



Topic 8

Data Hazards

Hazards

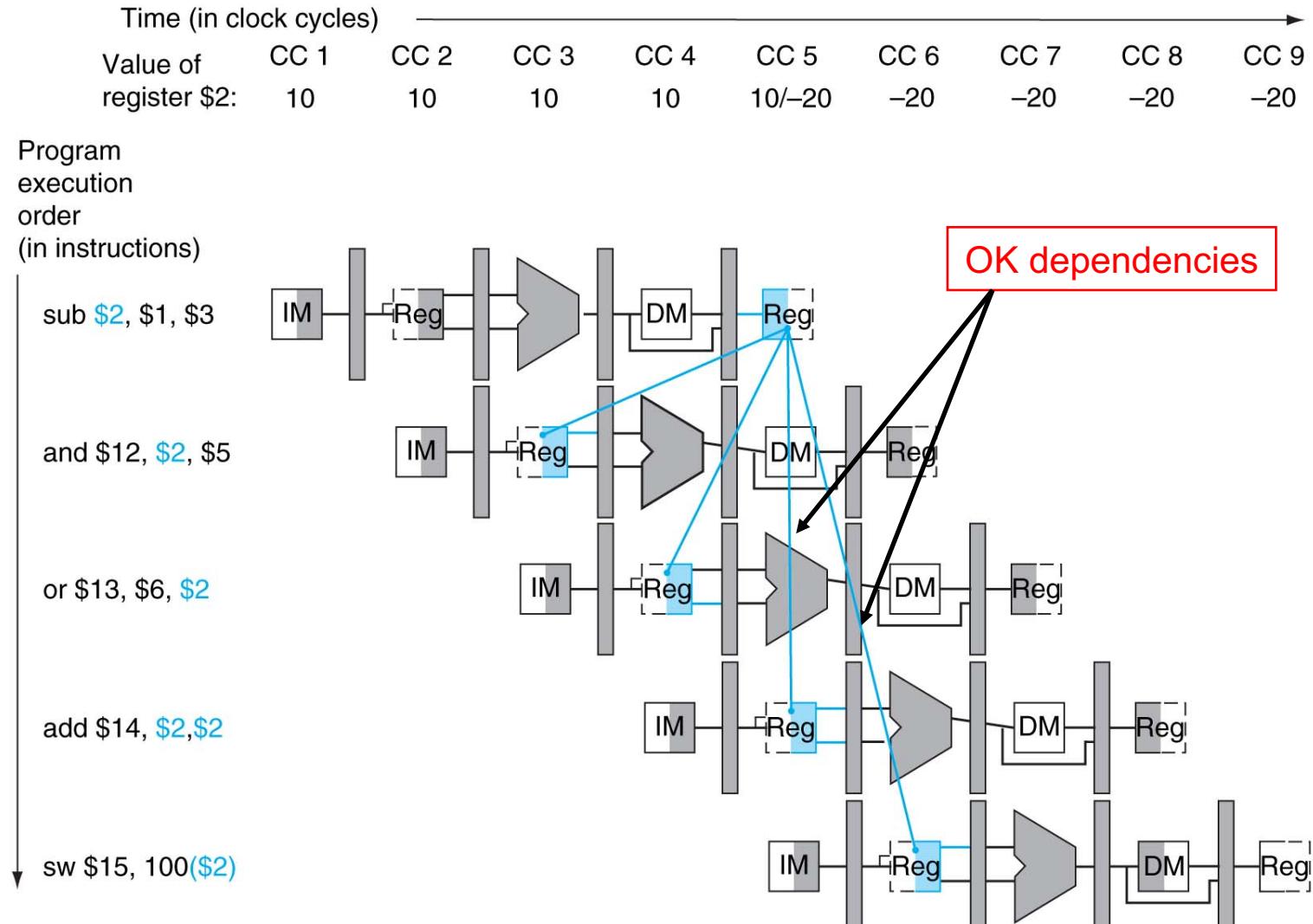
- Situations that prevent starting the next instruction in the next cycle
- Data hazard
 - Need to wait for previous instruction to complete its data read/write
- Control hazard
 - Decision on control action depends on previous instruction
- Structure hazards
 - A required resource is busy

Data Hazards in ALU Instructions

- Consider this sequence:

```
sub $2, $1, $3
and $12, $2, $5
or $13, $6, $2
add $14, $2, $2
sw $15, 100($2)
```

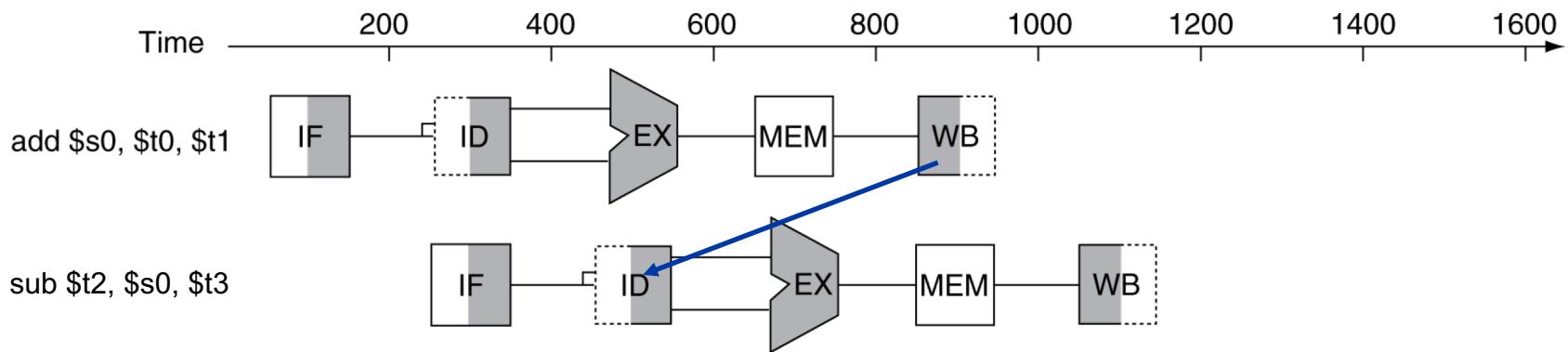
Data Dependencies



Data Hazards

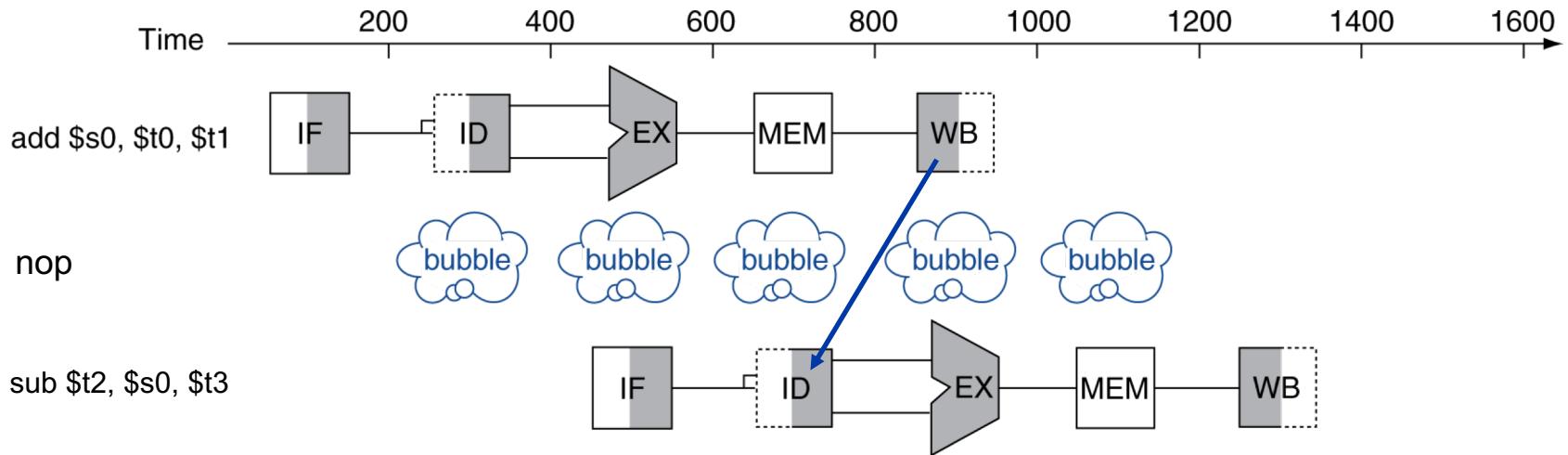
- An instruction must be delayed
 - If an instruction depends on completion of a data by a previous instruction
 - By inserting ***bubbles*** or ***stalls*** (**nop** instruction)
 - Example

