Design and Analysis of Algorithms

L11: Algo Design Ideas

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Resources

- Martin Richards
 - Notes on Data Structures and Algorithms
- Text book 2: Horowitz
- Text book I: Levitin
- https://visualgo.net/en

Ideas on Algorithm Design

- Recognize a variant on known problem
- Reduce to a simpler problem
- Divide and Conquer
- Estimation of cost by recurrence relation
- Dynamic Programming
- Greedy Algorithm
- Back tracking
- Hill Climbing
- Identify wasted work in a simple method
- Find a mathematical lower bound
- Million Monkey method

Variant on a Known Problem

- Try to identify a related problem whose solution is known
- Use solution of one problem to solve essentially tricky part of other problem.
- Solve the problem in totality
- Example I: multiply two numbers
 - Known variant: add two numbers
 - How to use addition of two numbers to achieve multiplication
- Example 2: Class assignment submission
 - variant: Identify the classmate who has done it
 - Copy the assignment with change of USN

Reduction to a Simpler Problem

- Typically works where induction methods are used to show correctness proof
- Recursive functions take this appproach
- Example: Compute n²

$$n^2 = (n-1)^2 + 2(n-1)$$

• Example: factorial n

$$n! = n*(n-1)!$$

Example: GCD computation

$$gcd(x,y) = gcd(r,x)$$
 where $y=kx+r$

Divide and Conquer Algo

- Divide (break) the problem (size n) into similar sub problems
 - Size of sub problems should be some factor of original e.g. n/c
 - When small enough, solve by brute force
- Conquer (Solve) the sub-problem
 - Use recursion to solve small problem
- Combine (Merge) the solution of sub-parts
- The cost is
 - cost of breaking
 - cost of solving subproblem
 - cost of combining

Cost Estimation by Recurrence

- Use recurrence relation to find the cost of operation
 - Once the cost is known, then algorithm framework is worked out
 - Design the algorithm on how to break the given problem
- Example: Mergesort

$$f(n) = 2f(n/2) + n$$

-> $f(n) = O(nlog_2n)$

• Example: Word search in English dictionary

$$f(n) = f(n/k) + c$$

-> $f(n) = log_k n$

Dynamic Programming

- Build a table of solutions to smaller versions of the problem
- Work towards the solution by using this table
- Example: Binomial Coefficient

• Example: Computing n2

Greedy Algorithms

- Useful when some sort of optimizations is involved.
- Basic idea:
 - Perform whatever operation contributes most towards the final goal
 - Next step will be same approach after the previous step
 - Note: Final solution may not always be optimal
- Example:
 - Eating food in restaurant: get whatever hot now
 - No waiting is to be done
 - Taking exam (you know all answers) but time is less
 - Which questions to start writing

Backtracking

- When algorithm involves search (explore)
 - backtracking is useful
- Split the search procedure in multiple paths(parts)
 - Start with one path
 - If solution not find and path ends,
 - Backtrack to previous starting point and
 - search on next path
- Example:
 - Finding a solution to the maze (childhood comicbooks)
 - Chessboard: queens placement
 - Chessboard: Knights traversal

Hill Climbing

- Generally, used in optimization problems
- Start with some feasible (non-optimal) solution
- Incrementally work towards improving the solution
- May not always find the optimal solution
- Analogy: Climbing to hilltop in hilly region
- Example:
 - Getting a job(placement)
 - Start with some placement (job)
 - Work towards a better placement (better company)

Identify Wasted Work

- Design a simple solution to the problem
- Analyze the solution and
 - identify critical costly part of the solution
- Attack the weakness of critical part to improve the solution
- Example (real life):
 - Phone converstation, and you are put on hold
 - can do some useful work while on call hold
- Example (CS): CPU Scheduling
 - when a processing is waiting for I/O
 - other process is assigned CPU

Seek Mathematical Lower Bound

- Establish a proof that some task must take at least certain minimum time
- Use this insight to design the algorithm
- A properly proved lower bound can prevent wasted time seeking improvement
- Example: Sorting n numbers:
 - min time: ceil(log2n!) comparisons
 - Consider n=5, min time = $log_2(120) = 7$
 - Design the algorithm

Million Monkey (MM) Method

- Problem: Give one typewriter each to million monkeys (random character press) and a million years, and one of the monkey will write the Shakespear play.
 - Idea: give a problem to a group of researchers and sufficient time, they will be able to find the solution of the problem.
- Example: expected output 'hello'.
 - Assume typewriter has 50 keys
 - Probability of not typing 'hello' = $1-(1/50)^5$
 - Probability of not typing hello by million monkeys in 10000 tries = $(1-(1/50)^5)^{10} = 1.2*10^{-14}$
 - Prob. of typing hello = $1-1.2*10^{-14}=1$

Fun Exercise of Game of 128 numbers

 A practical fun example of Data structures and Algorithm

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|-----|-----|-----|-----|-----|-----|-----|-----|
| 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 |
| 25 | 26 | 27 | 28 | 29 | 30 | 31 | 32 |
| 33 | 34 | 35 | 36 | 37 | 38 | 39 | 40 |
| 41 | 42 | 43 | 44 | 45 | 46 | 47 | 48 |
| 49 | 50 | 51 | 52 | 53 | 54 | 55 | 56 |
| 57 | 58 | 59 | 60 | 61 | 62 | 63 | 64 |
| 65 | 66 | 67 | 68 | 69 | 70 | 71 | 72 |
| 73 | 74 | 75 | 76 | 77 | 78 | 79 | 80 |
| 81 | 82 | 83 | 84 | 85 | 86 | 87 | 88 |
| 89 | 90 | 91 | 92 | 93 | 94 | 95 | 96 |
| 97 | 98 | 99 | 100 | 101 | 102 | 103 | 104 |
| 105 | 106 | 107 | 108 | 109 | 110 | 111 | 112 |
| 113 | 114 | 115 | 116 | 117 | 118 | 119 | 120 |
| 121 | 122 | 123 | 124 | 125 | 126 | 127 | 128 |

DAA/

Game:

- . Go thru a set of cards
- . Say Y/N if present or not
- . You will get you number graphically displayed to you

Q?:

Which algorithm we are discussing?

Aim: Can we find more such examples

RPR/

Game of 128 numbers - b

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|----|----|----|----|----|----|----|----|
| 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 |
| 25 | 26 | 27 | 28 | 29 | 30 | 31 | 32 |
| 33 | 34 | 35 | 36 | 37 | 38 | 39 | 40 |
| 41 | 42 | 43 | 44 | 45 | 46 | 47 | 48 |
| 49 | 50 | 51 | 52 | 53 | 54 | 55 | 56 |
| 57 | 58 | 59 | 60 | 61 | 62 | 63 | 64 |
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Exercise G

- Exercise G
 - Work out the remaining 6 cards

Summary

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