needs*: precondition of a_1 is mutex with a precondition with a_2 or vice versa

Notation: $(a_1, a_2) \in \mu A_i$

Example (not all mutex links shown):



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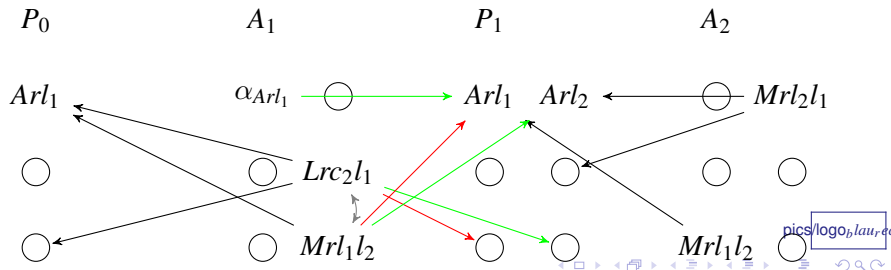
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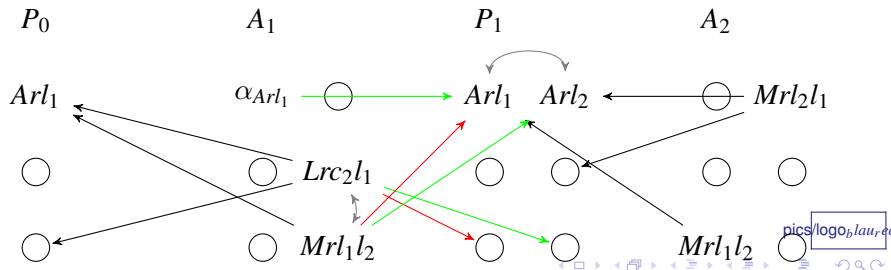
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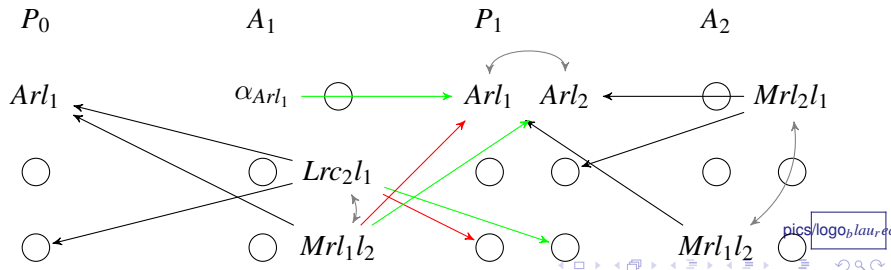
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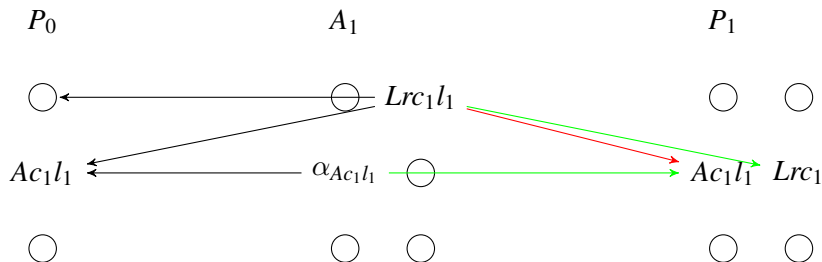
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Mutex Links for Literals:

Two literals p_1 and p_2 are mutex in P_i iff no nonmutex pair of actions (a_1, a_2) exists such that a_1 produces p_1 and a_2 produces p_2 .

Notation: $(p_1, p_2) \in \mu P_i$. We set $\mu P_0 = \emptyset$.

Example:



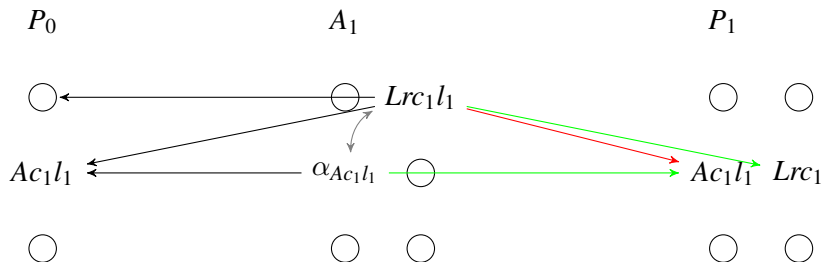
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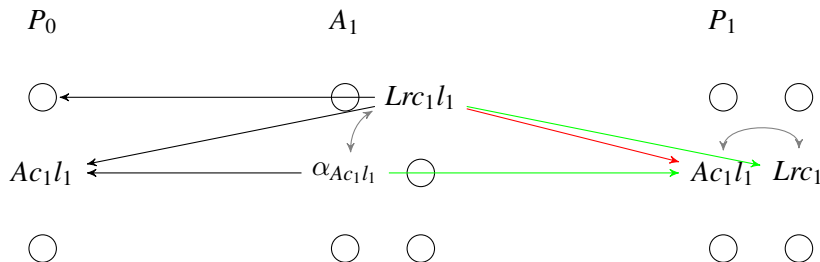
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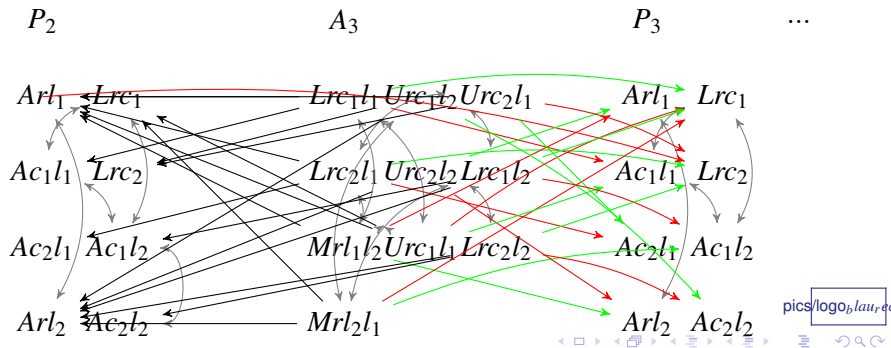


When a planning graph does not grow anymore, it has reached its fixed point.

Given any planning graph G . Then there is a smallest k such that all of the following conditions hold:

- ▶ *Fixed Proposition Layer*: number of propositions at layer k is equal to number of propositions at any layer $i > k$
- ▶ *Fixed Prop. Mutex Links*: number of mutual links at layer k is equal to number of mutual links at any layer $i > k$

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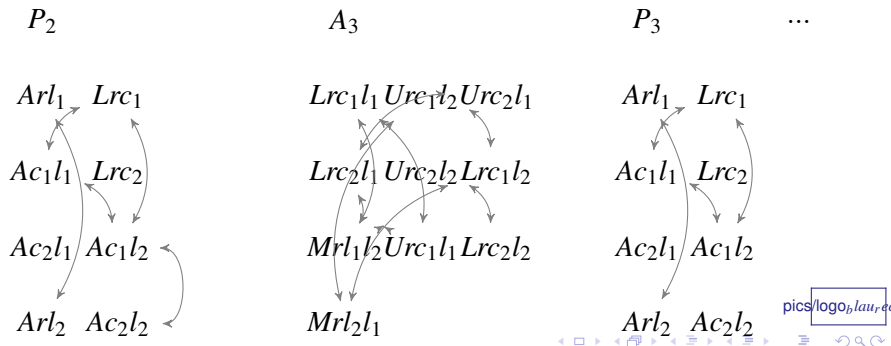


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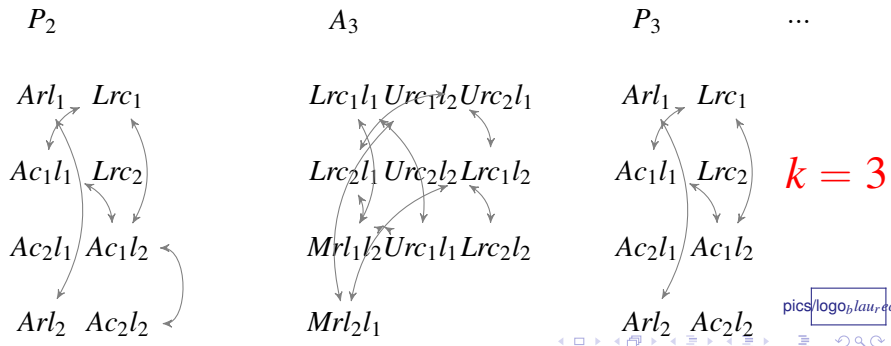