add \$s0, \$t0, \$t1
sub \$t2, \$s0, \$t3



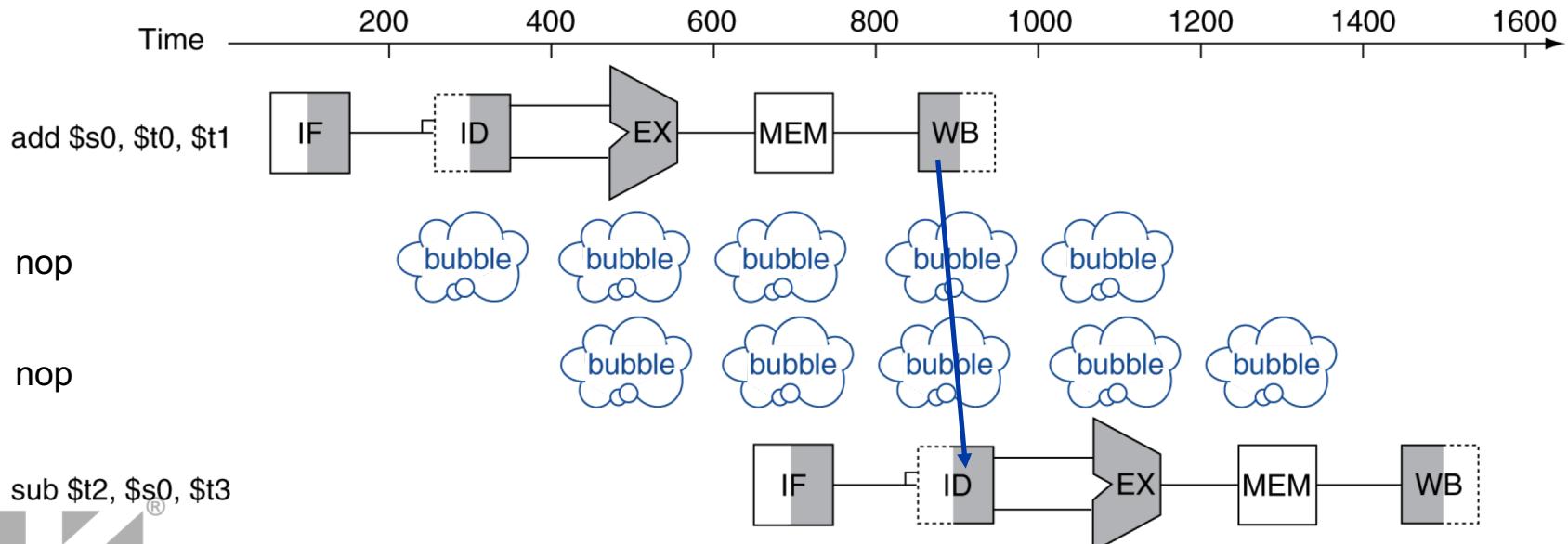
Data Hazards

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- ```
nop # no operation
 # machine code: 0x00000000
```



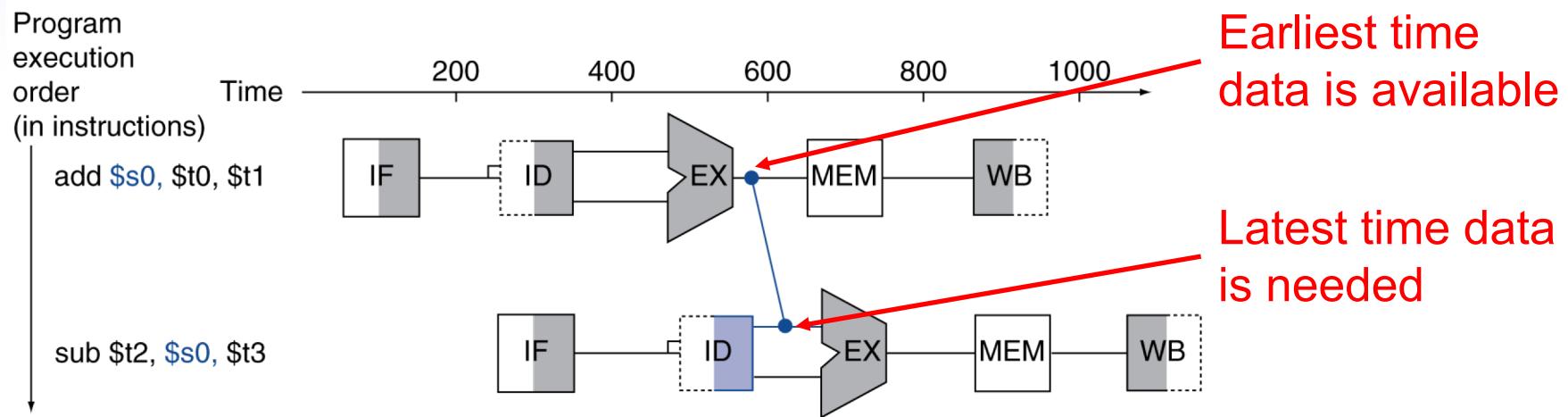
# Data Hazards

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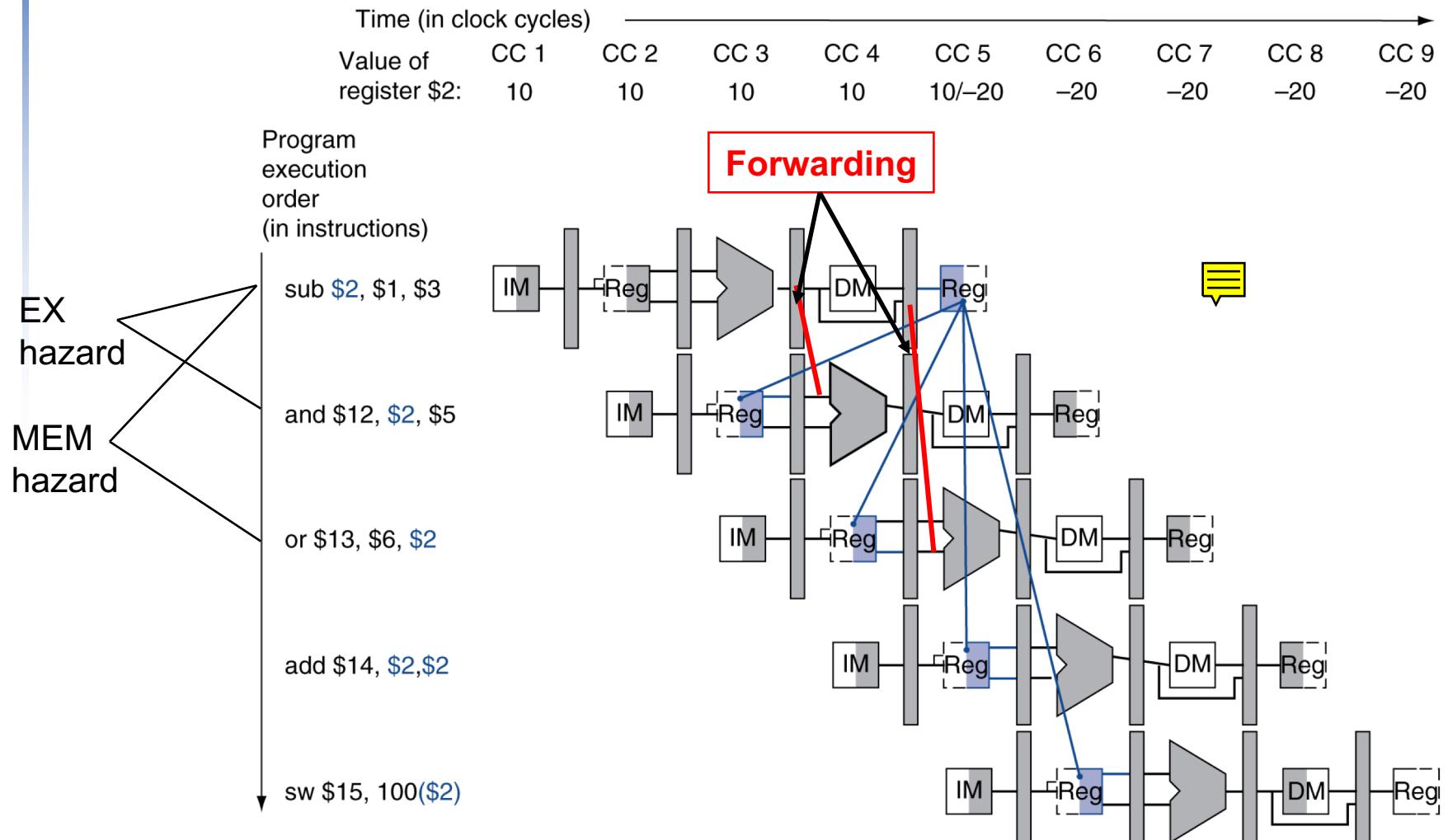
Forwarding (aka Bypassing)

- Inserting nop (delay) compromises performance
- Data hazard can be solved with **Forwarding**
- Use result as soon as it is available
 - Don't wait for it to be stored in a register
 - Requires extra hardware in the datapath

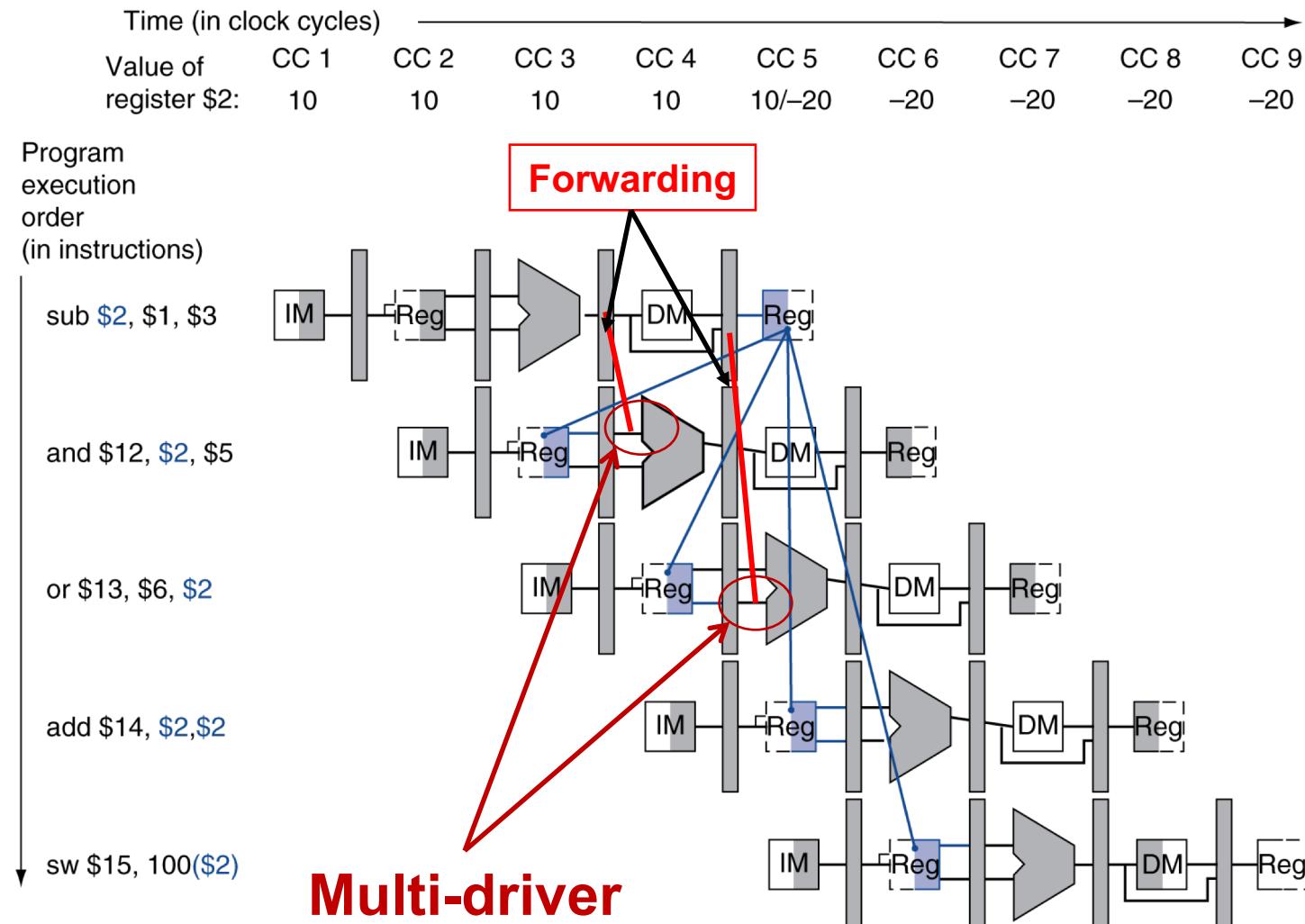


Can't connect input and output of a component directly

Forwarding from Registers

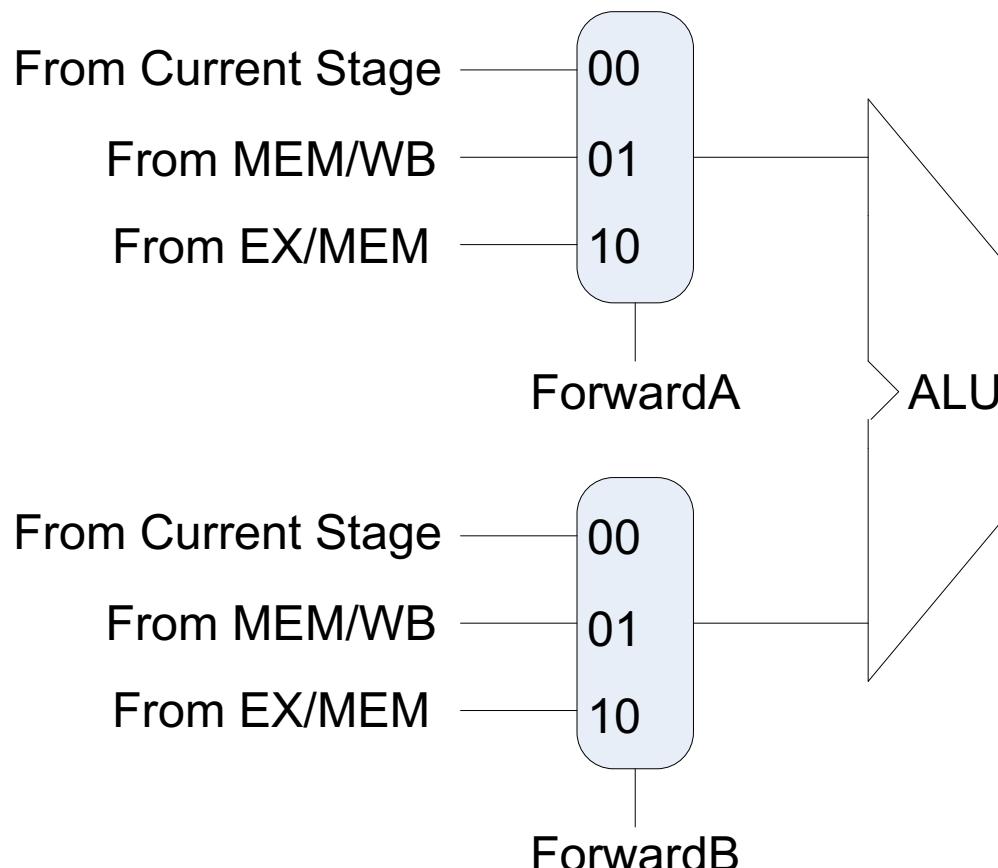


Dependencies & Forwarding

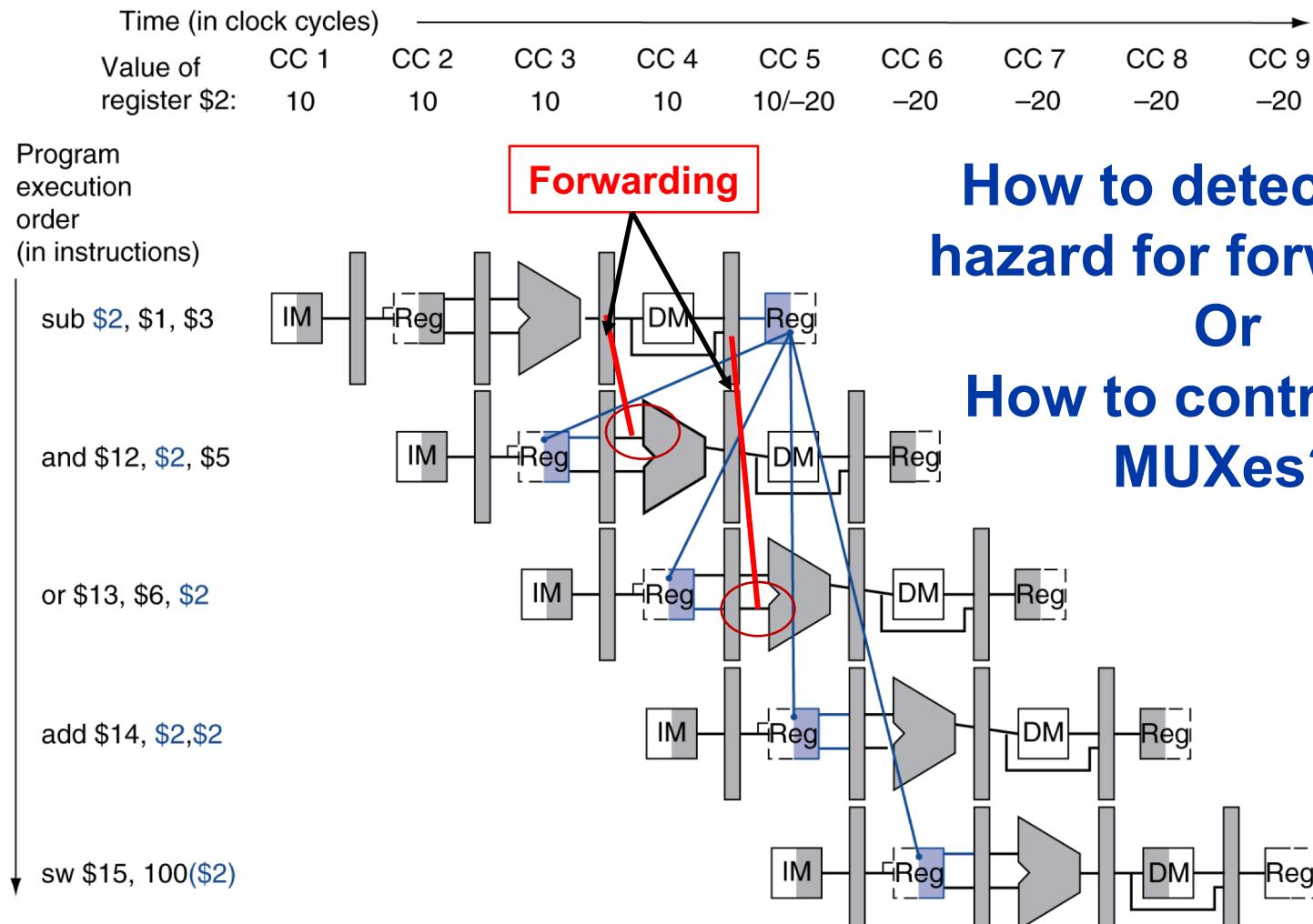


