Design and Analysis of Algorithms

L41: Intro to Dynamic Programming

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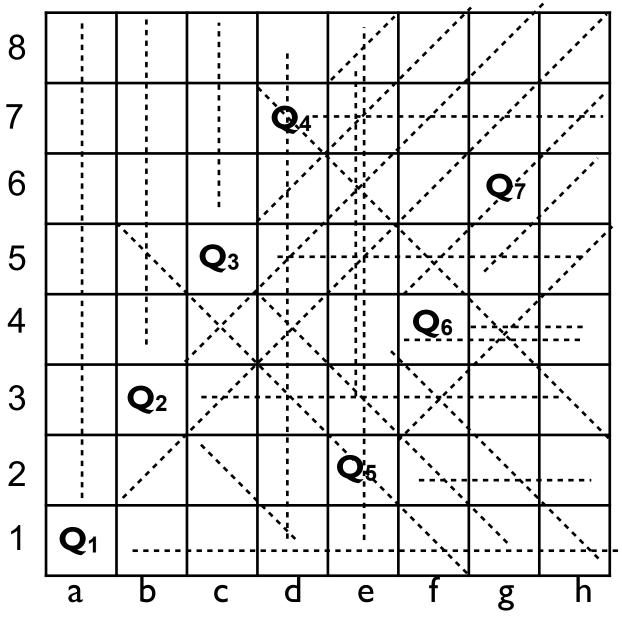
Resources

- Text book 2: Horowitz
 - -Sec 5.1, 5.2, 5.4, 5.8, 5.9
- Text book 1: Levitin
 - -Sec 8.2-8.4
- RI: Introduction to Algorithms
 - Cormen et al.
- https://en.wikipedia.org/wiki/Dynamic_programming
- https://www.codechef.com/wiki/tutorial-dynamic-programming
- https://www.hackerearth.com/practice/algorithms/dynamicprogramming/introduction-to-dynamic-programming-I/tutorial/

Overview of Backtracking

- Backtracking
 - Start from some solution.
 - Keep exploring for next part of solution
 - When exploration of solution stops (not possible to proceed further)
 - Resume back from the last point where decision was made to explore the current path.
 - Explore with the next path.

8-Queens Problem



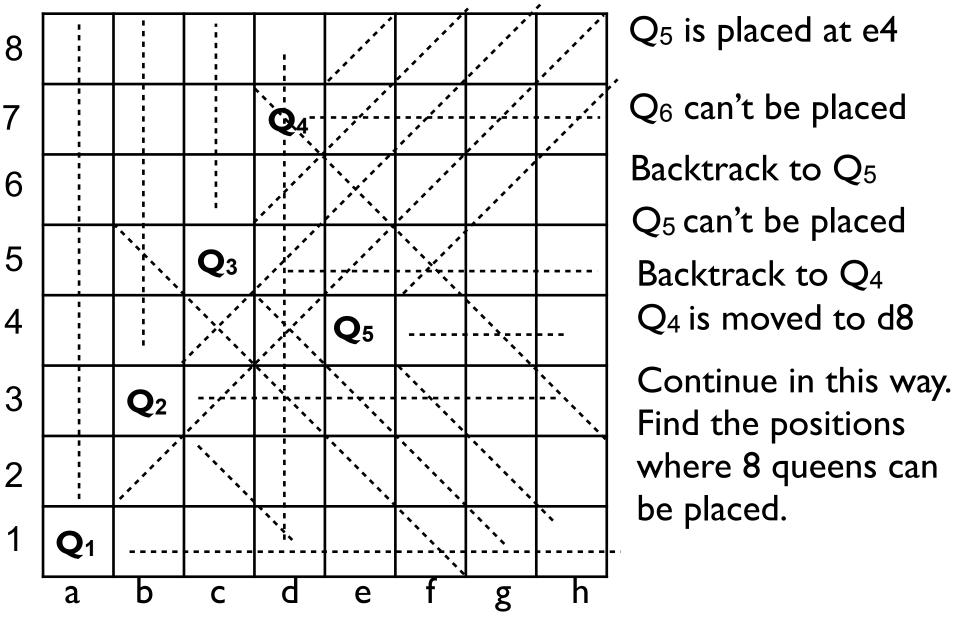
Q₈ can't be placed.

Backtrack to Q₇ which can't be placed too

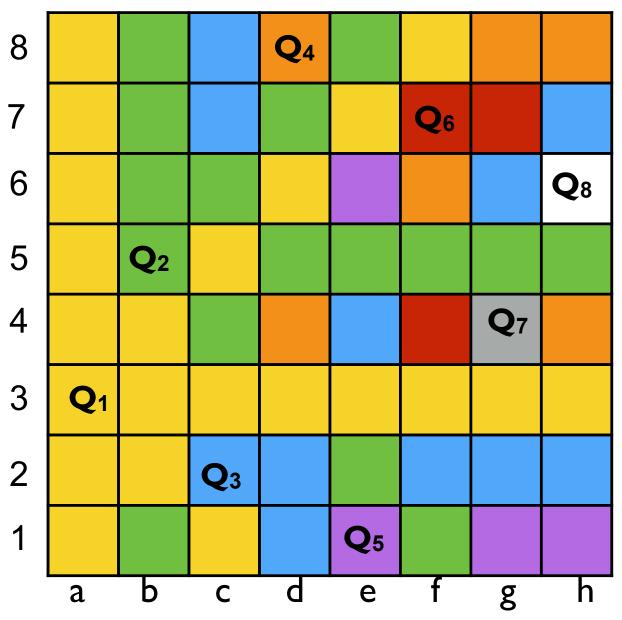
Backtrack to Q₆ which can't be placed too

