Design and Analysis of Algorithms

L15: Strassen Multiplication

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Resources

- Text book 2: Sec 3.8 Horowitz
- Text book 1: Sec 4.5 Levitin

Multiplication Example

• Multiply 2 numbers (of 2 digits) e.g. 98, 76

A= 98 =
$$10*9+8$$

B= 76 = $10*7+6$
C = A*B = 10^2 (9*7) +10 (8*7+9*6) +8*6
=6300+10 (56+54) +48 = 7448

- Multiplications: 4, additions: 3
- Alternatively

$$C = 10^{2}(9*7) + 8*6$$

+ $10((9+8)*(7+6) - 9*7 - 8*6))$
= $10^{2}*9*7 + 10*(8*7 + 9*6) + 8*6 = 7448$

- Multiplications: 3, Additions: 6

Multiplication Example

Multiply 2 numbers (of 2 digits)

$$-A = a_2a_1, B=b_2b_1$$

 $-C = A * B$
 $= 100c_3 + 10c_2 + c_1$

where

$$c_3=a_2b_2$$
, $c_2=(a_2b_1+a_1b_2)$, $c_1=a_1b_1$

- Multiplications: 4 Additions: 3
- Alternatively

$$c_3=a_2b_2$$
,
 $c_1=a_1b_1$
 $c_2=(a_2+a_1)*(b_2+b_1)-c_3-c_1$

- Multiplications: 3 Additions: 6

Multiplication Example

• Multiply 2 numbers (of 4 digits) e.g. 9876, 5432

A=
$$98*10^2+76$$
, B= $54*10^2+32$
C=A*B= 10^4 ($98*54$)
+ 10^2 ($98*32$ + $76*54$)
+ $76*32$

- Multiplications: 4, additions: 3
- Alternatively

$$C = 10^4 (98*54 + 76*32) + 10^2 ((98+76)*(54+32)-98*54-76*32))$$

- Multiplications: 3, additions: 6

Large Numbers Multiplication

- Problem:
 Given two large numbers with N digits, Multiply these in efficient way
- Solution: traditional way (high school mathematics)

- Efficiency: Multiplications: n2, Additions: 2n
- Complexity analysis: $\Theta(n^2)$

Large Numbers Multiplication

- Given two large numbers A, B with n digits,
 - multiply these in efficient way
- Let $A = A_1A_2$, and $B=B_1B_2$,
 - where A and B are n-digit numbers, and
- A_1 , A_2 , B_1 , and B_2 are (n/2)-digit numbers, then

$$A*B = A_1*B_1*10^n + (A_1*B_2+A_2*B_1) 10^{n/2} + A_2*B_2$$

- Efficiency : Multiplications: 4 each of n/2
- Recurrence relation

```
T(n) = 4T(n/2) + O(n) = 2^2T(n/2) + n
= 2^2[4T(n/4) + O(n/2)] + n
= 2^4T(n/2^2) + 2^2(n/2) + n = 2^4T(n/2^2) + 2n + n
= \Theta(n^2)
```

Large Numbers Multiplication

• Let $A=A_1A_2$, and $B=B_1B_2$, are n-digit numbers, and

-
$$A_1$$
, A_2 , B_1 , and B_2 are (n/2)-digit numbers, then

$$A*B = 10^n*A_1*B_1 + A_2*B_2 + 10^{n/2} ((A_1+A_2)*(B_1+B_2) - A_1*B_1 - A_2*B_2)$$

Recurrence relation using 3 multiplications

```
T (n) = 3T (n/2) + O (n) = 3<sup>1</sup>T (n/2) + cn

= 3<sup>1</sup> [3T (n/4) + cn/2] + cn

= 3<sup>2</sup>T (n/2<sup>2</sup>) + c [3<sup>1</sup>n/2<sup>1</sup>+n]

= ...

= 3<sup>k</sup>T (n/2<sup>k</sup>) + c [3<sup>k-1</sup>n/2<sup>k-1</sup>+...+n]

= \Theta (3<sup>1</sup>og<sub>2</sub>n) = \Theta (n<sup>1</sup>og<sub>2</sub>3)

\approx \Theta (n<sup>1</sup>·585)
```