Forwarding Paths

- Forwarding paths are created between stage pipeline registers and ALU inputs
 - By using MUXes



Dependencies & Forwarding



How to detect data hazard for forwarding
Or
How to control the MUXes?

Detect the Condition for Forwarding

- Denotation:
 - E.g. ID/EX.RegisterRs = register number for Rs residing in ID/EX pipeline register
- ALU operand register numbers in EX stage are given by
 - ID/EX.RegisterRs, ID/EX.RegisterRt

- Data hazard conditions include

Hazard
between
1&2

Hazard
between
1&3

{ 1a. EX/MEM.RegisterRd==ID/EX.RegisterRs }

{ 1b. EX/MEM.RegisterRd==ID/EX.RegisterRt }

{ 2a. MEM/WB.RegisterRd==ID/EX.RegisterRs }

{ 2b. MEM/WB.RegisterRd==ID/EX.RegisterRt }

Fwd from
EX/MEM
pipeline reg

Fwd from
MEM/WB
pipeline reg

Detecting the Need to Forward

- Only if an earlier instruction will write back to a register, determined by
 - EX/MEM.RegWrite $\equiv 1$
 - MEM/WB.RegWrite
- And only if Rd for that instruction is not \$zero (constant register), determined by
 - EX/MEM.RegisterRd $\neq 0$
 - MEM/WB.RegisterRd $\neq 0$

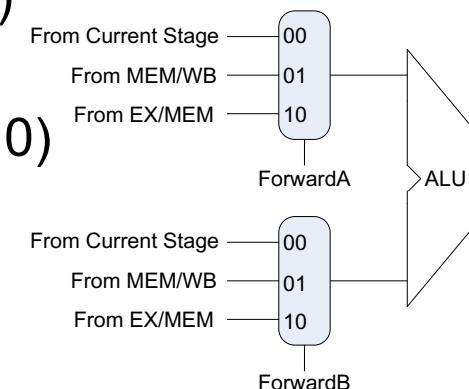
Revised Forwarding Conditions

EX hazard

- if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd == ID/EX.RegisterRs))
MUX select signal ForwardA = 10
- if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd == ID/EX.RegisterRt))
MUX select signal ForwardB = 10

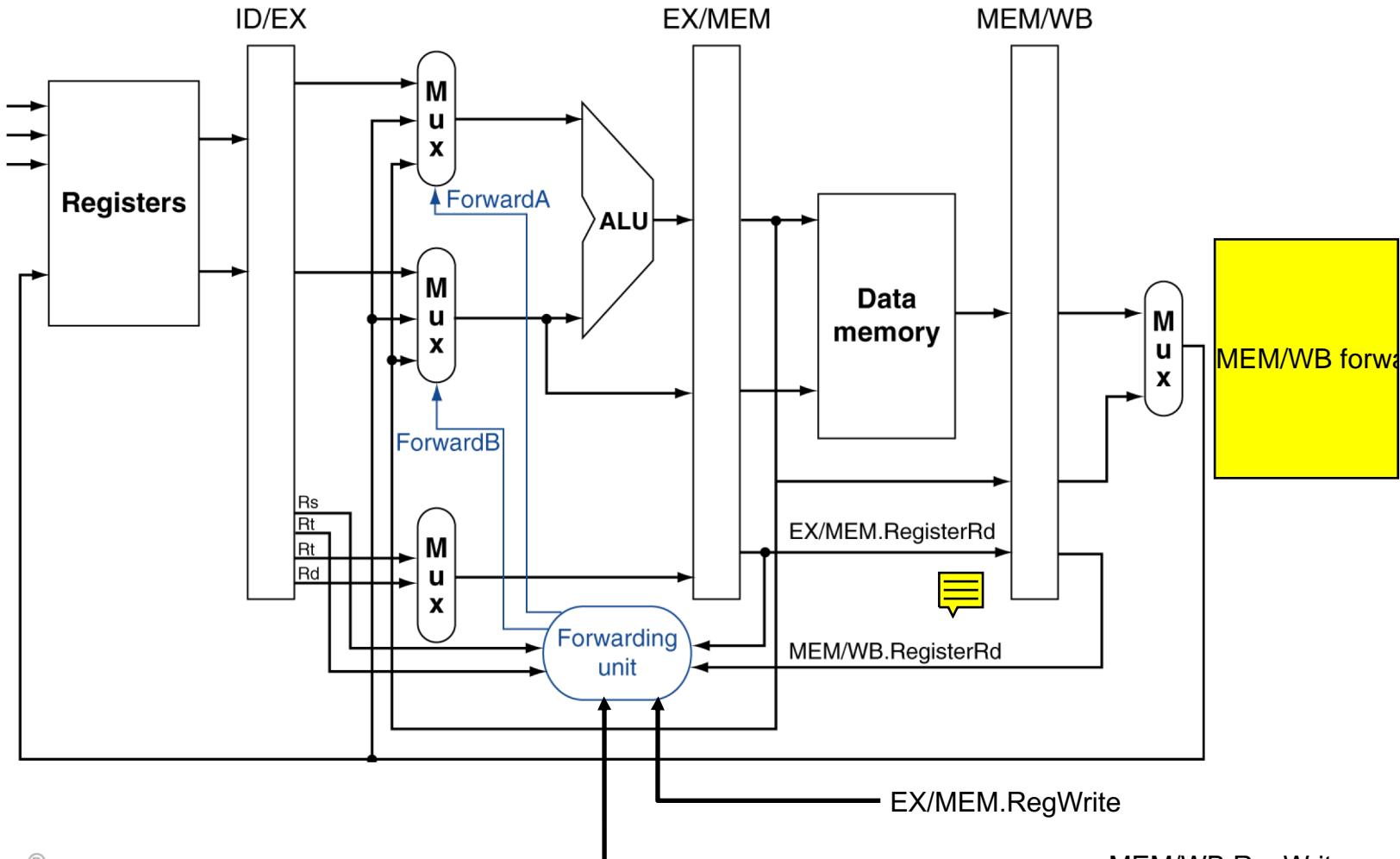
MEM hazard

- if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and (MEM/WB.RegisterRd == ID/EX.RegisterRs))
MUX select signal ForwardA = 01
- if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and (MEM/WB.RegisterRd == ID/EX.RegisterRt))