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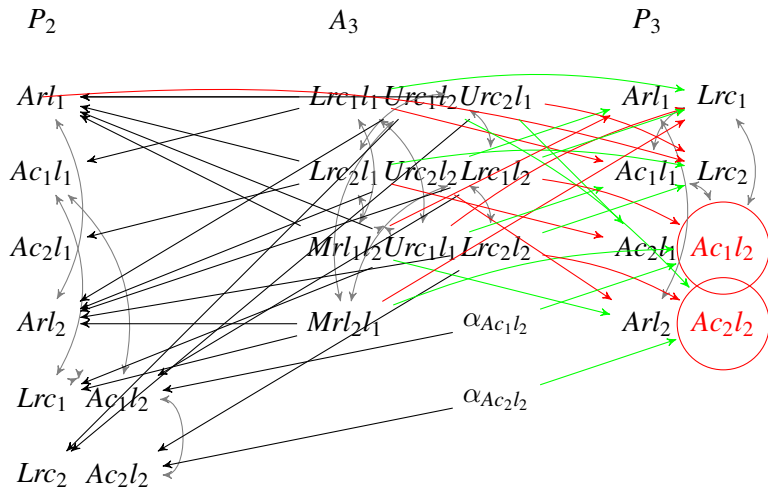
General

Forward Expansion (FE)

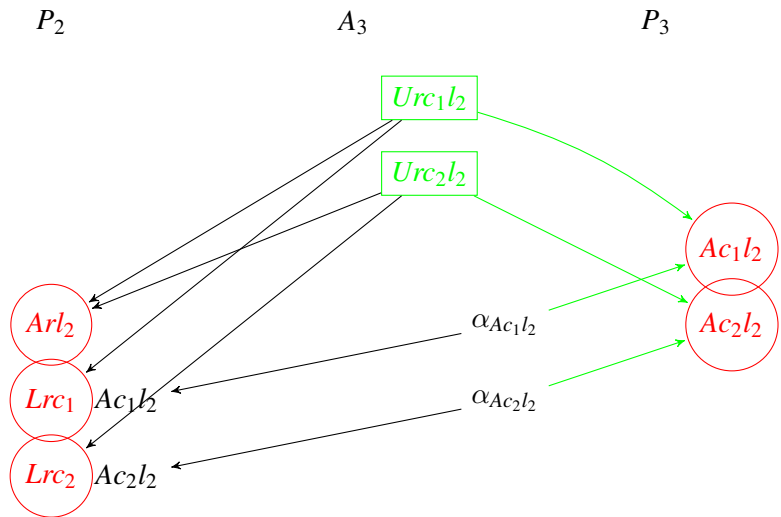
Backward Search (BS)

Conclusion

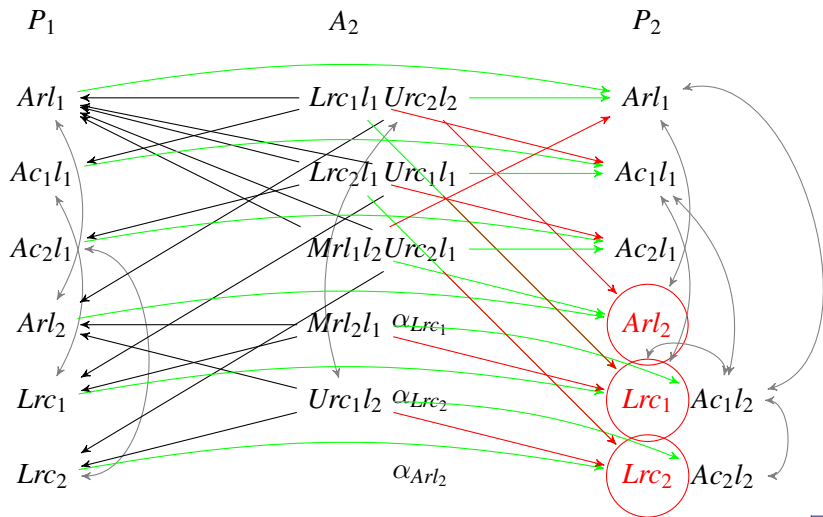
Graphplan recursively tracks the goal set down to the initial state.