Matrix Multiplication

Conventional matrix multiplication

$$a_{11} \ a_{12} \dots a_{1n}$$
 $a_{21} \ a_{22} \dots a_{2n}$
 $\vdots \qquad \vdots$
 $a_{n1} \ a_{n2} \dots a_{nn}$

$$a_{11} \ a_{12} \ ... \ a_{1n}$$
 | $b_{11} \ b_{12} \ ... \ b_{1n}$ | $c_{11} \ c_{12} \ ... \ c_{1n}$ | $c_{21} \ c_{22} \ ... \ c_{2n}$ | $c_{21} \ c_{22} \ ... \ c_{2n}$ | $c_{2n} \ c_{2n} \ c_{2n}$ | $c_{2n} \ c_{2n} \ c_{2n}$ | $c_{2n} \ c_$

where the element $C_{\dot{1}\dot{7}}$ is computed as

$$c_{ij} = a_{i1}b_{1j} + a_{i2}b_{2j} + a_{i3}b_{3j} + ... + a_{in}b_{nj}$$

- computations required for Cii
 - Multiplications: n
 - additions: n
- Total computations required for matrix multiplication:
 - 2n³ i.e.
 - Θ (n³)

Matrix Multiplication

Conventional matrix multiplication

$$\begin{bmatrix} C_1 & C_2 \\ C_3 & C_4 \end{bmatrix} = \begin{bmatrix} A_1 & A_2 \\ A_3 & A_4 \end{bmatrix} * \begin{bmatrix} B_1 & B_2 \\ B_3 & B_4 \end{bmatrix}$$

$$= \begin{bmatrix} A_1B_1 + A_2B_3 & A_1B_2 + A_2B_4 \\ A_3B_1 + A_4B_3 & A_3B_2 + A_4B_4 \end{bmatrix}$$

- Multiplications: 8 (=23), Additions: 4
- Recurrence relation:

$$T(n) = 8T(n/2) + 4(n/2) = 2^3T(n/2) + 2^2(n/2)$$

= $\Theta(n^3)$

• Regular matrix multiplication takes n3 computations

 Strassen discovered that product of two matrices can be computed as follows

$$\begin{bmatrix} C_{11} & C_{12} \\ C_{21} & C_{21} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} * \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$
$$= \begin{bmatrix} M_1 + M_4 - M_5 + M_7 & M_3 + M_5 \\ M_2 + M_4 & M_1 + M_3 - M_2 + M_6 \end{bmatrix}$$

- $M_1 = (A_{11} + A_{22}) * (B_{11} + B_{22})$
- $M_2 = (A_{21} + A_{22}) * B_{11}$
- $M_3 = A_{11} * (B_{12} B_{22})$
- $M_4 = A_{22} * (B_{21} B_{11})$
- $M_5 = (A_{11} + A_{12}) * B_{22}$
- $M_6 = (A_{21} A_{11}) * (B_{11} + B_{12})$
- $M_7 = (A_{12} A_{22}) * (B_{21} + B_{22})$

- $M_1 = (A_{11} + A_{22}) * (B_{11} + B_{22})$
- $M_2 = (A_{21} + A_{22}) * B_{11}$
- $M_3 = A_{11} * (B_{12} B_{22})$
- $M_4 = A_{22} * (B_{21} B_{11})$
- $M_5 = (A_{11} + A_{12}) * B_{22}$
- $M_7 = (A_{12} A_{22}) * (B_{21} + B_{22})$

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} * \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