MUX select signal ForwardB = 01

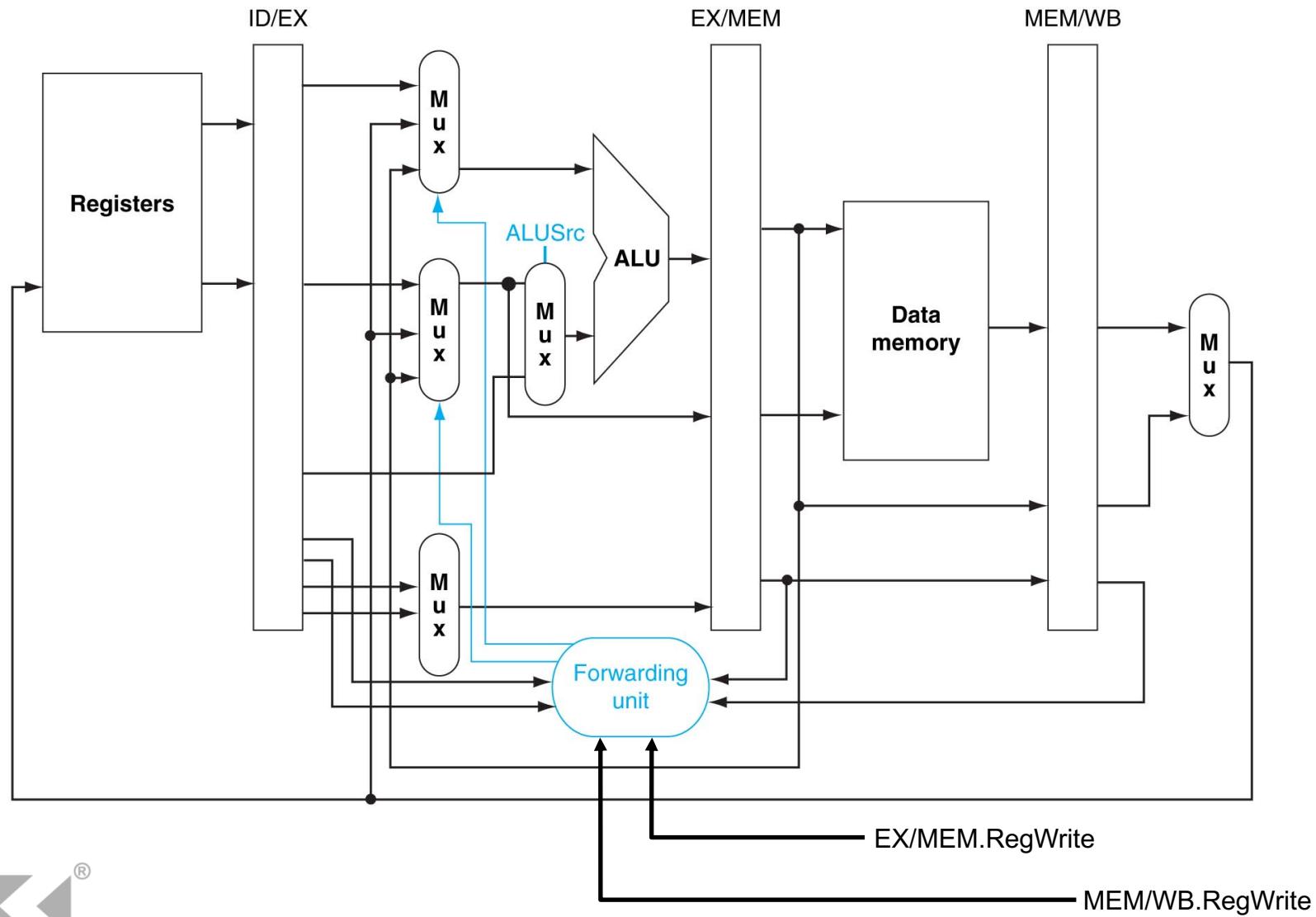


Forwarding Paths

Forwarding Unit in EX stage



Forwarding Paths with ALUSrc



MK®

Double Data Hazard

- Consider the sequence:

add \$1, \$1, \$2

add \$1, \$1, \$3

add \$1, \$1, \$4

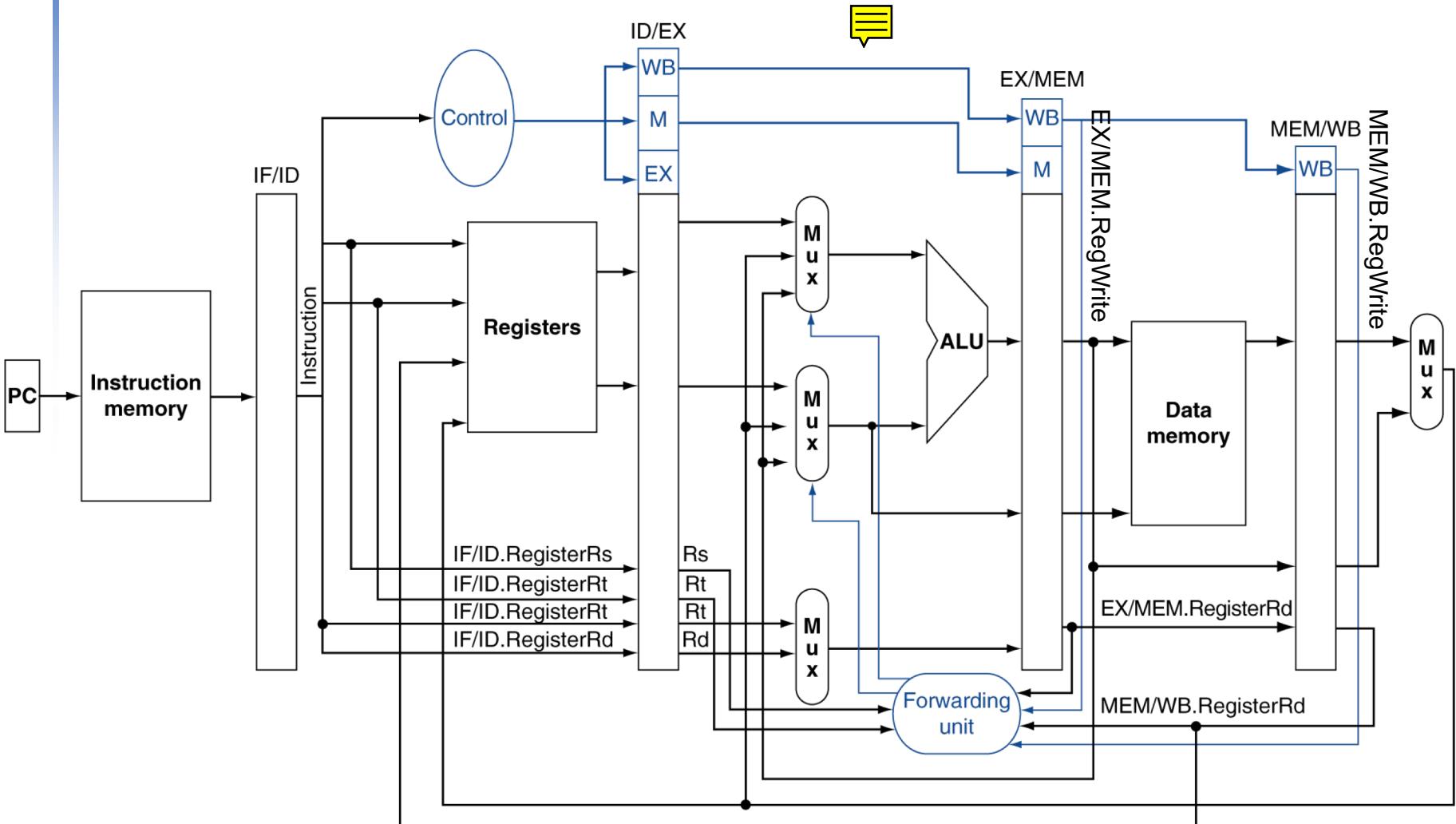
- Conditions for both hazards are satisfied
 - Want to use the most recent
 - Forwarding between 1st and 3rd – MEM hazard condition should be disabled
- Revise MEM hazard condition
 - Only fwd if EX hazard condition isn't true



Revised Forwarding Condition

- MEM hazard (not EX hazard)
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0) and (MEM/WB.RegisterRd = ID/EX.RegisterRs) and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0) and (EX/MEM.RegisterRd = ID/EX.RegisterRs)))
ForwardA = 01