Backtrack to Q₅ which can be placed at e4

8-Queens Problem

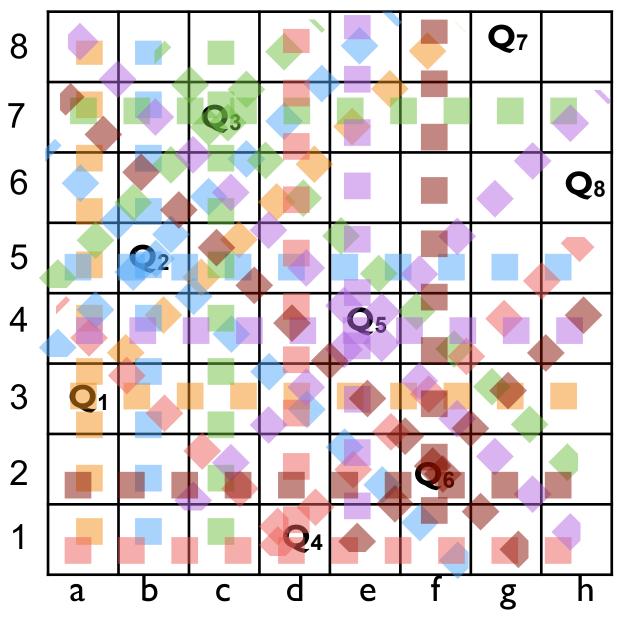


8-Queens Problem



Continuing in this way. positions of 8 queens.

8-Queens Problem: Soln 2



Continuing further another solution for 8 queens problem

3-Color Problem

5

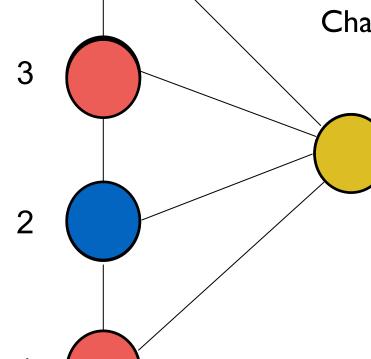
Consider three colors are: Red, Blue, Yellow Color the nodes with these 3 colors such that no two adjacent nodes have same color.

Can't color node 5, so backtrack Change 4 to Blue

> Can't color node 5, so backtrack Can't change 4 to Yellow Backtrack to 3

Solution: 1-R, 2-B, 3-R, 4-B, 5-Y

Are there other solutions?



Sum of Subset Problem

- Given a set S of numbers and a value m,
 - Find all subsets $S_i \subseteq S$ so that their sum of elements in S_i equals m.
 - An element in a subset is to be considered only once.
- Example

```
S = \{11, 13, 24, 7\}, and m = 31
```

Possible subsets are

```
S_1 = \{11, 23, 7\}

S_2 = \{24, 7\}
```

Backtracking: General Method

- General solution is an n-tuple $(x_1, ..., x_n)$, where
 - $-x_i$ is chosen from some finite set S_i .
 - While choosing x_i , it has to follow some constraints
 - or meet a criterion function $P(x_1, ..., x_n)$
- Suppose, the size of each set S_i is m_i
- Then, total number of possible tuples are $M=m_1*m_2*...*m_n$
- Identify those tuples that satisfies the contraints i.e.
 Criterion function.
- Backtracking approach provides the answer in far fewer trials than M.

Backtracking: 8-Queens Method

- Let queens are numbered 1 thru 8, i.e. $Q_1, ..., Q_8$
- Each queen must be on a separate column (and row)
 - For simplicity, let's say Q_i is placed on i^{th} column.
- Thus, solution can be represented by an 8-tuple

```
\{x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8\}
```

- where $x_1 \approx a$, $x_2 \approx b$, $x_3 \approx c$, $x_8 \approx h$
- Each queen must be on a separate row.
- Thus, each x_i can have a value from 1 to 8.
 - Thus, constraint is $x_i \in S_i = \{1, 2, 3, 4, 5, 6, 7, 8\}$
- Solution space size before and after the constraint
 - before: 88, after: 8!
- Representation for solution-1
 - **{**3,5,2,8,1,7,4,6}

Backtracking: Sum of Subsets

- Problem: $S = \{11, 13, 24, 7\}$, and m = 31
- Solution approach 1:
 - Consider 4-tuple $\{x_1, x_2, x_3, x_4\}$
 - where, $x_i \in S_i = \{0, 1\}$
 - Size of solution space: 2n
 - possible solutions
 - {1,1,0,1}
 - {0,0,1,1}
- Solution approach 2
 - Solution contains the index values of elements.
 - Solution is a vector of varying dimentions
 - Possible solutions
 - (1, 2, 4)
 - (3,4)

Backtracking: 3-Color problem

- Problem: $G = \{ \forall, E \}$, and 3 colors to color the graph
- Solution vector: n-tuple $(x_1, ..., x_n)$
 - where $x_i \in S_i = \{R, B, Y\}$
- Size of total solution space: 3n
 - An edge reduces solution space from 3^2 to 3*2=6
 - Any path of length k reduces the solution space from 3^{k+1} to $3*2^k$

Summary

- Overview of backtracking
- Problem examples for backtracking
 - 8-queens problems
 - Sum of subsets
 - 3-color problem
- Solution space
- Possible solution space.