•
$$M_4 = A_{22} + (B_{21} - B_{11})$$

• $M_5 = (A_{11} + A_{12}) * B_{22}$
• $M_6 = (A_{21} - A_{11}) * (B_{11} + B_{12})$

$$M_1 + M_4 - M_5 + M_7$$

$$M_2 + M_4$$

$$M_1 + M_3 - M_2 + M_6$$

$$C_{11}=M_{1}+M_{4}-M_{5}+M_{7}=(A_{11}+A_{22})*(B_{11}+B_{22})+A_{22}*(B_{21}-B_{11})$$

$$-(A_{11}+A_{12})*B_{22}+(A_{12}-A_{22})*(B_{21}+B_{22})$$

$$=A_{11}B_{11}+A_{22}B_{11}+A_{11}B_{22}+A_{22}B_{22}+A_{22}B_{21}$$

$$-A_{22}B_{11}-A_{11}B_{22}-A_{12}B_{22}+A_{12}B_{21}$$

$$-A_{22}B_{21}+A_{12}B_{22}-A_{22}B_{22}$$

$$=A_{11}B_{11}+A_{12}B_{21}$$

- $M_1 = (A_{11} + A_{22}) * (B_{11} + B_{22})$
- $M_2 = (A_{21} + A_{22}) * B_{11}$
- $M_3 = A_{11} * (B_{12} B_{22})$
- $M_4 = A_{22} * (B_{21} B_{11})$
- $M_5 = (A_{11} + A_{12}) * B_{22}$
- $M_6 = (A_{21} A_{11}) * (B_{11} + B_{12})$
- $M_7 = (A_{12} A_{22}) * (B_{21} + B_{22})$

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} * \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

$$M_1+M_4-M_5+M_7$$
 M_3+M_5 M_2+M_4 $M_1+M_3-M_2+M_6$

$$C_{12}=M_3+M_5=A_{11}*(B_{12}-B_{22})+(A_{11}+A_{12})*B_{22}$$

$$=A_{11}B_{12}-A_{11}B_{22}+A_{11}B_{22}+A_{12}B_{22}$$

$$=A_{11}B_{12}+A_{12}B_{22}$$

$$C_{21}=M_2+M_4=(A_{21}+A_{22})*B_{11}+A_{22}*(B_{21}-B_{11})$$

$$=A_{21}B_{11}+A_{22}B_{11}+A_{22}B_{21}-A_{22}B_{11}$$

$$=A_{21}B_{11}+A_{22}B_{21}$$

- $M_1 = (A_{11} + A_{22}) * (B_{11} + B_{22})$
- $M_2 = (A_{21} + A_{22}) * B_{11}$
- $M_3 = A_{11} * (B_{12} B_{22})$
- $M_4 = A_{22} * (B_{21} B_{11})$
- $M_5 = (A_{11} + A_{12}) * B_{22}$
- $M_6 = (A_{21} A_{11}) * (B_{11} + B_{12})$
- $M_7 = (A_{12} A_{22}) * (B_{21} + B_{22})$

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} * \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

$$\begin{array}{|c|c|c|c|c|c|}\hline & M_1 + M_4 - M_5 + M_7 & M_3 + \overline{M_5} \\ & M_2 + M_4 & M_1 + M_3 - M_2 + M_6 \\ \hline \end{array}$$

$$C_{22}=M_{1}+M_{3}-M_{2}+M_{6}=(A_{11}+A_{22})*(B_{11}+B_{22})+A_{11}*(B_{12}-B_{22})$$

$$-(A_{21}+A_{22})*B_{11}+(A_{21}-A_{11})*(B_{11}+B_{12})$$

$$=A_{11}B_{11}+A_{22}B_{11}+A_{11}B_{22}+A_{22}B_{22}+A_{11}B_{12}-A_{11}B_{22}$$

$$-A_{21}B_{11}-A_{22}B_{11}+A_{21}B_{11}-A_{11}B_{11}+A_{21}B_{12}-A_{11}B_{12}$$

$$=A_{22}B_{22}+A_{21}B_{12}$$

$$=A_{21}B_{12}+A_{22}B_{22}$$

- $M_1 = (A_{11} + A_{22}) * (B_{11} + B_{22})$
- $M_2 = (A_{21} + A_{22}) * B_{11}$
- $M_3 = A_{11} * (B_{12} B_{22})$
- $M_4 = A_{22} * (B_{21} B_{11})$
- $M_5 = (A_{11} + A_{12}) * B_{22}$
- $M_6 = (A_{21} A_{11}) * (B_{11} + B_{12})$
- $M_7 = (A_{12} A_{22}) * (B_{21} + B_{22})$

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} * \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

$$M_1+M_4-M_5+M_7$$
 M_3+M_5 M_2+M_4 $M_1+M_3-M_2+M_6$

- Count number of Multiplications and Additions
 - Multiplications: 7, Additions: 18
- Recurrence equation

$$T(n) = 7T(n/2) + 18(n/2); T(1) = 1$$

= $\Theta(7^{\log_2 n})$
= $\Theta(n^{\log_2 7})$
= $\Theta(n^{2.807})$ vs. $\Theta(n^3)$ of brute force

• There exist algos with better efficiency, but more complexity

Summary

- Large number multiplication
- Matrix Multiplication