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0) and (MEM/WB.RegisterRd = ID/EX.RegisterRt) and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0) and (EX/MEM.RegisterRd = ID/EX.RegisterRt)))
ForwardB = 01

Datapath with Forwarding Unit



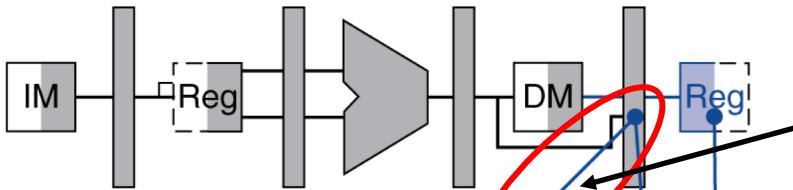
Load-Use Data Hazard

- Can't always avoid stalls by forwarding
 - If value not computed when needed
 - Can't forward back in time!

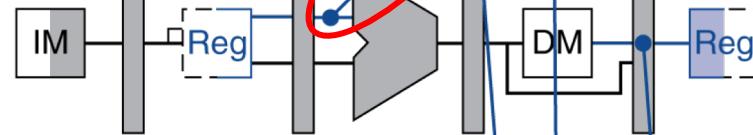
Load-Use Data Hazard

Program
execution
order
(in instructions)

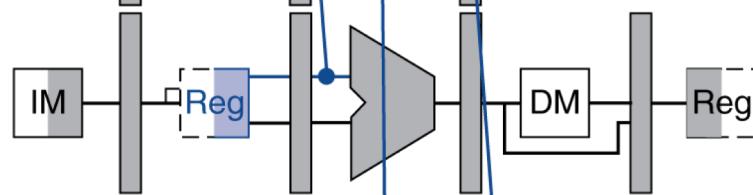
lw \$2, 20(\$1)



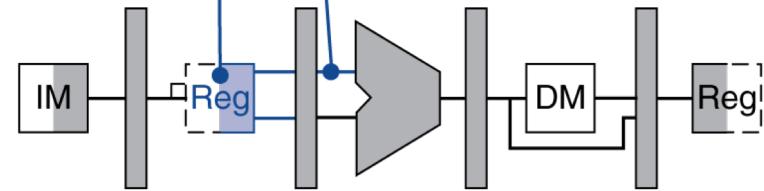
and \$4, \$2, \$5



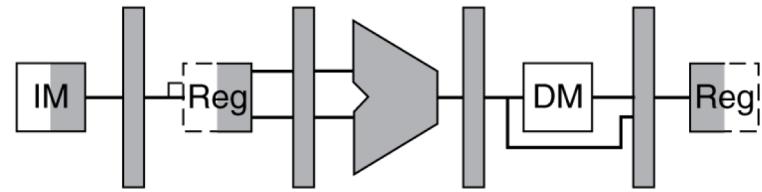
or \$8, \$2, \$6



add \$9, \$4, \$2



slt \$1, \$6, \$7



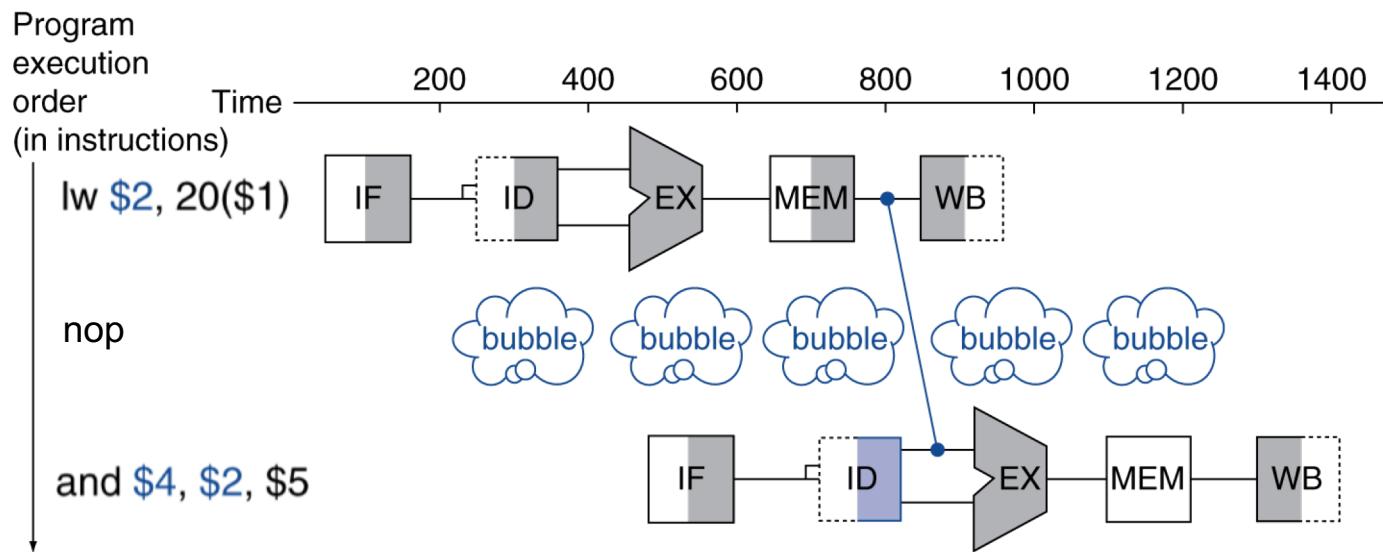
Referring data
in later time.
