Design and Analysis of Algorithms

L11: Algo Design Ideas

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Resources

- Martin Richards
 - Notes on Data Structures and Algorithms
- Text book 2: Horowitz
- Text book I: Levitin
- https://visualgo.net/en

Ideas on Algorithm Design

- Recognize a variant on known problem
- Reduce to a simpler problem
- Divide and Conquer
- Estimation of cost by recurrence relation
- Dynamic Programming
- Greedy Algorithm
- Back tracking
- Hill Climbing
- Identify wasted work in a simple method
- Find a mathematical lower bound
- Million Monkey method

Variant on a Known Problem

- Try to identify a related problem whose solution is known
- Use solution of one problem to solve essentially tricky part of other problem.
- Solve the problem in totality
- Example 1: multiply two numbers
 - Known variant: add two numbers
 - How to use addition of two numbers to achieve multiplication
- Example 2: Class assignment submission
 - variant: Identify classmate who has done similar
 - Use the solution with different input params

Reduction to a Simpler Problem

- Typically works where induction methods are used to show correctness proof
- Recursive functions take this appproach
- Example: Compute n²

$$n^2 = (n-1)^2 + ??$$

= $(n-1)^2 + 2n - 1$

• Example: factorial n

$$n! = n*(n-1)!$$

Example: GCD computation

$$gcd(y,x) = gcd(x,r)$$
 where $r=y%x$

Divide and Conquer Algo

- Divide (break) the problem (size n) into similar sub problems
 - Size of sub problems should be some factor of original e.g. n/c
 - When small enough, solve by brute force
- Conquer (Solve) the sub-problem
 - Use recursion to solve small problem
- Combine (Merge) the solution of sub-parts
- The cost is
 - cost of breaking
 - cost of solving subproblem
 - cost of combining

Cost Estimation by Recurrence

- Use recurrence relation to find the cost of operation
 - Once the cost is known, then algorithm framework is worked out
 - Design the algorithm on how to break the given problem
- Example: Mergesort

$$f(n) = 2f(n/2) + n$$

-> $f(n) = O(nlog_2n)$

• Example: Word search in English dictionary

$$f(n) = f(n/k) + c$$

-> $f(n) = log_k n$

Dynamic Programming

- Build a table of solutions to smaller versions of the problem
- Work towards the solution by using this table
- Example: Binomial Coefficient

• Example: Computing n²

Greedy Algorithms

- Useful when some sort of optimizations is involved.
- Basic idea:
 - Perform whatever operation contributes most towards the final goal
 - Next step will be same approach after the previous step
 - Note: Final solution may not always be optimal
- Example:
 - Eating in restaurant: get whatever available now
 - No waiting is to be done
 - Taking exam (you know all answers) but time is less
 - Which questions to start writing

Backtracking

- When algorithm involves search (explore)
 - backtracking is useful
- Split the search procedure in multiple paths(parts)
 - Start with one path
 - If solution not found, and path ends,
 - Backtrack to previous starting point and
 - search on next path
- Example:
 - Finding a solution to the maze (childhood comicbooks)
 - Chessboard: queens placement
 - Chessboard: Knights traversal

Hill Climbing

- Generally, used in optimization problems
- Start with some feasible (non-optimal) solution
- Incrementally work towards improving the solution
- May not always find the optimal solution
- Analogy: Climbing to hilltop in hilly region
- Example:
 - Getting a job(placement)
 - Start with some placement (job)
 - Work towards a better placement (better company)

Identify Wasted Work

- Design a simple solution to the problem
- Analyze the solution and
 - identify critical costly part of the solution
- Attack the weakness of critical part to improve the solution
- Example (real life):
 - Phone converstation, and you are put on hold
 - can do some useful work while on call hold
- Example (CS): CPU Scheduling
 - when a processing is waiting for I/O
 - other process is assigned CPU

Seek Mathematical Lower Bound

- Establish a proof that some task must take at least certain minimum time
- Use this insight to design the algorithm
- A properly proved lower bound can prevent wasted time seeking improvement
- Example: Sorting n numbers:
 - min time: ceil(log2 n!) comparisons
 - Consider n=4, min time = $log_2(24) = 5$
 - Consider n=5, min time = $log_2(120) = 7$
 - Design the algorithm

Million Monkey (MM) Method

- Problem: Give one typewriter each to million monkeys (random character press) and a million years, and one of the monkey will write the Shakespear play.
 - Idea: give a problem to a group of researchers and sufficient time, they will be able to find the solution of the problem.
- Example: expected output 'hello'.
 - Assume typewriter has 50 keys
 - Probability of not typing 'hello' = $1-(1/50)^5$
 - Probability of not typing 'hello' by million monkeys in 10000 tries is
 - $-(1-(1/50)^5)^{10} = 1.2*10^{-14}$
 - Probability of typing 'hello' = $1-1.2*10^{-14} \approx 1$

Exercises-A

- Given 4 numbers: a, b, c, d
 - Sort them using log₂ 4!

```
=log_2 24
```

=5 comparisons

Sort 4 numbers

```
abcd
              cabd
 abdc badc cadb
 acbd <del>bead</del> <del>cbad</del> <del>dbac</del>
 acdb <del>bcda cbda dbca</del>
adbe bdae cdab deab
adch bdca cdba dcba
Consider two numbers a, b
 if a < b
    12 combinations goes out
 else
   other 12 combinations goes out
 if c < d
   6 combinations goes out
```

Exercise-B

- Given 5 numbers: a, b, c, d, e
 - Sort them using log₂ 5!

```
=log_2 120
```

=7 comparisons

Summary

- Ideas for algorithm design
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