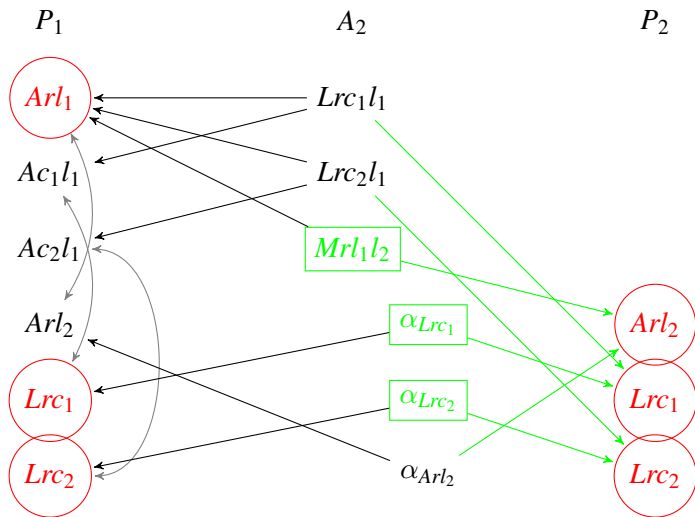
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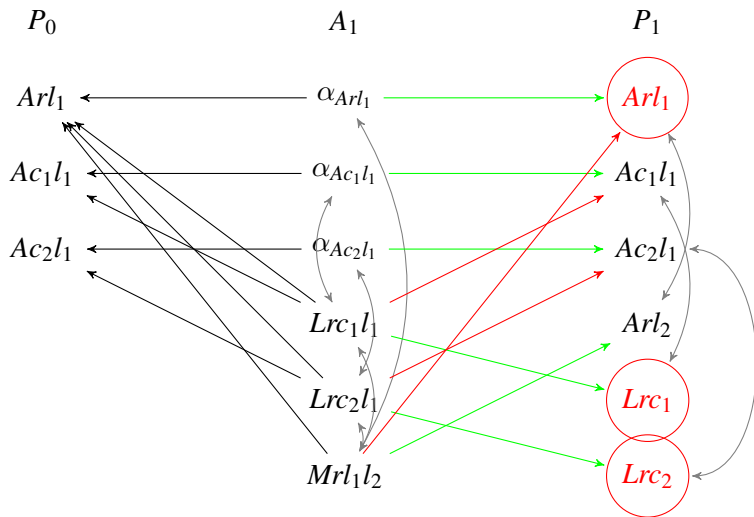
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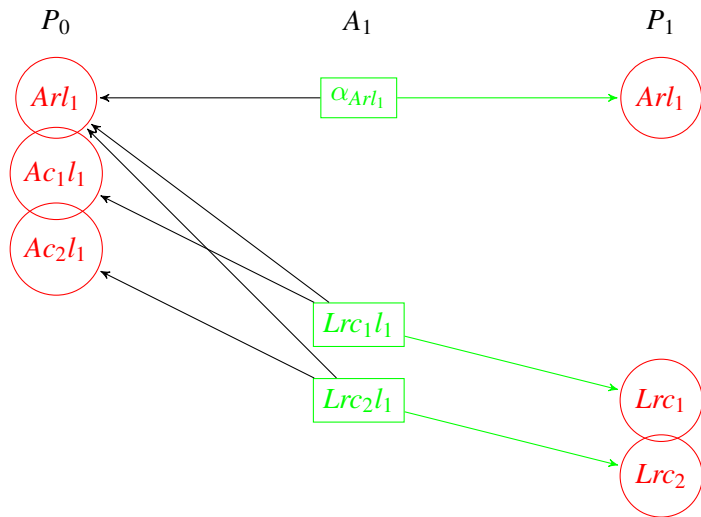
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Layered Solution Plan:

- ▶ contains sets of nonmutex actions that correspond to a specific layer in the graph
- ▶ has a strict order of elements, i.e. each set of nonmutex actions has to be performed sequentially (but the order within one set is arbitrary)
- ▶ the layered solution plan to the example problem instance:

$$\Pi = (\pi_1, \pi_2, \pi_3) = (\{Lrc_1l_1, Lrc_2l_1\}, \{Mrl_1l_2\}, \{Urc_1l_2, Urc_2l_2\}).$$

Graphplan and SATPlan have assets and drawbacks.

	Asset	Drawback
Graphplan	<ul style="list-style-type: none">▶ data structure	<ul style="list-style-type: none">▶ no good heuristics (e.g. for choosing producers in BS)
SATPlan	<ul style="list-style-type: none">▶ permanently new SAT solvers coming up, hot research	<ul style="list-style-type: none">▶ number of clauses might be unfeasible to ground

CP systems are deployed in critical situations.

- ▶ Satellite launch
- ▶ Hubble Space Telescope
- ▶ etc...