Need to stall
for one cycle

Load-Use Hazard Detection

- Need to check only for a *load* instruction, determined by
 - ID/EX.MemRead
- Load-use hazard when
 - ID/EX.MemRead and
$$((\text{ID/EX.RegisterRt} == \text{IF/ID.RegisterRs}) \text{ or } (\text{ID/EX.RegisterRt} == \text{IF/ID.RegisterRt}))$$
- Hazard Detection is in **ID** stage
- If detected, stall and insert bubble

Load-Use Data Hazard

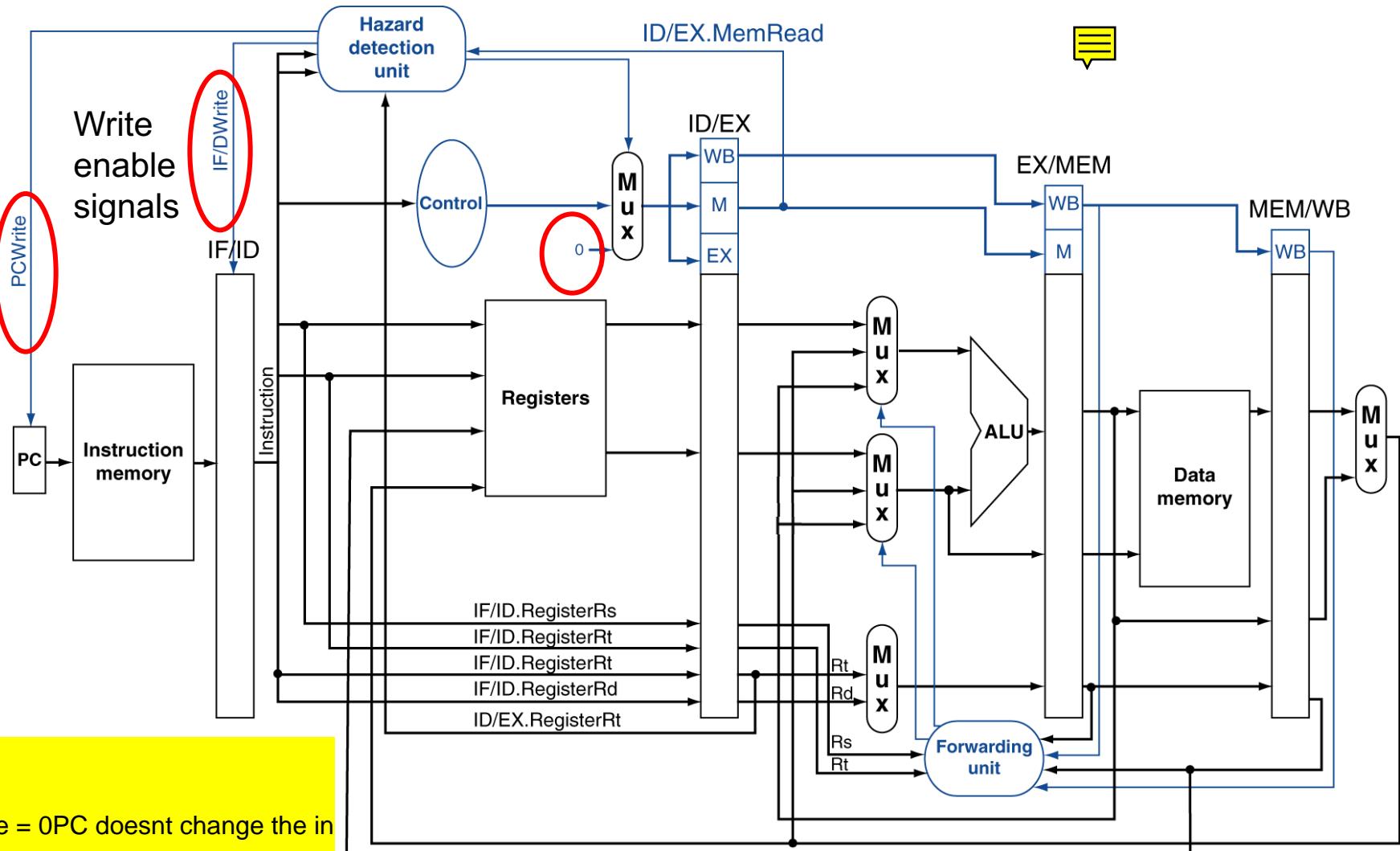
- Resolve by inserting stalls



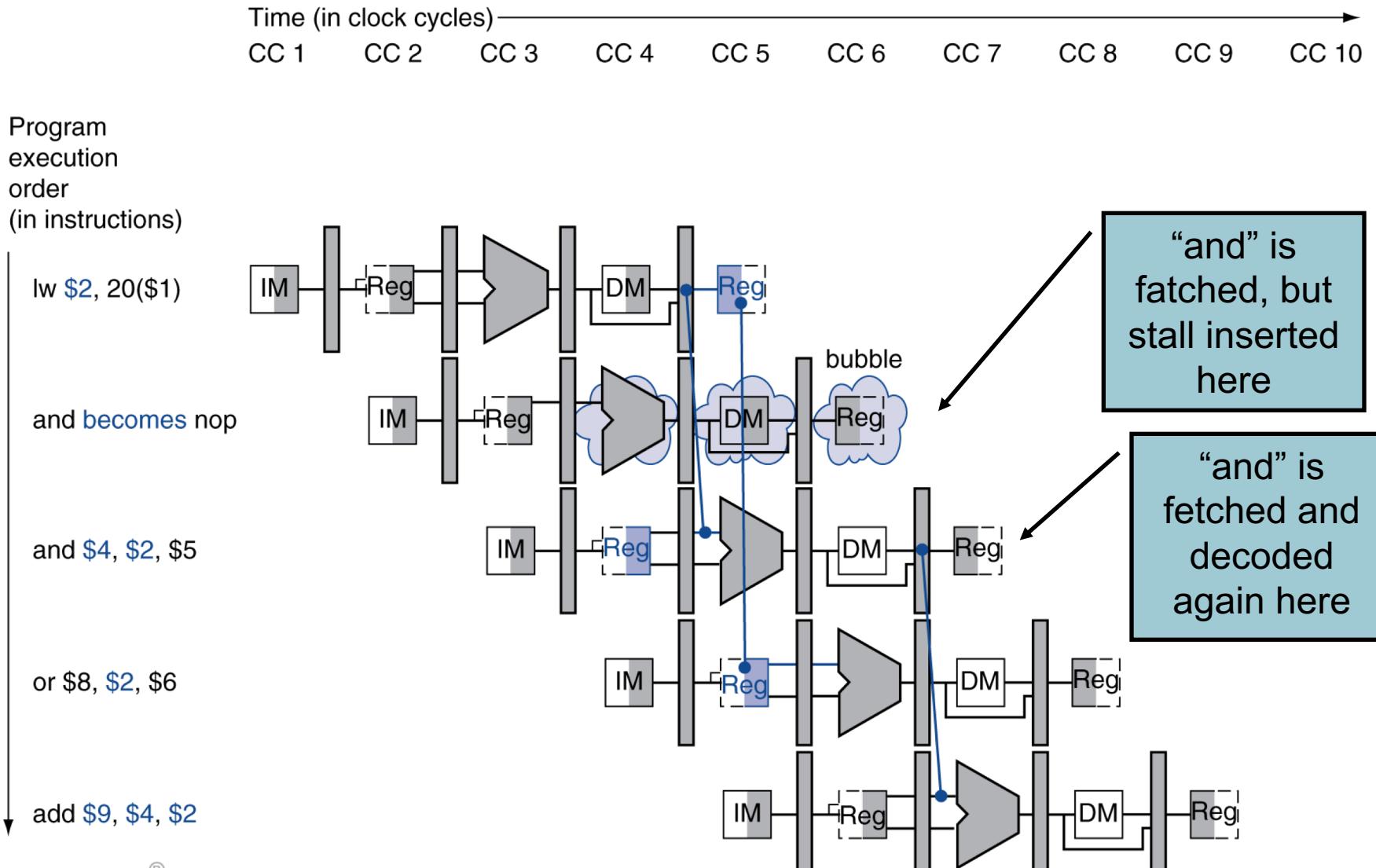
How to Stall the Pipeline?

- (how to insert bubbles?)
- Force control signals in ID/EX register to 0's
 - EX, MEM and WB in following cycles do no-operation (nop) with those 0 control signals
- Prevent updates of PC and IF>ID registers
 - Instruction in IF>ID is held and decoded again
 - PC is held, so the same instruction is fetched again

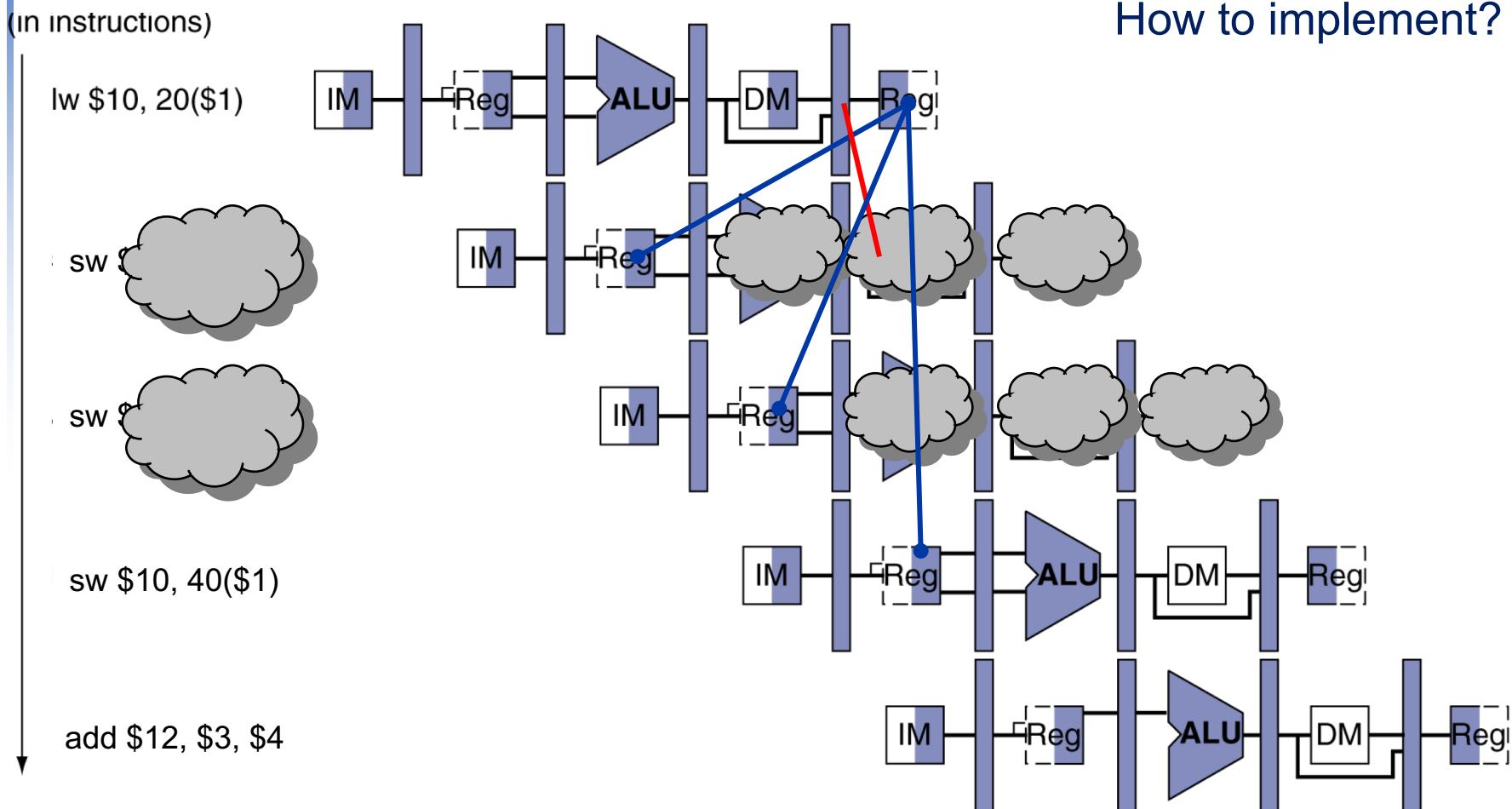
Datapath with Hazard Detection



Stall/Bubble in the Pipeline



Question

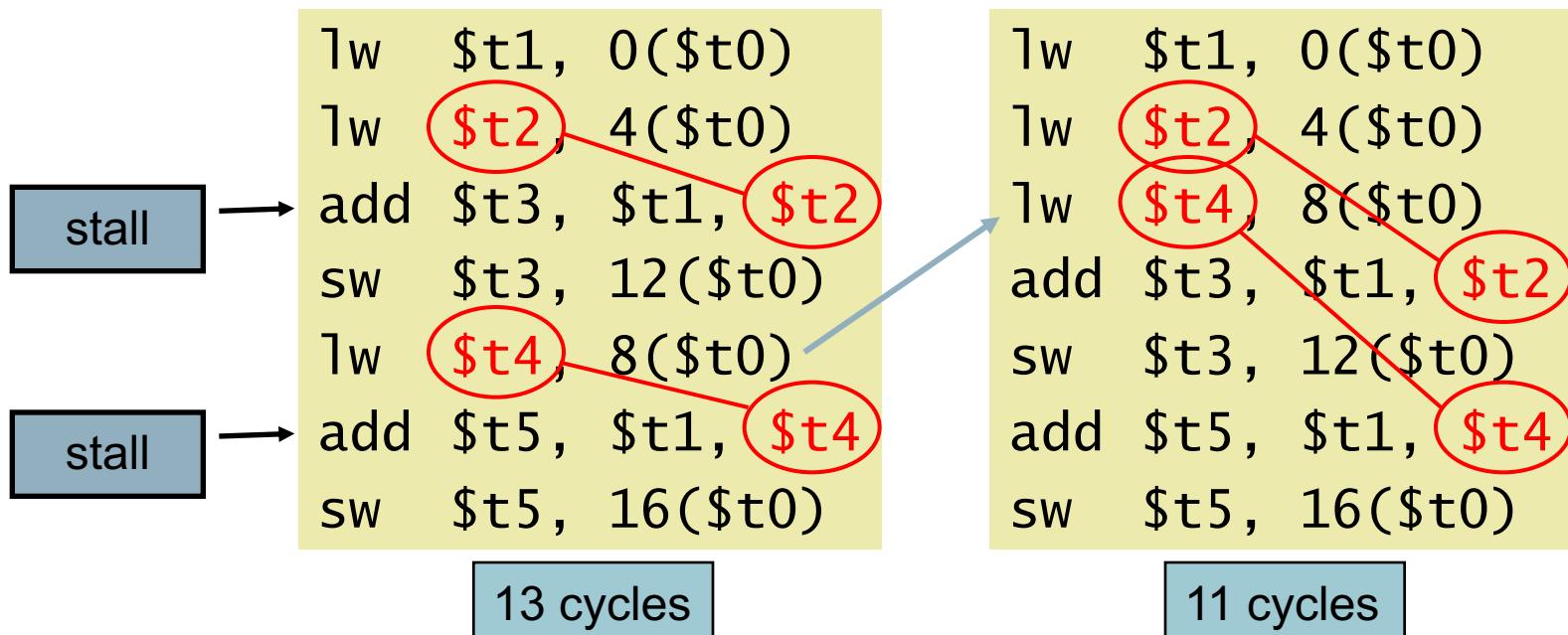


This delays the execution too much, how to fix?



Code Scheduling to Avoid Stalls

- Reorder code to avoid use of load result in the next instruction
- C code for $A = B + E; C = B + F;$



Structure Hazards

- Conflict for use of a resource
- If common memory for program and data
 - Load/store requires data access
 - Instruction fetch would have to *stall* for that cycle
 - Would cause a pipeline “bubble”
- Hence, pipelined datapaths require separate instruction/data memories
 - separate instruction/data caches (Harvard)
 - Common low-level memory (Von Neumann)

Stalls and Performance

The BIG Picture

- Stalls reduce performance
 - But are required to get correct results
- Hardware can be improved to take care of the hazards automatically
 - forwarding
- Compiler can arrange code to avoid hazards and stalls
 - Requires knowledge of the pipeline structure